

Group 29 -
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CS 325 -- Section 400
Project 1 Report

Theoretical Run-time Analysis:

Algorithm 1: Enumeration

Theoretical runtime for algo1: $O(n^3)$

```
Function algo1(A[0...n-1]):  
    Initialize maxSum = 0  
    If n = 1:  
        maxSum = A[0]  
        Return maxSum  
    For i = 0 ... (n-1):  
        For j = i ... (n-1):  
            Initialize tempSum = 0  
            For k = i ... j:  
                tempSum += A[k]  
                If tempSum > maxSum:  
                    maxSum = tempSum  
    Return maxSum
```

Algorithm 2: Better Enumeration

Theoretical runtime for algo2: $O(n^2)$

```
Function algo2(A[0...n-1]):  
    Initialize maxSum = 0  
    If n = 1:  
        maxSum = A[0]  
        Return maxSum  
    For i = 0 ... (n-1):  
        Initialize tempSum = A[i]  
        For j = (i+1) ... (n-1):  
            tempSum += A[j]  
            If tempSum > maxSum:  
                maxSum = tempSum  
    Return maxSum
```

Algorithm 3: Divide and Conquer

Theoretical runtime for algo3: $O(n \lg n)$

```
Function algo3(A[0...n-1]):  
    Initialize maxSum = minimum int
```

```

If n = 1:
    maxSum = A[0]
    Return maxSum
Initialize mid = n / 2
Left[] = A[0...mid-1]
Right[] = A[mid...n-1]

//Recursive calls for cases 1 and 2
left_maxSum = algo3(left)
right_maxSum = algo3(right)
Lr_maxSum = max(left_maxSum, right_maxSum)

Initialize tempSum = 0
Initialize l_sum = r_sum = minimum int

//For case 3:
For i = (mid-1) ... 0: //Loop backwards through left half of A
    tempSum += A[i]
    If tempSum > l_sum:
        l_sum = tempSum
For i = mid ... (n-1):
    tempSum += A[i]
    If tempSum > r_sum:
        r_sum = tempSum

//Determining which case to use
if Lr_maxSum >= l_sum + r_sum:
    //Case 1 or 2
    maxSum = Lr_maxSum
Else:
    //Case 3
    maxSum = l_sum + r_sum

Return maxSum

```

Algorithm 4 Pseudocode:

Theoretical Run-Time Analysis for Algorithm 4: $O(n)$

```

For i=0 to size-1
    (from the last iteration) Make subarray arr1 or max sum that ends with index i-1

    Find subarray arr2 or max sum that ends with index i. This will be arr1, or the
    subarray of just i.

```

(from the last iteration) Make subarray arr3 of max sum be between indexes 0 and i-1.

Find subarray of maxsum between indexes 0 and i. This will be arr2 or arr3.

Testing:

Testing the validity of the algorithms was done using the provided test set contained in MSS_TestProblems.txt. Arrays from the test set were fed to the algorithms and their output was compared with the known results contained in MSS_TestResults.txt. For each algorithm, their output matched perfectly with already known results demonstrating the correctness of the algorithms.

Experimental Analysis:

1. Average running times:

n	Algo1 Average Runtime (microseconds)
100	2698
200	16888
300	35834
400	71381
500	138665
600	238797
700	376276
800	550983
900	795115
1000	1086242

n	Algo2 Average Runtime (microseconds)
100	46
200	183
500	1133

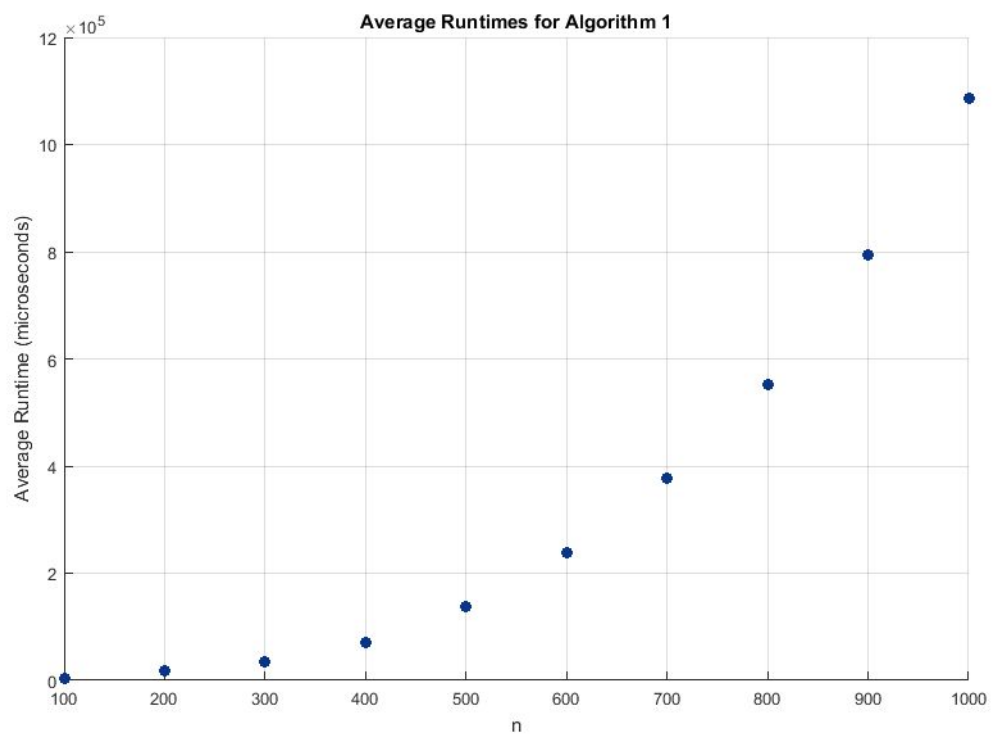
1000	4530
2000	25319
5000	117379
10000	454838
20000	1864108
50000	11750929
100000	46634559

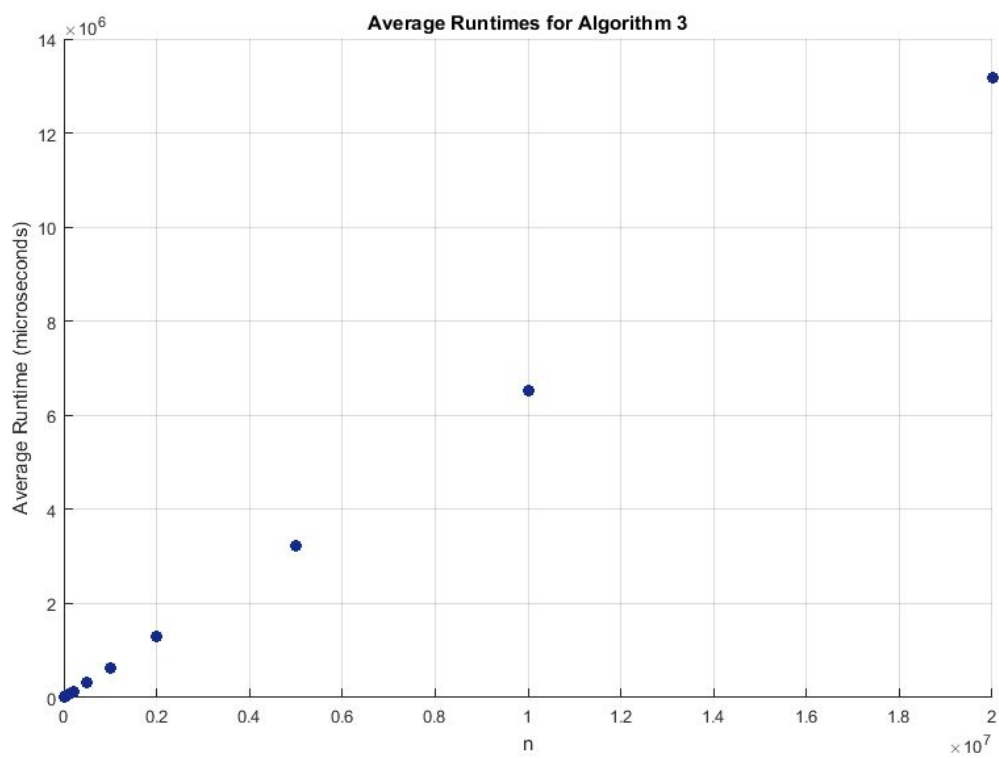
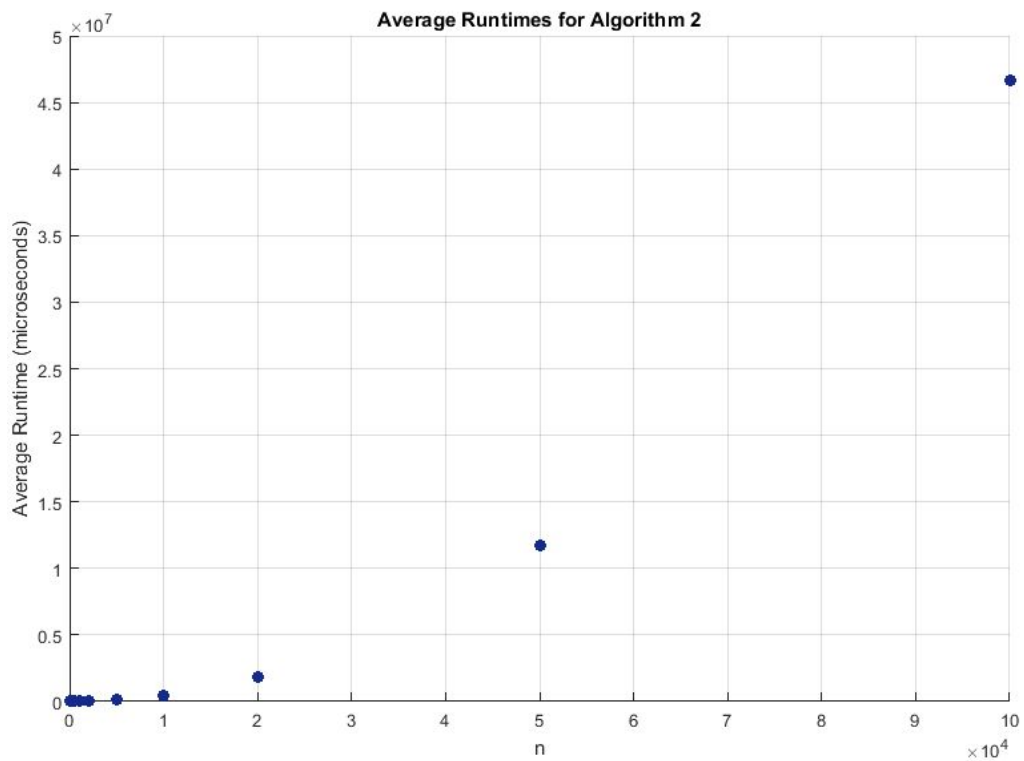
n	Algo3 Average Runtime (microseconds)
20000	22574
50000	35207
100000	72930
200000	123246
500000	312900
1000000	631114
2000000	1289305
5000000	3215504
10000000	6522413
20000000	13179061

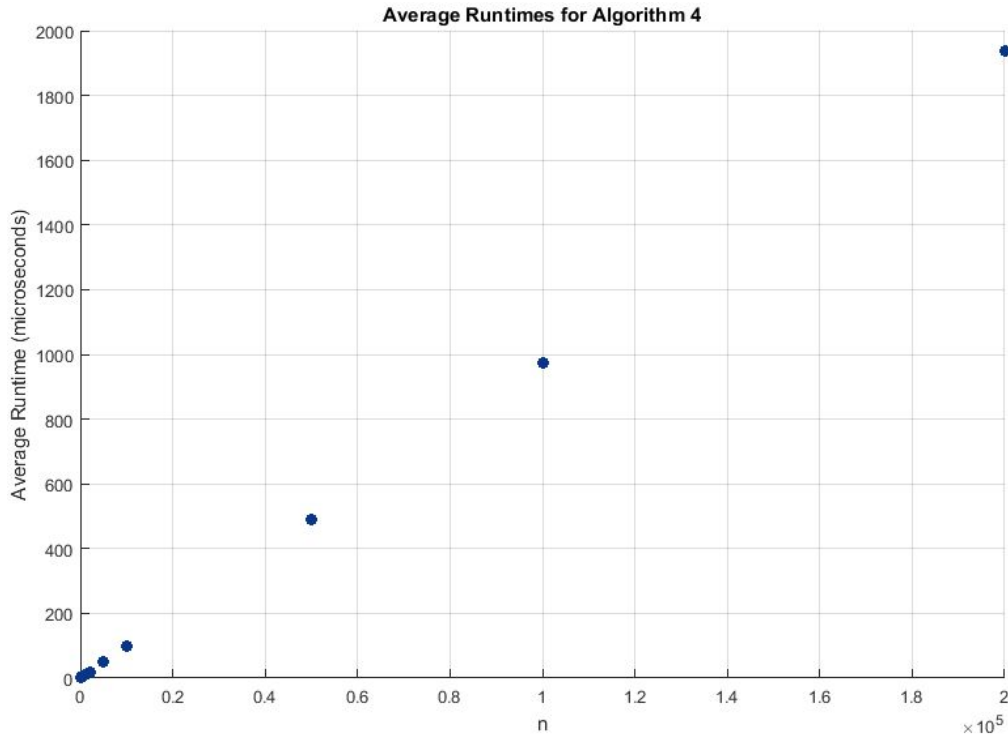
n	Algo4 Average Runtime (microseconds)
100	1
200	2
500	5
1000	10

2000	20
5000	50
10000	98
50000	488
100000	973
200000	1936

2. Running time plots:







3. Regression models:

- Algorithm 1:
 - $T_1(n) = 0.001152n^3 - 0.1136n^2 + 49.57n - 714.9$
 - $R^2 = 0.9999$
- Algorithm 2:
 - $T_2(n) = 0.004639n^2 + 2.628n - 8624$
 - $R^2 = 1.0000$
- Algorithm 3:
 - $T_3(n) = 0.02718n \cdot \lg(n) + 81910$
 - $R^2 = 0.9997$
- Algorithm 4:
 - $T_4(n) = 0.009686n + 1.073$
 - $R^2 = 1.0000$

4. Discrepancies:

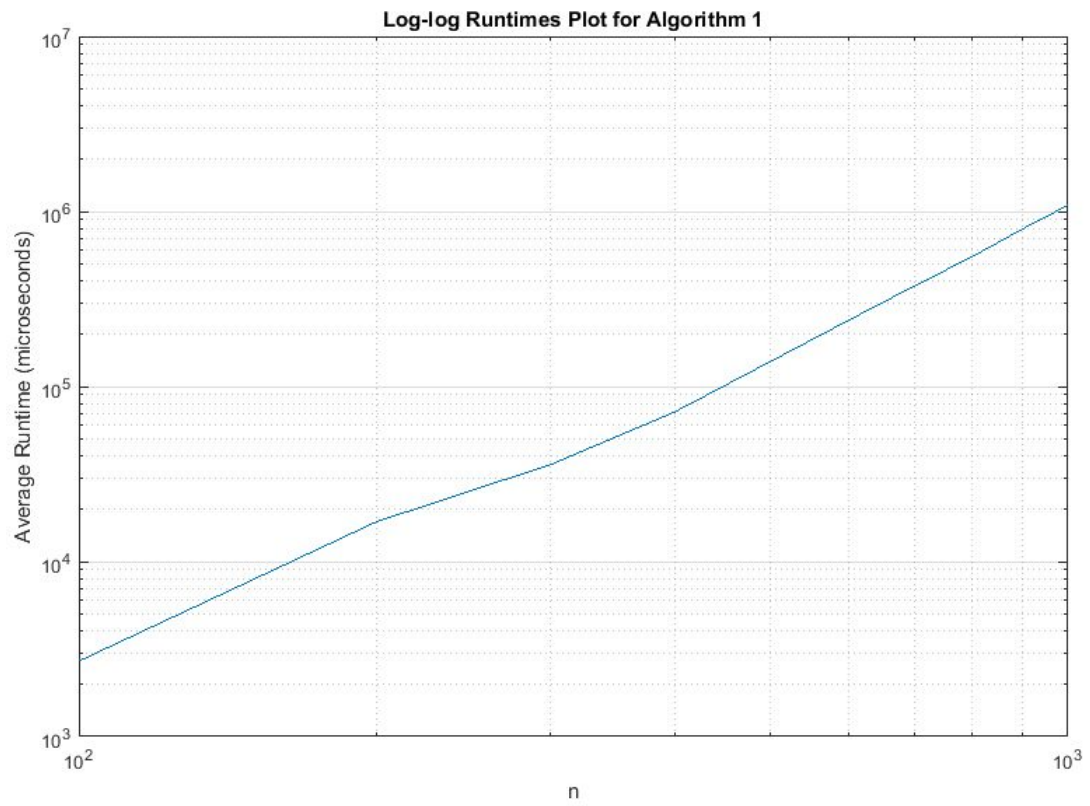
- **Algorithm 1:** For algorithm 1 there were not any apparent discrepancies between expected runtime and our results. We knew that it would be unable to efficiently run values of n larger than 10000 given its cubic runtime complexity. Our runtime test program would hang for several hours with values greater than 20000 making it impractical to test beyond n values of 1000. But for the values of n we did run the regression had a R^2 of .9999 indicating a good fit.
- **Algorithm 2:** Algorithm 2 did not have any discrepancies between expected runtime and our results. We knew it would be roughly $O(n^2)$ for the complexity the results of the regression reflect this as well for R^2 is equal to 1.00.

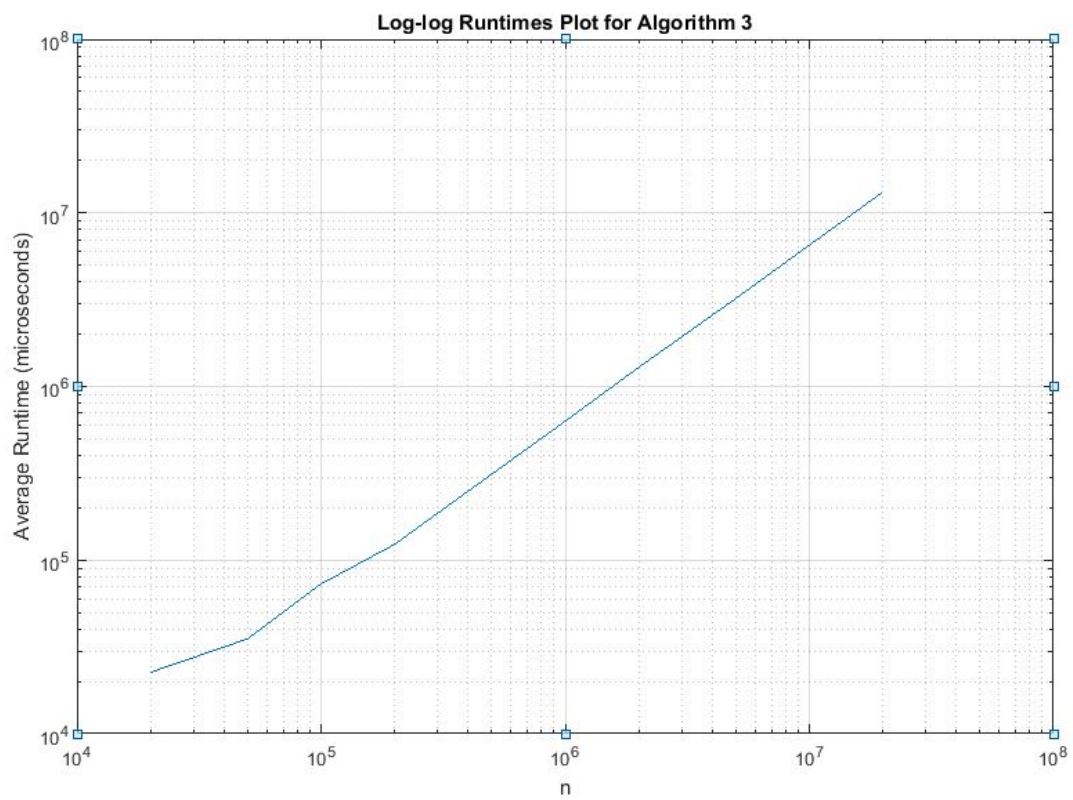
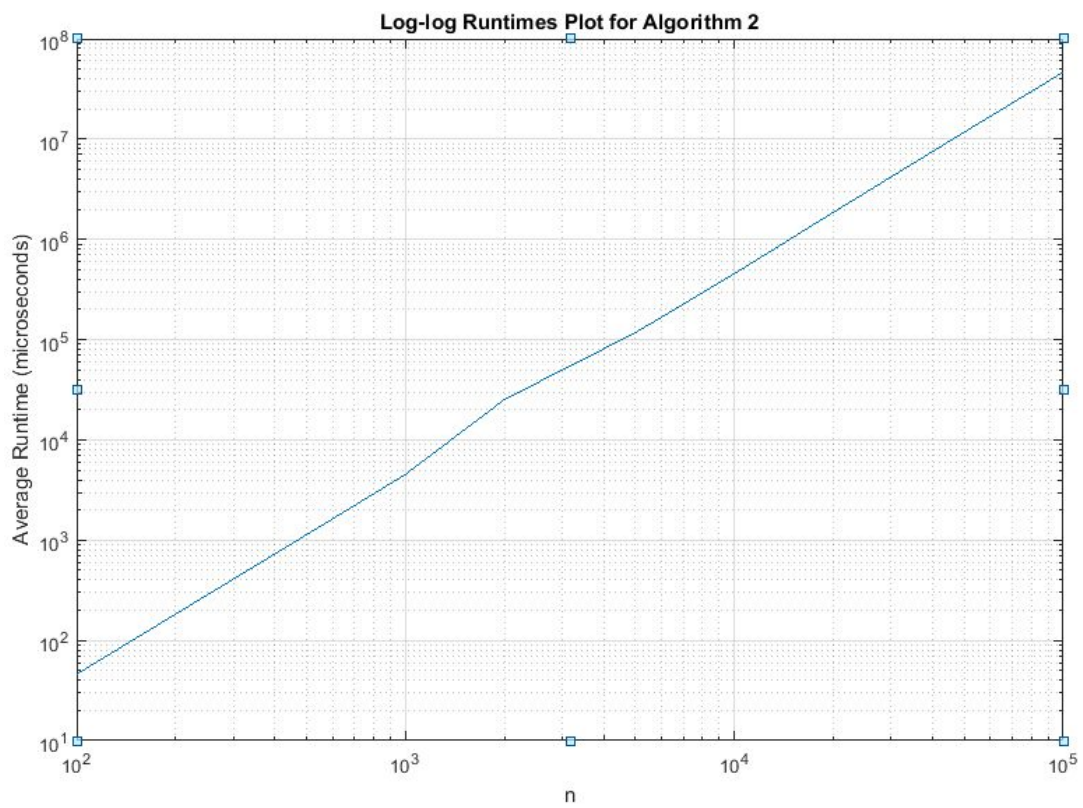
- **Algorithm 3:** Algorithm 3 did appear to have some mild discrepancies relative to the expected runtime, specifically it was expected to have a $O(n \lg n)$ runtime due to its divide and conquer approach for finding the maximum subarray, but had a runtime more closely resembling $O(n)$ when inspecting the plot. The regression for Algorithm 3 using $f(n) = a \cdot n \lg n + c$ as a model had an R^2 value of .9997 indicating that its runtime was in fact $f(n) = O(n \lg n)$, despite appearing as a straight line. This is okay because, Algorithm 3 is still bound by $O(n^2)$ from above and $\Omega(n)$ from below. The runtimes would likely begin to fit an $n \lg n$ curve even better at higher n , but it was not feasible to test them with any higher values of n than we already did.
- **Algorithm 4:** Algorithm 4 was expected to have a linear runtime and the regression analysis demonstrates that it does indeed have a linear time with the regression having an R^2 of 1.00.

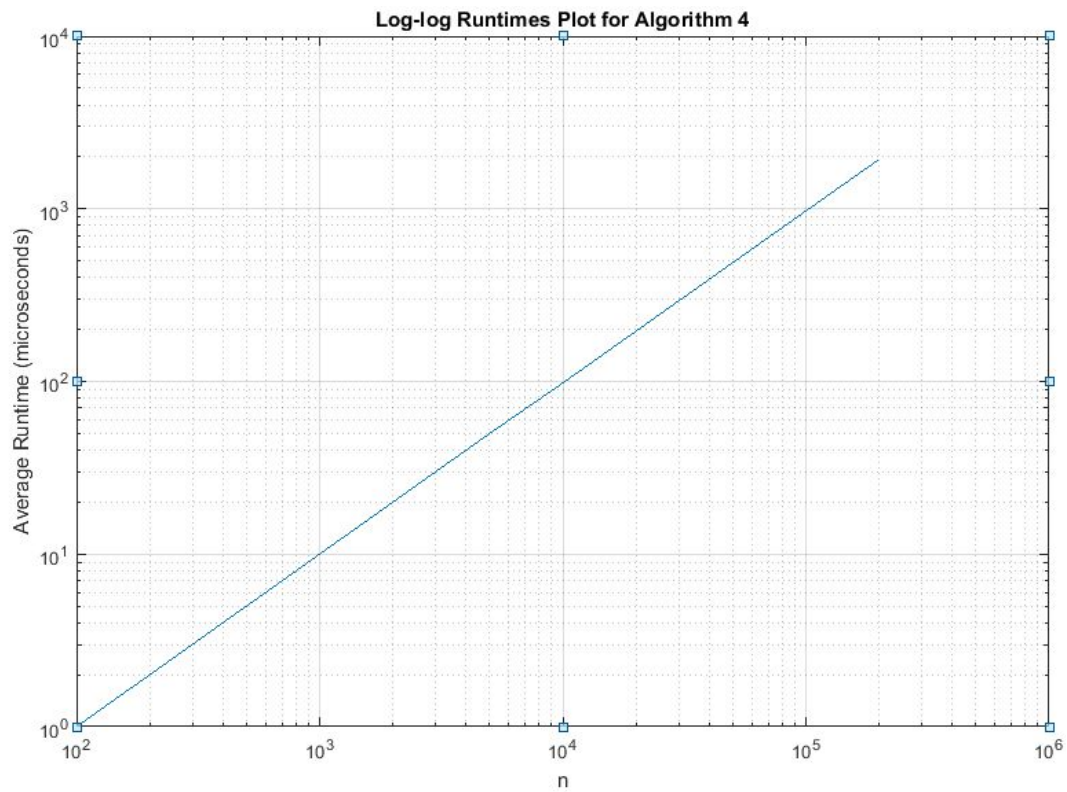
5. Largest input in given time:

	Largest value of n possible for Algorithm...			
Time (microseconds)	1	2	3	4
5000000	1655	32570	7897091	516208850
10000000	2081	46160	15289665	1032417811
60000000	3763	113447	83758161	6194507425

6. Log-log plots:







7. Combined plot:

