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# Bimodal Reference Manual

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*A Logic for Tense and Modality*

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Primary Reference:

Brast-McKie, *The Construction of Possible Worlds*, manuscript, 2025.

## **Abstract**

This reference manual provides the formal specification of the Bimodal TM logic as implemented in the ProofChecker project. TM is a bimodal logic combining S5 metaphysical necessity with linear temporal operators for past and future. The soundness theorem is fully proven; completeness infrastructure is in place but the core lemmas remain axiomatized.

## **Contents**

## 1 Introduction

Bimodal TM is a logic combining S5 metaphysical necessity ( $\Box$ ) with linear temporal operators for past ( $H$ ) and future ( $G$ ). The logic provides a framework for reasoning about necessary truths across time, with applications to metaphysics, action theory, and formal philosophy.

The semantics is based on *task frames*, which extend Kripke frames with temporal structure. A task frame consists of world states connected by a *task relation* indexed by temporal durations. World histories are temporal slices through a task frame, representing the unfolding of a world over time.

### Implementation Status

- **Syntax:** Complete (6 primitive constructors, derived operators)
- **Semantics:** Complete (task frames, world histories, truth conditions)
- **Proof System:** Complete (14 axiom schemata, 7 inference rules)
- **Soundness:** Fully proven (all axioms valid, rules preserve validity)
- **Completeness:** Infrastructure only (Lindenbaum lemma, canonical model axiomatized)

### Source Code

The Lean 4 implementation is in the `Bimodal/` directory:

- `Syntax/Formula.lean` – Formula type and operators
- `Semantics/` – Task frames, world histories, truth conditions
- `ProofSystem/` – Axioms and derivation trees
- `Metalogic/` – Soundness and completeness
- `Theorems/` – Perpetuity principles and modal theorems

## 2 Syntax

### 2.1 Formulas

Formulas are defined inductively with six primitive constructors.

**Definition 2.1** (Formula). The type `Formula` is defined by:

$\varphi, \psi ::= p$	(atomic proposition)
$\mid \perp$	(falsity)
$\mid \varphi \rightarrow \psi$	(implication)
$\mid \Box \varphi$	(necessity)
$\mid H\varphi$	(always past)
$\mid G\varphi$	(always future)

where  $p$  ranges over propositional atoms (type `String`).

## 2.2 Notation

Symbol	Name	Lean	Type
$p, q, r$	Atom	<code>atom s</code>	<code>Formula</code>
$\perp$	Falsity	<code>bot</code>	<code>Formula</code>
$\varphi \rightarrow \psi$	Implication	<code>imp</code>	<code>Formula → Formula → Formula</code>
$\Box \varphi$	Necessity	<code>box</code>	<code>Formula → Formula</code>
$H\varphi$	Always past	<code>all_past</code>	<code>Formula → Formula</code>
$G\varphi$	Always future	<code>all_future</code>	<code>Formula → Formula</code>

## 2.3 Derived Operators

The following operators are defined in terms of the primitives.

**Definition 2.2** (Propositional).

$\neg \varphi := \varphi \rightarrow \perp$	(negation)
$\varphi \wedge \psi := \neg(\varphi \rightarrow \neg \psi)$	(conjunction)
$\varphi \vee \psi := \neg \varphi \rightarrow \psi$	(disjunction)

**Definition 2.3** (Modal).

$\Diamond \varphi := \neg \Box \neg \varphi$	(possibility)
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**Definition 2.4** (Temporal).

$P\varphi := \neg H \neg \varphi$	(sometime past)
$F\varphi := \neg G \neg \varphi$	(sometime future)
$\triangle \varphi := H\varphi \wedge \varphi \wedge G\varphi$	(always)
$\nabla \varphi := P\varphi \vee \varphi \vee F\varphi$	(sometimes)

## 2.4 Temporal Duality

The `swap_temporal` function exchanges past and future operators.

**Definition 2.5** (Temporal Swap).

$$\begin{aligned}
 \text{swap}(p) &= p \\
 \text{swap}(\perp) &= \perp \\
 \text{swap}(\varphi \rightarrow \psi) &= \text{swap}(\varphi) \rightarrow \text{swap}(\psi) \\
 \text{swap}(\Box \varphi) &= \Box \text{swap}(\varphi) \\
 \text{swap}(H\varphi) &= G\text{swap}(\varphi) \\
 \text{swap}(G\varphi) &= H\text{swap}(\varphi)
 \end{aligned}$$

**Theorem 2.6** (Involution).  $\text{swap}(\text{swap}(\varphi)) = \varphi$

## 3 Semantics

### 3.1 Task Frames

Task frames are the fundamental semantic structures for TM.

**Definition 3.1** (Task Frame). A **task frame** over temporal type  $T$  (with totally ordered abelian group structure) is a triple  $\mathbf{F} = (W, T, R)$  where:

- $W$  is a type of world states
- $T$  is the temporal duration type with `LinearOrderedAddCommGroup` instance
- $R : W \times T \times W \rightarrow \text{Prop}$  is the task relation

satisfying two constraints:

1. **Nullity:**  $\forall w. R(w, 0, w)$
2. **Compositionality:**  $\forall w u v x y. R(w, x, u) \wedge R(u, y, v) \rightarrow R(w, x + y, v)$

The nullity constraint ensures zero-duration tasks are identity. Compositionality ensures sequential tasks compose with additive time.

Lean Field	Type	Description
<code>WorldState</code>	Type	World state type
<code>task_rel</code>	$\text{WorldState} \rightarrow T \rightarrow \text{WorldState} \rightarrow \text{Prop}$	Task relation
<code>nullity</code>	$\forall w, \text{task\_rel } w \ 0 \ w$	Identity constraint
<code>compositionality</code>	(see below)	Composition constraint

### 3.2 World Histories

A world history is a temporal slice through a task frame.

**Definition 3.2** (World History). A **world history** in task frame  $\mathbf{F}$  is a structure  $h$  with:

- $\text{dom} : T \rightarrow \text{Prop}$  — a convex temporal domain
- $\text{states} : (t : T) \rightarrow \text{dom}(t) \rightarrow W$  — world state at each time
- $\text{respects\_task}$  — states respect the task relation

**Definition 3.3** (Convex Domain). A domain  $\text{dom} : T \rightarrow \text{Prop}$  is **convex** if:

$$\forall a b c. \text{dom}(a) \wedge \text{dom}(c) \wedge a \leq b \wedge b \leq c \rightarrow \text{dom}(b)$$

**Definition 3.4** (Respects Task). A history respects the task relation if for all  $s, t$  in the domain:

$$R(\text{states}(s), t - s, \text{states}(t))$$

Lean Field	Type	Description
domain	$T \rightarrow \text{Prop}$	Temporal domain
convex	Convex domain	Domain convexity
states	$(t : T) \rightarrow \text{domain } t \rightarrow \text{WorldState}$	State function
respects_task	(see code)	Task relation respect

### 3.3 Task Models

A task model adds valuation to a task frame.

**Definition 3.5** (Task Model). A **task model** over frame  $\mathbf{F}$  is a structure  $\mathbf{M}$  with valuation:

$$V : W \rightarrow \text{String} \rightarrow \text{Prop}$$

where  $V(w, p)$  holds iff atomic proposition  $p$  is true at world state  $w$ .

### 3.4 Truth Conditions

Truth is defined relative to a model, history, and time.

**Definition 3.6** (Truth at a Point). For model  $\mathbf{M}$ , history  $h$ , time  $t \in \text{dom}(h)$ :

$$\begin{aligned}
\mathbf{M}, h, t \models p &\iff V(\text{states}(t), p) \\
\mathbf{M}, h, t \models \perp &\iff \text{False} \\
\mathbf{M}, h, t \models \varphi \rightarrow \psi &\iff \mathbf{M}, h, t \models \varphi \rightarrow \mathbf{M}, h, t \models \psi \\
\mathbf{M}, h, t \models \Box \varphi &\iff \forall h'. \text{states}'(t) = \text{states}(t) \rightarrow \mathbf{M}, h', t \models \varphi \\
\mathbf{M}, h, t \models H\varphi &\iff \forall s \leq t. s \in \text{dom} \rightarrow \mathbf{M}, h, s \models \varphi \\
\mathbf{M}, h, t \models G\varphi &\iff \forall s \geq t. s \in \text{dom} \rightarrow \mathbf{M}, h, s \models \varphi
\end{aligned}$$

The modal operator  $\Box$  quantifies over all histories agreeing at the current world state. The temporal operators  $H$  and  $G$  quantify over past and future times within the history's domain.

### 3.5 Time-Shift

The time-shift operation translates a history by a temporal offset.

**Definition 3.7** (Time-Shift). Given history  $h$  and offset  $d : T$ :

$$\text{time\_shift}(h, d) = h'$$

where:

- $\text{dom}'(t) \iff \text{dom}(t + d)$
- $\text{states}'(t) = \text{states}(t + d)$

**Theorem 3.8** (Time-Shift Preserves Convexity). *If  $h$  has convex domain, so does  $\text{time\_shift}(h, d)$ .*

**Theorem 3.9** (Time-Shift Preserves Task Respect). *If  $h$  respects the task relation, so does  $\text{time\_shift}(h, d)$ .*

The time-shift construction is essential for proving the temporal axioms MF and TF.

### 3.6 Validity

**Definition 3.10** (Validity). A formula  $\varphi$  is **valid** (written  $\vdash \varphi$ ) if:

$$\forall T. \forall \mathbf{F} : \text{TaskFrame } T. \forall \mathbf{M}. \forall h. \forall t \in \text{dom}. \mathbf{M}, h, t \models \varphi$$

where  $T$  ranges over all types with `LinearOrderedAddCommGroup` instance.

**Definition 3.11** (Semantic Consequence).  $\Gamma \models \varphi$  (semantic consequence) holds if in every model where all formulas in  $\Gamma$  are true,  $\varphi$  is also true.

**Definition 3.12** (Satisfiability). A context  $\Gamma$  is **satisfiable** in temporal type  $T$  if there exists a model  $\mathbf{M}$ , history  $h$ , and time  $t$  where all formulas in  $\Gamma$  are true.

**Theorem 3.13** (Valid iff Empty Consequence).  $\vdash \varphi \iff \Box \models \varphi$

**Theorem 3.14** (Monotonicity). If  $\Gamma \subseteq \Delta$  and  $\Gamma \models \varphi$ , then  $\Delta \models \varphi$ .

## 4 Proof Theory

### 4.1 Axiom Schemata

The TM proof system has 14 axiom schemata.

#### 4.1.1 Propositional Axioms

**Axiom 4.1** (K (Distribution)).  $((\varphi \rightarrow (\psi \rightarrow \chi)) \rightarrow ((\varphi \rightarrow \psi) \rightarrow (\varphi \rightarrow \chi)))$

**Axiom 4.2** (S (Weakening)).  $(\varphi \rightarrow (\psi \rightarrow \varphi))$

**Axiom 4.3** (EFQ (Ex Falso)).  $(\perp \rightarrow \varphi)$

**Axiom 4.4** (Peirce).  $((\varphi \rightarrow \psi) \rightarrow \varphi) \rightarrow \varphi$

#### 4.1.2 Modal Axioms (S5)

**Axiom 4.5** (MT (Modal T)).  $(\Box \varphi \rightarrow \varphi)$

**Axiom 4.6** (M4 (Modal 4)).  $(\Box \varphi \rightarrow \Box \Box \varphi)$

**Axiom 4.7** (MB (Modal B)).  $(\varphi \rightarrow \Box \Diamond \varphi)$

**Axiom 4.8** (M5 (Modal 5 Collapse)).  $(\Diamond \Box \varphi \rightarrow \Box \varphi)$

**Axiom 4.9** (MK (Modal K Distribution)).  $(\Box(\varphi \rightarrow \psi) \rightarrow (\Box \varphi \rightarrow \Box \psi))$

#### 4.1.3 Temporal Axioms

**Axiom 4.10** (TK (Temporal K Distribution)).  $(G(\varphi \rightarrow \psi) \rightarrow (G\varphi \rightarrow G\psi))$

**Axiom 4.11** (T4 (Temporal 4)).  $(G\varphi \rightarrow GG\varphi)$

**Axiom 4.12** (TA (Temporal A)).  $(\varphi \rightarrow GP\varphi)$

**Axiom 4.13** (TL (Temporal L)).  $(\Delta \varphi \rightarrow GH\varphi)$

#### 4.1.4 Modal-Temporal Interaction

**Axiom 4.14** (MF (Modal-Future)).  $(\Box\varphi \rightarrow \Box G\varphi)$

**Axiom 4.15** (TF (Temporal-Future)).  $(\Box\varphi \rightarrow G\Box\varphi)$

Axiom	Lean Constructor	Pattern
K	Axiom.prop_k	Distribution
S	Axiom.prop_s	Weakening
EFQ	Axiom.ex_falso	Explosion
Peirce	Axiom.peirce	Classical
MT	Axiom.modal_t	Reflexivity
M4	Axiom.modal_4	Transitivity
MB	Axiom.modal_b	Symmetry
M5	Axiom.modal_5_collapse	S5 collapse
MK	Axiom.modal_k_dist	Modal distribution
TK	Axiom.temp_k_dist	Temporal distribution
T4	Axiom.temp_4	Temporal transitivity
TA	Axiom.temp_a	Connectedness
TL	Axiom.temp_l	Introspection
MF	Axiom.modal_future	Modal-future
TF	Axiom.temp_future	Temporal-modal

## 4.2 Inference Rules

The proof system has 7 inference rules.

**Definition 4.16** (Axiom Rule). If  $\varphi$  matches an axiom schema, then  $\Gamma \vdash \varphi$ .

**Definition 4.17** (Assumption Rule). If  $\varphi \in \Gamma$ , then  $\Gamma \vdash \varphi$ .

**Definition 4.18** (Modus Ponens).

$$\frac{\Gamma \vdash \varphi \rightarrow \psi \quad \Gamma \vdash \varphi}{\Gamma \vdash \psi}$$

**Definition 4.19** (Necessitation).

$$\frac{\vdash \varphi}{\vdash \Box\varphi}$$

Applies only to theorems (empty context).

**Definition 4.20** (Temporal Necessitation).

$$\frac{\vdash \varphi}{\vdash G\varphi}$$

Applies only to theorems (empty context).

**Definition 4.21** (Temporal Duality).

$$\frac{\vdash \varphi}{\vdash \text{swap}(\varphi)}$$

Applies only to theorems (empty context).

**Definition 4.22** (Weakening).

$$\frac{\Gamma \vdash \varphi \quad \Gamma \subseteq \Delta}{\Delta \vdash \varphi}$$

Rule	Lean Constructor	Context Requirement
Axiom	<code>DerivationTree.axiom</code>	Any
Assumption	<code>DerivationTree.assumption</code>	Any
Modus Ponens	<code>DerivationTree.modus_ponens</code>	Any
Necessitation	<code>DerivationTree.necessitation</code>	Empty only
Temp. Necessitation	<code>DerivationTree.temporal_necessitation</code>	Empty only
Temporal Duality	<code>DerivationTree.temporal_duality</code>	Empty only
Weakening	<code>DerivationTree.weakening</code>	Any

### 4.3 Derivation Trees

Derivations are represented as inductive trees.

**Definition 4.23** (Derivation Tree). `DerivationTree`  $\Gamma \varphi$  (written  $\Gamma \vdash \varphi$ ) is an inductive type representing a derivation of  $\varphi$  from context  $\Gamma$ .

**Definition 4.24** (Height). The height of a derivation tree:

- Base cases (axiom, assumption): height 0
- Unary rules: height of subderivation + 1
- Modus ponens: max of both subderivations + 1

The height measure enables well-founded recursion in metalogical proofs.

## 4.4 Notation

Notation	Lean
$\Gamma \vdash \varphi$	<code>DerivationTree <math>\Gamma</math> <math>\varphi</math></code>
$\vdash \varphi$	<code>DerivationTree [] <math>\varphi</math></code>

## 5 Metalogic

### 5.1 Soundness

The soundness theorem establishes that derivability implies semantic validity.

**Theorem 5.1** (Soundness). *If  $\Gamma \vdash \varphi$  then  $\Gamma \models \varphi$ .*

The proof proceeds by induction on the derivation structure:

- **Axioms:** Each of the 14 axiom schemata is proven valid
- **Assumptions:** Assumed formulas are true by hypothesis
- **Modus ponens:** Validity preserved under implication elimination
- **Necessitation:** Valid formulas become necessarily valid
- **Temporal necessitation:** Valid formulas become always-future valid
- **Temporal duality:** Past-future swap preserves validity
- **Weakening:** Adding premises preserves semantic consequence

Axiom Validity	Lean Theorem	Technique
prop_k_valid	Propositional K	Propositional reasoning
prop_s_valid	Propositional S	Propositional reasoning
ex_falso_valid	EFQ	Vacuous implication
peirce_valid	Peirce	Classical case analysis
modal_t_valid	MT	Reflexivity of accessibility
modal_4_valid	M4	Transitivity of accessibility
modal_b_valid	MB	Symmetry of accessibility
modal_5_collapse_valid	M5	S5 equivalence structure
modal_k_dist_valid	MK	Distribution
temp_k_dist_valid	TK	Temporal distribution
temp_4_valid	T4	Transitivity of time
temp_a_valid	TA	Temporal connectedness
temp_l_valid	TL	Always implies recurrence
modal_future_valid	MF	Time-shift invariance
temp_future_valid	TF	Time-shift invariance

The MF and TF axioms use time-shift invariance (via `WorldHistory.time_shift`) to relate truth at different times.

## 5.2 Deduction Theorem

**Theorem 5.2** (Deduction Theorem). *If  $A :: \Gamma \vdash B$  then  $\Gamma \vdash A \rightarrow B$ .*

The proof uses well-founded induction on derivation height, handling each rule:

- **Axiom:** Use S axiom to weaken
- **Assumption:** Identity if same, S axiom if different
- **Modus ponens:** Use K axiom distribution
- **Weakening:** Case analysis on assumption membership
- **Modal/temporal rules:** Do not apply (require empty context)

## 5.3 Consistency

**Definition 5.3** (Consistent). A context  $\Gamma$  is **consistent** if  $\Gamma \not\vdash \perp$ .

**Definition 5.4** (Maximal Consistent). A context  $\Gamma$  is **maximal consistent** if it is consistent and for all  $\varphi \notin \Gamma$ , the context  $\varphi :: \Gamma$  is inconsistent.

## 5.4 Completeness Infrastructure

The completeness proof uses canonical model construction.

**Lemma 5.5** (Lindenbaum). *Every consistent context can be extended to a maximal consistent context.*

**Definition 5.6** (Canonical Frame). The canonical frame has:

- World states: Maximal consistent sets
- Times: Integers ( $\mathbb{Z}$ )
- Task relation: Defined via modal/temporal formula transfer

**Definition 5.7** (Canonical Valuation). An atom  $p$  is true at world  $\Gamma$  iff  $p \in \Gamma$ .

**Lemma 5.8** (Truth Lemma). *In the canonical model,  $\varphi \in \Gamma$  iff  $\varphi$  is true at  $\Gamma$ .*

**Theorem 5.9** (Weak Completeness (Axiomatized)). *If  $\vdash \varphi$  then  $\vdash \varphi$ .*

**Theorem 5.10** (Strong Completeness (Axiomatized)). *If  $\Gamma \models \varphi$  then  $\Gamma \vdash \varphi$ .*

## 5.5 Implementation Status

Component	Status	Lean
Soundness	Fully proven	soundness
Deduction Theorem	Fully proven	deduction_theorem
Lindenbaum Lemma	Axiomatized	lindenbaum
Canonical Frame	Axiomatized	canonical_frame
Truth Lemma	Axiomatized	truth_lemma
Weak Completeness	Axiomatized	weak_completeness
Strong Completeness	Axiomatized	strong_completeness

## 6 Theorems

### 6.1 Perpetuity Principles

The perpetuity principles establish deep connections between modal necessity ( $\Box$ ) and temporal operators ( $\Delta$ ,  $\nabla$ ).

**Theorem 6.1** (P1: Necessity Implies Always).  $\vdash \Box\varphi \rightarrow \Delta\varphi$

**Theorem 6.2** (P2: Sometimes Implies Possible).  $\vdash \nabla\varphi \rightarrow \Diamond\varphi$

**Theorem 6.3** (P3: Necessity of Perpetuity).  $\vdash \Box\varphi \rightarrow \Box\Delta\varphi$

**Theorem 6.4** (P4: Possibility of Occurrence).  $\vdash \Diamond\forall\varphi \rightarrow \Diamond\varphi$

**Theorem 6.5** (P5: Persistent Possibility).  $\vdash \Diamond\forall\varphi \rightarrow \Delta\Diamond\varphi$

**Theorem 6.6** (P6: Occurrent Necessity is Perpetual).  $\vdash \forall\Box\varphi \rightarrow \Box\Delta\varphi$

All six perpetuity principles are fully proven in the Lean implementation.

Principle	Lean Theorem	Key Lemmas
P1	perpetuity_1	MF, TF, MT
P2	perpetuity_2	Contraposition of P1
P3	perpetuity_3	P1, box_mono
P4	perpetuity_4	Contraposition
P5	perpetuity_5	modal_5, temporal K
P6	perpetuity_6	P5, bridge lemmas

## 6.2 Modal S5 Theorems

**Theorem 6.7** (T-Box-to-Diamond).  $\vdash \Box\varphi \rightarrow \Diamond\varphi$

**Theorem 6.8** (Box Distributes Over Disjunction).  $\vdash (\Box\varphi \vee \Box\psi) \rightarrow \Box(\varphi \vee \psi)$

**Theorem 6.9** (Box Preserves Contraposition).  $\vdash \Box(\varphi \rightarrow \psi) \rightarrow \Box(\neg\psi \rightarrow \neg\varphi)$

**Theorem 6.10** (K Distribution for Diamond).  $\vdash \Box(\varphi \rightarrow \psi) \rightarrow (\Diamond\varphi \rightarrow \Diamond\psi)$

**Theorem 6.11** (S5 Collapse).  $\vdash \Diamond\Box\varphi \leftrightarrow \Box\varphi$

**Theorem 6.12** (Box-Conjunction Biconditional).  $\vdash \Box(\varphi \wedge \psi) \leftrightarrow (\Box\varphi \wedge \Box\psi)$

**Theorem 6.13** (Diamond-Disjunction Biconditional).  $\vdash \Diamond(\varphi \vee \psi) \leftrightarrow (\Diamond\varphi \vee \Diamond\psi)$

## 6.3 Modal S4 Properties

The following S4 properties are derived from the TM axiom system.

**Theorem 6.14** (Modal 5).  $\vdash \Diamond\varphi \rightarrow \Box\Diamond\varphi$

**Theorem 6.15** (Diamond 4).  $\vdash \Diamond\Diamond\varphi \rightarrow \Diamond\varphi$

**Theorem 6.16** (Box Monotonicity). *If  $\vdash \varphi \rightarrow \psi$  then  $\vdash \Box\varphi \rightarrow \Box\psi$ .*

**Theorem 6.17** (Diamond Monotonicity). *If  $\vdash \varphi \rightarrow \psi$  then  $\vdash \Diamond\varphi \rightarrow \Diamond\psi$ .*

### 6.4 Propositional Theorems

**Theorem 6.18** (Identity).  $\vdash \varphi \rightarrow \varphi$

**Theorem 6.19** (Double Negation Introduction).  $\vdash \varphi \rightarrow \neg\neg\varphi$

**Theorem 6.20** (Double Negation Elimination).  $\vdash \neg\neg\varphi \rightarrow \varphi$

**Theorem 6.21** (Contraposition). *If  $\vdash \varphi \rightarrow \psi$  then  $\vdash \neg\psi \rightarrow \neg\varphi$ .*

**Theorem 6.22** (De Morgan (Disjunction)).  $\vdash \neg(\varphi \vee \psi) \leftrightarrow (\neg\varphi \wedge \neg\psi)$

**Theorem 6.23** (De Morgan (Conjunction)).  $\vdash \neg(\varphi \wedge \psi) \leftrightarrow (\neg\varphi \vee \neg\psi)$

### 6.5 Combinator Infrastructure

The combinator infrastructure provides Hilbert-style proof tools.

**Theorem 6.24** (B Combinator).  $\vdash (B \rightarrow C) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$

**Theorem 6.25** (Implication Transitivity). *If  $\vdash A \rightarrow B$  and  $\vdash B \rightarrow C$  then  $\vdash A \rightarrow C$ .*

**Theorem 6.26** (Pairing).  $\vdash A \rightarrow (B \rightarrow (A \wedge B))$

**Theorem 6.27** (Classical Merge).  $\vdash (P \rightarrow Q) \rightarrow ((\neg P \rightarrow Q) \rightarrow Q)$

### 6.6 Generalized Necessitation

**Theorem 6.28** (Generalized Modal Necessitation). *If  $\Gamma \vdash \varphi$  then  $\Box\Gamma \vdash \Box\varphi$  where  $\Box\Gamma = [\Box\psi \mid \psi \in \Gamma]$ .*

**Theorem 6.29** (Generalized Temporal Necessitation). *If  $\Gamma \vdash \varphi$  then  $G\Gamma \vdash G\varphi$  where  $G\Gamma = [G\psi \mid \psi \in \Gamma]$ .*

### 6.7 Module Organization

Module	Contents
Perpetuity.lean	P1-P6 principles
Modals5.lean	S5 characteristic theorems
Modals4.lean	S4 properties (modal_5, diamond_4)
Propositional.lean	Classical propositional theorems
Combinators.lean	B, I, S combinators, imp_trans
GeneralizedNecessitation.lean	Context-level necessitation

## 7 Notes

### 7.1 Implementation Status

Component	Status	Notes
Syntax	Complete	6 primitives, derived operators
Semantics	Complete	Task frames, world histories, truth
Proof System	Complete	14 axioms, 7 inference rules
Soundness	Fully Proven	All 14 axioms valid, 7 rules sound
Deduction Theorem	Fully Proven	Well-founded recursion on height
Completeness	Infrastructure Only	Axiomatized (Lindenbaum, truth lemma)
Perpetuity Principles	Fully Proven	P1-P6 all proven

### 7.2 Discrepancy Notes

This section documents differences between the paper “The Construction of Possible Worlds” and the Lean implementation.

#### 7.2.1 Terminology

- The paper uses “perpetuity principles” for P1-P6; the Lean code uses the same terminology.
- The paper’s notation  $\triangle$  and  $\nabla$  for “always” and “sometimes” is preserved in the Lean implementation as `always` and `sometimes`.

#### 7.2.2 Axiom Naming

Paper Name	Lean Name	Notes
MT (Modal T)	<code>Axiom.modal_t</code>	$\Box\varphi \rightarrow \varphi$
M4 (Modal 4)	<code>Axiom.modal_4</code>	$\Box\varphi \rightarrow \Box\Box\varphi$
MB (Modal B)	<code>Axiom.modal_b</code>	$\varphi \rightarrow \Box\Diamond\varphi$
MK	<code>Axiom.modal_k_dist</code>	K distribution
TK	<code>Axiom.temp_k_dist</code>	Temporal K distribution
T4	<code>Axiom.temp_4</code>	Temporal transitivity
TA	<code>Axiom.temp_a</code>	Temporal connectedness
TL	<code>Axiom.temp_l</code>	Temporal introspection
MF	<code>Axiom.modal_future</code>	Modal-future interaction
TF	<code>Axiom.temp_future</code>	Temporal-future interaction

### 7.2.3 M5 Collapse Axiom

The implementation includes an explicit M5 collapse axiom (`Axiom.modal_5_collapse`):

$$\Diamond \Box \varphi \rightarrow \Box \varphi$$

This is derivable from the other S5 axioms (MB + M4) but is included as a primitive for proof convenience in the S5 collapse theorem.

### 7.2.4 Temporal Type Generalization

The paper uses a fixed temporal type  $T = \mathbb{Z}$  (integers). The implementation generalizes to any type  $T$  with `LinearOrderedAddCommGroup` instance, allowing for integers, rationals, reals, or custom bounded time structures.

### 7.2.5 Completeness Status

The paper proves completeness via canonical model construction. The Lean implementation has the completeness infrastructure (Lindenbaum lemma statement, canonical frame types, truth lemma statement) but the core lemmas remain axiomatized. Estimated effort to complete: 70-90 hours of focused development.

## 7.3 Source Files

Directory	Contents
Bimodal/Syntax/	Formula, Context, temporal swap
Bimodal/Semantics/	TaskFrame, WorldHistory, Truth, Validity
Bimodal/ProofSystem/	Axioms, Derivation
Bimodal/Metalogic/	Soundness, DeductionTheorem, Completeness
Bimodal/Theorems/	Perpetuity, ModalS5, Propositional, Combinators