

# Design and Analysis of Algorithms: Exercise 18.2-5

Ben Chaplin

Exercise 18.2-5 from *Introduction to Algorithms* by Cormen, Leiserson, Rivest & Stein (3rd ed)

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## 1 Problem statement

In B-trees, since leaf nodes require no pointers to children, they could conceivably use a different (larger)  $t$  value than internal nodes for the same disk page size. Show how to modify the procedures for creating and inserting into a B-tree to handle this variation

## 2 Improved B-tree

### 2.1 Data structure

Given  $t$ , we can assume our disk page size had room for  $2t$  pointers to children. Let  $p$  be the size of a pointer,  $k$  be the size of a key, and  $c$  be a constant such that  $cp = k$ . Then we can support  $pt = \frac{k}{c}t$  new keys.

### 2.2 B-tree creation

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**Algorithm 1** Original B-tree create

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**Input**

$T$  pointer to the tree

- 1:  $x = \text{ALLOCATE-NODE}()$
  - 2:  $x.\text{leaf} = \text{True}$
  - 3:  $x.n = 0$
  - 4:  $\text{DISK-WRITE}(x)$
  - 5:  $T.\text{root} = x$
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Nothing needs to be changed to support the improved B-tree.

## 2.3 B-tree insertion

(No changes required to B-TREE-SPLIT-CHILD)

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**Algorithm 2** B-TREE-INSERT

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**Input**

$T$  pointer to the tree

$k$  key to insert

```
1:  $r = T.root$ 
2: if  $r.n == 2t - 1$  or  $(r.leaf$  and  $r.n == (2t - 1) + pt)$  then
3:    $s = \text{ALLOCATE-NODE}()$ 
4:    $T.root = s$ 
5:    $s.leaf = \text{False}$ 
6:    $s.n = 0$ 
7:    $s.c_1 = r$ 
8:   B-TREE-SPLIT-CHILD( $s, 1$ )
9:   B-TREE-INSERT-NONFULL( $s, k$ )
10: else
11:   B-TREE-INSERT-NONFULL( $r, k$ )
12: end if
```

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**Algorithm 3** B-TREE-INSERT-NONFULL

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**Input**

$x$  node under which to insert

$k$  key to insert

```
1:  $i = x.n$ 
2: if  $x.leaf$  then
3:   while  $i \geq 1$  and  $k < x.key_i$  do
4:      $x.key_{i+1} = x.key_i$ 
5:      $i = i - 1$ 
6:   end while
7:    $x.key_{i+1} = k$ 
8:    $x.n = x.n + 1$ 
9:   DISK-WRITE( $x$ )
10: else
11:   while  $i \geq 1$  and  $k < x.key_i$  do
12:      $i = i - 1$ 
13:   end while
14:    $i = i + 1$ 
15:   DISK-READ( $x.c_i$ )
16:   if  $x.c_i.n == 2t - 1$  or  $(x.c_i.leaf$  and  $x.c_i.n == (2t - 1) + pt)$  then
17:     B-TREE-SPLIT-CHILD( $x, i$ )
18:     if  $k > x.key_i$  then
19:        $i = i + 1$ 
20:     end if
21:   end if
22:   B-TREE-INSERT-NONFULL( $x.c_i, k$ )
23: end if
```

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