

## CSE 101 – Nov 8, 2019 (Week 6)

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### Lemma 5

Let  $G$  be connected with  $n$  vertices and  $m$  edges. Assume also  $m = n - 1$ . Then  $G$  is acyclic, hence a tree

**Proof** Assume, to get a contradiction, that  $G$  is not acyclic, containing a cycle  $C$ .

Let  $e$  be any edge on  $C$  and remove it. Call the result:

$$G - e$$

Observe  $G - e$  is connected.

By lemma 3:

$$|E(G - e)| \geq |V(G - e)| - 1$$

$$\therefore m - 1 \geq n - 1$$

$$\therefore m \geq n$$

But also:  $m = n - 1$

$$\therefore n - 1 \geq n \text{ (contradiction!)}$$

This contradiction shows our assumption must be false, so no such cycle exists.

$\therefore G$  is acyclic.

### Lemma 6

Let  $G$  be acyclic with  $n$  vertices and *edges* Suppose also that  $m = n - 1$ . Then  $G$  is also connected.

**Proof** Let  $k$  be the # of connected components in  $G$ . By lemma 2, we have  $m = n - k$ . But also, we have  $m = n - 1$ , so  $n - 1 = n - k$ . Therefore,  $k = 1$ , hence  $G$  is connected.

### Lemma 7

Let  $G$  be a connected graph with  $n$  vertices and  $m$  edges. Suppose also  $m = n$ . Then  $G$  contains exactly 1 cycle. (called **unicyclic**).

## Proof Handout:

### Consider 3 Properties

1.  $G$  is connected
2.  $G$  is acyclic
3.  $m = n - 1$

**Note Lemma 1:** (i) and (ii)  $\rightarrow$  (iii)

**Lemma 5:** (i) and (iii)  $\rightarrow$  (ii)

**Lemma 6:** (ii) and (iii)  $\Rightarrow$  (i)

### Logical possibilities

i	ii	iii	Possible?
F	F	F	?
F	F	T	?
F	T	F	?
F	T	T	No, by Lemma 6
T	F	F	?
T	F	T	No, by Lemma 5
T	T	F	No, by Lemma 1
T	T	T	Yes, any tree!

**Exercise:** Show all ? are true; find smallest example

### Theorem (Treeness theorem)

The following are equivalent.

- a)  $G$  is a tree
- b)  $G$  contains a **unique**  $x - y$  path for all  $x, y \in V(G)$
- c)  $G$  is connected, but if any edge  $e$  is removed, the result  $G - e$  is disconnected
- d)  $G$  is connected and  $m = n - 1$
- e)  $G$  is acyclic and  $m = n - 1$
- f)  $G$  is acyclic, but if any edge  $e$  is added (joining any 2 non-adjacent vertices), then the result  $G + e$  contains a unique cycle

*Read proof in book or handout*

## Directed graphs

### Definition

A **directed graph (digraph)** is a pair of sets

$$G = (V, E)$$

where

$$V \neq \emptyset$$

$$E \subseteq V^2 = V * V$$

( $x$  origin,  $y$  terminus)

### Example

*Drawing*

$$V = \{1, 2, 3, 4\}$$

$$E = \{(1, 2), (1, 4), (3, 1), (4, 2), (4, 3)\}$$

**Now allowed:**  $(x, x)$ ,

$$(x, y) \neq (y, x)$$

### Representation

- Incidence Matrix
- Adjacency Matrix
- Adjacency List