



sit with anyone that
looks friendly

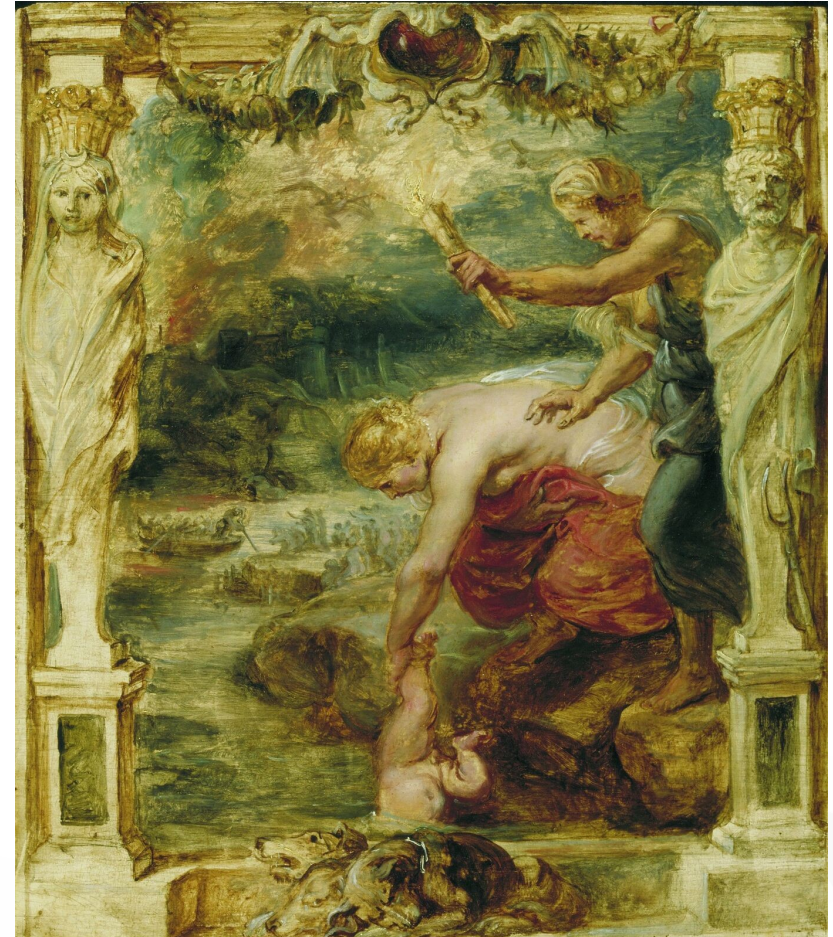
CMPS 12B/M

Introduction to Data Structures

▼ Instructor: Nathan Whitehead

Achille's Heel

- ▼ Goddess Thetis dips her baby in river Styx in Hades
 - ▼ Achilles becomes invulnerable except on his heel where held
- ▼ What is merge sort's heel?
- ▼ Number of comparisons
 - ▼ $O(n \log(n))$ *good*
- ▼ Space required
 - ▼ Needs extra n workspace *bad*



More Visualizations

 Insertion	 Selection	 Bubble	 Shell	 Merge	 Heap	 Quick	 Quick3
 Random							
 Nearly Sorted							
 Reversed							
 Few Unique							

<http://www.sorting-algorithms.com/>

Moral of Story

tradeoffs

- ▼ There is not one perfect sort that is optimal for all problems
- ▼ Best sort depends on
 - ▼ which operations are expensive
 - ▼ how the data is ordered coming in
 - ▼ usually not entirely random (random is rare)
 - ▼ how easy it is to get working
 - ▼ buggy sort is no use to anyone

Why not just use the library sort built-in?

- ▼ **YES, YOU SHOULD!**
- ▼ But not always possible
 - ▼ API might not match way you want to sort
 - ▼ Structure might not all fit into memory
 - ▼ Might be badly optimized for your actual data ordering
 - ▼ There might be *many* sorts in libraries, how to choose?

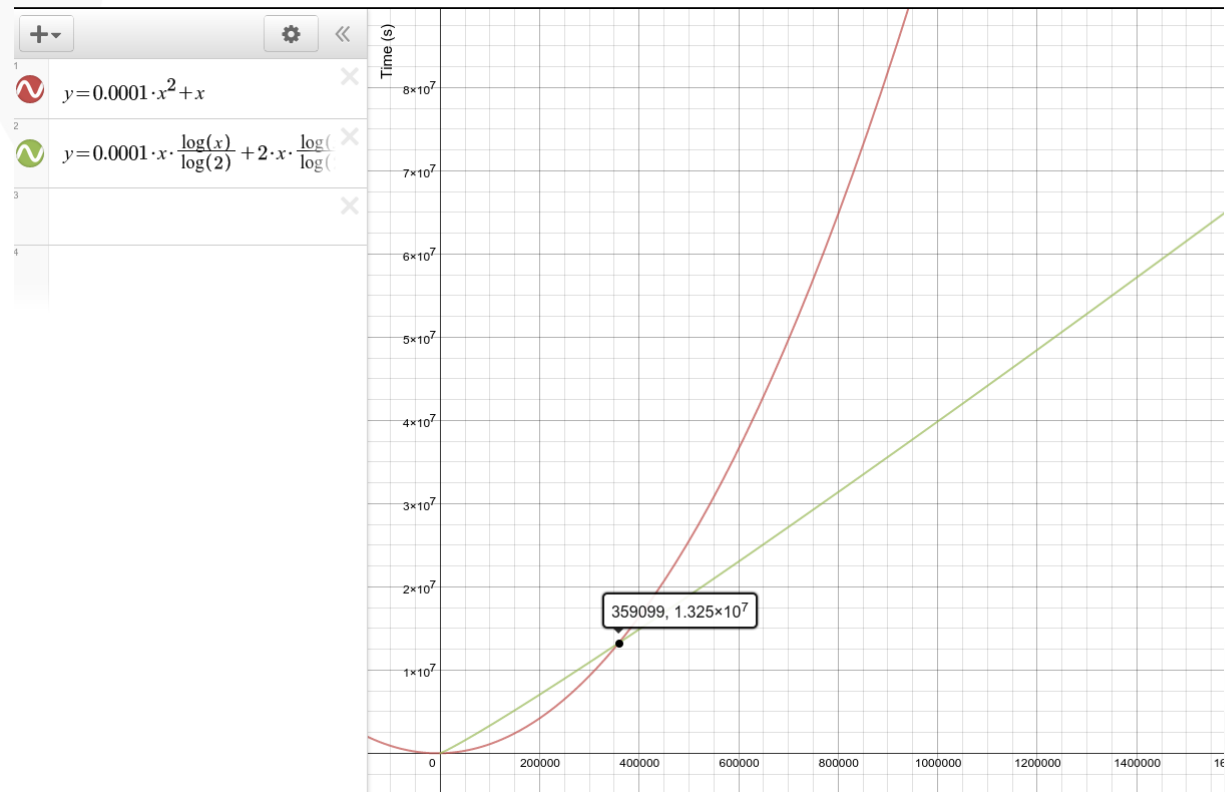
Sorting Pictures in Flash Storage

- ▼ **Quiz problem 3** – more details
- ▼ Suppose
 - ▼ Comparisons take 0.1 ms (read from cache)
 - ▼ Writing one picture to flash takes 1 s (slow!)
- ▼ *Merge sort* versus *Selection sort*
 - ▼ Merge sort has better big O...
 - ▼ Selection sort is linear in swaps...
- ▼ About how big does n have to be before merge sort wins?

Times

- ▼ Selection sort
 - ▼ Total comparisons
 - ▼ About n^2
 - ▼ Total moves/swaps
 - ▼ Exactly n
- ▼ Merge sort
 - ▼ Total comparisons
 - ▼ About $n \log_2(n)$
 - ▼ Total moves/swaps
 - ▼ About $2 n \log_2(n)$
- ▼ Crossover
 - ▼ $n=360000$, time= $1.3e7$ seconds = 150 days

<https://www.desmos.com/calculator/kyks5ivrlj>



Adding up Numbers

- ▼ Let's add up numbers using code
- ▼ Given n , find sum of numbers 1 through n
- ▼ OK, we can do this...

Imperative (i.e. normal)

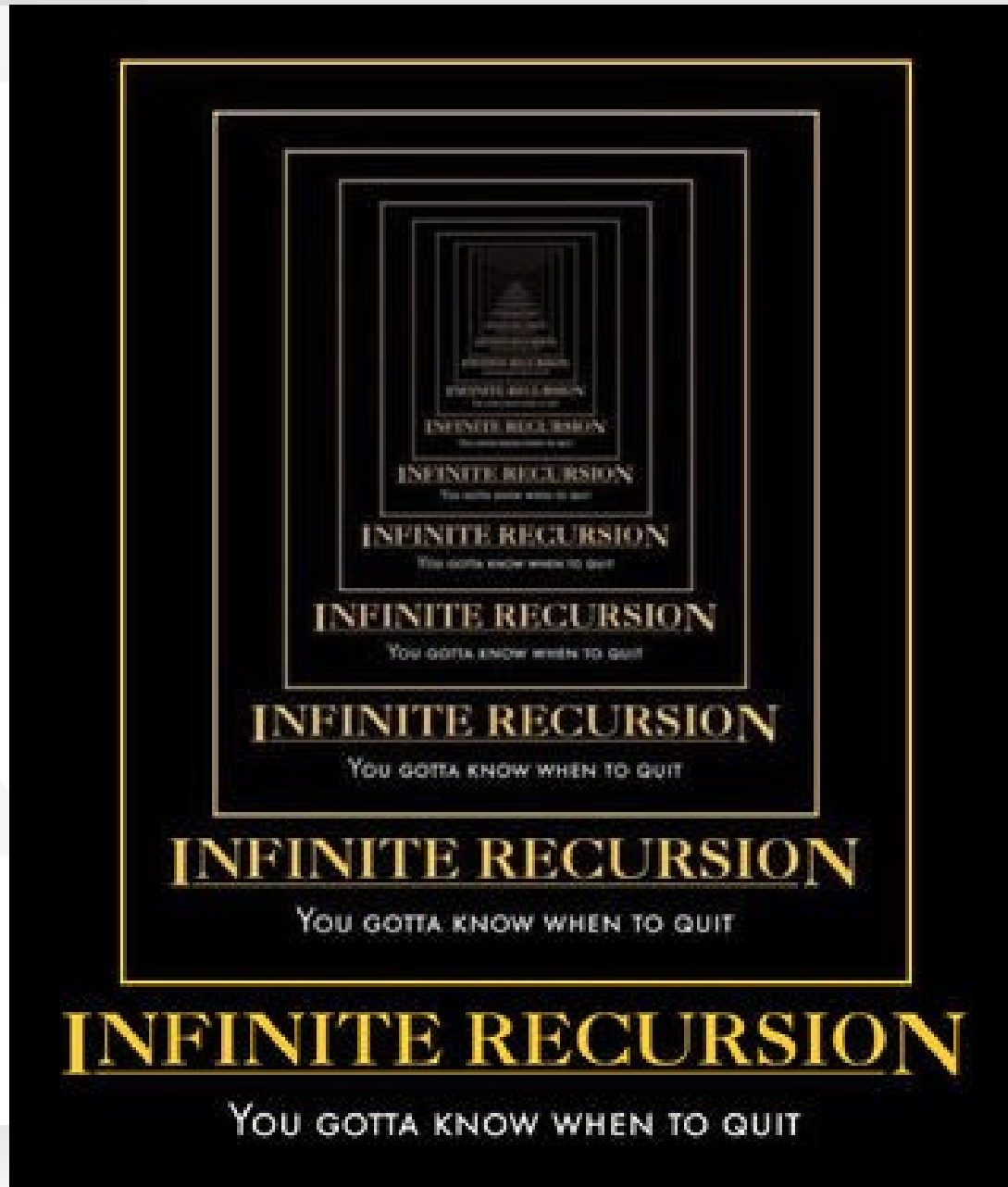
```
public static int sum(int n) {  
    int total = 0;  
    for(int i = 1; i <= n; i++) {  
        total += i;  
    }  
    return total;  
}
```

Recursive Version

- ▼ Want to write function sum
- ▼ What if we could ask the genie for solutions to any smaller problem?
 - ▼ How would that help?
 - ▼ How could we use smaller solutions to build bigger solutions?
- ▼ Suppose we know $\text{sum}(n-1)$
 - ▼ How does that help us get $\text{sum}(n)$?

$$\text{sum}(n) = \text{sum}(n-1) + n$$

Knowing When to Quit



Recursive Sum Facts

$$\text{sum}(n) = \text{sum}(n-1) + n$$

▼ When to quit?

▼ How about $\text{sum}(1) = 1$

If $n > 1$ then

$$\text{sum}(n) = \text{sum}(n-1) + n$$

else

$$\text{sum}(1) = 1$$

Recursive Sum Code

```
public static int sum(int n) {  
    if (n > 1) {  
        return sum(n - 1) + n;  
    } else {  
        return 1;  
    }  
}
```

sum(4)
=(sum(3) + 4)
=((sum(2) + 3) + 4)
=(((sum(1) + 2) + 3) + 4)
=(((1+2)+3)+4)
=((3+3)+4)
=(6+4)
=10

The Smart Way...

```
public static int sum(int n) {  
    return n * (n + 1) / 2;  
}
```

math for the win!

When to use recursion?

- ▼ It takes memory to keep track of where you are in recursion
 - ▼ Cheap but not free
- ▼ It takes time to do function calls
 - ▼ Cheap but not free
- ▼ Makes many problems *conceptually simpler*
 - ▼ Small overhead is worth simpler code!

Stacks and Queues

- ▼ An Englishman, even if he is alone, forms an orderly queue of one.
-*George Mikes*

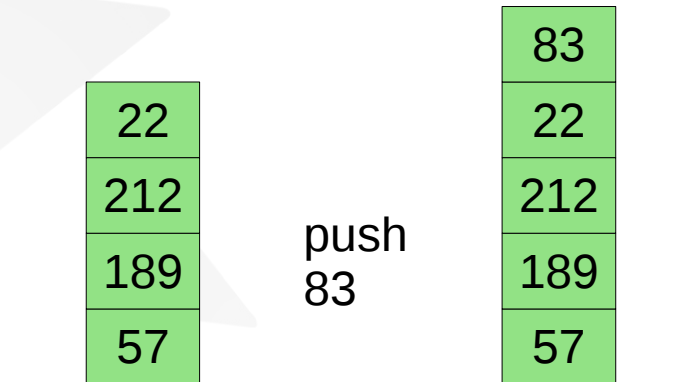
Different Kind of Structure

- ▼ Stacks and queues restrict access
 - ▼ Less capabilities compared to arrays
 - ▼ More suited to programming tasks than data tasks

interface matters

Stacks

- ▼ Only have access to top of stack
 - ▼ *Push* – add an element to top
 - ▼ *Pop* – remove top element
 - ▼ *Peek* – look at top without removing



Stack Code

▼ `Examples/Chap04/Stack/stack.java`

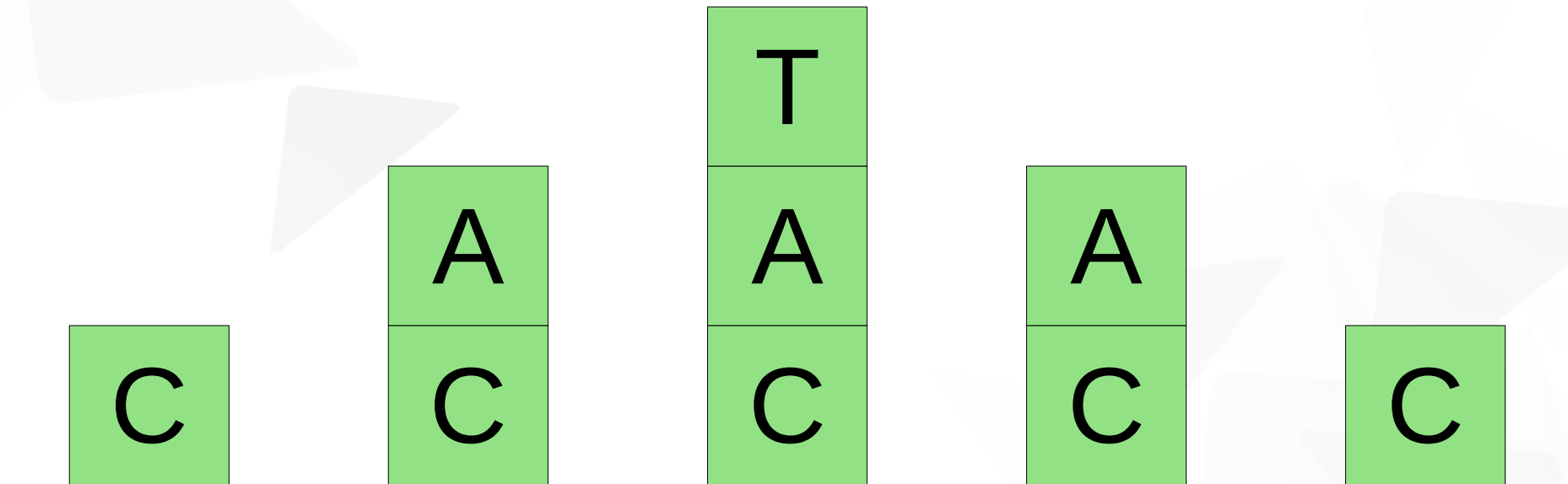
Reversing Items

- ▼ Suppose you want to process items in reverse order
- ▼ We've solved this
 - ▼ Put all the elements in an array
 - ▼ Do a loop, swapping elements to reverse the array
 - ▼ Process elements in the array
- ▼ Easier
 - ▼ Put all elements in an array
 - ▼ Loop through array in reverse order

Reversing with a Stack

C A T

T A C



Matching Parentheses and Brackets

a [b (a)]

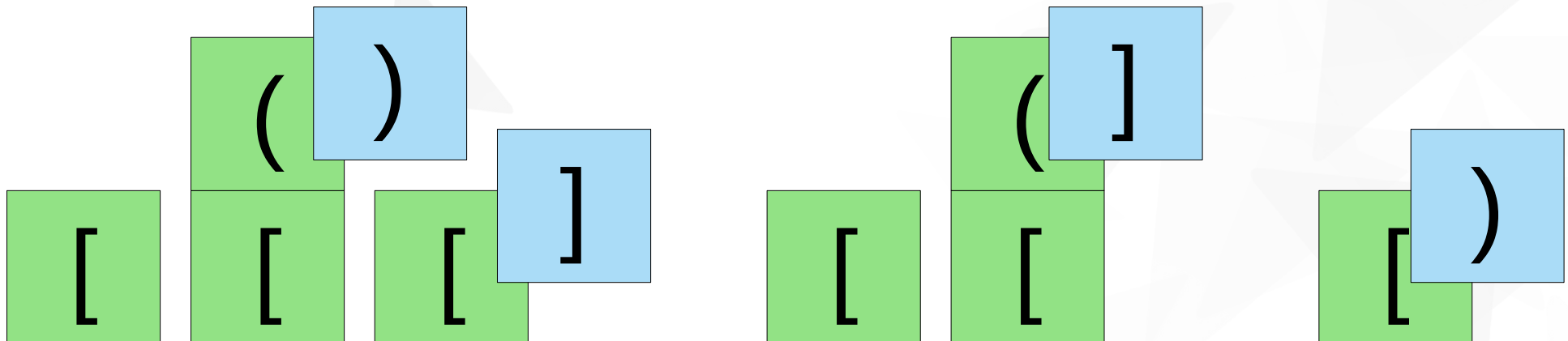
✓

a [b (a]]

✗

a [b) a (]

✗



Bracket Checker Code

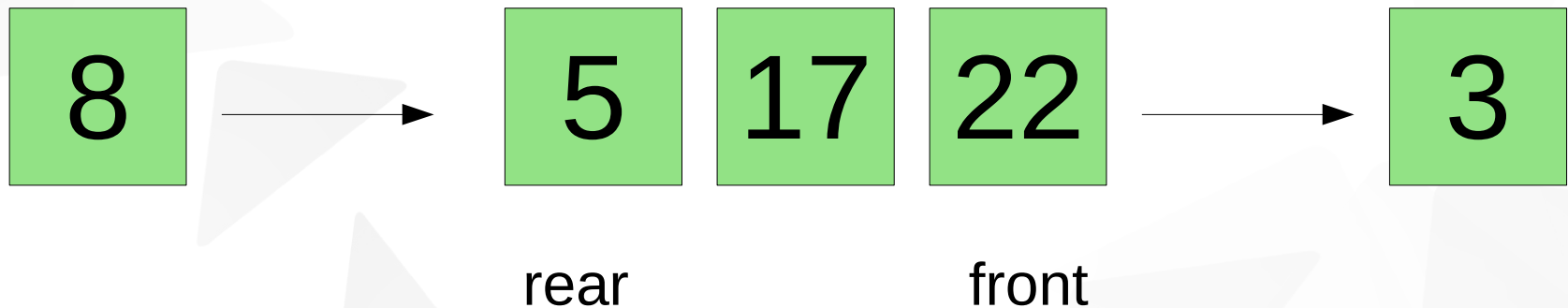
- ▼ `Examples/Chap04/Brackets/brackets.java`

Stack Efficiency

- ▼ How efficient are stacks?
 - ▼ Push, pop, peek
 - ▼ No comparisons needed
 - ▼ One array access
 - ▼ Doesn't depend on how many elements in stack
 - ▼ $O(1)$ time for all operations

Queue

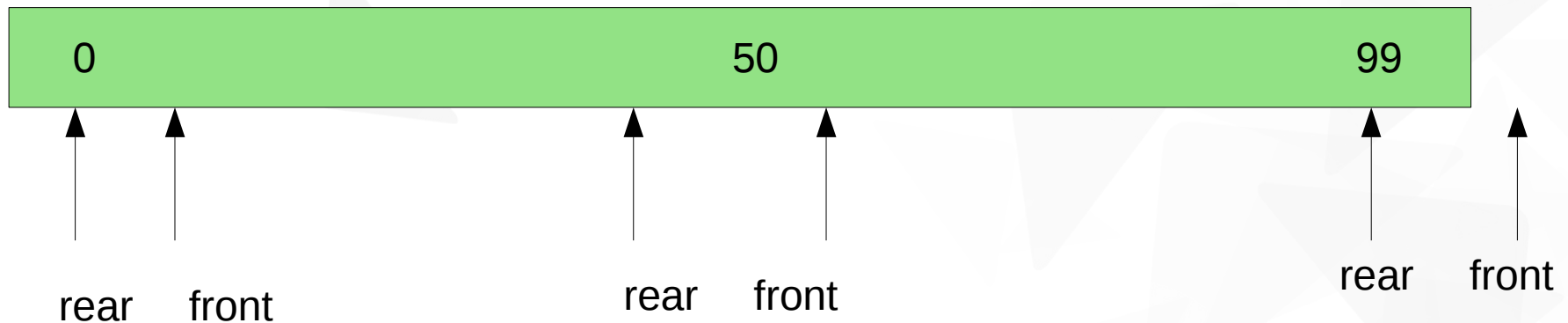
- ▼ Waiting in line
 - ▼ Arrive at the rear
 - ▼ Leave at the front
 - ▼ FIFO (first-in-first-out)



incredibly useful for keeping
track of events, work to do, etc.

Queue

- ▼ Methods
 - ▼ *Insert* – add to rear of queue
 - ▼ *Remove* – remove from front of queue
 - ▼ *Peek* – look at front of queue
- ▼ Implementation
 - ▼ Maybe an array?
 - ▼ Front and rear keep increasing over time



Circular Buffer

7	
6	
5	D
4	C
3	B
2	A
1	
0	

insert E
insert F

7	F
6	E
5	D
4	C
3	B
2	A
1	
0	

insert G

7	F
6	E
5	D
4	C
3	B
2	A
1	
0	G

remove
remove
insert H

7	F
6	E
5	D
4	C
3	
2	
1	H
0	G

wraparound

Circular Queue Code

▼ `Examples/Chap04/Queue/queue.java`

Queue Efficiency

▼ How efficient are circular queues?

▼ Insert

- ▼ No comparisons
- ▼ One array access
- ▼ $O(1)$

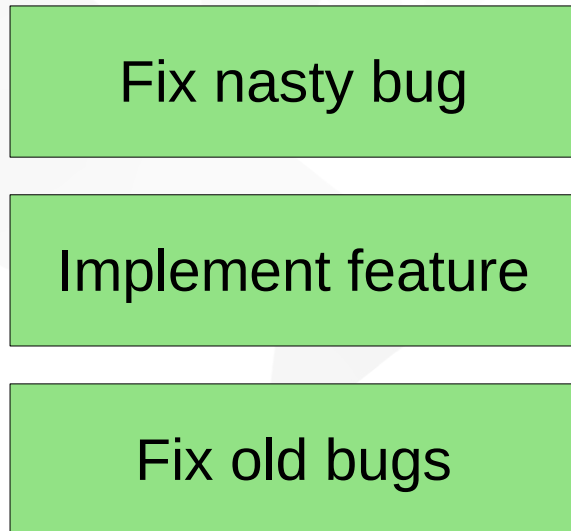
▼ Remove

- ▼ No comparisons
- ▼ One array access
- ▼ $O(1)$

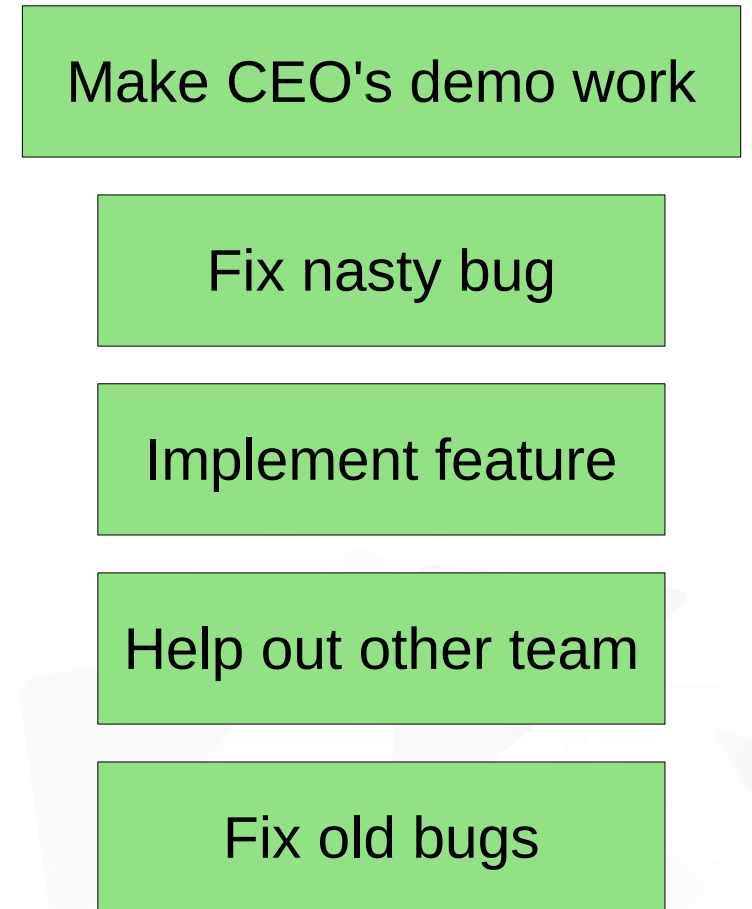
Deque

- ▼ With circular techniques we can insert and remove from front or rear of buffer
 - ▼ Left and right indexes
 - ▼ insertLeft, removeLeft, insertRight, removeRight
 - ▼ insertRight and removeRight == stack
 - ▼ insertLeft and removeRight == queue
- ▼ Deques are generalization of stacks and queues
 - ▼ Commonly provided by libraries
 - ▼ Usually a stack or queue is what problem needs

Priority Queue



new work
→



Priority Queue

- ▼ Items go into queue
- ▼ Items come out of queue lowest first
 - ▼ Not FIFO
- ▼ One implementation
 - ▼ Keep sorted array
 - ▼ *Remove* – take out element in front (it is smallest)
 - ▼ *Insert* – find proper place to put new item, make room for it

Priority Queue

7	
6	
5	
4	
3	22
2	57
1	189
0	212

Remove

7	
6	
5	
4	
3	
2	57
1	189
0	212

Insert
81

7	
6	
5	
4	
3	57
2	81
1	189
0	212

Insert
302

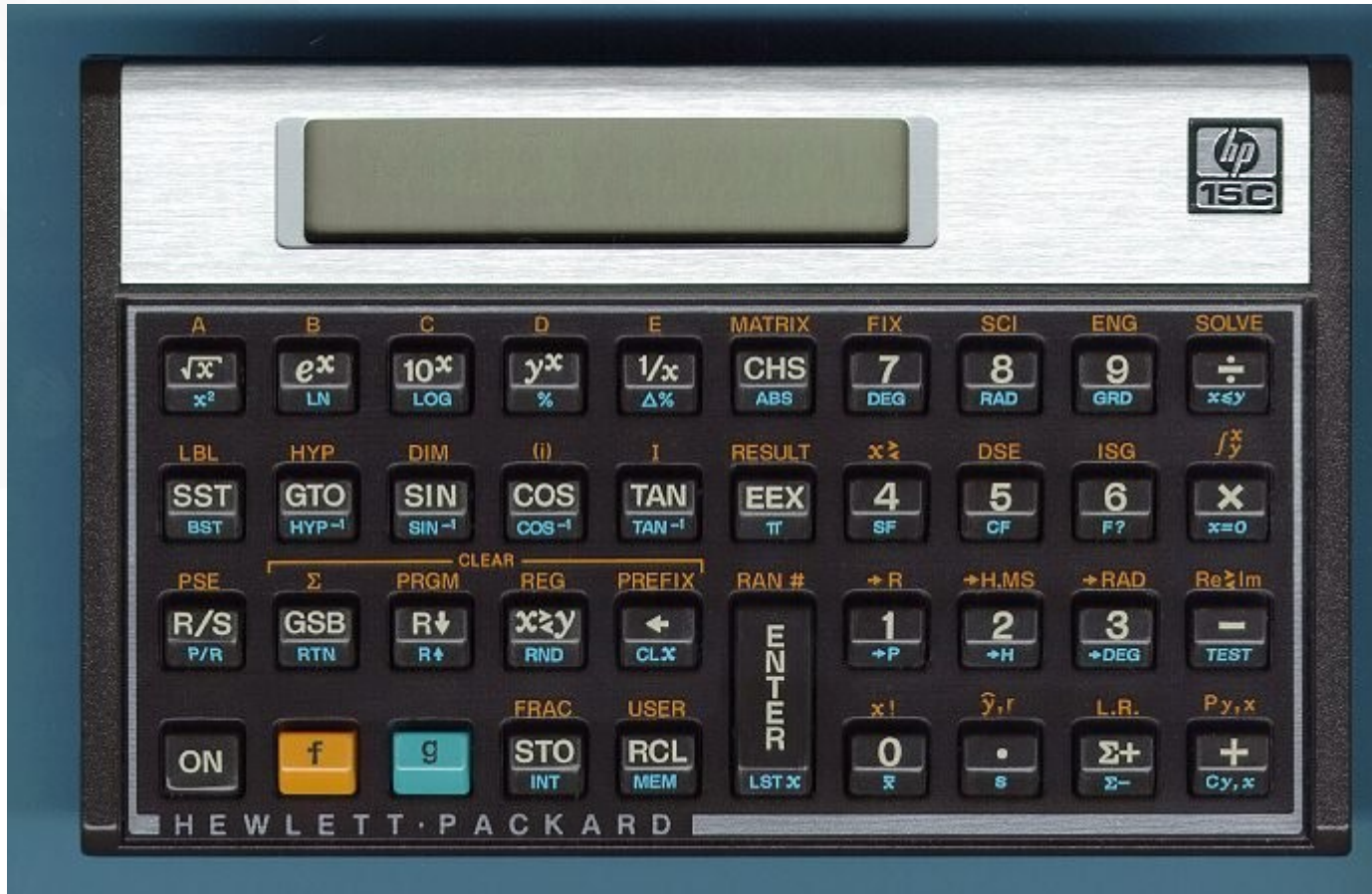
7	
6	
5	
4	57
3	81
2	189
1	212
0	302

Efficiency of Priority Queue

- ▼ Our first implementation
 - ▼ Remove
 - ▼ No comparisons
 - ▼ No copying
 - ▼ One array access
 - ▼ $O(1)$
 - ▼ Insert
 - ▼ Find where to put it
 - ▼ Shift elements out of the way
 - ▼ Worst case requires shifting n elements
 - ▼ $O(n)$ comparisons and copies

there is another way
we'll see later

Parsing Arithmetic Expressions



<http://hp15c.com/>



The End