

MONASH INFORMATION TECHNOLOGY

FIT2004 Algorithms and Data Structures

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Referencing materials by Nathan Companez, Aamir Cheema, Arun Konagurthu and Lloyd Allison





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COMMONWEALTH OF AUSTRALIA

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Ready?

Agenda

Directed Acyclic Graph (DAG)



Agenda

- Directed Acyclic Graph (DAG)
- Topological Sort
 - Kahn's algorithm
 - Depth-First Search (DFS) modification





Let us begin...



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
- What is a DAG?



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
- What is a DAG?
 - Directed



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
- What is a DAG?
 - Directed
 - Acyclic



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
- What is a DAG?
 - Directed
 - Acyclic
 - and of course it is a Graph...



- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
 - Can you give me an example of a real world DAG?
- What is a DAG?
 - Directed
 - Acyclic
 - and of course it is a Graph...

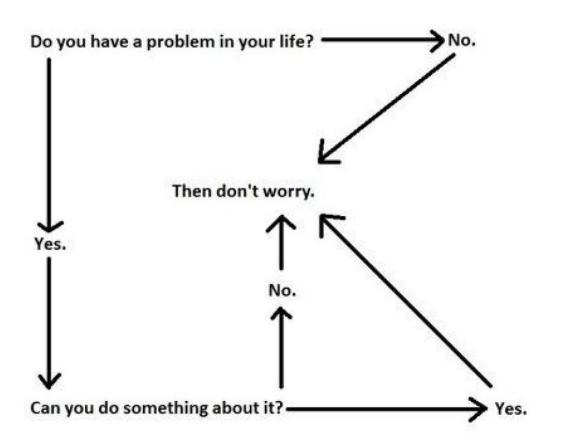


- Graphs are very commonly used to model real world scenario
 - One of which is a DAG
 - Can you give me an example of a real world DAG?
 - Can you give me an example of a real world non-DAG?
- What is a DAG?
 - Directed
 - Acyclic
 - and of course it is a Graph...

What is it?



A real world DAG



What is it?

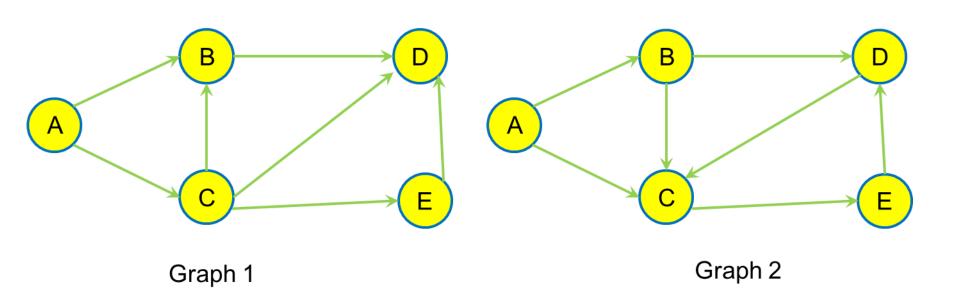
A real world not DAG







- Which graph is a DAG?
 - And where is the cycle?





Questions?

What is it for?



Prerequisite mapping

What is it for?



- Prerequisite mapping
 - If I have a directed edge from A to B, this means I need A before B

What is it for?



- Prerequisite mapping
 - If I have a directed edge from A to B, this means I need A before B
 - Common in project management
 - Common in talent/skill trees!



What is it for?



Prerequisite mapping

- If I have a directed edge from A to B, this means I need A before B
 - Common in project management
 - Common in talent/skill trees!
 - Your university units!
 - Pass 1045
 - Pass 1008
 - Pass 2004
 - Pass 3155

What is it for?



So for an edge <A,B>

What is it for?



- So for an edge <A,B>
 - A is a prerequisite for B
 - A is an ancestor of B
 - A is the subset of B
 - A is ordered before B

What is it for?



- So for an edge <A,B>
 - A is a prerequisite for B
 - A is an ancestor of B
 - A is the subset of B
 - A is ordered before B
 - This enable us to sort a DAG



Questions?

Ordering of Vertices



- A topological sort
 - Permutation of vertices in a DAG
 - Vertex U will appear before vertex V if we have edge <U,V>

Ordering of Vertices



A topological sort

- Permutation of vertices in a DAG
- Vertex U will appear before vertex V if we have edge <U,V>
- Vertex U will appear before vertex W is we have edge <U,W>

Ordering of Vertices



A topological sort

- Permutation of vertices in a DAG
- Vertex U will appear before vertex V if we have edge <U,V>
- Vertex U will appear before vertex W is we have edge <U,W>
- But if we don't have edge <V,W> then V and W are of the same order

Ordering of Vertices



A topological sort

- Permutation of vertices in a DAG
- Vertex U will appear before vertex V if we have edge <U,V>
 - U<V
- Vertex U will appear before vertex W is we have edge <U,W>
 - U<W
- But if we don't have edge <V,W> then V and W are of the same order
 - V==W

Ordering of Vertices

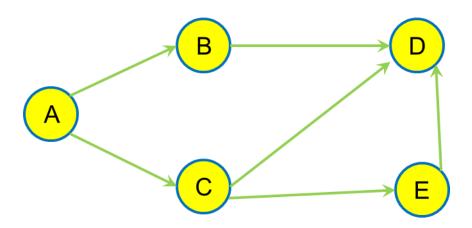


- A topological sort
 - Permutation of vertices in a DAG
 - Vertex U will appear before vertex V if we have edge <U,V>
 - U<V
 - Vertex U will appear before vertex W is we have edge <U,W>
 - U<W
 - But if we don't have edge <V,W> then V and W are of the same order
 - V==W
- So we have a DAG of your units
- Topological sort of this DAG gives you the order of units to take!

Ordering of Vertices



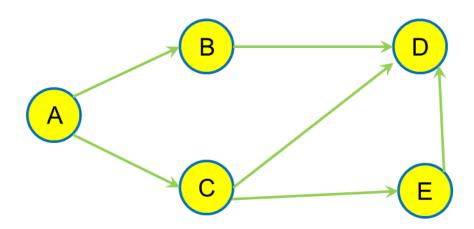
Which one of these are not a valid topological sort of the DAG?



Ordering of Vertices



- Which one of these are not a valid topological sort of the DAG?
 - 1. A, B, C, E, D
 - 2. A ,C, B, E, D
 - 3. A, C, E, B, D
 - 4. A, B, E, C, D



Ordering of Vertices



- Topological sort can be done via
 - Kahn's algorithm
 - A modified DFS



Questions?

Kahn's Algorithm

For topological sort



What is the concept like?

Kahn's Algorithm

For topological sort



- What is the concept like?
 - Start with vertices without incoming edges



- What is the concept like?
 - Start with vertices without incoming edges
 - Delete all outgoing edges from the vertex



- What is the concept like?
 - Start with vertices without incoming edges
 - Delete all outgoing edges from the vertex
 - Add in vertices without incoming edges



- What is the concept like?
 - Start with vertices without incoming edges
 - Delete all outgoing edges from the vertex
 - Add in vertices without incoming edges
 - Repeat!

For topological sort



Algorithm as follow

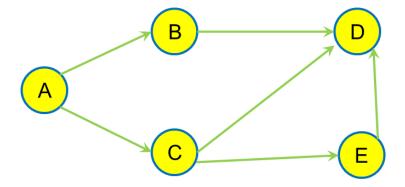
```
sorted_list = []
  process = []
  add all vertices without incoming edges to process
□ while len(process) > 0:
      vertex u = process.pop()
      sorted list.append(vertex u)
      for edge in vertex_u.edges:
          remove edge from graph
          if edge.vertex v has no incoming edges:
              process.append(vertex v)
∃ if graph still has edges:
      print("Error. Not a cycle")
∃ else:
      print(sorted list)
```



- Algorithm as follow
 - Let us try it out

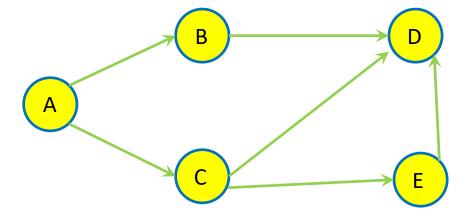
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∃ if graph still has edges:
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- Algorithm as follow
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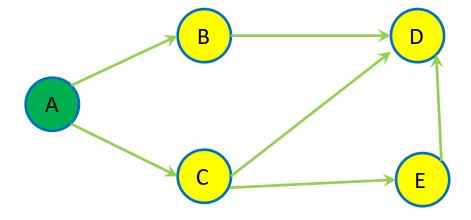
```
process process
```

- Algorithm as follow
 - Let us try it out



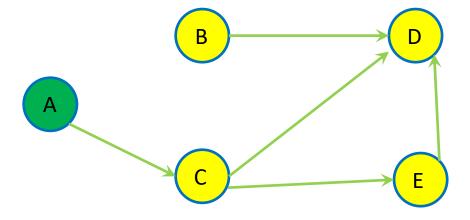
sorted			
process	А		

- Algorithm as follow
 - Let us try it out



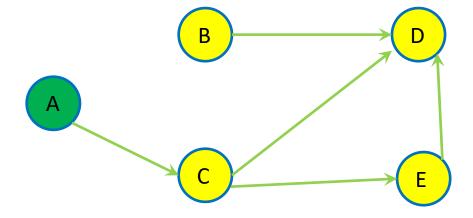
sorted	A		
process			

- Algorithm as follow
 - Let us try it out



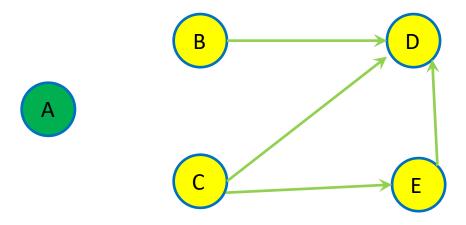
```
sorted A process
```

- Algorithm as follow
 - Let us try it out



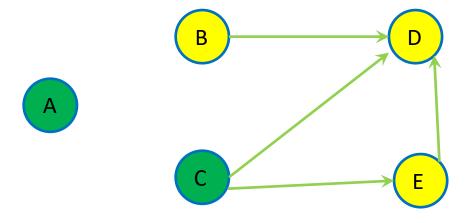
```
sortedAprocessB
```

- Algorithm as follow
 - Let us try it out



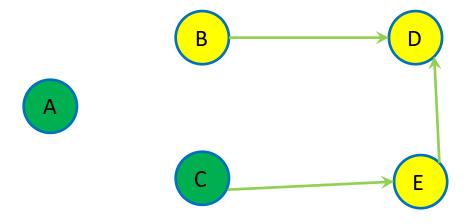
sorted	Α			
process	В	С		

- Algorithm as follow
 - Let us try it out



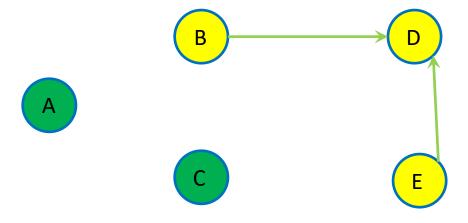
```
sortedACprocessB
```

- Algorithm as follow
 - Let us try it out



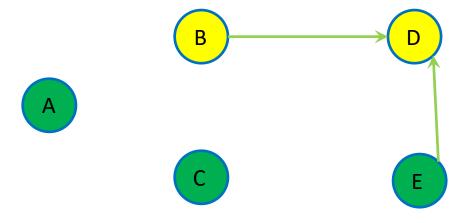
```
sortedACprocessB
```

- Algorithm as follow
 - Let us try it out



```
sortedACprocessBE
```

- Algorithm as follow
 - Let us try it out



```
sortedACEprocessB
```

- Algorithm as follow
 - Let us try it out









sorted	A	С	E	
process	В			

- Algorithm as follow
 - Let us try it out









sorted	A	С	E	В	
process					

- Algorithm as follow
 - Let us try it out











sorted	Α	С	E	В	
process	D				

- Algorithm as follow
 - Let us try it out













- Algorithm as follow
 - Seemed simple right?
 - Let us zoomed in to the algorithm more



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 - Seemed simple right?
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```
# for output
  sorted list = []
 # tracks number of incoming edges
  incoming edges = [0] * len(vertices)

        ∃ for edge in graph:

                               # edge <u,v>
      incoming edges[vertex_v] += 1
  # process queue or stack
  process = []

    □ for vertex v in incoming edges:
      if incoming edges[vertex v] == 0:
          process.append(vertex v)
  # kahn's
⊟ while len(process) > 0:
      vertex u = process.pop()
      sorted list.append(vertex u)
      for edge in vertex u.edges:
          incoming edges[edge.vertex v] -= 1
          if incoming edges[vertex_v] == 0:
              process.append(vertex v)
  # results
☐ if graph still has edges:
      print("Error. Not a cycle")
∃ else:
      print(sorted list)
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- Algorithm as follow
 - Seemed simple right?
 - Let us zoomed in to the algorithm more
- Complexity?

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          incoming edges[edge.vertex v] -= 1
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               process.append(vertex v)
  # results

☐ if graph still has edges:
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```



- Algorithm as follow
 - Seemed simple right?
 - Let us zoomed in to the algorithm more
- Complexity?
 - O(V+E) time space

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```



Questions?

Modified for topological sorting



We saw the complexity of Kahn's being O(V+E)

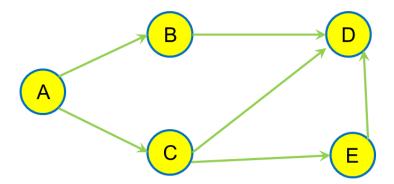


- We saw the complexity of Kahn's being O(V+E)
- Can we modify DFS to do the same?



- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see

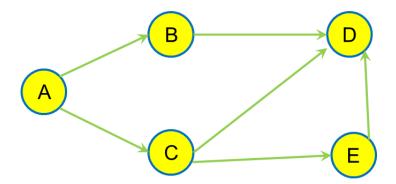
```
28  def dfs_topological(vertex_u):
29     vertex_u.visited = True
30  for edge in vertex_u.edges:
31  fedge.vertex_v.visited == False:
32     dfs_topological(vertex_v)
```





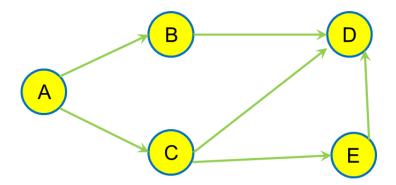
- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A

```
28  def dfs_topological(vertex_u):
29     vertex_u.visited = True
30  for edge in vertex_u.edges:
31  def dfs_topological(vertex_v)
```



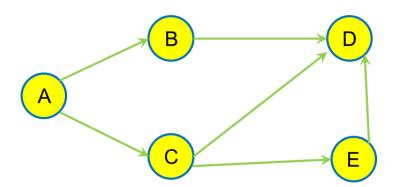


- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B



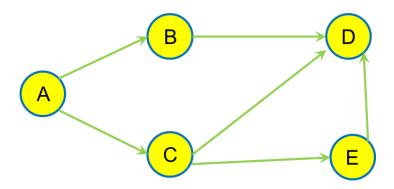


- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B
 - Go to D



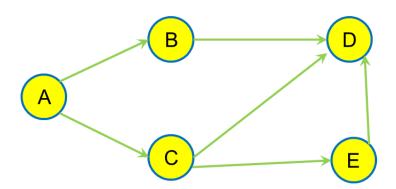


- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B
 - Go to D
 - Go to C



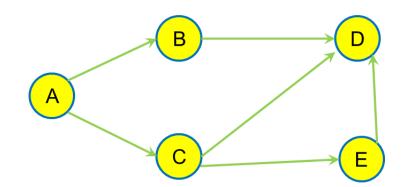


- Can we modify DFS to do the same?
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 - Go to B
 - Go to D
 - Go to C
 - Go to E





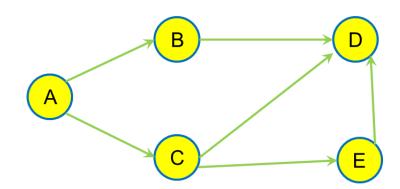
- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B
 - Go to D
 - Go to C
 - Go to E
 - So we have A, B, D, C, E





- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B
 - Go to D
 - Go to C
 - Go to E
 - So we have A, B, D, C, E
 - Any other DFS order?

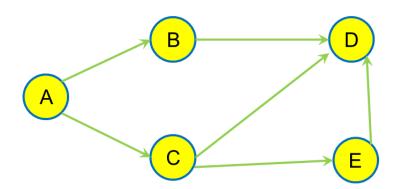
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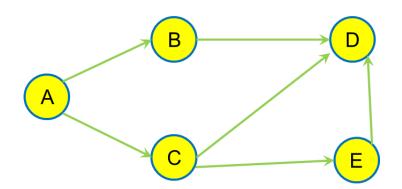
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 - So we have A, B, D, C, E
 - Any other DFS order?
 - A, C, D, E, B

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28  def dfs_topological(vertex_u):
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```





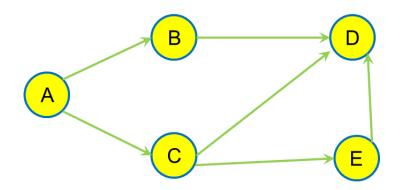
- Can we modify DFS to do the same?
 - Let us run DFS on the following graph and see
 - Start from A
 - Go to B
 - Go to D
 - Go to C
 - Go to E
 - So we have A, B, D, C, E
 - Any other DFS order?
 - A, C, D, E, B
 - A, C, E, D, B





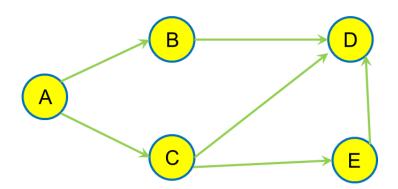
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 - Go to D
 - Go to C
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 - So we have A, B, D, C, E
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 - A, C, D, E, B
 - A, C, E, D, B
 - A, C, E, B, D?

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28  def dfs_topological(vertex_u):
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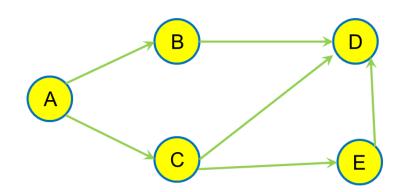


- Can we modify DFS to do the same?
 - Any other DFS order?
 - A, B, D, C, E
 - A, C, D, E, B
 - A, C, E, D, B





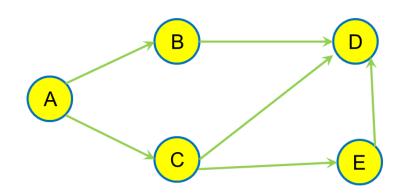
- Can we modify DFS to do the same?
 - Any other DFS order?
 - A, B, D, C, E
 - A, C, D, E, B
 - A, C, E, D, B
 - A possible topological sort
 - A, B, C, E, D
 - A, C, B, E, D





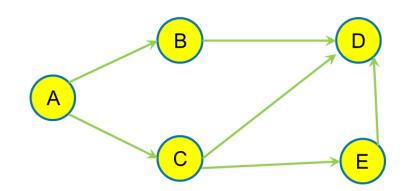
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 - A, C, B, E, D
 - Notice something?

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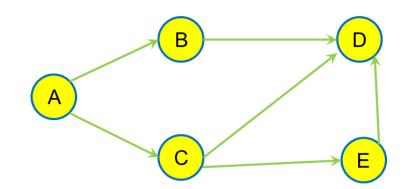


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 - A, C, D, E, B
 - A, C, E, D, B
 - A possible topological sort
 - A, B, C, E, D
 - A, C, B, E, D
 - Notice something?
 - When we reach the end of the DFS, we go back to an earlier vertex but this vertex should be early in topological sort (such as vertex B or C)



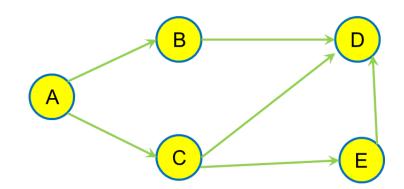


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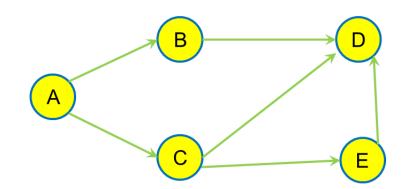
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- Can we modify DFS to do the same?
 - Any other DFS order?
 - A, B, D, C, E
 - A, C, D, E, B
 - A, C, E, D, B
 - A possible topological sort
 - A, B, C, E, D
 - A, C, B, E, D
 - Notice something?
 - When we reach the end of the DFS, we go back to an earlier vertex but this vertex should be early in topological sort (such as vertex B or C)
 - So how do we arrange it?

```
28  def dfs_topological(vertex_u):
29     vertex_u.visited = True
30  for edge in vertex_u.edges:
31  def dfs_topological(vertex_v)
```



Modified for topological sorting



Here's how we modify with a stack!

```
□ def dfs topological(vertex u):
      # start for result
      stack = []
      # run DFS
      vertex u.visited = True
      for edge in vertex u.edges:
          if edge.vertex v.visited == False:
              dfs topological aux(vertex v,stack)
      # output
      print(stack)
□ def dfs topological aux(vertex u,stack):
      vertex u.visited = True
      for edge in vertex u.edges:
          if edge.vertex v.visited == False:
              dfs topological aux(vertex v)
      # add to stack
      stack.push(vertex u)
```

Modified for topological sorting



Complexity?

```
□ def dfs_topological(vertex_u):
      # start for result
      stack = []
      # run DFS
      vertex u.visited = True
      for edge in vertex u.edges:
          if edge.vertex v.visited == False:
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      # output
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      for edge in vertex_u.edges:
          if edge.vertex v.visited == False:
              dfs topological aux(vertex v)
      # add to stack
      stack.push(vertex u)
```



- Complexity?
 - O(V+E) since we only added a stack

```
□ def dfs topological(vertex u):
      # start for result
      stack = []
      # run DFS
      vertex u.visited = True
      for edge in vertex u.edges:
          if edge.vertex v.visited == False:
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      for edge in vertex_u.edges:
          if edge.vertex v.visited == False:
              dfs topological aux(vertex v)
      # add to stack
      stack.push(vertex u)
```



Questions?



Thank You