School of Information Technology, Deakin University

SIT320 — Advanced Algorithms

Module Six — Graphs II

Overview

ChatGPT Link

GitHub Link

Task 0: Overview

Key Points (Lecture notes)

- Shortest Path Problem If there are no weight we can just comment about how many edges were traversed. Usually used to find the sum of the weights of the shortest path from node A to node B.
 - Single Source Shortest Path Find the shortest path from a single source to all other vertices in the graph.
 - * Dijkstra's Algorithm Greedy algorithm that finds the shortest path from a single source to all other vertices in the graph. Does not work when there is negative weights in the graph. I used a simple list to store the distances from the source to each vertex.
 - * Bellman-Ford Algorithm Dynamic programming algorithm that finds the shortest path from a single source to all other vertices in the graph. Loops through all the edges in the graph and relaxes them. Relaxation is the process of updating the distance of a vertex from the source vertex. The analogy is when a spring is relaxed it is at its shortest distance. A second pass is done to detect negative cycles. If a negative cycle is detected then the algorithm will return a message that there is a negative cycle in the graph. Slower than Djikstra's algorithm and does not use a labelling mechanism for which vertices have been visited. O(m*n) where m and n are the number of edges and vertices respectively.
 - All Pairs Shortest Path Find the shortest path between all pairs of vertices in the graph.
 - * Floyd-Warshall Algorithm Dynamic programming algorithm that finds the shortest path between all pairs of vertices in the graph. Uses an **inefficient** triple nested loop. Can handle negative weights and cycles, but **may produce the wrong output** for the latter.
- Interaction with the Lecturer/Peers

- Trying to active participate in the chat
- Sometimes I would ask questions and sometimes I would answer questions. Not always the best asnwers though!
- I want to answer more questions in the future, but I feel the pressure of answering on the spot =)
- Thanks to Adrian for his input. His explanations were very clear and concise and helped me understand the material better.

Task 1: Turning Dijkstra's into a Bellman-Ford

Key Points

- The challenge of turning Djikstra's into a Bellman-Ford is that Djikstra's is a greedy algorithm and Bellman-Ford is a dynamic programming algorithm
- The first step was to understand what problem Bellman-Ford was trying to solve
 - Bellman-Ford solves the problem of a graph with negative weights and cycles
- The second step was to understand how Bellman-Ford works
 - Bellman-Ford works by relaxing all the edges in the graph n-1 times where n is the number of vertices in the graph Djikstra Vs Bellman-Ford
- The main difference between Djikstra's and Bellman-Ford is that Djikstra's uses a Greedy design whereas Bellman-Ford uses a Dynamic Programming design Test Cases
- The main thing to test for was graphs with negative weights and cycles Algorithm Analysis
- The algorithm is O(m*n) where m and n are the number of edges and vertices respectively
- This is worse than Djikstra's which is $O((n+m)\log(n))$ where n and m are the number of vertices and edges respectively

Task 2: Floyd-Warshall Algorithm - The most adaptable but inefficient of the shortest path algorithms

Key Points

- The Floyd-Warshall algorithm is a dynamic programming algorithm that finds the shortest path between all pairs of vertices in the graph
- The key point here is **ALL PAIRS** of vertices
- All pairs means that the algorithm uses a systematic approach to find the shortest path between all pairs of vertices Floyd-Warshall Algorithm
- the fact that Floyd-Warshall find all pairs means that it is a triple nested loop
- This negatively impacts the efficiency of the algorithm Handling Negative Weights and Cycles

- Negative weights and cycles are **no problem** for Floyd-Warshall
- The algorithm will still produce the correct output Algorithm Implementation
- The algorithm is implemented using a triple nested loop
- The first loop is used to iterate through the vertices
- The second loop is used to iterate through the rows of the adjacency matrix
- The third loop is used to iterate through the columns of the adjacency matrix **Test Cases**
- I testing with a simple graph that had no negative weights or cycles
- I then tested the algorithm with a graph with negative weights and cycles **Algorithm Analysis**
- The running time of the algorithm is $O(n^3)$ where n is the number of vertices in the graph
- This is **not good** for large graphs
- A real life usage of the algorithm is in the routing of packets in a network

Readings

- Algorithms Illuminated Part 2 & 3 Tim Roughgarden
 - Chapter 9: Dijkstra's Shortest-Path Algorithm
 - Chapter 13: Introduciton to Greedy Algorithms
 - Chapter 16: Introduction to Dynamic Programming
 - Chapter 18: Shortest Paths Revisited
- Introduction to Algorithms CLRS
 - Chapter 15: Dynamic Programming
 - Chapter 16: Greedy Algorithms
 - Chapter 24: Single-Source Shortest Paths
 - Chapter 25: All-Pairs Shortest Paths