9. Semantics of QL

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a. Arguments and validity in QL

Validity of arguments

Valid?

Everyone is either good or evil.

Not everyone is a villain.

Only villains are evil.

$\therefore$ Some heroes are good.

Validity in QL

- Want to capture validity **in virtue of the meanings of the connectives and the quantifiers**
  (but ignoring the meanings of predicate symbols)

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  (but ignoring the meanings of predicate symbols)
- So we want to ignore any restrictions the predicate symbols place on their **extensions**
- Hence: allow **any** extension in a potential counterexample
- An argument is **QL-valid** if there is **no interpretation** in which the premises are true and the conclusion false

Forms of arguments

Everyone is either good or evil.

Not everyone is a villain.

Only villains are evil.

∴ Some heroes are good.

$(\forall x)(Gx \vee Ex)$

$\sim(\forall x)\, Vx$

$(\forall x)(Ex \supset Vx)$

$\therefore\ (\exists x)(Hx \,\&\, Gx)$

(In)validity of arguments

$$(\forall x)(Gx \vee Ex)$$
$$\sim(\forall x)\, Vx$$
$$(\forall x)(Ex \supset Vx)$$
$$\therefore\ (\exists x)(Hx \,\&\, Gx)$$

Domain: the inner planets (Mecury, Venus, Mars, Earth)

Gx: x is smaller than Earth

Ex: x is inhabited

Vx: x has a moon

Hx: x has rings

9. Semantics of QL

b. Interpretations

Interpretations

- Domain: collection of objects (not empty)

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- Relations between each pair of objects
  - **Extension** of each 2–place predicate symbol:
    all pairs of objects standing in that relation

Extensions

Domain: the inner planets

Gx: x is smaller than Earth

Ex: x is inhabited

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Domain: Mercury, Venus, Earth, Mars

Gx: Mercury, Venus, Mars

Ex: Earth

Vx: Earth, Mars

Hx: —

(In)validity of arguments

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\begin{array}{rl}
& (\forall x)(Gx \vee Ex) \\
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Domain: 1, 2, 3, 4

Gx: 1, 2, 4

Ex: 3

Vx: 3, 4

Hx: —

(In)validity of arguments

$$
\begin{array}{rl}
 & (\forall x)(Gx \lor Ex) \\
 & \sim(\forall x)\, Vx \\
 & (\forall x)(Ex \supset Vx) \\
\therefore & (\exists x)(Hx \,\&\, Gx)
\end{array}
$$

Domain: 1

Gx: 1

Ex: —

Vx: —

Hx: —

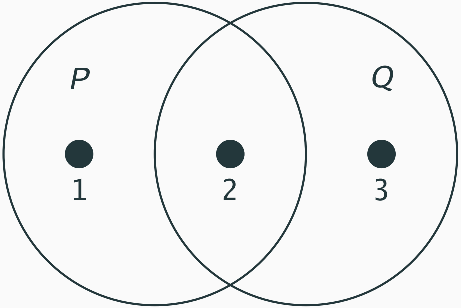
Extensions of predicates

Domain: 1, 2, 3

Px: 1, 2

Qx: 2, 3

Rx: —



$R = \varnothing$

(In)validity of arguments

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Domain: 1, 2

Gx: 1

Ex: 2

Vx: 2

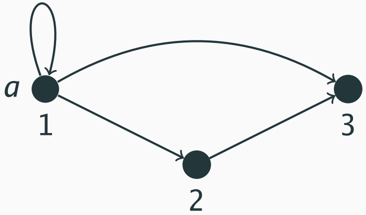
Hx: 2

Extensions of predicates

Domain: 1, 2, 3

a: 1

Axy: $\langle 1, 1\rangle$, $\langle 1, 2\rangle$, $\langle 1, 3\rangle$, $\langle 2, 3\rangle$



9. Semantics of QL

c. Truth of sentences of QL

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- $\mathcal{A} \supset \mathcal{B}$ is true iff $\mathcal{A}$ is false or $\mathcal{B}$ is true

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- $(\forall x)\,\mathcal{A}x$ is true iff $\mathcal{A}x$ is **satisfied** by **every** object in the domain

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- $(\exists x)(\mathcal{A}x \,\&\, \mathcal{B}x)$ is true iff some object satisfies '$\mathcal{A}x \,\&\, \mathcal{B}x$'

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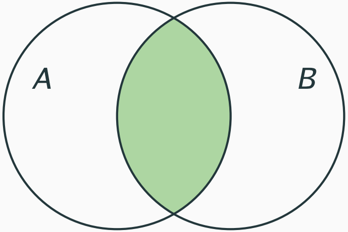
      or
    - o does satisfy $\mathcal{B}x$

Making "Some *A*s are *B*s" true

- $(\exists x)(Ax \mathbin{\&} Bx)$
- Extension of A and B must have something in common. (Filled area must contain at least one object)
- A and B can overlap, be equal, or be contained.
- Same situations make "No *A*s are *B*s" **false**.



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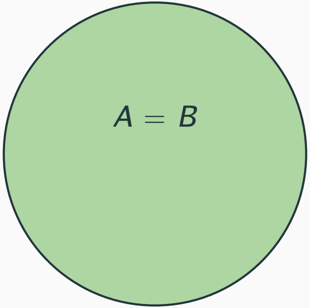

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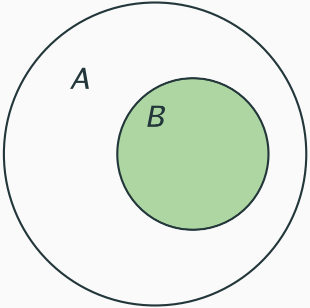

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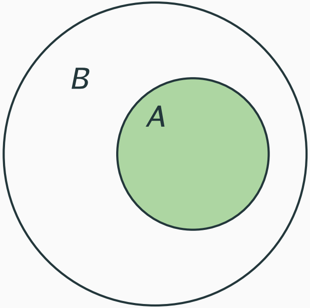

Making “Some *A*s are *B*s” false

- $\sim(\exists x)(Ax \,\&\, Bx)$
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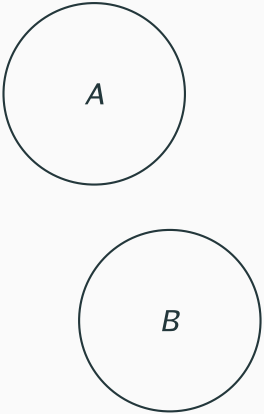

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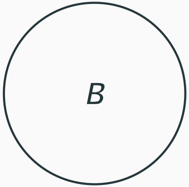

$A = \varnothing$

Making “Some *A*s are *B*s” false

- $\sim(\exists x)\,(Ax \;\&\; Bx)$
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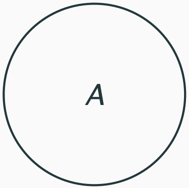


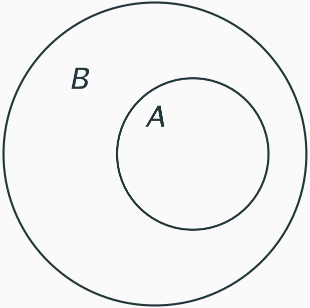
Making “Some As are Bs” false

- $\sim(\exists x)\,(Ax \,\&\, Bx)$
- Extension of A and B must have nothing in common.
- A and B don’t overlap, or one or both empty.
- Same situations make “No As are Bs” **true**.

$$A = B = \varnothing$$

Making “All *A*s are *B*s” true

- $(\forall x)(Ax \supset Bx)$
- Extension of *A* must be contained in extension of *B*.
- Extensions of *A* and *B* can be the same.
- Extension of *A* can be empty.
- Same situations make ...
  - “Only *B*s are *A*s” **true**.
  - “Some *A*s are not *B*s” **false**.



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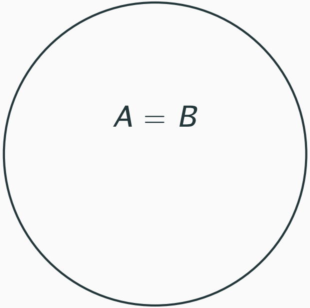

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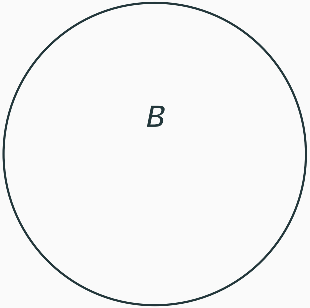

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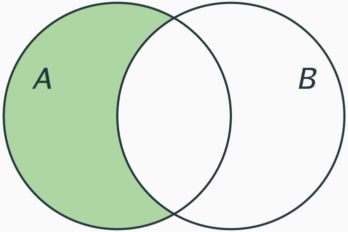

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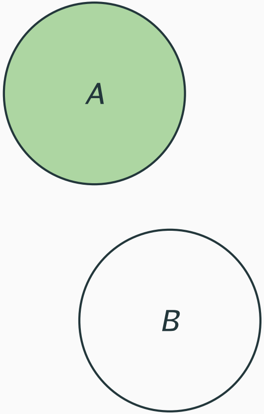

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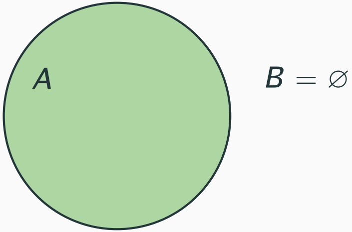

9. Semantics of QL

d. Testing for validity

Arguments involving quantifiers

1. If an action x is morally wrong then A is blameworthy for freely doing x.

–John Skorupski, *Ethical Explorations*, 2000 (link)

Arguments involving quantifiers

1. If an action x is morally wrong then A is blameworthy for freely doing x.

2. If x is rationally optimal (there is no action which A has reason to think there is more reason for A to do), then A is not blameworthy for freely doing x.

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Arguments involving quantifiers

1. If an action x is morally wrong then A is blameworthy for freely doing x.

2. If x is rationally optimal (there is no action which A has reason to think there is more reason for A to do), then A is not blameworthy for freely doing x.

3. Therefore, if x is morally wrong, then x is not rationally optimal. (Principle of moral categoricity.)

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Symbolizing Skorupski

1. If an action x is morally wrong then A is blameworthy for freely doing x.

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

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$(\forall x)(Wx \supset Bx)$

2. If x is rationally optimal, then A is not blameworthy for freely doing x.

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

Ox: x is rationally optimal

$(\forall x)(Wx \supset Bx)$

Symbolizing Skorupski

2. If x is rationally optimal, then A is not blameworthy for freely doing x.

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

Ox: x is rationally optimal

$(\forall x)(Wx \supset Bx)$

$(\forall x)(Ox \supset \sim Bx)$

Symbolizing Skorupski

3. Therefore, if x is morally wrong, then x is not rationally optimal.

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

Ox: x is rationally optimal

$(\forall x)(Wx \supset Bx)$

$(\forall x)(Ox \supset {\sim}Bx)$

Symbolizing Skorupski

3. Therefore, if x is morally wrong, then x is not rationally optimal.

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

Ox: x is rationally optimal

$(\forall x)(Wx \supset Bx)$

$(\forall x)(Ox \supset \sim Bx)$

$\therefore (\forall x)(Wx \supset \sim Ox)$

Symbolizing Skorupski

Domain: actions

Wx: x is morally wrong

Bx: A is blameworthy for freely doing x

Ox: x is rationally optimal

$(\forall x)(Wx \supset Bx)$

$(\forall x)(Ox \supset \sim Bx)$

$\therefore (\forall x)(Wx \supset \sim Ox)$

All Ws are Bs

No Os are Bs (iff No Bs are Os)

$\therefore$ No Ws are Os

Determining validity

- Make conclusion $(\forall x)(Wx \supset {\sim}Ox)$ false.

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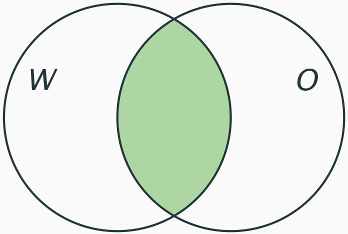
- Make $(\exists x)(Wx \,\&\, Ox)$ true.



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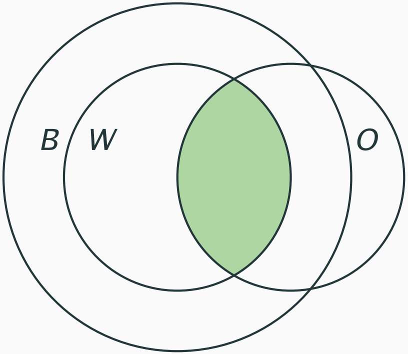

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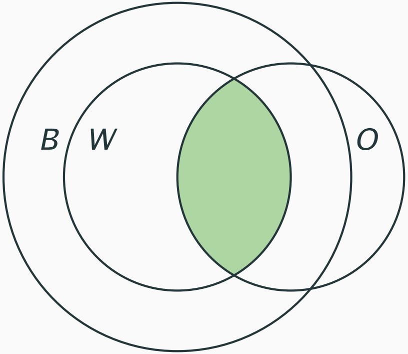
- Make conclusion $(\forall x)(Wx \supset {\sim}Ox)$ false.
- Make $(\exists x)(Wx \,\&\, Ox)$ true.
- Make $(\forall x)(Wx \supset Bx)$ true.
- $(\exists x)(Ox \,\&\, Bx)$ is now forced to be true.



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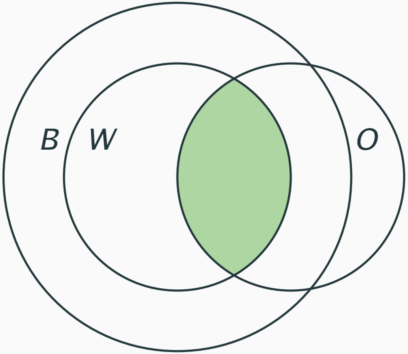
- Make $(\exists x)(Wx \,\&\, Ox)$ true.
- Make $(\forall x)(Wx \supset Bx)$ true.
- $(\exists x)(Ox \,\&\, Bx)$ is now forced to be true.
- So, $(\forall x)(Ox \supset {\sim}Bx)$ is false.



Determining validity

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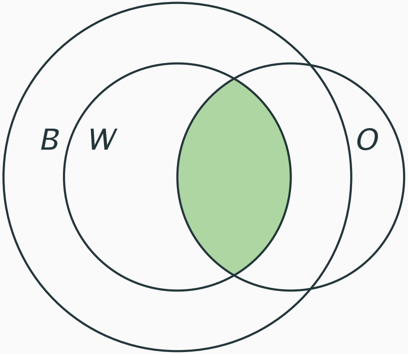
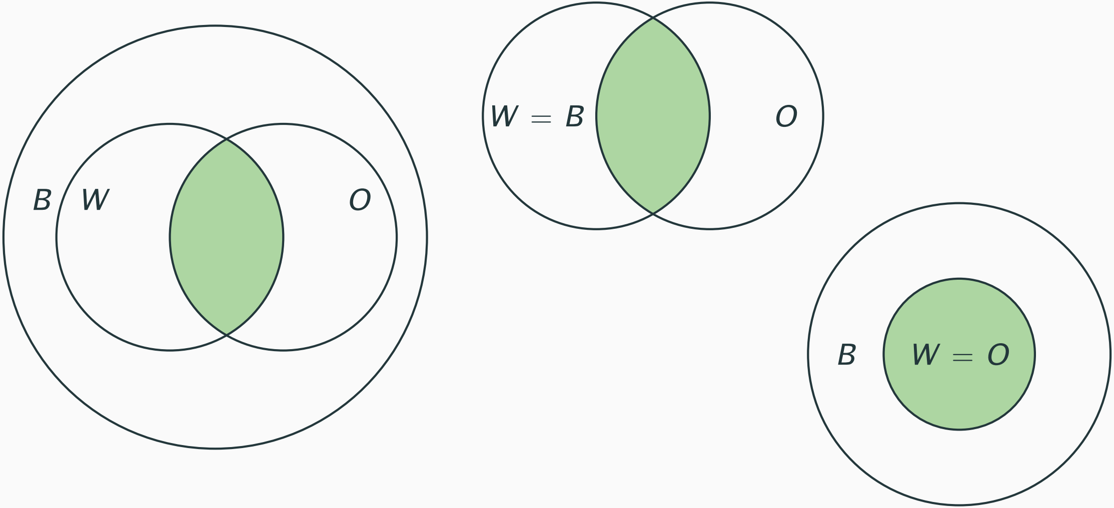
- $(\exists x)(Ox \,\&\, Bx)$ is now forced to be true.
- So, $(\forall x)(Ox \supset \sim Bx)$ is false.
- But those are not the only possibilities!



Other configurations



9. Semantics of QL

e. Semantic notions in QL

Semantics notions in QL

- $\mathcal{P}_1, \ldots, \mathcal{P}_n \vDash \mathcal{Q}$ if no interpretation makes all of $\mathcal{P}_1, \ldots, \mathcal{P}_n$ true and $\mathcal{Q}$ false.

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- $\mathcal{P}$ and $\mathcal{Q}$ are **equivalent in QL** if no interpretation makes one true but the other false.
- $\mathcal{P}_1, \ldots, \mathcal{P}_n$ are **jointly satisfiable in QL** if some interpretation makes all of them true

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Examples

- $(\forall x)(Ax \vee Bx)$ and $(\forall x)\, Ax \vee (\forall x)\, Bx$ are not equivalent.

Test solutions on this week's practice problems!

Examples

- $(\forall x)(Ax \vee Bx)$ and $(\forall x)\, Ax \vee (\forall x)\, Bx$ are not equivalent.
- $(\forall x)(Ax \supset Bx), (\forall x)(Ax \supset {\sim}Bx)$ are jointly satisfiable.

Test solutions on this week's practice problems!

Examples

- $(\forall x)(Ax \lor Bx)$ and $(\forall x)\, Ax \lor (\forall x)\, Bx$ are not equivalent.
- $(\forall x)(Ax \supset Bx), (\forall x)(Ax \supset {\sim}Bx)$ are jointly satisfiable.
- $(\forall x)({\sim}Ax \supset Bx), (\exists x)(Bx \,\&\, Cxb) \nvDash (\exists x)({\sim}Ax \,\&\, Cxb)$.

Test solutions on this week's practice problems!

Examples

- $(\forall x)(Ax \vee Bx)$ and $(\forall x)\, Ax \vee (\forall x)\, Bx$ are not equivalent.
- $(\forall x)(Ax \supset Bx), (\forall x)(Ax \supset {\sim}Bx)$ are jointly satisfiable.
- $(\forall x)({\sim}Ax \supset Bx), (\exists x)(Bx \,\&\, Cxb) \not\models (\exists x)({\sim}Ax \,\&\, Cxb)$.
- $\not\models (\exists x)\, Aax \supset (\exists x)\, Axx$.

Test solutions on this week's practice problems!

9. Semantics of QL

f. Arguing about interpretations

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- That's because there is no way to inspect all possible interpretations.
- But we can show that arguments are valid, by:
  - a formal proof (a future topic!)
  - an informal argument
- The informal argument makes use of the **truth conditions** for sentences of QL.
- Analogous to arguing about valuations in SL.

Example

$$((\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x) \models (\forall x)(\mathcal{A}x \lor \mathcal{B}x)$$

- Suppose an interpretation makes premise $(\forall x)\,\mathcal{A}x \lor (\forall x)\,\mathcal{B}x$ true.

Example

$$((\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x) \vDash (\forall x)(\mathcal{A}x \lor \mathcal{B}x)$$

- Suppose an interpretation makes premise $(\forall x)\,\mathcal{A}x \lor (\forall x)\,\mathcal{B}x$ true.
- By truth conditions for $\lor$, it makes either $(\forall x)\mathcal{A}x$ or $(\forall x)\mathcal{B}x$ true.

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- Suppose it's the first, i.e., $(\forall x)\,\mathcal{A}x$ is true.

Example

$$((\forall x)\mathcal{A}x \vee (\forall x)\mathcal{B}x) \vDash (\forall x)(\mathcal{A}x \vee \mathcal{B}x)$$

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- By truth conditions for $\vee$, it makes either $(\forall x)\mathcal{A}x$ or $(\forall x)\mathcal{B}x$ true.
- Suppose it's the first, i.e., $(\forall x)\,\mathcal{A}x$ is true.
    - By truth conditions for $\forall$, every object in the domain satisfies $\mathcal{A}x$.

Example

$$((\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x) \vDash (\forall x)(\mathcal{A}x \lor \mathcal{B}x)$$

- Suppose an interpretation makes premise $(\forall x)\,\mathcal{A}x \lor (\forall x)\,\mathcal{B}x$ true.
- By truth conditions for $\lor$, it makes either $(\forall x)\mathcal{A}x$ or $(\forall x)\mathcal{B}x$ true.
- Suppose it's the first, i.e., $(\forall x)\,\mathcal{A}x$ is true.
  - By truth conditions for $\forall$, every object in the domain satisfies $\mathcal{A}x$.
  - By the truth conditions for $\lor$, every object satisfies $\mathcal{A}x \lor \mathcal{B}x$

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$$((\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x) \vDash (\forall x)(\mathcal{A}x \lor \mathcal{B}x)$$

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- By truth conditions for $\lor$, it makes either $(\forall x)\mathcal{A}x$ or $(\forall x)\mathcal{B}x$ true.
- Suppose it's the first, i.e., $(\forall x)\,\mathcal{A}x$ is true.
  - By truth conditions for $\forall$, every object in the domain satisfies $\mathcal{A}x$.
  - By the truth conditions for $\lor$, every object satisfies $\mathcal{A}x \lor \mathcal{B}x$
  - So, by the truth conditions for $\forall$, $(\forall x)(\mathcal{A}x \lor \mathcal{B}x)$ is true.

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  - So, by the truth conditions for $\forall$, $(\forall x)(\mathcal{A}x \lor \mathcal{B}x)$ is true.
- Suppose it's the second, i.e., $(\forall x)\mathcal{B}x$ is true: Similarly...

Example

$$((\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x) \vDash (\forall x)(\mathcal{A}x \lor \mathcal{B}x)$$

- Suppose an interpretation makes premise $(\forall x)\mathcal{A}x \lor (\forall x)\mathcal{B}x$ true.
- By truth conditions for $\lor$, it makes either $(\forall x)\mathcal{A}x$ or $(\forall x)\mathcal{B}x$ true.
- Suppose it's the first, i.e., $(\forall x)\mathcal{A}x$ is true.
  - By truth conditions for $\forall$, every object in the domain satisfies $\mathcal{A}x$.
  - By the truth conditions for $\lor$, every object satisfies $\mathcal{A}x \lor \mathcal{B}x$
  - So, by the truth conditions for $\forall$, $(\forall x)(\mathcal{A}x \lor \mathcal{B}x)$ is true.
- Suppose it's the second, i.e., $(\forall x)\mathcal{B}x$ is true: Similarly...
- These are the only possibilities: so any interpretation that makes the premise true must also make the conclusion true.