

ASSIGNMENT-1

1. Use Mathematical Induction to prove:-

For any natural number n , $n^3 - n$ is divisible by 3.

→ Let $n = 1$; $n^3 - n = 1^3 - 1 = 0 \Rightarrow$ divisible by 3 $= 3(0)$

$n = 2$; $n^3 - n = 2^3 - 2 = 8 - 2 = 6 \Rightarrow$ divisible by 3 $= 3(2)$

According to mathematical induction,

assume the above will be true for $n = k$

Let $n = k$; $n^3 - n = k^3 - k = 3(m)$... which is divisible by 3
↳ ①

Let $n = k + 1$, if $n = k$ is true then $n = k + 1$ will be true

$$(n^3 - n) = (k+1)^3 - (k+1) = k^3 + 3k^2 + 3k + 1 - k - 1$$

$$= k^3 + 3k^2 - 2k$$

$$= (k^3 - k) + 3(k^2 + k)$$

∴ Substituting $k^3 - k = 3m$ from ①,

$$= 3m + 3(k^2 + k)$$

$$= 3(m + k^2 + k)$$

..... which is divisible by 3.

∴ According to the above proof of mathematical induction,
 $n^3 - n$ is divisible by 3.

2. Write algorithm of calculating Factorial by following the requirements:-

a) Write the algorithm using loop (as opposed to recursion)

→ Algorithm for finding factorial:-

① Start

② Declare variable n, i , factorial

③ Read number n from user

④ Initialise variable factorial = 1 and $i = 1$

⑤ Repeat until $i \leq n$

then factorial = factorial * i

$i = i + 1$

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⑥ Print factorial

⑦ Stop

Pseudo code for calculating factorial:-

read n

factorial = 1

i = 1

WHILE $i \leq n$

factorial = factorial * i

i = i + 1

END WHILE

WRITE factorial

b) Explain what "loop invariant" is and prove the algorithm is correct.

→ Loop invariant is a condition that is true each time the loop test is evaluated. It is used to check the correctness of an algorithm. A loop invariant is a good one when it satisfies three properties:-

① Initialization

② Maintenance

③ Termination.

Proving correctness using mathematical induction:-

For any natural number n,

$$n! = n * (n-1) * (n-2) * \dots * 3 * 2 * 1 = n * (n-1)!$$

Let $n = 1, 2, 3$

$$1! = 1! = 1 * (1-1) = 1$$

$$2! = 2! = 2 * (2-1) = 2 * 1 = 2$$

$$3! = 3! = 3 * (3-1) = 3 * 2 = 6$$

According to mathematical induction, from the previous steps, assume, this will be true for $n = k$.

Let $n = k$,

$$n! = k! = k * (k-1) * \dots * 2 * 1 = k * (k-1)! \longrightarrow \textcircled{1}$$

Let $n = k+1$, if $n = k$ is true then $n = k+1$ will be true,

$$\begin{aligned} n! &= (k+1)! = (k+1) * k * (k-1) * \dots * 2 * 1 \\ &= (k+1) * k! \quad \dots \dots \text{Substituted } k! = k * (k-1) * \dots * 2 * 1 \\ &\quad \text{from eq. } \textcircled{1} \end{aligned}$$

\therefore According to above proof of mathematical induction, $n! = n * (n-1)!$ is true.

c) Calculate the running time of this algorithm (you can assume any constant)

$$\begin{aligned} \rightarrow T(n) &= \theta(1) + \theta(1) + \theta(1) + \theta(n) + \theta(n) + \theta(n) \\ &= \theta(3n) = \theta(n) \end{aligned}$$

Explanation:-

read n	$\theta(1)$
factorial = 1	$\theta(1)$
$i = 1$	$\theta(1)$
WHILE $i \leq n$	$\theta(n)$
factorial = factorial * i	$\theta(n)$
$i = i + 1$	$\theta(n)$

d) Write algorithm using recursion.

```

→ read n
  if n = 1
    return 1
  return n * fact (n-1)
    
```


3. Solve the "Maximum Ascending Subarray Sum" problem (no, 1800) on leetcode.com, then follow the requirements:

- a. Make a successful submission. Attach the screenshot image of your submission

Note. your submission image should be something like this (that is, pass all test cases, instead of an accept page of a particular test case)

Success Details >

Runtime: **0 ms**, faster than **100.00%** of Java online submissions for Maximum Ascending Subarray Sum.

Memory Usage: **36.6 MB**, less than **63.64%** of Java online submissions for Maximum Ascending Subarray Sum.

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Time Submitted	Status	Runtime	Memory	Language
01/07/2022 16:23	Accepted	0 ms	36.6 MB	java

3. Solve the "Maximum Ascending Subarray Sum" problem.

b) Calculate the running time of your algorithm as a function of your input size.

→

```
int temp = nums[0];           O(1)
int result = nums[0];         O(1)
for (int i = 1; i < nums.length; i++)   O(n)
    if (nums[i] > nums[i-1])             O(n)
        temp = temp + nums[i]           O(1)
    else
        temp = nums[i]                 O(1)
result = Math.max(temp, result)         O(n)
return result
```

$\therefore T(n) = O(1) + O(1) + O(n) + O(n) + O(1) + O(1) + O(n)$
 $= O(3n) = O(n)$