

On the Solution of Systems of Equations by the Epsilon Algorithm of Wynn

By E. Gekeler

Abstract. The ϵ -algorithm has been proposed by Wynn on a number of occasions as a convergence acceleration device for vector sequences; however, little is known concerning its effect upon systems of equations. In this paper, we prove that the algorithm applied to the Picard sequence $\mathbf{x}_{i+1} = F(\mathbf{x}_i)$ of an analytic function $F: \mathbb{R}^n \supset D \rightarrow \mathbb{R}^n$ provides a quadratically convergent iterative method; furthermore, no differentiation of F is needed. Some examples illustrate the numerical performance of this method and show that convergence can be obtained even when F is not contractive near the fixed point. A modification of the method is discussed and illustrated.

1. Introduction. The ϵ -algorithm is a nonlinear method for accelerating the convergence of sequences; in its simplest form, it is identical with the δ^2 transformation of Aitken [1]. The determinantal formulae upon which it is based were given by Jacobi [6], Schmidt [11], and Shanks [12]; Wynn [13] developed it and examined it thoroughly in connection with various sequences and series [14]–[17]. The ϵ -algorithm provides higher (integer) order methods for the computation of a fixed point of an analytic function $f: \mathbb{C} \supset D \rightarrow \mathbb{C}$ [4]. Using the generalized matrix inverse of Moore [8] and Penrose [9], the method has recently been applied to sequences of matrices and vectors as they arise, for example, in the solution of linear systems of equations [5], [7], [10], [18], [21], [22], [23]. Wynn points out that the algorithm also provides good results in the numerical solution of nonlinear systems [18], [19], [21], [22]. But, until now, nothing is known concerning convergence. In this paper, we examine the behaviour of the ϵ -algorithm when applied to the Picard sequence of an analytic function $F: \mathbb{R}^n \supset D \rightarrow \mathbb{R}^n$ with fixed point \mathbf{z} . With the help of a theorem of McLeod [7], we show that the algorithm, used in a manner similar to Steffensen's method, is a quadratically convergent iterative method for the computation of \mathbf{z} (compare also Brezinski [2]*). Because of the complicated recursive relationships, the convergence considered is of local nature, and Landau symbols are used in the proof. A short discussion of numerical properties of the method follows at the end of the paper.

We use certain standard notations: $i \in \mathbb{N}$ means that i is a nonnegative integer; lower (upper) case bold face letters denote vectors (matrices); $\|\mathbf{x}\|$ is the Euclidean norm $(\mathbf{x}^* \mathbf{x})^{1/2}$ of the n -dimensional column vector $\mathbf{x} \in \mathbb{C}^n$; $O(\|\mathbf{x}\|^i)$ denotes a vector-valued function of the vector \mathbf{x} whose norm remains bounded as $\|\mathbf{x}\| \rightarrow 0$ after division by $\|\mathbf{x}\|^i$; $O\{\|\mathbf{x}\|^i\}$ denotes a real valued function with the same properties.

We also make use of the concept of an analytic function of a vector and of a vector-valued Taylor series. Let D be an open subset of \mathbb{R}^n , then $F: \mathbb{R}^n \supset D \rightarrow \mathbb{R}^n$ is called

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analytic if, for every point $\mathbf{a} \in D$, there is an open polycylinder $P = \{\mathbf{x} \in \mathbb{R}^n, |x_i - a_i| < r_i, 0 < r_i, 1 \leq i \leq n\} \subset D$, such that in P , $F(\mathbf{x})$ is equal to the sum of an absolutely summable power series in the n variables $x_i - a_i$ ($1 \leq i \leq n$). An analytic function is indefinitely differentiable, and, if the segment joining \mathbf{x} and $\mathbf{x} + \mathbf{y}$ is in D , we have, for $r \in \mathbb{N}$,

$$F(\mathbf{x} + \mathbf{y}) = F(\mathbf{x}) + \sum_{k=1}^{r-1} \frac{1}{k!} F^{(k)}(\mathbf{x}) \cdot \mathbf{y}^{(k)} + \left(\int_0^1 \frac{(1-t)^{r-1}}{(r-1)!} F^{(r)}(\mathbf{x} + t\mathbf{y}) dt \right) \cdot \mathbf{y}^{(r)},$$

where $\mathbf{y}^{(k)}$ stands for $(\mathbf{y}, \mathbf{y}, \dots, \mathbf{y})$ (k times). For further details, we refer to the famous book of Dieudonné [3].

2. Picard Sequences. We consider some iterative schemes for determining a fixed point \mathbf{z} of the equation $\mathbf{x} = F(\mathbf{x})$. If \mathbf{s}_p ($p \in \mathbb{N}$, $0 \leq p$) is near \mathbf{z} , we have, using a Taylor expansion for $F(\mathbf{z})$,

$$(1) \quad \mathbf{z} = F(\mathbf{s}_p) + F'(\mathbf{s}_p)(\mathbf{z} - \mathbf{s}_p) + O(\|\mathbf{z} - \mathbf{s}_p\|^2).$$

Thus, when using the simple iteration scheme

$$(2) \quad \mathbf{s}_{p+1} = F(\mathbf{s}_p) \quad (0 \leq p),$$

we have

$$\mathbf{z} - \mathbf{s}_{p+1} = F'(\mathbf{s}_p)(\mathbf{z} - \mathbf{s}_p) + O(\|\mathbf{z} - \mathbf{s}_p\|^2).$$

Hence, the simple scheme (2) is, in general, at best linearly convergent; whether it converges or not depends upon the magnitudes of the eigenvalues of the Jacobian matrices $F'(\mathbf{s}_p)$ ($0 \leq p$) in the neighbourhood of \mathbf{z} . We can, however, devise a quadratically convergent scheme based upon the solution of the linear system

$$\hat{\mathbf{s}}_{p+1} = F(\hat{\mathbf{s}}_p) + F'(\hat{\mathbf{s}}_p)(\hat{\mathbf{s}}_{p+1} - \hat{\mathbf{s}}_p) \quad (0 \leq p)$$

or

$$(3) \quad (\mathbf{I} - F'(\hat{\mathbf{s}}_p))\hat{\mathbf{s}}_{p+1} = F(\hat{\mathbf{s}}_p) - F'(\hat{\mathbf{s}}_p)\hat{\mathbf{s}}_p \quad (0 \leq p)$$

for $\hat{\mathbf{s}}_{p+1}$. For, replacing \mathbf{s}_p in formula (1) by $\hat{\mathbf{s}}_p$, we now have

$$\mathbf{z} - \hat{\mathbf{s}}_{p+1} = F'(\hat{\mathbf{s}}_p)(\mathbf{z} - \hat{\mathbf{s}}_{p+1}) + O(\|\mathbf{z} - \hat{\mathbf{s}}_p\|^2) \quad (0 \leq p),$$

i.e.,

$$(\mathbf{I} - F'(\hat{\mathbf{s}}_p))(\mathbf{z} - \hat{\mathbf{s}}_{p+1}) = O(\|\mathbf{z} - \hat{\mathbf{s}}_p\|^2) \quad (0 \leq p)$$

or, again subject to certain assumptions concerning the eigenvalues of $F'(\mathbf{x})$ in the neighbourhood of \mathbf{z} ,

$$\mathbf{z} - \hat{\mathbf{s}}_{p+1} = O(\|\mathbf{z} - \hat{\mathbf{s}}_p\|^2) \quad (0 \leq p).$$

The second scheme, although yielding quadratic convergence, involves evaluation of a Jacobian matrix and the solution of a linear system at each stage. However, by use of the ϵ -algorithm one can, as we shall show, obtain quadratic convergence without

the computation of the derivatives occurring in the Jacobian matrix, and without the solution of a linear system.

3. The Algorithm. The ϵ -algorithm [13], [22] is a computational procedure in which successive columns of an array $(\epsilon_q^{(p)})_{0 \leq p, 0 \leq q}$ with row index p are obtained by use of the formula

$$(4) \quad \epsilon_{q+1}^{(p)} = \epsilon_{q-1}^{(p+1)} + (\epsilon_q^{(p+1)} - \epsilon_q^{(p)})^{-1} \quad (0 \leq p, 0 \leq q),$$

starting from the initial conditions

$$(5) \quad \epsilon_{-1}^{(p)} = 0, \quad \epsilon_0^{(p)} = s_p \quad (0 \leq p).$$

If the inverse of a nonzero vector $\mathbf{x} \in \mathbb{C}^n$ is defined, by [8], [9],

$$(6) \quad \mathbf{x}^{-1} = (\mathbf{x}^* \mathbf{x})^{-1} \bar{\mathbf{x}},$$

then we can apply the algorithm to sequences $\{s_p\}_{0 \leq p}$ of vectors and have the fundamental theorem [7], [23] which we need later:

THEOREM 1. *Let $\{s_p\}_{0 \leq p}$ be a sequence of vectors with complex coefficients which satisfy the irreducible linear recursion*

$$(7) \quad \sum_{r=0}^m c_r s_{p+r} = \left(\sum_{r=0}^m c_r \right) s \quad (0 \leq p),$$

where s is fixed and

$$(8) \quad \sum_{r=0}^m c_r \neq 0, \quad c_r \in \mathbb{R}.$$

If then the elements of the array $(\epsilon_q^{(p)})$ are determined by using (4), (5), and (6), and if all $\epsilon_q^{(p)}$ with $p + q \leq 2m$ exist, then

$$\epsilon_{2m}^{(0)} = s.$$

Following a conjecture of Wynn [24] and Greville [5], Theorem 1 remains true if relations (7), (8) hold for complex scalars only, but this has not yet been proved. In conclusion, we get

COROLLARY. *Let \mathbf{z} be the unique solution of the linear system $\mathbf{x} = \mathbf{A}\mathbf{x} + \mathbf{c}$ with real coefficients and let m be the degree of the minimal polynomial of the matrix \mathbf{A} for $\mathbf{y} = \mathbf{x}_0 - \mathbf{z}$. If the ϵ -algorithm is applied to the Picard sequence $\{\mathbf{x}_p; \mathbf{x}_{p+1} = \mathbf{A}\mathbf{x}_p + \mathbf{c}\}_{0 \leq p}$ and if all $\epsilon_q^{(p)}$ with $p + q \leq 2m$ exist, then*

$$\epsilon_{2m}^{(0)} = \mathbf{z}.$$

Proof. Let $p(x) = \sum_{r=0}^m a_r x^r$ be the minimal polynomial of \mathbf{A} for \mathbf{y} , then

$$\sum_{r=0}^m a_r \mathbf{x}_{p+r} = \left(\sum_{r=0}^m a_r \right) \mathbf{z} + \left(\sum_{r=0}^m a_r \mathbf{A}^{p+r} \right) \mathbf{y} = \left(\sum_{r=0}^m a_r \right) \mathbf{z},$$

because $\mathbf{x}_p = \mathbf{z} + \mathbf{A}^p \mathbf{y}$ holds. By assumption, we have $\sum_{r=0}^m a_r \neq 0$, since 1 is not eigenvalue of \mathbf{A} (the equation $\mathbf{x} = \mathbf{A}\mathbf{x} + \mathbf{c}$ has a unique solution), and the Corollary results from Theorem 1.

4. The Application of the Epsilon Algorithm to Picard Sequences. The general strategy adopted in deriving our main result is this: we first consider the behaviour of the vectors $\tilde{\epsilon}_q^{(p)}$ ($p + q \leq 2n$) derived by means of the ϵ -algorithm from the sequence $\tilde{s}_p = z + A^p y$ ($0 \leq p$), where $y, z \in \mathbb{R}^n$ and A is a real $n \times n$ matrix, for small values of $\|y\|$ (we know from the above Corollary that, subject to certain conditions, $\tilde{\epsilon}_{2n}^{(0)} = z$). We then consider the behaviour of corresponding vectors derived from the sequence $s_p = \tilde{s}_p + \delta_p$, where $\delta_p = O(\|y\|^2)$ ($0 \leq p$). Finally, we use these results with $A = F'(z)$ and

$$s_{p+1} = F(s_p) = z + F'(z)(s_p - z) + O(\|s_p - z\|^2) \quad (0 \leq p)$$

to examine the behaviour of the vectors $\epsilon_q^{(p)}$ produced from this iterative scheme when s_0 is near a fixed point z and, in particular, to show that repeated use of the vector $\epsilon_{2n}^{(0)}$ in place of s_0 results in a quadratically convergent process for determining the fixed point in question. *In the sequel, let $Q_m(A) \subset \mathbb{R}^n$ be the set of vectors x for which m is the degree of the minimal polynomial of A .*

LEMMA 1. *For a given z , let $\tilde{\epsilon}_q^{(p)}$ be the vectors obtained by means of the ϵ -algorithm from the sequence $\{\tilde{s}_p; \tilde{s}_p = z + A^p y\}_{0 \leq p}$. If there is a neighbourhood U of 0 such that all $\tilde{\epsilon}_q^{(p)}$ with $p + q \leq 2m$ exist for all $y \in U \cap Q_m(A)$, then*

$$\begin{aligned} \tilde{\epsilon}_q^{(p)} &= z + O(\|y\|), & q \text{ even,} \\ \tilde{\epsilon}_q^{(p)} &= O(\|y\|^{-1}), & q \text{ odd,} \end{aligned}$$

for $y \in Q_m(A)$ and $p + q \leq 2m$.

Proof. Let $m > 0$, $p \leq 2m - q$, and $\Delta_p \tilde{\epsilon}_q^{(p)} = \tilde{\epsilon}_q^{(p+1)} - \tilde{\epsilon}_q^{(p)}$. For $q = 1$, we get $\Delta_p \tilde{\epsilon}_0^{(p)} = A^p(A - I)y = B_p y$, and $B_p y \neq 0$ for $y \in Q_m(A)$, by assumption. Hence,

$$\begin{aligned} \|\tilde{\epsilon}_1^{(p)}\| &= \|(y^* B_p^* B_p y)^{-1} B_p y\| \\ &= \frac{1}{\|y\|} \frac{y^* y}{y^* B_p^* B_p y} \frac{1}{\|y\|} \|B_p y\| \leq \frac{1}{\|y\|} \frac{\|B_p\|}{\lambda_{\min}}, \end{aligned}$$

where $0 < \lambda_{\min}$ is the smallest eigenvalue of $B_p^* B_p$. Let now $k \in \mathbb{N}$, $k < m$, $y \in Q_m(A)$, and let the statement be true for all $q \leq 2k$. By assumption, we have $\Delta_p \tilde{\epsilon}_{2k}^{(p)} = O(\|y\|) \neq 0$, thus

$$\begin{aligned} (\Delta_p \tilde{\epsilon}_{2k}^{(p)})^* (\Delta_p \tilde{\epsilon}_{2k}^{(p)}) &= O(\|y\|^2), \\ \tilde{\epsilon}_{2k+1}^{(p)} &= \tilde{\epsilon}_{2k-1}^{(p+1)} + [(\Delta_p \tilde{\epsilon}_{2k}^{(p)})^* (\Delta_p \tilde{\epsilon}_{2k}^{(p)})]^{-1} \Delta_p \tilde{\epsilon}_{2k}^{(p)} \\ &= O(\|y\|^{-1}) + O(\|y\|^{-2}) O(\|y\|) = O(\|y\|^{-1}). \end{aligned}$$

$\Delta_p \tilde{\epsilon}_{2k+1}^{(p)} \neq 0$, since, by assumption, all $\tilde{\epsilon}_q^{(p)}$ which contribute to $\tilde{\epsilon}_{2m}^{(0)}$ exist. Therefore,

$$(\Delta_p \tilde{\epsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\epsilon}_{2k+1}^{(p)}) = O(\|y\|^{-2}),$$

and

$$\begin{aligned} \tilde{\epsilon}_{2k+2}^{(p)} &= \tilde{\epsilon}_{2k}^{(p+1)} + [(\Delta_p \tilde{\epsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\epsilon}_{2k+1}^{(p)})]^{-1} \Delta_p \tilde{\epsilon}_{2k+1}^{(p)} \\ &= z + O(\|y\|) + O(\|y\|^2) O(\|y\|^{-1}) = z + O(\|y\|), \end{aligned}$$

and the assertion of the lemma follows by induction.

LEMMA 2. *Let $\{\delta_p\}_{0 \leq p}$ be a sequence of analytic functions $\delta_p(y) = O(\|y\|^2)$. For a given z , let $\epsilon_q^{(p)}$ be the vectors obtained by means of the ϵ -algorithm from the sequence*

$\{s_p; s_p = z + A^p y + \delta_p(y)\}_{0 \leq p}$. If there is a neighbourhood U of 0 such that all $\varepsilon_q^{(p)}$, $\tilde{\varepsilon}_q^{(p)}$ with $p + q \leq 2m$ exist for all $y \in U \cap Q_m(A)$, then

$$\varepsilon_q^{(p)} = \tilde{\varepsilon}_q^{(p)} + O(\|y\|^2), \quad q \text{ even},$$

$$\varepsilon_q^{(p)} = \tilde{\varepsilon}_q^{(p)} + O(1), \quad q \text{ odd},$$

for $y \in Q_m(A)$ and $p + q \leq 2m$.

Proof. Let $m > 0$ and $p \leq 2m - q$. For $q = 1$, we have $\Delta_p \varepsilon_0^{(p)} = \Delta_p \tilde{\varepsilon}_0^{(p)} + O(\|y\|^3) \neq 0$ and $\Delta_p \tilde{\varepsilon}_0^{(p)} \neq 0$ for $y \in Q_m(A)$, by assumption. Then

$$\begin{aligned} (\Delta_p \varepsilon_0^{(p)})^* (\Delta_p \varepsilon_0^{(p)}) &= (\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)}) + 2(\Delta_p \tilde{\varepsilon}_0^{(p)})^* O(\|y\|^2) + O\{\|y\|^4\} \\ &= (\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)}) \left[1 + 2 \frac{(\Delta_p \tilde{\varepsilon}_0^{(p)})^* O(\|y\|^2)}{(\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)})} + \frac{O\{\|y\|^4\}}{(\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)})} \right]. \end{aligned}$$

$\Delta_p \varepsilon_0^{(p)} = O(\|y\|)$ and hence,

$$(\Delta_p \varepsilon_0^{(p)})^* (\Delta_p \varepsilon_0^{(p)}) = (\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)}) (1 + O\{\|y\|\}).$$

Since $\Delta_p \varepsilon_0^{(p)}$ is an analytic function, we get

$$[(\Delta_p \varepsilon_0^{(p)})^* (\Delta_p \varepsilon_0^{(p)})]^{-1} = [(\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)})]^{-1} [1 + O\{\|y\|\}]$$

and

$$\begin{aligned} \varepsilon_1^{(p)} &= \tilde{\varepsilon}_1^{(p)} + [(\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)})]^{-1} O\{\|y\|\} \Delta_p \tilde{\varepsilon}_0^{(p)} \\ &\quad + [(\Delta_p \tilde{\varepsilon}_0^{(p)})^* (\Delta_p \tilde{\varepsilon}_0^{(p)})]^{-1} [1 + O\{\|y\|\}] O(\|y\|^2) \\ &= \tilde{\varepsilon}_1^{(p)} + O(1). \end{aligned}$$

Let now $k \in \mathbb{N}$, $k < m$, $y \in Q_m(A)$, and let the statement be true for all $q \leq 2k$. By assumption, we have $\Delta_p \varepsilon_{2k}^{(p)} = \Delta_p \tilde{\varepsilon}_{2k}^{(p)} + O(\|y\|^2) \neq 0$ and $\Delta_p \tilde{\varepsilon}_{2k}^{(p)} \neq 0$. According to the proof for $q = 1$, we get, by use of Lemma 1,

$$[(\Delta_p \varepsilon_{2k}^{(p)})^* (\Delta_p \varepsilon_{2k}^{(p)})]^{-1} \Delta_p \varepsilon_{2k}^{(p)} = [(\Delta_p \tilde{\varepsilon}_{2k}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k}^{(p)})]^{-1} \Delta_p \tilde{\varepsilon}_{2k}^{(p)} + O(1)$$

and hence,

$$\varepsilon_{2k+1}^{(p)} = \tilde{\varepsilon}_{2k+1}^{(p)} + O(1).$$

$\Delta_p \varepsilon_{2k+1}^{(p)} = \Delta_p \tilde{\varepsilon}_{2k+1}^{(p)} + O(1)$ and $\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)}$ are equally supposed to be different from zero and, therefore, we get, by use of Lemma 1,

$$\begin{aligned} &(\Delta_p \varepsilon_{2k+1}^{(p)})^* (\Delta_p \varepsilon_{2k+1}^{(p)}) \\ &= (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)}) \left[1 + 2 \frac{(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* O(1)}{(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})} + \frac{O\{1\}}{(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})} \right] \\ &= (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)}) (1 + O\{\|y\|\}). \\ \varepsilon_{2k+2}^{(p)} &= \tilde{\varepsilon}_{2k+2}^{(p)} + O(\|y\|^2) \\ &\quad + [(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})]^{-1} [1 + O\{\|y\|\}] [\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)} + O(1)] \\ &= \tilde{\varepsilon}_{2k+2}^{(p)} + O(\|y\|^2) + [(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})]^{-1} O\{\|y\|\} \Delta_p \tilde{\varepsilon}_{2k+1}^{(p)} \\ &\quad + [(\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})^* (\Delta_p \tilde{\varepsilon}_{2k+1}^{(p)})]^{-1} [1 + O\{\|y\|\}] O(1) \\ &= \tilde{\varepsilon}_{2k+2}^{(p)} + O(\|y\|^2). \end{aligned}$$

In conclusion, we have the following result:

THEOREM 2. *Let $F: \mathbb{R}^n \supset D \rightarrow \mathbb{R}^n$ be an analytic function with fixed point $z \in \overset{\circ}{D}$ and let $Q_m(F'(z)) \subset \mathbb{R}^n$ be the set of vectors x for which m is the degree of the minimal polynomial of $F'(z)$. Further, let $\epsilon_q^{(p)}$ and $\tilde{\epsilon}_q^{(p)}$ be the vectors obtained by means of the ϵ -algorithm from the sequences*

$$\{s_p; s_{p+1} = F(s_p)\}_{0 \leq p}, \quad \text{and} \quad \{\tilde{s}_p; \tilde{s}_p = z + (F'(z))^p (s_0 - z)\}_{0 \leq p},$$

respectively. Assume that

(i) *1 is not an eigenvalue of $F'(z)$,*

(ii) *the vectors $\epsilon_q^{(p)}$, $\tilde{\epsilon}_q^{(p)}$, $p + q \leq 2m$, exist for all s_0 sufficiently close to z with $s_0 - z \in Q_m(F'(z))$.*

Set

$$(9) \quad \epsilon_{2m}^{(0)} = G(s_0, \dots, s_{2m}) = H_F(s_0),$$

then the computational procedure

$$x_{i+1} = H_F(x_i) \quad (0 \leq i)$$

is, for x_0 sufficiently close to z and $x_0 - z \in Q_m(F'(z))$, a quadratically convergent iterative method for the computation of z .

Proof. By the corollary and Lemma 2, we have

$$H_F(x_0) = \epsilon_{2m}^{(0)} = z + O(\|x_0 - z\|^2)$$

for $x_0 - z \in Q_m(F'(z))$.

5. A Modification of the Method. When a system of equations $x = F(x)$ of order n is to be solved by the ϵ -algorithm, the way of doing this is normally to put $m = n$. Then, we need, for each step of iteration, $4n^3 + 2n^2$ multiplications, $2n^2 + n$ divisions, $6n^3 - n^2$ additions/subtractions and the computation of $s_p = F(s_{p-1})$ for $1 \leq p \leq 2n$. The computation of the vectors s_p rather quickly produces a characteristic overflow if the eigenvalues of the Jacobian matrix $F'(x)$ are greater in absolute value than unity near the fixed point z . This disadvantage can possibly be eliminated by replacing the Picard sequence $s_{p+1} = F(s_p)$ by

$$s_{p+1} = F_\alpha(s_p) = (1 - \alpha)s_{p-1} + \alpha F(s_p) \quad (0 \leq p)$$

with a suitable α , $0 < \alpha < 1$; in this way, the rate of growth of the components of the vectors s_p is reduced. If we have, for example, $\rho(F'(z)) = 2$ for the spectral radius ρ of $F'(z)$, we get $\rho(F'_\alpha(z)) = 3/2$ for $\alpha = 1/2$. Those eigenvalues λ of $F'(z)$ for which $|\lambda| < 1$ are thereby increased, but they remain smaller than one in absolute value. Apart from this, convergence is slow if the eigenvalues of $F'(x)$ approach one near z .

The rounding errors affect the computation severely. Perhaps, it is possible that the numerical properties can be improved if a modification proposed by Wynn [20] is applied. If the eigenvalues λ of $F'(x)$ with $|\lambda| < 1$ predominate, we can indicate a modification of the method, by giving up the (theoretic) quadratic convergence, which considerably reduces the amount of work. To achieve this, we replace $2m$ by $2[(m+1)/2]$ in (9) and obtain for the basic formula of the algorithm

$$(9^*) \quad \epsilon_n^{(0)} = G(s_0, \dots, s_n) = H_F^*(s_0)$$

in the case $m = n$ even. We need now, per step of iteration, only

$$(n^3 + 8n^2 - 4n)/8$$

multiplications/divisions,

$$(6n^3 - 2n^2)/8$$

additions/subtractions and the computation of $s_p = F(s_{p-1})$ for $1 \leq p \leq n$.

6. Numerical Examples. Let $F: \mathbb{R}^4 \rightarrow \mathbb{R}^4$. In order to illustrate the method of Theorem 2 and its modifications, we consider some systems of quadratic equations $\mathbf{x} = F(\mathbf{x})$ with fixed point $\mathbf{z} = (1, 1, 1, 1)^T$:

$$(10) \quad F(\mathbf{x}) = \mathbf{z} + F'(\mathbf{z})(\mathbf{x} - \mathbf{z}) + \frac{1}{2}F''(\mathbf{z})(\mathbf{x} - \mathbf{z})^{(2)}.$$

For the Taylor series (10), we write briefly

$$(11) \quad F(\mathbf{x}) = \mathbf{z} + \mathbf{A}(\mathbf{x} - \mathbf{z}) + Q(\mathbf{x} - \mathbf{z})$$

and choose for \mathbf{A} (linear) and Q various mappings. The fixed point \mathbf{z} of the systems given in that manner is computed by means of single-precision arithmetic with ten decimal digits. In detail, let $P^{(i)}(\mathbf{x}) = (p_1^{(i)}(\mathbf{x}), \dots, p_4^{(i)}(\mathbf{x}))^T$ and

$$\begin{aligned} p_1^{(1)}(\mathbf{x}) &= -(x_1^2 + x_1x_4)/2, & p_1^{(2)}(\mathbf{x}) &= -x_1^2/4, \\ p_2^{(1)}(\mathbf{x}) &= -x_2^2/2, & p_2^{(2)}(\mathbf{x}) &= -x_2^2/4, \\ p_3^{(1)}(\mathbf{x}) &= -x_3^2/2, & p_3^{(2)}(\mathbf{x}) &= -x_3^2/4, \\ p_4^{(1)}(\mathbf{x}) &= -(x_4x_1 + x_4^2)/2, & p_4^{(2)}(\mathbf{x}) &= -x_4^2/4. \end{aligned}$$

Furthermore, let

$$\mathbf{D}_1 = (0.9, 0.8, 0.7, 0.6),$$

$$\mathbf{D}_2 = (1.5, 0.8, 0.7, 0.6),$$

$$\mathbf{D}_3 = (2.0, 0.8, 0.7, 0.6)$$

be diagonal matrices and

$$\mathbf{U}_1 = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}, \quad \mathbf{U}_2 = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 2 & 3 & 4 \\ 1 & 3 & 6 & 10 \\ 1 & 4 & 10 & 20 \end{pmatrix}.$$

We remark that \mathbf{U}_1 is orthogonal, whereas \mathbf{U}_2 is the ill-conditioned Pascal matrix of order four having an integer-valued inverse. It should be pointed out that

$$\left(\frac{\partial P^{(j)}(\mathbf{x} - \mathbf{z})}{\partial \mathbf{x}} \right) \bigg|_{\mathbf{x}=\mathbf{z}} = 0 \quad (\text{Matrix}) \quad (j = 1, 2);$$

hence, choosing $Q = P^{(j)}$ in eq. (11), we get, indeed, $F'(\mathbf{z}) = \mathbf{A}$. Now, if $\mathbf{A} = \mathbf{U}_m \mathbf{D}_l \mathbf{U}_m^{-1}$ ($l = 1, 2, 3; m = 1, 2$), then \mathbf{D}_l is the matrix of eigenvalues and \mathbf{U}_m is the matrix of eigenvectors of $F'(\mathbf{z})$.

In Examples I–VI, \mathbf{z} is computed by the method proposed in Theorem 2.

I	2.0	$1.2 \cdot 10^{-2}$	$1.0 \cdot 10^{-5}$	$1.4 \cdot 10^{-1}$	$6.8 \cdot 10^{-2}$	$8.4 \cdot 10^{-3}$	$7.5 \cdot 10^{-5}$	$2.5 \cdot 10^{-8}$
II	$7.4 \cdot 10^{-1}$	$6.6 \cdot 10^{-1}$	$4.5 \cdot 10^{-1}$					
III	$6.0 \cdot 10^{-1}$	$5.4 \cdot 10^{-5}$						
IV	0.9	$8.2 \cdot 10^{-2}$	$2.7 \cdot 10^{-6}$					
V	2.0	$9.9 \cdot 10^{-1}$	$2.8 \cdot 10^{-6}$					
VI	1.8	$1.5 \cdot 10^{-1}$	$3.7 \cdot 10^{-2}$					
VII	1.9	$8.6 \cdot 10^{-2}$	$5.5 \cdot 10^{-3}$	$1.1 \cdot 10^{-2}$	$4.4 \cdot 10^{-3}$	$2.2 \cdot 10^{-3}$	$1.1 \cdot 10^{-3}$	$3.6 \cdot 10^{-4}$
VIII	$7.9 \cdot 10^{-1}$	$6.5 \cdot 10^{-1}$	$4.3 \cdot 10^{-1}$	$5.0 \cdot 10^{-5}$	$5.2 \cdot 10^{-8}$			
IX	$6.0 \cdot 10^{-1}$	$6.0 \cdot 10^{-3}$	$4.0 \cdot 10^{-6}$	$1.3 \cdot 10^{-1}$	$4.6 \cdot 10^{-2}$	$1.6 \cdot 10^{-3}$	$5.1 \cdot 10^{-5}$	$4.8 \cdot 10^{-8}$
X	$8.9 \cdot 10^{-1}$	$1.1 \cdot 10^{-1}$	$3.2 \cdot 10^{-4}$	$2.4 \cdot 10^{-7}$				
XI	$3.8 \cdot 10^{-1}$	$5.1 \cdot 10^{-1}$	$1.1 \cdot 10^{-1}$	$3.8 \cdot 10^{-4}$	$1.1 \cdot 10^{-6}$			
XII	1.5	$3.2 \cdot 10^{-1}$	$1.1 \cdot 10^{-1}$	$4.6 \cdot 10^{-2}$	$2.1 \cdot 10^{-2}$	$1.0 \cdot 10^{-2}$	$9.5 \cdot 10^{-4}$	$6.5 \cdot 10^{-4}$

Example I: $F'(z) = U_1 D_1 U_1^{-1}$, $Q = P^{(1)}$, initial vector $x_0 = 2z$;

Example II: $F'(z) = U_1 D_2 U_1^{-1}$, $Q = P^{(1)}$, $x_0 = 0$;

Example III: as Example II but using $x_0 = 2z$;

Example IV: $F'(z) = U_2 D_2 U_2^{-1}$, $Q = P^{(2)}$, $x_0 = 0.5z$;

Example V: as Example IV but using $x_0 = 1.5z$;

Example VI: $F'(z) = U_1 D_3 U_1^{-1}$, $Q = P^{(1)}$, $x_0 = 2z$, using the modified Picard sequence $s_{p+1} = F_\alpha(s_p)$ with $\alpha = 1/2$.

The Examples VII–XII are the same as Examples I–VI, respectively, but z is computed using formula (9*) instead of (9).

The above table contains in column i ($1 \leq i \leq 8$) the values $\|x_i - x_{i-1}\|$ (compare Theorem 2) with rounded mantissae; values for which $\|z - x_i\| < 5.0 \cdot 10^{-9}$ (the process has then terminated) are omitted. Generally speaking, we have found that the algorithm produces better results if the Jacobian matrix of the given system $x = F(x)$ is symmetric. Finally, it should be mentioned that it seems to be impossible at the moment to say more about the error than that it is of quadratic order.

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Universität Mannheim
Lehrstuhl für Mathematik IV
Mannheim 6800
Germany

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