AN INTERPRETATIVE FLOATING DECIMAL SYSTEM FOR THE IBM TYPE 650

G. R. Trimble, Jr.

E. C. Kubie

Introduction

١.

This floating decimal system will perform 18 basic operations using a floating decimal number system by means of interpretative programming. It was designed with coding convenience as a prime objective. This is illustrated by the fact that one of the basic operations is essentially a vector by vector multiplication.

Floating decimal instructions, that is, instructions which are to be interpreted, have a negative sign. A floating decimal instruction consists of a two digit operation code, a 4 digit address specifying the location of the first factor, a 4 digit address specifying the location of the second factor (if required) and another 4 digit address specifying the location of the third factor (if required). It is a variable address system in that only as many addresses as are needed for the particular operation are required. Thus, some operations use only one address, some require two addresses and others require 3 addresses.

Floating decimal instructions will be taken from consecutive memory locations. If an instruction requires one or two addresses, the instruction is stored in one memory location. If three addresses are required, the third address is stored in the memory location immediately following the one containing the instruction.

Should a positive instruction appear, it is interpreted as a normal 650 instruction and subsequent instructions are not interpreted. Thus, upon occurrence of a positive instruction the 650 operates in its usual D-I mode. This continues until an instruction address of 0026 is given which causes control to return to the interpretative routine. The program will return to the floating decimal mode of operation at the point of departure.

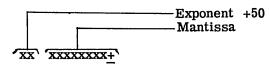
For example, consider the following sequence:

Location	Contents
n	floating point instruction (-)
n+1	floating point instruction (-)
n+2	floating point instruction (-)
n+3	normal 650 instruction (+)

(The contents of n+3 are interpreted as a normal 650 instruction including the instruction address for sequencing. Normal execution of instructions continues until 0026 is used as an instruction address, at which time the program returns to the floating decimal mode of operation beginning with the instruction in location n+4).

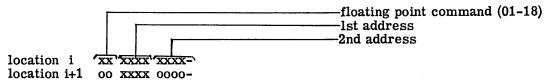
Number Form

The number form is as follows:



Instruction Form

The instruction form is as follows:



i+1 has this form only if the third address is needed, otherwise, the form of i+1 is the same as the form of i.

Floating Decimal Accumulator

A floating point accumulator referred to as K is used for accumulation and accumulative multiplication. It does not need to be addressed in operations which make use of it. For example, operation 01 (A+B —)K) uses two addresses, A and B. The sum is automatically stored in K as the result of this operation.

Operations

Below is a list of the operations with their codes, addresses required and estimated average time of execution.

Code and Addresses	Operation	Estimated Average Time
01, A, B	$A+B\longrightarrow K$	62.8
02, A	$A+K\longrightarrow K$	54.7
03, A, B	A- B-→K	63.1
04, A	$K-A\longrightarrow K$	59.5
05, A, B, C	$A+B\longrightarrow C$	81.8

Code and Addresses	Operation	Estimated Average Time
06, A, B, C	A-B-→C	96.2
07, A, B, C	AxB—→ C	84.8
08, A, B	AxB—→ K	78.4
09, A, B, C	$A/B \longrightarrow C$	92.2
10, A, B	A/B → K	72.2
11, A, B	BR MIN A	28.4
12, A	BR	18.7
13, A, B	BRNZ A	30.8
14, A, B	(AxB)+K→K	102.5
15, A, B	K-(AxB)→K	102.5
16, A, B	$\sqrt{A} \longrightarrow B$	157.9
17, A, B, C	$\sqrt{A^2+B^2+C^2}$ \rightarrow K	416.3
18, A ₁ , B ₁ , n	$ \sum_{i=1}^{n} A_{i}B_{i} \rightarrow K $	(32.4+92.3n)

Explanation of Programs

1. General Interpretation

The general interpretation routine takes instructions from consecutive memory locations and analyzes them to see if they are normal 650 instructions or if they must be interpreted. If the instruction is a normal 650 instruction, it is executed as such. If the instruction is to be interpreted, the routine obtains the factor at address A and stores it in location 0037. Control is then transferred to the proper sub-routine. The constant in location 0193 facilitates use of the translating routine TR1. The amount of translation is placed in the instruction address positions of this constant. This amount is then added to the operation code and control is transferred to the translated sub-routine.

2. Addition Sub-Routine

The addition sub-routine adds the factors in locations 0037, in which A is stored, and 0057, which is the floating decimal accumulator K. The result is normally stored in the floating decimal accumulator K.

3. Multiplication Sub-Routine

The multiplication sub-routine multiplies the factors in 0037, which is A, and 0089, in which B is placed. The product is normally stored in 0037.

4. Division Sub-Routine

The division sub-routine divides the factor in 0037 which is A, by the factor B, which had previously been placed in the lower accumulator. The result is normally stored in 0057 which is the floating decimal accumulator.

The following sub-routines interpret the second and third addresses if required, obtain the necessary factors and modify the addition, multiplication and division sub-routines so as to store the result in the desired location. In some cases, they also modify portions of other sub-interpretative routines.

5. Interpretation of 01

This sub-interpretative routine obtains the factor B and transfers control to the addition sub-routine.

6. Interpretation of 02

Since the factors are already in place all that is necessary for this subinterpretative routine is to transfer control to the addition sub-routine.

7. Interpretation of 03

This sub-interpretative routine obtains the factor B, reverses its sign and transfers control to the addition sub-routine.

8. Interpretation of 04

The sign of the factor A is reversed and control is transferred to the addition sub-routine.

9. Interpretation of 05

Since the addition sub-routine makes use of the floating decimal accumulator K, the contents of K must be temporarily transferred to another location, and returned after the operation has been completed. This sub-interpretative routine obtains B, modifies the addition sub-routine to store the result in location C and transfers control to the addition sub-routine. After completion of the operation, control is returned to the sub-interpretative routine and the addition sub-routine is restored to normal.

10. Interpretation of 06

This sub-interpretative routine modifies the sub-interpretative routine, "Interpretation of 05" so that the sign of factor B is reversed.

11. Interpretation of 07

This sub-interpretative routine obtains the factor B, modifies the multiply sub-routine to store the product in location C and transfers control to the multiply sub-routine. After completion of the multiplication, control is returned to this sub-interpretative routine and the multiply sub-routine is restored to normal.

12 Interpretation of 08

The factor B is obtained, the last instruction of the multiply sub-routine is modified to store the product in K and control is transferred to the multiply sub-routine. After completion of the multiplication, control is returned to the sub-interpretative routine "Interpretation of 07" which restores the last instruction of the multiply sub-routine to normal.

13. Interpretation of 09

The factor B is obtained, the last instruction of the divide sub-routine is modified to store the quotient in C and control is transferred to the divide sub-routine. After completion of the division, control is returned to this sub-interpretative routine and the last instruction of the divide sub-routine is restored to normal.

14. Interpretation of 10

The factor B is obtained and control is transferred to the divide sub-routine.

15. Interpretation of 11

The sign of the factor in location A is examined. If the sign is minus, control is transferred to location B by modifying the general interpretation routine. Sequencing of instructions then continues in the normal fashion beginning with B. If the factor in location A is positive or its mantissa is zero, sequencing of instructions continues in the normal fashion.

16. Interpretation of 12

The general interpretation routine is modified so that the next instruction is taken from location A. Sequencing of instructions then continues in the normal fashion beginning with A.

17. Interpretation of 13

The factor in A is examined and if its mantissa is not zero, the general interpretation routine is modified to take the next instruction from location B. Sequencing of instructions then continues in the normal fashion beginning with B. If the factor in A is zero, sequencing of instructions continues in the normal manner.

18. Interpretation of 14

The factor B is obtained, the last instruction in the multiply sub-routine is modified to store the product in 0037, which is A, and control is transferred to the multiply sub-routine. After completion of the multiplication, control is returned to the sub-interpretative routine, the last instruction of the multiply sub-routine is restored to normal and control is transferred to the addition sub-routine which computes the desired sum.

19. Interpretation of 15

The sign of the factor B is reversed and control is transferred to the sub-interpretative routine "Interpretation of 14".

20. Interpretation of 16

This sub-interpretative routine computes the required square root and stores it in location B. The initial approximation is $X_0 = (1+4A)/(4+A)$, obtained from the RAND Corporation "Approximations in Numerical Analysis" form 15S, Notes.

21. Interpretation of 17

Factor B is obtained and temporarily stored in 0265. Factor A is squared and stored in 0057 by the multiplication sub-routine. Factor B is then squared and added to the square of A. Factor C is then obtained, squared and added to $A^2 + B^2$. Finally, control is transferred to the square root sub-routine and $\sqrt{A^2+B^2+C^2}$ is computed and stored in 0057 which is K.

22. Interpretation of 18

Operation 18 is a vector by vector multiplication. Address A_1 and address B_1 are the addresses of the first elements of each of the vectors. Succeeding elements of the vectors are then taken from consecutive memory locations starting with the initial locations A_1 and B_1 . The third address, n, indicates the number of elements in each vector.

Use is made of the sub-interpretative routine "Interpretation of 14" to perform accumulative multiplication. As each pair of factors are multiplied together and added to the previous sum, n is reduced by 1 and zero tested. When n has been reduced to zero, the last two factors have been multiplied and the operation is complete.

Storage

This system requires 466 storage locations. It uses locations 0000 through 0465. Every location within this block is used, thus making it easy to incorporate this program with other programs. The translating routine, TR1 may be used to translate this program to any desired block of memory locations. It should be translated by an even amount however to preserve the even-odd conditions.

Localium 0017, 0173, 1212, 10353, 0372 4 0420 an inter-20 Se freshing 0017 -0024 an inter-to-cities Apr 1 25 20 freshing of American Se School De addition of 6 Mars of Section of American

FORM NO. 22-6181-0

IBM Trade-Mark

IBM TYPE 650 PROGRAM SHEET

PROBLEM: General Interpretation WRITTEN BY:_____

LOCATION	OPERAT	ION	ADDR	ESS	REMARKS
OF NSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	
0026	RAU	60	0029	0092	
0092	AU	ĬŎ	0099	0118	Increase i to i+1
0118	ST U	21	0029	8001	
8001	RAL	65	i+1	0439	Pick up instruction i+1
0439	BRMIN	46	0025	8002	Is instruction i+1 to be interpreted?
0025	LD	69	0028	0031	
0031	ST DA	22	0035	8001	Put data at address A in location
8001	LD	69	A	0455	0037.
0455	SLT	35	0002	0334	
0334	ST D	24	0037	0040	
	AU	10	0193	8003	Transfer to proper sub-routine.
8003	NO OP	00	0000	00fc	
0000	7.4.7	CE		0439	
0029	RAL	65	1 0001		Constants
0099		00	0001	0000	Constants
0028	LD	69	0000		
0193		00	0000	0000	

PROBLEM: (A) + (K) \longrightarrow K Subroutine WRITTEN BY: WRITTEN BY: WRITTEN BY:

NSTRUCTION ABBRV. CODE DATA INSTRUCTION REMARKS	LOCATION	OPERAT		ADDRI		i, ii- a, mammar
0051 RAL 0061 SLT 35 0002 0067 0067 STU 21 0073 0027 0027 STL 20 0081 0034 0034 RAL 65 0037 0041 0041 SLT 35 0002 0055 0048 STU 21 0062 0048 0048 STU 21 0062 0055 0062 RAABL 0073 0077 0069 SABL 18 0073 0077 0077 SLT 35 0004 0087 0077 SLT 35 0004 0087 0078 SABL 18 0073 0077 12 12 12 12 12 12 12 12 12 12 12 12 12 1						REMARKS
0061 SLT 35 0002 0067			65			0057• K
OBST						1
0027						2
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						A → 0081
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						0037: A
0048 STU 21 0052 0055						<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0048					a ₂ → 0052
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		STL		0059	0062	$A_2 \rightarrow 0059$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0062	RAABL	67	8003	0069	\
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0069	SABL		0073		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$				0004	0087	for case where $0 < a_2 - a_1 < 10$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0087	LD	69			
10056 RAL	0043					1
0078						$ls a_2 = a_1 ?$
0038 BRNZU 44 0091 0044 Is $A_{\frac{1}{2}} > 10$? 0091 SRD 31 0003 0101 $A_{\frac{1}{2}} > 10$. Thus $A_3 = 0.1$ $A_{\frac{1}{2}} > 0058 0058 0058 0058 0058 0111 0112 0065 0070 $						$a_2 = a_1$. Thus $A_3 = A_2 + A_1$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0078					
0101 SLT 35 0002 0058	0038					IS AS DIV
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						$ A_3 \ge 10$. Thus $A_3 = 0.1 A_3$
$\begin{array}{cccccccccccccccccccccccccccccccccccc$						* 4 <00
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$						Is A ₃ < 0?
0127 SLT 35 0008 0045 Put exponent in normal positions 0045 AL 15 8003 0053 (2 high order digit positions) 0112 AL 15 0065 0120 A3 > 0. Increase exponent by 1 0120 AABL 17 0073 0127 0044 BRNZ 45 0049 0053 $ A_3^1 \le 10$. Is $A_3^1 = 0$? 0049 RAU 60 8002 0066 SCT 36 0000 0088 0 $< A_3^1 \le 10$. Its $A_3^1 = 0$? 0066 SCT 36 0000 0088 0 $< A_3^1 \le 10$. Thus $A_3^1 = 0$? 0105 RAL 65 8003 0116 shifted and counted. 0105 RAL 65 8003 0116 Is $A_3 < 0$? 0104 I.D 69 0057 0060 1s $A_3 < 0$? 0074 BRNZ 45 0080 0079 1s $A_3 < 0$? 0080 BR MIN 46 0033 0084 $ A_3 > 0 <$						A3<0. Increase exponent by I
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0070					Dut ownered in normal positions
0112 AL 15 0065 0120 A ₃ >0. Increase exponent by 1 0120 AABL 17 0073 0127 0044 BRNZ 45 0049 0053 $ A_3 \le 10$. Is $A_3^1 = 0$? 0049 RAU 60 8002 0066 0066 SCT 36 0000 0088 0 $ A_3 \le 10$. Thus A_3 equals A_3^1 0088 SU 11 8002 0105 shifted and counted. 0105 RAL 65 8003 0116 0116 BR MIN 46 0070 0120 Is $A_3 \le 0$? 0104 LD 69 0057 0060 0060 SRT 30 0005 0074 $ A_3 \le 1$ Is $ A_3 = 1$ Is $ A_3 = 1$ 10? 0060 SRT 30 0005 0074 $ A_3 = 1$ Is $ A_3 = 1$ 10? 0074 BRNZ 45 0080 0079 0084 $ A_3 = 1$ Is $ A_3 = 1$ 10. Is $ A_3 = 1$ 10. O084 RAL 65 0037 0053 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 110. O084 RAL 65 0037 0053 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 110. O084 RAL 65 0059 0064 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0079 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0079 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0079 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0074 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0074 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0074 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0074 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0075 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 10. O084 SRT 30 0002 0074 $ A_3 = 1$ 10. Thus $ A_3 = 1$ 11 11 11 11 11 11 11 11 11 11 11 11						(2 high order digit positions)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0112					A > 0 Increase exponent by 1
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0120					113 U. Increase exponent by 1
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0044					$ A_{3} < 10$, Is $A_{3}^{1} = 0$?
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$) 13
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			36		0088	$0 \le Ab \le 10$. Thus A3 equals A5
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0088			8002		shifted and counted.
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$		RAL	65	8003		
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	0116	BR MIN	46		0120	Is A ₂ < 0 ?
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$				0057	0060	
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$a_2 \neq a_1$. Is $ a_2 - a_1 \ge 10$?
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$				0080		· ·
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$ a_2 - a_1 \ge 10$. Is $a_2 > a_1$?
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$a_1 \ge a_2 + 10$. Thus $a_3 A_3 = a_1 A_1$
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$a_{\overline{2}} > a_1 + 10$. Thus $a_3 A_3 = a_2 A_2$
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$0 < a_2 - a_1 < 10$. Is $a_2 > a_1$?
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$		1				
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						$\frac{1}{1}$ $\frac{2}{2}$ $\frac{2}{1}$ $\frac{2}$
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$					0098	K
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						Compute Al
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						Compare A3
0068 ST D 24 0073 0076	0039		69		0068	K
0076 RAL 65 0081 0036 $\begin{cases} a_2 > a_1$. Thus $A_3^2 = A_2 + \text{shifted } A_1$ 0036 SRT 30 0002 0046 Also, use a_2 to compute a_3						
0036 SRT 30 0002 0046 Also, use as to compute as						$a_0 > a_1$. Thus $A_0^1 = A_0 + \text{shifted } A_0$
					1	
						/

PROBLEM: (A) + (K) - K Subroutine

WRITTEN BY: _____

	(A) = a	L A1.	(K) = a	$A_2 a = 1$	xx, $A = x$, $xxxxxxx$ +
LOCATION	OPERAT	TION	ADDRESS		REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0053	ST L	20	0057	0026	Store sum in K
0090	: 31 000	0 007	1		Constants
0065	: 00 000	000.00			
1					

	$(A) = m_1 M_1, (B) = m_2 M_2, B$		12 W12 , I	11=XX, 1V1=X, XXXXXXX +	
LOCATION	OPERAT	TION	ADDR	ESS	
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0134	RAL	65	0037	0141	0037: A. Separate exponent and
0141	SLT	35	0002	0097	mantissa exponent and
0097	STL	20	0151	0106	M ₁ →0151 M ₁ =x. xxxxxxxx00
0106	RAU	60	8003	0063	K m1_20121 m1 =x.xxxxxxx00
0063	SRT	30	0002	0119	(m. 60
0119	RAABL	67	8002	0128	\ m ₁ - 50
0113	SL	16		T	
0085	AU	10	0131 0089	0085 0093	0089: B. Test to see if M ₂ <0
0093					Outs: B. Test to see II M2
0095	BR MIN	11	0096 8001	0147 0103	
0103	SU SL	16	8001	0103	1 10
0161	SLT			0101	$M_2 < 0$
		35 10	0002		Put M ₂ in upper accumulator and M ₃ = m ₁ + m ₂ -50 in lower accumulator. Acc: odMMMMMMMM/mmooooooo
0117 0075	AU	61	8002	0075	$M_3 = m_1 + m_2 - 30 \text{ in lower accumulator.}$
	RSU		8003	0138	ACC: OCIVITATIATIATIATIATIATA UTUTOOOOOOO
0138	SRT	30	0002	0095	
0147	SŪ	11	8001	0155	
0155	AL	15	8001	0113	$M_2 \geqslant 0$.
0113	SLT	35	0002	0121	Put M ₂ in upper accumulator and
0121	AU	10	8002	0115	Put M ₂ in upper accumulator and m ₃ = m ₁ + m ₂ -50 in lower accumulator
0115	RAU	60	8003	0138	<i>y</i>
_0095	MPY	19	0151	0082	$M_3' = M_1 \times M_2$
0082	BR MIN		0135	0086	
0135	SL	16	0139	0094	
0086		15	0139	0094	Round M3 in 7th decimal place
0094	LD	69	_8003	0100	
0100	SLT	35	0002	0107	/
0107	SRT	30	0001	0163	Test to see if M ₂ is of the form
0163	SCT	36	0001	0170	ox. xxxxxxx or xx. xxxxxxxx
0170	BROV	47	0123	0126	
0123	SRT	30	0009	0181	
0181	SRD	31	0002	0102	
0102	ST DA	22	0205	0108	
0108	ST IA	$\bar{2}\bar{3}$	0211	0114	Ma is of the form ox. xxxxxxxx
0114	RAL	65	8001	0171	Round Ma in 8th decimal place
0171	BR MIN	46	0124	0125	Mg is of the form ox.xxxxxxxx Round Mg in 8th decimal place to get Mg.
0124	AL	15	0178	0083	<u> </u>
0125	SL	16	0178	0083	
0083	ST L	20	0037	0026	
0126	RAL	65	8001	0042	M_3^1 is of the form xx.xxxxx. Thus, $M_3^2 = M_3 + 1$
0042		46	0000	0050	$m_0^3 = m_0 + 1$
0000	SL	16	0109	0083	
0050	AT.	15	0109	0083	J
0131	• 500000	იიიი			
0139	: 500000		1		Constants
0178	000000	hoos	1		/ WIII COLLEGE
0109	: 010000				<i>)</i>
0109	ATOOOD!	μουσ			
1	10 F	4.:		11	-1 1/1 // 2

all 9's. The will smootidate the 2 subsequent sters Addresses instructions.

PROBLEM: (A)/(B) \longrightarrow K Subroutine WRITTEN BY: (A) = d_1D_1 , (B) = d_2D_2 d=xx, D= x.xxxxxx

LOCATION	OPERAT		ADDRI		, D- A. AAAAAA
OF	UPERAI				REMARKS
INSTRUCTION	ABBRV.	CODE		INSTRUCTION	
0213	SLT	35	0002	0169	
	STU	21	0174	0177	d2 → 0174 D2 → 0081
0177	STL	20	0081	0184	$D_2 \longrightarrow 0.081$
0184	RAL	65	0037	0191	
	SLT	35	0002	0148	
0148	STU	21	0059	0263	d ₁ → 0059
0263	RAU	60	8002	0221	
0221	SRT	30	0001	0227	$D_3 = D_1/D_2$
	DIV RU	64	0081	0464	
0464	BR MIN	46	0159	0209	
0159	SL	16	0162	0167	Round Di in 8th decimal place.
	AL	15	0162	0167	
0167	SLT	35	0001	0223	Test to see if rounded D ₂ is of the
0223	BRNZU	44	0130	0228	Test to see if rounded D ₂ is of the form x.xxxxxxxxx or .xxxxxxxxx
0228	BR MIN		0132	0182	}
0132	AL	15	0185	0189	$D_{\mathbf{Q}}^{1}$ is of the form o.xxxxxxxxx, thus
0182	SL	16	0185	0189	D_3^1 is of the form o.xxxxxxxxx, thus round D_3^1 in the 9th decimal place
0189	SRT SLT	30	0002	0145	to get D3.
0145	SLT	35	0002	0152	<i>J</i>
0152	BR MIN	46	0255	0156	
0255	SABL	18	0059	0164	
0164	SL	16	0217	0321	1
0321	AABL	17	0174	0179	$d_3 = d_1 - d_2 + 50 - 1$
0156	AABL	17	0059	0214	<u> </u>
0214	AL	15	0217	0371	
0371	SABL	18	0174	0179	
0130	SRT	.30	0003	0140	D_3^1 is of the form x.xxxxxxxxx, thus
0140	SLT	35	0002	0149	$D_3 = D_3^{\frac{1}{3}}$
0149	BR MIN	46	0153	0154) 3 -3
0153	SABL	18	0059	0313	
0313	SL	16	0166	0321	$d_3 = d_1 - d_2 + 50$
0154	AABL	17	0059	0363	1 3 - 2
0363	AL	15	0166	0371	
0179	RAU	60	8002	0137	
0137	SRT	_30	0002	0143	Put d ₃ in normal position.
0143	AU	10	8002	0202	
0202	STU	21	0057	0026	
0162	:000000				
0185	:0000000				Constants
0217	:000000				
0166	:000000	b050			<u> </u>
	322300			[[
			 		
				 	
			ļ		
		ļ	_		
L		L	<u> </u>	Li	

France-Mark						
PROBLEM	1: INTE	RPRE'	TATION_	OF 01:	_ WRITTEN BY:	
	(A) +	(B) —	→ K			
LOCATION	OPERAT	-	ADDR	ESS		REMARKS
OF	ABBRV.	CODE	DATA	INSTRUCTION	1	
0001	SLT	35	0002	0157		
0157	LD	69	0110	0165		
0165	ST DA	22	0073	8001	Put B in 0057	
8001	RAL	65	<u>B</u>	0453	}	
0453	STL	20	0057	0054		
			<u> </u>	 		
				 		
0110	: 65000	00453				
0110						
IBM			IBM TY	PE 650 F	PROGRAM SHEET	FORM NO. 22-6181 Printed in U.S.A.
Trade-Mark						
	A. TNITTI	rpppi	מרוא שורה א	OF 02.	_ WRITTEN BY:	
PRUBLEN	/ι: <u></u> ://	· (K) -	-K	OF UL.	_ •••••••••••••••••••••••••••••••••••••	
	OPERAT			ESS		
LOCATION OF INSTRUCTION				INSTRUCTION		REMARKS
0002	NO OP	0005	DATA 0000	0054		, i
0002	NO OP	00	UUUU	VVUX	,	
IBM			IBM TY	PE 650 I	PROGRAM SHEET	FORM NO. 22-6181 PRINTED IN U.S.A
Trade-Mark						PRINTED IN U.S.A.
					MOITTEN DV	• · · · · ·
PROBLEM				OF 03:	_ WHILEN BY:	`
		(B) —				
LOCATION	OPERA	TION	ADDR			REMARKS
INSTRUCTION	ABBRV.	CODE		INSTRUCTION		
0003	SLT	35	0002	0160		
	LD	69	0215	0168	D.+ D:- 0057	
0168	ST DA	22	0073	8001	Put -B in 0057	
8001	RSL	66	B	0453		,
0453	STL	20	0057	0054		
		1				
0215	: 66000	00453	<u> </u>	 		
L			<u> </u>	<u> </u>		
TDM			IBM TY	PF 650	PROGRAM SHEET	FORM NO. 22-618
IBM Trade-Mark						PRINTED IN U.S.A
PROBLE	Ν: <u>ΙΝΤΙ</u>	<u>ERPRI</u>	ETATION	OF 04:	_ WRITTEN BY:	
	(K) -	- (À)—	→K			
LOCATION	OPERA	TION	ADDF	RESS	ì	REMARKS
INSTRUCTION	T	CODE	DATA	INSTRUCTION	1	
0004	RSL	66	0037	0133	Put -A in 0037	

PROBLEM: INTERPRETATION OF 05. WRITTEN BY: ______

	(A) +	(B) —	7		
LOCATION	OPERAT	ION	ADDR	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARAS
0005	SLT	35	0002	0212	
0212	LD	69	0216	0220	Modify instruction to get B
0220	ST DA	22	0273	0176	
0176	RAU	60	0029	0183	
0183	AU	10	0186	0192	Increase i+1 to i+2
0192	AL	15	8003	0199	
0199	LD	69	0252	0206	
0206	ST DA	22	0059	8001	Get third address (C) from j+2
8001	AABL	17	i+2	0463	` '
0463	LD	69	0057	0262	Put K in 0265 temporarily
0262	STD	24	0265	0218	
0218	SL	16	8003	0175	Store i+2 in 0029
0175	ST D	24	0029	0187	
0187	LD	69	0142	0196	Modify add routine to store sum
0196	ST DA	22	0053	0273	in C.
0273	RAL	65	В	0453	Get (B) and transfer to add
0453	ST L	20	0057	0054	routine
				ļ	
0188	T.D	69	0265	0195	Return K to 0057
0195	ST D	24	0057	0194	
0194	LD	69	0197	0150	Restore last instruction of Add
0150	ST D	24	0053	0026	routine to normal.
	·	ļ		ļ	
				 	
0010		150		 	
0216	: 650000			 	}
0186	<u> </u>				1
0252	: 170000				Constants
0142	: 200000	<u> </u>			
0197	: 200057	P026			
	ļ	l		11	

IBM

IBM TYPE 650 PROGRAM SHEET

FORM NO. 22-6181-0 PRINTED IN U.S.A.

PROBLEM: INTERPRETATION OF 06: WRITTEN BY:

(A) - (B) C										
LOCATION	OPERAT	TION	ADDR	ESS	REMARKS					
INSTRUCTION	ABBRV.	CODE		INSTRUCTION						
0006	SLT	35	0002	0264	Modify interpretation of 05 to					
0264	LD	69	0267	0220	obtain -B					
				<u> </u>						
	00000	00459								
0267	: 00000	00453		 						
<u> </u>										

PROBLEM: INTERPRETATION OF 07: WRITTEN BY: _____

LOCATION	OPERATION ADDRESS		ecc l		
INSTRUCTION	OFERA				REMARKS
		CODE	DATA	INSTRUCTION	
0007	SLT	35	0002	0266	
0266	RSL	66	8002	0225	
0225	AU	10	0029	0233	Modify instruction to get B
0233	I.D	69	0236	0190	
0190	ST DA	22	0243	0246	
0246	RAL	65	8003	0203	
0203	AL	15	0207	0261	
0261	LD	69	0314	0268	Increase i+1 to i+2. Get C and
0268	ST DA	22	0073	8001	
8001	SU	11	i+2	0456	modify last instruction of multiply
0456	ST L	20	0029	0239	
0239	AU	10	0242	0198	routine to store product in C
0198	ST U	21	0083	0243) I daime to bloze product in o
0243	RAL	65	B1	77.50456	Get B and store it in 0089. Go
0456	ST L	20	0089	0134	to multiply sub-routine.
- 西暦年間					
0454	LD	69	0226	0230	Restore last instruction of multiply
0230	ST D	24	0083	0026	routine to normal
		L		[
0236:	_6500000				
0207:	-0000010				
0314:	1100000				Constants
0242:	2000000	454			
0226:	2000370	026		/	

IBM

IBM TYPE 650 PROGRAM SHEET

PRINTED IN U.S.A.

LOCATION	OPERAT	TION	ADDF	RESS	OSMANIA
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
8000	SLT	35	0002	0315	
0315	LD	69	0219	0172	
0172	ST DA	22	0277	8001	Get B and store it in 0089
8001	LD	69	В	0462	
0462	ST D	24	_0089_	0244	
0244	LD	69	0158	0129	Modify last inst. of mult. routine
0129	ST D	24	0083	0134	to store product in K, then go to
				<u> </u>	0454 to restore last inst. of mult.
				<u> </u>	routine to normal.
0219:	6900000				
0158:	20005704	154		<u> </u>	Constants

IBM TYPE 650 PROGRAM SHEET

FORM NO. 22-6181-0 PRINTED IN U.S.A.

IBM

PROBLEM: <u>INTERPRETATION OF 09:</u> WRITTEN BY: ____

	(11)/(11)				
LOCATION	OPERAT	ION	ADDR	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	
0009	SLT	35	0002	0317	
0317	RSL	66	8002	0276	
0276	AU	10	0029	0234	Modify instruction to get B
0234	LD	69	0237	0240	and the second s
0240	ST DA	22	0243	0248	
0248	AII	10	0201	0208	
0208	RAL	65	8003	0316	Increase i+1 to i+2, get C and
0316	AU	10	0269	0224	Increase i+1 to i+2, get C and modify last instruction of divide
0224	ST L	20	0029	0232	routine to store quotient in C.
0232	LD	69	0285	0238	
0238	ST DA	22	0241	8001	
8001	SU	11	i+2	0461	
0461	ST U	21	0202	0243	
0243	RAL	65	B	0213	Get B and go to divide routine.
0293	LD	69	0296	0249	Restore last instruction of divide
0249	ST D	24	0202	0026	routine to normal.
0237:	65000002	13		ļ	
0201:	00000100	000			
0269	21000002	93			Constants
0285:	11000004	61		ļ	
0296:	21005700		ļ		/
1		l			

IBM

IBM TYPE 650 PROGRAM SHEET

FORM NO. 22-6181-0 PRINTED IN U.S.A.

PROBLEM: INTERPRETATION OF 10: WRITTEN BY: ____

	(// (/	,			
LOCATION			ADDR	ESS	REMARKS
NSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0010	SLT	35	0002	0222	Y 23
0222	LD	69	0237	0290	Put B in lower accumulator and go
0290	ST DA	22	0243	8001	to divide routine
8001	RAL	65	В	0213	
				<u>'</u>	
0237	650000	0213			Constant
	000000	0213			Omstant

PROBLEM		RPRE'		OF 11:	WRITTEN BY:
LOCATION	OPERA	TION	ADC	RESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	TE MATERIA

LOCATION	OPERAT	ION	ADDR	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	TI CONTAIN O
0011	SRT	30	0002	0417	
0417-	AU	1.0	0037	0378	Test for minus. If plus, do not branch
0378	BR MIN	1 46	0382	0026	
0382	SLT	35	0002	0339	
0339	SRT	30	0002	0397	Test for zero. If zero, do not branch
0397	BRNZU	44_	0352_	0026	
0352	_SLT	35_	0004	0420	
0420	LD	69	0373	0426	Modify 0029 and get next
0426	ST DA	22_	0029	8001	instruction
8001	RAL	65_	B	0439	2
1 27	RBL	6.5	7.7.		
	99		1.0		
0373:	650000	439			Constant

IBM TYPE 650 PROGRAM SHEET

FORM NO. 22-6181-0 PRINTED IN U.S.A.

PROBLEM: INTERPRETATION OF 12: WRITTEN BY: _____

LOCATION	OPERAT	ION	ADD	RESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	
0012	SRT	30	0002	0420	
0420	LD	69	0373	0426	Put A in 0029
0426	ST DA	22	0029	8001	
0373:	65000004	39			Constant
				ļ i	-

ibm

IBM TYPE 650 PROGRAM SHEET

FORM NO. 22-6181-0 PRINTED IN U.S.A.

		1 11011			
LOCATION	OPERAT	ION	ADDRI	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	
0013	SRT	30	0002	0423	\
-0423	-A-U	-10	0037	- 0431	
0431	BR MIN	46	0382	0336	Put - A in upper accumulator
0336	SU	11	8001	0443	
0443	SU	11	8001	0382	
0382	SLT	35	0002	0339	
0339	SRT	30	0002	0397	Test for zero. If zero, do not branch
0397	BRNZU	44	0352	0026	
0352	SLT	35	0004	0420	
0420	LD	69	0373	0426	Modify 0029 and get next
0426	ST DA	22	0029	8001	instruction
8001	RAL	65	В	0439	<u> </u>
00/2	\$5.7		<u> </u>	27	
7	wert.	1. 4.	2.55 × 2.5	6.221	
3.5.7	7771	10		11111	
0373;	6500000	439			Constant
			<u> </u>	1	

IBM TYPE 650 PROGRAM SHEET

PROBLEM: INTERPRETATION OF 14: WRITTEN BY: _____

LOCATION	OPERAT	TION	ADDR	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0014	SLT	35	0002	0422	
0422	LD	69	0375	0278	
0278	ST DA	22	0081	8001	Put B in 0089
8001	RAL	65	В	0358	
0358	ST L	20	0089	0392	
0392	LD	69	0398	0357	Modify last instruction of multiply
0357	STD	24	0083	0134	routine to store product in 0037
					•
0399	LD	69	0402	0355	Restore last instruction of multiply
0355	ST D	24	0083	0054	routine to normal
		<u> </u>			
0055					
0375:	650000				
0398:	200037			L	Constants
0402:	200037	0026			

IBM

IBM TYPE 650 PROGRAM SHEET

PROBLEM: <u>INTERPRETATION OF 15:</u> WRITTEN BY: ____

LOCATION OF INSTRUCTION	OPERATION		ADDRESS		REMARKS
	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0015	SLT	35	0002	0421	Modify "Interpretation of 14" to
0421	LD	69	0424	0278	obtain -B
V.2					
0424:	660000	358			Constant

PROBLEM: INTERPRETATION OF 16: WRITTEN BY: _____

LOCATION	OPERA	TION	ADDI	RESS	
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	REMARKS
0016	LD	69	0326	0329	
0329	SLT	35	0002	0335	Modify instruction to store square
0335	ST DA	22	0241	0306	root in B
0306	RAL	65	0037	0323)
0323	SLT	35	0002	0279	Separate at A1. Store A1 in 0283
0279	ST L	20	0283	0286	beparate at Al. Store Al in UZ65
0286	RAL	65	8003	0343	
0343	AU	10	8001	0251	
0251	AU	10	8001	0210	1/2 a ₁
0210	AII	10	8003	0367	1/2 4]
0367	AU	10	8002	0325	
0325	SRT	30	0001	0281	
0281	SU	11	8003	0289	Is a ₁ even or odd?
0289	SLT	35	0003	0245	is at even or odd?
0245	BRNZU	44	0250	0200	
0250	RAU	60	8001	0307	
0307	AU	10	0260	0415	0. over 00-1/9 e.d. 95
0415	SRT	30	0002	0322	a_1 even. $a_2 = 1/2 a_1 + 25$
0322	ST L	20	0277	0330	a2=0277
0330	RAII	60	0283	0287	a2=0211
0287	SRT.	30	0002	0294	
0200	RAII	60	8001	0257	
0257	AU	10	0260	0365	2
0365	SRT	30	0200	0365	$a_1 \text{ odd.} a_2 = 1/2 \ a_1 + 25$
0272	ST L	20	0277	0272	20077
0280	RAU	60	0283	0337	a ₂ → 0277
0337	SRT	30	0001	0294	
0294	AL	15	8003	0301	
0301	AL	15	0204	0259	Store A in 0265.
	ST U	21	0265	0318	Store 4+A in 0273.
	ST L	20	0273	0376	Store 4+A III 0273.
0376	RAU	60	8003	0284	
0284	AU	10	8001	0291	Compute 1+4A
0291	AU	10	8003	0299	Compute 144A
0299	AU	10	0253	0258	
	SRT	30	0001	0270	$X_0 = (1+4A)/(4+A)$
0270	DIV RU	64	0273	0440	$\Delta_0 = \frac{(1+4A)/(4+A)}{4}$
0440	AU	10	0265	0369	
	SL	16	8002	0303	Test to see if A is zero.
	BRNZ	45	0380	0274	TEST TO SEE II A 18 ZEFO.
		64	8001	0460	
	AI,	15	8001	0413	
	SRT	30	0001	0320	
	AU	10	8002	0379	V 1/9 /V . A /V.)
		15	8001	0379	$X_1 = 1/2 (X_0 + A/X_0)$
	AL	15	8002	0295	
		15	8003	0303	
	RAT.	65	8002	0303	
	AII	10	0265		
0370	SL.	16	8002	0370	722
			OUUZ	0381	

PROBLEM: INTERPRETATION OF 16: WRITTEN BY: _____

	√(A) —	♣B				
LOCATION	OPERAT	ION	ADDR	ESS		REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION		
0381	DIV RU	64	8001	0451		
0451	AL	15	8001	0304	1_	
0304	SRT	30	0001	0312		
0312	AU	10	8002	0374		$X_2 = 1/2 (X_1 + A/X_1)$
0374	AL	15	8001	0282		
0282	AL	15	8002	0292	┸	
0292	AL	15	8003	0300		
0300	RAL	65	8002	0310	_	
0310	AU	10	0265	0419		
0419	SL	16	8002	0377		
0377	DIV RU	64_	8001	0452	1_	
0452	AL	_15	8001	0256	4	
0256	SRT	30	0001	0364	1_	$X_3=1/2 (X_2+A/X_2)$
_0364	AU	_10_	8002	0324	1	
0324	AL	_15	8001	0331	<u> </u>	
0331	AL	15	8002	0340		
0340	AL	15	8003	0450		
0450	RAL	65	8002	0409		
0409	AU	_10	0265	0180	\	
0180	SL	16	8002	0144	1	
0144 0458	DIV RU AL	64 15	8001 8001	0458 0459	-	4 /0 / 1 / 1
	AL				4	$X_4 = 1/2 (X_3 + A/X_3)$
0459	SRT	30	0001	0235	_	
0235	AU	10	8002	0146	1	
0146	AL	15	8001	0405	 	
0405	AL	15_	8002	0231	-	
0231	AL	15	8003	0247	-	
0247	RAU	60	8002	0305	<u> </u>	N alima the management and
0305	SCT	36	0000	0361	-	Normalize the square root and reset overflow circuit.
0361	BROV	47	0366	0366	(-	reset overnow circuit.
0366	SRT	30_	0002	0274	-	Det
0274	AU	10	0277	0241		Put exponent in normal position
0241	STU	21_	В	0026	/	and store result in B
0236:	91,0000	000	1	+	\leftarrow	
	2100000		- 1.76	24 / 722 /	Η-	Constants
0260: 0204:	-0000000 040000		1		1	Constants
0253:	0100000				- [
U493:	1 010000	עטען	 	 	/	
L		L			L	

PROBLEM: INTERPRETATION OF 17: $\sqrt{(A)^2 + (B)^2 + (C)^2} \longrightarrow K$

WRITTEN BY: _____

LOCATION	OPERAT	ION	ADDRI	ESS	REMARKS
INSTRUCTION	ABBRV.	CODE		INSTRUCTION	N2 IIIA IIIA
0017	SLT	35	0002	0275	
0275	LD	69	0037	0136	
0136	ST D	24	0089	0342	
0342	LD	69	0345	0298	
0298	ST DA	22	0351	8001	Prepare to compute (A) ² and
8001	RAL	65	В	0408	
0408	LD	69	0427	0430	put (A) ² in 0057
0430	ST D	24	0083	0346	1
0346	ST L	20	0265	0134	temporarily put (B) in 0265.
0360	RAL	65	0265	0333	
0333	ST L	20	0089	0297	Prepare to compute (B) ² and
0297	LD	69	0254	0308	
0308	ST D	24	0083	0141	put (B) ² in 0037
					· ' '
0341	LD	69	0395	0349	Prepare to compute (A) ² + (B) ² and store it in 0057.
0349	ST D	24	0053	0054	store it in 0057.
0362	RAL	65	0029	0383	
0383	AL	15	0186	0391	
0391	LD	69	0394	0347	
0347	ST DA	22	0351	0354	Increase i+1 to i+2. Get C
0354	AU	10	8001	0411	
0411	STL	20	0029	8003	and prepare to compute (C) ²
8003	RAL	65	i+2	0466	
0466	SL	16	0338	8002	and put it in 0037
8002	RAL	65		0437	
0437	ST L	20	0089	0396	
0396	LD	69	0302	0368	
0368	ST D	24	0083	0141	
0390	LD	69	0393	0348	1
0348	ST D	24	0053	0309	Prepare to compute $(A)^2+(B)^2+(C)^2$
0309	_LD	69_	0226	0429	and store it in 0037
0429	ST D	24	0083	0054	<u>/</u>
		l			
0328	LD	69	0332	0288	Restore last instruction of add
0288	ST D	24	0241	0344	routine to normal
0344	_LD	69	0197	0350	Prepare to compute $\sqrt{(A)^2+(B)^2+(C)^2}$
0350	ST D	24	0053	0306	<u> </u>
		L	 	ļ	
0345:	650000		ļ		
0427:	200057	0360	<u> </u>	ļ	
0254:	200037	0341		 	
0395:	200057		ļ	ļ	
0394:	650000		ļ	 	
0338:	650000		<u> </u>	ļ	Constants
0302:	200037		ļ	<u> </u>	
0393:	200037			ļ	/
0332:	210057	0026	1	<u></u>	

	Σ (Ai)	$x (B_i)$	≻ K		
LOCATION	L=1 OPERATION		ADDRESS		REMARKS
INSTRUCTION	ABBRV.	CODE	DATA	INSTRUCTION	NEMANAS
0018	SRT	30	0002	-0425	
0425	SU	11	0029	0384	Store (18 A ₁ B ₁) in 0389
0384	ST L	20	0389	0442	` ' '
0442	RSL	66	್ಷ:8003	0401	JET 445 A.)
0401	ST U	21	0057	0412	Put zeros in K
0412	AL_	15	0416	0432	
0432	LD	69	0385	0388	Increase i+1 to i+2 and get n
0388	ST DA	22	0241	8001	
8001	SU	11	i+2	0457	
0457	ST L	20	0029	0433	
0433	SU	11	0386	0441	Store n-1 in 0386
0441	LD	69	0444	0447	Modify last instruction of add
0447	ST D	24	0053	0356	routine to return control to 410
0356	ST U	21	0414	0434	
0434	_RAL	65	0389	0445	Prepare to get B ₁
0445	SLT	35	0004	0422	
0410	RAU	60_	0414	0418	Test to see if n has been reduced
0418	BRNZ	45	0122	0438	to zero.
0122	_SU	11	0428	0436	
0436	AABL	17	0389	0446	Decrease n by 1. Increase A; and
0446	AL	15	0449	0359	B _i by 1.
0359	STU	21	0414	0435	
0435	STL	20	0389	0448	
0448	LD	69	0403	0406	
0406	ST DA	22	0059	8001	Get A _{i+1} and prepare to get B _{i+1}
8001	LD	69	Aj+1 0037	0407	
0407	STD	24	0037	0404	
0404	SLT	35	0004	0422	
0438	LD	69	0197	0400	Restore last instruction of add
0400	ST D	24	0053	0026	routine to normal.
		ļ			
0410	000001	,			
0416: 0385:	000001			ļ	
0386:	110000 000001				<u> </u>
0444:	200057				Constants
0444:	000001			 	Constants
0449:	000001				
0449:	690000	0001 0407			
V400.	00000	D-101			
L	1	<u> </u>	L		