

Completeness of $\text{FOL}^=$

LOGIC I

Benjamin Brast-McKie

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Basic Lemmas

- L9.1** If $\hat{a}(\alpha) = \hat{c}(\alpha)$ for all free variables α in a wff φ , then $\mathcal{V}_{\mathcal{I}}^{\hat{a}}(\varphi) = \mathcal{V}_{\mathcal{I}}^{\hat{c}}(\varphi)$.
- L11.5** $\mathcal{V}_{\mathcal{I}}^{\hat{a}}(\varphi) = \mathcal{V}_{\mathcal{I}}^{\hat{a}}(\varphi[\beta/\alpha])$ if $\mathfrak{v}_{\mathcal{I}}^{\hat{a}}(\alpha) = \mathfrak{v}_{\mathcal{I}}^{\hat{a}}(\beta)$ and β is free for α in φ .
- L11.6** If \mathcal{M} and \mathcal{M}' have the same domain \mathbb{D} where $\mathcal{I}(\mathcal{F}^n) = \mathcal{I}'(\mathcal{F}^n)$ and $\mathcal{I}(\alpha) = \mathcal{I}'(\alpha)$ for every n -place predicate \mathcal{F}^n and constant α that occurs in a wff φ , then $\mathcal{V}_{\mathcal{I}}^{\hat{a}}(\varphi) = \mathcal{V}_{\mathcal{I}'}^{\hat{a}}(\varphi)$ for any v.a. \hat{a} defined over \mathbb{D} .
- L12.1** If α is a constant and X is an $\text{FOL}^=$ derivation in which the constant β does not occur, then $X[\beta/\alpha]$ is also an $\text{FOL}^=$ derivation.
- L12.3** If $\Lambda \cup \{\varphi\}$ is inconsistent, then $\Lambda \vdash \neg\varphi$.
- L12.4** If $\Lambda \vdash \varphi$ and $\Pi \cup \{\varphi\} \vdash \psi$, then $\Lambda \cup \Pi \vdash \psi$.
- L12.6** If $\Lambda \cup \{\varphi\}$ and $\Lambda \cup \{\neg\varphi\}$ are both inconsistent, then Λ is inconsistent.
- L12.9** If $\Lambda \vdash \varphi$ and $\Lambda \vdash \neg\varphi$, then Λ is inconsistent.
- L12.11** If $\Lambda \vdash \varphi$, then $\Lambda \cup \Pi \vdash \varphi$.

Completeness

T12.1 Every consistent set of $\mathcal{L}^=$ wfss Γ is satisfiable.

Completeness: If $\Gamma \models \varphi$, then $\Gamma \vdash \varphi$.

- Assuming $\Gamma \models \varphi$, we know $\Gamma \cup \{\neg\varphi\}$ is unsatisfiable (check).
- So $\Gamma \cup \{\neg\varphi\}$ is inconsistent by **T12.1**.
- So $\Gamma \vdash \neg\neg\varphi$ by **L12.3**, and so $\Gamma \vdash \varphi$ by DN and **L12.4**.

Assume: Let Γ be a set of $\mathcal{L}^=$ wfss that is consistent in $\text{FOL}^=$.

Saturation

Extension: Let $\mathcal{L}_{\mathbb{N}}^{\equiv}$ include the extra constants \mathbb{N} .

- Let \mathbf{C} be the set of all constants in $\mathcal{L}_{\mathbb{N}}^{\equiv}$.

Free: Let $\varphi(\alpha)$ be a wff of $\mathcal{L}_{\mathbb{N}}^{\equiv}$ with at most one free variable α .

Saturated: A set of wfss Σ is SATURATED in $\mathcal{L}_{\mathbb{N}}^{\equiv}$ just in case for each wff $\varphi(\alpha)$ of $\mathcal{L}_{\mathbb{N}}^{\equiv}$, there is a constant β where $(\exists \alpha \varphi \rightarrow \varphi[\beta/\alpha]) \in \Sigma$.

L12.2 Γ is consistent in $\text{FOL}_{\mathbb{N}}^{\equiv}$.

- We will hence forth take ‘consistent’ to mean ‘consistent in $\text{FOL}_{\mathbb{N}}^{\equiv}$ ’.

Free Enumeration: Let $\varphi_1(\alpha_1), \varphi_2(\alpha_2), \varphi_3(\alpha_3), \dots$ enumerate all wffs of $\mathcal{L}_{\mathbb{N}}^{\equiv}$ with at most one free variable.

Witnesses: $\theta_1 = (\exists \alpha_1 \varphi_1 \rightarrow \varphi_1[n_1/\alpha_1])$ where $n_1 \in \mathbb{N}$ is the first constant not in φ_1 .

$\theta_{k+1} = (\exists \alpha_{k+1} \varphi_{k+1} \rightarrow \varphi_{k+1}[n_{k+1}/\alpha_{k+1}])$ where $n_{k+1} \in \mathbb{N}$ is the first constant not in φ_{k+1} or θ_j for any $j \leq k$.

Saturation: $\Sigma_0 = \Gamma$,
 $\Sigma_{n+1} = \Sigma_n \cup \{\theta_{n+1}\}$, and
 $\Sigma_{\Gamma} = \bigcup_{i \in \mathbb{N}} \Sigma_i$.

L12.5 Σ_{Γ} is consistent and saturated in $\mathcal{L}_{\mathbb{N}}^{\equiv}$.

Base: $\Sigma_0 = \Gamma$ is consistent.

- Immediate from **L12.2**.

Induction: Assume Σ_m is consistent.

- Assume $\Sigma_{m+1} = \Sigma_m \cup \{\theta_{m+1}\}$ is inconsistent for contradiction.
- So $\Sigma_m \vdash \neg \theta_{m+1}$ by **L12.3**, and so $\Sigma_m \vdash \neg(\exists \alpha_{m+1} \varphi_{m+1} \rightarrow \varphi_{m+1}[n_{m+1}/\alpha_{m+1}])$.
- So $\Sigma_m \vdash \exists \alpha_{m+1} \varphi_{m+1}$ and $\Sigma_m \vdash \neg \varphi_{m+1}[n_{m+1}/\alpha_{m+1}]$ by derived PL rules.
- So $\Sigma_m \vdash \forall \alpha_{m+1} \neg \varphi_{m+1}$ by $\forall I$ since n_{m+1} is not in $\forall \alpha_{m+1} \neg \varphi_{m+1}$ or Σ_m .
- So $\Sigma_m \vdash \neg \exists \alpha_{m+1} \varphi_{m+1}$ by $\forall \neg$, and so Σ_m is inconsistent by **L12.9**.
- It follows by *reductio* that Σ_{m+1} is consistent.
- By weak induction, we know that Σ_k is consistent for all $k \in \mathbb{N}$.

Limit: If Σ_{Γ} is inconsistent, then X derives \perp from Σ_{Γ} in $\text{FOL}_{\mathbb{N}}^{\equiv}$.

- Since X is finite, $\Sigma_m \vdash \perp$ for some $m \in \mathbb{N}$ including all premises in X .
- So Σ_m is inconsistent, contradicting the above.
- By *reductio*, Σ_{Γ} is consistent.

Maximization

Maximal: A set of wfss Δ is MAXIMAL in $\mathcal{L}_{\mathbb{N}}^{\bar{=}}$ just in case either $\psi \in \Delta$ or $\neg\psi \in \Delta$ for every wfs ψ in $\mathcal{L}_{\mathbb{N}}^{\bar{=}}$.

Full Enumeration: Let $\psi_0, \psi_1, \psi_2, \dots$ enumerate all wfss in $\mathcal{L}_{\mathbb{N}}^{\bar{=}}$.

Maximization: $\Delta_0 = \Sigma_{\Gamma}$,

$$\Delta_{n+1} = \begin{cases} \Delta_n \cup \{\psi_n\} & \text{if } \Delta_n \cup \{\psi_n\} \text{ is consistent} \\ \Delta_n \cup \{\neg\psi_n\} & \text{otherwise.} \end{cases},$$

$$\Delta_{\Sigma} = \bigcup_{i \in \mathbb{N}} \Delta_n.$$

L12.7 $\Delta = \Delta_{\Sigma_{\Gamma}}$ is maximal in $\mathcal{L}_{\mathbb{N}}^{\bar{=}}$ and consistent.

Base: $\Delta_0 = \Sigma_{\Gamma}$ is consistent by **L12.5**.

Induction: Assume Δ_n is consistent.

- Want to show that Δ_{n+1} is consistent.

Case 1: If $\Delta_n \cup \{\psi_n\}$ is consistent, then $\Delta_{n+1} = \Delta_n \cup \{\psi_n\}$ is consistent.

Case 2: If $\Delta_n \cup \{\psi_n\}$ is inconsistent, then $\Delta_{n+1} = \Delta_n \cup \{\neg\psi_n\}$.

- Assume $\Delta_n \cup \{\neg\psi_n\}$ is inconsistent for contradiction.
- So Δ_n is inconsistent by **L12.6**.
- So Δ_{n+1} is consistent in both cases.

Limit If Δ is inconsistent, then Y derives \perp from Δ in $\text{FOL}_{\mathbb{N}}^{\bar{=}}$.

- Since Y is finite, $\Delta_m \vdash \perp$ for some $m \in \mathbb{N}$ including all premises in Y .
- This contradicts the above, and so Δ is consistent by *reductio*.

Maximal: Let ψ be any wfs of $\mathcal{L}_{\mathbb{N}}^{\bar{=}}$, and so $\psi = \psi_k$ for some $k \in \mathbb{N}$.

- By construction, $\psi_k \in \Delta_{k+1}$ or $\neg\psi_k \in \Delta_{k+1}$.
- Generalizing on ψ shows that Δ is maximal.

L12.8 $\Gamma \subseteq \Sigma_{\Gamma} \subseteq \Delta$ where Δ is saturated.

- Immediate from the definitions.

L12.10 $\varphi \in \Delta$ whenever $\Delta \vdash \varphi$.

- Assuming $\Delta \vdash \varphi$, we know $\Delta \not\vdash \neg\varphi$ by **L12.9**.
- So $\neg\varphi \notin \Delta$ since otherwise $\Delta \vdash \neg\varphi$.
- Thus $\varphi \in \Delta$ by maximality.