This is a two-in-one summary.

The first paper, *BI* as an Assertion Language [1], (apparently) introduces the magic wand operator —* and gives a semantics for freeing memory. Magic wands are implications with a fresh piece of memory as their premise.¹ The semantics for freeing memory came as a consequence of giving a classical model for Reynolds' then-unnamed intuitionistic logic about heaps & pointers.

The second paper, Reynolds's survey on separation logic [2], distills \sim 2 years of research into a smooth introduction with many examples. It contains no theorems or "new" results, but rather demonstrates separation logic specifications & simple proofs. Also, the paper coins the term "separation logic", emphasizing that splitting the heap (or generally, resource) into disjoint sections is the key insight.

Strengths

- Separation logic is very important. If I have to talk about memory, this is the specification langauge I want to use. And these papers started it all. (I hope today's presentation gives an update on how much of Reynolds 3-page future work section has been checked off, and what new ideas have come up along the way.)
- The BI paper is going over my head now, but the LICS paper is very easy to understand.

Weaknesses

- Equating datatypes (lists & trees) with propositions about the heap layout upset me. Different types
 can have similar in-memory representations! Equating these types in separation logic will break my
 type safety guarantees.
 - My takeway from those sections is that I hope separation logic become obselete. It's very good that we have it today, to validate the many C programs that are born and exist, but in the future I hope we rely on semantics-preserving compilers and do our formal methods in a higher-level language.
- Can we use separation logic to talk about the physical layout of machines? For instance, reasoning about address spaces and pages swapped out to disk—giving * a physical semantics. This seems like it would be an easy connection to make after extolling modular reasoning and I'm surprised neither paper mentions it.
- Now that I think of locality and caches, I remember that coherence is an important issue. Separation logic does not allow disjoint & overlapping memory, but this would be important to model in multicore or distributed settings—"coherent" overlaps with a specification for resolving conflicts.
 - This also sounds very hard. We suddenly have two unique stores (references in two caches) that aren't exactly shared, but happen to correspond to one store in the main memory.

I can't resist making two typesetting comments—one weakness, one strength:

- What happened to the margins in the BI paper?
- John Reynolds was very good at making diagrams.

References

- [1] Samin Ishtiaq and Peter O'Hearn. BI as an assertion language for mutable data structures. In POPL, 2001.
- [2] John C. Reynolds. Separation logic: A logic for shared mutable data structures. In LICS, 2002.

¹The authors of both papers especially note that A →*B does not imply $A \Rightarrow B$ i.e. this is a novel operator and not just a specialization of linear logic's $-\infty$.