

# Programming Paradigms Compendium

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1. Find a Haskell expression whose type is  $(\text{Ord } a1, \text{Eq } a2) \Rightarrow a2 \rightarrow a2 \rightarrow (a1, a1) \rightarrow a1$

**Comment** The type is a function type. The function must be curried and takes three arguments, the third of which is a pair. There are type class constraints for  $a1$  and  $a2$ . In this case, comparisons regarding  $<$  (due to  $\text{Ord}$ ) and  $=$  (due to  $\text{Eq}$ ) can be placed in a condition.

```
plop' x y (u,v) = if (x == y) && (u > v) then u else u
```

2. Use list abstraction to define a function `flop` that, when given a list of pairs, returns a list of pairs whose components are swapped. The list can be empty. For example, `flop [(1,'a'),(3,'r'),(9,'e')]` should return the list `[(1,'a'),(3,'r'),(9,'e')]`.

**Comment** Here we have

```
flop xs = [(y,x) | (x,y) <- xs]
```

3. A list  $[a_1, a_2, \dots, a_n]$  is *decreasing* if  $a_1 \geq a_2 \geq \dots \geq a_n$ . Write a function `descending` that returns `True` if the list it receives as an argument is decreasing, and `False` otherwise. For example, `descending [6,5,5,1]` should return `True`, and `descending ["plip","pli","ppp"]` should return `False`.

**Comment** The solution is

```
descending (x:y:ys) | x >= y = descending (y:ys)
                  | x < y = False
descending xs = True
```

Note which patterns are needed.

4. Implement the function `sumrows`. The function takes a list of lists of numbers and returns a one-dimensional list of numbers, where each number is equal to the sum of the corresponding row in the input list. If a list is empty, its sum is `0`. For example, `sumrows [[1,2], [3,4]]` should give us `[3, 7]`, and `sumrows [[],[],[1]]` should give us `[0,0,1]`.

**Comment** This is a classic use of `map` together with `sum`.

```
sumrows xs = map sum xs
```

5. Define a Haskell datatype `Aexp` for arithmetic expressions with addition, multiplication, numbers, and variables. The construction rules in the abstract syntax are

$$E ::= n \mid x \mid E_1 + E_2 \mid E_1 \cdot E_2$$

Assume that variables  $x$  are strings, and that numbers  $n$  are integers.

**Comment** The construction rules can be translated directly. Each case in the construction rules should have its own term constructor.

```
data Aexp = N Int | V String | Plus Aexp Aexp | Mult Aexp Aexp
```

6. Define a function `thesame` that takes a list of pairs `xs` and returns a list of pairs where the first and second components are equal. For example,

```
thesame [(1,2),(4,4),(6,7),(17,17)]
```

should return

```
[(4,4),(17,17)]
```

What should the type of `thesame` be?

**Comment** The type class `Eq` (book, page 31) is the family of types that have an implementation of equality. And from the description of `thesame`, we can see that the components of a pair must be able to be checked for equality. So

```
thesame :: Eq a => [(a,a)] -> [(a,a)]
```

7. Use recursion to define a Haskell value `letters`, which is a sequence of actions that does the following:

- Receives a string
- Prints the characters in the string one at a time, each followed by a newline

What we want is, for example:

```
*Main> letters
dingo
d
i
n
g
o
*Main>
```

**Comment** Here it is:

```
letters = do
  x <- getLine
  each x
  return ()
where
  each [] = do
    return ()
  each (x:xs) = do
    putChar x
    putChar '\n'
    each xs
```

8. Here is a type declaration for simple expressions.

```
data Exp a = Var a | Val Integer | Add (Exp a) (Exp a) | Mult (Exp a) (Exp a) deriving
Show
```

Show how to make this type an instance of `Functor`. When would it be useful to consider `Exp a` as a functor? Think of a good example!

**Comment** Here we need to define `fmap`, and it is a good idea to first see what type this function should have. It should be

```
fmap :: (a -> b) -> Exp a -> Exp b
```

So, from an expression where variables have type `a`, we must construct an expression where variables have type `b`. The solution is therefore

```
instance Functor Exp where
  fmap f (Var x) = Var (f x)
  fmap f (Val n) = Val n
  fmap f (Add e1 e2) = Add (fmap f e1) (fmap f e2)
  fmap f (Mult e1 e2) = Mult (fmap f e1) (fmap f e2)
```

A good example of when it is useful to think of this type as a functor is *renaming variables*. Renaming variables consists of replacing variable occurrences with other variable occurrences. A more general example is *substitution* (of which renaming variables is a special case).

9. Here is a function `readNumber` that takes a `s` of type `String` and returns a value of type `Maybe Int`. It evaluates to `Nothing` if `s` does not represent a valid integer.

```
readNumber :: String -> Maybe Int
readNumber s = case reads s of
  [(n, "")] -> Just n
  _          -> Nothing
```

For example, `readNumber "484000"` evaluates to `Just 484000`, and `readNumber "plip"` evaluates to `Nothing`. Use `do`-notation to create a function `readAndAdd` that takes two arguments of type `String`. If the two arguments represent numbers, the function should return a value of type `Maybe Int`. If one or both arguments are `Nothing`, the result should be `Nothing`. The function should behave as follows.

```
readAndAdd "5" "3"
Just 8
readAndAdd "5" "abc"
Nothing
```

Your solution must *not* use pattern matching or local declarations, only the usual monad constructs.

**Comment** Here is a solution.

```
readAndAdd :: String -> String -> Maybe Int
readAndAdd s1 s2 = do
  x <- readNumber s1
  y <- readNumber s2
  return (x+y)
```

10. Give a definition of the list `naturals` of natural numbers, using *recursion*. The smallest natural number is 1.

**Comment** We have

```
naturals = from 1
  where
    from n = n : (from (n+1))
```