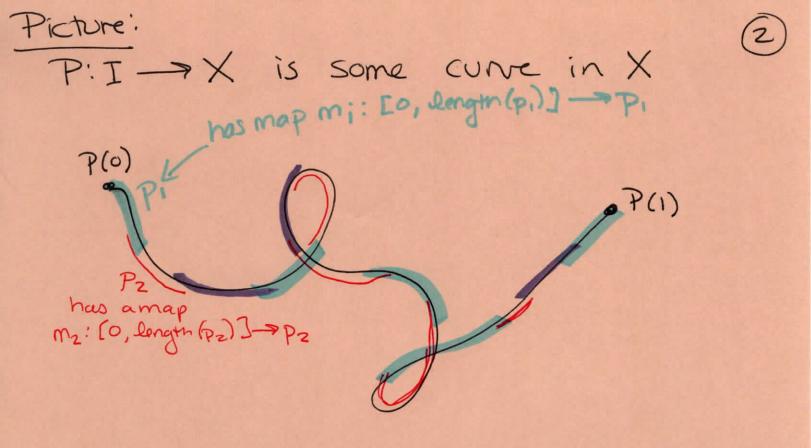
A fun inductive problem's Gen probi. comparing graphs embedded in R<sup>2</sup> there comes a new idea: lenghth-presenting (fréchet) dist. Given two metric spaces (X,dx), (Y,dy), we say that a continuous for 1: (x,dx) -> (Y,dy) is "length-preserving" iff \Xx,y \xx 1 dx (x,y) = dy (f(x), f(y)) e.g., 8: I - R2 be a curre parametrized by arclingth.

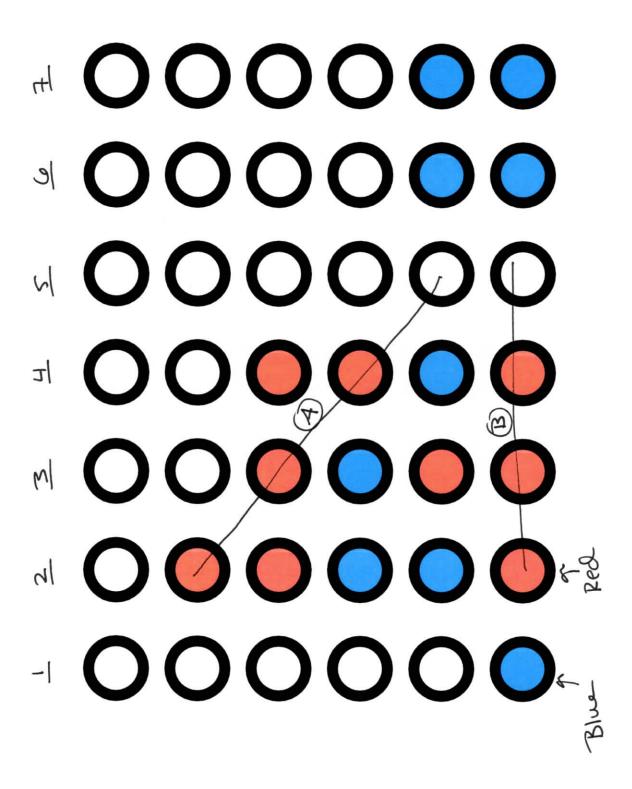
8(0)
8(t) + 8 has traveled distance t.

Claim: Let P:I -> (X,dx) be a path in (X,dx) such that I at cover of P town where Zpi CPZ where I maps m;: [a, a;] -> Pi that are length-presenting Bijections.

Then, I a reparametrization X:I -> TopI such mat Then, The reparametrization X:I -> TopI such that are length-presenting.



Challenge: use induction to prove this





Unless you've seen this game before<sup>3</sup>, you probably don't have any idea how to play it well. Nevertheless, there is a relatively simple backtracking algorithm that can play this game—or any two-player game without randomness or hidden information that ends after a finite number of moves—*perfectly*. That is, if we drop you into the middle of a game, and it is *possible* to win against another perfect player, the algorithm will tell you how to win.

A **state** of the game consists of the locations of all the pieces and the identity of the current player. These states can be connected into a *game tree*, which has an edge from state x to state y if and only if the current player in state x can legally move to state y. The root of the game tree is the initial position of the game, and every path from the root to a leaf is a complete game.

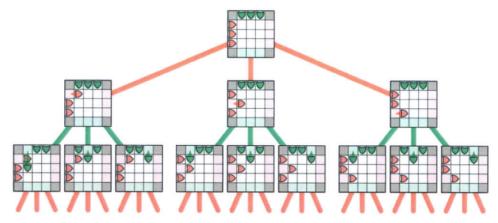


Figure 2.5. The first two levels of the fake-sugar-packet game tree.

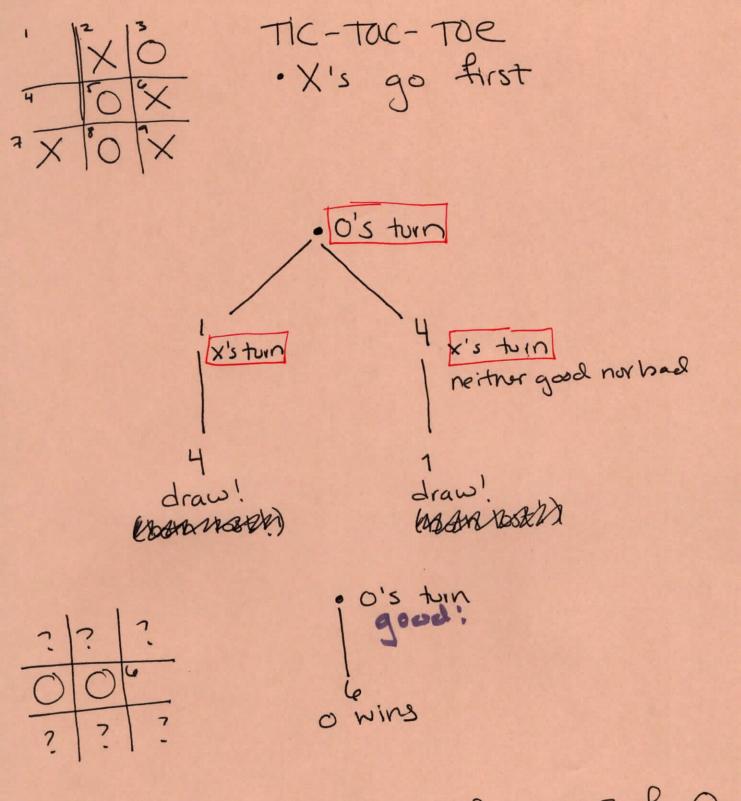
To navigate through this game tree, we recursively define a game state to be **good** or **bad** as follows:

- A game state is good if either the current player has already won, or if the current player can move to a bad state for the opposing player.
- A game state is bad if either the current player has already lost, or if every available move leads to a good state for the opposing player.

Equivalently, a non-leaf node in the game tree is good if it has at least one bad child, and a non-leaf node is bad if all its children are good. By induction, any player that finds the game in a good state on their turn can win the game, even if their opponent plays perfectly; on the other hand, starting from a bad state, a player can win only if their opponent makes a mistake. This recursive definition was proposed by Ernst Zermelo in 1913.<sup>4</sup>

<sup>3</sup>If you have, please tell me where!

<sup>&</sup>lt;sup>4</sup>In fact, Zermelo considered the more subtle class of games that have a finite number of states, but that allow infinite sequences of moves. (Zermelo defined infinite play to be a draw.)



In groups, play 3 moves for X+3 for O then draw the game tree (list all possibilities!) + marke if each node in the game tree is good or bod (or neither) —> leaf rodes can be considered as toin.

## cont:



- 1) If d draw means X mins, what the good/bad states of your gome tree?
- 2) If a draw means O wirs, what are the good / bad states?

## Notes:

- → When encountering new del'n/algo, try it out on small examples + play w it to get a feel for it.
- drawing trees choose labels wirely in general, what labeling conventions can you establish to show what you are doing clearly