



BSV Training

Eg03: Concurrent Bubblesort

A simple concurrent Bubblesort example.

```

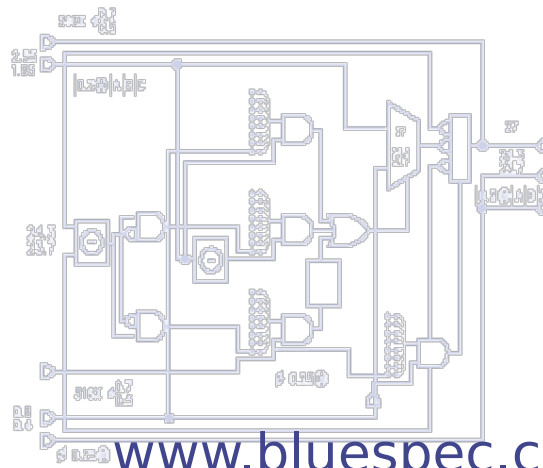
import P2P2:
  typedef BitN(24) (uint);
  module ex_hdl_csr2_bsp;

    Integer nfa_depth = 32;

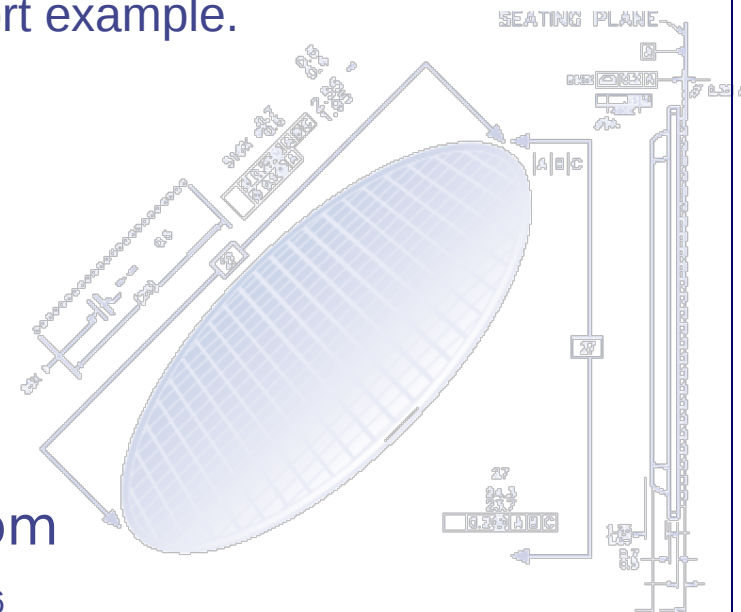
    function BitN(24) determine_pump@dataTest;
      return (0);
    endfunction

    P2P2@dataTest in_bond;
    ex_hdl_csr2_bsp@dataTest in_bond;
    P2P2@dataTest out_bond;
    ex_hdl_csr2_bsp@dataTest out_bond;
    P2P2@dataTest out_bond;
    ex_hdl_csr2_bsp@dataTest out_bond;

    rule end (True);
      dataTest in_bond = in_bond;
      P2P2@dataTest out_bond =
        determine_pump@dataTest = 0 ? out_bond : out_bond;
      out_bond;
    endrule;
  endmodule;
  
```



www.bluespec.com



Generalization via 5 versions

The accompanying code demonstrates several versions:

Eg03a_Bubblesort/	Sorts 5 items, each of type <code>`Int#(32)`</code> . Completely sequential, to show correspondence with conventional software implementation.
Eg03b_Bubblesort/	Parallel version. Uses <code>'maxBound'</code> (largest <code>Int#(32)</code>) as a “special” value different from any of the sorted values (like “+infinity”).
Eg03c_Bubblesort/	Generalizes <code>'5'</code> items to <code>'n'</code> items.
Eg03d_Bubblesort/	Generalizes items of type <code>'Int#(32)'</code> to items of arbitrary type <code>'t'</code> (i.e., makes the program polymorphic).
Eg03e_Bubblesort/	Uses the <code>'Maybe#(t)'</code> type to eliminate the need for a <code>'maxBound'</code> special value.

Note: this is a pedagogical introductory example focusing on concurrency and modularity, and is not intended as an example of efficient sorting!
(See `'Eg06_Mergesort'` for more efficient and scalable sorting.)

1st version: directory Eg03a_Bubblesort/

Examine the two source files: src_BSV/Testbench.bsv src_BSV/Bubblesort.bsv

We suggest that you take multiple passes through the files, incrementally increasing your understanding. Initially, just try to understand the syntactic structure, making frequent reference to the lecture slides in “Lec_Basic_Syntax” in the “Reference” directory.

The following excerpts show an outline of the syntactic structure of the two source files.

```
package Testbench;
...
import ...
import Bubblesort :: *;
...
Int#(32) n = 5;
...
module mkTestbench (Empty);
  Reg #(Int#(32)) rg_j1 <- mkReg (0);
  ...
  Sort_IFC sorter <- mkBubblesort;
  ...
  rules
```

```
package Bubblesort;
...
import ...
...
interface Sort_IFC;
  ...
  module mkBubblesort (Sort_IFC);
    Reg #(UInt #(3)) rg_j <- mkReg (0);
    ...
    rules
    ...
    method definitions
```

Examine the two source files: src_BSV/Testbench.bsv src_BSV/Bubblesort.bsv

```
package Testbench;
...
import ...
import Bubblesort :: *;
...
Int#(32) n = 5;
...
module mkTestbench (Empty);
  Reg #(Int#(32)) rg_j1 <- mkReg (0);
  ...
  Sort_IFC sorter <- mkBubblesort;
  ...
  rules
```

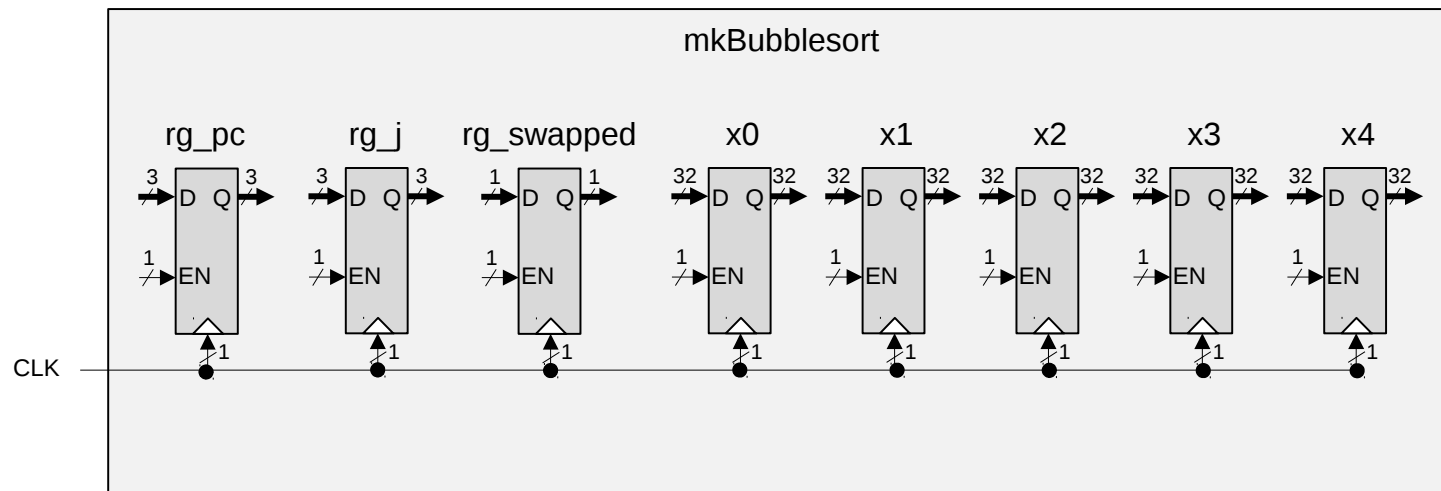
```
package Bubblesort;
...
import ...
...
interface Sort_IFC;
  ...
  module mkBubblesort (Sort_IFC);
    Reg #(UInt #(3)) rg_j <- mkReg (0);
    ...
    rules
    ...
    method definitions
```

Notes on the syntactic structure:

- The top-level module is mkTestbench; it has an Empty interface (an interface with no methods).
- It instantiates several sub-modules: a few registers, a pseudo-random number generator (LFSR, for Linear Feedback Shift Register), and the mkBubblesort module.
- The mkBubblesort module, in turn, instantiates some registers as its sub-modules.
- The rules in the testbench invoke the put and get methods in the bubblesort module's interface.
- All the rules and methods have Boolean conditions indicating when they are “ready”.

In module mkBubblesort:

- The register `rg_pc` is used to sequence the sorting actions. The name is suggestive of “Program Counter”. 3 bits are enough to distinguish all the states.
- The register `rg_j` is used to count 5 incoming values. Thus, it is typed `UInt#(3)`, i.e., an unsigned integer of 3 bits, capable of holding values 0..7. It’s reset value (initial value) is 0.
- The register `rg_swapped` remembers whether any swap occurred in the current pass.
- The registers `x0..x4` hold the 5 values to be sorted. Each value is of type `Int#(32)`, i.e., a signed integer of 32 bits. Each register’s reset value is `maxBound`, a symbolic name representing the largest `Int#(32)`, i.e., $+2^{31-1}$. We assume all input values to be sorted will be strictly $< \text{maxBound}$



The module `mkBubblesort`, showing the registers in isolation

The rules in module mkBubblesort essentially encode a sequential algorithm similar to that shown below in a pseudo-C notation:

```
swapped = False;
while (True)
  for (pc = 1; pc < 5; pc++) {
    if (x[pc-1] > x[pc] {
      swap them;
      swapped = True;
    }
    if (! swapped) break;
  }
```

Each rule has:

- A condition that says when it can fire, and
- A body that says what happens if it fires.

Building and running the codes

Each variation is built and run in the same way:

- In the Build/ directory you can use the 'Makefile' for building and running Bluesim or Verilog sim:

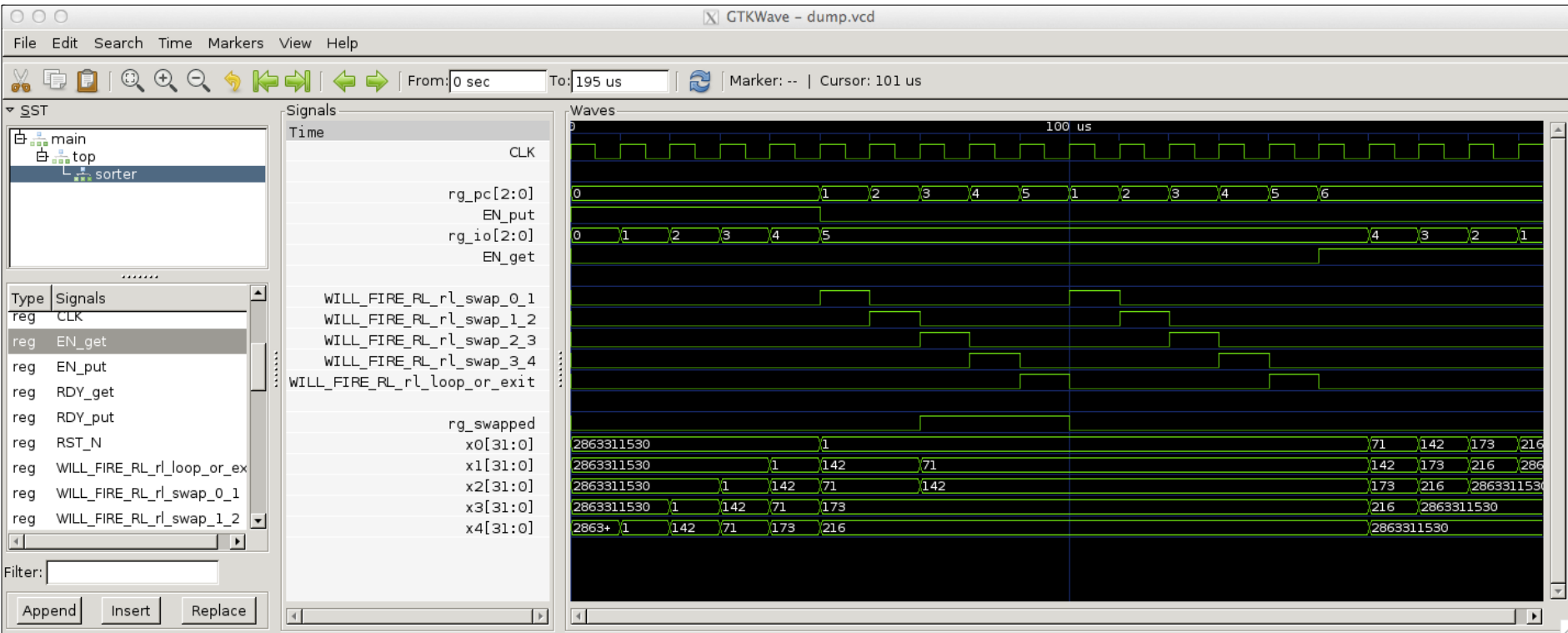
```
% make compile link simulate      // for Bluesim
% make verilog v_link v_simulate  // for Verilog sim
```

Note: When building for Bluesim, compiler-temporary files are created in the build_bsim/ directory. When building for Verilog, temporaries are in build_v/ and Verilog files are in verilog_dir/. 'make clean' and 'make full_clean' will clean up these directories.

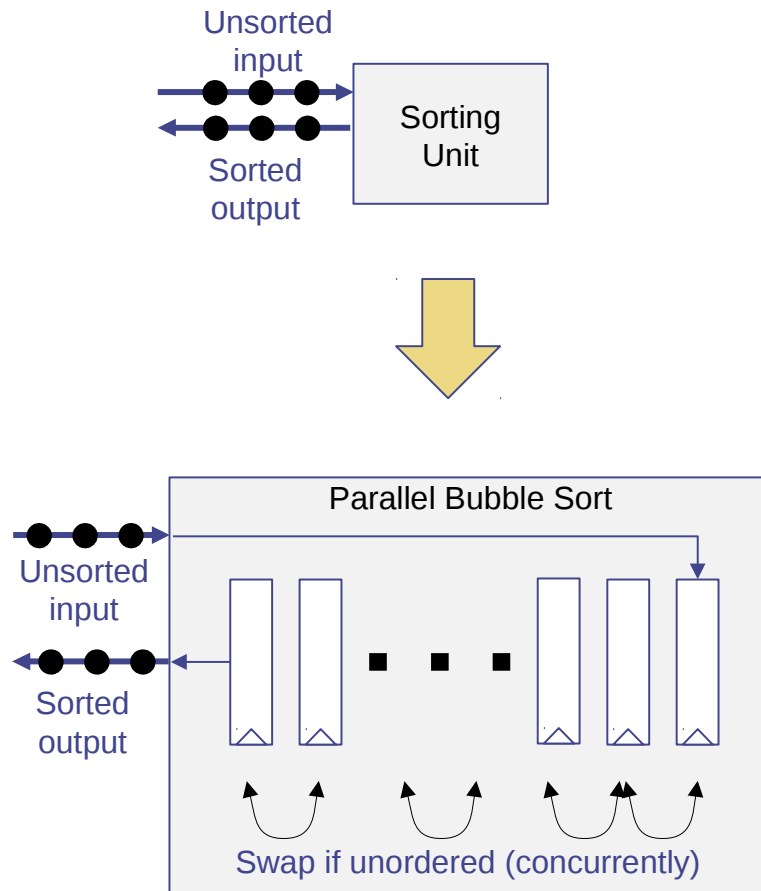
Build and run the 1st version

- In the Build/ directory, practice building and running the program in both ways:
 - Using 'make' commands, for Bluesim.
 - Using 'make' commands, for Verilog sim.
- Observe the inputs and outputs and verify that they are reasonable (there are 5 inputs and outputs, and the outputs are a sorted version of the inputs).
- When you run your simulation, a file “dump.vcd” is created, containing waveforms, which you can view in your favorite waveform viewer.
 - (It is created because of the -V flag for Bluesim or the +bscvcd flag for Verilog sim.)
 - The best way to view waveforms is from BDW (Bluespec Development Workstation), because it can customize the view to show data using the more expressive BSV source code types, instead of the less expressive Verilog bits.
 - The picture “Waves_screenshot.tiff” is a screenshot of such a view.
- Study the waveform and make sure you understand how it reflects the behavior of the BSV code.

Waveforms from the 1st version



Eg03: Basic concurrency and modularity



Algorithm/Architecture for remaining versions:

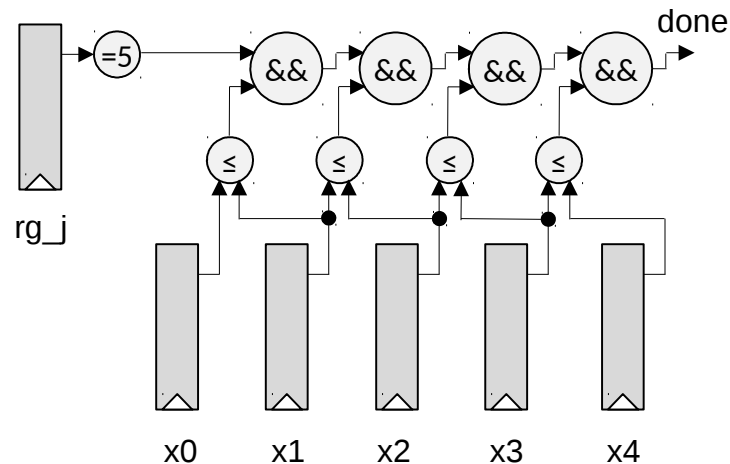
- The module accepts a stream of 'n' input items (unsorted).
- These are shifted in to 'n' registers.
- Concurrently (and even while inputs are arriving), whenever two adjacent registers contain values in the wrong order, we swap their contents.
- When 'n' inputs have been received, and all of them are sorted, the module yields a stream of 'n' sorted outputs by shifting them out.
- When 'n' outputs have been streamed out, the module is ready for its next set of 'n' inputs.

In module mkBubblesort:

- The Bool function “done” defines the condition that all values have been received and are sorted:

```
// Test if array is sorted
function Bool done ();
  return ((rg_j == 5) && (x0 <= x1) && (x1 <= x2)
          && (x2 <= x3) && (x3 <= x4));
endfunction
```

- Note: in BSV, a zero-argument function like this is identical to a constant
- In BSV, value-expressions like this just represent combinational circuits:

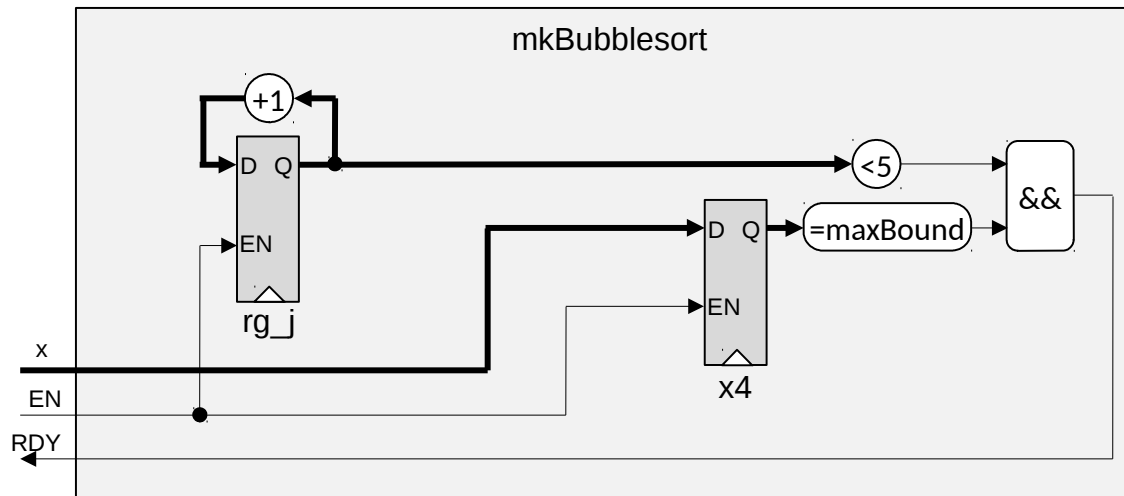


```
// Inputs: feed input values into x4
method Action put (Int#(32) x) if ((rg_j < 5) && (x4 == maxBound));
    x4 <= x;
    rg_j <= rg_j + 1;
endmethod
```

In module mkBubblesort:

- The “put” method is of type Action, i.e., it’s body is an expression of type Action.
 - The body has two sub-Actions: one places a new input value x into register x4, and the other increments the input count rg_j.
 - The method condition ensures two things:
 - “(rg_j < 5)” ensures that we stop after receiving 5 inputs
 - “(x4==maxBound)” ensures that we don’t over-write a meaningful value already in x4, i.e., it can place a new value in x4 only after the swap rules have moved the previous value out and replaced it with maxBound.

Hardware for put method, in isolation



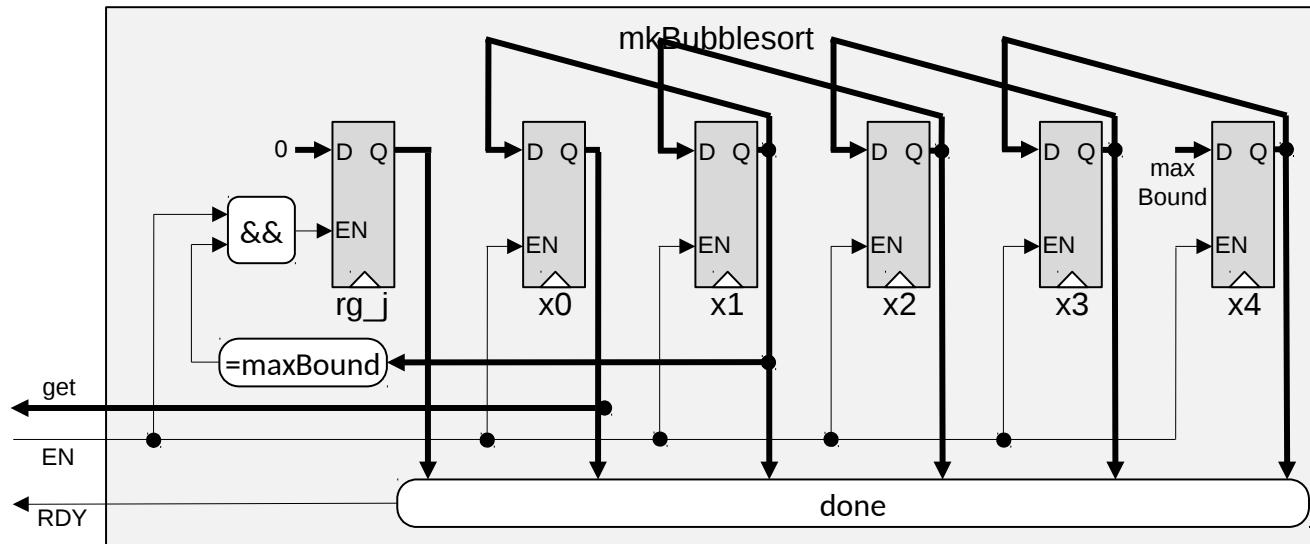
```
// Outputs: drain by shifting them out of x0
method ActionValue#(Int#(32)) get () if (done);
  x0 <= x1;
  x1 <= x2;
  x2 <= x3;
  x3 <= x4;
  x4 <= maxBound;
  if (x1 == maxBound) rg_j <= 0;
  return x0;
endmethod
```

In module mkBubblesort:

- The “get” method is of type ActionValue #(Int#(32)), i.e., it’s body is an Action, and it also returns a value of type Int#(32).
- The body returns x0, shifts all remaining values in x1..x4 down into x0..x3, respectively, and shifts the value maxBound into x4.
- The method condition of the “get” method:
 - Is only enabled when “done” is true, i.e., all inputs have been received and all values are sorted
 - Is only enabled until we’ve returned the 5th value, after which x0 contains maxBound
- When we return the 5th value (from x0), x1 will be maxBound; at this time we can reset rg_j to 0
- Note that after returning 5 values, the module is again in it’s original state—rg_j contains 0, and x0..x4 contain maxBound—and so it is ready to receive the next 5 values to be sorted

```
// Outputs: drain by shifting them out of x0
method ActionValue#(Int#(32)) get () if (done);
  x0 <= x1;
  x1 <= x2;
  x2 <= x3;
  x3 <= x4;
  x4 <= maxBound;
  if (x1 == maxBound) rg_j <= 0;
  return x0;
endmethod
```

Hardware for “get” method, in isolation



```

rule r1_swap_0_1 (x0 > x1);
  x0 <= x1;
  x1 <= x0;
endrule

```

```

rule r1_swap_1_2 (x1 > x2);
  x1 <= x2;
  x2 <= x1;
endrule

```

```

rule r1_swap_2_3 (x2 > x3);
  x2 <= x3;
  x3 <= x2;
endrule

```

```

(* descending_urgency =
"r1_swap_3_4, r1_swap_2_3,
r1_swap_1_2, r1_swap_0_1" *)
rule r1_swap_3_4 (x3 > x4);
  x3 <= x4;
  x4 <= x3;
endrule

```

In module mkBubblesort:

- In each of the four rules:
 - The rule condition (an expression of type Bool), e.g., “(x0>x1)” tests if two adjacent values are in the wrong order. This is a prerequisite for the rule to “fire”, i.e., to execute its body.
 - The rule body (an expression of type Action) is composed of two sub-Actions whose effect is to swap the contents of the corresponding two registers. Note:
 - Expressions/methods/functions of type “Action” or “ActionValue#(t)” (potentially) change the state of the system (e.g., write to a register, write to memory, enqueue onto a FIFO, etc.). We also say that such expressions have “side-effects”.
 - Expressions that are not of type Action or ActionValue (e.g., Bool) *never* change the state of the system (this is guaranteed by BSV’s type-checking rules). We also say that these are “pure” or “value” expressions.
 - All Actions in a rule are considered instantaneous and simultaneous. The two assignments are not read sequentially (as is typical in software programs), and this explains why there is no need for a “temporary” intermediate variable (as is typical in software programs)

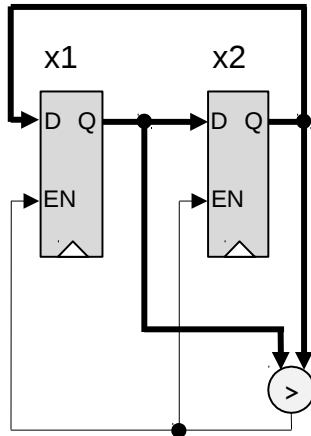
```
rule rl_swap_0_1 (x0 > x1);
  x0 <= x1;
  x1 <= x0;
endrule
```

```
rule rl_swap_1_2 (x1 > x2);
  x1 <= x2;
  x2 <= x1;
endrule
```

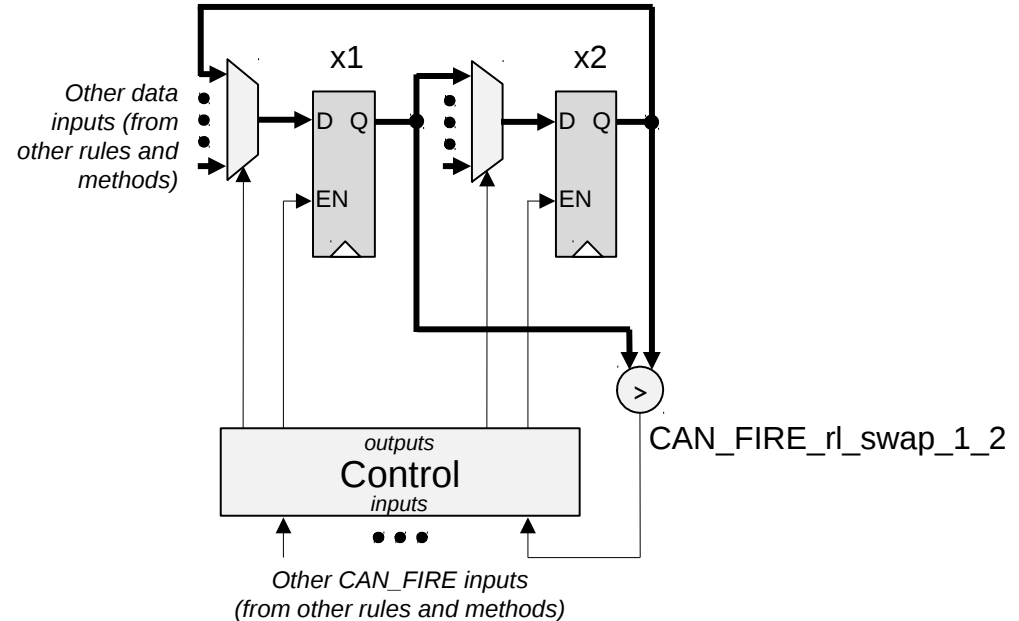
```
rule rl_swap_2_3 (x2 > x3);
  x2 <= x3;
  x3 <= x2;
endrule
```

```
(* descending_urgency =
"rl_swap_3_4, rl_swap_2_3,
rl_swap_1_2, rl_swap_0_1" *)
rule rl_swap_3_4 (x3 > x4);
  x3 <= x4;
  x4 <= x3;
endrule
```

Hardware for rl_swap_1_2, in isolation



... and in more detail



In module mkTestbench:

- rg_j1 and rg_j2 are used to count 5 inputs and 5 outputs respectively. They are somewhat arbitrarily typed at Int#(32) (since they count 0..4, even Int#(3) or UInt#(3) would have been ok, but we usually don't worry so much about minimizing state bits in a testbench).
- mkLFSR_8 instantiates, from the BSV library, a pseudo-random number generator of 8-bit values (see Sec.C.8.1 in the Reference Guide).
- In rule "rl_feed_inputs":
 - The rule condition ensures that is only enabled to generate 5 inputs
 - lfsr.value() is a "value method" that returns a Bit#(8); zeroExtend() extends this to Bit#(32); unpack() converts this to an Int#(32). lfsr.next() is an Action method that tells the LFSR to advance to its next random number
 - "sorter.put(x)" sends this new input into the bubblesort module
 - Recall that the "put" method in mkBubblesort has a method condition that prevents overwriting a previous value. Thus, rl_feed_inputs can only fire when that condition is true.
- In rule "rl_drain_outputs"
 - The rule condition ensures that it only retrieves 5 outputs
 - "sorter.get()" yields the next output, whenever the method condition on the "get" method allows, and this value is displayed.

Summary of behavior:

- Rule “rl_feed_inputs” in the testbench module feeds 5 random values into the sorter module using the “put” method.
- The sorter module receives the 5 inputs via the “put” method. Sorting begins as soon as the first input arrives (the first step sorts the first input against its +infinity neighbor), and continues as more inputs arrive. Sorting may continue for some time after all inputs arrive, until all 5 registers are sorted.
- When all 5 registers are sorted, the “get” method is enabled, and the testbench drains out the 5 values in order. The sorter module yields these 5 values in order by shifting them through the registers.
- After yielding 5 outputs, the sorter module is back in its initial state, and ready for another 5 inputs (although our testbench quits after the first set).

This behavior can be seen in the waveforms for this design created during simulation. We will shortly describe simulation and waveform generation but, for your convenience, the directory includes a screen shot (“waves_screenshot.tiff”) which you can view directly.

Build and run the 2nd version

- In the Build/ directory, practice building and running the program in both ways:
 - Using 'make' commands, for Bluesim.
 - Using 'make' commands, for Verilog sim.
- Observe the inputs and outputs and verify that they are reasonable (there are 5 inputs and outputs, and the outputs are a sorted version of the inputs).
- Note that an input is supplied only on *every other* clock. Why?
 - (Hint: see previous remarks on method condition of the "put" method.)
- When you run your simulation, a file "dump.vcd" is created, containing waveforms, which you can view in your favorite waveform viewer.
 - (It is created because of the -V flag for Bluesim or the +bscvcd flag for Verilog sim.)
 - The best way to view waveforms is from BDW (Bluespec Development Workstation), because it can customize the view to show data using the more expressive BSV source code types, instead of the less expressive Verilog bits.
 - The picture "waves_screenshot.tiff" is a screenshot of such a view. Note that it shows BSV source code types Int#(32), Bool, ..., and not raw Verilog bits.
- Study the waveform and make sure you understand how it reflects the behavior of the BSV code.

Suggested exercises

In this and future examples, we suggest extra exercises to deepen your understanding of BSV, which you can pursue on your own (may not be covered during classroom training)

- Examine the generated Verilog file mkBubblesort.v (in the “verilog/” directory).
 - Look at the input and output ports, and understand how they correspond to the BSV interface Sort_IFC and its methods.
 - Skim the interior of the Verilog module, and notice correspondences with the BSV source module (registers, rules, rule and method conditions, ...).
- Analyze and explain in detail what would happen if the testbench supplied ‘maxBound’ as one (or more) of the input values. If you wish, modify the testbench to test this.
- Modify the testbench to sort two (or three, or more) rounds of 5-input sequences, instead of quitting after the first round.
- Modify the testbench and the bubblesort modules to sort sequences of some other length (say, 10) instead of 5.

Time-out to reinforce some concepts

Before moving on with the examples, please study the lecture: [Lec_Rule_Semantics](#) to understand the concepts behind Rules:

- Semantics of individual rules (and the methods they call)
 - “simultaneous/parallel actions”, “instantaneous”
- Semantics of concurrency of a set of rules in each clock
 - “concurrency”
 - “ordering constraints” between methods in different rules
 - “rule schedules”

The 3rd version generalizes the 2nd version from sorting 5 numbers to sorting n numbers

Examine the two source files: `src_BSV/Testbench.bsv` `src_BSV/Bubblesort.bsv`

It is useful to compare these side-by-side with the corresponding two files in `Eg03b_Bubblesort/src_BSV/`, to see how they have changed.

Notice that the interface type is now parameterized by n :

```
interface Sort_IFC #(numeric type n_t);
```

In BSV, certain types and type parameters can be *numeric types*.

Note that *numeric types* (which are only meaningful at compile time) are distinct from *numeric values* (which can of course occur at run time).

BSV is very strict in type-checking—there are no automatic conversions
The code uses these explicit conversions:



The 2nd version's explicit registers x0..x4 have here been replaced by a vector of registers:

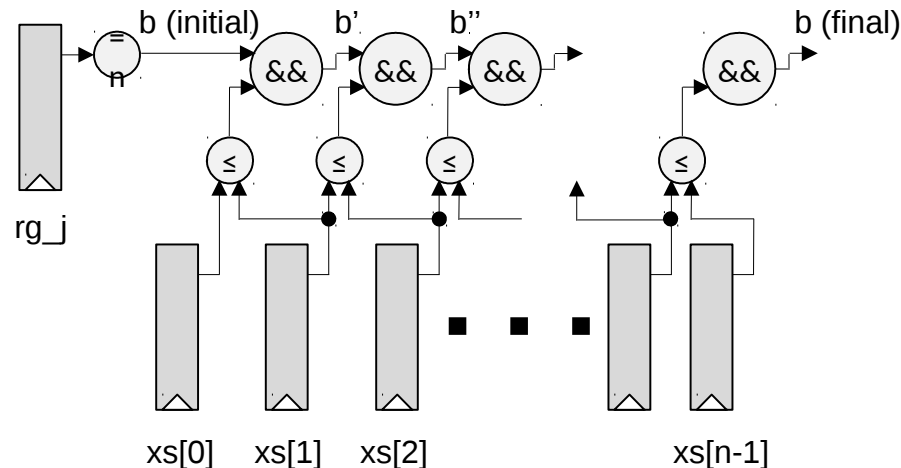
```
Vector #(n_t, Reg #(Int #(32))) xs <- replicateM (mkReg (maxBound));
```

“replicateM” repeatedly applies “mkReg(maxBound)” to create them.

The 2nd version's explicit “done” test is here computed in a for-loop:

```
function Bool done ();
  Bool b = (rg_j == fromInteger (n));
  for (Integer i = 0; i < n-1; i = i + 1)
    b = b && (xs[i] <= xs[i+1]);
  return b;
endfunction
```

Note: this for-loop is also “statically elaborated” (i.e., unfolded by the compiler) into a circuit of combinational logic



Observe that, unlike traditional languages, there is no storage location corresponding to variables like *i* and *b*. The former has disappeared completely, and the latter exists in multiple versions (*b*, *b'*, *b''*, ...) each of which just names a wire.

The 2nd version's separately-written rules are here replaced by a for-loop generating them:

```
for (Integer i = 0; i < n-1; i = i+1)
  rule rl_swap_i (xs [i] > xs [i+1]);
    xs [i]    <= xs [i+1];
    xs [i+1] <= xs [i];
  endrule
```

Note: this for-loop is “statically elaborated” (i.e., unfolded by the compiler), and is equivalent to writing out n-1 rules explicitly.

In the “put” method, the following line:

```
writeVReg (xs, shiftInAtN (readVReg (xs), maxBound));
```

uses a number of BSV library functions on vectors:

- readVReg() returns a vector of values corresponding to the contents of a vector of registers
- shiftInAtN() takes a vector of values x_1, \dots, x_N and a new value y , and returns a vector of values x_2, \dots, x_N, y (discards x_1)
- writeVReg() writes a vector of values into a vector of registers

The above line could also be have been written like this:

```
for (Integer i = 0; i < n-1; i = i+1)
  xs[i] <= xs[i+1];
xs[n-1] <= maxBound;
```

Note: this for-loop is also “statically elaborated” (i.e., unfolded by the compiler)

Time-out to reinforce some concepts

Before moving on with the examples, please study the lecture: `Lec_Types` to understand the concepts behind types, polymorphism, and numeric types.

Please also *skim* through Section C.3 (“Vectors”) in the Reference Guide.

Synthesis hierarchy

- In the 1st and 2nd versions, in front of 'module mkBubblesort' line, there is a (*synthesize*) attribute:

```
(* synthesize *)  
module mkBubblesort(Sort_IFC);  
...  
endmodule
```

- When generating Verilog, this creates a 'mkBubblesort.v' file with a 'mkBubblesort' Verilog module
- In the 3rd version, this (*synthesize*) attribute is removed.
- This is because BSV cannot separately synthesize *polymorphic* modules; they can only be inlined into a parent module
- Instead, in Testbench.bsv, we have created a specific instance of the module (for $n=20$); since this is no longer polymorphic, it can be separately synthesized. This is the module actually instantiated in the module mkTestbench:

```
typedef 20 N_t;  
...  
(* synthesize *)  
module mkBubblesort_nt (Sort_IFC #(N_t));  
    Sort_IFC #(N_t) m <- mkBubblesort;  
    return m;  
endmodule
```

- The above is a common idiom in BSV code, for creating a separately synthesized instance of a polymorphic module

Build and run the 3rd version

- In the Build directory, build and run using the 'make' commands, with Bluesim and/or with Verilog sim, as described earlier
- Observe the inputs and outputs and verify that they are reasonable (that there are 20 inputs and outputs, and the outputs are a sorted version of the inputs)
- In Testbench.bsv, change the 1 line that defines the size of the problem:

```
typedef 20 N_t;
```

Rebuild and re-run to test it.

The 4th version generalizes the 3rd version from sorting values of type `Int#(32)` to sorting values of any type “`t`”

Examine the two source files: `src_BSV/Testbench.bsv` `src_BSV/Bubblesort.bsv`

It is useful to compare these side-by-side with the corresponding two files in the previous version, to see how they have changed.

Notice that the interface type is now further parameterized by `t`:

```
interface Sort_IFC #(numeric type n_t, type t);
```

The “module `mkBubblesort`” line is now extended with *provisos*:

```
provisos (Bits #(t, wt),  
          Ord #(t),  
          Eq #(t),  
          Bounded #(t));
```

These are assertions, respectively, that:

- values of type `t` have a bit representation (with bit-width `wt`), since we need to store them in registers
- values of type `t` can be compared with the `<=` operator (ordering)
- values of type `t` can be compared with the `==` operator (equality)
- there exists a ‘`maxBound`’ value of type `t`

Without these assertions, the compiler would emit a type-checking error, since all these properties are used inside the module.

Time-out to reinforce some concepts

Please study the lecture: `Lec_Typeclasses`
to understand the concepts behind `Bits`, `Ord`, `Eq`, `Bounded`.

Build and run the 4th version

- In the Build directory, build and run using the 'make' commands, with Bluesim and/or with Verilog sim, as described earlier
- Observe the inputs and outputs and verify that they are reasonable (that there are 20 inputs and outputs, and the outputs are a sorted version of the inputs)
- In Testbench.bsv, change the two lines that define the size of the problem and the type of values being sorted:

```
typedef 20 N_t;  
typedef UInt #(24) MyT;
```

Rebuild and re-run to test it.

The 5th version eliminates the dependency on the existence of a separate ‘maxBound’ value that is separate from legitimate input values.

Before proceeding, please review the lecture: Lec_Types and, specifically, the section on “Maybe” types.

Examine the two source files: src_BSV/Testbench.bsv src_BSV/Bubblesort.bsv

It is useful to compare these side-by-side with the corresponding two files in the previous version, to see how they have changed.

Notice that the “Bounded#(t)” proviso has been removed from module mkBubblesort.

The vector of registers now contain values of type “Maybe#(t)” instead of “t”:

```
Vector #(n_t, Reg #(Maybe #(t))) xs <- replicateM (mkReg (tagged Invalid));
```

These are initialized/reset to the value “tagged Invalid”.

This plays the role previously played by “maxBound”.

In the “put” method, instead of inserting x into the register, we insert:

```
xs[jMax] <= tagged Valid x;
```

After the module mkBubblesort, there is some new code

```
instance Ord #(Maybe #(t))
  provisos (Ord #(t));
...
```

A value of type `Maybe#(t)` has two forms:

- tagged Invalid
- tagged Valid x

Conceptually, if values of type `t` are represented in `n` bits, then a value of type `Maybe#(t)` is represented in `n+1` bits. The extra bit is called a tag. When the tag is 0 (Invalid), the remaining `n` bits don't matter. When the tag is 1 (Valid), the remaining `n` bits represent a legitimate value of type `t`.

The “instance” declaration defines how to compare two values of type `Maybe#(t)`:

- tagged Invalid <= tagged Invalid
- tagged Valid x < tagged Invalid for any x
- tagged Invalid > tagged Valid y for any y
- tagged Valid x <= tagged Valid y if (x <= y)

Please see the lecture: `Lec_Typeclasses`
for a more detailed explanation of these concepts.

Build and run the 5th version

- In the Build directory, build and run using the 'make' commands, with Bluesim and/or with Verilog sim, as described earlier
- Observe the inputs and outputs and verify that they are reasonable (that there are 20 inputs and outputs, and the outputs are a sorted version of the inputs)
- In Testbench.bsv, change the two lines that define the size of the problem and the type of values being sorted:

```
typedef 20 N_t;  
typedef UInt #(24) MyT;
```

Rebuild and re-run to test it.

Time-out to reinforce some concepts

Please study the lecture: `Lec_Types`
to understand the concepts behind the `Maybe` type.

Suggested exercises

- Change the program to sort in *descending* order instead of ascending order.
- Instead of sorting scalar values, change the testbench to sort struct values:
 - Define a struct type with at least two fields.
 - Generate structs with random values for these fields.
 - Define the “<” operator on this struct type so as to compare only one field (you will have to declare an “Ord” instance for this struct type).
 - Test your program.
- Study Section C.3.10 (“Fold functions”) in the Reference Guide. In the sorter module, redefine the “done” function to use foldl(), foldr() or fold() in place of the explicit for-loop. What is the circuit structure using the for-loop, foldl, foldr and fold? What is the advantage of the circuit structure using fold?
- The “done” function is a combinational circuit (using any of the above implementations). For large n , this can limit the clock speed at which the circuit can be synthesized. Split the circuit into multiple smaller stages, by adding registers at suitable points and introducing rules to propagate values.
 - This means that detecting “done” will be delayed by a few cycles. Does this matter?
 - How should you reset these registers during final output in the “get” method?

Summary

These examples have provided a basic familiarity with many of the concepts in the BSV language:

- File and package structure
- Module structure, module instantiation, interfaces and methods
- Rules and methods, and their semantics
- Types and typeclasses, numeric types
- Static elaboration
- Parameterization on sizes
- Parameterization on types (polymorphism)

It has also provided practice in using various Bluespec tools:

- Building and executing in Bluesim
- Generating Verilog, and executing in Verilog sim
- Using Makefiles
- Generating waveforms and viewing them
- BDW workstation

End

```

import PFCtrl;

typedef BitN(28) DataT;

module ex_hdl_ctrl_fsm_top;

  Integer ffa_depth = 32;

  function BitN(28) determine_pump(DataT val);
    return (val[0]);
  endfunction

  PFCtrl(DataT) inboud;
  addSendPFCtrl(ffa_depth) ffa_inboud(inboud);
  PFCtrl(DataT) outboud;
  addSendPFCtrl(ffa_depth) ffa_outboud(outboud);
  PFCtrl(DataT) outboud;
  addSendPFCtrl(ffa_depth) ffa_outboud(outboud);

  rule end (True);
    DataT in_data = inboud.first;
    PFCtrl(DataT) out_data =
      determine_pump(in_data) == 0 ? outboud : outboud;
    out_data.send(ffa_outboud);
  endrule;
endmodule : ex_hdl_ctrl_fsm

```

