

Hashing

- Idea: put elements into a fixed-size array by taking the remainder
 - $A = [1, 5, 6, 3, 16, 20]$, with size of hash table = 6
 - This gives hash table $H = [6, 1, 20, 3, 16, 5]$
- Okay in the above example, but what about $A = [6, 12, 1, 13, 2, 14]$? Now have multiple elements that hash to same thing: $6 \bmod 6 = 12 \bmod 6 = 0$.
- Solution: use a list to keep track of all elements that hash to same index, and instead just store a pointer to the list in H if needed:
 - $H = [[6, 12], [1, 13], [2, 14], -1, -1, -1]$ (fill unused slots with -1)

Hash function effectiveness

- Downside of simple dividing/remainder hash function: can be bad if data has lots of elements that have the same remainder mod (size of H):
 - $A = [1, 1, 1, 1, 1, 1, 1, 1]$
 - $H = [[1, 1, 1, 1, 1, 1, 1, 1], -1, -1, -1, -1, -1]$
 - This puts us back at the $O(n)$ time to find an element in a list (not good)
- Solution: use more complex hash function or methodology to distribute items more evenly. Good hash functions will expect a constant number of elements (so, $O(1)$ lookup time) in each bucket at most, not n elements.
- Result: **expected** $O(1)$ lookup time.