

#### MONASH INFORMATION TECHNOLOGY

# FIT2004 Algorithms and Data Structures

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Referencing materials by Nathan Companez, Aamir Cheema, Arun Konagurthu and Lloyd Allison





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Ready?

## **Agenda**

- Sorting Algorithms
  - Comparison based
    - Selection
    - Insertion
  - Non-comparison based (the IMBA ones)
    - Counting
    - Radix





Let us begin...



- We are back to sorting!
  - Bubble
  - Insertion
  - Selection
  - Merge
  - Quick



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  - Selection
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  - Quick



Janelle Shane @Janelle CShane · 14 Apr



For example, there was an algorithm that was supposed to sort a list of numbers. Instead, it learned to delete the list, so that it was no longer technically unsorted.



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- All of these are known as comparison based sorting.
   Why? Because we compare between items to know if a < b or b > a
- Now let us analyze them based on what we have learnt!



## Questions?



- Correctness
- Complexity



- Correctness
  - Loop invariant
  - Termination
- Complexity
  - Time
  - Space



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#### Selection Sort



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  - my\_list[0...i-1] is sorted
  - my\_list[0...i-1] <= my\_list[i...N]</p>
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- Termination
  - i and j always increment and both reach the end of the list
- So why is it working then?
  - i keep increment till n and we know from invariant 0...i-1 is sorted, thus we will sort the entire list!



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  - Time
    - Best = O(N^2)
    - Worst = O(N^2)
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    - Best = O(N^2) because no matter what we have to find the minimum and cant terminate earlier!
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    - O(N) for the input list
    - Auxiliary?



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    - O(N) for the input list
    - Auxiliary? O(1) in place



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      - Like comparing between words, you need to compare the alphabets



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      - We know complexity is based on comparison O(N^2) comparisons...



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    - Worst = O(N^2)
    - But what if I tell you comparing the items have a cost of O(k)
      - Like comparing between words, you need to compare the alphabets
      - We know complexity is based on comparison O(N^2) comparisons...
      - So our final complexity is O(kN^2)



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- Stable?



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  - Relative ordering doesn't change



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  - Is it stable?



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  - [4a, 2, 3, 4b, 1]



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  - [4a, 2, 3, 4b, 1]
  - Minimum is 1, so we swap
  - [1, 2, 3, 4b, 4a]



- Correctness
- Complexity
- Stable?
  - Relative ordering doesn't change
  - Is it stable? No! but why?
  - [4a, 2, 3, 4b, 1]
  - Minimum is 1, so we swap
  - [1, 2, 3, 4b, 4a]
  - Now we see that 4a is behind 4b!



# Questions?



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- Complexity

#### **Insertion Sort**



- Correctness
- Complexity

**Problem 1.** Write psuedocode for insertion sort, except instead of sorting the elements into non-decreasing order, sort them into non-increasing order. Identify a useful invariant of this algorithm.



- Correctness
- Complexity

```
def insertion_sort(my_list):
    for i in range(1, len(my_list)):
        key = my_list[i]
        j = i - 1
        # keep shifting to left if left is greater
        while j >= 0 and key < my_list[j]:
            my_list[j+1] = my_list[j]
            j = j - 1
        my_list[j+1] = key</pre>
```



- Correctness
  - Loop invariant
  - Termination
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- Correctness
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- Correctness
  - Loop invariant
    - my\_list[0...i-1] sorted
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- Correctness
- Complexity
  - Best O(N) comparison
    - Each loop only look and compare with left item once
  - Worst

```
def insertion_sort(my_list):
    for i in range(l, len(my_list)):
        key = my_list[i]
        j = i - l
        # keep shifting to left if left is greater
        while j >= 0 and key < my_list[j]:
            my_list[j+1] = my_list[j]
            j = j - l
        my_list[j+1] = key</pre>
```



- Correctness
- Complexity
  - Best O(N) comparison
    - Each loop only look and compare with left item once
  - Worst O(N^2)
    - Each loop keep look left, compare and swap till beginning of list

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  - What about space?

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    - Each loop keep look left, compare and swap till beginning of list
  - So if O(k) is the comparison cost, when we have O(kN^2) worst case!

def insertion sort(my list):

for i in range(l, len(my\_list)):

- What about space?
  - O(N) for the input list
  - O(1) auxiliary cause it is in-place

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key = my_list[i]
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- Correctness
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- Stability
  - Yes
  - Don't swap if value is the same
  - Most shifting will ensure stability

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# Questions?

# Sorting



|                | Best               | Worst   | Average            | Stable? | In-<br>place? |
|----------------|--------------------|---|--------------------|---------|---------------|
| Selection Sort | O(N <sup>2</sup> ) | $O(N^2)$  | O(N <sup>2</sup> ) | No      | Yes           |
| Insertion Sort | O(N)               | $O(N^2)$  | O(N <sup>2</sup> ) | Yes     | Yes           |
| Heap Sort      | O(N log N)         | O(N log N)  | O(N log N)         | No      | Yes           |
| Merge Sort     | O(N log N)         | O(N log N)  | O(N log N)         | Yes     | No            |
| Quick Sort     | O(N log N)         | O(N <sup>2</sup> ) – can<br>be made<br>O(N log N) | O(N log N)         | Depends | No            |

# Sorting



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| Selection Sort | O(N <sup>2</sup> ) | $O(N^2)$  | O(N <sup>2</sup> ) |          |                    |
| Insertion Sort | O(N)               | O(N <sup>2</sup> )                                | O(N <sup>2</sup> ) |          |                    |
| Heap Sort      | O(N log N)         | O(N log N)  | O(N log N)         | WHAT THE | menneyenerator.net |
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# **Auxiliary for Recursion**



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  - If I have recursion log N times, then I take O(log N) space for the recursion alone!
  - If each recursion is k, then my total space is O(k log N)!!!

# **Auxiliary for Recursion**



- The recursion stack takes up memory!!!
  - So that is why it isn't in-place!
    - Iterative is easier to get in-place
  - If I have recursion log N times, then I take O(log N) space for the recursion alone!
  - If each recursion is k, then my total space is O(k log N)!!!

# Sorting



|                | Best               | Worst   | Average            | Stable? | In-<br>place? |
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| Selection Sort | O(N <sup>2</sup> ) | O(N <sup>2</sup> )                                | O(N <sup>2</sup> ) | No      | Yes           |
| Insertion Sort | O(N)               | O(N <sup>2</sup> )                                | O(N <sup>2</sup> ) | Yes     | Yes           |
| Heap Sort      | O(N log N)         | O(N log N)  | O(N log N)         | No      | Yes           |
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# **Complexity**

#### Time – Lower Bound



- So... what is the lower bound for the sorting algorithms that we have learnt?
  - Bubble
  - Insertion
  - Selection
  - Quick
  - Merge
- These are all comparison based
- Ω(N log N)
- We will see more of this later



# Questions?



Thank you