Activity 3: Outer products, operators, and circuits

Resolution of the identity

In two dimensions

Suppose that $|i\rangle$, $|j\rangle$ form an orthonormal basis for \mathbb{C}^2 . Show that

$$|i\rangle\langle i| + |j\rangle\langle j| = I$$

where I is the identity operator.

Note: this identity generalizes to d dimensions. If $\{|n\rangle\}_{n=0}^{d-1}$ is an orthonormal basis for \mathbb{C}^d , then

$$\sum_{n=0}^{d-1} |n\rangle\langle n| = I$$

Solution

A linear operator is completely defined by its action on any orthonormal basis. $\{|i\rangle, |j\rangle\}$ is orthonormal implies that

$$\langle i|i\rangle = \langle j|j\rangle = 1$$

 $\langle i|j\rangle = 0$

Multiplication of kets is distributive, so we have

$$(|i\rangle\!\langle i|+|j\rangle\!\langle j|)\,|i\rangle=|i\rangle\underbrace{\langle i|i\rangle}_{-1}+|j\rangle\underbrace{\langle j|i\rangle}_{-0}=I\,|i\rangle$$

Similarly,

$$(|i\rangle\!\langle i|+|j\rangle\!\langle j|)\,|j\rangle=|j\rangle\underbrace{\langle i|j\rangle}_{=0}+|j\rangle\underbrace{\langle j|j\rangle}_{=1}=I\,|j\rangle$$

Since they act the same way on an orthonormal basis,

$$(|i\rangle\langle i| + |j\rangle\langle j|) = I$$

Unitary operators

Unitary operators preserve inner products

We defined a unitary operator as a linear operator that conserves total probability, i.e.

$$\left\langle \psi\right|U^{\dagger}U\left|\psi\right\rangle =\|U\left|\psi\right\rangle \|^{2}=\|\left|\psi\right\rangle \|^{2}=1$$

Using the fact that this is true for an arbitrary state $|\psi\rangle$, prove that $U^{\dagger}U=I$, where I is the identity operator.

Hint: Consider the states $\frac{|0\rangle+|1\rangle}{\sqrt{2}}$ and $\frac{|0\rangle+i|1\rangle}{\sqrt{2}}$.

Solution

Let $|\phi\rangle = (|0\rangle + |1\rangle)/\sqrt{2}$. We can check that $||\phi\rangle||^2 = 1$. By the probability-conserving property,

$$\langle \phi | U^{\dagger} U | \phi \rangle = 1$$

We can also expand this in terms of $|0\rangle$, $|1\rangle$:

$$\langle \phi | U^{\dagger} U | \phi \rangle = \left(\frac{\langle 0 | + \langle 1 |}{\sqrt{2}} \right) U^{\dagger} U \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}} \right) \tag{1}$$

$$=\frac{\langle 0|U^{\dagger}U|0\rangle + \langle 1|U^{\dagger}U|0\rangle + \langle 0|U^{\dagger}U|1\rangle + \langle 0|U^{\dagger}U|0\rangle}{2}$$
(2)

By the probability-conserving property, we have $\langle 0|U^{\dagger}U|0\rangle = 1$ and $\langle 1|U^{\dagger}U|1\rangle = 1$. As a general property of the adjoint operation, we also have

$$\langle 1 | U^{\dagger} U | 0 \rangle = (\langle 0 | U^{\dagger} U | 1 \rangle)^*$$

Where * denotes the complex conjugate. This gives

$$\frac{\langle 0|U^{\dagger}U|0\rangle + \langle 1|U^{\dagger}U|0\rangle + \langle 0|U^{\dagger}U|1\rangle + \langle 1|U^{\dagger}U|1\rangle}{2} = \frac{2 + 2\operatorname{Re}(\langle 0|U^{\dagger}U|1\rangle)}{2} = 1 \tag{3}$$

This top equation implies that $\operatorname{Re}(\langle 0|U^{\dagger}U|1\rangle) = 0$. We can repeat the same process with $|\psi\rangle = (|0\rangle + i|1\rangle)/\sqrt{2}$. We get

$$\langle \psi | U^{\dagger} U | \psi \rangle = \left(\frac{\langle 0 | -i \langle 1 | \rangle}{\sqrt{2}} \right) U^{\dagger} U \left(\frac{|0\rangle + i |1\rangle}{\sqrt{2}} \right) \tag{4}$$

$$=\frac{\left\langle 0\right|U^{\dagger}U\left|0\right\rangle -i\left\langle 1\right|U^{\dagger}U\left|0\right\rangle +i\left\langle 0\right|U^{\dagger}U\left|1\right\rangle +\left\langle 1\right|U^{\dagger}U\left|1\right\rangle }{2}\tag{5}$$

$$=\frac{2+2i\operatorname{Im}(\langle 0|U^{\dagger}U|1\rangle)}{2}=1\tag{6}$$

The last equation implies that $\operatorname{Im}(\langle 0|U^\dagger U|1\rangle)=0$. Together, we have $\langle 0|U^\dagger U|1\rangle=0$ and $\langle 1|U^\dagger U|0\rangle=(\langle 0|U^\dagger U|1\rangle)^*=0$. This implies that $U^\dagger U|0\rangle=\alpha|0\rangle$ and $U^\dagger U|1\rangle=\beta|1\rangle$, where α and β are constants. Since $\langle 0|U^\dagger U|0\rangle=\alpha=1$ and $\langle 1|U^\dagger U|1\rangle=\beta=1$, we have $U^\dagger U|0\rangle=|0\rangle$ and $U^\dagger U|1\rangle=|1\rangle$. This shows that $U^\dagger U=I$.

Pauli-Euler identity

Consider again the Pauli-X operator, defined by $X|0\rangle = |1\rangle$ and $X|1\rangle = |0\rangle$. Recall from last time that the X-basis is defined as

$$|+\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}}$$
 $|-\rangle = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$ (7)

Since $X \mid + \rangle = \mid + \rangle$ and $X \mid - \rangle = - \mid - \rangle$, we can consistently define $e^{i\theta X} \mid + \rangle = e^{i\theta} \mid + \rangle$ and $e^{i\theta X} \mid - \rangle = e^{-i\theta} \mid - \rangle$. Using this, show that $e^{i\theta X} = I \cos(\theta) + iX \sin(\theta)$

where I is the identity operator.

Note: This holds for any operator that is self-inverse because these operators have eigenvalues ± 1 , including Y, Z, and $n_x X + n_y Y + n_z Z$ where $||(n_x, n_y, n_z)|| = 1$.

Solution

We can specify a linear operator by its action on an orthonormal basis, and $|+\rangle$, $|-\rangle$ forms such a basis. Using the Euler formula, we have

$$(I\cos(\theta) + iX\sin(\theta)) |+\rangle = (\cos(\theta) + i\sin(\theta)) |+\rangle = e^{i\theta} |+\rangle$$

and

$$\left(I\cos(\theta)+iX\sin(\theta)\right)\left|-\right\rangle = \left(\cos(\theta)-i\sin(\theta)\right)\left|-\right\rangle = e^{-i\theta}\left|-\right\rangle$$

This shows that $e^{i\theta X} = I\cos(\theta) + iX\sin(\theta)$.

The trace

The trace is a useful operation in quantum mechanics. The trace of an operator \hat{O} is defined

$$Tr\{\mathcal{O}\} = \langle 0 | \mathcal{O} | 0 \rangle + \langle 1 | \mathcal{O} | 1 \rangle$$

If we write \mathcal{O} using outer products,

$$\mathcal{O} = \mathcal{O}_{00} |0\rangle\langle 0| + \mathcal{O}_{10} |1\rangle\langle 0| + \mathcal{O}_{01} |0\rangle\langle 1| + \mathcal{O}_{11} |1\rangle\langle 1|$$

we can see that $\text{Tr}\{\mathcal{O}\} = \mathcal{O}_{00} + \mathcal{O}_{11}$, which is just the standard definition of the trace as the sum of diagonal elements. In d dimensions, if $\{|n\rangle\}_{n=0}^d$ is an orthonormal basis for \mathbb{C}^d , then the trace can be similarly defined as

$$\operatorname{Tr}\{\mathcal{O}\} = \sum_{n} \langle n | \mathcal{O} | n \rangle$$

Cyclic property of the trace

Let A and B both be operators on a d-dimensional complex vector space. Using the resolution of the identity, show that

$$Tr{AB} = Tr{BA}$$

Solution

Note that $\langle n|A|m\rangle$ is just a scalar, so it commutes with everything. Then using the defintion of the trace and resolution of the identity, we have

$$\operatorname{Tr}\{AB\} = \sum_{n=0}^{d-1} \langle n | AB | n \rangle \tag{8}$$

$$= \sum_{n=0}^{d-1} \langle n | A \sum_{m=0}^{d-1} | m \rangle \langle m | B | n \rangle$$

$$(9)$$

$$= \sum_{m=0}^{d-1} \sum_{n=0}^{d-1} \langle n | A | m \rangle \langle m | B | n \rangle$$

$$(10)$$

$$= \sum_{m=0}^{d-1} \sum_{n=0}^{d-1} \langle m | B | n \rangle \langle n | A | n \rangle \tag{11}$$

$$= \sum_{m=0}^{d-1} \langle m | B \sum_{n=0}^{d-1} | n \rangle \langle n | A | m \rangle$$
 (12)

$$= \sum_{m=0}^{d-1} \langle m | BA | m \rangle = \text{Tr}\{BA\} \tag{13}$$

Unitary invariance of trace

If $\{|i\rangle\}_{i=0}^{d-1}$ and $\{|j\rangle\}_{j=0}^{d-1}$ are two different orthonormal bases, then the operator U such that $U|i\rangle = |j\rangle$ is unitary. Use this to show that the trace is invariant with respect to the choice of basis:

$$\operatorname{Tr}\{\mathcal{O}\} = \sum_{i} \langle i | \mathcal{O} | i \rangle = \sum_{j} \langle j | \mathcal{O} | j \rangle$$

Note: This is a slight abuse of notation because the labels i and j are both denoting indices and distinguishing between the two bases.

Solution

Let $\text{Tr}\{O\} = \sum_i \langle i | \mathcal{O} | i \rangle$, and take $|j\rangle = U | i \rangle$. Since U is unitary, we have $U^{\dagger}U = UU^{\dagger} = I$. By the cyclic property of the trace,

$$\sum_{i} \langle j | \mathcal{O} | j \rangle = \sum_{i} \langle i | U^{\dagger} \mathcal{O} U | i \rangle = \text{Tr} \{ U^{\dagger} \mathcal{O} U \} = \text{Tr} \{ \mathcal{O} U U^{\dagger} \} = \text{Tr} \{ \mathcal{O} \}$$

Hello Quantum World

The Hadamard gate

The Hadamard gate is used to create superpositions, and it is defined as

$$H\left|0\right\rangle = \left|+\right\rangle \tag{14}$$

$$H|1\rangle = |-\rangle \tag{15}$$

Make the following circuit in Qiskit: $|0\rangle$ —H— \subset . Recall that the definition of the Z operator is $Z|0\rangle = |0\rangle$ and $Z|1\rangle = -|1\rangle$, and the expectation value is defined as $\langle Z\rangle = \langle \psi | Z | \psi \rangle$. Can you predict the expectation value $\langle Z\rangle$? Confirm your prediction by simulating in Qiskit.

Solution

For a generic state $\psi = \alpha |0\rangle + \beta |1\rangle$, the expectation value of Z is

$$\begin{split} \langle Z \rangle &\equiv \langle \psi | \, Z \, | \psi \rangle = (\alpha^* \, \langle 0 | + \beta^* \, \langle 1 |) Z (\alpha \, | 0 \rangle + \beta \, | 1 \rangle) \\ &= |\alpha|^2 \underbrace{\langle 0 | \, Z \, | 0 \rangle}_{=1} + \alpha^* \beta \underbrace{\langle 0 | \, Z \, | 1 \rangle}_{=0} + \alpha \beta^* \underbrace{\langle 1 | \, Z \, | 0 \rangle}_{=0} + |\beta|^2 \underbrace{\langle 1 | \, Z \, | 1 \rangle}_{=-1} \\ &= |\alpha|^2 - |\beta|^2 \end{split}$$

By applying a Hadamard gate to the $|0\rangle$ state, we are left in the state $|+\rangle$, which has $\alpha = \beta = \frac{1}{\sqrt{2}}$. The expectation value of Z is then

$$|\alpha|^2 - |\beta|^2 = \frac{1}{2} - \frac{1}{2} = 0$$

We also recognize that $|\alpha|^2$ is the probability given by the Born rule of measuring $|\psi\rangle$ in the state $|0\rangle$ and $|\beta|^2$ is the probability of measuring $|\psi\rangle$ in the state $|1\rangle$. In order to compute $\langle Z\rangle$ experimentally, we can run the experiment many times and average the number of 0's and 1's to estimate $|\alpha|^2$ and $|\beta|^2$. The code in Qiskit is the following:

```
from qiskit import QuantumCircuit, execute
from qiskit.providers.aer import AerSimulator
sim = AerSimulator()
qc = QuantumCircuit(1,1)
qc.h(0)
qc.measure(0,0)
shots = 10000
counts = execute(qc, sim, shots = shots).result().get_counts()
expec_Z = (counts.get("0",0)-counts.get("1",0))/shots
```

Rotation gates

Create the following circuit in Qiskit: $|0\rangle - R_x(\theta)$, with the definition

$$R_x(\theta) \equiv e^{-iX\theta/2}$$

What will be the expectation value of Z as a function of θ ? Make a prediction and confirm it by measuring $\langle Z \rangle$ for values of θ between 0 and 2π .

Solution

We have the Pauli-Euler identity, which gives

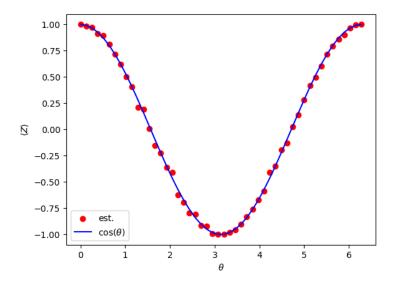
$$e^{-iX\theta/2}\left|0\right\rangle = \cos(\theta/2)\left|0\right\rangle - i\sin(\theta/2)\left|1\right\rangle$$

Using the formula $\langle Z \rangle = |\alpha|^2 - |\beta|^2$, we get

$$\langle Z \rangle = \cos^2(\theta/2) - \sin^2(\theta/2) = \cos(\theta)$$

The code we can use to test this in Qiskit is the following:

```
import numpy as np
def get_expec(theta):
    qc = QuantumCircuit(1,1)
    qc.rx(theta, 0)
    qc.measure(0,0)
    shots = 1000
    counts = execute(qc, sim, shots = shots).result().get_counts()
    return (counts.get("0",0)-counts.get("1",0))/shots
thetalist = np.linspace(0, 2*np.pi, 50)
expects = list(map(get_expec, thetalist))
```



The Bell Basis

We define the Bell basis as

$$\begin{split} |\Phi_1\rangle &= \frac{|00\rangle + |11\rangle}{\sqrt{2}} \\ |\Phi_3\rangle &= \frac{|01\rangle + |10\rangle}{\sqrt{2}} \\ |\Phi_4\rangle &= \frac{|01\rangle - |11\rangle}{\sqrt{2}} \end{split}$$

Come up with a circuit to prepare $|\Phi_1\rangle$. Using the definitions

$$X = |1\rangle\langle 0| + |0\rangle\langle 1| \qquad \qquad Y = i\,|1\rangle\langle 0| - i\,|0\rangle\langle 1| \qquad \qquad Z = |0\rangle\langle 0| - |1\rangle\langle 1|$$

- (a) Show that $X_1I_2|\Phi_1\rangle = |\Phi_3\rangle$, $Y_1I_2|\Phi_1\rangle = -i|\Phi_4\rangle$, and $Z_1I_2|\Phi_1\rangle = |\Phi_2\rangle$.
- (b) Create a circuit that maps $|00\rangle$ to $|\Phi_1\rangle$. Now invert this circuit. What does the inverse do to $|\Phi_2\rangle, |\Phi_3\rangle$, and $|\Phi_4\rangle$?

(c) Confirm this prediction in Qiskit.

Property (a) reflects a symmetry of $|\Phi_1\rangle$ which makes it a useful resource for superdense coding and quantum teleportation, due to the fact that local (single-qubit) operators are shown to produce nonlocal (multi-qubit) changes to the state.

Part a

Using the definitions of X, Y, and Z,

$$X_1 I_2 |\Phi_1\rangle = \frac{X_1 I_2 |00\rangle + X_1 I_2 |11\rangle}{\sqrt{2}}$$

= $\frac{|10\rangle + |01\rangle}{\sqrt{2}} = |\Phi_3\rangle$

$$Y_1 I_2 |\Phi_1\rangle = \frac{Y_1 I_2 |00\rangle + Y_1 I_2 |11\rangle}{\sqrt{2}}$$

= $\frac{i |10\rangle - i |01\rangle}{\sqrt{2}} = -i |\Phi_4\rangle$

$$Z_1 I_2 |\Phi_1\rangle = \frac{Z_1 I_2 |00\rangle + Z_1 I_2 |11\rangle}{\sqrt{2}}$$
$$= \frac{|00\rangle - |11\rangle}{\sqrt{2}} = |\Phi_2\rangle$$

Part b

The circuit we want performs the following operation:

$$|00\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)|0\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |10\rangle) \mapsto \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

The first operation is a Hadamard gate on the first qubit, and the second operation is a CNOT gate controlled on the first qubit and targeting the second qubit. As a circuit,

$$\begin{array}{c|c} |0\rangle \hline H \\ |0\rangle \hline \end{array}$$

Since $H^{-1} = H$ and $CNOT^{-1} = CNOT$, inverting the circuit is as simple as switching the order of the gates:

We can then work out the action of this inverse on the three remaining bell states:

$$CNOT |\Phi_2\rangle = \frac{|00\rangle - |10\rangle}{\sqrt{2}} = \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) |0\rangle = |-\rangle |0\rangle$$

$$H_1I_2 |-\rangle |0\rangle = |10\rangle$$

$$CNOT |\Phi_{3}\rangle = \frac{|01\rangle + |11\rangle}{\sqrt{2}} = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) |1\rangle = |+\rangle |1\rangle$$

$$H_{1}I_{2} |+\rangle |1\rangle = |01\rangle$$

$$CNOT |\Phi_{4}\rangle = \frac{|01\rangle - |11\rangle}{\sqrt{2}} = \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) |1\rangle = |-\rangle |1\rangle$$

$$H_{1}I_{2} |-\rangle |1\rangle = |11\rangle$$

So the inverse maps $|\Phi_2\rangle \mapsto |10\rangle$, $|\Phi_3\rangle \mapsto |01\rangle$, and $|\Phi_4\rangle \mapsto |11\rangle$. In this way, a two-bit message can be encoded by applying only an operation to the first qubit. This is a quantum algorithm known as 'superdense coding.'

Part c

Note: Qiskit follows "little endian" order, so the qubits are ordered $|q_1q_0\rangle$ instead of $|q_0q_1\rangle$ as above.

```
qc = QuantumCircuit(2)
#prepare |Phi_1>
qc.h(0)
qc.cx(0,1)
#invert
qc.cx(0,1)
qc.h(0)
qc.measure_all()
print(execute(qc, sim).result().get_counts())
#"{'00': 1024}"
gc = QuantumCircuit(2)
#prepare |Phi_2>
qc.h(0)
qc.cx(0,1)
qc.z(0)
#invert
qc.cx(0,1)
qc.h(0)
qc.measure_all()
print(execute(qc, sim).result().get_counts())
#"{'01': 1024}"
qc = QuantumCircuit(2)
#prepare |Phi_3>
qc.h(0)
qc.cx(0,1)
qc.x(0)
#invert
qc.cx(0,1)
qc.h(0)
qc.measure_all()
print(execute(qc, sim).result().get_counts())
#"{'10': 1024}"
```

```
qc = QuantumCircuit(2)
#prepare |Phi_4>
qc.h(0)
qc.cx(0,1)
qc.y(0)
#invert
qc.cx(0,1)
qc.h(0)
qc.measure_all()
print(execute(qc, sim).result().get_counts())
#"{'11': 1024}"
```