

```
... object to mirror...
mirror_mod.mirror_object = ...
operation = "MIRROR_X":
mirror_mod.use_x = True
mirror_mod.use_y = False
mirror_mod.use_z = False
operation = "MIRROR_Y":
mirror_mod.use_x = False
mirror_mod.use_y = True
mirror_mod.use_z = False
operation = "MIRROR_Z":
mirror_mod.use_x = False
mirror_mod.use_y = False
mirror_mod.use_z = True

...selection at the end -add...
..._ob.select= 1
..._ob.select=1
...context.scene.objects.active
...selected = st (mod if)
..._ob.select=0
...py.context.selected_obj
...data.objects[one.name].selec

... (pls select exactly
... OPERATOR CLASSES -----
... types.Operator):
... X mirror to the selected
... _ob.mirror_mirror_x"
```

2020-21 COMP3234A

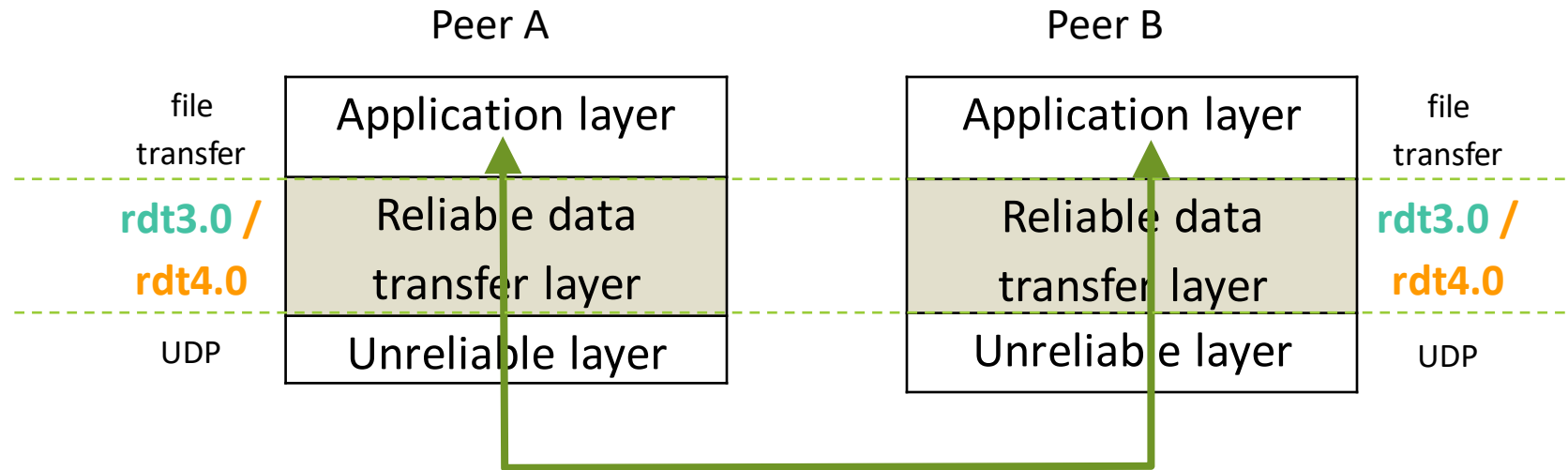
Programming Project

STOP-AND-WAIT (RDT3.0) ARQ AND

EXTENDED-STOP-AND-WAIT (RDT4.0) ARQ

Overview

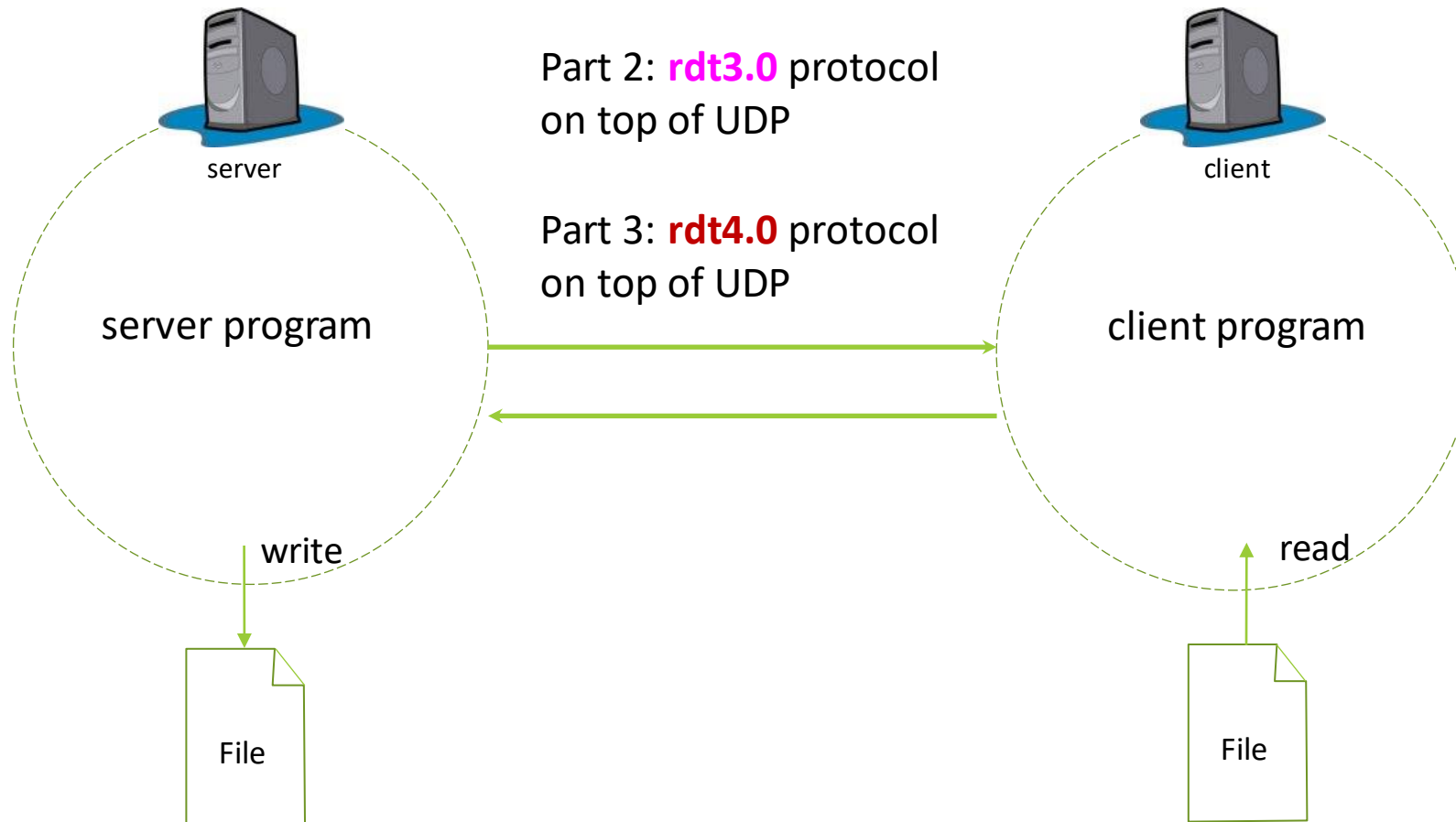
Our RDT protocols support **connectionless** reliable **duplex** data transfer on top of unreliable UDP



Objectives

- To get better understanding of the principles behind Stop-and-Wait protocol;
- To understand the performance difference between a pipelined protocol and the Stop-and-Wait protocol;
- To gain experience in using socket functions to implement a real-life protocol
- The project is divided into **three parts** to make it more attainable
- An assessment task related to ILO4 [Implementation]

Application



Service Interface of the RDT layer

<code>rdt_network_init()</code>	To set up the simulation environment.
<code>rdt_socket()</code>	To create a RDT socket.
<code>rdt_bind()</code>	To assign address info used by this RDT socket.
<code>rdt_peer()</code>	To inform the system the address info of a remote peer.
<code>rdt_send()</code>	To reliably transmit an application message to the targeted remote peer through this RDT socket.
<code>rdt_rcv()</code>	To block and wait for a message from the targeted remote peer.
<code>rdt_close()</code>	To close this RDT socket.

Note: Our RDT layer offers a slightly different Service Interface as compared to TCP and UDP; in particular, we have the functions `rdt_network_init()` and `rdt_peer()` that do not exist in standard socket interface.

Structure of the Project

Part 1 Warm up

Examine rdt1.py

- **Assume UDP is reliable**
- Implement the reliable layer directly on top of UDP without adding extra functionality to UDP

rdt_send(),
rdt_rcv(),
rdt_close()

Part 2 [9 points]

Implement rdt3.py

- **UDP is unreliable** with packet losses and corruptions
- Implement the reliable layer using **Stop-and-Wait** (rdt3.0) ARQ on top of UDP

rdt_send(),
rdt_rcv(),
rdt_close()

Part 3 [9 points]

Implement rdt4.py

- **UDP is unreliable** with packet losses and corruptions
- Implement the reliable layer using **Extended-Stop-and-Wait** (rdt4.0) on top of UDP to improve performance

rdt_send(),
rdt_rcv(),
rdt_close()

rdt_socket(), rdt_bind(), rdt_peer()

Part 1 – rdt1.py

Download Part1.zip

Examine and Test

- rdt1.py
 - There are six rdt_xxxxx() functions
 - rdt_socket(), rdt_bind(), rdt_peer(), rdt_send(), rdt_recv(), & rdt_close()
 - rdt_send() and rdt_recv() use `__udt_send()` and `__udt_recv()` for all communications
 - We have implemented these two internal functions, which consists of the main logic that simulates the underlying unreliable network
- Take note of the difference between rdt_bind() and rdt_peer()
 - rdt_bind() is for setting the address info ("mailbox address") of the UDP socket; so the socket can be used for receiving message
 - rdt_peer() is for specifying the address info of the remote peer, so that we can reference to the peer's address info easily during communication

Part 1 – rdt1.py

- Testing platform – any platform with python3 installed
 - To run the server: `python3 test-server1.py localhost`
 - To run the client: `python3 test-client1.py localhost <<filename>>`
 - Always **start server process first** before executing client program **otherwise**, you **may experience intermittent transmission errors** in the client process if server process is missing
- Test cases
 - small file (around 30 KB)
 - medium size (around 500 KB)
 - large file (around 10 MB)
- Script files
 - run-simulation1.bat, run-simulation1-OSX.sh, run-simulation1-Ubuntu.sh

Submission Deadline:
April 7 17:00

Part 2 – rdt3.py

Download Part2-template.zip

test-client2.py, test-server2.py
run-simulation2 script files

rdt3.py

Part2-sample-output.pdf

Task

- Complete rdt3.py
 - For rdt_socket(), rdt_bind(), & rdt_peer(), you can reuse Part 1 code. But you may have to make minor changes to fit for rdt3.0 logic.
 - Add the reliable logic (rdt3.0) to
 - rdt_send(), rdt_recv(), & rdt_close()
 - **must use __udt_send() and __udt_recv() for all communications**

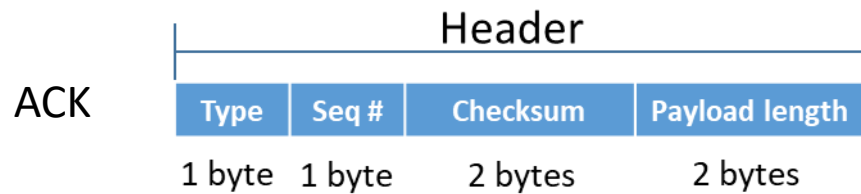
Part 2 – Simulate Losses & Errors

```
def __udt_send(sockd, peer_addr, byte_msg):
    :
    :
    else:
        #Simulate packet loss
        drop = random.random()
        if drop < __LOSS_RATE:
            #simulate packet loss of unreliable send
            print("WARNING: udt_send: Packet lost in unreliable layer!!")
            return len(byte_msg)

        #Simulate packet corruption
        corrupt = random.random()
        if corrupt < __ERR_RATE:
            err_bytearr = bytearray(byte_msg)
            pos = random.randint(0, len(byte_msg)-1)
            val = err_bytearr[pos]
            if val > 1:
                err_bytearr[pos] -= 2
            else:
                err_bytearr[pos] = 254
            err_msg = bytes(err_bytearr)
            print("WARNING: udt_send: Packet corrupted in unreliable layer!!")
            return sockd.sendto(err_msg, peer_addr)
        else:
            return sockd.sendto(byte_msg, peer_addr)
```

You are required to use `__udt_send()` and `__udt_rcv()` for all message communications in the `rdt_send()`, `rdt_rcv()`, and `rdt_close()` functions.

Part 2 – Message format



Type	ACK = 11, DATA = 12	
Sequence no.	0 or 1	
Checksum	Use __IntChksum() to calculate the checksum for the whole packet (Header + Payload)	
Payload length	0 to 1000	Network-byte order

You learn how to assemble a header with binary data thru Workshop 3

Part 2 – Checksum Calculation

```
def __IntChksum(byte_msg):
```

```
    total = 0
    length = len(byte_msg)  #length of the byte message object
    i = 0
    while length > 1:
        total += ((byte_msg[i+1] << 8) & 0xFF00) + ((byte_msg[i]) & 0xFF)
        i += 2
        length -= 2

    if length > 0:
        total += (byte_msg[i] & 0xFF)

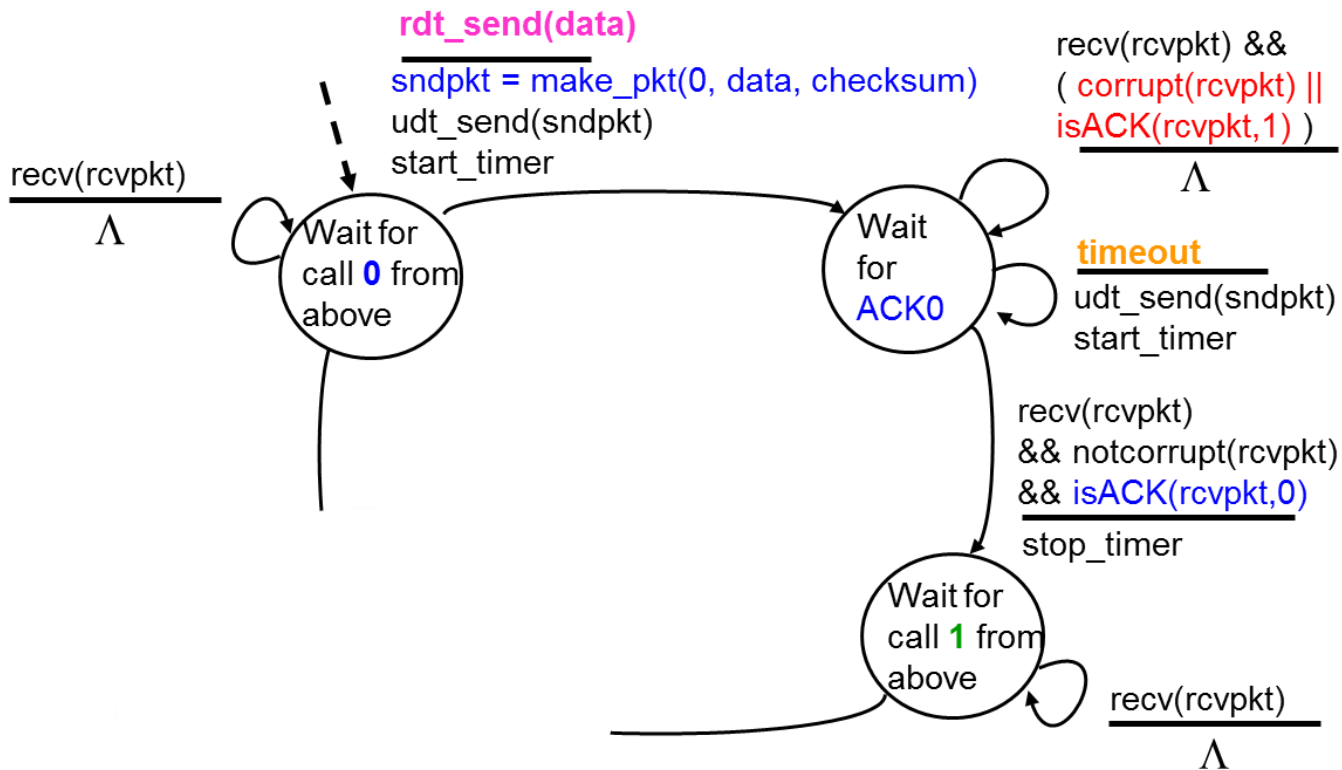
    while (total >> 16) > 0:
        total = (total & 0xFFFF) + (total >> 16)

    total = ~total
    return total & 0xFFFF
```

This function treats the whole message as a sequence of bytes and calculates the 16-bit checksum value.

This function is also being used at the receiving end to check whether the received message is unimpaired.

Part 2 – rdt_send()



Make the packet

- assemble the packet header
- set checksum field to zero
- copy application data to payload field
- calculate checksum for whole packet
- store checksum value in packet header

udt_send(packet)

do

wait for ACK or timeout

if is timeout

retransmit the packet

endif

take appro. action if packet is corrupted

if is ACK

check for correctness of ACK

else is DATA

take appropriate action

endif

repeat until received expected ACK

“wait for ACK or timeout”

How to do that?

Wait for ACK

- The process has to call `__udt_rcv()` to wait for incoming packet
- `__udt_rcv()` is a blocking call
- How can it return after waiting for a fixed duration?

Use socket timeout mode or `select.select()`

- We learn the technique thru Workshop 2

Part 2 – rdt_rcv()

```
recv(rcvpkt) && notcorrupt(rcvpkt)
&& has_seq0(rcvpkt)
sndpkt = make_pkt(ACK0, checksum)
udt_send(sndpkt)
extract(rcvpkt, data)
return data to rdt_rcv()
```

```
recv(rcvpkt) &&
(corrupt(rcvpkt) ||
has_seq1(rcvpkt))
sndpkt = make_pkt(ACK1,
checksum)
udt_send(sndpkt)
```



```
recv(rcvpkt) && notcorrupt(rcvpkt)
&& has_seq1(rcvpkt)
sndpkt = make_pkt(ACK1, checksum)
udt_send(sndpkt)
extract(rcvpkt, data)
return data to rdt_rcv()
```

do

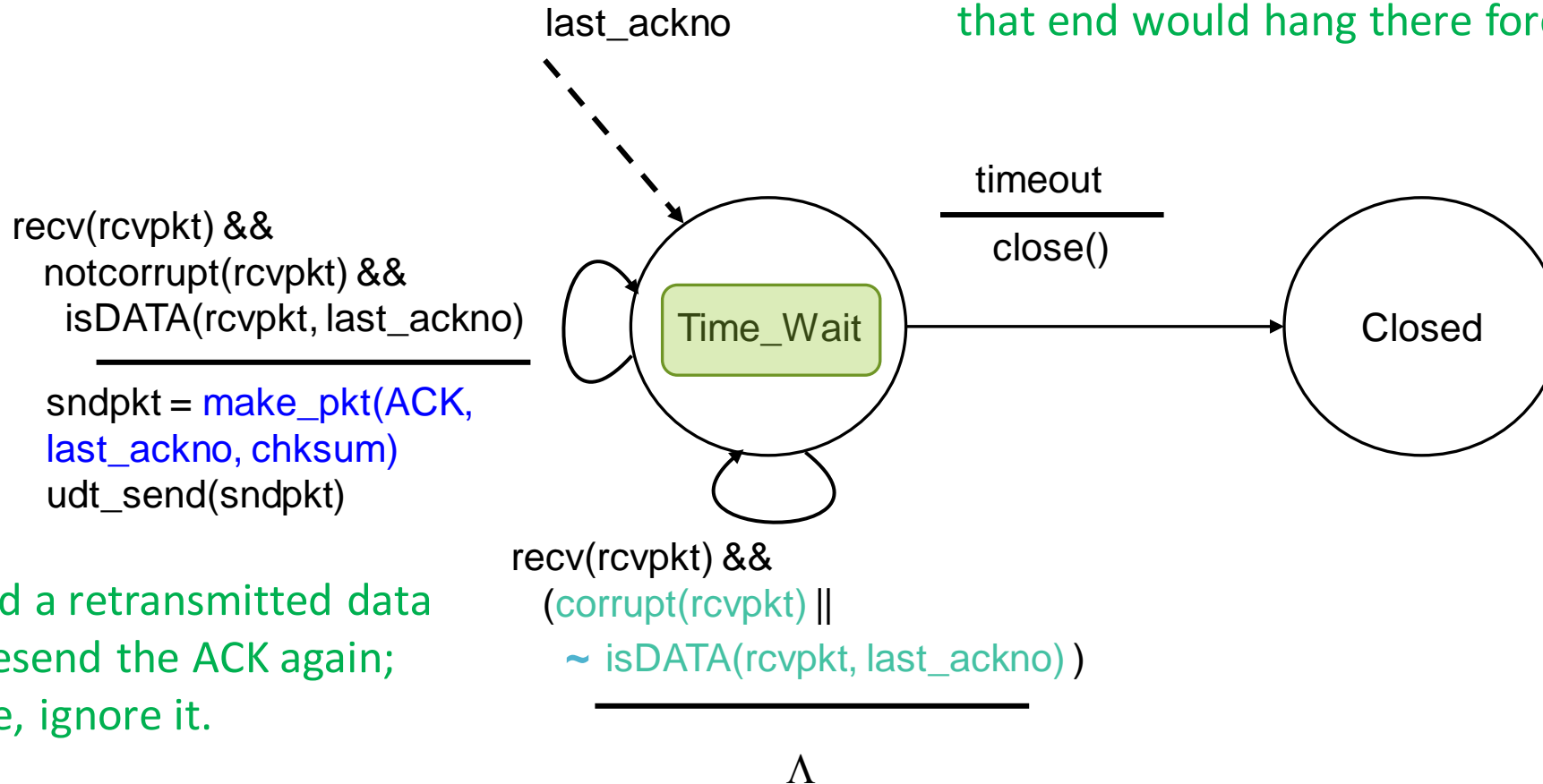
```
receive(packet)
take appro. action if packet is corrupted
if is DATA
    if is expected packet
        send ACK
        extract message and return it to upper layer
    else
        take appropriate action
endif
else is ACK
    take appropriate action
endif
```

repeat until received expected DATA

```
recv(rcvpkt) &&
(corrupt(rcvpkt) ||
has_seq0(rcvpkt))
sndpkt = make_pkt(ACK0,
checksum)
udt_send(sndpkt)
```

Part 2 – rdt_close()

The last ACK sent by this peer may be lost or corrupted; if this peer closes its socket and leaves, nobody is going to handle the retransmitted packet from the other end and that end would hang there forever !!!



If received a retransmitted data packet, resend the ACK again; otherwise, ignore it.

Implementation Requirements

Cannot use TCP

Require to simulate losses and errors by using the `__udt_send()` function to transmit all outgoing packets

Zero mark will be given if we find that

- Your implementation makes use of TCP
- Your implementation does not call `__udt_send()`

Part 2 – rdt3.py

Test

- Testing platform – any platform with python3
 - To run the server: `python3 test-server1.py localhost <<loss rate>> <<error rate>>`
 - To run the client: `python3 test-client1.py localhost <<filename>> <<loss rate>> <<error rate>>`

- Test cases

- small file (around 30 KB)
- large file (around 10 MB)
- different combinations of PACKET LOSS RATE and PACKET ERROR RATE

- Script files

- run-simulation2.bat, run-simulation2-OSX.sh, run-simulation2-Ubuntu.sh

PACKET LOSS RATE	PACKET ERROR RATE
0.0	0.0
0.2	0.0
0.0	0.2
0.2	0.2
0.3	0.3

Output Display

Very important – that helps you to check the correctness of your logic as well as to aid the debugging

Recommendation

- Generate an output statement whenever the RDT layer sends or receives a packet
- Generate an output statement to identify the type of packet and some control information
- Generate an output statement whenever the RDT layer detects or experiences an expected event or unexpected event or error situation

Please refer to [Part2-sample-output.pdf](#) for the sample output

*Submission Deadline:
May 3 17:00*

Part 3 – rdt4.py

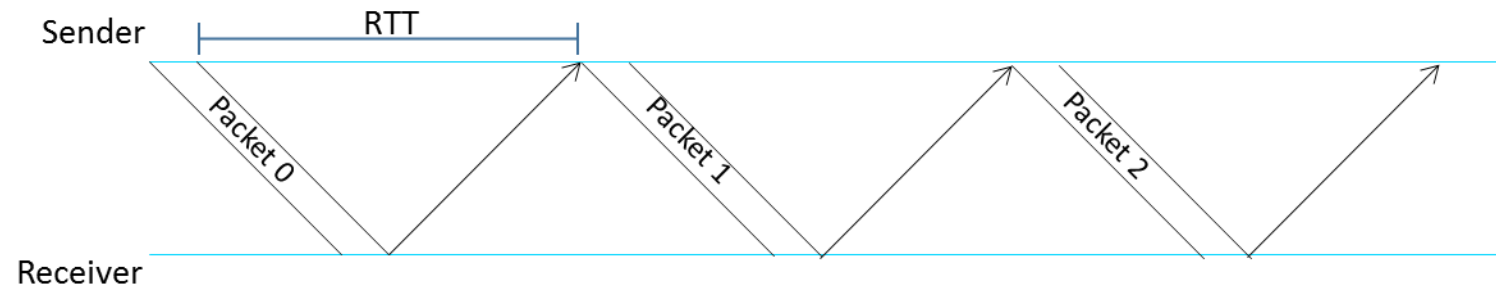
Download Part3-template.zip

Task

- Complete rdt4.py
 - For rdt_socket(), rdt_bind(), & rdt_peer(), you can reuse Part 2 implementation
 - Enhance rdt3.0 logic to include the Extended-Stop-and-Wait (rdt4.0) logic in
 - rdt_send(), rdt_recv()
 - For rdt_close()
 - The behavior of this function is the same as in Part 2 except that a peer may receive retransmitted packets with different sequence numbers within previous window

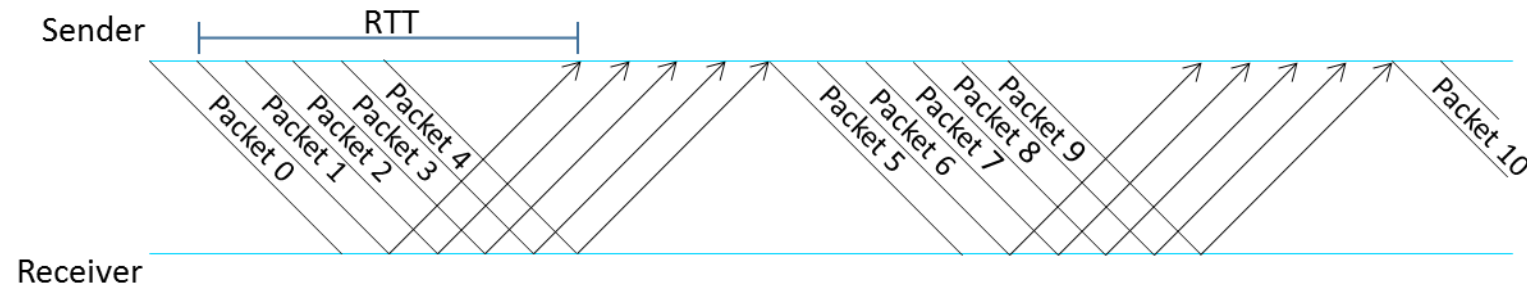
Extended-Stop-and-Wait (rdt4.0)

Stop-and-Wait (rdt3.0)

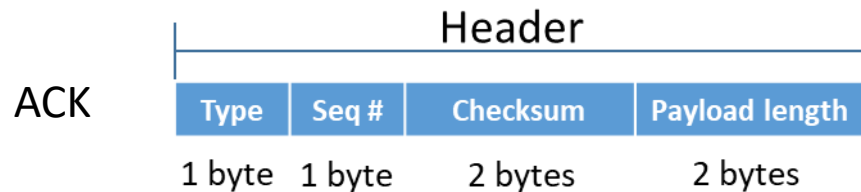


Extended-Stop-and-Wait (rdt4.0)

Window size $W = 5$



Part 3 – Message format (Same as Part 2)



Type	ACK = 11, DATA = 12
Sequence no.	0 to 255
Checksum	Use __IntChksum() to calculate the checksum for the whole packet (Header + Payload)
Payload length	0 to 1000

Part 3 – rdt_send()

Application process calls this function to transmit a message (up to a limit of $\text{PAYLOAD} \times W$ bytes) to targeted remote process

rdt_send(data)

```
N = count_pkt(data)
S = nextseqnum
for i = 1 to N {
    sndpkt[i] = make_pkt(nextseqnum,
                          data, checksum)
    udt_send(sndpkt[i])
    nextseqnum++
}
```

start_timer

nextseqnum=0

Wait for
call from
above

rcv(rcvpkt) &&
(corrupt(rcvpkt) ||
 $\sim \text{isACKbetween}(\text{rcvpkt}, S, S+N-1)$)

Λ

Wait for
ACKs

rcv(rcvpkt) &&
notcorrupt(rcvpkt) &&
 $\text{isACKbetween}(\text{rcvpkt}, S, S+N-2)$
k = getACKnum(rcvpkt)
set all sndpkt[] between S to k as acked

Cumulative acknowledgment

timeout

retransmit all unacked sndpkt[]
start_timer

rcv(rcvpkt)
&& notcorrupt(rcvpkt)
&& $\text{isACK}(\text{rcvpkt}, S+N-1)$
stop_timer

rcv(rcvpkt)
&& notcorrupt(rcvpkt)
&& $\text{isDATA}(\text{rcvpkt})$
drop the packet
take appropriate action if needed

Count no. of packets that will be generated
Compose and send all packets; each with unique sequence number
do

wait for ACK or timeout

if is timeout

retransmit all unACKed packets

endif

take appropriate action if packet is corrupted

if is ACK

check for correctness of ACK and take appro. action

else is DATA

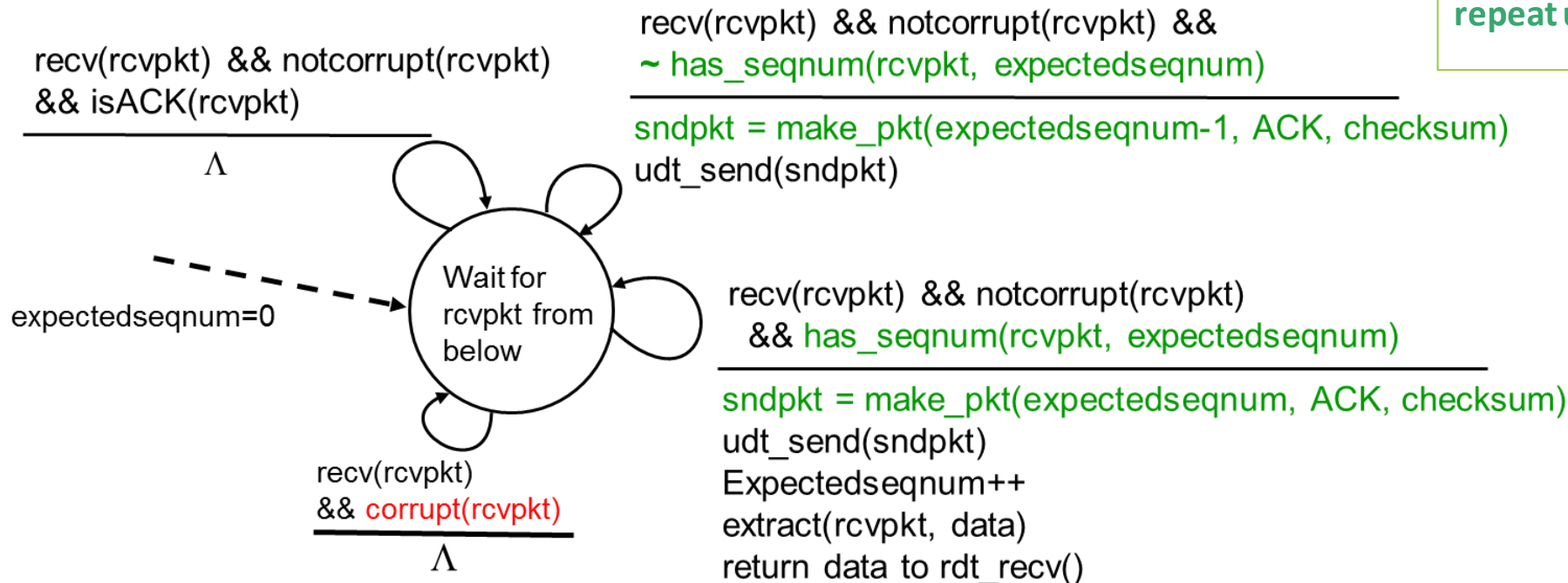
take appropriate action

endif

repeat until received all ACKs

Part 3 – rdt_rcv()

Please note that the receiver at RD layer does not know how many packets are coming. Thus, it simply accepts one packet at a time and passes it to upper layer.



do

```
receive(packet)
take appro. action if packet is corrupted
if is DATA
    if is expectedseqnum
        send ACK
        return message to upper layer
    else
        take appropriate action
    endif
else is ACK
    take appropriate action
endif
```

repeat until received the expected DATA

Part 3 – rdt4.py

Testing platform – any platform with python3

- To run the server: `python3 test-server3.py localhost <<loss rate>> <<error rate>> <<window size>>`
- To run the client: `python3 test-client3.py localhost <<filename>> <<loss rate>> <<error rate>> <<window size>>`

Test cases

- small file (around 30 KB)
- large file (around 10 MB)
- different combinations of W, LOSS_RATE and ERR_RATE

Script files

- run-simulation3.bat, run-simulation3-OSX.sh, run-simulation3-Ubuntu.sh

Please refer to Part3-sample-output.pdf for the sample output

W	PACKET LOSS RATE	PACKET ERROR RATE
1	0.0	0.0
5	0.0	0.0
9	0.0	0.0
1	0.1	0.1
5	0.1	0.1
9	0.1	0.1
1	0.3	0.3
5	0.3	0.3
9	0.3	0.3

Submissions

Part 2

- Deadline – 5:00pm, **April 7** (Wednesday)
- Submit file: rdt3.py

Part 3

- Deadline – 5:00pm, **May 3** (Monday)
- Submit file: rdt4.py

Late submission policy:

- At most 3 days with 10% penalty for each day of delay.

Grading Policy

<p>Part 2 (9 points)</p>	<ul style="list-style-type: none"> • The program can transfer data and terminate correctly in an environment without packet loss and corruption. [2.5/9] • The program can transfer data and terminate correctly in an environment with packet loss but no corruption ($0.0 < \text{LOSS} \leq 0.3$, $\text{ERROR} = 0.0$). [2/9] • The program can transfer data and terminate correctly in an environment with packet corruption but no loss ($\text{LOSS} = 0.0$, $0.0 < \text{ERROR} \leq 0.3$). [2/9] • The program can transfer data and terminate correctly in an environment with packet loss and corruption ($0.0 < \text{LOSS} \leq 0.3$, $0.0 < \text{ERROR} \leq 0.3$). [2/9] • Documentation [0.5/9] <ul style="list-style-type: none"> • Include necessary documentation to clearly indicate the logic of the program; include required student's info at the beginning of the program
<p>Part 3 (9 points)</p>	<ul style="list-style-type: none"> • The program can transfer data and terminate correctly with W=1 in an environment without packet loss and corruption [1/9] and in an environment with loss and corruption [1.5/9]. • The program can transfer data and terminate correctly with 1 < W ≤ 10 in an environment without packet loss and corruption [2/9] and in an environment with loss and corruption ($0.0 < \text{LOSS} \leq 0.3$, $0.0 < \text{ERR} \leq 0.3$) [4/9]. • Documentation [0.5/9] <ul style="list-style-type: none"> • Include necessary documentation to clearly indicate the logic of the program; include required student's info at the beginning of the program