

# Lambda Calculus and Combinators

Syntax	Name
$x$	Variable
$(\lambda x.M)$	Abstraction
$(M\ N)$	Application

  

Tape

R/W head

Program

# Housekeeping

- Cameras on for interviews
- Get your code ready before the interview
- Observables due this week (together with this week's tute) - marked next week

# Lambda Calculus

$I = \lambda x . x$       lambda calculus expression       $(\lambda x . x) y$

$I = x \Rightarrow x$       JavaScript       $(x \Rightarrow x) (y)$

# Alpha Equivalence

- expressions are equivalent if their variables are renamed
- $\lambda x y . x z = \lambda a b . a c$

# Beta Reduction

- Application of functions involves substituting the argument into the expression

$$\begin{aligned} &= (\lambda x y . x) (\lambda x . x) \\ &= (\lambda x [x := \lambda x . x] y) \quad \Leftarrow \text{Beta reduction} \\ &= \lambda y . (\lambda x . x) \\ &= \lambda y x . x \quad \Leftarrow \text{Equivalent due to currying} \\ &= \lambda x y . y \quad \Leftarrow \text{Alpha equivalence} \end{aligned}$$

- The same thing in javascript!

```
> const f = (x => y => x)(x => x)
< undefined
> f(2)(3)
< 3
```

# Eta Conversion

- Wrapping a simple lambda around an expression does not change the expression

$$\lambda x . M x = M$$



Eta conversion

# Combinators

Combinators are functions which are expressions of only their parameters

They let us combine and transform other functions in various ways

`const`

`I = x => x`

Identity: more useful than you might think!  
e.g. test map: `a.map(I) == a`

,

`K = x => y => x`

Ignore the second parameter.  
Where have we seen this before?