Replicated Block Storage

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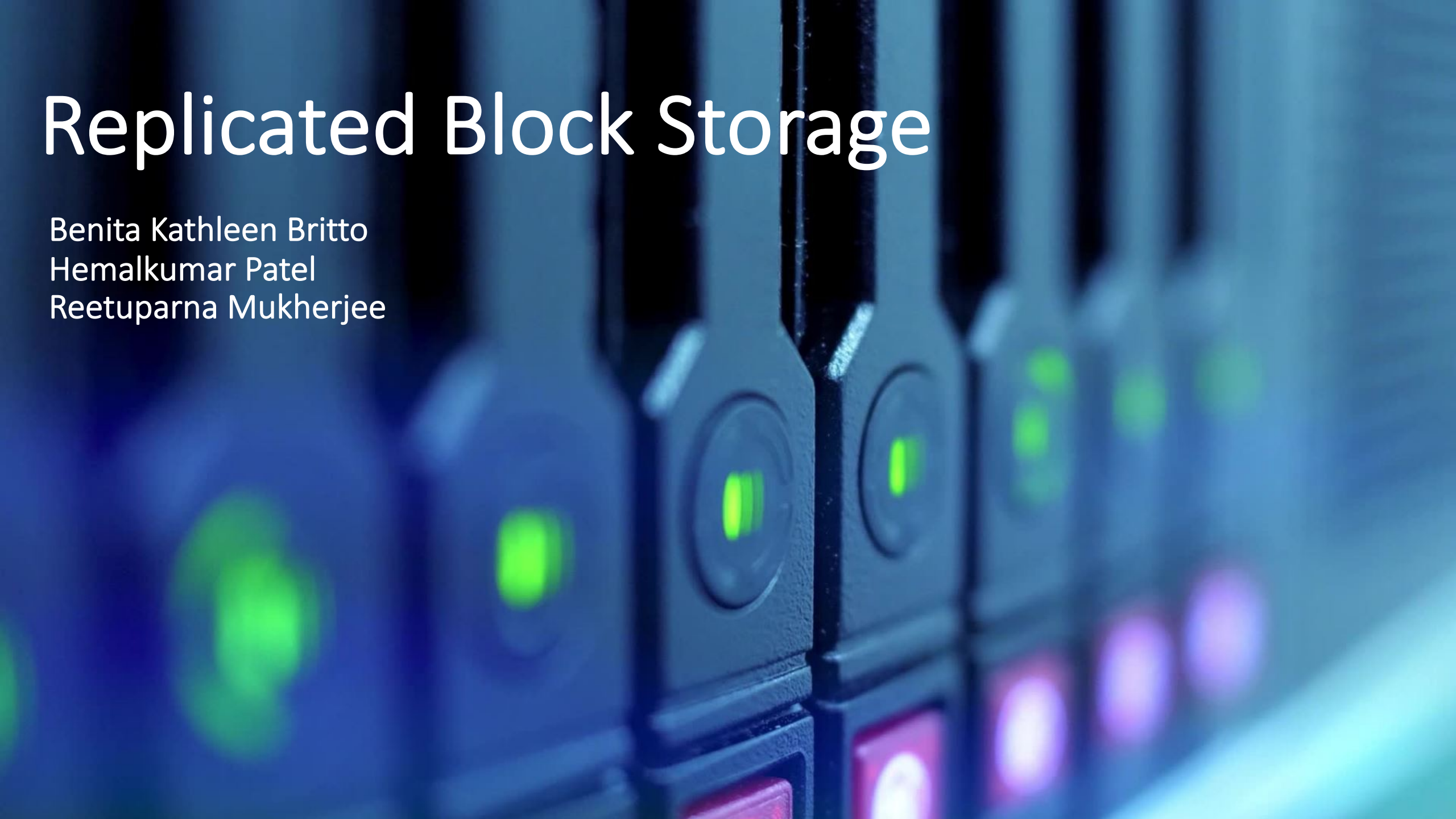
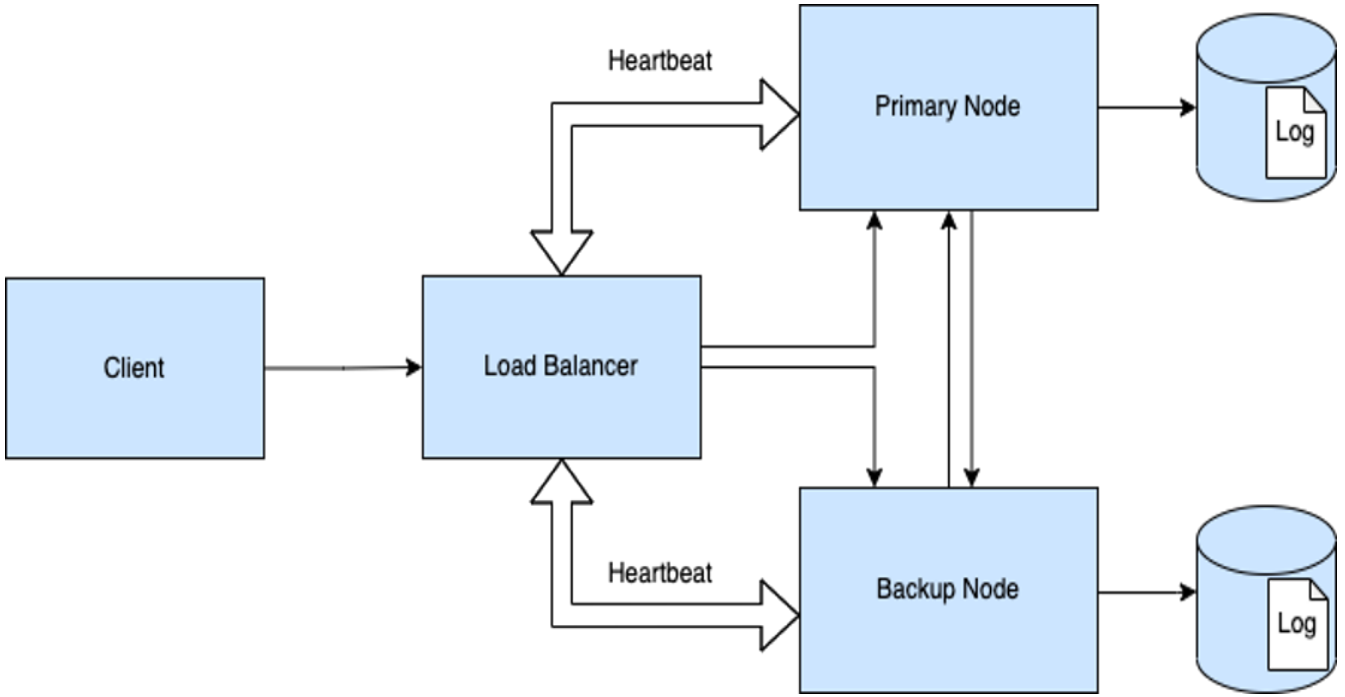


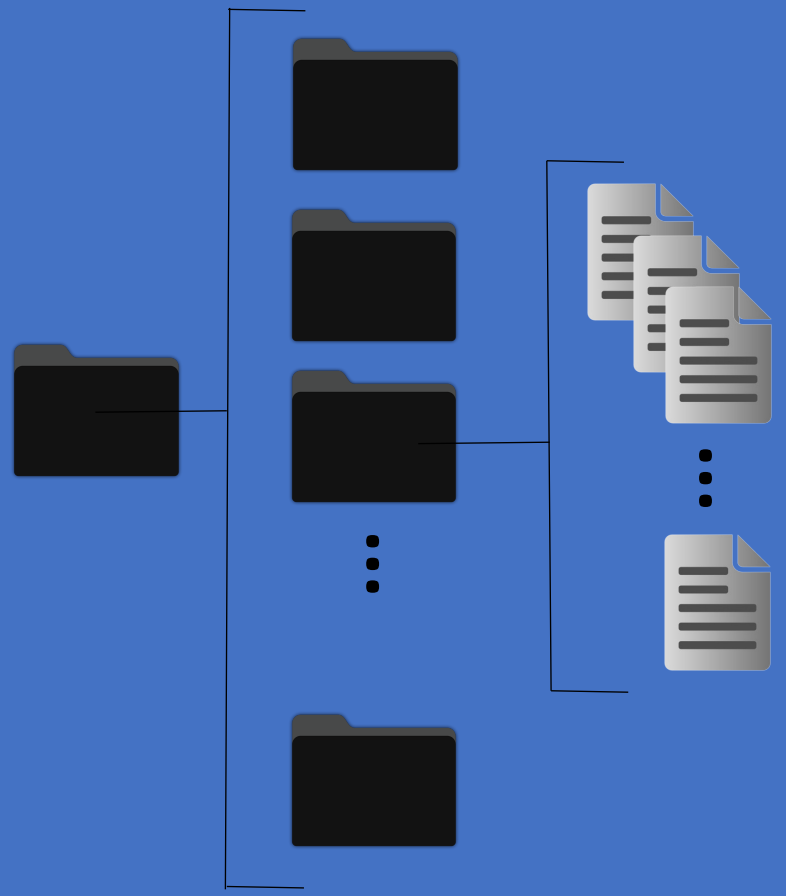
Key Assumptions

- Non-Byzantine Failures
- Stable Storage – No disk corruption
- Stable Network – No packet drop
  - No response means node has crashed
- One Node Always Up
  - Backup becomes primary
  - No failures during crash recovery

System Architecture



Storage Structure



Block Storage Size: 1GB

1024 directories, each having 256 files of 4KB size

Easy scale up/down (not included in this project)

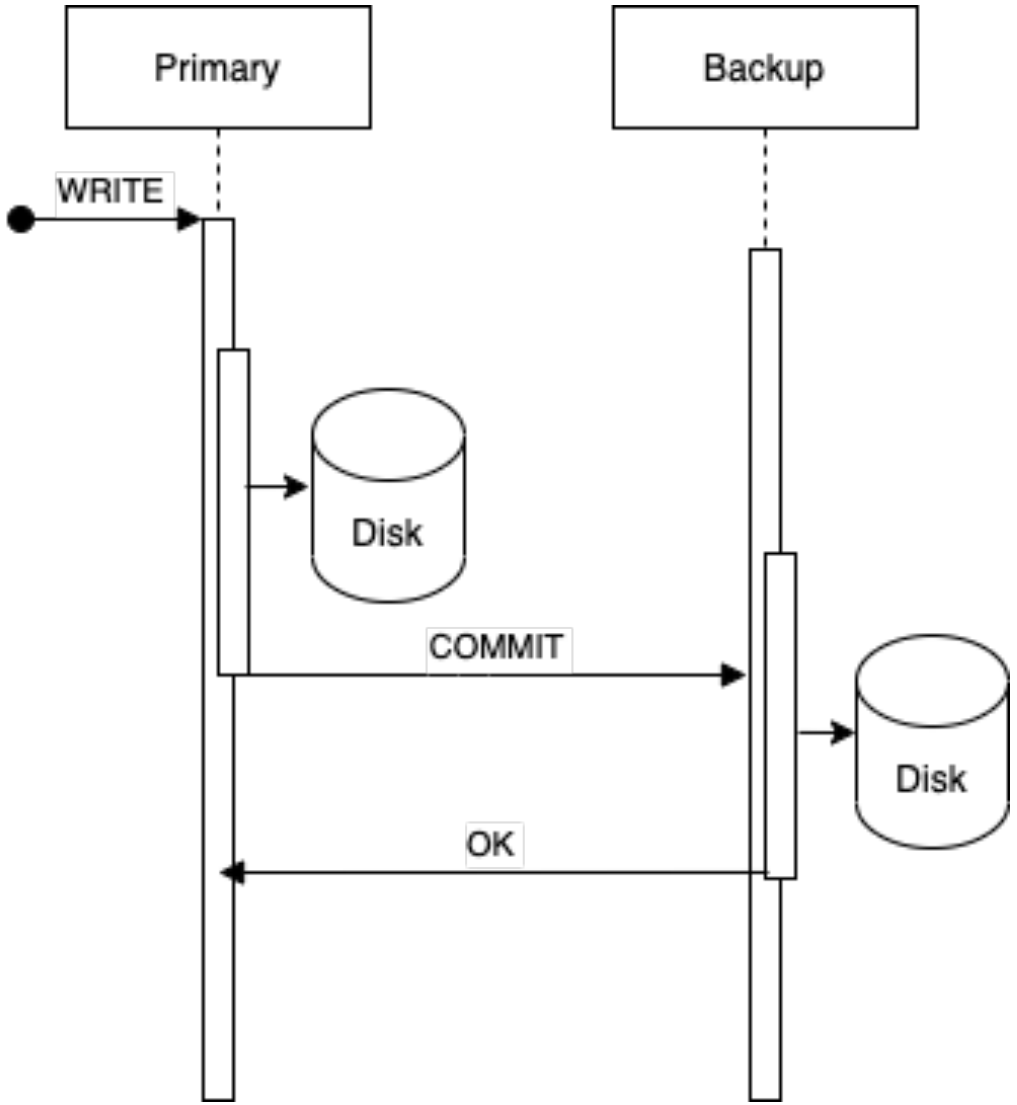
Can partition data across storage (not included in this project)

Replication Strategy

Approach #1: Primary Commits First, Backup Second

On write():

1. Primary writes (commits) data first
2. Primary contacts backup to write
3. Backup writes data



Approach #1: Primary Commits First, Backup Second

Problem:

What happens if the

primary crashes before contacting

the backup?

- Backup has no knowledge of

the request

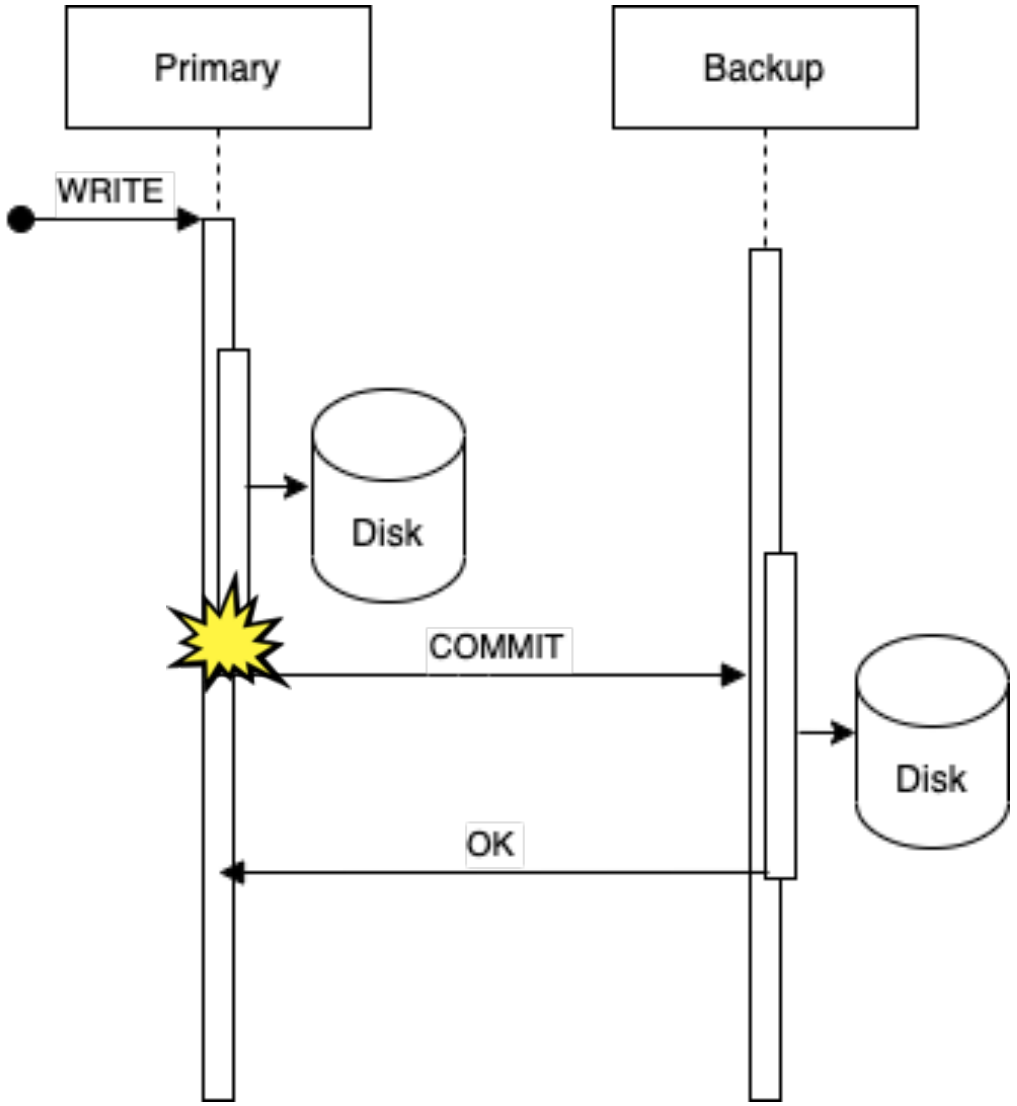
- Client will see the old data and assume that the request failed

Solution:

When primary reboots,

it must reconcile with backup for

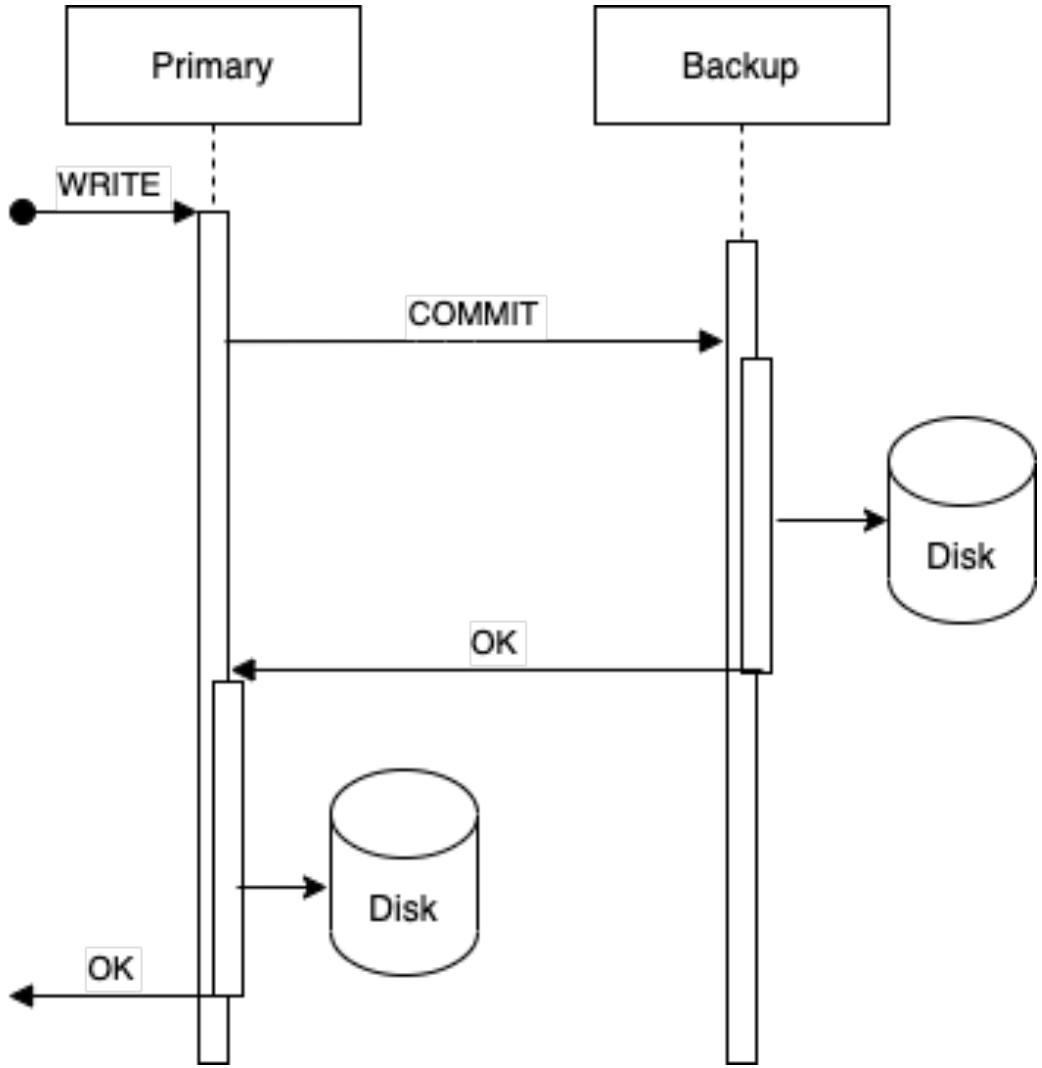
all requests – Expensive!



Approach #2: Backup Commits First, Primary Second

On write():

1. Primary contacts backup to write
2. Backup writes data
3. Primary writes (commits) data on receiving an OK



Approach #2: Backup Commits First, Primary Second

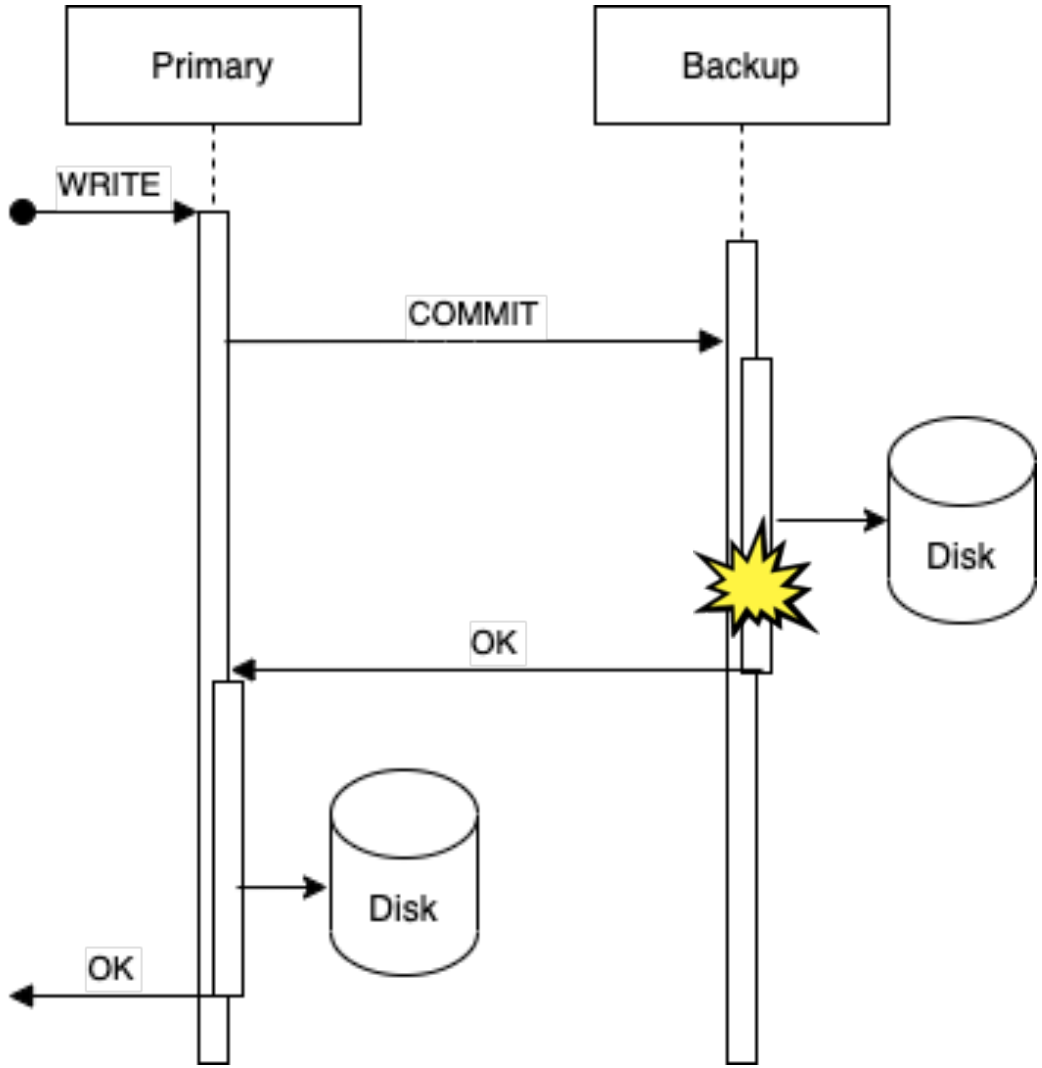
Problem:

What happens if backup crashes before responding to the primary?

- Primary can commit/abort the txn

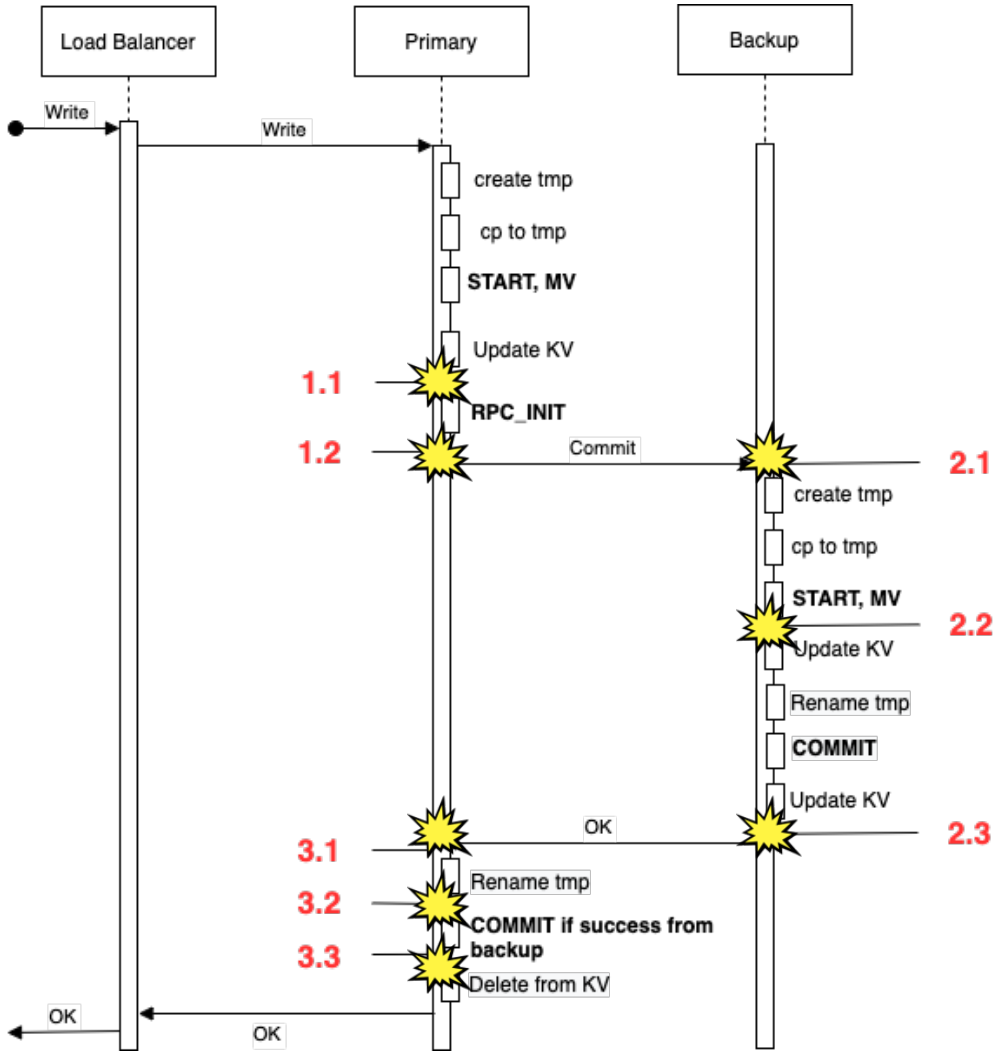
Solution:

When backup reboots,
it must reconcile with primary
for all txns – Expensive!



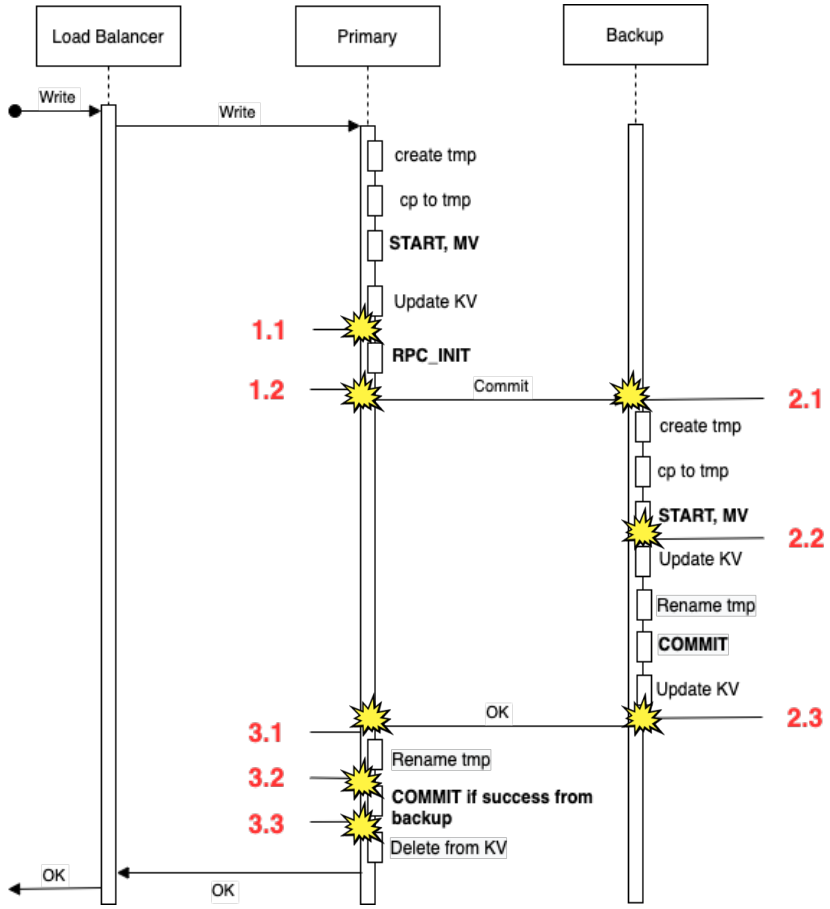
Approach #3 (Our Design): Primary Prepares, Backup Commits, Primary Commits

Crash	Primary	Backup	Client
1.1	Reconcile the txn status Fetches all pending txns from primary on reboot	Assumes Primary role	Retry
1.2			Retry



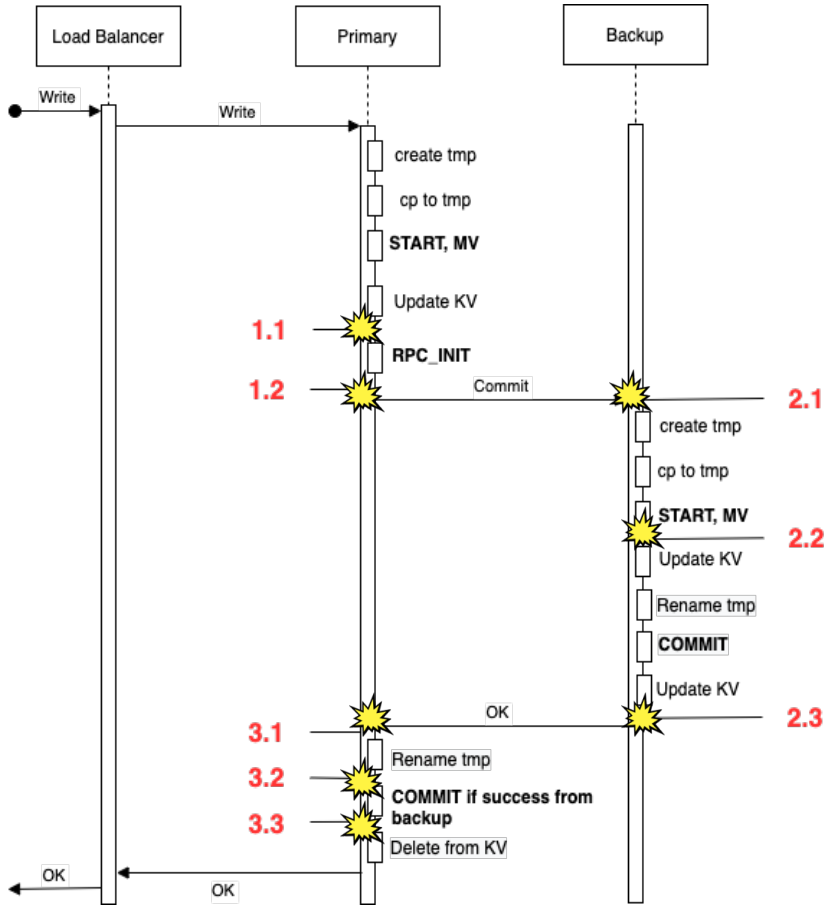
Approach #3 (Our Design): Primary Prepares, Backup Commits, Primary Commits

Crash	Primary	Backup	Client
2.1	Logs replication pending Keep txn in memory Commits txn	GetState(txnId) from primary on reboot and applies changes Apply pending txns	-
2.2			
2.3		-	



Approach #3 (Our Design): Primary Prepares, Backup Commits, Primary Commits

Crash	Primary	Backup	Client
3.1	GetState(txnId) from primary on reboot and applies changes Apply pending txns	-	-
3.2			
3.3	-	-	-



System Behavior on Crash

		Scenario	Behavior
Case- 1		When both nodes are up and no failures	• Reads: Primary/Backup • Writes: Only Primary
Case-2	i	Backup crashed	Reads/Writes: Only Primary
	ii	Primary crashed	Backup->Primary, Reads/Writes: Only Primary
Case - 3		If one node is alive and the other is recovering	Ongoing write requests are sent to the live node and future writes are rejected until the recovering node is stable, reads are sent to the live node and finally, we go back to Case #1/#2 once we reach a stable state for the recovering node. We assume that there are no failures during the recovery process

Design Decisions

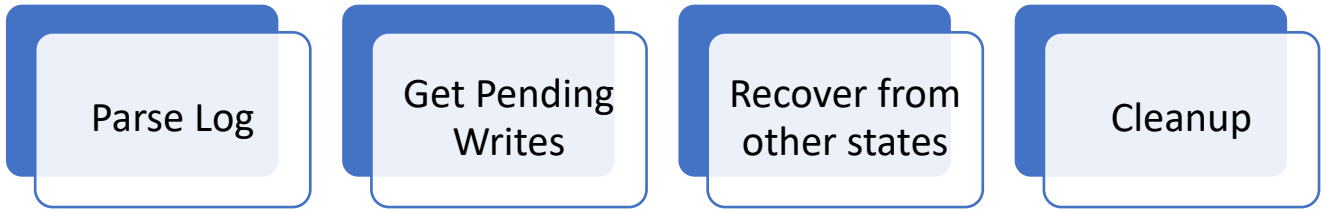
Retry mechanism on client

Lock file during writes

RPC status UNAVAILABLE is used to determine crash

Primary stops processing writes when backup is recovering

Crash Recovery Protocol



Durability (ALICE)

Write(Addr)

…

creat(tmp)

write(tmp)

fsync(tmp)

…

[rename(tmp,file)]

…

Testing Strategy

CRC for comparing block content

Macros for predefined crash points

system(kill) and dereferencing NULL to simulate a crash

Measurements

System Specifications

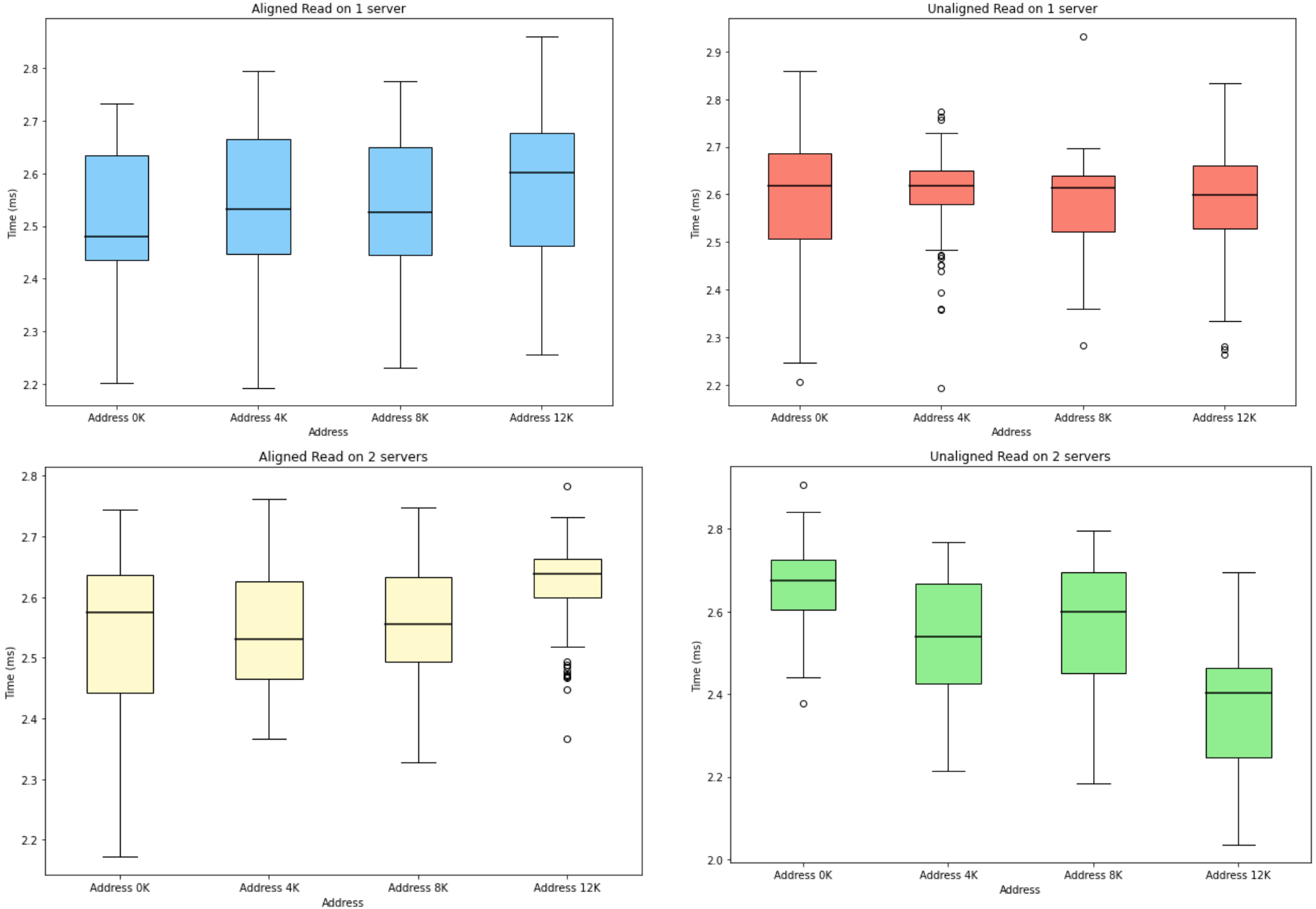
Cloud Provider: CloudLab

Architecture: x86_64

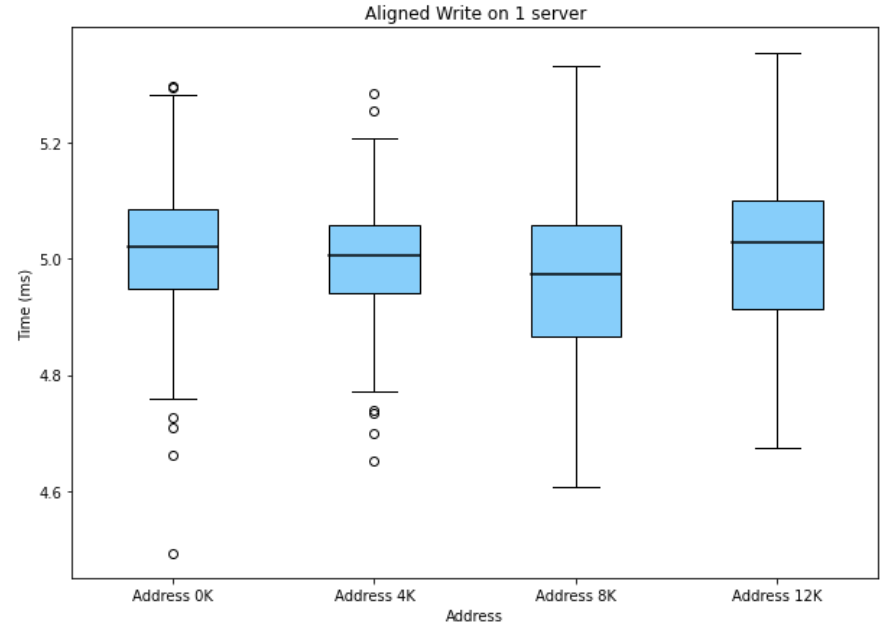
CPU(s): 40

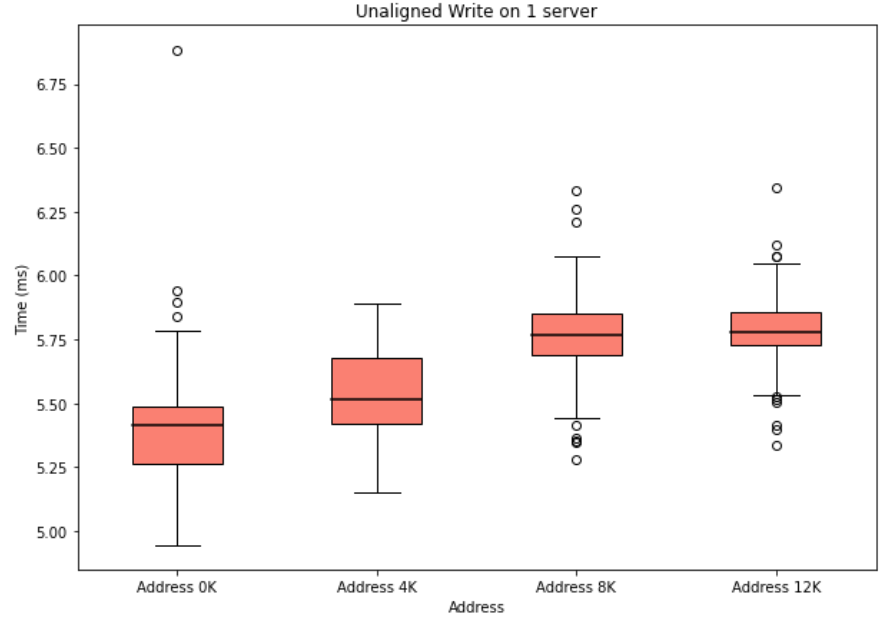
OS: Ubuntu 18 LTS

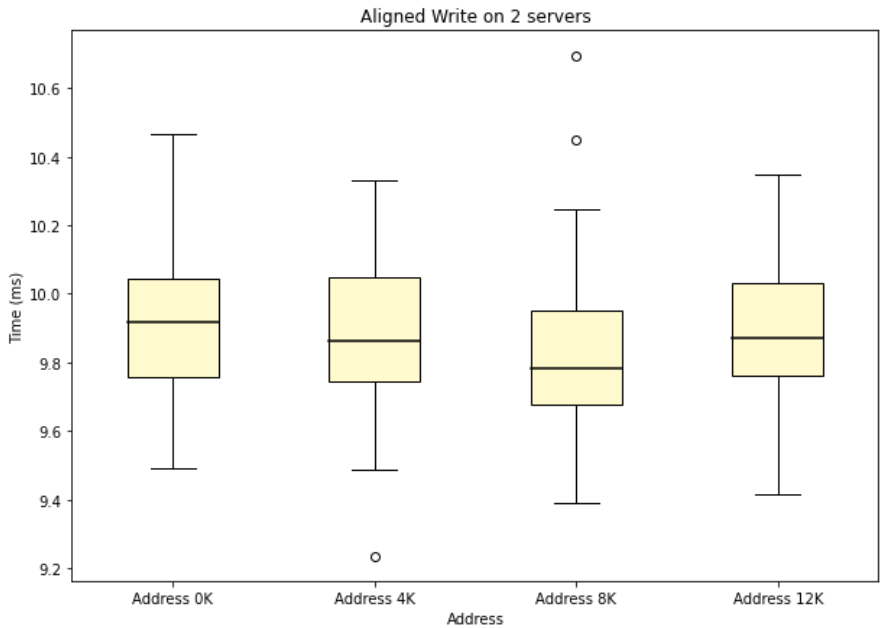
Read Latency without Crash

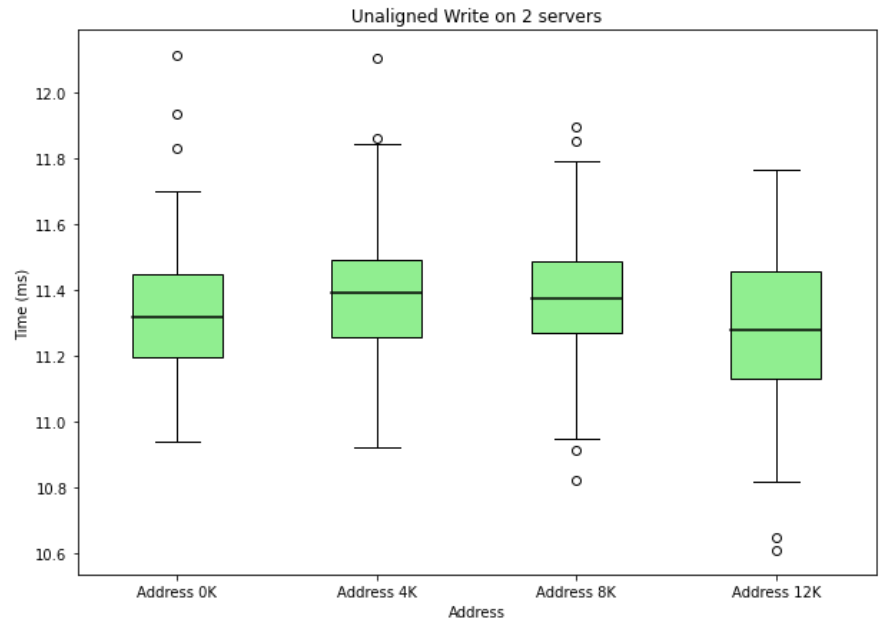

Aligned Read on 1 server
Time (ms)
2.8
2.7
2.6
2.5
2.4
2.3
2.2
Address 0K
Address 4K
Address 8K
Address 12K
Address
Unaligned Read on 1 server
Time (ms)
2.9
2.8
2.7
2.6
2.5
2.4
2.3
2.2
Address 0K
Address 4K
Address 8K
Address 12K
Address
Aligned Read on 2 servers
Time (ms)
2.8
2.7
2.6
2.5
2.4
2.3
2.2
Address 0K
Address 4K
Address 8K
Address 12K
Address
Unaligned Read on 2 servers
Time (ms)
2.8
2.6
2.4
2.2
2.0
Address 0K
Address 4K
Address 8K
Address 12K
Address
1S - Aligned
1S - Unaligned
2S - Aligned
2S - Unaligned

Write Latency without Crash


Aligned Write on 1 server
Time (ms)
5.2
5.0
4.8
4.6
Address 0K
Address 4K
Address 8K
Address 12K
Address


Unaligned Write on 1 server
Time (ms)
6.75
6.50
6.25
6.00
5.75
5.50
5.25
5.00
Address 0K
Address 4K
Address 8K
Address 12K
Address


Aligned Write on 2 servers
Time (ms)
10.6
10.4
10.2
10.0
9.8
9.6
9.4
9.2
Address 0K
Address 4K
Address 8K
Address 12K
Address

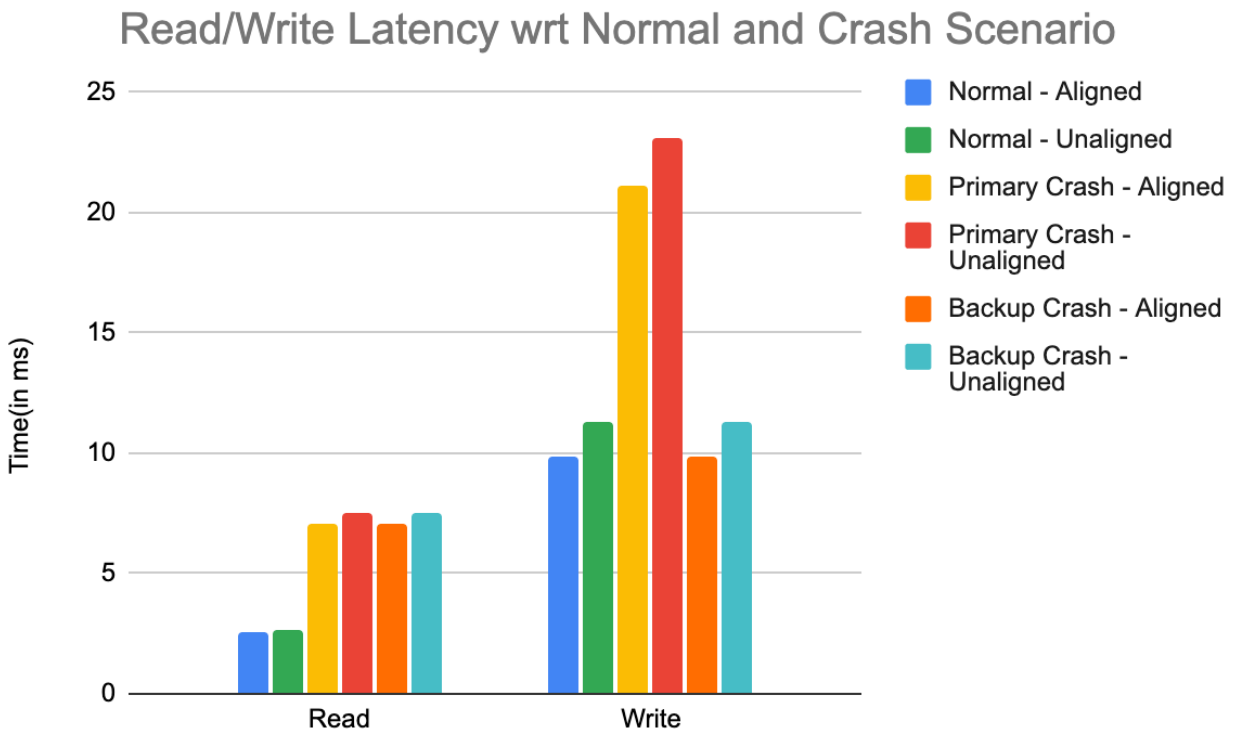

Unaligned Write on 2 servers
Time (ms)
12.0
11.8
11.6
11.4
11.2
11.0
10.8
10.6
Address 0K
Address 4K
Address 8K
Address 12K
Address


1S - Aligned
1S - Unaligned
2S - Aligned
2S - Unaligned

Read and Write Latency with Crash

Key Observations:

- Normal = 2 servers up
- For read, latency for primary crash is almost double that of reads without crash, plus 1 exponential backoff
- For writes, latency for primary crash is (two server + one server + heartbeat latency + 1 exponential backoff
- Backup crashes behave the same as normal



Configuration for Concurrent Operations

Spin up X number of worker threads. Each threads work on different directories.

Each worker thread starts from specific offset in that directory and jumps within that directory uniformly.

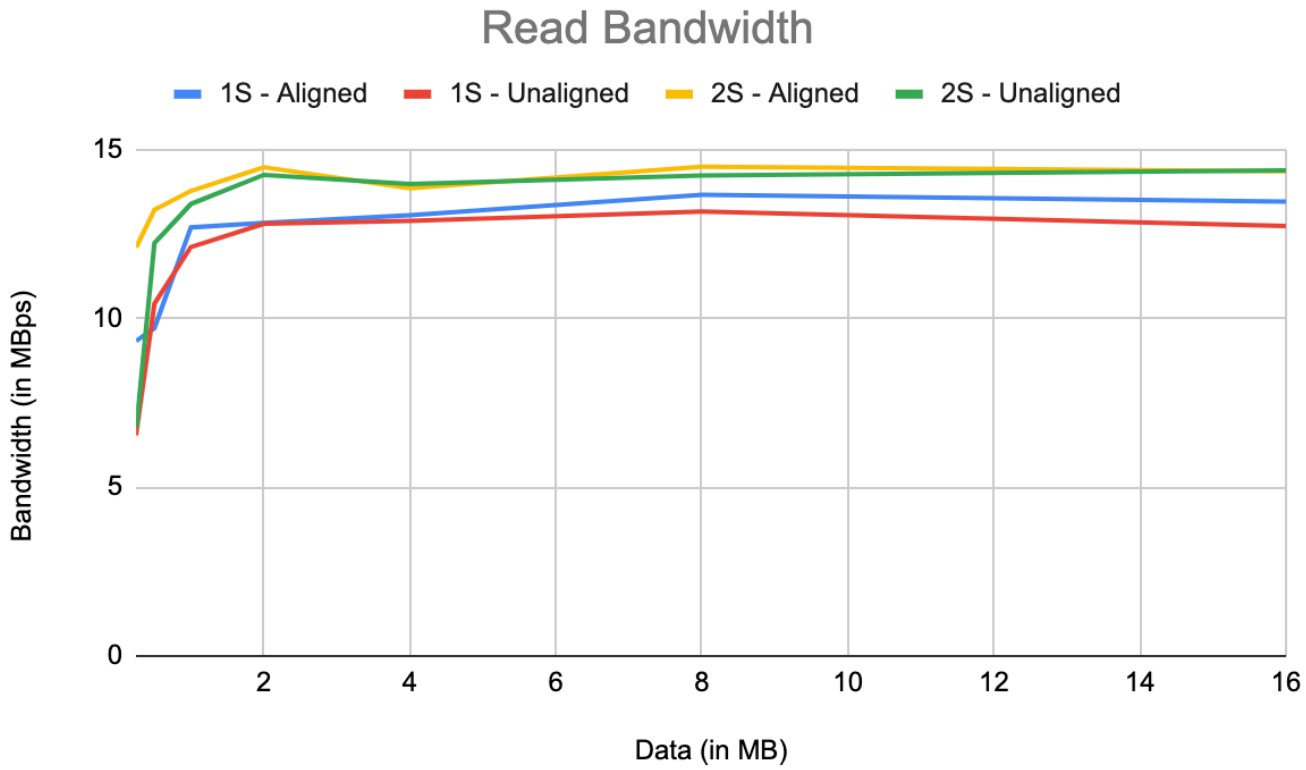
Aligned Address: Starts from 0th offset

Unaligned Address: Starts from (4K+2K)th offset

Read Scalability

Key Observations:

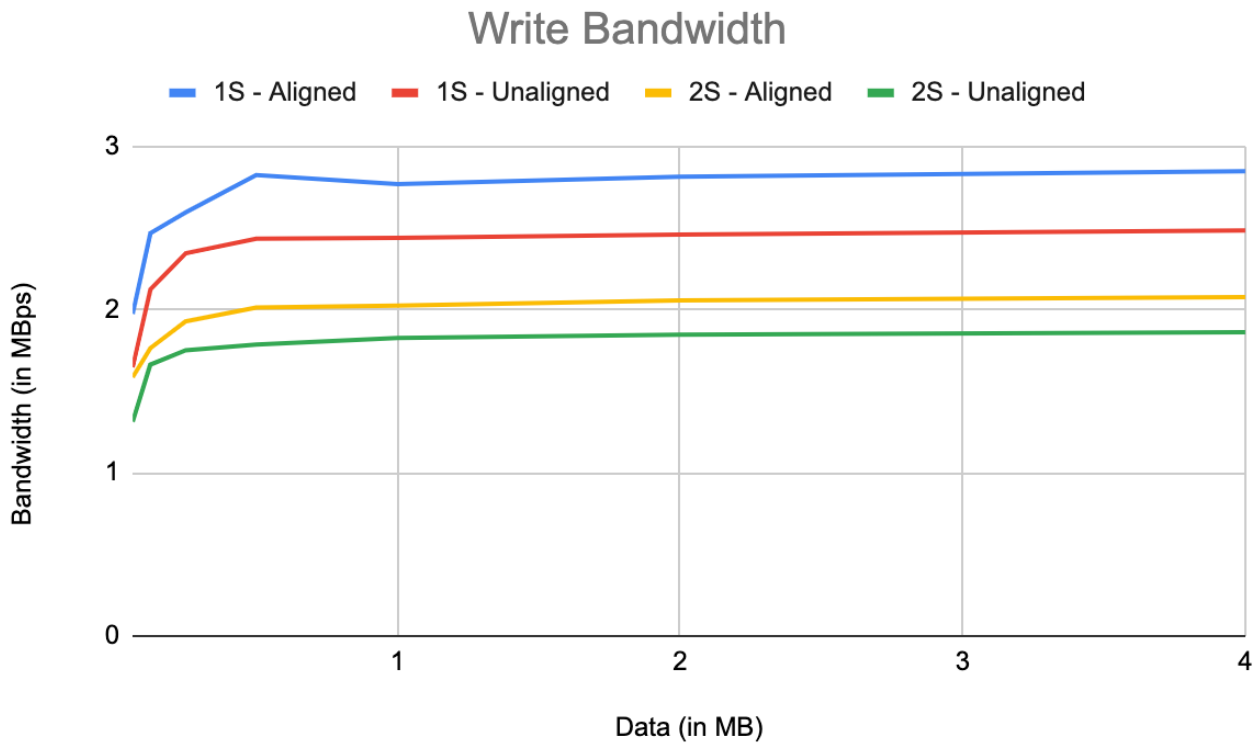
- For 1 server, Aligned vs Unaligned difference is more noticeable when number of request increases (Reason: 2x file reads per request).
- For 2 Servers, Aligned vs Unaligned difference decreases when number of request increases (Reason: Load Balancer + load distribution over 2 servers).
- Bandwidth: 2 Server > 1 Server; Aligned > Unaligned



Write Scalability

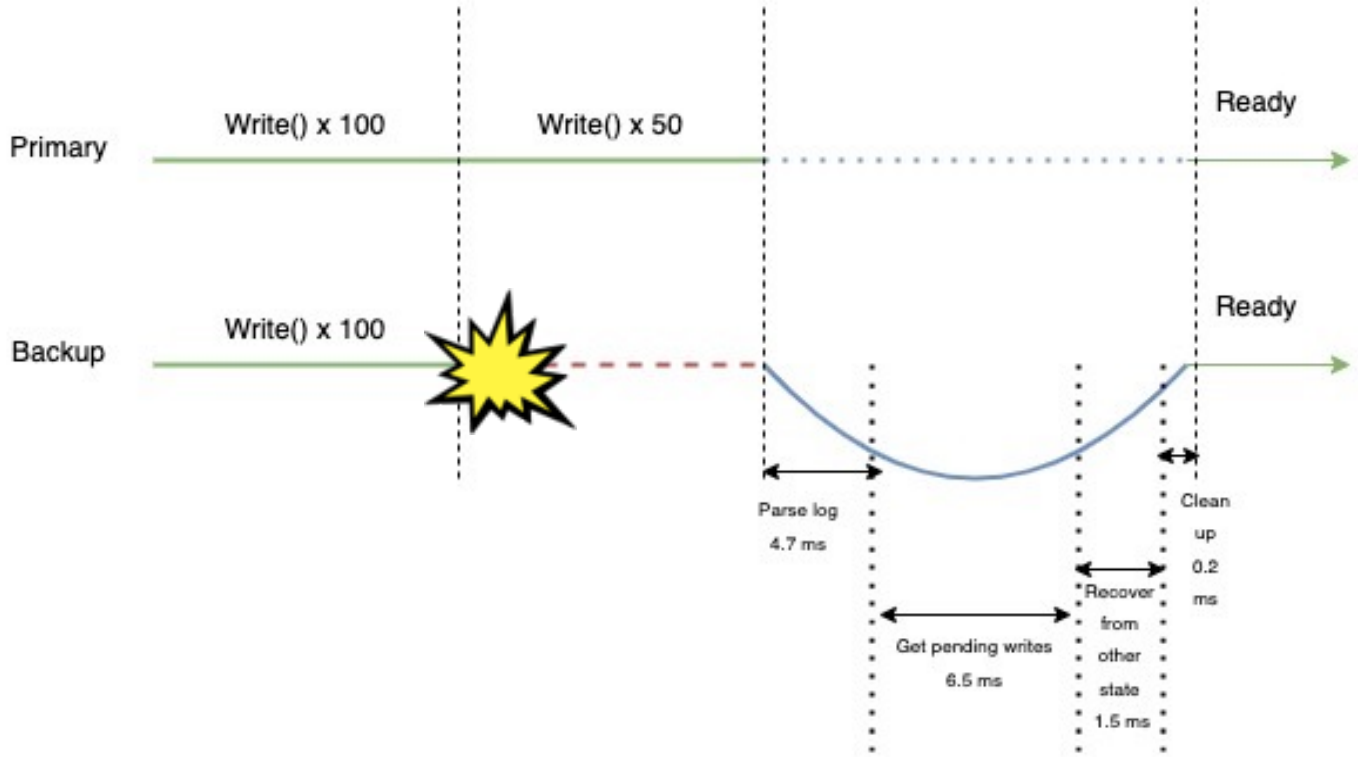
Key Observations:

- Bandwidth: 1 Server > 2 Server and Aligned > Unaligned



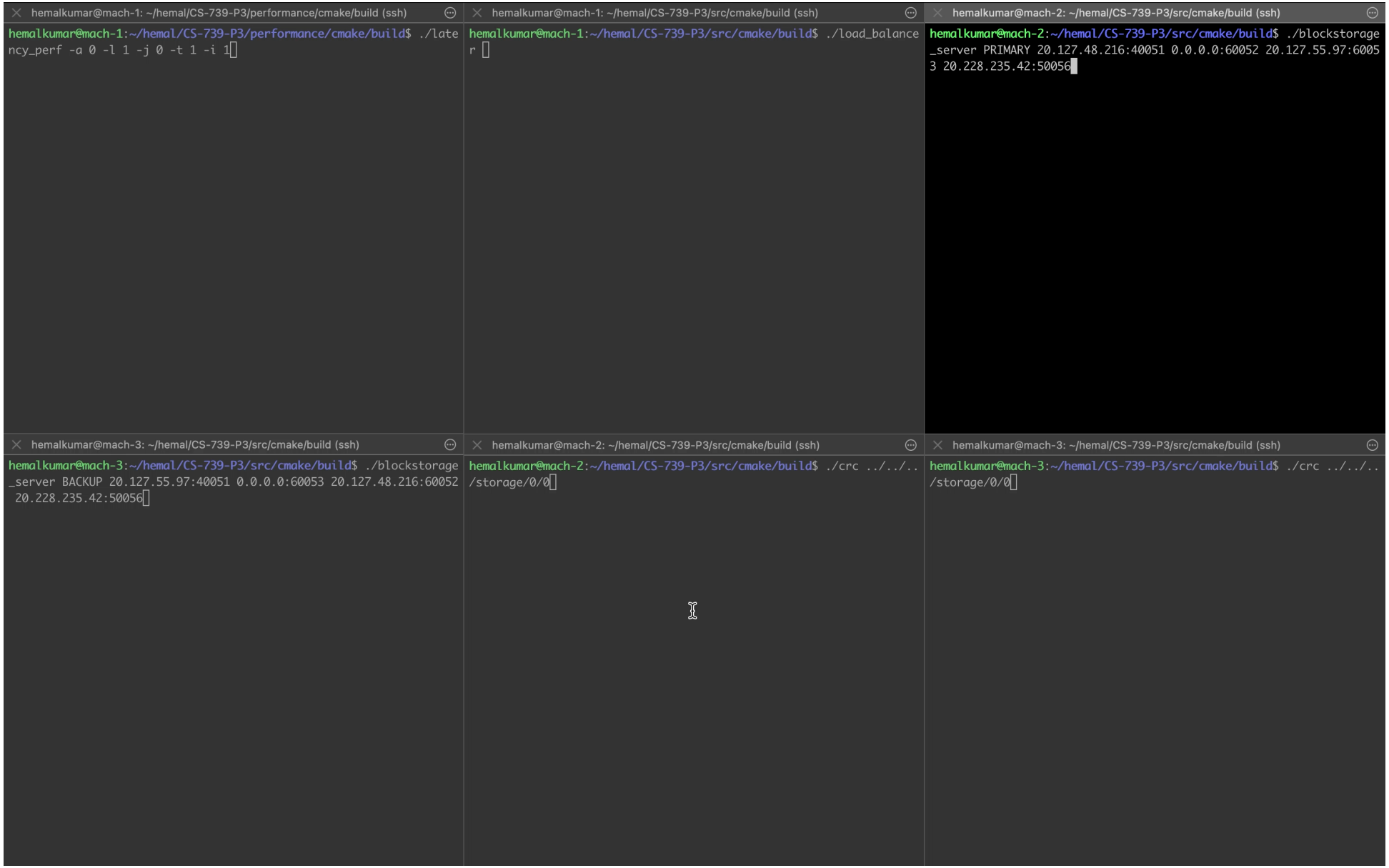
Recovery Time

- 4.7 ms for parsing 100 writes
- 6.5 ms for 50 pending writes
- 1.5 ms for recovering from other states i.e. looping through the 100 states from the 100 writes when the backup was alive
- 0.2 ms to delete the in-memory map of the log and to truncate the log file

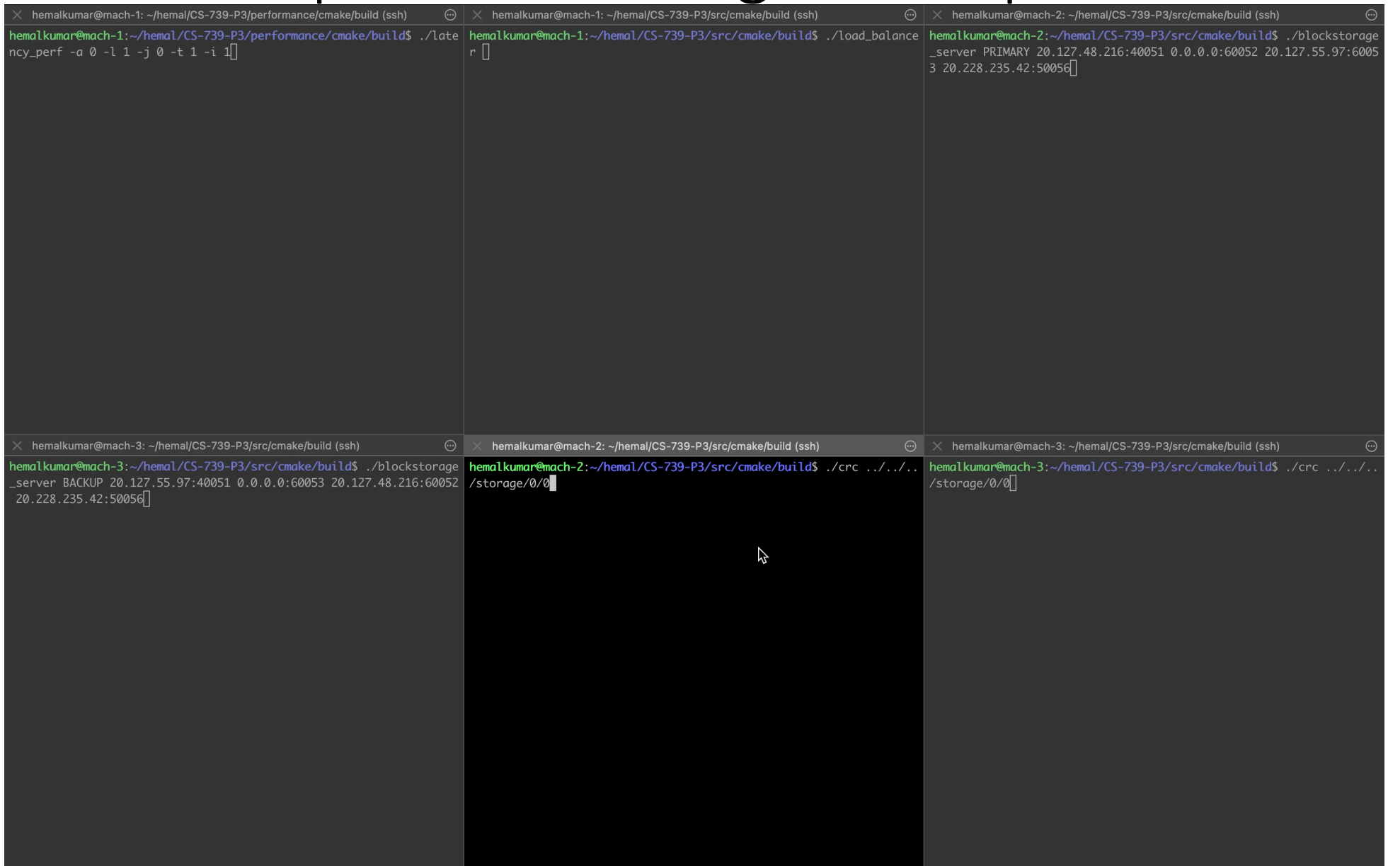


Demos

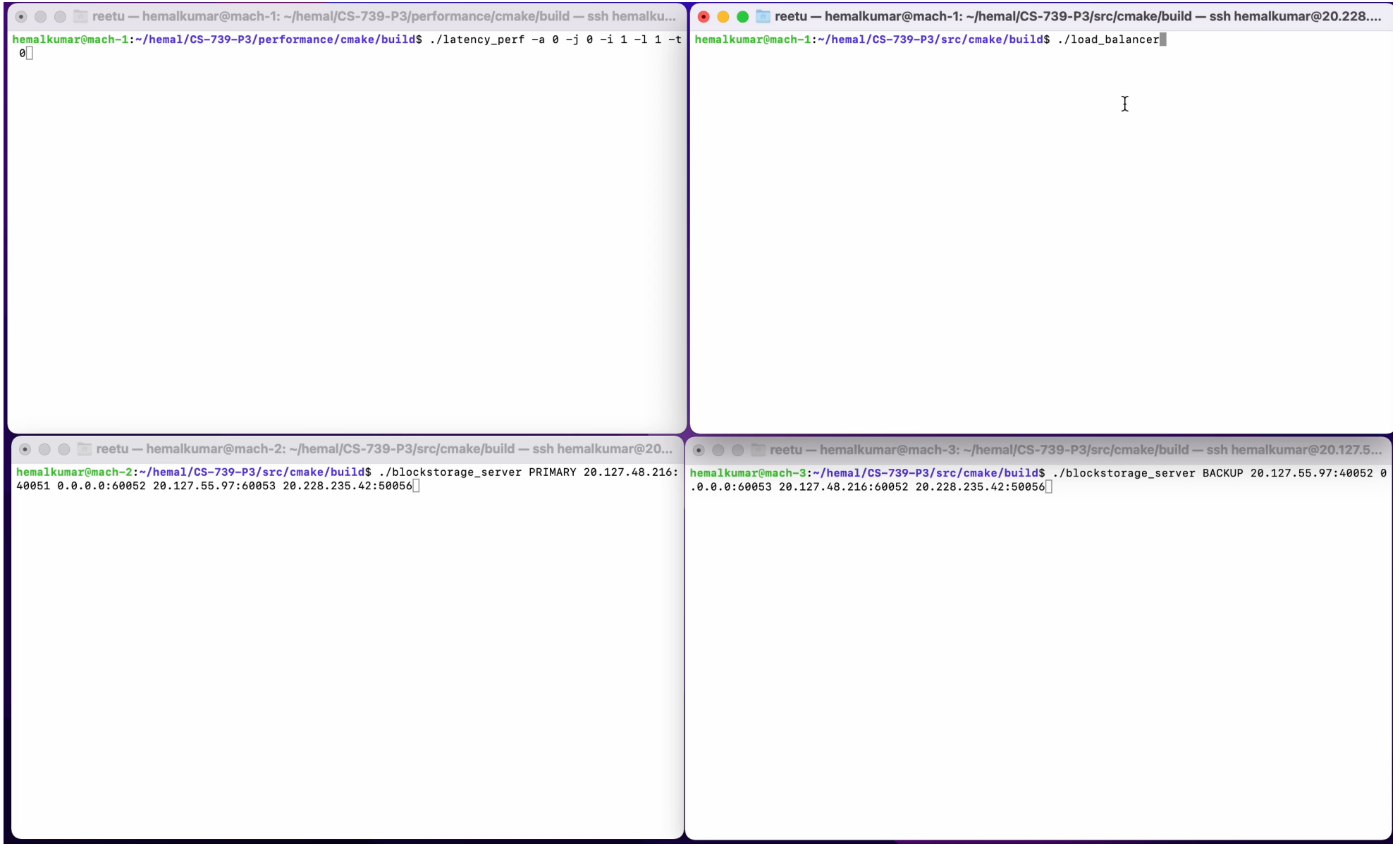
Failover


hemalkumar@mach-1: ~/hemal/CS-739-P3/performance/cmake/build (ssh)
hemalkumar@mach-1:~/hemal/CS-739-P3/performance/cmake/build$ ./latency_perf -a 0 -l 1 -j 0 -t 1 -i 1
hemalkumar@mach-1: ~/hemal/CS-739-P3/src/cmake/build (ssh)
hemalkumar@mach-1:~/hemal/CS-739-P3/src/cmake/build$ ./load_balancer
hemalkumar@mach-2: ~/hemal/CS-739-P3/src/cmake/build (ssh)
hemalkumar@mach-2:~/hemal/CS-739-P3/src/cmake/build$ ./blockstorage_server PRIMARY 20.127.48.216:40051 0.0.0.0:60052 20.127.55.97:60053 20.228.235.42:50056
hemalkumar@mach-3: ~/hemal/CS-739-P3/src/cmake/build (ssh)
hemalkumar@mach-3:~/hemal/CS-739-P3/src/cmake/build$ ./blockstorage_server BACKUP 20.127.55.97:40051 0.0.0.0:60053 20.127.48.216:60052 20.228.235.42:50056
hemalkumar@mach-2: ~/hemal/CS-739-P3/src/cmake/build (ssh)
hemalkumar@mach-2:~/hemal/CS-739-P3/src/cmake/build$ ./crc ../../../storage/0/0
hemalkumar@mach-3: ~/hemal/CS-739-P3/src/cmake/build (ssh)
hemalkumar@mach-3:~/hemal/CS-739-P3/src/cmake/build$ ./crc ../../../storage/0/0

Backup failure during write operation

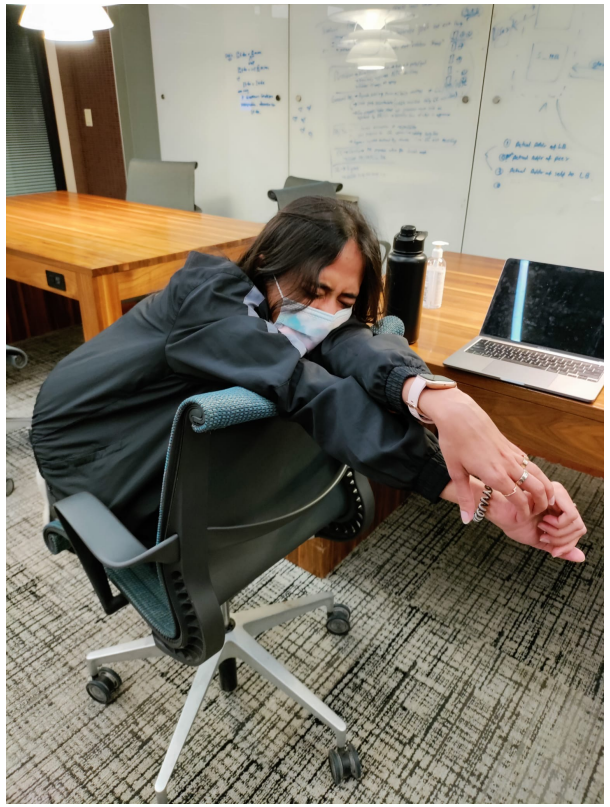


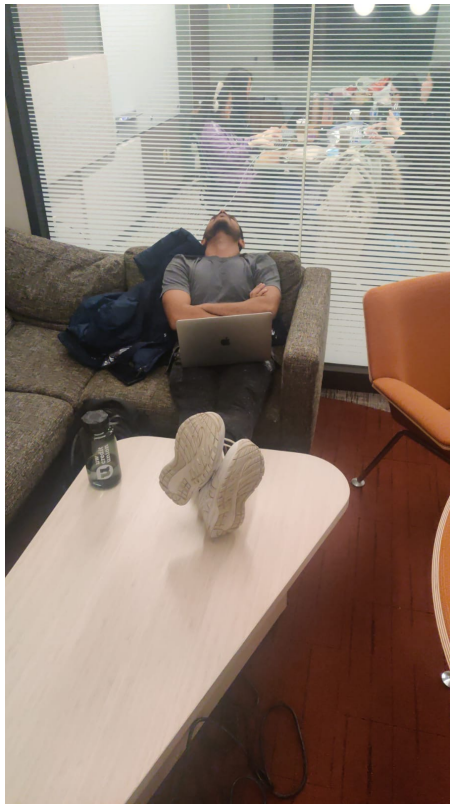
Read with crash



Thank You!

No request to serve...





Crash!!!!

Recovery