

# Data Privacy in the Digital World

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# What's Wrong with Our Data?

- ▶ 2 pieces of information for example: 1) there is only 1 person in a distant town A who has a certain disease; 2) Bob is from town A and checked in to the hospital
- ▶ Oftentimes, just removing sensitive information is not enough to protect the privacy of the data
- ▶ The sensitive information is identifiable when linked with another data set.
- ▶ 87% of individuals living in the US can be uniquely identified by using 3 data features: birth date, zip code, and gender

# Differential Privacy - Definition in Practice

- ▶ I'm doing a survey about mental health that requires my sensitive information
- ▶ My college will release the data set for research but remove the sensitive information
- ▶ Still not private enough. My college will modify the data set
  - ▶ Captures the characteristics of the original data set while also making my information unidentifiable
  - ▶ Utility vs. privacy trade off
- ▶ To ensure unidentifiability: let the existence of one single answer make no difference on the probability of getting the released data set

## Differential Privacy - Mathematical Definition

- ▶ Let  $I$  be the population whose data are collected
- ▶  $d_i$  be the information given by person  $i$
- ▶  $D_I = d_i | i \in I$  be the data set collected from all people in  $I$
- ▶  $Q$  be the privatized query run on a data set, and  $R = Q(D_I)$  be the resultant modified data set released to the public.
- ▶ Ideally, since whether one person is in the data set does not impact the answers or data set released, we have

$$Q(D_{I-me}) = Q(D_I)$$

This should hold whenever, meaning the probability of  $Q(D_{I-me})$  being equal to  $Q(D_I)$  should be similar. Thus  $\epsilon$ -differential privacy is defined as:

$$\frac{\text{Prob}(Q(D_I) = R)}{\text{Prob}(Q(D_{I \pm i}) = R)} \leq e^\epsilon, \text{ for small } \epsilon \geq 0$$

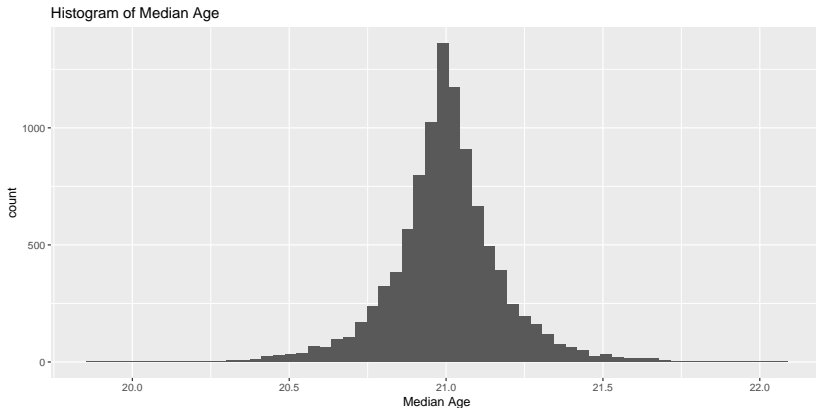
## Differential Privacy Methods

## Method 1: Laplacian Noise

- ▶ Add to the true answers noises drawn from the Laplace distribution:  $TrueValue \pm Noise$
- ▶ Parameters:  $Noise \sim Lap(\mu = 0, b = \frac{\Delta F}{\epsilon})$
- ▶ Tune the parameters to be differentially private
- ▶ Global sensitivity:  $\Delta F = \max_{(D_1, D_2)} |F(D_1) - F(D_2)|$ , which means max difference in answers that adding or removing any individual from the data set can cause
- ▶ The released answers will have a Laplace distribution  
 $Prob(R = x|D) = \frac{\epsilon}{2\Delta F} e^{-\frac{|x - F(D)|\epsilon}{\Delta F}}$  with  $\epsilon$ -differential privacy

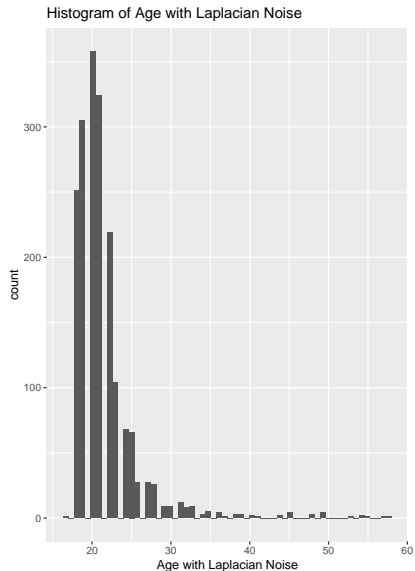
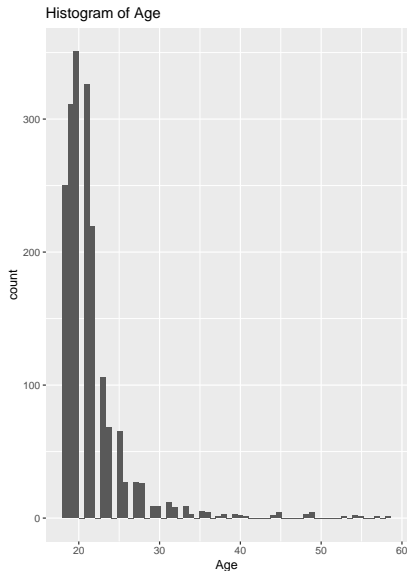
# Method 1: Laplacian Noise - Binary Data & A Statistic

- ▶ Useful when releasing the count, mean, median,...
- ▶ Example: Query = Median Age in a data set. True median = 21.



# Method 1: Laplacian Noise - Numeric Data

- ▶ add noise to each age and round to a whole number



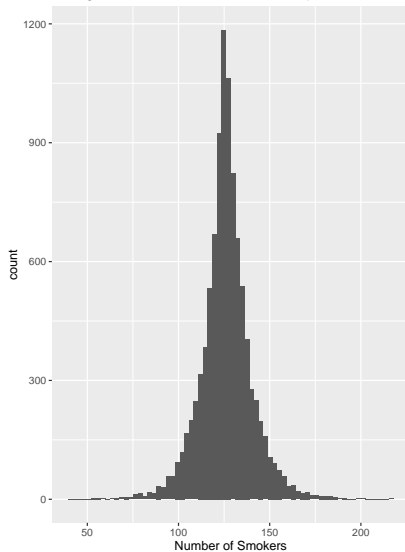


## Method 2: Randomized Response

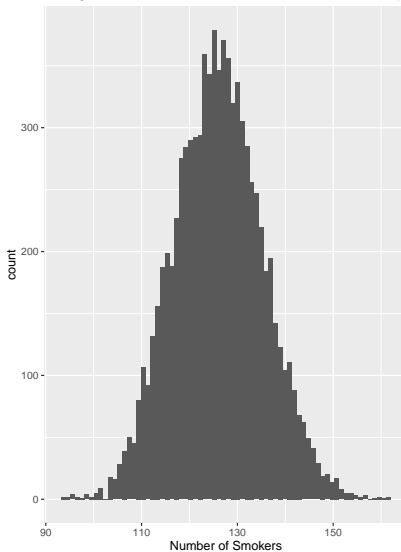
- ▶ Useful when releasing the count in a binary data set
- ▶ Flip a biased coin with probability of heads  $\alpha$ . If heads, then answer truthfully with  $d$ . If tails, flip a coin with probability of heads  $\beta$  and then answer with one response if heads and the other response if tails.
- ▶ We control the parameters  $\alpha$  and  $\beta$  that satisfy differential privacy using an extreme case of definition:  
$$\frac{P[Q(d_{yes}, \alpha, \beta) = \text{Yes}]}{P[Q(d_{no}, \alpha, \beta) = \text{Yes}]} \leq e^\epsilon$$
 from which we can get  
$$\ln\left(\frac{\alpha + (1-\alpha)\beta}{1 - (\alpha + (1-\alpha)\beta)}\right) \leq \epsilon.$$
- ▶ For numeric data, we can also flip one coin with probability  $\alpha$ , and report with a Laplacian noise if tail.

## Method 2: Randomized Response

Histogram of Number of Smokers – Laplace



Histogram of Number of Smokers – Randomized Resp



# Evaluation: Utility vs. Privacy

- ▶ Take differentially private output of the mean for example
- ▶ Utility:  $\frac{|output - real|}{\epsilon\sqrt{n}}$ , the percentage of the outputs that are useful
- ▶ Exponential Mechanism:  $Pr[M_q^\epsilon(D) = o] = \frac{\exp(\frac{\epsilon q(D,o)}{2\Delta q})}{\sum_{o' \in O} \exp(\frac{\epsilon q(D,o')}{2\Delta q})}$ ,  
the probability of an output

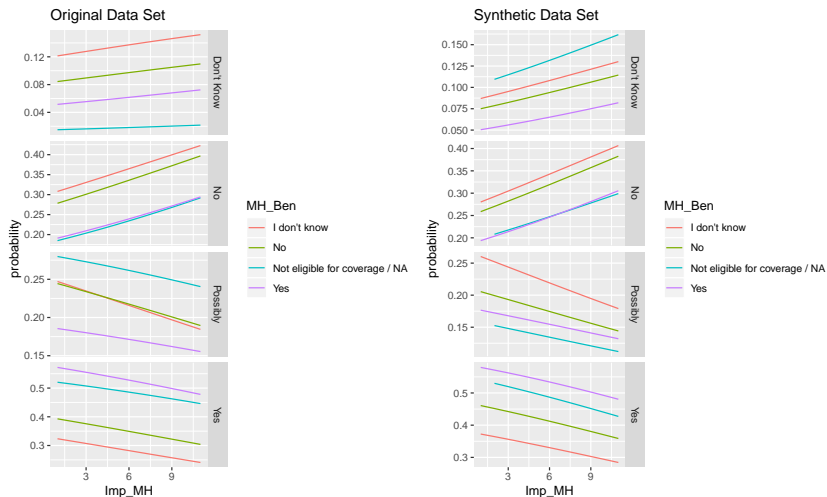
## Synthetic Data

# Synthesizing Open Sourcing Mental Illness Dataset

- ▶ Using `synthpop`, we can create a synthetic data set
- ▶ Control: order, method, restrictions. . .
- ▶ does not ensure  $\epsilon$ -differential privacy

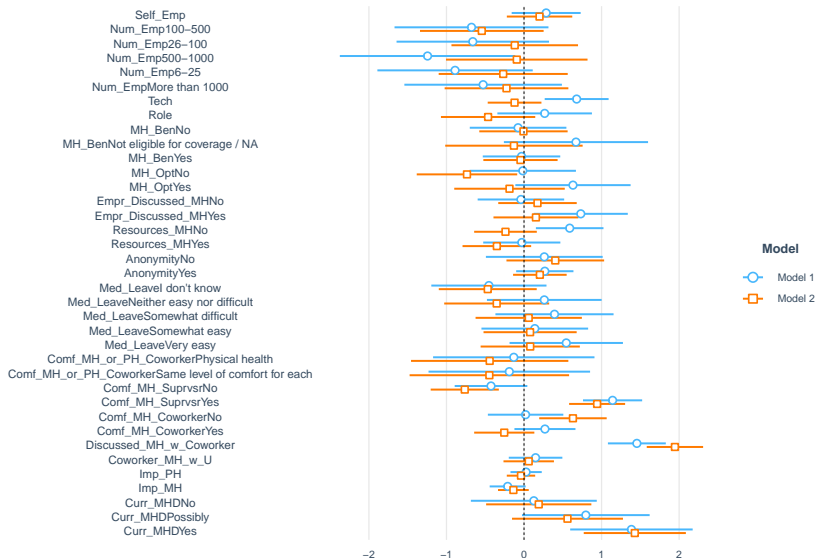
# Synthesizing Open Sourcing Mental Illness Dataset

- ▶ The probability one currently has a mental health disorder vs. how much the company values mental health



# Synthesizing Open Sourcing Mental Illness Dataset

- Confidence intervals of parameters of Logistic Regressions using original (Model 1) and synthetic (Model 2) data set



# Wake Forest University Healthy Minds Dataset

- ▶ Mostly categorical variables that are not necessarily correlated, thus `synthpop` does not work well
- ▶ Synthesize a subset of variables and move on to the whole data set.



## Reference

- ▶ [1] Open Sourcing Mental Illness Ltd. OSMI Mental Health in Tech Survey. 2017, 2018. url: <https://osmihelp.org/research>.
- ▶ [2] David McClure and Jerome P. Reiter. Differential Privacy and Statistical Disclosure Risk Measures: An Investigation with Binary Synthetic Data“. In: Trans. Data Privacy 5.3 (Dec. 2012), pp. 535-552. issn: 1888-5063. url: <http://dl.acm.org/citation.cfm?id=2423656.2423658>.
- ▶ [3] Christine Task. Privacy-preserving Datamining: Differential Privacy And Applications. June 2014.
- ▶ [4] Roberto Agostino Vitillo. Differential Privacy for Dummies. July 2016. url: <https://robertovitillo.com/2016/07/29/differential-privacy-for-dummies/>.