11. Multiple quantifiers &

The Identity Predicate

- 1. Multiple quantifiers & The Identity Predicate
- 1.1 Two quantifiers
- 1.2 Multiple Determiners
- 1.3 The identity predicate
- 1.4 Numerical quantification
- 1.5 Both 'both' and 'neither'
- 1.6 'The' Definite Description
- 1.7 Using quantifiers to express properties
- 1.8 Multiple determiners: worked example
- 1.9 Quantifier scope ambiguity
- 1.10 Donkey sentences

a. Two quantifiers

11. Multiple quantifiers &

The Identity Predicate

#### Formulas expressing relations

- ightharpoonup A formula Ax with one free variable expresses a **property**
- ightharpoonup A formula  $\mathcal{B}xy$  with two free variables expresses a **relation**
- ▶  $(\forall x)(\forall y) \mathcal{B}xy$  is a sentence:
- ▶ It's true iff every pair of objects  $\alpha$ ,  $\beta$  stand in the relation expressed by  $\mathcal{B}xy$ .
- ▶  $(\exists x)(\exists y) \mathcal{B}xy$  is a sentence:
- ▶ It's true iff at least one pair of objects  $\alpha$ ,  $\beta$  stand in the relation expressed by  $\mathcal{B}xy$ .

#### Multiple uses of a single quantifier: $\forall$

- ightharpoonup Axy: x admires y.
- $\blacktriangleright$   $(\forall x)(\forall y)$  Axy: for every pair  $\langle \alpha, \beta \rangle$ ,  $\alpha$  admires  $\beta$ .
- ► In other words: everyone admires everyone.
- ▶ Note: "every pair" includes pairs  $\langle \alpha, \alpha \rangle$ , i.e.,
- ▶  $(\forall x)(\forall y) Axy$  is true only if all pairs  $\langle \alpha, \alpha \rangle$  satisfy Axy.
- ► That means, everyone admires themselves, in addition to everyone else.
- ► So:  $(\forall x)(\forall y)$  Axy does **not** symbolize "everyone admires everyone **else**." (To handle that, we'll need identity!)

### Multiple uses of single quantifier: $\exists$

- ▶  $(\exists x)(\exists y) Axy$ : for at least one pair  $\langle \alpha, \beta \rangle$ ,  $\alpha$  admires  $\beta$ .
- ▶ In other words: at least one person admires at least one person.
- ▶ Note: includes pairs  $\langle \alpha, \alpha \rangle$ , i.e.,
- ▶  $(\exists x)(\exists y) \ Axy$  is already true if a single pair  $\langle \alpha, \alpha \rangle$  satisfies Axy.
- ▶ That means, we could just have one person admiring themselves.
- ► So:  $(\exists x)(\exists y)$  Axy does **not** symbolize "someone admires someone **else**." (again, for that, we'll need the identity predicate)

#### Alternating quantifiers

- 1.  $(\forall x)(\exists y) Axy$ Everyone admires someone (possibly themselves)
- 2.  $(\forall y)(\exists x) Axy$ Everyone is admired by someone (not necessarily the same person)
- 3.  $(\exists x)(\forall y) Axy$ Someone admires everyone (including themselves)
- 4.  $(\exists y)(\forall x) Axy$ Someone is admired by everyone (including themselves)

#### Convergence vs. uniform convergence

► A function f is point-wise continuous if

$$(\forall \epsilon)(\forall x)(\forall y)(\exists \delta)(|x-y|<\delta\supset |f(x)-f(y)|<\epsilon)$$

► A function f is uniformly continuous if

$$(\forall \epsilon)(\exists \delta)(\forall x)(\forall y)(|x-y|<\delta\supset |f(x)-f(y)|<\epsilon)$$

The Identity Predicate

11. Multiple quantifiers &

b. Multiple Determiners

## "Determiner phrases" say what?

- ► Determiners: quantifiers and indefinite or definite articles (also possessives and demonstratives)
- ▶ e.g. many, some, a, the, his, their, this, that
- Determiner phrases: combine a determiner with a (possibily modified) noun:
- ► 'all heroes'; 'a cape'
- 'some woman'; 'the donkey'

#### Symbolizing multiple determiners

- ▶ What if your sentence contains more than one determiner phrase?
- ▶ Deal with each determiner separately!
- ► Think of determiner phrase as replaced with name or variable—result has one less determiner.
- ► When you're down to one determiner, apply known methods for single quantifiers.
- ► This results in formulas that express properties or relations, but themselves contain quantifiers.

#### Two separate determiner phrases

- ► All heroes wear a cape
- ► All heroes satisfy "x wears a cape"

$$(\forall x)(Hx\supset "x \text{ wears a cape"})$$

► x wears a cape

$$(\exists y)(Ey \& Rxy)$$

► Together:

$$(\forall x)(Hx\supset (\exists y)(Ey \& Rxy))$$

#### Determiner within determiner phrase

- ► All heroes who wear a cape admire Greta.
- ► All things that satisfy "x is a hero who wears a cape" admire Greta.

$$(\forall x)$$
 ("x is a hero who wears a cape"  $\supset Axg$ )

► x is a hero who wears a cape

$$Hx \& (\exists y) (Ey \& Rxy)$$

► Together:

$$(\forall x)((Hx \& (\exists y)(Ey \& Rxy)) \supset Axg)$$

## "Any" is sometimes existential

► Any (every) cape is worn by a hero:

$$(\forall x)(Ex\supset (\exists y)(Hy \& Ryx))$$

► No hero wears any cape:

$$(\forall x)(Hx \supset \sim (\exists y)(Ey \& Rxy))$$
  
 
$$\sim (\exists x)(Hx \& (\exists y)(Ey \& Rxy))$$

► No hero wears every cape:

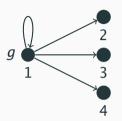
$$(\forall x)(Hx \supset \sim(\forall y)(Ey \supset Rxy))$$
  
 
$$\sim(\exists x)(Hx \& (\forall y)(Ey \supset Rxy))$$

11. Multiple quantifiers &

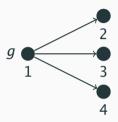
c. The identity predicate

The Identity Predicate

## Greta admires everyone (else)



Greta admires everyone.  $(\forall x) Aqx$ 



Greta admires

everyone else.  $(\forall x)("x \text{ is not Greta"} \supset Agx)$   $(\forall x)(\sim x = g \supset Agx)$ 

#### The identity predicate

- ► A new, special two-place predicate: =
  - Written between arguments, without parentheses.
  - Needs no mention in symbolization key.
  - Always interpreted the same: extension of '=' is all pairs  $\langle \alpha, \alpha \rangle$ .
- ightharpoonup 'a = b' true iff 'a' and 'b' name one and the same object.
- ightharpoonup x = y satisfied by all and only the pairs  $\langle \alpha, \alpha \rangle$ .
- $ightharpoonup \sim x = y$  is satisfied by a pair  $\langle \alpha, \beta \rangle$  iff  $\alpha$  and  $\beta$  are different objects.

## MISTAKES! Ungrammatical expressions with identity

- $ightharpoonup x = \sim y$  is not grammatical.
  - $\sim$  can only go in front of a formula, and y is not one.
- $ightharpoonup \sim (x = y)$  is also not grammatical.
  - (x = y)' is also not a formula.
- ► Carnap will not tolerate this nonsense! Take heed!

#### 'Something else' and 'everything else'

- ► Remember: different variables do NOT entail different objects.
- $\blacktriangleright$   $(\exists x)(\exists y)$  Axy doesn't mean that someone admires someone else.
- ► It just means that someone admires someone (possibly themselves).
- ► To symbolize "someone else" add  $\sim x = y$ :

$$(\exists x)(\exists y)(\sim x=y \& Axy)$$

- $(\forall x)(\forall y)$  Axy says that everyone admires everyone (including themselves).
- ▶ To symbolize "everyone admires everyone else" add  $\sim x = y$ :

$$(\forall x)(\forall y)(\sim x=y\supset Axy)$$

#### 'Something else' and 'everything else'

► The closest quantifier (typically) determines whether you should use & or ⊃:

```
(\forall x)(\exists y)(\sim x=y \& Axy) vs. (\exists x)(\forall y)(\sim x=y\supset Axy) Everyone admires someone else vs. Someone admires everyone else
```

- ► If you have mixed domains, it works the same way:
- ► Recall predicate 'Px': "x is a person"
- Everyone admires someone else:

$$(\forall x)(Px\supset (\exists y)((Py\&\sim x=y)\&Axy))$$

Someone admires everyone else:

$$(\exists x)(Px \& (\forall y)((Py \& \sim x=y) \supset Axy)$$

#### Other than, except

► "Someone other than Greta is a hero":

$$(\exists x)(\sim x = g \& Hx)$$

- "Everyone other than Greta is a hero"; same as:
- "Everyone except Greta is a hero":

$$(\forall x)(\sim x = g \supset Hx)$$

## 'No-one other than' vs. Singular "only"

► "No-one other than Greta is a hero":

$$\sim (\exists x)(Hx \& \sim x = g)$$
$$(\forall x)(Hx \supset x = g)$$

- ► "Only Greta is a hero":
- ► Content: No-one other than Greta is a hero, AND Greta is a hero:

$$(\forall x)(Hx \supset x=g) \& Hg$$
  
 $(\forall x)(Hx \equiv x=g)$ 

#### Uniqueness

Non-unique: "There is at least one hero":

$$(\exists x) Hx$$

- ► Unique: "There is exactly one hero":
  - There's at least one hero, AND
  - There are no others:

$$(\exists x) (Hx \& \sim (\exists y) (\sim y = x \& Hy))$$
  
$$(\exists x) (Hx \& (\forall y) (Hy \supset x = y))$$

• Or more succinctly:  $(\exists x)(\forall y)(Hy \equiv x=y)$ 

## d. Numerical quantification

The Identity Predicate

11. Multiple quantifiers &

#### Numerical Quantification: n-many as 'at least n'

- Cardinal numbers can be determiners:
  - Three heroes wear capes.
- Not always clear if "three heroes" means exactly vs. at least three hero
- ► We'll assume the latter.
  - Do you have two dollars? Yes, I have two dollars.
     (Uncontroversially true even if you have more than \$2)
- ▶ Using QL, we can express the following kinds of sentences:
  - At least n people are . . .
  - Exactly *n* people are ...
  - At most n people are ...
- ▶ i.e. we can count on QL!

#### At least n

► At least 1 hero is inspiring:

$$(\exists x)(Hx \& Ix)$$

At least 2 heroes are inspiring:

$$(\exists x)(\exists y)(\sim x=y \& ((Hx \& Ix) \& (Hy \& Iy)))$$

► At least 3 heroes are inspiring:

$$(\exists x)(\exists y)(\exists z)\Big((\sim x = y \& (\sim y = z \& \sim x = z)) \& \\ ((Hx \& Ix) \& ((Hy \& Iy) \& (Hz \& Iz)))\Big)$$

#### At least n

► There are at least *n* As, i.e. " $(\exists^{\geq n} x) Ax$ ":

$$(\exists x_1) \dots (\exists x_n) \Big( (\sim x_1 = x_2 \& (\sim x_1 = x_3 \& \dots \& (\sim x_1 = x_n \& (\sim x_2 = x_3 \& \dots \& (\sim x_2 = x_n \& \dots \& (\sim x_1 = x_n \& ($$

#### At least n

► Note: must state that **every pair** of variables is different, e.g.,

$$(\exists x_1)(\exists x_2)(\exists x_3)((\sim x_1 = x_2 \& \sim x_2 = x_3) \& (Hx_1 \& (Hx_2 \& Hx_3)))$$

only says "There are at least two heroes"!

- Take extension of Hx to be: 1, 2
- Then 1 can play role of  $x_1$  and  $x_3$ , 2 role of  $x_2$ .
- Both " $\sim 1 = 2$ " and " $\sim 2 = 3$ " are true.
- ► At least *n B*s are *C*s: substitute '*Bx* & *Cx*' for '*Ax*':

$$(\exists^{\geq n} x)(Bx \& Cx)$$

### Exactly one (i.e. Uniqueness; see above)

► There is exactly one hero:

$$(\exists x)(Hx \& \sim (\exists y)(Hy \& \sim x=y))$$

► This is equivalent to:

$$(\exists x)(Hx \& (\forall y)(Hy \supset x=y))$$

► In general: "g has property A uniquely":

$$Ag \& (\forall y)(Ay \supset g=y)$$
 or just:  $(\forall y)(Ay \equiv g=y)$ 

#### Exactly n

▶ There are exactly *n* As, i.e. " $(\exists^{=n}x)$  Ax":  $(\exists x_1) \dots (\exists x_n) ( (\sim x_1 = x_2 \& (\sim x_1 = x_3 \& \dots \& (\sim x_1 = x_n \& \dots ) ) )$  $(\sim X_2 = X_3 \& \dots \& (\sim X_2 = X_n \&$  $\sim X_{n-1} = X_n \ldots$  $(Ax_1 \& (Ax_2 \& ... \& Ax_n)...)) \&$  $(\forall y)(Ay\supset(y=x_1\vee\cdots\vee y=x_n))$ 

Exactly *n B*s are *C*s:

$$(\exists^{=n}x)(Bx \& Cx)$$

## Exactly n

▶ There are exactly *n* As, i.e. " $(\exists^{=n}x)$  Ax":  $(\exists x_1) \dots (\exists x_n) ( (\sim x_1 = x_2 \& (\sim x_1 = x_3 \& \dots \& (\sim x_1 = x_n \& \dots ) ) )$  $(\sim X_2 = X_3 \& \dots \& (\sim X_2 = X_n \&$  $\sim X_{n-1} = X_n ) \dots ) \&$ 

 $(\forall y)(Ay \equiv (y = x_1 \lor \cdots \lor y = x_n))$ 

► Exactly *n B*s are *C*s:

 $(\exists^{=n} x)(Bx \& Cx)$ 

11 d 6

#### At most n

► There are at most n As  $\Leftrightarrow$  There are not at least n+1 As

$$(\exists^{\leq n} x) Ax \Leftrightarrow \sim (\exists^{\geq (n+1)} x) Ax$$

► For instance: There are at most two heroes:

$$\sim (\exists x)(\exists y)(\exists z)((Hx \& (Hy \& Hz)) \& (\sim x = y \& (\sim x = z \& \sim y = z)))$$
  
 $(\forall x)(\forall y)(\forall z)((Hx \& (Hy \& Hz)) \supset (x = y \lor (x = z \lor y = z)))$ 

ightharpoonup  $\sim (\exists^{\geq (n+1)}x) Ax$  is equivalent to:

$$(\forall x_1) \dots (\forall x_{n+1}) ((Ax_1 \& \dots \& Ax_{n+1}) \supset (x_1 = x_2 \lor (x_1 = x_3 \lor \dots \lor (x_1 = x_{n+1} \lor (x_2 = x_3 \lor \dots \lor (x_2 = x_{n+1} \lor (x_n + x_n))))$$

 $X_n = X_{n+1} \ldots )))$ 

# The Identity Predicate

11. Multiple quantifiers &

e. Both 'both' and 'neither'

#### Schematizing 'Both'

"Both heroes inspire": this means that There are exactly 2 heroes, and both inspire:

$$(\exists x)(\exists y)\big(((\sim x = y \& (Hx \& Hy)) \& (\forall z)(Hz \supset (z = x \lor z = y))) \& (Ix \& Iy)\big)$$

- Note: "Both heroes inspire" implies "There are exactly two inspiring heroes", but not vice versa!
- e.g. if there are exactly two inspiring heros and one (or more) not-inspiring hero(s)

#### Schematizing 'Neither'

► "Neither hero inspires": this means that

There are exactly 2 heroes, and neither of them inspires:

$$(\exists x)(\exists y) \Big( ((\sim x = y \& (Hx \& Hy)) \& (\forall z)(Hz \supset (z = x \lor z = y))) \& (\sim Ix \& \sim Iy) \Big)$$

f. 'The' Definite Description

11. Multiple quantifiers &

The Identity Predicate

# Definite descriptions

- ► Definite description: the so-and-so
- Russell's analysis of definite description: to say

"The *A* is B"

is to say:

- ► There is something, which:
  - is *A*,
  - is the **only** A (i.e. the unique thing that is A),
  - is *B*.
- ► In QL:

$$(\exists x)(Ax \& (\forall y)(Ay \supset x=y) \& Bx)$$

or more succinctly:

$$(\exists x)(\forall y)\big((Ay\equiv x=y) \& Bx\big)$$

# Example: 'The author of Waverley is blah'

- Schematize "The author of Waverley is Scottish":
- Use the following symbolization key:
- ► Ax: x is an author; Wxz: x wrote z; Sx: x is Scottish;  $\ell$ : Waverley

 $(\exists x)(Ax \& Wx\ell \& (\forall y)((Ay \& Wy\ell) \supset x=y) \& Sx)$ 

# "The" vs. "exactly one"

- ► Compare:
  - 1. The hero inspires:

$$(\exists x)(Hx \& (\forall y)(Hy \supset x=y) \& Ix)$$

2. There is exactly one inspiring hero:

$$(\exists x)(Hx \& Ix \& (\forall y)((Hy \& Iy) \supset x=y))$$

- ▶ (2) can be true without (1), but not vice versa.
- ► (Namely when there is exactly one inspiring hero, but also a non-inspiring hero.)
- ► So (1) entails (2), but not vice versa.

## Strawson's analysis (presuppositional theories)

- According to Russell, "The hero wears a cape" is **false** if there is no hero, or if there is more than one.
- Consider: "the present King of France is bald."
- P. F. Strawson disagrees with these truth conditions. Rather, we only succeed in making a statement if there is a unique hero (or a unique king of France).
- "There is a unique hero" is not part of what is said by a definite description, but is only presupposed.

# Singular possessive (a definite description)

- ► Singular possessives form noun phrases, e.g., "Joe's cape"
- ► They work like definite descriptions: Joe's cape is the cape Joe owns. E.g.:
  - "Autumn wears Joe's cape" symbolizes the same as:
     "Autumn wears the cape that Joe owns":

$$(\exists x) \Big( ((Ex \& Ojx) \& \\ (\forall y) ((Ey \& Ojy) \supset x = y)) \& \\ Wax \Big)$$

# Singular vs. plural possessive

- ► Compare plural possessives: those are '∀'s':
  - "Autumn wears Joe's capes" symbolizes the same as:

"Autumn wears every cape that Joe owns":

$$(\forall x)((Ex \& Ojx) \supset Wax)$$

► So plural possessives are NOT definite descriptions.

# g. Using quantifiers to express

11. Multiple quantifiers &

The Identity Predicate

properties

# Our symbolization key

Domain:	people alive in 2022 and items of clothing
a:	Autumn
g:	Greta
Px:	$\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$
Lx:	$\underline{}_{x}$ is an item of clothing.
Ex:	$\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$
Rxy:	<sub>x</sub> wears <sub>y</sub>
Hx:	$\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$
Ix:	x inspires
Yxy:	$\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$
Axy:	$\underline{\hspace{1cm}}_{x}$ admires $\underline{\hspace{1cm}}_{y}$
Oxy:	v ownsv

## Expressing properties, revisited

► One-place predicates express properties, e.g.,

Hx expresses property "being a hero"

 Combinations of predicates (with connectives, names) can express derived properties, e.g.,

Axg expresses "x admires Greta"

Hx & Cx expresses "x is a hero who wears a cape"

 Using quantifiers, we can express even more complex properties, e.g.,

 $(\exists y)(Py \& Axy)$  expresses "x admires someone"

# Finding, using properties expressed

- $\blacktriangleright$  If you can say it for Greta, you can say it for x.
  - Greta admires a hero.

 $(\exists v)(Hv \& Aqv)$ x admires a hero.

 $(\exists y)(Hy \& Axy)$ 

- ▶ If you can say it for x, you can say it for Greta.
  - x wears a cape.

 $(\exists y)(Ey \& Rxy)$ 

• Greta wears a cape.  $(\exists v)(Ev \& Rav)$ 

Ex: x is a capeRxy: \_\_\_\_\_\_ wears \_\_\_\_\_\_\_

## Examples

- $\triangleright$  x wears a cape.  $(\exists v)(Ev \& Rxv)$
- $\triangleright$  x is admired by everyone.

$$(\forall y)(Py\supset Ayx)$$

- $\triangleright$  x admires a hero.
  - $(\exists y)(Hy \& Axy)$
- $\triangleright$  x admires only heroes.
  - $(\forall y)(Axy\supset Hy)$
- ightharpoonup x is unclothed (i.e. naked).

$$\sim (\exists y)(Ly \& Rxy)$$

$$(\forall y)(Ly\supset \sim Rxy)$$

Px \_\_\_\_\_\_x is a person Lx \_\_\_\_\_x is an item of clothing Ex \_\_\_\_\_x is a cape Rxy \_\_\_\_\_x wears \_\_\_\_y

## The Identity Predicate

example

\_\_\_\_\_

11. Multiple quantifiers &

h. Multiple determiners: worked

## Mary Astell, 1666-1731



- ► British political philosopher
- ► Some Reflections upon Marriage (1700)
- ► In preface to 3rd ed. 1706 reacts to William Nicholls' claim (in The Duty of Inferiors towards their Superiors, in Five Practical Discourses (London 1701), Discourse IV: The Duty of Wives to their Husbands), that women are naturally inferior to men.

### Astell TL;DR

- What can Nicholls possibly mean by "women are naturally inferior to men"?
- ► It can't be that some woman is inferior to some man, since that's "no great discovery."
- ► After all, surely some men are inferior to some women.
- ► The obviously intended meaning must be: all women are inferior to all men.
- ▶ But that can't be right, for then "the greatest Queen ought not to command but to obey her Footman."
- ► It can't even be just: all women are inferior to some men.
- ► Since "had they been pleased to remember their Oaths of Allegiance and Supremacy, they might have known that *One* Woman is superior to *All* the Men in these Nations."

11.h.2

# Symbolizing Astell

- ► Some woman is superior to every man
- ► Some woman satisfies "x is superior to every man"

$$(\exists x)(Wx \& "x \text{ is superior to every man"})$$

► *x* is superior to every man

$$(\forall y)(My\supset Sxy)$$

► Together:

$$(\exists x)(Wx \& (\forall y)(My \supset Sxy))$$

# Formalizing Astell

► Some woman is superior to some man.

$$(\exists x)(Wx \& (\exists y)(My \& Sxy))$$

► Every woman is superior to every man.

$$(\forall x)(Wx\supset (\forall y)(My\supset Sxy))$$

Every woman is superior to some man.

$$(\forall x)(Wx\supset (\exists y)(My \& Sxy))$$

► Some woman is superior to every man.

$$(\exists x)(Wx \& (\forall y)(My \supset Sxy))$$

## i. Quantifier scope ambiguity

11. Multiple quantifiers &

The Identity Predicate

# More scope ambiguity

- ► "Autumn and Greta admire Isra or Luisa."
- ► Two logically distinct, natural readings:
- 1) Autumn admires Isra or Luisa, and so does Greta.

$$(Aai \lor Aal) \&$$
  
 $(Agi \lor Agl)$ 

2) Autumn and Greta both admire Isra, or they both admire Luisa.

```
(Aai \& Agi) \lor (Aal \& Agl)
```

# Negation and the quantifiers

- "All heroes don't inspire"
  - Denial of "all heroes inspire". Ask: "Do all heroes inspire (Answer: No, it's not the case that all heroes inspire")

$$\sim (\forall x)(Hx \supset Ix)$$
$$(\exists x)(Hx \& \sim Ix)$$

All heroes are not inspiring, i.e.,
 No heroes inspire

$$(\forall x)(Hx\supset\sim Ix)$$
$$\sim (\exists x)(Hx \& Ix)$$

# Multiple quantifiers and ambiguity

- "All heroes wear a cape"
  - "A cape" in the scope of "all heroes", i.e.,
     "For every hero, there is a cape they wear"

```
(\forall x)(Hx \supset (\exists y)(Ey \& Rxy))(\forall x)(\exists y)(Hx \supset (Ey \& Rxy))
```

"All heroes" in scope of "a cape", i.e.,
 "There is a cape which every hero wears"

$$(\exists y)(Ey \& (\forall x)(Hx \supset Rxy))$$
$$(\exists y)(\forall x)(Ey \& (Hx \supset Rxy))$$

► A (probably bad) joke: "Every day, a tourist is mugged on the streets of New York. He's going through a lot of wallets."

#### •

11. Multiple quantifiers &

The Identity Predicate

**Donkey sentences** 

## Happy farmers

"Every farmer who owns a donkey is happy"

- ► Step-by-step symbolization: "All As are Bs"
- ightharpoonup x is a farmer who owns a donkey ...

$$Fx \& (\exists y)(Dy \& Oxy)$$

Every farmer who owns a donkey is happy

$$(\forall x)((Fx \& (\exists y)(Dy \& Oxy))\supset Hx)$$

Notice how 'a donkey' is bound by an existential here

## Unhappy donkeys :(

"Every farmer who owns a donkey beats it"

- ► Step-by-step symbolization: "All As are Bs"
- $\triangleright$  x is a farmer who owns a donkey ...

$$Fx \& (\exists y)(Dy \& Oxy)$$

Every farmer who owns a donkey beats it:

$$(\forall x)((Fx \& (\exists y)(Dy \& Oxy))\supset Bxy)$$

▶ PROBLEM: 'y' is unbound! So this is not a QL sentence. Gasp!

## Save the donkeys: a failed attempt

► This was our problem: a donkey lay beaten and unbound:

$$(\forall x)((Fx \& (\exists y)(Dy \& Oxy))\supset Bxy)$$

► Can we simply extend the scope of the existential?

$$(\forall x)(\exists y)((Fx \& (Dy \& Oxy)) \supset Bxy)$$

- ▶ 'y' is now bound, but alas, this sentence is trivially true:
- Provided at least one thing in our UD is not a donkey, that thing makes the antecedent of the conditional false, making the conditional trivially true, for any x.

In particular, our farmer is not a donkey.

But he still sounds like kind of a jack@\$\$!

# Symbolizing donkey sentences

"Every farmer who owns a donkey beats it"

When is it false that every farmer who owns a donkey beats it? If there's a farmer who owns a donkey but doesn't beat it. Deny that!

$$\sim (\exists x)(Fx \& (\exists y)(Dy \& Oxy \& \sim Bxy))$$

► For every farmer and every donkey they own: the farmer beats the donkey.

$$(\forall x)(\forall y)((Fx \& (Dy \& Oxy)) \supset Bxy)$$

Every farmer beats every donkey they own.

$$(\forall x)(Fx\supset (\forall y)((Dy \& Oxy)\supset Bxy))$$

▶ But what about the case where at least one farmer with a donkey beats only one of his donkeys? #Quitting
11.j.4