## Input:

- Set  $\mathcal{R} = (R_1, R_2, ..., R_m)$  representing available resources, where each resource  $R_k \in \mathcal{R}$  has capacity  $c_k$ .
- Value  $c^{(GPU)}$  representing GPU capacity.
- Set of tasks  $\mathcal{T} = (T_1, T_2, ..., T_n)$ .
- Values  $p_{i,k}$  and  $E_{i,k}$  representing processing time and energy consumption of task  $T_i$  on resource  $R_k$ . If task  $T_i$  cannot be executed on resource  $R_k$ , then  $p_{i,k} = \infty$ .
- Values  $p_{i,k}^{(GPU)}$  and  $E_{i,k}^{(GPU)}$  representing processing time and energy consumption of task  $T_i$  on resource  $R_k$  if the task is offloaded to GPU. If task  $T_i$  cannot be executed on resource  $R_k$ , or cannot be offloaded to GPU, then  $p_{i,k}^{(GPU)} = \infty$ .
- Value h representing major frame length.

## ILP model:

$$\min \sum_{T_i \in \mathcal{T}} \sum_{W_j \in \mathcal{W}} \sum_{R_k \in \mathcal{R}} \left( a_{i,j,k} \cdot E_{i,k} + a_{i,j,k}^{(GPU)} \cdot E_{i,k}^{(GPU)} \right) \text{ subject to:}$$
 (1)

$$\sum_{W_i \in \mathcal{W}} \sum_{R_k \in \mathcal{R}} (a_{i,j,k} + a_{i,j,k}^{(GPU)}) = 1 \quad \forall T_i \in \mathcal{T}$$
 (2)

$$\sum_{W_j \in \mathcal{W}} l_j \le h \tag{3}$$

$$l_{j} \geq \frac{a_{i,j,k} \cdot p_{i,k} + a_{i,j,k}^{(GPU)} \cdot p_{i,k}^{(GPU)}}{0.6} \quad \forall T_{i} \in \mathcal{T}, \forall W_{j} \in \mathcal{W}, \forall R_{k} \in \mathcal{R}$$
 (4)

$$\sum_{T_i \in \mathcal{T}} (a_{i,j,k} + a_{i,j,k}^{(GPU)}) \le c_k \quad \forall W_j \in \mathcal{W}, \forall R_k \in \mathcal{R}$$
 (5)

$$\sum_{T_i \in \mathcal{T}} \sum_{R_k \in \mathcal{R}} a_{i,j,k}^{(GPU)} \le c^{(GPU)} \quad \forall W_j \in \mathcal{W}$$
 (6)

$$a_{i,j,k} = 0 \quad \forall T_i \in \mathcal{T}, \forall W_j \in \mathcal{W}, \forall R_k \in \mathcal{R} : p_{i,k} = \infty$$
 (7)

$$a_{i,j,k}^{(GPU)} = 0 \quad \forall T_i \in \mathcal{T}, \forall W_j \in \mathcal{W}, \forall R_k \in \mathcal{R} : p_{i,k}^{(GPU)} = \infty$$
 (8)

$$a_{i,j,k} \in \{0,1\}, a_{i,j,k}^{(GPU)} \in \{0,1\} \quad \forall T_i \in \mathcal{T}, \forall W_j \in \mathcal{W}, \forall R_k \in \mathcal{R}$$
 (9)

$$l_i \in \mathbb{N} \quad \forall W_i \in \mathcal{W}$$
 (10)