

# 6.215/6.255J/15.093J/IDS.200J Optimization Methods

## Lecture 9: Network Optimization I

October 7, 2021

# Today's Lecture

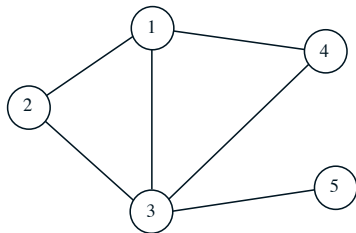
## Outline

- Networks - introduction
- Network flow problems
  - Shortest path problem and applications
  - Maximum flow problem and applications
  - Minimum cost flow problem and applications

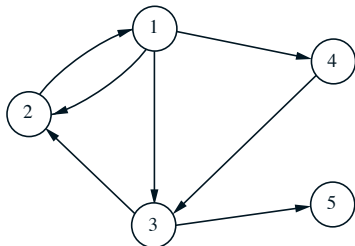
# Networks

What are they?

Formally networks are (undirected or directed) graphs with additional information on nodes and/or on edges or arcs.



undirected  $G = (N, E)$



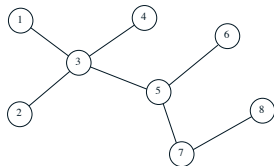
directed  $G = (N, A)$

In our contexts, “self-loops” or “self-arcs” will not be considered.

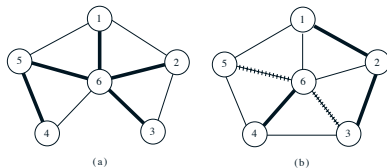
# Networks

## Some special undirected graphs

- A *tree* is an undirected graph that is connected and has no cycles.

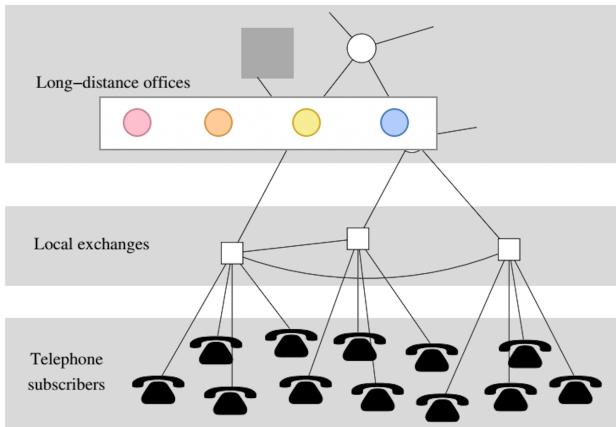


- A *spanning tree* of a graph  $G$  is a subgraph that is a tree and contains all nodes of  $G$ .



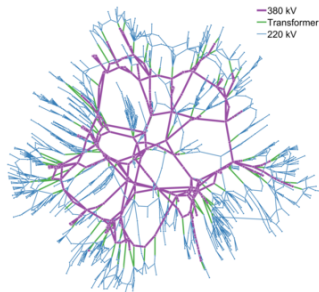
# Networks

## Illustrations - (three-tiers) telephone networks



# Networks

Illustrations - electrical & power networks



# Networks

## Illustrations - road networks



# Networks

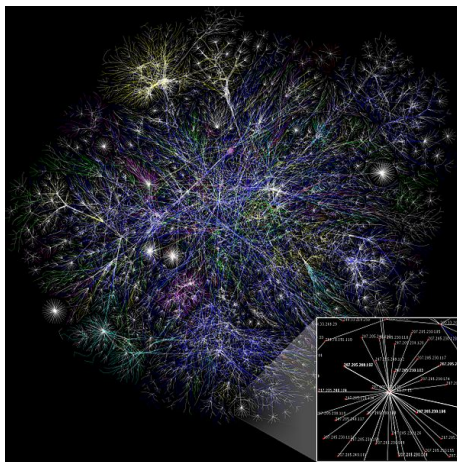
## Illustrations - airline routes





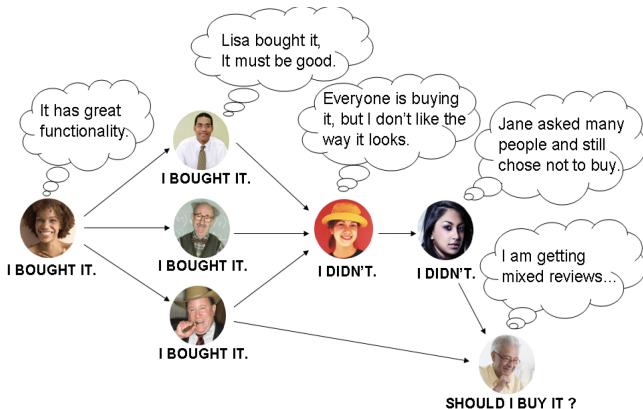
# Networks

## Illustrations - internet backbone



# Networks

## Illustrations - social networks



# Network optimization

## Common thrust

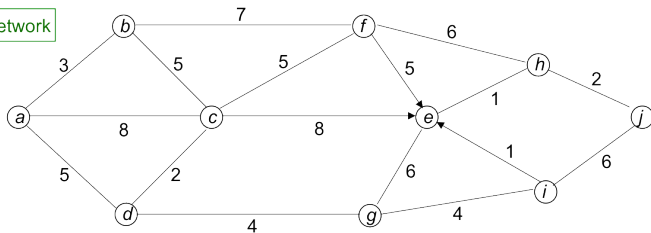
Move some entity (electricity, a consumer product, a person, a vehicle, a message, etc.) from one point to another in the underlying network, as efficiently as possible.

- Note: Many problems are network problems in disguise!
- Today: Learn how to model application settings as network flow problems.
- Next lecture: Study ways to solve the resulting models.

# Shortest Path

## Description

Mixed network



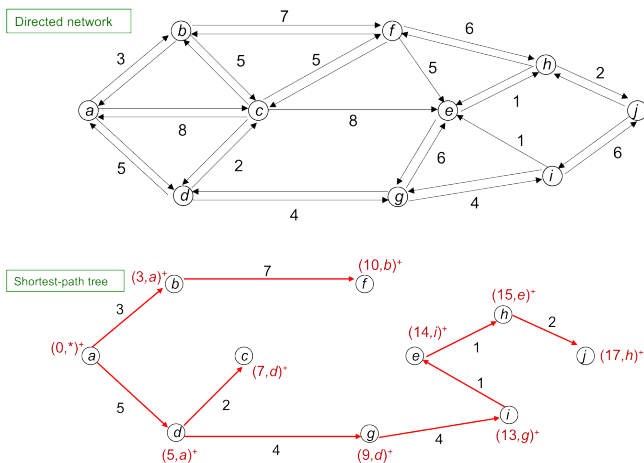
**Problem:** Identify a shortest path from a given source node  $s$  to a given sink node  $t$ .

### Examples:

- a path of minimum length.
- a path taking minimum time.
- a path of maximum reliability.

# Shortest Path

## Shortest path tree solution



# Shortest Path

## Linear optimization formulation

- Graph  $G = (N, A)$ , source node  $s$ , sink node  $t$
- Arc lengths  $c : A \rightarrow \mathbb{Z}$

$$\begin{aligned} \min \quad & \sum_{(i,j) \in A} c_{ij} x_{ij} \\ \text{s.t.} \quad & \sum_{j:(i,j) \in A} x_{ij} - \sum_{j:(j,i) \in A} x_{ji} = 0 \quad \text{for all } i \in N \setminus \{s, t\} \\ & \sum_{j:(s,j) \in A} x_{sj} = 1 \\ & \sum_{j:(j,t) \in A} x_{jt} = 1 \\ & x_{ij} \geq 0 \quad \text{for all } (i,j) \in A \end{aligned}$$

# Shortest Path Formulations

## Illustration 1 - Interword spacing in $\text{\LaTeX}$

The spacing between words and characters is normally set automatically by  $\text{\LaTeX}$ . Interword spacing within one line is uniform.  $\text{\LaTeX}$  also attempts to keep the word spacing for different lines as nearly the same as possible.

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# Shortest Path Formulations

## Illustration 1 - Interword spacing in $\text{\LaTeX}$

- The paragraph consists of  $n$  words, indexed by  $1, 2, \dots, n$ .
- $c_{ij}$  is the attractiveness of a line if it begins with  $i$  and ends with  $j - 1$ .
- $\text{\LaTeX}$  uses a (complicated) formula to compute the value of each  $c_{ij}$ . For instance,

$$c_{12} = -10,000$$

$$c_{13} = -1,000$$

$$c_{14} = -100$$

$$c_{1,37} = -100,000$$

...

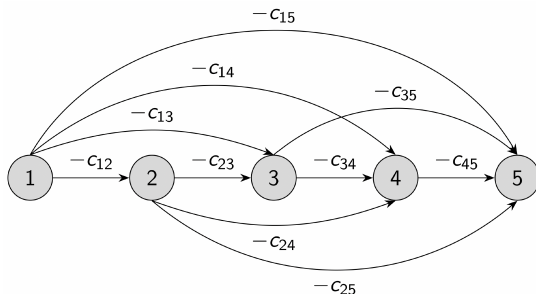
- The problem of decomposing a paragraph into several lines of text to maximize total attractiveness can be formulated as a shortest path problem.
- Nodes? Arcs? Costs?



# Shortest Path Formulations

## Illustration 1 - Interword spacing in $\text{\LaTeX}$

- Introduce a node for each word / syllable (and one terminal node).
- For all  $i < j$  we introduce an edge between word / syllable  $i$  and word / syllable  $j$  with length  $-c_{ij}$ .



# Shortest Path Formulations

## Illustration 2 - Project management

- A project consists of a set of jobs and a set of precedence relations
- In particular, we are given a set  $A$  of job pairs  $(i, j)$  indicating that job  $i$  cannot start before job  $j$  is completed.
- $c_i$  duration of job  $i$
- Find the least possible duration of the project

# Shortest Path Formulations

## Illustration 2 - Project management

### Formulation

- Introduce two artificial jobs  $s$  and  $t$ , of zero duration, that signify the beginning and the completion of the project
- Add  $(s, i)$  and  $(i, t)$  to  $A$
- $p_i$ : time that job  $i$  begins
- $(i, j) \in A$ :  $p_j \geq p_i + c_i$
- Project duration:  $p_t - p_s$

# Shortest Path Formulations

## Illustration 2 - Project management

- Linear optimization formulation:

$$\begin{array}{ll}\min & p_t - p_s \\ \text{s.t.} & p_j - p_i \geq c_i, \quad \forall (i,j) \in A.\end{array}$$

- Dual:

$$\begin{array}{ll}\max & \sum_{(i,j) \in A} c_i x_{ij} \\ \text{s.t.} & \sum_{\{j | (j,i) \in A\}} x_{ji} - \sum_{\{j | (i,j) \in A\}} x_{ij} = b_i \\ & x_{ij} \geq 0\end{array}$$

where  $b_s = -1$ ,  $b_t = 1$ , and  $b_i = 0$  for  $i \neq s, t$ .

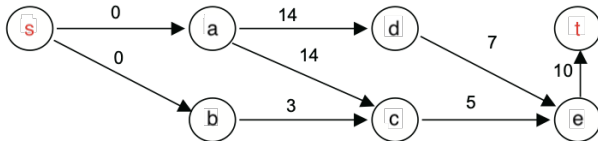
- Shortest path problem, where each precedence relation  $(i,j) \in A$  corresponds to an arc with cost of  $-c_i$ .

# Shortest Path Formulations

## Illustration 2 - Project management

Example:

Activity	Immediate Predecessor	Time ( $c_i$ )
<b>s</b>		0
a	s	14
b	s	3
c	a,b	5
d	a	7
e	c,d	10
<b>t</b>	e	0



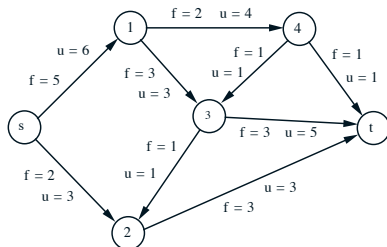
# Maximum Flow

## Description

**Problem:** Determine the maximum flow that can be sent from a given source node  $s$  to a given sink node  $t$  in a capacitated network.

**Examples:** Find maximum steady-state flow of

- petroleum products in a pipeline network
- cars in a road network
- messages in a telecommunication network
- electricity in an electrical network



# Maximum Flow

## Linear optimization formulation

- Network  $G = (N, A)$ , source node  $s$ , sink node  $t$
- Arc capacities  $u : A \rightarrow \mathcal{N}$

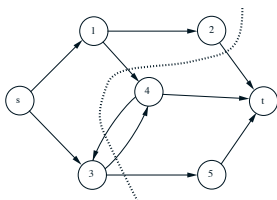
$$\begin{aligned} \max \quad & \sum_{j:(j,t) \in A} x_{jt} \\ \text{s.t.} \quad & \sum_{j:(i,j) \in A} x_{ij} - \sum_{j:(j,i) \in A} x_{ji} = 0 \quad \text{for all } i \in N \setminus \{s, t\} \\ & x_{ij} \leq u_{ij} \quad \text{for all } (i,j) \in A \\ & x_{ij} \geq 0 \quad \text{for all } (i,j) \in A \end{aligned}$$

# Maximum Flow

## Duality - Max Flow Min Cut Theorem

An  $(s, t)$ -cut in a network  $G = (N, A)$  is a partition of  $N$  into two disjoint subsets  $S$  and  $T$  such that  $s \in S$  and  $t \in T$ .

The capacity of an  $(s, t)$ -cut  $(S, T)$  is  $\sum_{i \in S} \sum_{j \in T} u_{ij}$ .



## Theorem

*The value of a maximum  $(s, t)$ -flow = the capacity of a minimum  $(s, t)$ -cut.*



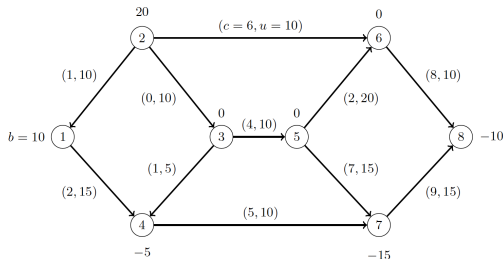
# Min-Cost Flow

## Description

**Problem:** Determine a least cost shipment of a commodity through a network in order to satisfy demands at certain nodes from available supplies at other nodes. Arcs have capacities and cost associated with them

## Examples:

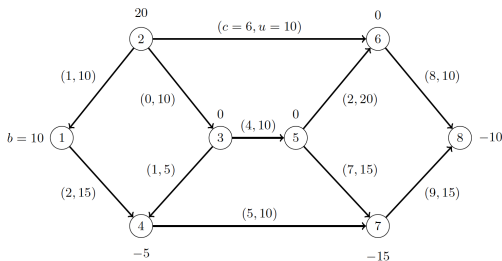
- Distribution of products
- Flow of items in a production line
- Routing of cars through street networks
- Routing of telephone calls



# Min-Cost Flow

## Linear optimization formulation

- Network  $G = (N, A)$ .
- Arc costs  $c : A \rightarrow \mathcal{Z}$ .
- Arc capacities  $u : A \rightarrow \mathcal{N}$ .
- Node balances  $b : N \rightarrow \mathcal{Z}$ .

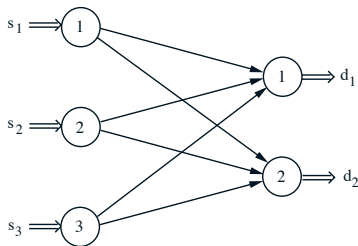


$$\begin{aligned} \min \quad & \sum_{(i,j) \in A} c_{ij} x_{ij} \\ \text{s.t.} \quad & \sum_{j: (i,j) \in A} x_{ij} - \sum_{j: (j,i) \in A} x_{ji} = b_i \quad \text{for all } i \in N \\ & x_{ij} \leq u_{ij} \quad \text{for all } (i,j) \in A \\ & x_{ij} \geq 0 \quad \text{for all } (i,j) \in A \end{aligned}$$

# Min-Cost Flow

## Some special cases

### Transportation problem:



**Assignment problem:** A transportation problem where:

- number of suppliers = number of customers
- each supplier has unit supply
- each customer has unit demand

# Min-Cost Flow Formulation

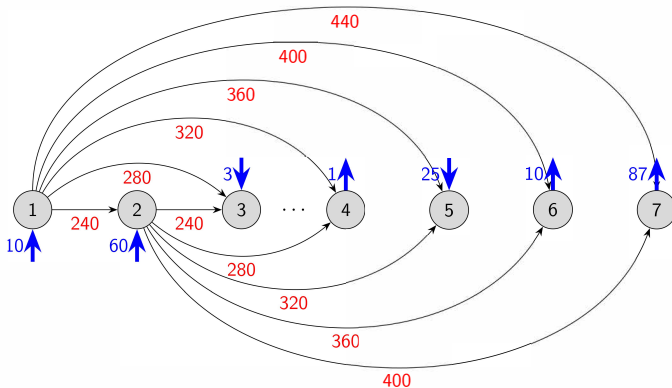
## Illustration: Passenger routing

- United Airlines has seven daily flights from BOS to SFO, every two hours, starting at 7am.
- Capacities are 100, 100, 100, 150, 150, 150, and  $\infty$ .
- Passengers suffering from overbooking are diverted to later flights.
- Delayed passengers get \$200 plus \$20 for every hour of delay.
- Suppose that today the first six flights have 110, 160, 103, 149, 175, and 140 confirmed reservations.

Determine the most economical passenger routing strategy!

# Min-Cost Flow Formulation

Illustration: Passenger routing



# Further Network Problems

- Minimum spanning tree problems,
- Matching problems,
- Postman problems,
- Generalized flow problems,
- Multicommodity flow problems,
- Constrained shortest path problems,
- Unsplittable flow problems,
- Network design problems,
- ...