

sit with your clan
if you can

CMPS 12B/M

Introduction to Data Structures

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How fast is HighArray?

- ▼ How many steps does it take on average to do these operations for HighArray?
 - ▼ (Assume no duplicates allowed)
 - ▼ Let n be number of elements in array.
 - ▼ Insert
 - ▼ Delete
 - ▼ Find

What are we discovering?

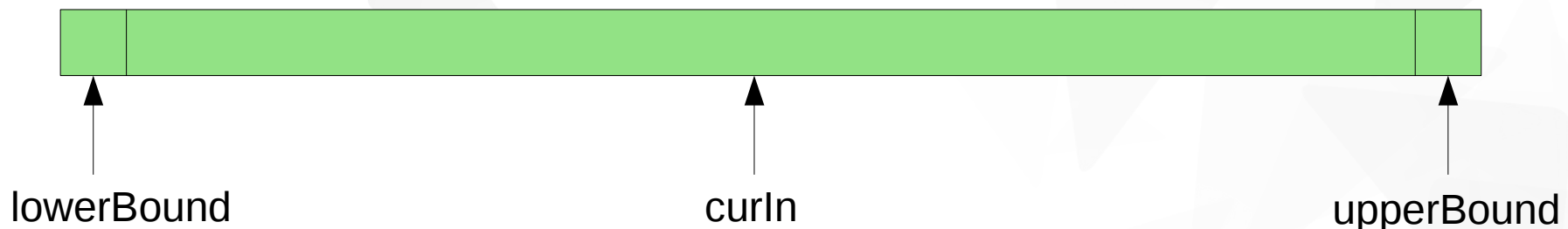
- ▼ Need to carefully define the thing we're counting
 - ▼ "step" is ambiguous
 - ▼ "comparison" or "array access" are better
- ▼ Insert is faster than delete
- ▼ Find is linear in size of array

How fast is OrderedArray?

- ▼ How many steps does it take on average to do these operations for OrderedArray?
 - ▼ (Assume no duplicates allowed)
 - ▼ Let n be number of elements in array.
 - ▼ Insert
 - ▼ Delete
 - ▼ Find

Binary Search

- ▼ You're looking for a needle in a haystack
- ▼ Compare needle to middle hay
 - ▼ If it matches, you're done
 - ▼ If it is smaller, then find the needle in the left haystack
 - ▼ If it is larger, then find the needle in the right haystack
- ▼ To keep track of where we're looking:
 - ▼ lowerBound, curIn, upperBound
 - ▼ Keep going until $\text{lowerBound} > \text{upperBound}$



Questions

- ▼ How many comparisons will it take?
- ▼ Does it matter if we pick the "upper" middle or the "lower" middle?
- ▼ Experiment time
 - ▼ Run the algorithm for $n=10$
 - ▼ Each group pick randomly: upper/lower middle
 - ▼ Use seat row last digit for which number to search for, in:
 - ▼ 0 1 2 3 4 5 6 7 8 9

Formula for Comparisons

▼ What function is the number of comparisons as n gets bigger?

▼ Worse case:

▼ $n=10$	4
▼ $n=100$	7
▼ $n=1000$	10
▼ $n=10000$	14
▼ $n=100000$	17

What function looks like this?

$$\log_2(n)$$

Recurrence Relation

- ▼ Another way of thinking about the question...
 - ▼ Let $T(n)$ be the function we're trying to find
- ▼ What happens when you double the number of businesses?
 - ▼ For OrderedArray:
 - ▼ $T(2n) = \text{___} T(n) \text{___}$

You Can Solve Recurrences

- ▼ Techniques
 - ▼ Guess and check (the most powerful technique)
 - ▼ Look for patterns (linear, quadratic, logarithmic, exponential)
 - ▼ Ask Wolfram Alpha



Another Recurrence

- ▼ What about HighArray?
- ▼ How many comparisons does HighArray take?
 - ▼ $T(2n) = \underline{\hspace{1cm}} T(n) \underline{\hspace{1cm}}$
 - ▼ Solve it using guess & check (since you already know the answer)

Comparing Algorithms

- ▼ If we know the exact problem size n
 - ▼ Can directly compare $T_1(n)$ and $T_2(n)$ for different algorithms
- ▼ How do we compare algorithms more abstractly?
 - ▼ Different "classes" of algorithms
 - ▼ Grow bigger in n in different ways
 - ▼ Don't care so much about specific ratio with fixed problem size

Some Classes of Algorithms

- ▼ Examples:
 - ▼ *Linear* – doubling the n makes it twice as big
 - ▼ Searching unordered array (average or worst case)
 - ▼ *Constant* – changing n doesn't make it change
 - ▼ Inserting into an array without holes
- ▼ Our latest class:
 - ▼ *Logarithmic* – doubling n makes it one bigger

Ignoring the Constant

- ▼ All linear algorithms have a formula like:
 - ▼ $T(n) = K n$
 - ▼ Average number of comparisons to find an element in an unordered array: $\frac{1}{2} n$
 - ▼ Worst case number of comparisons to find an element in an unordered array: n
- ▼ Key idea of Big O notation
 - ▼ ***Ignore the constant***
 - ▼ All linear algorithms are in $O(n)$
 - ▼ $O(\text{something})$ is a set of functions that grow as n gets bigger like *something*

Big O

- ▼ Some algorithms and running times
 - ▼ Linear search $O(n)$
 - ▼ Binary search $O(\log n)$
 - ▼ Insert into unordered array $O(1)$
 - ▼ Insert into ordered array $O(N)$
 - ▼ Delete in unordered array $O(N)$
 - ▼ Delete in ordered array $O(N)$

Simplifying into Big O

- ▼ Suppose an algorithm always takes 3 comparisons no matter the size of the input.

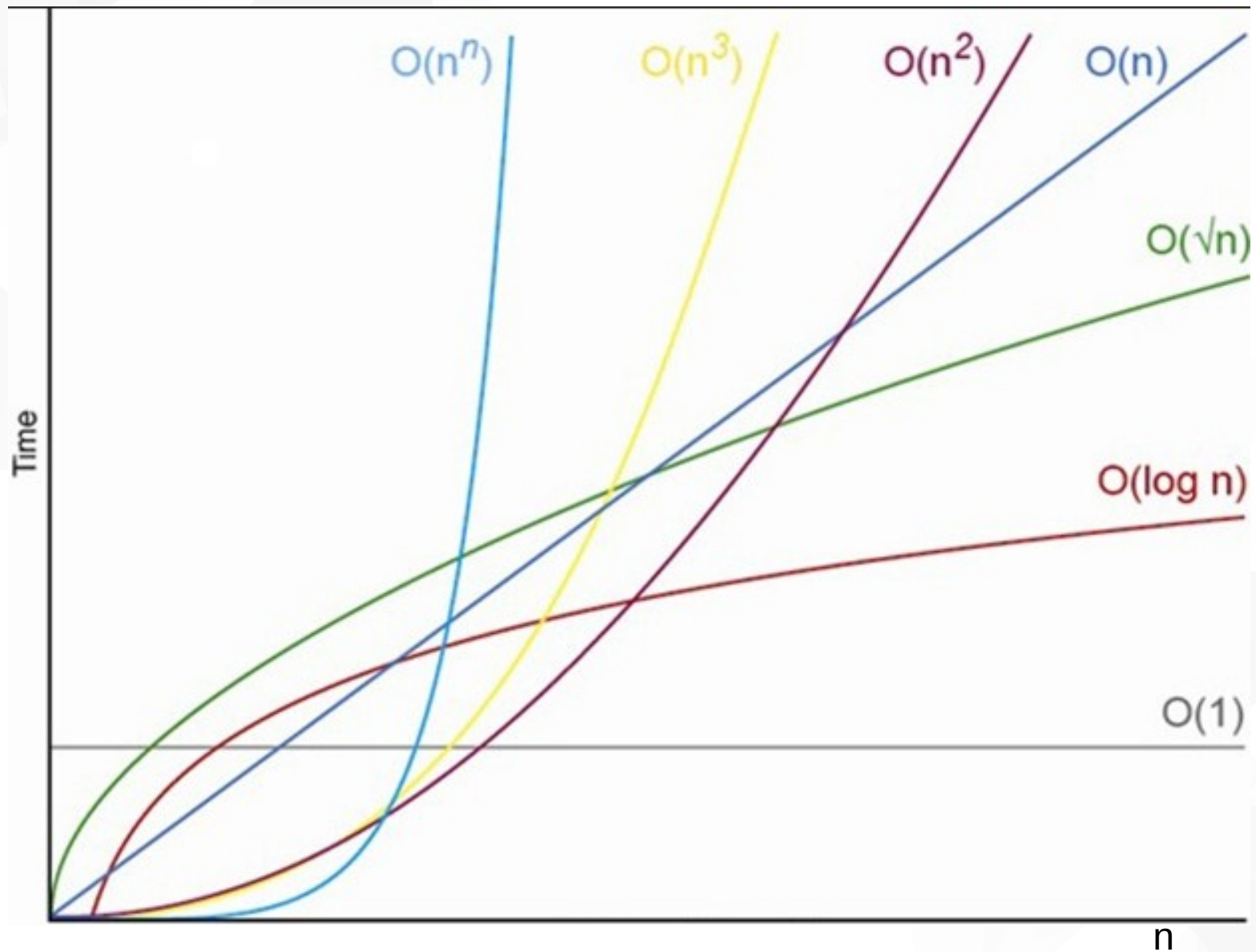
▼ $T(n) = 3$ $O(3)$ $O(1)$

- ▼ Suppose an algorithm takes n comparisons to scan through an entire array, then does another n comparisons to scan it *again*

▼ $T(n) = 2n$ $O(2n)$ $O(n)$

ignore any multiplicative constant

Complexities Graphed



Cheatsheet

Data Structure	Time Complexity					
	Average				Worst	
	Indexing	Search	Insertion	Deletion	Indexing	Search
Basic Array	$O(1)$	$O(n)$	-	-	$O(1)$	$O(n)$
Dynamic Array	$O(1)$	$O(n)$	$O(n)$	$O(n)$	$O(1)$	$O(n)$
Singly-Linked List	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$	$O(n)$
Doubly-Linked List	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$	$O(n)$
Skip List	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(n)$	$O(n)$

<http://bigocheatsheet.com/>

A Business Problem

- ▼ Suppose we have a text file
 - ▼ One line per business
 - ▼ Business name, colon, phone number
- ▼ Need a program that:
 - ▼ Starts up, reads all businesses
 - ▼ Waits for queries that are business names
 - ▼ Prints out phone number for matching business
 - ▼ OR shows "NOT FOUND" if no business found
- ▼ *Which data structure should we use?*
 - ▼ How do we decide?
 - ▼ Does it even matter?

Simplest Thing

- ▼ Start with simplest thing
 - ▼ Put everything in an array

```
class Record {
    public String name;
    public String number;
}

class BusinessDatabase {

    // Database of all businesses
    private Record[] businesses;

    public Record parse(String txt) {
        // Parsing code here
        ...
    }

    // Parse the input lines
    // Assuming one business per line
    public void parseData(String[] lines) {
        int numBusinesses = lines.length;
        businesses = new Record[numBusinesses];

        for(int i = 0; i < lines.length; i++) {
            Record rec = parse(lines[i]);
            businesses[i] = rec;
        }
    }
}
```

What about search?

▼ To find a business:

```
// Search for a business
public String getPhoneNumber(String name) {
    for (int i = 0; i < businesses.length; i++) {
        if (name.equals(businesses[i].name)) {
            return businesses[i].number;
        }
    }
    return "NOT FOUND";
}
```

How many comparisons does it take
to find a business if there are 1,000,000?

Which is better for problem?

- ▼ Decide between unordered and ordered array
 - ▼ Once data is all loaded and processed, which is better for searching?
 - ▼ For small numbers of businesses, probably doesn't matter
 - ▼ Logarithmic growth is slower as n increases
 - ▼ For 1,000,000 businesses, go with ordered array
- ▼ But how do we get the data into order?
 - ▼ Inserting each business would take $O(n)$, for each business...

how do we order the data?

Simple Sorting

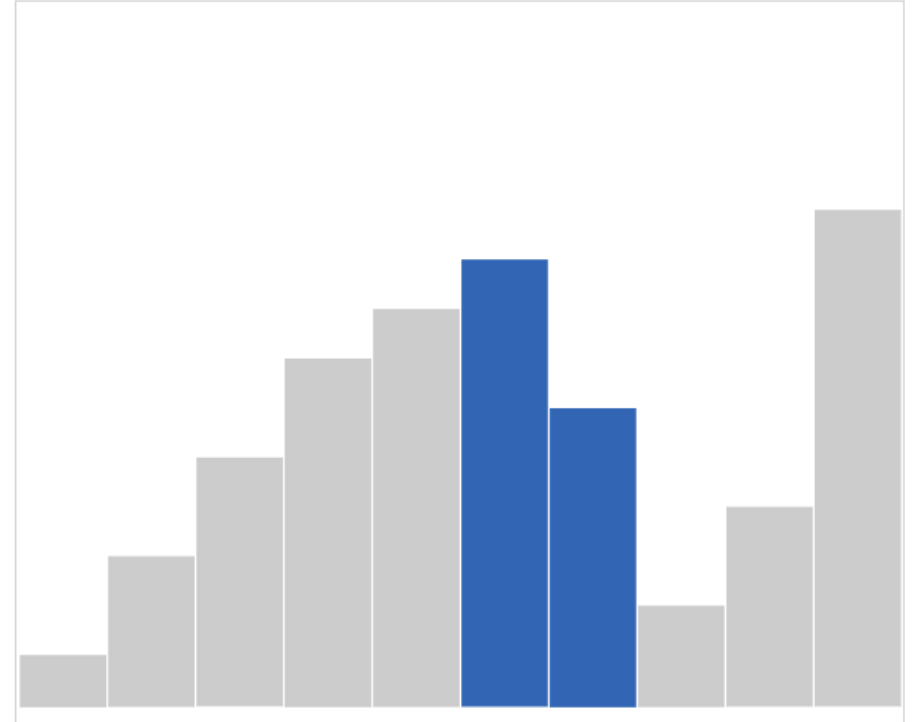
- ▼ Well if you're sure, better be... GRYFFINDOR!

Sort Problem

- ▼ Suppose you had an unordered array, and wanted an ordered array
- ▼ Operations allowed:
 - ▼ Compare two items
 - ▼ Swap two items, or copy one item to a temp location
- ▼ All sorting algorithms do the same basic things
 - ▼ Produce ordered output using compares/swaps
 - ▼ Difference in the details

Bubble Sort

- ▼ Scan through left to right
 - ▼ Compare adjacent entries
 - ▼ If they are wrong way around, swap them
- ▼ If any swaps were made, loop back again
- ▼ No swaps made means you're done



<http://www.algomatic.com/player?algorithm=54162b10f3166302000a91f2>

Bubble Sort in Java

- ▼ `Examples/Chap03/BubbleSort/bubbleSort.java`
- ▼ Observation:
 - ▼ Each run through, we know we have the biggest element to the right
 - ▼ So stop scanning once we hit the finished section on the right

Running Time of BubbleSort

- ▼ How many comparisons does optimized bubble sort make in the worst case?
 - ▼ For $n=10$:
 - ▼ Our variation: $9+9+9+9+9+9+9+9+9$
 - ▼ With observation: $9+8+7+6+5+4+3+2+1$



The End