CSE 101 – Oct 7, 2019 (Week 2)

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PA 1 Questions

Recall $f(n) = \Omega(g(n)) \iff \exists c, n_0 \text{ s.t. } \forall n \geq n_0$

$$0 > c * q(n) > f(n)$$

holds

Diagram

If f(n), g(n) are asymptotically positive, then we have a picture

$$0 < c \ge \frac{f(n)}{g(n)}$$

Diagram

Fact: $an^2 + bn + c = \Omega(dn + e)$

Note: Both sides must be positive

Theorem

$$f(n) = O(g(n)) \iff g(n) = \Omega(f(n))$$

Proof (\Rightarrow)

Assume f(n) = O(g(n)). Then there exist positive c_1, n_1 such that for all $n \ge n_1$:

(1) $0 \le f(n) \le c_1 g(n)$

is true. We must show that $g(n) = \Omega(f(n))$, i.e. that there exist positive c_2, n_2 such that for all $n \geq n_2$:

 $0 \le c_2 f(n) \le g(n)$

is true. Let $c_2 = \frac{1}{c_1}$ and $n_2 = n_1$. Then, if $n \ge n_2$, we have $n \ge n_1$ and therefore (1)is true. Hence

$$0 \le \frac{1}{c_1} f(n) \le g(n)$$

Since $c_2 = \frac{1}{c_1}$, (3) gives us that (2) is true.

$$\therefore f(n) = O(g(n)) \Rightarrow g(n) = \Omega(f(n))$$

Exercise: Do ←

Analogy

f(n) = O(g(n)) is analogous to $x \le y$

 $f(n) = \Omega(g(n))$ is analogous to $x \ge y$

 $f(n) = \Theta(g(n))$ is analogous to x = y

f(n) = o(g(n)) is analogous to x < y

 $f(n) = \omega(g(n))$ is analogous to x > y ### Definition

$$\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$$

Equivalently:

 $f(n) = \Theta(g(n)) \iff$ there exist positive c_1, c_2, n_0 such that for all $n \geq n_0$

$$0 \le c_1 g(n) \le f(n) \le c_2 g(n)$$

We say: g(n) is a tight asymptotic bound on f(n)

If f(n), g(n) are asymptotically positive, then

$$0 \le c_1 \le \frac{f(n)}{g(n)} \le c_2$$

Exercise: Prove $f(n) = \Theta(g(n)) \iff g(n) = \Theta(f(n))$

Exercise: Let c > 0 and g(n) be a function. Prove that

$$c * g(n) = O(g(n))$$

$$c * g(n) = \Omega(g(n))$$

$$c * g(n) = \Theta(g(n))$$

Example

$$\sqrt{n+10} = \Theta(\sqrt{n})$$

Proof

We must find positive c_1, c_2, n_0 such that for all $n \geq n_0$

$$0 \le c_1 \sqrt{n} \le \sqrt{n+10} \le c_2 \sqrt{n}$$