

# 6.215/6.255J/15.093J/IDS.200J Optimization Methods

## Lecture 7: Sensitivity Analysis

September 30, 2021

# Today's Lecture

## Outline

- Last word on duality (for now)
- Sensitivity analysis: Motivation
- Global sensitivity analysis
- Local sensitivity analysis
  - Changes in  $\mathbf{b}$
  - Changes in  $\mathbf{c}$
  - A new variable is added
  - A new constraint is added
  - Changes in  $\mathbf{A}$
- Detailed example

# Duality revisited

- So far: Simplex  $\longrightarrow$  Duality  $\longrightarrow$  Farkas lemma
  - specialized to LP, relied on a particular algorithm (simplex)
- Alternative: Separation theorem  $\longrightarrow$  Farkas lemma  $\longrightarrow$  Duality
  - purely geometric, generalizes to some nonlinear problems, more fundamental

# Background - elements of real analysis

- **Continuous function:** A function  $f : S \rightarrow \mathbb{R}$  is continuous at  $\mathbf{x} \in S$  iff

$$f(\mathbf{x}^i) \rightarrow f(\mathbf{x})$$

for all converging sequences  $\{\mathbf{x}^i\} \in S$  with limit  $\mathbf{x}$ .

- **Closed set:** A set  $S \subset \mathbb{R}^n$  is closed iff when  $\mathbf{x}^1, \mathbf{x}^2, \dots$  is a sequence of elements of  $S$  that converges to some  $\mathbf{x} \in \mathbb{R}^n$ , then  $\mathbf{x} \in S$ .
- **Remarks:**
  - Level sets  $\{\mathbf{x} : f(\mathbf{x}) \leq \alpha\}$ ,  $\{\mathbf{x} : f(\mathbf{x}) \geq \alpha\}$ ,  $\{\mathbf{x} : f(\mathbf{x}) = \alpha\}$  of continuous functions are closed sets for any  $\alpha$ .
  - Every halfspace is closed.
  - Intersections of closed sets are closed.
  - Every polyhedron is closed.

# Weierstrass' theorem

## Theorem

*If  $f : \mathbb{R}^n \mapsto \mathbb{R}$  is a continuous function, and if  $S$  is a nonempty, closed, and bounded subset of  $\mathbb{R}^n$ , then there exists some  $\mathbf{x}^* \in S$  such that  $f(\mathbf{x}^*) \leq f(\mathbf{x})$  for all  $\mathbf{x} \in S$ . Similarly, there exists some  $\mathbf{y}^* \in S$  such that  $f(\mathbf{y}^*) \geq f(\mathbf{x})$  for all  $\mathbf{x} \in S$ .*

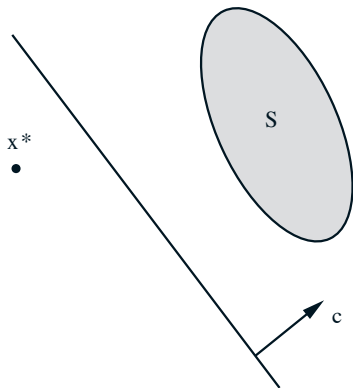
Note: Weierstrass' theorem is not valid if the set  $S$  is not closed.

(for example, consider  $S = \{x \in \mathbb{R} \mid 0 < x \leq 1\}$ ,  $f(x) = x$ )

# Separating hyperplane theorem

## Theorem

Let  $S$  be a nonempty closed convex subset of  $\mathbb{R}^n$  and let  $\mathbf{x}^* \in \mathbb{R}^n$  such that  $\mathbf{x}^* \notin S$ . Then there exists a vector  $\mathbf{c} \in \mathbb{R}^n$  such that  $\mathbf{c}^T \mathbf{x}^* < \mathbf{c}^T \mathbf{x} \quad \forall \mathbf{x} \in S$ .



# Sensitivity Analysis - Motivation

## Questions

$$\begin{array}{ll} z = \min & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{array}$$

## Sensitivity Analysis:

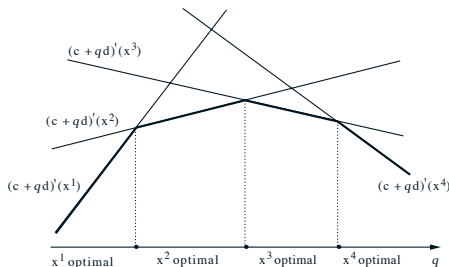
- How does  $z$  depend globally on  $\mathbf{c}$ ? on  $\mathbf{b}$ ?
- How does  $z$  change locally if either  $\mathbf{b}$ ,  $\mathbf{c}$ ,  $\mathbf{A}$  change?
- How does  $z$  change if we add new constraints, introduce new variables?
- Importance: Insight about linear optimization and practical relevance

# Global sensitivity analysis

Dependence on  $\mathbf{c}$

$$\begin{aligned} G(\mathbf{c}) = \min \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A}\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{aligned}$$

- let  $\mathbf{x}^1, \mathbf{x}^2, \dots, \mathbf{x}^N$  be the BFSs in the feasible set (assumed nonempty)
- $\Rightarrow G(\mathbf{c}) = \min_{i=1, \dots, N} \mathbf{c}^T \mathbf{x}^i$  is a **concave** piecewise-linear function of  $\mathbf{c}$





# Global sensitivity analysis

Dependence on  $\mathbf{b}$

Easier to see from the dual:

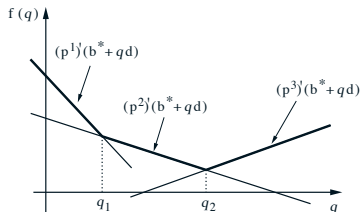
primal

$$\begin{aligned} F(\mathbf{b}) = \min \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A}\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{aligned}$$

dual

$$\begin{aligned} F(\mathbf{b}) = \max \quad & \mathbf{p}^T \mathbf{b} \\ \text{s.t.} \quad & \mathbf{p}^T \mathbf{A} \leq \mathbf{c}^T \end{aligned}$$

- let  $\mathbf{p}^1, \mathbf{p}^2, \dots, \mathbf{p}^N$  be the extreme points of the dual feasible set
- $\Rightarrow F(\mathbf{b}) = \max_{i=1, \dots, N} \{(\mathbf{p}^i)^T \mathbf{b}\}$  is a **convex** piecewise-linear function of  $\mathbf{b}$



# Local sensitivity analysis

- Linear optimization in standard form:

$$\begin{aligned} z = \min \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{aligned}$$

- What does it mean that a basis  $\mathbf{B}$  is optimal?
  - feasibility conditions:  $\mathbf{B}^{-1}\mathbf{b} \geq 0$
  - optimality conditions:  $\mathbf{c}^T - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A} \geq 0$
- Suppose that there is a change in either  $\mathbf{b}$  or  $\mathbf{c}$ .  
How do we find whether  $\mathbf{B}$  is still optimal?
  - check whether the feasibility and optimality conditions are satisfied

# Local sensitivity analysis

## Changes in $\mathbf{b}$

$b_i$  becomes  $b_i + \delta$ , i.e.  $\mathbf{b} \rightarrow \mathbf{b} + \delta \mathbf{e}_i$

$$\begin{array}{ll} (P) & \min \quad \mathbf{c}^T \mathbf{x} \\ & \text{s.t.} \quad \mathbf{Ax} = \mathbf{b} \\ & \quad \mathbf{x} \geq 0 \end{array} \quad \rightarrow \quad \begin{array}{ll} (P') & \min \quad \mathbf{c}^T \mathbf{x} \\ & \text{s.t.} \quad \mathbf{Ax} = \mathbf{b} + \delta \mathbf{e}_i \\ & \quad \mathbf{x} \geq 0 \end{array}$$

- $\mathbf{B}$  optimal basis for  $(P)$
- Is  $\mathbf{B}$  optimal for  $(P')$ ?

# Local sensitivity analysis

Changes in  $\mathbf{b}$

Need to check:

- ① Feasibility:  $\mathbf{B}^{-1}(\mathbf{b} + \delta \mathbf{e}_i) \geq 0$
- ② Optimality:  $\mathbf{c}^T - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A} \geq 0$

Observations:

- ① Changes in  $\mathbf{b}$  affect feasibility
- ② Optimality conditions are not affected

# Local sensitivity analysis

Changes in  $\mathbf{b}$

Feasibility condition:

$$\mathbf{B}^{-1}(\mathbf{b} + \delta \mathbf{e}_i) \geq 0$$

Define

$$\beta_{ij} = [\mathbf{B}^{-1}]_{ij}, \quad \bar{b}_j = [\mathbf{B}^{-1}\mathbf{b}]_j$$

Thus,

$$[\mathbf{B}^{-1}\mathbf{b}]_j + \delta[\mathbf{B}^{-1}\mathbf{e}_i]_j \geq 0 \quad \Rightarrow \quad \bar{b}_j + \delta\beta_{ji} \geq 0 \quad \Rightarrow$$

$$\max_{\beta_{ji} > 0} \left( -\frac{\bar{b}_j}{\beta_{ji}} \right) \leq \delta \leq \min_{\beta_{ji} < 0} \left( -\frac{\bar{b}_j}{\beta_{ji}} \right)$$

# Local sensitivity analysis

Changes in  $\mathbf{b}$

$$\underline{\delta} \leq \delta \leq \bar{\delta}$$

Within this range

- Current basis  $\mathbf{B}$  is optimal
- Optimal cost is affine in  $\delta$ :

$$z = \mathbf{c}_B^T \mathbf{B}^{-1}(\mathbf{b} + \delta \mathbf{e}_i) = \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{b} + \delta p_i$$

- What if  $\delta > \bar{\delta}$ ?

# Local sensitivity analysis

Changes in  $\mathbf{b}$

$$\underline{\delta} \leq \delta \leq \bar{\delta}$$

Within this range

- Current basis  $\mathbf{B}$  is optimal
- Optimal cost is affine in  $\delta$ :

$$z = \mathbf{c}_B^T \mathbf{B}^{-1}(\mathbf{b} + \delta \mathbf{e}_i) = \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{b} + \delta p_i$$

- What if  $\delta > \bar{\delta}$ ?

$\Rightarrow$  current solution is primal infeasible, but satisfies optimality conditions (so dual feasible)  $\Rightarrow$  use dual simplex method

# Local sensitivity analysis

## Changes in $\mathbf{c}$

After perturbation,  $c_j$  becomes  $c_j + \delta$ , i.e.  $\mathbf{c} \rightarrow \mathbf{c} + \delta \mathbf{e}_j$ .

Is current basis  $\mathbf{B}$  still optimal?

Need to check:

- Feasibility:  $\mathbf{B}^{-1}\mathbf{b} \geq 0$ , unaffected
- Optimality:  $\mathbf{c}^T - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A} \geq 0$ , affected

There are two cases:

- $x_j$  nonbasic
- $x_j$  basic



# Local sensitivity analysis

Changes in  $\mathbf{c}$ ;  $x_j$  nonbasic

$x_j$  nonbasic:

- Then  $\mathbf{c}_B$  are unaffected
- Solution remains optimal if

$$(c_j + \delta) - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A}_j \geq 0 \quad \Rightarrow \quad \bar{c}_j + \delta \geq 0$$

so solution still optimal if  $\delta \geq -\bar{c}_j$

- What if  $\delta < -\bar{c}_j$ ?

$\Rightarrow$  apply the primal simplex to the current BFS

# Local sensitivity analysis

Changes in  $\mathbf{c}$ ;  $x_j$  basic

$x_j$  basic:

- Then  $\mathbf{c}_B \rightarrow \hat{\mathbf{c}}_B = \mathbf{c}_B + \delta \mathbf{e}_j$
- Solution optimal if for all  $i \neq j$ :

$$[\mathbf{c}^T - \hat{\mathbf{c}}_B^T \mathbf{B}^{-1} \mathbf{A}]_i \geq 0 \quad \Rightarrow \quad c_i - [\mathbf{c}_B + \delta \mathbf{e}_j]^T \mathbf{B}^{-1} \mathbf{A}_i \geq 0$$

- Defining  $\bar{a}_{ji} := [\mathbf{B}^{-1} \mathbf{A}]_{ji}$

$$\bar{c}_i - \delta \bar{a}_{ji} \geq 0 \quad \Rightarrow \quad \max_{\bar{a}_{ji} < 0} \frac{\bar{c}_i}{\bar{a}_{ji}} \leq \delta \leq \min_{\bar{a}_{ji} > 0} \frac{\bar{c}_i}{\bar{a}_{ji}}$$

- What if  $\delta$  is outside this range?

$\Rightarrow$  apply primal simplex to current BFS

# Local sensitivity analysis

A new variable is added

$$\begin{array}{ll} \min & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{array} \quad \rightarrow \quad \begin{array}{ll} \min & \mathbf{c}^T \mathbf{x} + c_{n+1} \mathbf{x}_{n+1} \\ \text{s.t.} & \mathbf{Ax} + \mathbf{A}_{n+1} \mathbf{x}_{n+1} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{array}$$

In the new problem is  $x_{n+1} = 0$  or  $x_{n+1} > 0$ ? (i.e., is the new activity profitable?)

# Local sensitivity analysis

A new variable is added

Current basis  $\mathbf{B}$ . Is solution  $\mathbf{x} = \mathbf{B}^{-1}\mathbf{b}, x_{n+1} = 0$  optimal?

- Feasibility conditions are satisfied
- Optimality conditions:

$$c_{n+1} - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A}_{n+1} \geq 0 \quad \Rightarrow \quad c_{n+1} - \mathbf{p}^T \mathbf{A}_{n+1} \geq 0?$$

- If yes, solution  $\mathbf{x} = \mathbf{B}^{-1}\mathbf{b}, x_{n+1} = 0$  optimal
- Otherwise, use primal simplex

# Local sensitivity analysis

A new constraint is added

$$\begin{array}{ll} \min & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{array} \rightarrow \begin{array}{ll} \min & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{a}_{m+1}^T \mathbf{x} = b_{m+1} \\ & \mathbf{x} \geq 0 \end{array}$$

If current solution feasible, it is optimal; otherwise, apply dual simplex

# Local sensitivity analysis

## Changes in $\mathbf{A}$

- Suppose  $a_{ij} \rightarrow a_{ij} + \delta$
- Assume  $\mathbf{A}_j$  does not belong in the basis
  - Feasibility conditions:  
 $\mathbf{B}^{-1}\mathbf{b} \geq 0$ , unaffected
  - Optimality conditions:  $c_\ell - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A}_\ell \geq 0$ ,  $\ell \neq j$ , unaffected
  - Optimality condition:  $c_j - \mathbf{p}^T (\mathbf{A}_j + \delta \mathbf{e}_i) \geq 0 \Rightarrow \bar{c}_j - \delta p_i \geq 0$
- What if  $\mathbf{A}_j$  is basic?

see BT exercise 5.3

# Example

A furniture company

- A furniture company makes desks, tables, chairs
- The production requires wood, finishing labor, carpentry labor

	Desk	Table	Chair	Avail.
Profit	60	30	20	-
Wood (ft)	8	6	1	48
Finish hrs.	4	2	1.5	20
Carpentry hrs.	2	1.5	0.5	8

# Example

## A furniture company - formulation

Decision variables:

$x_1 = \#$  desks,  $x_2 = \#$  tables,  $x_3 = \#$  chairs

LO formulation:

$$\begin{array}{llll} \max & 60x_1 + 30x_2 + 20x_3 & & \\ \text{s.t.} & 8x_1 + 6x_2 + x_3 & \leq 48 & \text{(Wood)} \\ & 4x_1 + 2x_2 + 1.5x_3 & \leq 20 & \text{(Finishing)} \\ & 2x_1 + 1.5x_2 + 0.5x_3 & \leq 8 & \text{(Carpentry)} \\ & x_1, x_2, x_3 & \geq 0 & \end{array}$$



# Example

A furniture company - formulation standard form

Decision variables:

$x_1 = \#$  desks,  $x_2 = \#$  tables,  $x_3 = \#$  chairs

$$\begin{array}{llll} \max & 60x_1 + 30x_2 + 20x_3 & & \\ \text{s.t.} & 8x_1 + 6x_2 + x_3 & \leq 48 & \text{(Wood)} \\ & 4x_1 + 2x_2 + 1.5x_3 & \leq 20 & \text{(Finishing)} \\ & 2x_1 + 1.5x_2 + 0.5x_3 & \leq 8 & \text{(Carpentry)} \\ & x_1, x_2, x_3 & \geq 0 & \end{array}$$

In standard form:

$$\begin{array}{llll} z := \min & -60x_1 - 30x_2 - 20x_3 & & \\ \text{s.t.} & s_1 + 8x_1 + 6x_2 + x_3 & = 48 & \text{(Wood)} \\ & s_2 + 4x_1 + 2x_2 + 1.5x_3 & = 20 & \text{(Finishing)} \\ & s_3 + 2x_1 + 1.5x_2 + 0.5x_3 & = 8 & \text{(Carpentry)} \\ & s_1, s_2, s_3, x_1, x_2, x_3 & \geq 0 & \end{array}$$

Note: profit =  $-z$

# Example

A furniture company - simplex tableau(s)

Initial tableau:

		$s_1$	$s_2$	$s_3$	$x_1$	$x_2$	$x_3$
	0	0	0	0	-60	-30	-20
$s_1 =$	48	1	0	0	8	6	1
$s_2 =$	20	0	1	0	4	2	1.5
$s_3 =$	8	0	0	1	2	1.5	0.5

Final tableau:

		$s_1$	$s_2$	$s_3$	$x_1$	$x_2$	$x_3$
	280	0	10	10	0	5	0
$s_1 =$	24	1	2	-8	0	-2	0
$x_3 =$	8	0	2	-4	0	-2	1
$x_1 =$	2	0	-0.5	1.5	1	1.25	0

# Example

A furniture company - information in the final tableau

Problem in standard form:

$$\begin{array}{llll} \min & -60x_1 - 30x_2 - 20x_3 & & \\ \text{s.t.} & s_1 + 8x_1 + 6x_2 + x_3 & = & 48 \\ & s_2 + 4x_1 + 2x_2 + 1.5x_3 & = & 20 \\ & s_3 + 2x_1 + 1.5x_2 + 0.5x_3 & = & 8 \\ & s_1, s_2, s_3, x_1, x_2, x_3 & \geq & 0 \end{array}$$

$$A = \begin{pmatrix} 1 & 0 & 0 & 8 & 6 & 1 \\ 0 & 1 & 0 & 4 & 2 & 1.5 \\ 0 & 0 & 1 & 2 & 1.5 & 0.5 \end{pmatrix}$$

Final tableau - What is  $B$ ,  $B^{-1}$ ?

	$s_1$	$s_2$	$s_3$	$x_1$	$x_2$	$x_3$
280	0	10	10	0	5	0
$s_1 =$	24	1	2	-8	0	-2
$x_3 =$	8	0	2	-4	0	-2
$x_1 =$	2	0	-0.5	1.5	1	1.25

$$B = \begin{pmatrix} 1 & 1 & 8 \\ 0 & 1.5 & 4 \\ 0 & 0.5 & 2 \end{pmatrix}$$

$$B^{-1} = \begin{pmatrix} 1 & 2 & -8 \\ 0 & 2 & -4 \\ 0 & -0.5 & 1.5 \end{pmatrix}$$

What is the dual optimal solution  $p$

$$p^T = c_B^T B^{-1} = (0, -20, -60) \begin{pmatrix} 1 & 2 & -8 \\ 0 & 2 & -4 \\ 0 & -0.5 & 1.5 \end{pmatrix} = (0, -10, -10)$$

# Example

Changes in **b** - furniture company - marginal cost/shadow price

Why is the optimal dual solution  $p_2 = -10$  called a marginal cost or shadow price for the finishing hours constraints?

- Suppose that finishing hours become 21 (from 20).
- Currently only desks ( $x_1$ ) and chairs ( $x_3$ ) are produced
- Finishing and carpentry hours constraints are tight
- Does this modification (from 20 to 21) leave current basis optimal?

New solution:	$8x_1 + x_3 + s_1 = 48$	$\Rightarrow$	New	Previous
	$4x_1 + 1.5x_3 = 21$		$s_1 = 26$	24
	$2x_1 + 0.5x_3 = 8$		$x_1 = 1.5$	2
			$x_3 = 10$	8

Solution change:

$$Z_{\text{new}} - Z_{\text{old}} = (-60 \cdot 1.5 - 20 \cdot 10) - (-60 \cdot 2 - 20 \cdot 8) = -10$$

An increase profit of 10

# Example

Changes in  $\mathbf{b}$  - furniture company - budget range for finishing hours

- Suppose finishing hours change by  $\delta$  becoming  $(20 + \delta)$
- Basic variables

$$\mathbf{x}_B = \mathbf{B}^{-1} \begin{pmatrix} 48 \\ 20 + \delta \\ 8 \end{pmatrix} = \begin{pmatrix} 1 & 2 & -8 \\ 0 & 2 & -4 \\ 0 & -0.5 & 1.5 \end{pmatrix} \begin{pmatrix} 48 \\ 20 + \delta \\ 8 \end{pmatrix} = \begin{pmatrix} 24 + 2\delta \\ 8 + 2\delta \\ 2 - 0.5\delta \end{pmatrix}$$

- Current basis optimal if and only if  $\mathbf{x}_B \geq 0$  or

$$-4 \leq \delta \leq 4$$

- Note that even if current basis is optimal, optimal solution values do change:

$$s_1 = 24 + 2\delta$$

$$x_3 = 8 + 2\delta$$

$$x_1 = 2 - 0.5\delta$$

$$z = -60(2 - 0.5\delta) - 20(8 + 2\delta) = -280 - 10\delta$$

# Example

Changes in **b** - furniture company - reduced cost

- What does it mean that the reduced cost for  $x_2$  is 5?
- Suppose you are forced to produce  $x_2 = 1$  (1 table)
- How much will the profit decrease?

$$\begin{array}{rclcl} 8x_1 + x_3 + s_1 & + 6 \cdot 1 & = 48 & & s_1 = 26 \\ 4x_1 + 1.5x_3 & + 2 \cdot 1 & = 20 & \Rightarrow & x_1 = 0.75 \\ 2x_1 + 0.5x_3 & + 1.5 \cdot 1 & = 8 & & x_3 = 10 \end{array}$$

$$\text{new profit} - \text{old profit} = \Delta_{\text{profit}} = -(z_{\text{new}} - z_{\text{old}})$$

$$\Delta_{\text{profit}} = (30 \cdot 1 + 60 \cdot 0.75 + 20 \cdot 10) - (60 \cdot 2 + 20 \cdot 8) = 275 - 280 = -5$$

## Example

Changes in **b** - furniture company - reduced cost

$$\begin{aligned} (-\text{profit}) \leftarrow z := \min \quad & -60x_1 - 30x_2 - 20x_3 \\ \text{s.t.} \quad & s_1 + 8x_1 + 6x_2 + x_3 = 48 \quad (\text{Wood}) \\ & s_2 + 4x_1 + 2x_2 + 1.5x_3 = 20 \quad (\text{Finishing}) \\ & s_3 + 2x_1 + 1.5x_2 + 0.5x_3 = 8 \quad (\text{Carpentry}) \\ & s_1, s_2, s_3, x_1, x_2, x_3 \geq 0 \end{aligned}$$

Another way to calculate the same thing using optimal dual variables:

$$\mathbf{p}^T = (0, -10, -10)$$

If  $x_2 = 1$

Profit effect due to producing one table	$-(-30) = +30$
Profit effect due to decrease wood by -6	$-(-6 * 0) = 0$
Profit effect due to decrease finishing hours by -2	$-(-2 * -10) = -20$
Profit effect due to decrease carpentry hours by -1.5	$-(-1.5 * -10) = -15$
Total profit effect	-5

Suppose profit from tables increases from \$30 to \$34.

Should it be produced? At \$35? At \$36?

# Example

Changes in  $\mathbf{c}$  - furniture company

- Suppose profit from desks becomes  $60 + \delta$ .
- For what values of  $\delta$  does current basis remain optimal?
- Optimality conditions

$$c_j - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A}_j \geq 0 \iff \mathbf{c} - \mathbf{p}^T \mathbf{A} \geq 0$$

for dual vector

$$\begin{aligned} \mathbf{p}^T &= \mathbf{c}_B^T \mathbf{B}^{-1} = [0, -20, -(60 + \delta)] \begin{bmatrix} 1 & 2 & -8 \\ 0 & 2 & -4 \\ 0 & -0.5 & 1.5 \end{bmatrix} \\ &= [0, -10 + 0.5\delta, -10 - 1.5\delta] \end{aligned}$$

- We have  $c_i - \mathbf{p}^T \mathbf{A}_i = 0$  for all basic variables.
- We need to check for all non-basic variables  $c_j - \mathbf{p}^T \mathbf{A}_j \geq 0$ .



# Example

Changes in  $c$  - furniture company

- Variables  $s_1, x_3, x_1$  are basic
- Reduced costs of non-basic variables

$$\begin{bmatrix} \bar{c}_{x_2} \\ \bar{c}_{s_2} \\ \bar{c}_{s_3} \end{bmatrix} = \begin{bmatrix} 30 \\ 0 \\ 0 \end{bmatrix} - [0, -10 + 0.5\delta, -10 - 1.5\delta] \begin{bmatrix} 6 & 0 & 0 \\ 2 & 1 & 0 \\ 1.5 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} \bar{c}_{x_2} \\ \bar{c}_{s_2} \\ \bar{c}_{s_3} \end{bmatrix} = \begin{bmatrix} 5 + 1.25\delta \\ 10 - 0.5\delta \\ 10 + 1.5\delta \end{bmatrix}$$

- Current basis optimal if  $-4 \leq \delta \leq 20$

## Example

A new variable is added - furniture company

- Suppose the company has the opportunity to produce stools
- Profit \$ 15; requires 1 ft of wood, 1 finishing hour, 1 carpentry hour.
- Should the company produce stools?

$$\begin{array}{rcccccccl} \max & 60x_1 & +30x_2 & +20x_3 & +15x_4 & & & \\ & 8x_1 & +6x_2 & +x_3 & +x_4 & +s_1 & & = 48 \\ & 4x_1 & +2x_2 & +1.5x_3 & +x_4 & & +s_2 & = 20 \\ & 2x_1 & +1.5x_2 & +0.5x_3 & +x_4 & & & +s_3 = 8 \\ & x_1, & x_2, & x_3, & x_4, & & & \geq 0 \end{array}$$

- Reduced cost  $c_4 - \mathbf{c}_B^T \mathbf{B}^{-1} \mathbf{A}_4 = -15 - (0, -10, -10) \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} = 5 \geq 0$
- Current basis still optimal. Do not produce stools