

MONASH INFORMATION TECHNOLOGY

FIT2004 Algorithms and Data Structures

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Referencing materials by Rafael Dowsley, Nathan Companez, Aamir Cheema, Arun Konagurthu and Lloyd Allison





Faculty of Information Technology, Monash University

COMMONWEALTH OF AUSTRALIA

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Ready?

Agenda

- Circulation with Demands
- Circulation with Demands and Lower Bound
- Applications Example
 - Survey Design
 - Airline Scheduling





Let us begin...

Recall...



- You have learnt Graph
- You have learnt Network Flow

Recall...



- You have learnt Graph
 - How many problems can be modelled as a Graph, then be solved
- You have learnt Network Flow

Recall...



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 - How many problems can be modelled as a Graph, then be solved.
- You have learnt Network Flow
 - Likewise, we have explore the simple Bipartite Matching problem.

Recall...

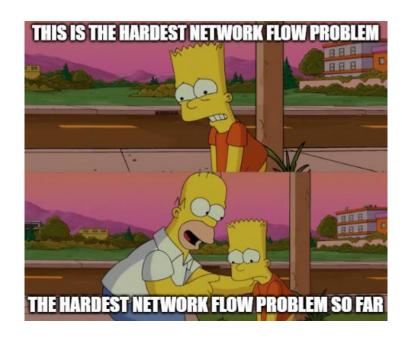


You have learnt Graph

 How many problems can be modelled as a Graph, then be solved.

You have learnt Network Flow

- Likewise, we have explore the simple Bipartite Matching problem.
- ... now let us push 1 step further!





Questions?

A Feasibility Problem...







Recall the 2 concepts from Network Flow

A Feasibility Problem...



- Recall the 2 concepts from Network Flow
 - Capacity Constraint
 - Flow <= Capacity for an edge
 - Flow conservation
 - Incoming flow to a vertex == outgoing flow from the vertex

A Feasibility Problem...

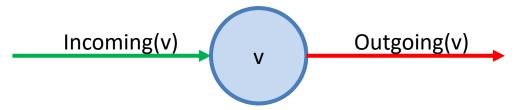


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 - Now what if we tweak this rule?

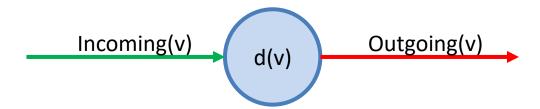
A Feasibility Problem...



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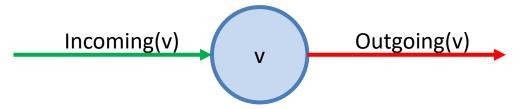
Now what if we tweak this rule?

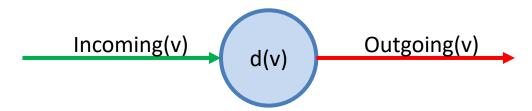


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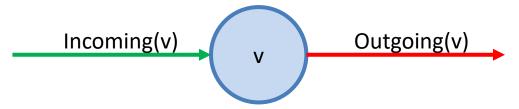


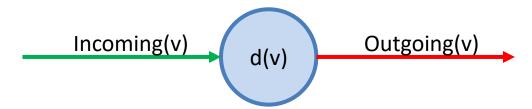


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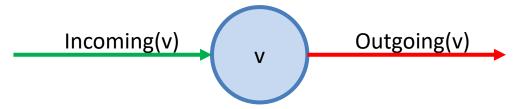


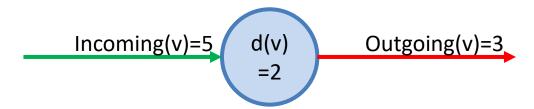


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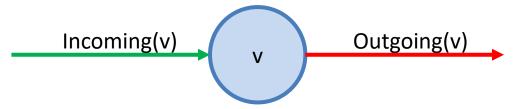


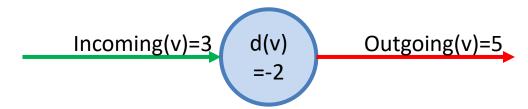


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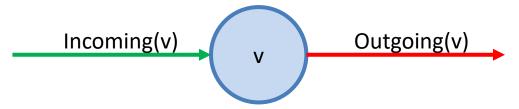


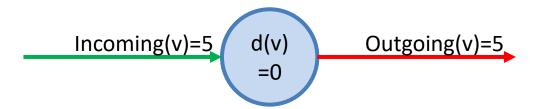


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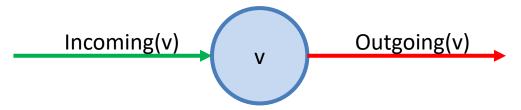


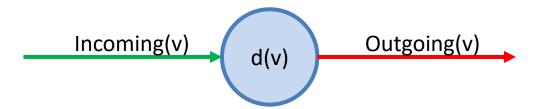


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A Feasibility Problem...

- Recall the 2 concepts from Network Flow
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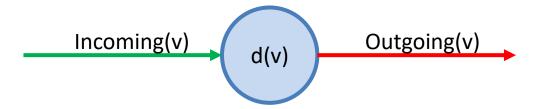




A Feasibility Problem...



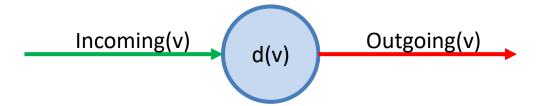
- Recall the 2 concepts from Network Flow
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 - Demand Constraint
 - Now what if we tweak this rule? incoming(v) outgoing(v) = demand(v)



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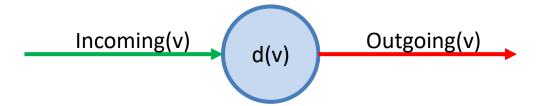


Circulation with Demands is a feasibility problem

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A Feasibility Problem...

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 - Now what if we tweak this rule? incoming(v) outgoing(v) = demand(v)



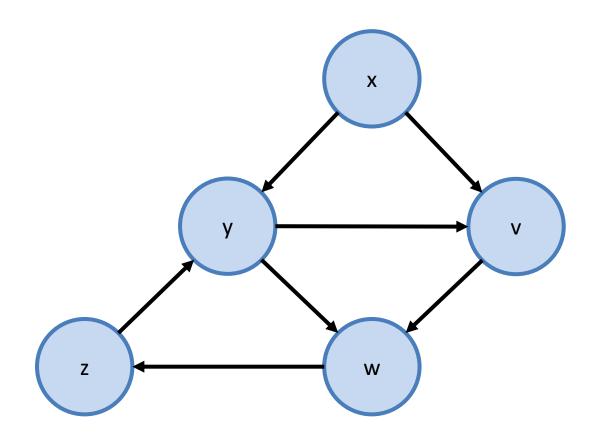
Circulation with Demands is a feasibility problem that satisfy both of the above!



Questions?

A Feasibility Problem...





A Feasibility Problem...



An example below...

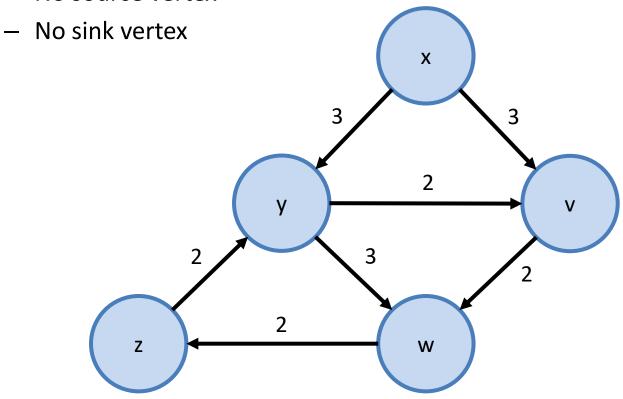
 With the usual capacity Χ 2 W

A Feasibility Problem...



An example below...

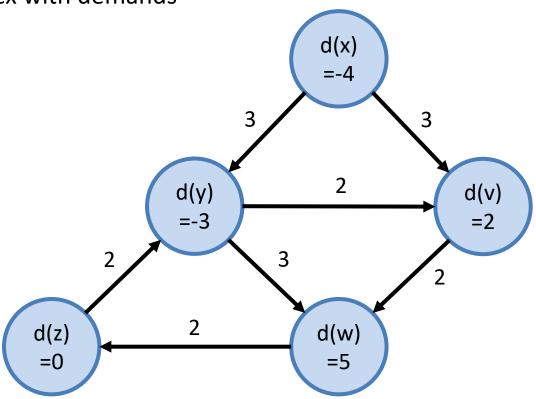
No source vertex



A Feasibility Problem...



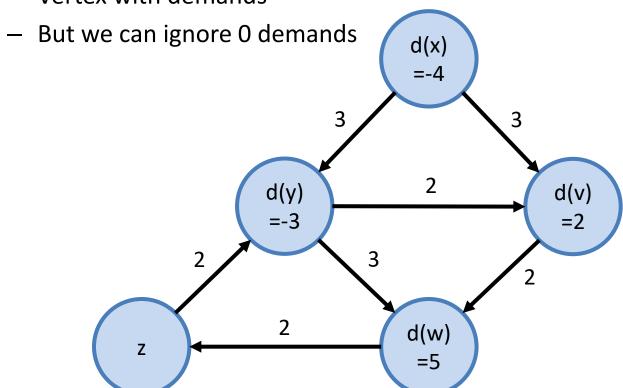
- An example below...
 - Vertex with demands



A Feasibility Problem...

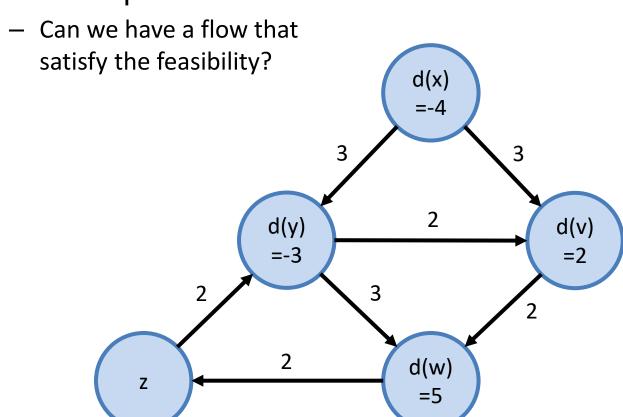


- An example below...
 - Vertex with demands



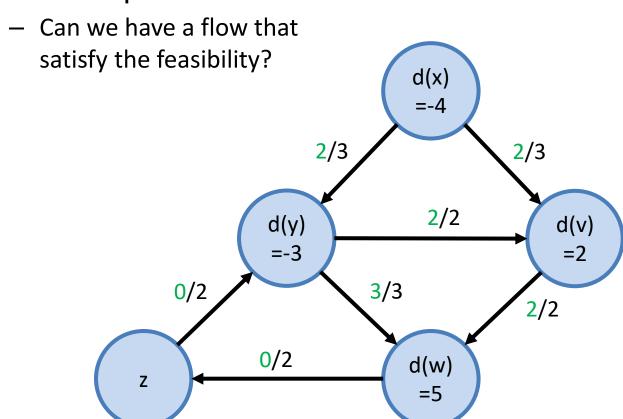
A Feasibility Problem...





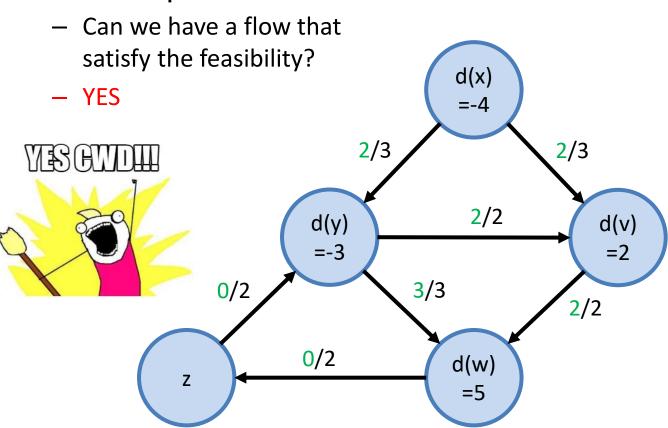
A Feasibility Problem...





A Feasibility Problem...





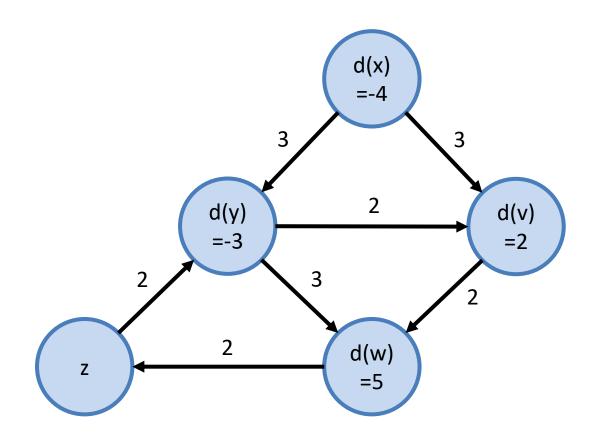


Questions?

How to Check Feasibility...



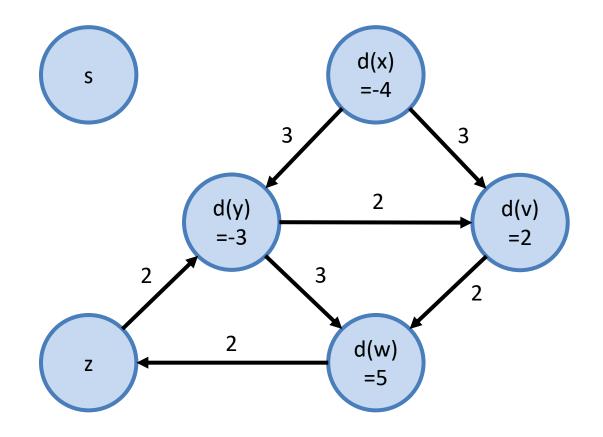
Given this... How do we solve this?



How to Check Feasibility...



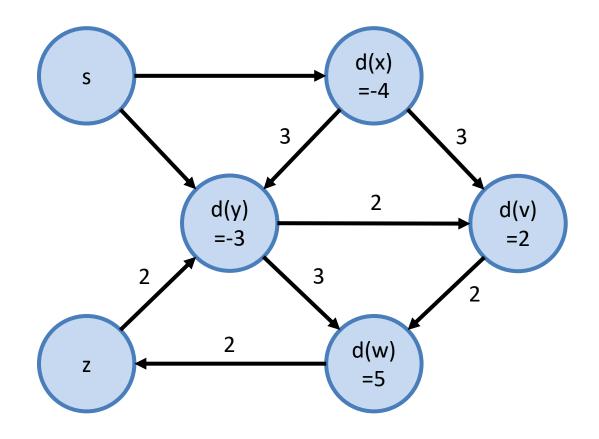
Make a source



How to Check Feasibility...



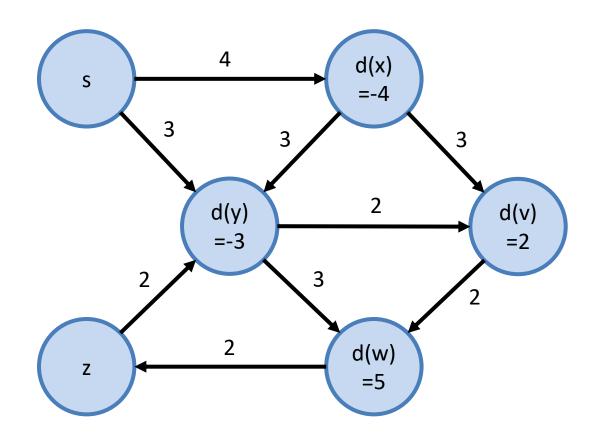
Make a source, link to all negative demand



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How to Check Feasibility...

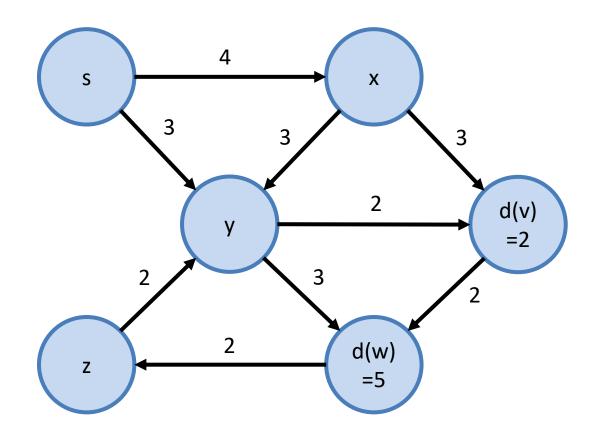
Make a source, link to all negative demand, weighted



How to Check Feasibility...



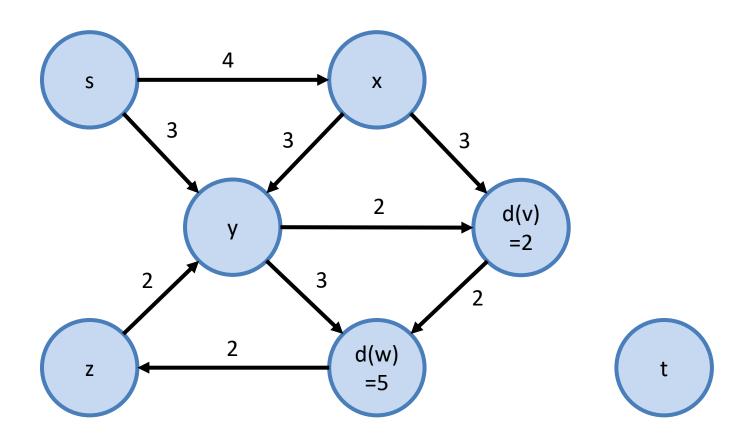
We are done! for source >.<</p>



How to Check Feasibility...



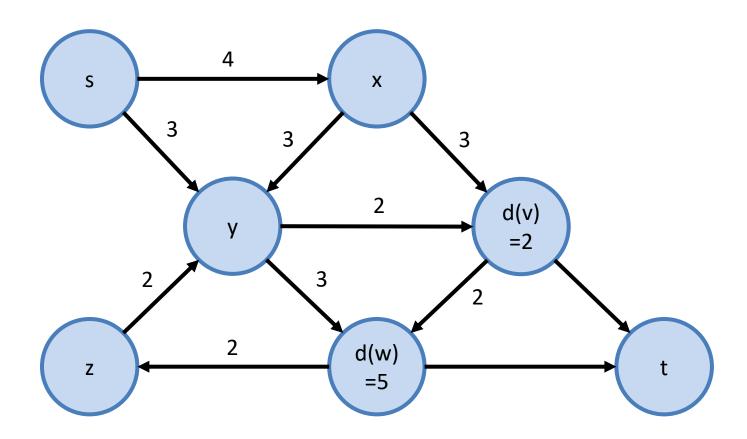
Make a sink



How to Check Feasibility...



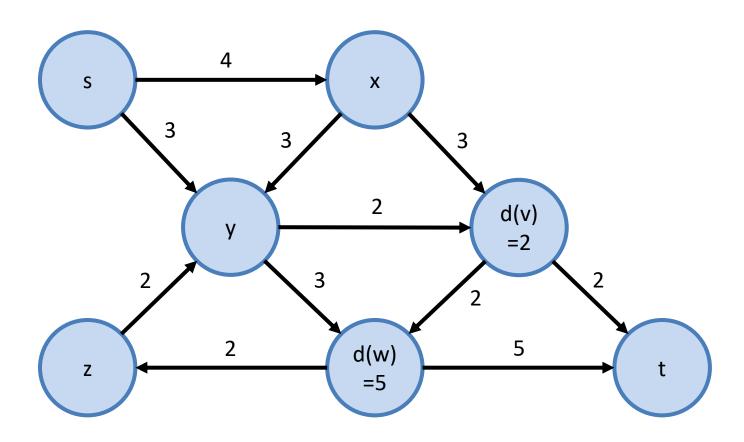
Make a sink, link from positive demand





How to Check Feasibility...

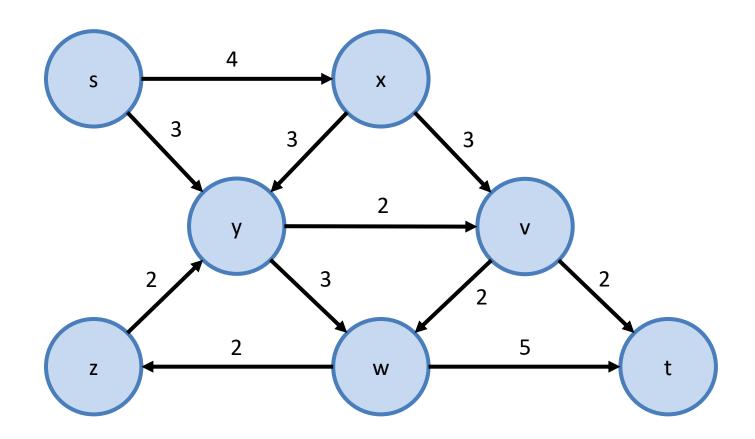
Make a sink, link from positive demand, weighted



How to Check Feasibility...

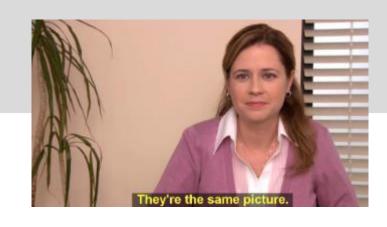


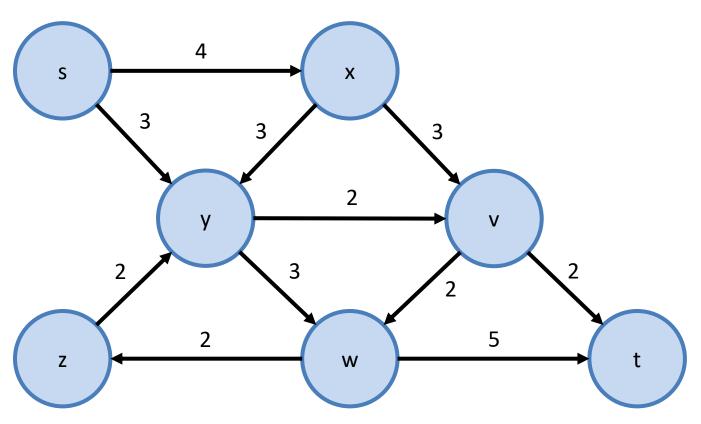
We are done! for sink >.<</p>



How to Check Feasibility...

Now same as network flow!





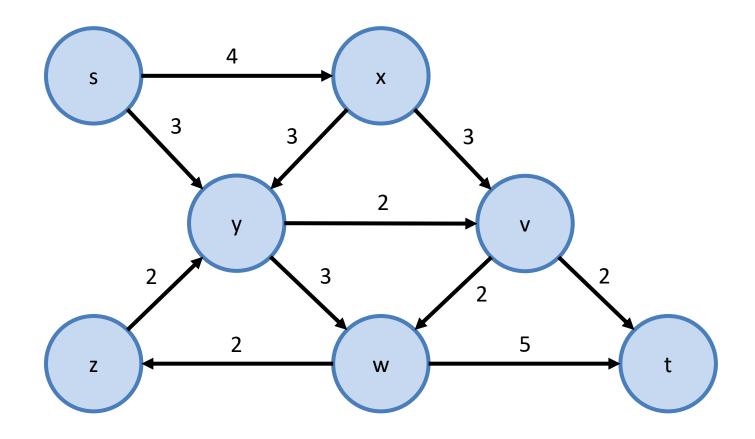


Questions?

How to Check Feasibility...



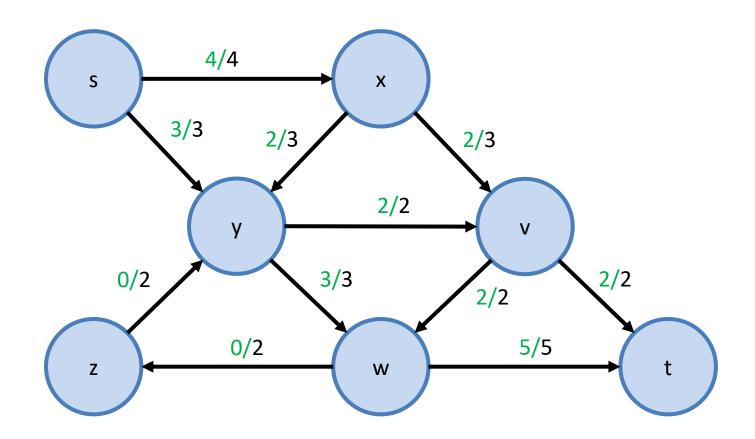
Then we run Ford-Fulkerson



How to Check Feasibility...



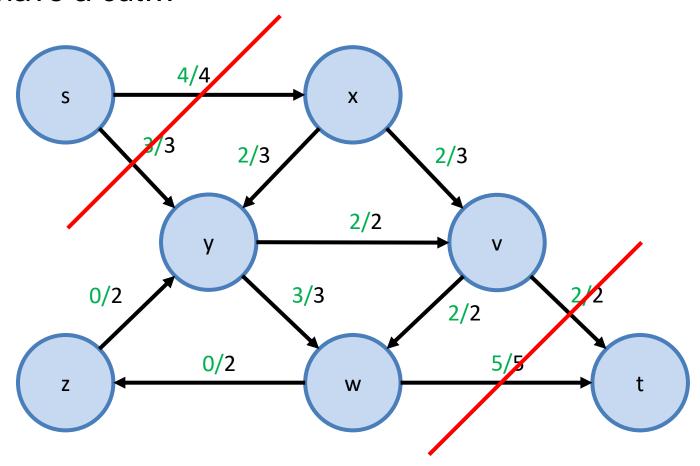
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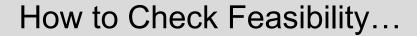


How to Check Feasibility...



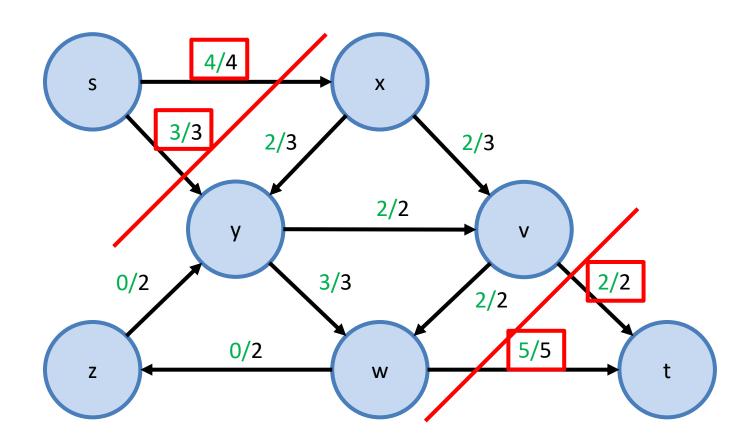
If we have a cut...







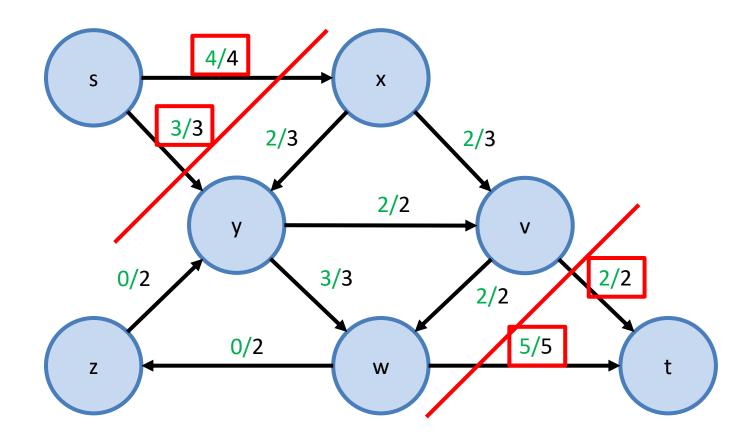
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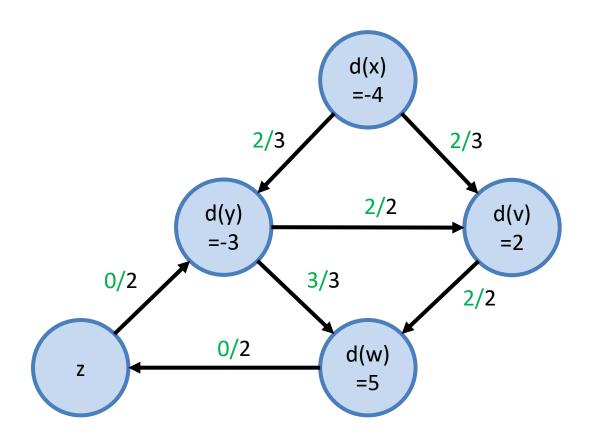
If we have a cut... it is feasible!



How to Check Feasibility...



Then we just clean it up as the solution



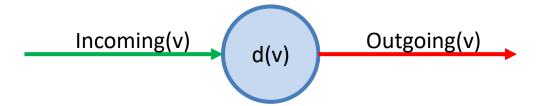


Questions?



A Feasibility Problem...

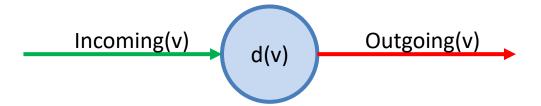
- Recall the 2 concepts from Circulation with Demands
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 - Demand Constraint
 - Now what if we tweak this rule? incoming(v) outgoing(v) = demand(v)



 Circulation with Demands is a feasibility problem that satisfy both of the above!



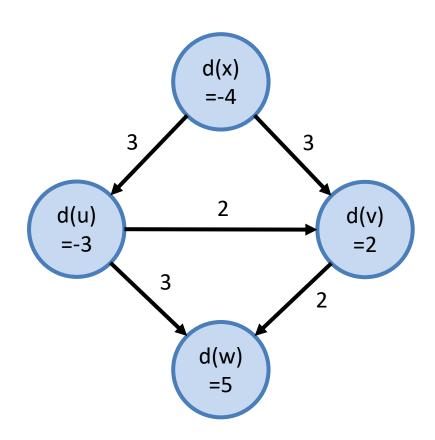
- Recall the 2 concepts from Circulation with Demands
 - Capacity Constraint
 - Lower bound for an edge <= Flow <= Capacity for an edge</p>
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 Circulation with Demands is a feasibility problem that satisfy both of the above!



Consider the following...

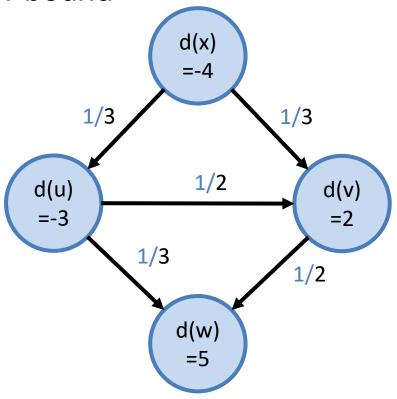




A Feasibility Problem...

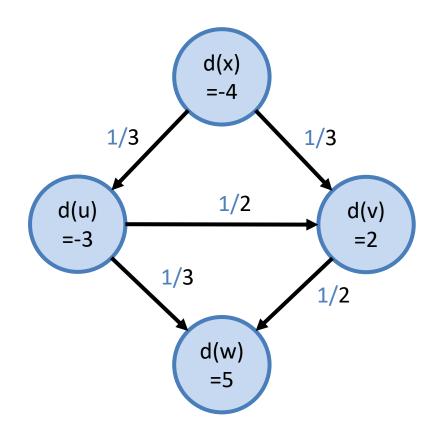
 Consider the following... edges have lower bound

of 1



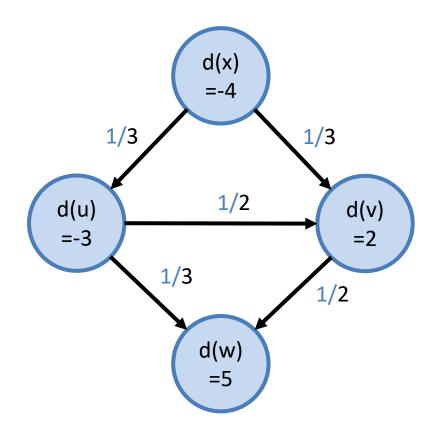


Can we find a feasible solution?





Can we find a feasible solution? Of course!

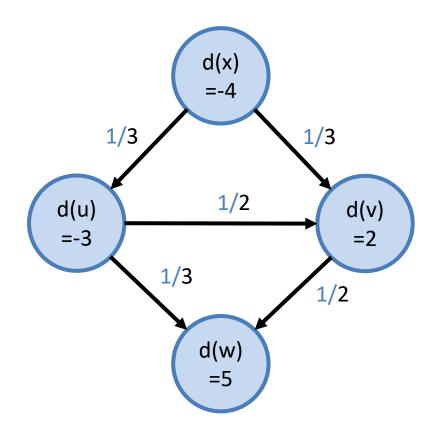




Questions?

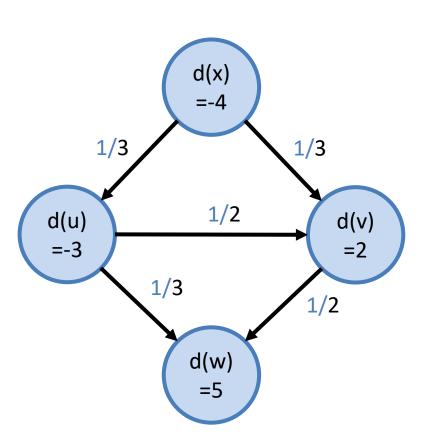


We will need to make some transformation...



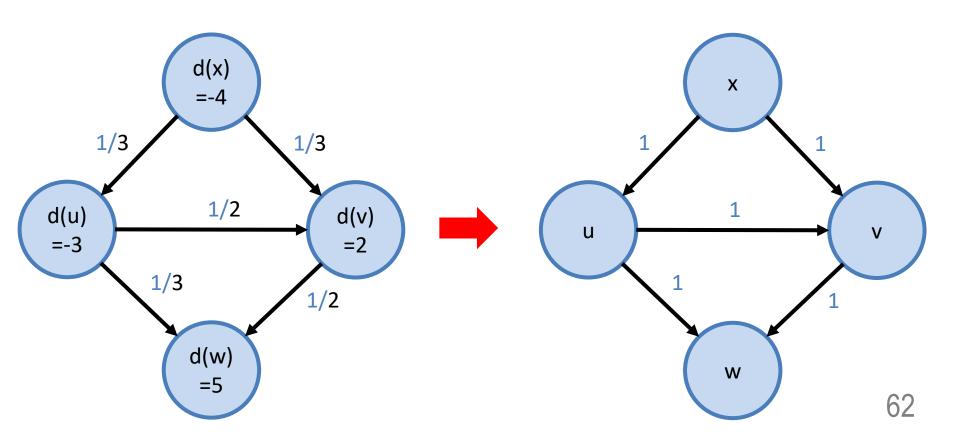


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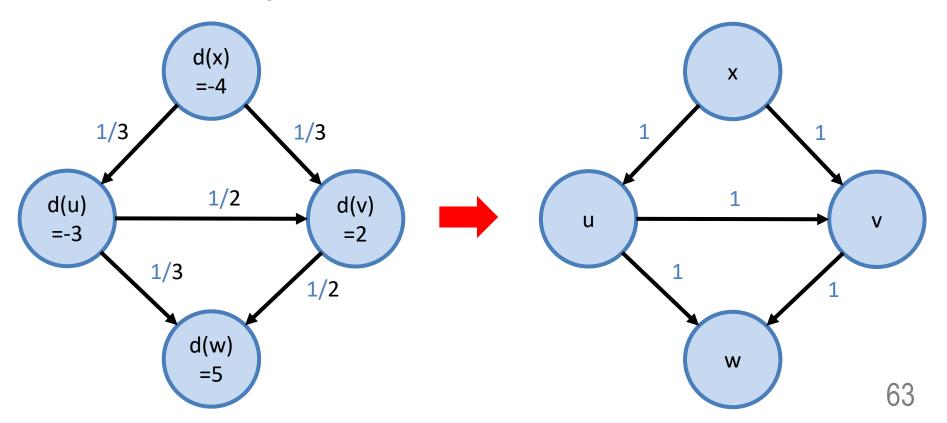


We will need to make some transformation by removing the lower bound to a temp network...



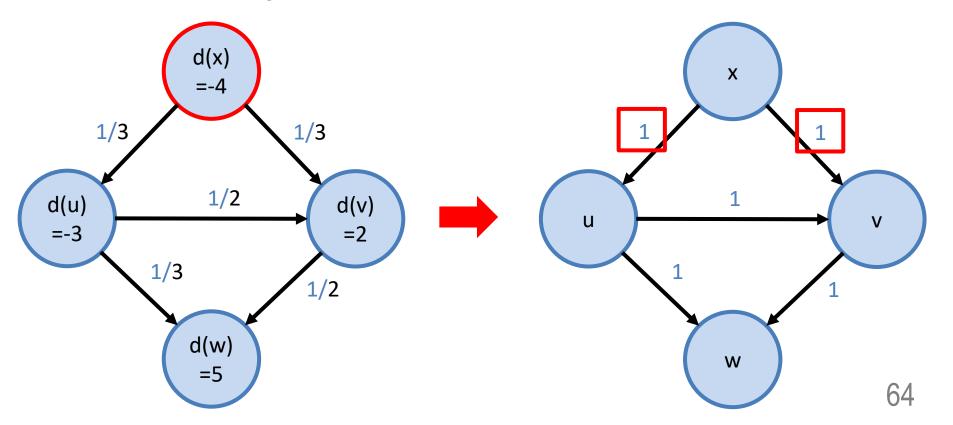


- We will need to make some transformation by removing the lower bound to a temp network...
- Thus, the original reduced...



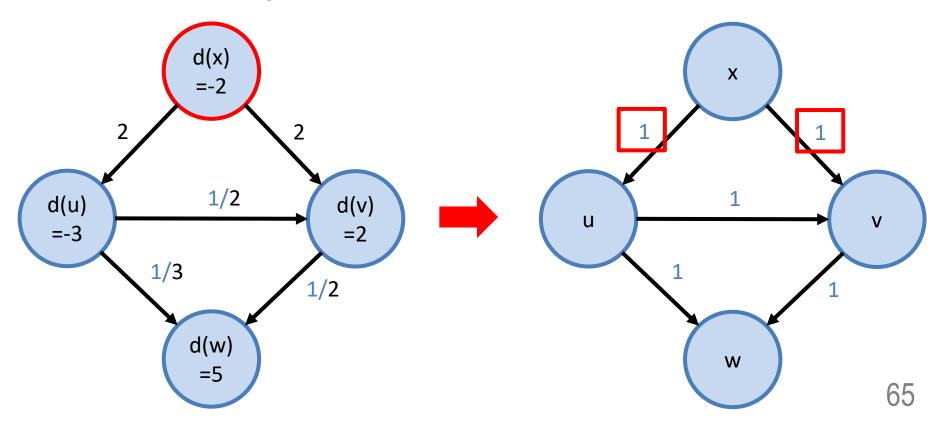


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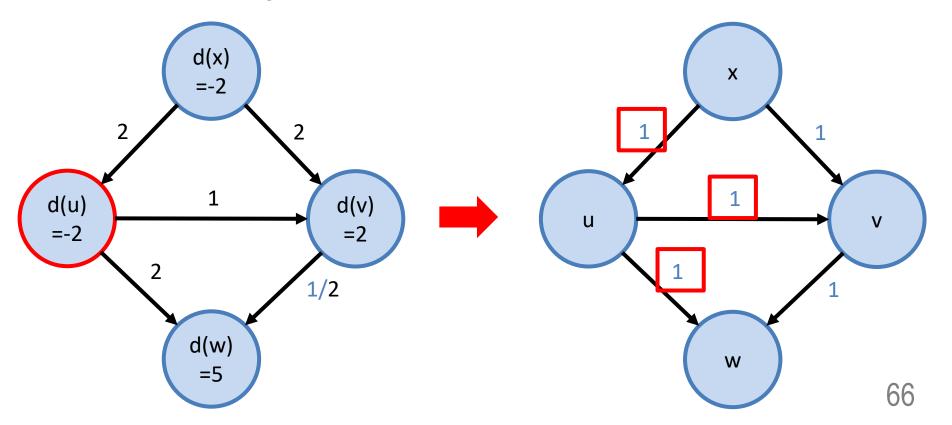


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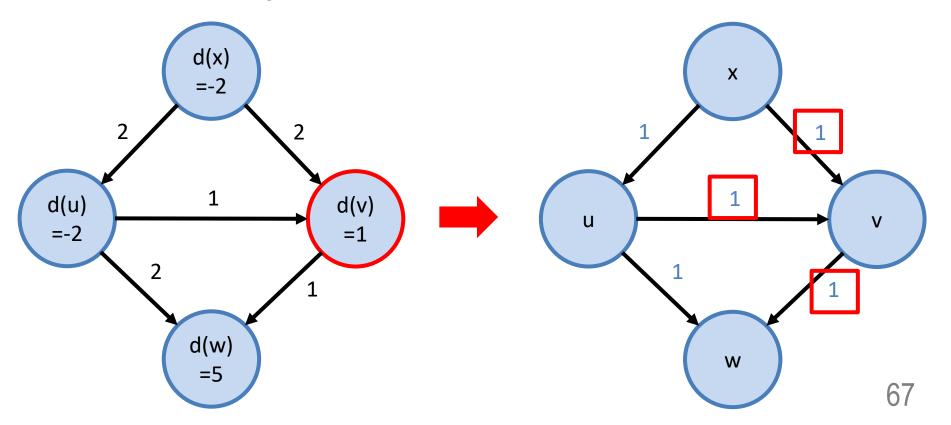


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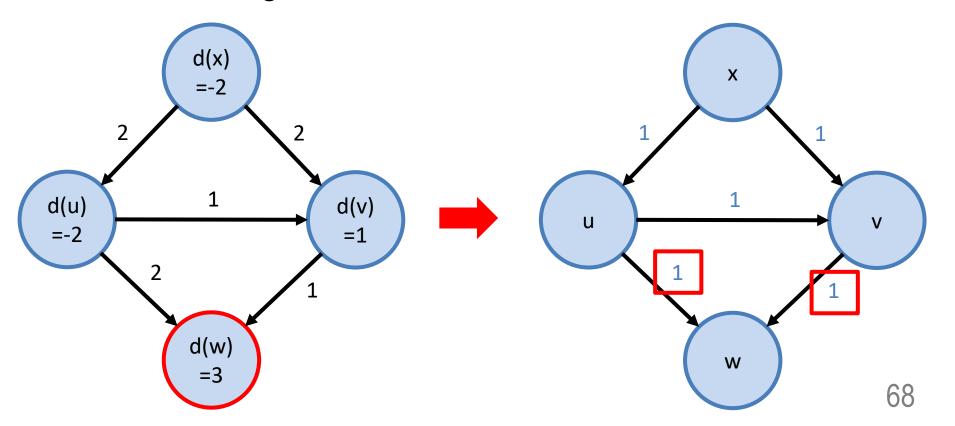


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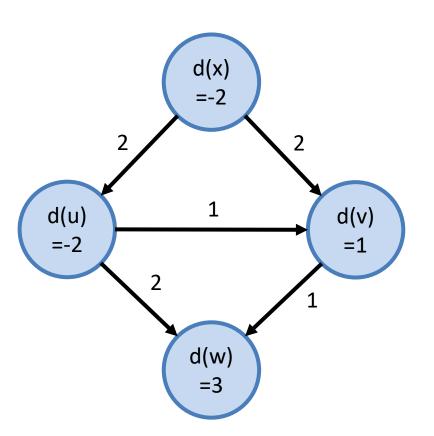
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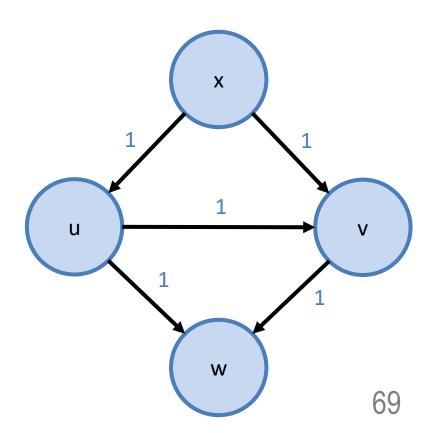




A Feasibility Problem...

Thus, the original reduced...

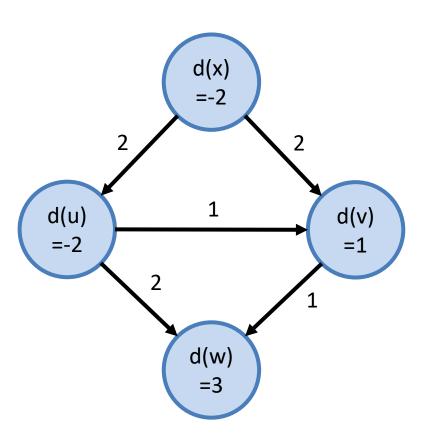


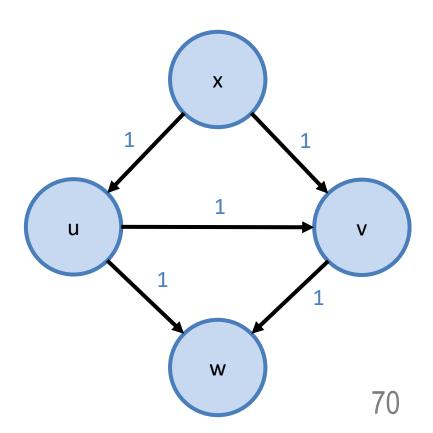




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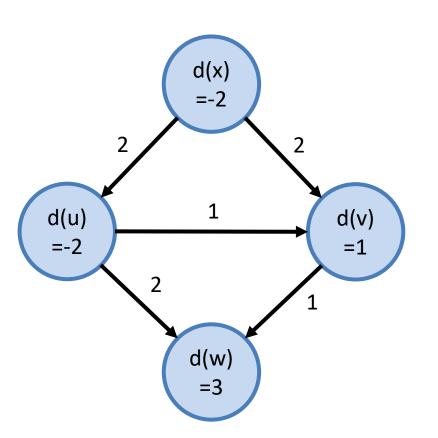


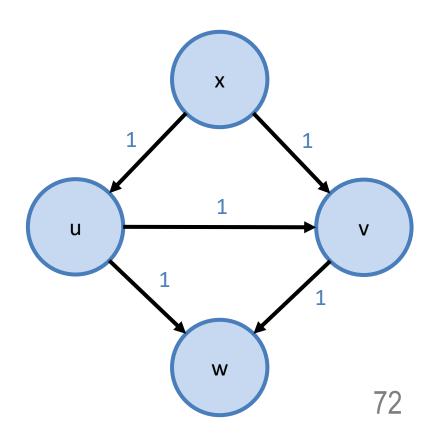
Questions?



A Feasibility Problem...

 Then we follow the same Circulation with Demands as earlier for the reduced network...

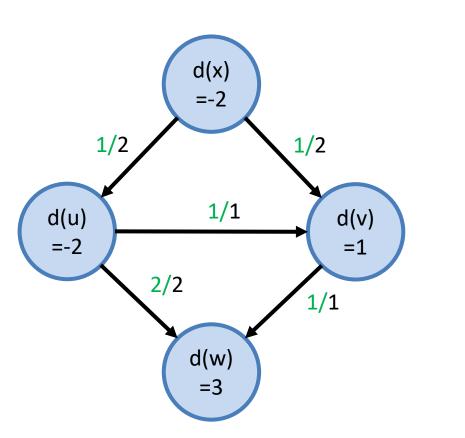


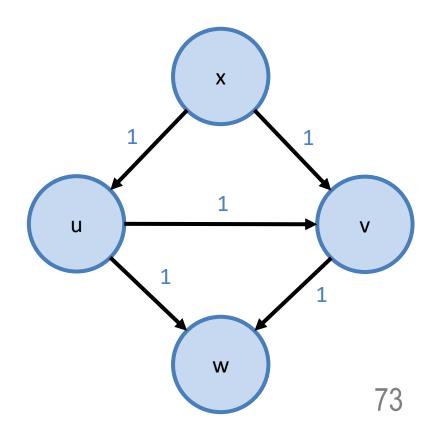




A Feasibility Problem...

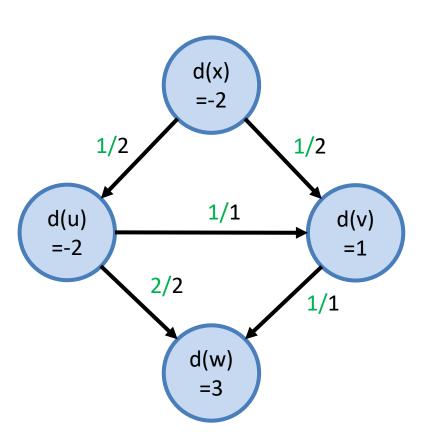
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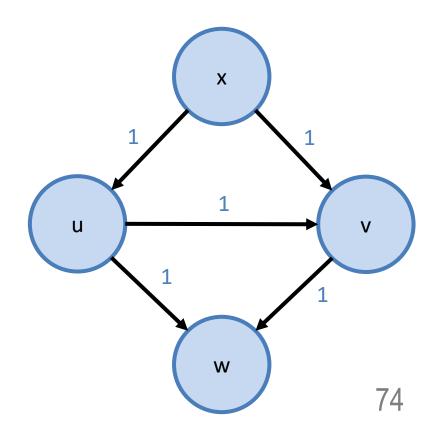




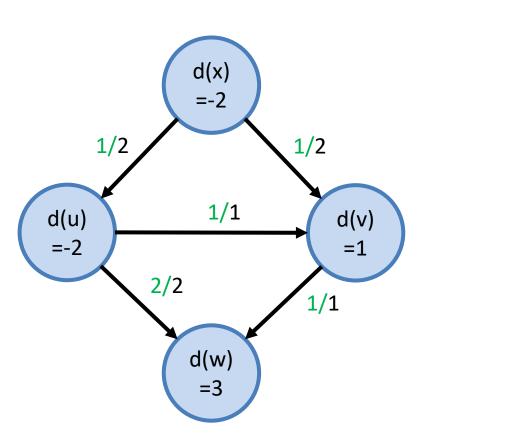


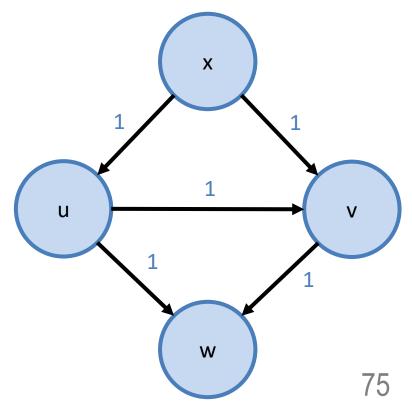
Then we follow the same Circulation with Demands as earlier for the reduced network... It is feasible!



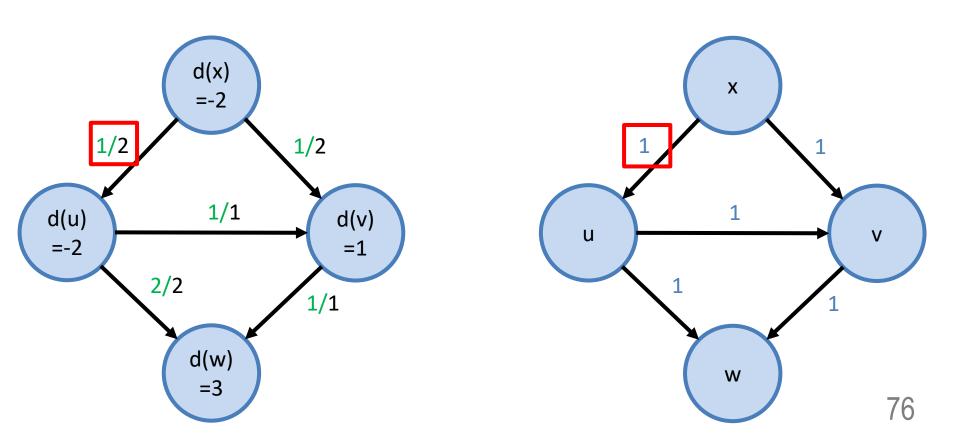




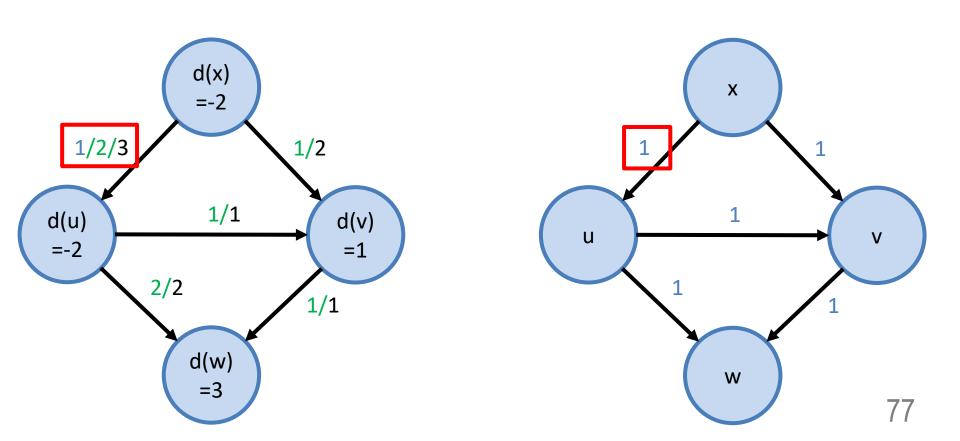




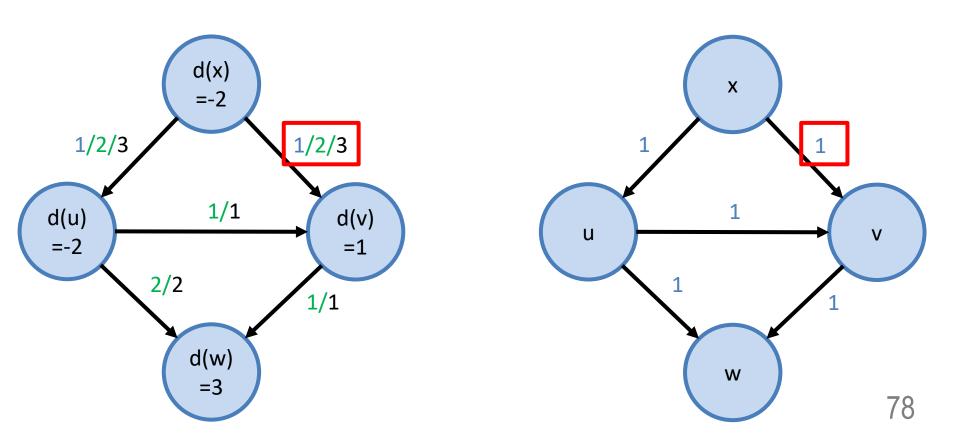




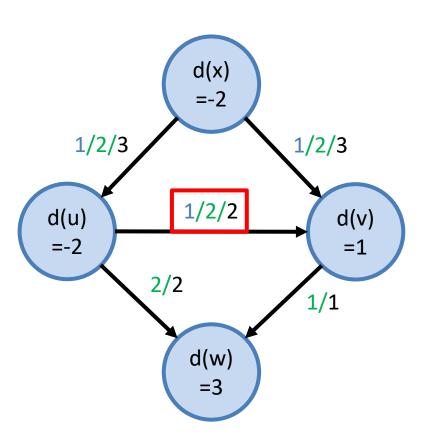


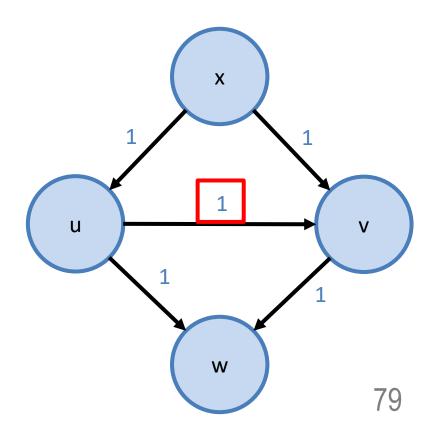




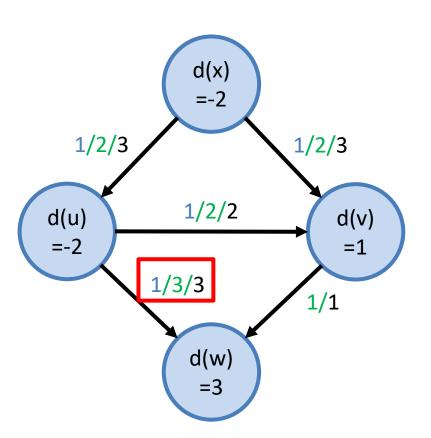


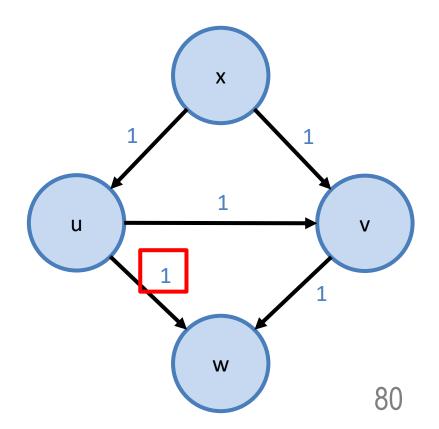




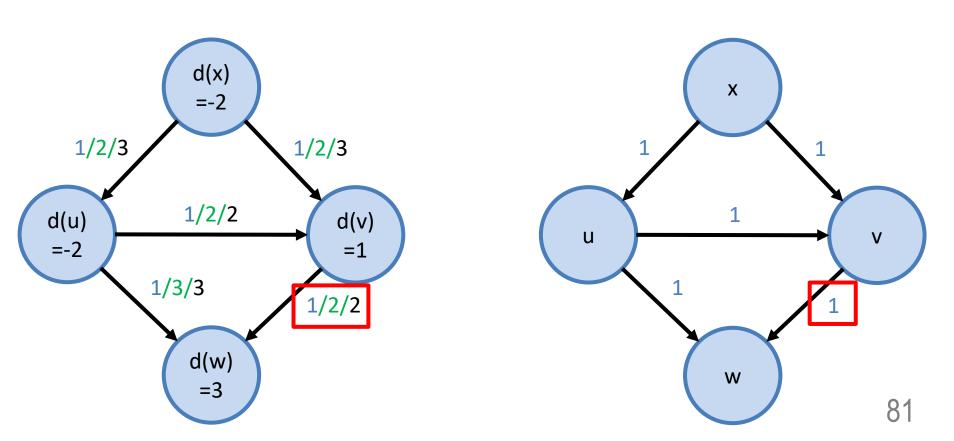






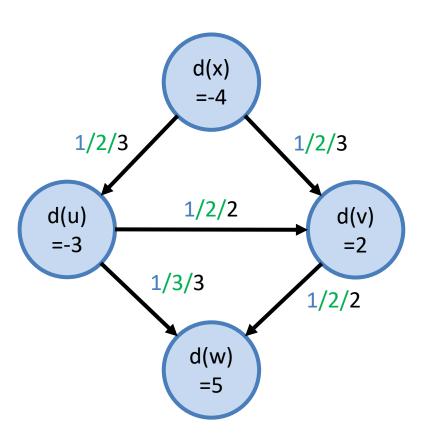


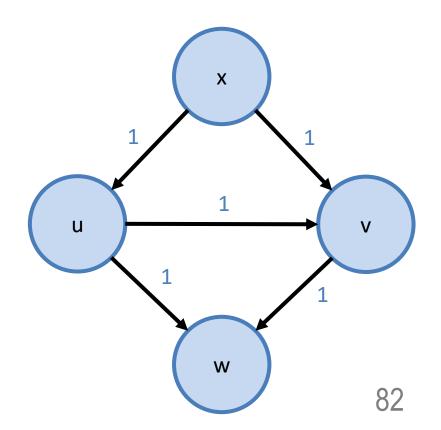






Don't forget the demand of the vertices as well!

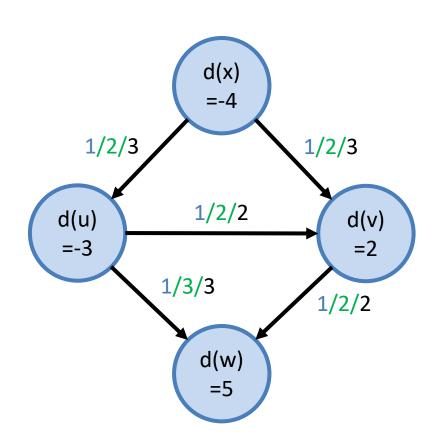






A Feasibility Problem...

And we are done!



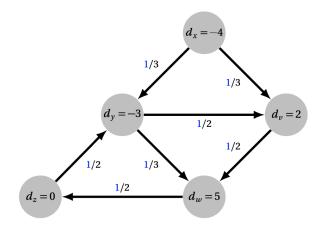


Questions?



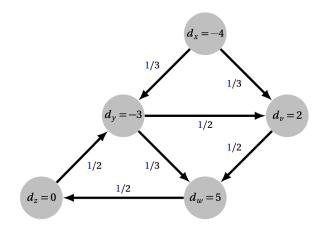
A Feasibility Problem...

 The following is an example from the Clayton campus



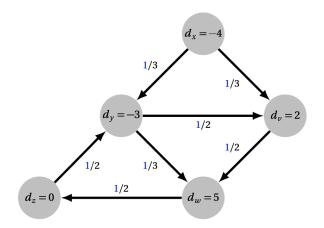


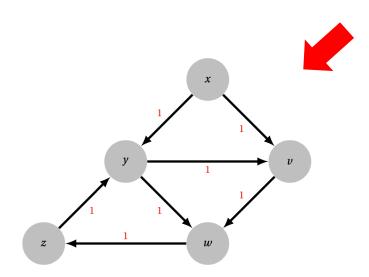
- The following is an example from the Clayton campus
- Work it out on your own to see if it is feasible before we discuss in class





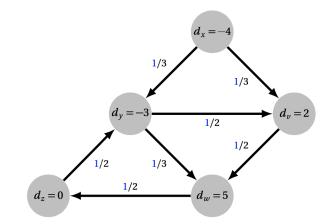
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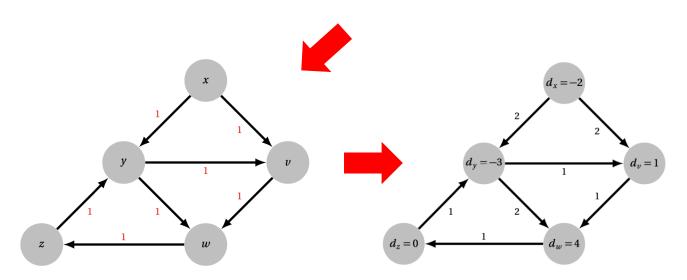






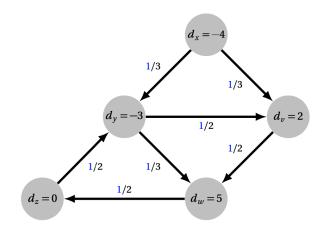
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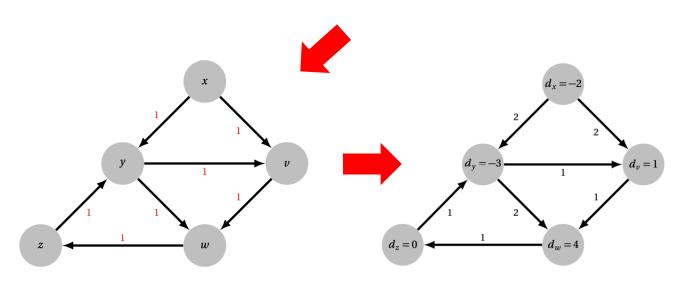






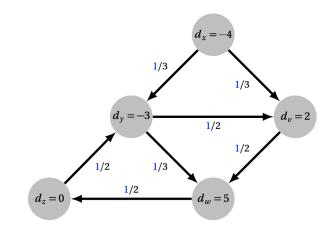
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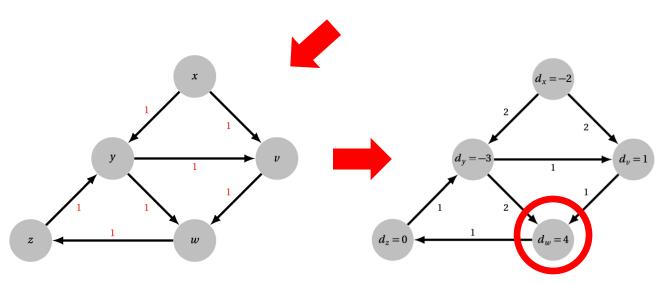






- The following is an example from the Clayton campus
- Work it out on your own to see if it is feasible before we discuss in class. It is it not!







Questions?



- We know what is a flow network.
- We know how to design flow network for bipartite matching.
- We know how to design flow network for circulation with demands and lower bound?



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- We know how to design flow network for circulation with demands and lower bound? We shall see them now...



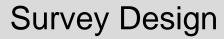
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- We know what is a flow network.
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- We know how to design flow network for circulation with demands and lower bound? We shall see them now...
- Not that we only deal with integers, to make it simpler



Questions?







- You have C customers who have used the product
- You have P products



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- You want to conduct a survey, but...



- You have C customers who have used the product
- You have P products
- You want to conduct a survey, but...
 - You do not want to ask the customer for too many reviews.
 - You do want to ask the customer for at least some reviews.
 - Each product needs to have at least some reviews.
 - Each product do not require more than some reviews.



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 - Of course, each customer can only give a review per product.
- Let us add some notations



- You have C customers who have used the product, $c_1, c_2, c_3, \ldots, c_n$
- You have P products, $p_1, p_2, p_3, ..., p_m$
- You want to conduct a survey, but...
 - You do not want to ask the customer c_i for too many reviews c_i^+ .
 - You do want to ask the customer c_i for at least some reviews c_i .
 - Each product p_j needs to have at least some reviews p_j^- .
 - Each product p_i do not require more than some reviews p_i^+ .
 - Of course, each customer can only give a review per product.
- Let us add some notations

Applications with Network FlowSurvey Design

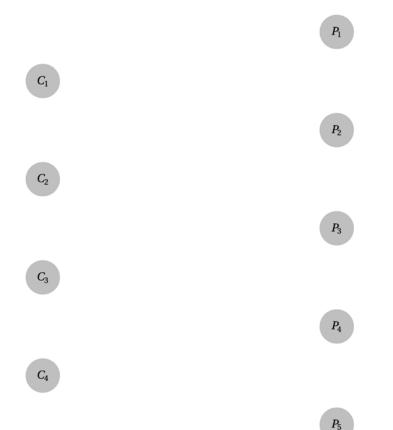


Now let us go through 1 by 1

Applications with Network Flow Survey Design



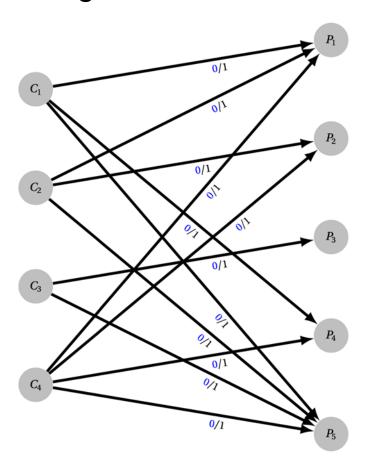
The customer and the product.



Survey Design



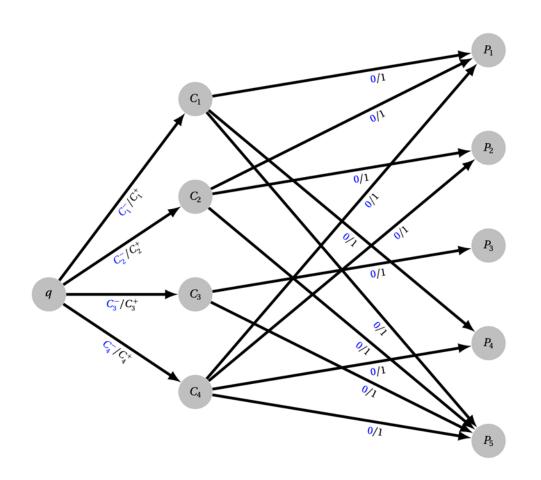
Each customer can give 0 or 1 review for the product



Survey Design



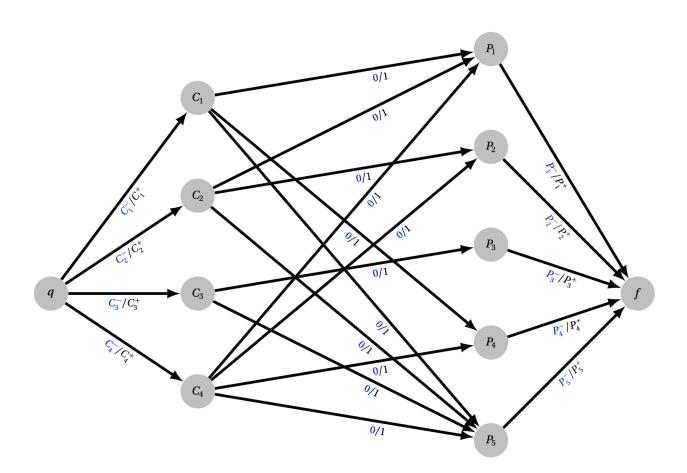
There is range of review expected from the customer



Survey Design



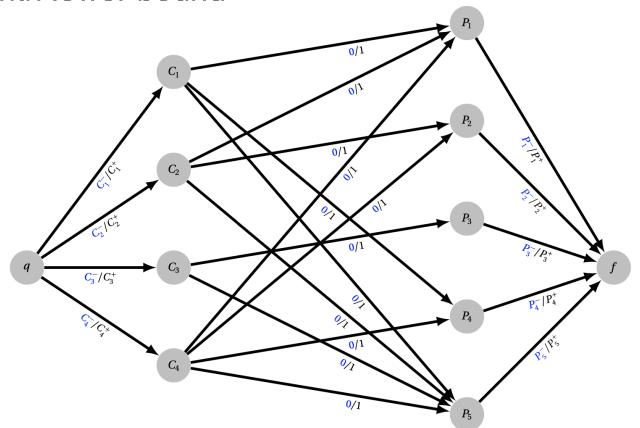
There is range of review expected from of the product



Survey Design



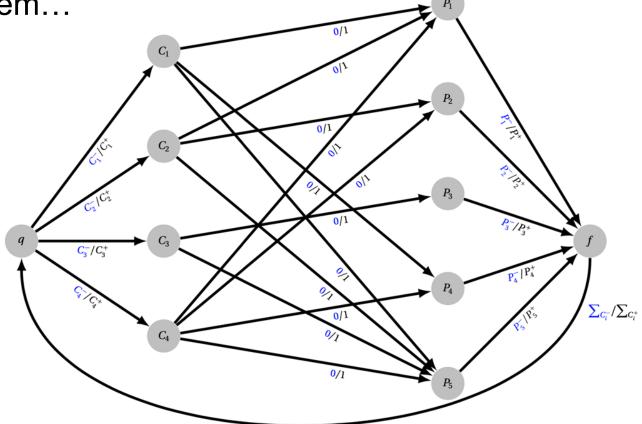
 In a way, it do look like a bipartite matching problem but with lower bound



Survey Design



In a way, it do look like a bipartite matching problem but with lower bound, but since it is a circulation problem...





Questions?



Airline Scheduling



You have a collection of airplanes



- You have a collection of airplanes
- You have a list of routes
- Some of the routes are very profitable, thus you want to fly the routes
 - Departure location
 - Departure time
 - Arrival location
 - Arrival time



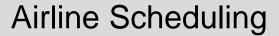
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- You have a collection of airplanes, k
- You have a list of routes, r₁, r₂,..., r_n
- Some of the routes are very profitable, thus you want to fly the routes
 - Departure location
 - Departure time
 - Arrival location
 - Arrival time
- The airplanes can start flying from any location



- Imagine you have the following routes:
 - Route 1: SYD 6am MEL 7am
 - Route 2: CBR 8am SYD 9am
 - Route 3: MEL 11am BNE 1pm
 - Route 4: PER 11am SYD 7pm





- Imagine you have the following routes:
 - Route 1: SYD 6am MEL 7am
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- Can you cover these 4 vital routes, using only 2 planes?

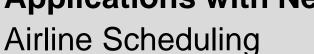


Airline Scheduling

First, we list down the routes imagine the x-axis as time...





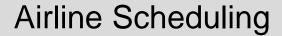




- First, we list down the routes imagine the x-axis as time...
 - Since they are vital flights, we want to always fly and thus lower-bound is set to 1.

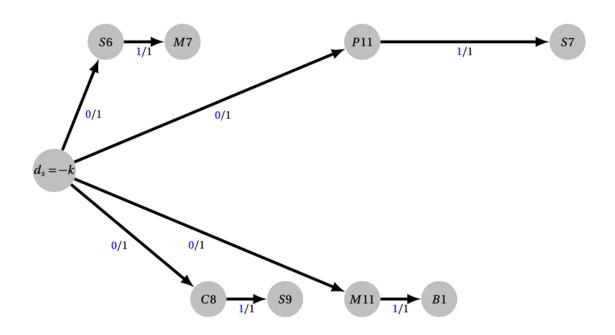


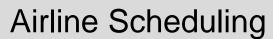






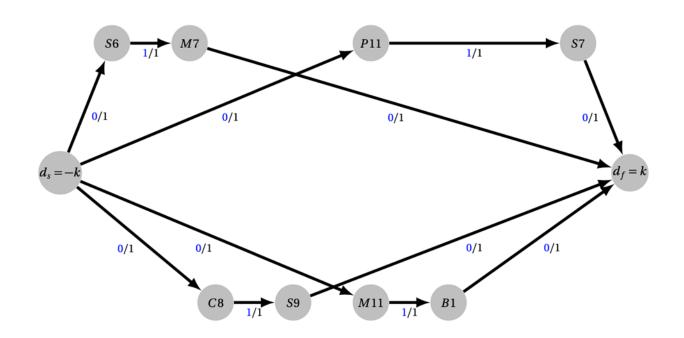
- Then we add a source, which we can place our planes from in any of the locations.
 - Lower bound is 0 because there is not requirement to be placed at which location

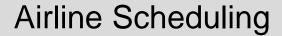






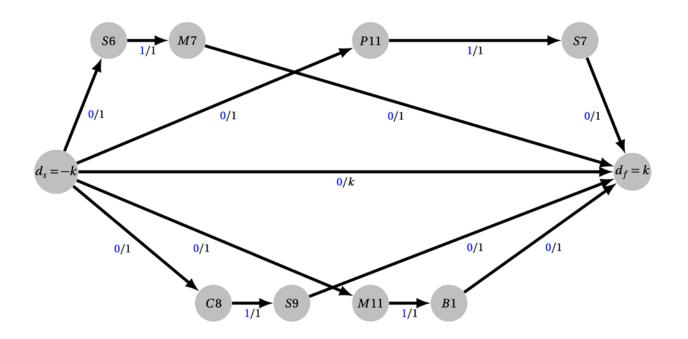
- Next, we add a sink which we can retire our planes at any of the locations.
 - No requirement for the planes to retire from any location

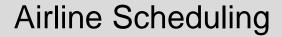






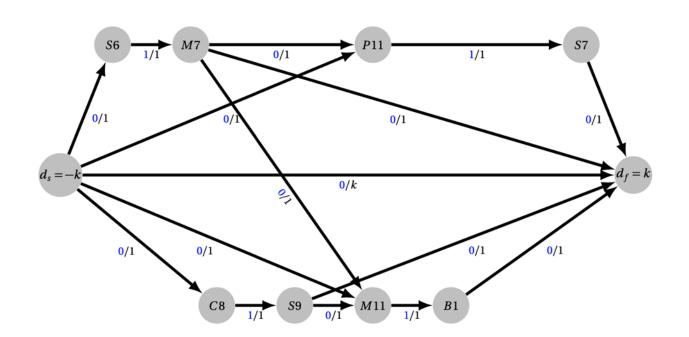
- But what if we don't need all our planes to cover all vital routes?
 - Thus, they can go from start to retire directly







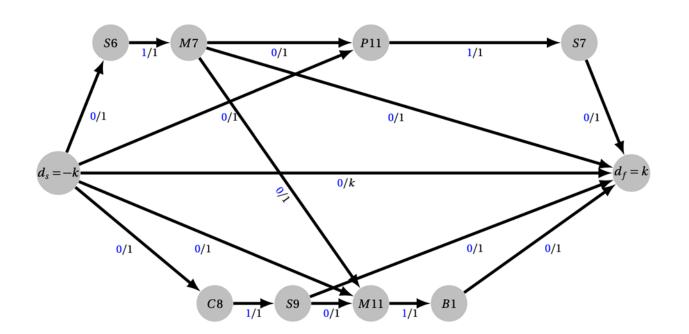
- Since it is possible for a plane to follow a route, then go to another route instead of retiring...
 - We add the edge, again it is optional



Airline Scheduling

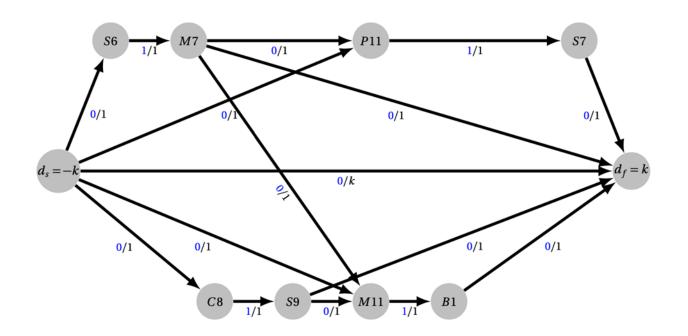


Then we just solve this as it is!





- Then we just solve this as it is!
- Answer is yes with 2 planes
 - Plane1: Route1 (S6->M7), then Route4 (P11->S7)
 - Plane2: Route2 (C8->S9), then Route3 (M11->B1)





Questions?

Other Examples?



Applications with Network Flow Other Examples?



Several examples in the studio

Other Examples?



- Several examples in the studio
 - Choosing profitable projects
 - Determining if teams/ players can progress in a tournament

Other Examples?



Several examples in the studio

- Choosing profitable projects
- Determining if teams/ players can progress in a tournament

... and many more

- Open-pit mining
- Image segmentation (e.g., background/foreground segmentation)
- Network connectivity
- Data mining
- Distributed computing
- Network intrusion detection
- Edge-disjoint paths in graphs
- Network reliability
- Multi-camera scene reconstruction
- Gene function prediction



Questions?



Thank You