

1 November 2021

an MST / the MST

↑  
recognize it might  
not be unique

↑ there is only one

Why is it ok to assume  $\exists!$  MST?

→ If we fix an (arbitrary) order on the vertices,  
then we can break ties using the product  
order, ~~order~~

So, this breaks all ties  $\Rightarrow$  a unique poset of  
edges, all with "different" weights.

→ See the beg. of Ch 7 for a full explanation

MST problem:

- given: weighted graph  $G = (V, E, \omega)$
- want:  $T \subseteq G$  such that  $T$  is a tree  
w/ vertex set  $V$  and among all such trees,  
 $T$  has minimum weight, where the weight of  
a tree is:  $\omega(T) := \sum_{e \in E_T} \omega(e)$

many sol'n's build  $T$  iteratively + a given pt  
have a spanning forest  $F \subseteq T$ . Here, note  $\nexists$  has  
3 types of edges:

- ① useless: edge that has 2 endpts in  
same component. (Lem 7.3)
- ② safe: for each cc (a tree), look at all outgoing  
edges. The min weight one is safe. (Lem. 7.2)
- ③ undecided

①

3 spins on the algorithm  
 $F = (V, \emptyset)$

① Kruskal's Algorithm

Sort edges by increasing weight. Scan through  
+ add safe edges as we encounter them.

② Jarník / Prim Algorithm

repeatedly add safe edges. (one at a time)

③ Borůvka's Algorithm

add all safe edges + repeat

In groups,  $\rightarrow$  walk through ② and ③.

$\rightarrow$  how much better is ③? How do  
we even describe this?

Initially,  $F = (V, \emptyset)$

what are our safe edges?  $S = \{1, 2, 3, 4, 5\}$

useless edges?  $\emptyset$

unsure edges:  $E \setminus S$

(after Prim: add edge 3

safe edges:  $S = \{1, 2, 4, 5\}$

useless:  $\emptyset$

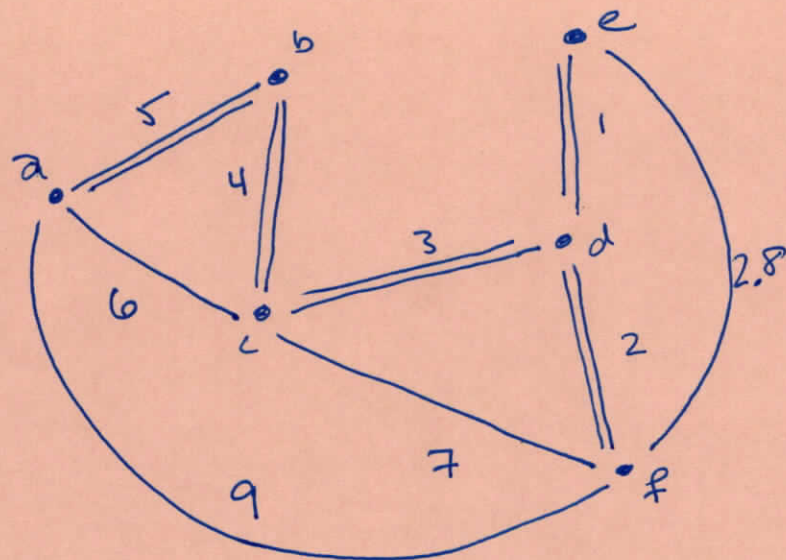
unsure:  $E \setminus (S \cup E_F)$

What can happen in general? New safe edges would  
be out of the cc containing the new edge!

②



Example:



① Kruskal

edges, sorted:  $[1, 2, 3, 4, 5, 6, 7, 9]$   
 $\uparrow \uparrow \uparrow$

init  $F = (V, \emptyset)$

$a(5)$      $b(4)$      $c(3)$      $d(1)$      $e(1)$      $f(2)$   
 $\bullet$      $\bullet$      $\bullet$      $\odot$      $\odot$      $\bullet$

1st: edge 1 connects  $[e]$  w/  $[d]$   
 $\textcircled{1} 8$      $\textcircled{1, 2, 3}$   
 $\checkmark$  1 is safe     $\checkmark$  1 is safe

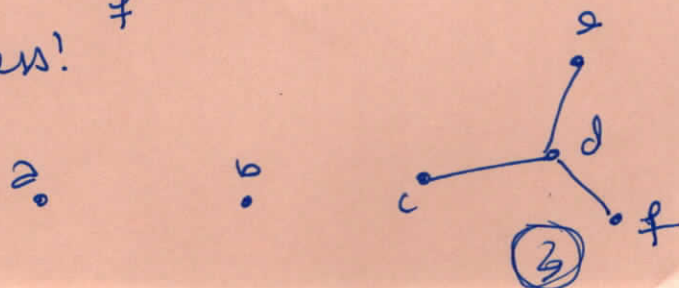
$F$ :  $a(5)$      $b(4)$      $c(3)$      $d$      $e$      $f(2)$   
 $\bullet$      $\bullet$      $\bullet$      $\bullet$      $\bullet$      $\bullet$

Next in scan, look at edge 2.

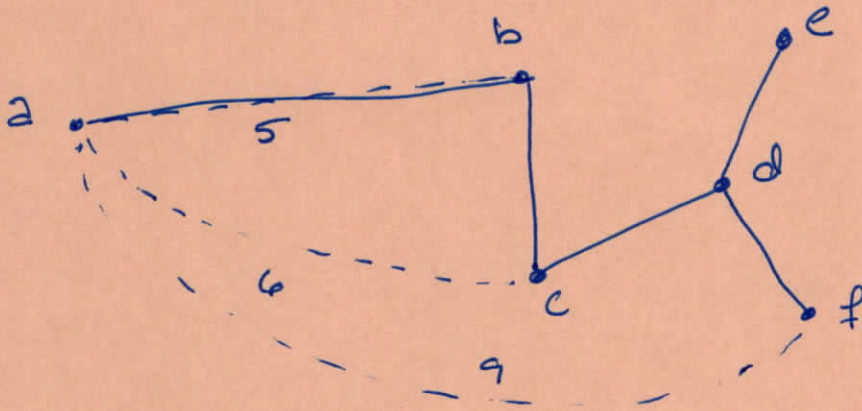
Safe, so add to ~~MST~~  $F$



Next up: edge 2.8 is useless!  
 edge 3: safe, so add it



Next: edge 4 is safe, so F becomes:



Next: edge 5 is safe, so add it

Next: edge 6 useless  
7 useless  
9 useless

} once I've established my spanning tree, everything else is useless.