



Integrating Intent Understanding and Optimal Behavior Planning for Behavior Tree Generation from Human Instructions

Xinglin Chen*, Yishuai Cai *, Yunxin Mao, Minglong Li †, Wenjing Yang, Weixia Xu, Ji Wang,

College of Computer Science and Technology, National University of Defense Technology {chenxinglin, caiyishuai, maoyunxin, liminglong10, wenjing.yang, xuweixia, wj}@nudt.edu.cn

汇报人姓名:王雅斌

汇报日期: 2024.10.15

步骤



目录

- 背景
- 动机
- 核心难点
- 技术方法
- 效果

研究背景



背景

行为树

自然语言

成功

机器人在家庭或工业环境中执行任务时,基本上需要具备适应性和可靠性。行为树(BT)因其模块化和反应性而成为这些场景的合适控制架构。然而,现有的BT生成方法要么不涉及自然语言的解释,要么无法理论上保证BT的成功。





解决问题





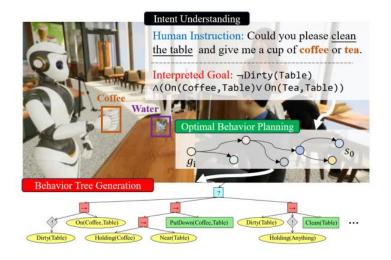
结合了LLM解决意图理解,展示OBTEA在各种指标上 优于基线BT扩展算法的优势(具有更低的成本和更少的条 件节点计时),并最终确认了框架的实际可部署性。





技术难点

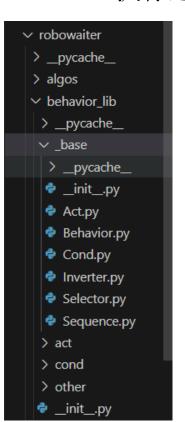
- 1. 由于自然语言的开放性,机器人必须理解用户的意图,明确任务的目标,并获得足够的高级技能来完成这些任务。
- 2. 机器人必须自动生成可靠的计划,并安全、稳健地执行这些计划。



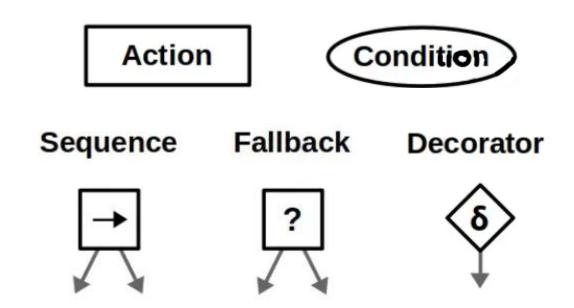


前置知识: 行为树

行为树(BT)是一种有向根树结构,其中叶节点控制机器人的感知和动作,而内部节点则处理这些叶节点的逻辑结构。行为树从其根节点开始执行,根据环境状态的不同,执行过程在树中创建控制流,最终决定机器人将执行的动作。



在本文中,我们关注五种典型的行为树节点:





前置知识:环境描述

谓词+对象

三元组 (O, Pc, Pa)

每个谓词的对象是O的子集,环境描述可以通过 pc(o1,...,ok)以 (合取)、 (析取)和 - (否定)的 形式组合表示。



前置知识: 行为规划

```
@classmethod
def get_info(cls,arg):
    info = {}
    info["pre"]= {f'Holding(Nothing)',f'Not IsClean({arg})'}
    info["add"] = {f'IsClean({arg})'}
    info["del_set"] = {f'Not IsClean({arg})'}
    return info
```

定义为三元组 P(a) =< pre(a), add(a), del(a) > 状态转移 s'=s U add(a) \ del(a)

$$D(\mathcal{T}) = \sum_{i=1}^{n} D(a_i).$$

描述

转移

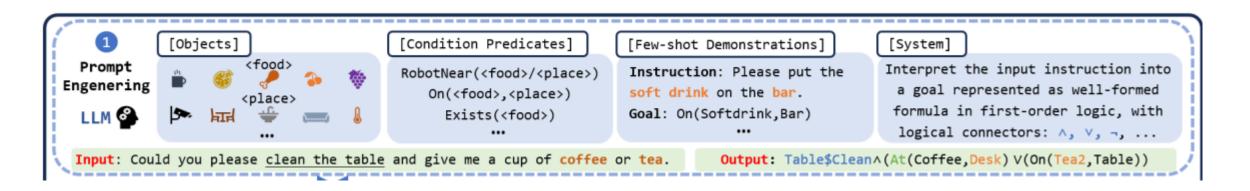
目标



LLM: 提示工程

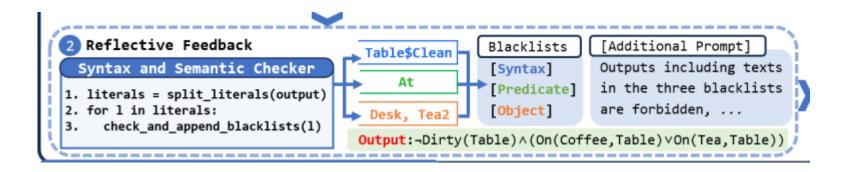
四个组成部分构成:

- (1) 对象。
- (2)条件谓词。
- (3) 少量示例演示。
- (4) 系统。





LLM: 反思性反馈



语法检查器评估输出是否构成良好公式

语义检查器则进一步验证谓词和对象的有效性。

增加了黑名单和额外的上下文,以指示LLM不要在输出中再次包含黑名单中的文本。



OBTEA

```
Output: (¬Dirty(Table) ∧ On(Coffee,Table))
(DNF) V
(¬Dirty(Table) ∧ On(Tea,Table))
```

```
Algorithm 1 OBTEA
Input: Goal G, init state s_0, action set A, transition rules P
Output: The optimal BT \mathcal{T}_*
 1: \mathcal{G}_{sub} \leftarrow ParseSubGoals(G)
                                                              b the set of sub goals
 2: for g_i \in \mathcal{G}_{sub} do
          \mathcal{T} \leftarrow Fallback(q_i)
                                                                           ⊳ init sub-BT
                                                \triangleright for \forall c, c \neq g_i, D(c) \leftarrow +\infty
          D(g_i) \leftarrow 0
         \mathcal{C}_U \leftarrow \{g_i\} \triangleright explored but unexpanded conditions
                                                            \mathcal{C}_E \leftarrow \emptyset
          while C_{IJ} \neq \emptyset do
              c \leftarrow \arg\min_{c \in \mathcal{C}_U} (D(c))
                                                           \triangleright explore and expand c
 8:
              for each a \in A do
 9:
                  if (c \cap (pre(a) \cup add(a) \setminus del(a)) \neq \emptyset) and (c \setminus a)
10:
                  del(a) = c) then
                      c_a \leftarrow pre(a) \cup c \setminus add(a)

    ► Exploration

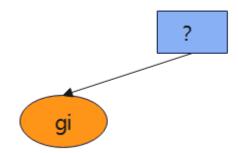
11:
                      if c_a \notin \mathcal{C}_E and D(c) + D(a) < D(c_a) then
12:
                          D(c_a) \leftarrow D(c) + D(a)
                                                                       \triangleright update D(c_a)
13:
                          \mathcal{C}_U \leftarrow \mathcal{C}_U \cup \{c_a\}
                                                                         \triangleright add c_a to \mathcal{C}_U
14:
                          \mathcal{M}(c_a) \leftarrow Sequence(c_a, a)
15:
                                                                \triangleright remove c from C_U
              \mathcal{C}_U \leftarrow \mathcal{C}_U \setminus \{c\}
16:
17:
              if c \neq q_i then
                  \mathcal{T} \leftarrow Fallback(\mathcal{T}, \mathcal{M}(c))
                                                                           ⊳ Expansion
18:
                 if c \subseteq s_0 then
19:
                      D(\mathcal{T}) \leftarrow D(c)
                                                                      \triangleright total cost of \mathcal{T}
20:
                      break \triangleright \mathcal{T} is the optimal BT for sub-goal q_i
21:
              C_E \leftarrow C_E \cup \{c\}
                                                                           \triangleright add c to \mathcal{C}_E
          \mathcal{T}_i \leftarrow Compact(\mathcal{T})
                                                                         ▷ Compaction
24: \mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{sub}|} = SortAscending(\mathcal{T}_1, \mathcal{T}_2, ..., \mathcal{T}_{|\mathcal{G}_{sub}|})
25: \mathcal{T}_* \leftarrow Fallback(\mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{sub}|})
                                                                             ⊳ Assembly
26: return T∗
```

输入: LLM最后的输出G,初始状态,动作谓词,状态转移规则

步骤1:初始化

目标: G=(|1^|2) \(\(|1^|3 \) Gsub=\{g1, g2\} g1=\{|1, |2\}, g2=\{|1, |3\}

T, D, CU, CE





OBTEA

```
Algorithm 1 OBTEA
Input: Goal G, init state s_0, action set A, transition rules P
Output: The optimal BT \mathcal{T}_*
 1: \mathcal{G}_{sub} \leftarrow ParseSubGoals(G)
                                                              b the set of sub goals
 2: for g_i \in \mathcal{G}_{sub} do
          \mathcal{T} \leftarrow Fallback(q_i)
                                                                           ⊳ init sub-BT
                                                \triangleright for \forall c, c \neq g_i, D(c) \leftarrow +\infty
          D(g_i) \leftarrow 0
                                  > explored but unexpanded conditions
          \mathcal{C}_U \leftarrow \{g_i\}
         \mathcal{C}_E \leftarrow \emptyset
                                                            while C_{IJ} \neq \emptyset do
              c \leftarrow \arg\min_{c \in \mathcal{C}_U} (D(c))
                                                            \triangleright explore and expand c
 8:
              for each a \in A do
 9:
                  if (c \cap (pre(a) \cup add(a) \setminus del(a)) \neq \emptyset) and (c \setminus a)
10:
                  del(a) = c) then
                      c_a \leftarrow pre(a) \cup c \setminus add(a)
11:

    ► Exploration

                      if c_a \notin \mathcal{C}_E and D(c) + D(a) < D(c_a) then
12:
                          D(c_a) \leftarrow D(c) + D(a)
                                                                       \triangleright update D(c_a)
13:
                          C_U \leftarrow C_U \cup \{c_a\}
                                                                        \triangleright add c_a to \mathcal{C}_U
14:
                          \mathcal{M}(c_a) \leftarrow Sequence(c_a, a)
15:
              \mathcal{C}_U \leftarrow \mathcal{C}_U \setminus \{c\}
                                                                \triangleright remove c from C_U
16:
17:
              if c \neq q_i then
                  \mathcal{T} \leftarrow Fallback(\mathcal{T}, \mathcal{M}(c))
                                                                           ⊳ Expansion
18:
                  if c \subseteq s_0 then
19:
                      D(\mathcal{T}) \leftarrow D(c)
                                                                      \triangleright total cost of \mathcal{T}
20:
                      break \triangleright \mathcal{T} is the optimal BT for sub-goal g_i
21:
              \mathcal{C}_E \leftarrow \mathcal{C}_E \cup \{c\}
                                                                           \triangleright add c to \mathcal{C}_E
          \mathcal{T}_i \leftarrow Compact(\mathcal{T})
                                                                         ▷ Compaction
24: \mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{sub}|} = SortAscending(\mathcal{T}_1, \mathcal{T}_2, ..., \mathcal{T}_{|\mathcal{G}_{sub}|})
25: \mathcal{T}_* \leftarrow Fallback(\mathcal{T}_1', \mathcal{T}_2', ..., \mathcal{T}_{|\mathcal{G}_{sub}|}')
                                                                            ▶ Assembly
26: return T∗
```

步骤2:探索

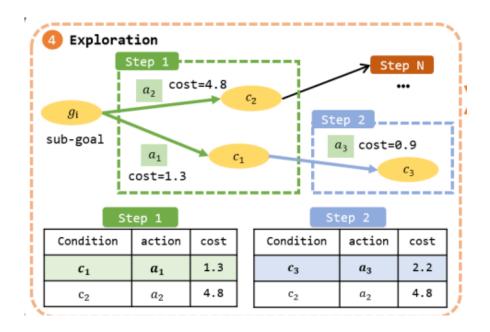
CU非空

邻近条件的生成Ca,忽略条件

 $1>ca \in CE$,

 $2>D(c) + D(a) \ge D(ca)$

否则更新 D(ca), CU





OBTEA

```
Algorithm 1 OBTEA
Input: Goal G, init state s_0, action set A, transition rules P
Output: The optimal BT \mathcal{T}_*
 1: \mathcal{G}_{sub} \leftarrow ParseSubGoals(G)
                                                              b the set of sub goals
 2: for g_i \in \mathcal{G}_{sub} do
          \mathcal{T} \leftarrow Fallback(q_i)
                                                                          ⊳ init sub-BT
                                               \triangleright for \forall c, c \neq g_i, D(c) \leftarrow +\infty
          D(q_i) \leftarrow 0
         \mathcal{C}_U \leftarrow \{g_i\} > explored but unexpanded conditions
                                                           \mathcal{C}_E \leftarrow \emptyset
          while C_{IJ} \neq \emptyset do
 8:
              c \leftarrow \arg\min_{c \in \mathcal{C}_U} (D(c))
                                                           \triangleright explore and expand c
              for each a \in A do
 9:
                  if (c \cap (pre(a) \cup add(a) \setminus del(a)) \neq \emptyset) and (c \setminus a)
10:
                 del(a) = c) then
                      c_a \leftarrow pre(a) \cup c \setminus add(a)
11:

    ► Exploration

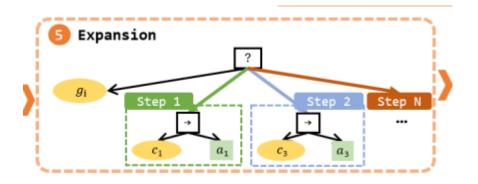
                      if c_a \notin \mathcal{C}_E and D(c) + D(a) < D(c_a) then
12:
                          D(c_a) \leftarrow D(c) + D(a) \triangleright update D(c_a)
13:
                         \mathcal{C}_U \leftarrow \mathcal{C}_U \cup \{c_a\}
                                                                        \triangleright add c_a to \mathcal{C}_U
14:
                          \mathcal{M}(c_a) \leftarrow Sequence(c_a, a)
15:
              \mathcal{C}_U \leftarrow \mathcal{C}_U \setminus \{c\}
                                                                \triangleright remove c from C_U
16:
17:
              if c \neq q_i then
                 \mathcal{T} \leftarrow Fallback(\mathcal{T}, \mathcal{M}(c))
                                                                          ⊳ Expansion
18:
                  if c \subseteq s_0 then
19:
                      D(\mathcal{T}) \leftarrow D(c)
                                                                      \triangleright total cost of \mathcal{T}
20:
                      break \triangleright \mathcal{T} is the optimal BT for sub-goal g_i
21:
              C_E \leftarrow C_E \cup \{c\}
                                                                          \triangleright add c to \mathcal{C}_E
          \mathcal{T}_i \leftarrow Compact(\mathcal{T})
                                                                        ▷ Compaction
24: \mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{sub}|} = SortAscending(\mathcal{T}_1, \mathcal{T}_2, ..., \mathcal{T}_{|\mathcal{G}_{sub}|})
25: \mathcal{T}_* \leftarrow Fallback(\mathcal{T}_1', \mathcal{T}_2', ..., \mathcal{T}_{|\mathcal{G}_{sub}|}')
                                                                            ⊳ Assembly
26: return T∗
```

步骤3: 拓展

添加子节点M(c)

如果未达到 sO, 我们将 c添加到扩展集合 CE 中

继续探索条件空间,直到扩展到s0





OBTEA

Algorithm 1 OBTEA

```
Input: Goal G, init state s_0, action set A, transition rules P
Output: The optimal BT \mathcal{T}_*

1: \mathcal{G}_{sub} \leftarrow ParseSubGoals(G) \triangleright the set of sub goals
```

```
b the set of sub goals
 2: for g_i \in \mathcal{G}_{sub} do
          \mathcal{T} \leftarrow Fallback(q_i)
                                                                           ⊳ init sub-BT
                                                \triangleright for \forall c, c \neq g_i, D(c) \leftarrow +\infty
          D(q_i) \leftarrow 0
         \mathcal{C}_U \leftarrow \{g_i\} > explored but unexpanded conditions
                                                            \mathcal{C}_E \leftarrow \emptyset
          while C_{IJ} \neq \emptyset do
 8:
              c \leftarrow \arg\min_{c \in \mathcal{C}_U}(D(c))
                                                            \triangleright explore and expand c
              for each a \in A do
 9:
                  if (c \cap (pre(a) \cup add(a) \setminus del(a)) \neq \emptyset) and (c \setminus a)
10:
                  del(a) = c) then
                      c_a \leftarrow pre(a) \cup c \setminus add(a)
11:

    ► Exploration

                      if c_a \notin \mathcal{C}_E and D(c) + D(a) < D(c_a) then
12:
                          D(c_a) \leftarrow D(c) + D(a) \triangleright update D(c_a)
13:
                          C_U \leftarrow C_U \cup \{c_a\}
                                                                         \triangleright add c_a to \mathcal{C}_U
14:
                          \mathcal{M}(c_a) \leftarrow Sequence(c_a, a)
15:
              \mathcal{C}_U \leftarrow \mathcal{C}_U \setminus \{c\}
                                                                 \triangleright remove c from C_U
16:
              if c \neq q_i then
17:
                  \mathcal{T} \leftarrow Fallback(\mathcal{T}, \mathcal{M}(c))
                                                                           ⊳ Expansion
18:
                  if c \subseteq s_0 then
19:
                      D(\mathcal{T}) \leftarrow D(c)
20:
                                                                      \triangleright total cost of \mathcal{T}
                      break \triangleright \mathcal{T} is the optimal BT for sub-goal q_i
21:
              C_E \leftarrow C_E \cup \{c\}
                                                                           \triangleright add c to \mathcal{C}_E
          \mathcal{T}_i \leftarrow Compact(\mathcal{T})
                                                                         ▷ Compaction
24: \mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{sub}|} = SortAscending(\mathcal{T}_1, \mathcal{T}_2, ..., \mathcal{T}_{|\mathcal{G}_{sub}|})
```

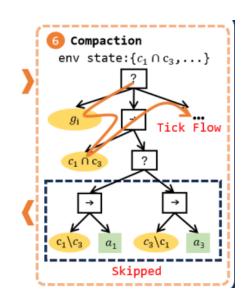
⊳ Assembly

25: $\mathcal{T}_* \leftarrow Fallback(\mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_{|\mathcal{G}_{-1}|})$

26: **return** *T*∗

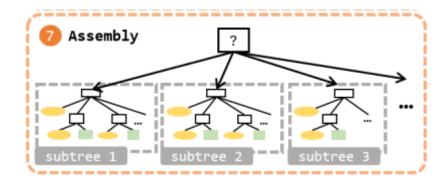
步骤4:压缩

压缩树结构并提高运行时效率。



步骤5:组装

所有子树生成完毕,我们首先 按照总成本的升序对所有子树 进行排序

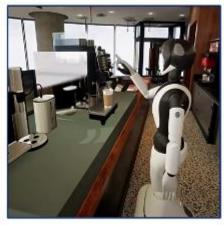




部署场景









酒店行业场景中部署了该框架,其中机器人充当咖啡馆的服务员

- 80 种可用物体
- 8个条件谓词: Pc = {RobotNear, Holding, On, Closed, Exists, Dirty, Activate, Low}
- 6 个动作谓词 Pa = {Make, MoveTo, PickUp, PutDown, Clean, Turn}



LLM

Difficulty	Instructions	Examples	Goals	Demonstrations	GA-NF(%)	GA-1F(%)	GA-5F(%)	IA(%)
Easy	30		RobotNear(Bar)	Zero-Shot	51.7	54.9	64.9	26.7
		Please come to the bar.		One-Shot	88.3	93.3	98.3	83.3
				Five-Shot	92.7	99.9	99.9	91.9
Medium	30	Please make sure there's enough water and the floor stays clean.	Exist(Water) ¬Dirty(Floor)	Zero-Shot	49.1	54.3	71.3	20.4
				One-Shot	81.1	88.9	98.9	75.6
				Five-Shot	91.1	99.9	99.9	83.3
Hard		I'm at table three, bring some	On(Dessert, Table3) A	Zero-Shot	38.3	52.5	65.0	12.5
		dessert and yogurt. If there is no	(On(Yogurt,Table3) V	One-Shot	72.5	84.9	88.3	50.8
		yogurt, coffee is fine.	On(Coffee, Table3))	Five-Shot	87.1	92.3	94.9	76.9

评估数据集:三个难度级别进行了分类。

两个指标进行评估:语法准确性(GA)和解释准确性(IA)。

研究了两个因素对我们方法的影响:示例的数量和最大重试次数。





OBTEA

Case	Scenario Configuration			Problem		Total Costs		Condition Node Ticks				
	$ \mathcal{O} $	$ \mathcal{P}_c $	$ \mathcal{P}_a $	MAC	Literals	States	Actions	Baseline	OBTEA	Baseline	OBTEA	OBTEA-NC
0	100	10	10	0	995.4	475.5	509.5	71.3	58.5	31.7	23.8	33.1
1	100	10	50	0	1000.0	4046.5	2495.9	52.0	26.8	25.7	22.7	33.8
2	500	50	50	0	25000.0	22468.6	12508.2	164.2	137.0	950.4	632.3	1779.6
3	100	10	10	5	995.1	466.1	527.3	71.7	26.1	54.0	28.5	54.7
4	100	30	10	5	2984.4	501.0	530.4	170.9	81.1	597.1	358.2	1005.1
5	100	50	10	5	4974.6	497.7	524.4	241.4	124.8	1581.2	1210.5	3578.8
6	100	50	30	5	5000.0	2496.0	1527.4	189.4	97.7	1391.9	1147.9	3111.3
7	100	50	50	5	5000.0	4501.6	2531.5	169.4	85.1	1584.4	971.8	2848.0
8	300	50	50	5	15000.0	13550.7	7560.9	164.2	82.5	1302.4	892.8	2577.2
9	500	50	50	5	25000.0	22427.9	12514.2	165.9	82.7	1436.6	909.8	2581.2

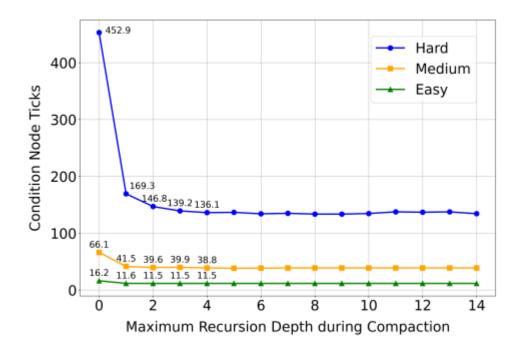
动作成本是基于人类经验设定的。

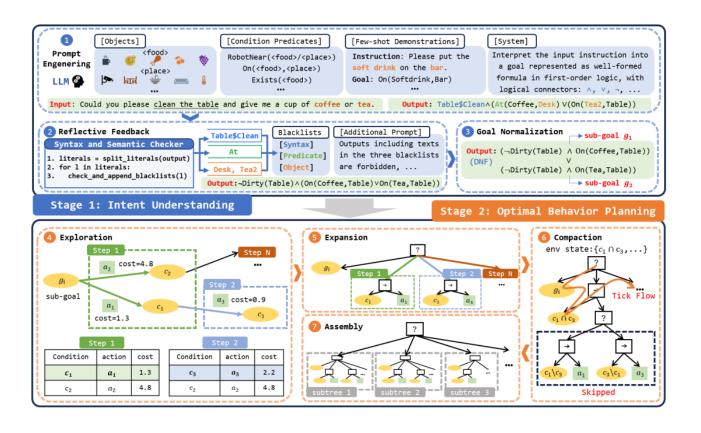
关注两个关键指标: BT的总成本和BT执行过程中条件节点的计数。



OBTEA

对压缩过程中最大递归深度的影响进行了消融分析





在第一阶段,输入指令通过大语言模型(LLM)与提示工程和反思反馈转化为逻辑表达的目标。该目标随后被规范化为合取范式(DNF)并分解为子目标。在第二阶段,通过探索、扩展和压缩为每个子目标生成一个子树。这些子树最终被组装以创建最终的最优决策树

结束





- 1.传参时是否有把动作,环境信息传入
- 2.黑名单的使用

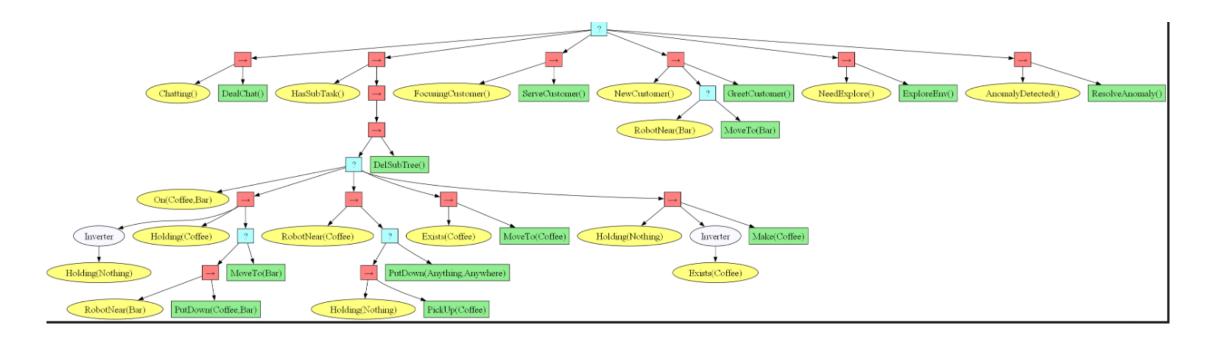
```
RobotNear_<items_place>, On_<items>_cplace>, Holding_<items>, Exists_<items>, IsClean_<furniture>, Active_<appliance>, Closed_<furnishing>, Low
     <items>=['Coffee', 'Water', 'Dessert', 'Softdrink', 'BottledDrink', 'Yogurt', 'ADMilk', 'MilkDrink', 'Milk', 'VacuumCup', 'Chips', 'NFCJuice',
     <place>=['Bar', 'Bar2', 'WaterStation', 'CoffeeStation', 'Table1', 'Table2', 'Table3', 'WindowTable6', 'WindowTable4', 'WindowTable5', 'QuietTabl
     <items place>=<items>+<place>
     <furniture>=['Table1', 'Floor', 'Chairs']
    <appliance>=['AC','TubeLight','HallLight']
     <furnishing>=['Curtain']
    <control>=['ACTemperature']
     [Examples]
     Instruction: "Would you be able to provide some chips at the third table?"
    Goal: On Chips Table3
     Instruction: If you have time, could you please hold the milk for me, or pull the curtains, or perhaps switch on the AC?
    Closed_Curtain | Active_AC | Holding_Milk
    Instruction: Please turn up the air conditioning and come to the bar counter.
    RobotNear Bar & ~Low ACTemperature
    Instruction: Please ensure the water is ready for service, and deliver the yogurt to table number one.
    Exists_Water & On_Yogurt_Table1
    Instruction: Let's not have the AC too cold. Could you turn off these tube lights and open the curtains for some natural light?
     ~Low ACTemperature & ~Active TubeLight & ~Closed Curtain
     [Condition] Lists all predicates representing conditions and their optional parameter sets.
     [Object] Lists all parameter sets.
     [Examples] Provide several examples of Instruction to Goal mapping.
40 Your task is to interpret customer instructions within a café setting and translate them into specific goals using logical expressions. Utilize
41 Apply logical operators (&, |, ~) to combine these elements appropriately.
42 & (AND Operator): Combines conditions such that the result is true only if both conditions are true.
43 (OR Operator): Combines conditions such that the result is true if at least one of the conditions is true.
44 ~ (NOT Operator): Negates or reverses the truth value of a single condition.
45 This process should accurately map each given instruction to a clear, actionable goal in the context of the café environment.
46 Please generate directly interpretable predicate formulas without additional explanations.
    The predicate formulas can be converted into disjunctive paradigms (DNFs) using the python package sympy.to dnf.
49 % 不要出现 [Blacklist] 中单词,若出现<Illegal Condition>中的单词,Please select the closest predicates from the [Condition] table to form the an
     [Blacklist]
    <Illegal Condition>=[]
55 <Illegal Object>=[Fries, Table6]
     <Other Illegal Words or Characters>=[]
     [Blacklist] Contains restricted elements.
    Do not include words from the [Blacklist] in the predicate formulas.
     If a word from <Illegal Condition> is encountered, select the most relevant predicates from the [Condition] table to construct the answer.
     If a word from <Illegal Object> is encountered, choose the nearest parameter from the [Object] table to formulate the answer.
```





3.生成的逻辑与优先级有关吗

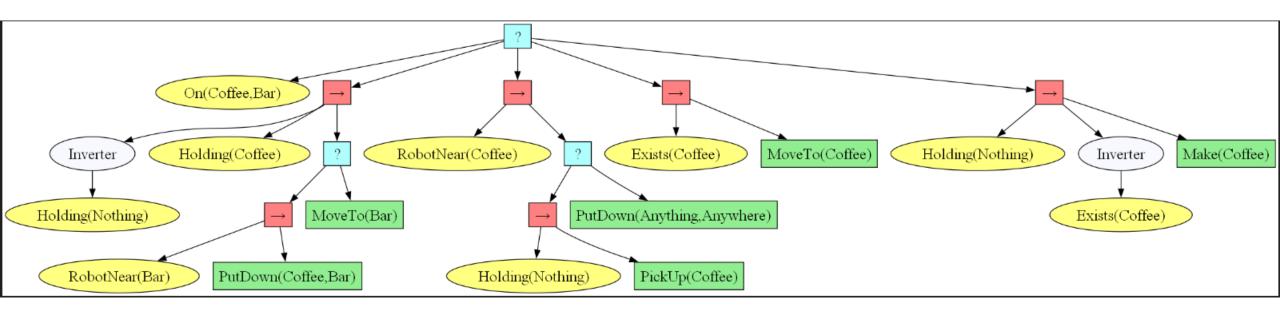
给我一杯咖啡







4.行为树的执行步骤







5.涉及到环境中没有的物体或者没有的动作他会怎么做(稳定性)

```
2024-10-18 15:10:19.612 | INFO | robowaiter.llm_client.multi_rounds:deal_response:171 - Function Call Response: {'name': 'create_sub_task', 'arguments': '{\n "goal": "{Prepare(Water)},{Prepare(Dessert)}, {Prepare(Coffee)}, '\n}'}

goal_set: [{',{Prepare(Dessert)', '{Prepare(Water)', ',{Prepare(Coffee)', ')'}}]
```

NameError: name 'Prepare' is not defined