

## CSE 101 – Oct 11, 2019 (Week 2)

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### PA1

- Run unit tests from Piazza
- Last extension of PA1

### Example

Show that  $\sqrt{n+10} = \Theta(\sqrt{n})$

**Proof:** **Goal:** Find positive  $c_1, c_2, n_0$  such that

$$0 \leq c_1\sqrt{n} \leq \sqrt{n+10} \leq c_2\sqrt{n}$$

for all  $n \geq n_0$ . Let

$$c_1 = 1, c_2 = \sqrt{2}, n_0 = 10$$

Then if  $n \geq n_0$ , we have

$$-10 \leq 0 \text{ and } 10 \leq n$$

$$\therefore -10 \leq (1-1)n \text{ and } 10 \leq (2-1)n$$

$$\therefore -10 \leq (1-c_1^2)n \text{ and } 10 \leq (c_2^2-1)n$$

$$\therefore c_1^2n \leq n+10 \text{ and } n+10 \leq c_2^2n$$

$$\therefore 0 \leq c_1^2n \leq n+10 \leq c_2^2n$$

$$\therefore 0 \leq c_1\sqrt{n} \leq \sqrt{n+10} \leq c_2\sqrt{n}$$

$$\therefore \sqrt{n+10} = \Theta(\sqrt{n})$$

### Exercise

Let  $a, b \in \mathbb{R}$  and  $b > 0$ . Show

$$(n+a)^b = \Theta(n^b)$$

**Theorem**

If  $h(n) = O(g(n))$  and  $f(n) \leq h(n)$  for all sufficiently large  $n$ , then

$$f(n) = O(g(n))$$

**Proof** By hypothesis, we have pos  $c_1, n_1$  such that

$$(1) \forall n \geq n_1, 0 \leq h(n) \leq c_1 g(n)$$

Also there exists pos  $n_2$  such that

$$(2) \forall n \geq n_2 : 0 \leq f(n) \leq h(n)$$

We must show there exist positive  $c, n_0$  such that

$$(3) \forall n \geq n_0 : 0 \leq f(n) \leq c g(n)$$

Let  $c = c_1$  and  $n_0 = \max(n_1, n_2)$

Then if  $n \geq n_0$ , we have both (1) and (2) true, therefore (3) is also true.

$$\therefore f(n) = O(g(n))$$

*~ Think of  $\Omega$  as the “greater than” function*

**Exercise**

If  $h(n) = \Omega(g(n))$  and  $f(n) \geq h(n)$  for sufficiently large  $n$ , then

$$f(n) = \Omega(g(n))$$

**Exercise**

If  $h_1(n) = \Omega(g(n))$  and  $h_2(n) = O(g(n))$  and  $h_1(n) \leq f(n) \leq h_2(n)$  for sufficiently large  $n$ , then

$$f(n) = \Theta(g(n))$$

**Example**

Let  $k \geq 0$ . Then

$$\sum_{i=1}^n i^k = \Theta(n^{k+1})$$

**Proof**  $\sum_{i=1}^n i^k \leq \sum_{i=1}^n n^k = n * n^k = n^{k+1} = O(n^{k+1})$

**Also**

$$\begin{aligned}
\sum_{i=1}^n i^k &\geq \sum_{i=\lceil \frac{n}{2} \rceil}^n i^k \\
&\geq \sum_{i=\lceil \frac{n}{2} \rceil}^n \lfloor \frac{n}{2} \rfloor^k \\
&= (n - \lceil \frac{n}{2} \rceil + 1) \lfloor \frac{n}{2} \rfloor^k \\
&= (\lfloor \frac{n}{2} \rfloor + 1) \lfloor \frac{n}{2} \rfloor^k \\
&\geq ((\frac{n}{2} - 1) + 1) (\frac{n}{2})^k \\
&= (\frac{n}{2})^{k+1} \\
&= (\frac{1}{2})^{k+1} * n^{k+1} \\
&= \text{const} * n^{k+1} = \Omega(n^{k+1}) \\
\therefore \sum_{i=1}^n i^k &= \Theta(n^{k+1})
\end{aligned}$$

**Definition**

$$o(g(n)) = \{f(n) | \forall c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq f(n) < cg(n)\}$$

“ $g(n)$  is a strict asymptotic upper-bound”

**Review**  $O(g(n)) = \{f(n) | \exists c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq f(n) \leq cg(n)\}$

Obviously:  $o(g(n)) \subseteq O(g(n))$

**Lemma**

$$f(n) = o(g(n)) \iff \lim_{n \rightarrow \infty} (\frac{f(n)}{g(n)}) = 0$$

$f(n) = o(g(n))$  is equivalent to

$$\forall c > 0, \exists n_0 > 0, \forall n \geq n_0 : 0 \leq \frac{f(n)}{g(n)} < c$$

This is the definition of

$$\lim_{n \rightarrow \infty} (\frac{f(n)}{g(n)}) = 0$$