$\ensuremath{\mathsf{TDT4165}}$ - Programming Languages

Assignment 3: Higher-Order Programming

Benjamin Zubača

October 2, 2024

a)

```
proc {QuadraticEquation A B C}
1
        RealSol
2
        Х1
 3
4
        Discriminant = B*B - 4.0*A*C
5
    in
6
        if Discriminant >= 0.0 then
7
           RealSol = true
           X1 = (~B + {Float.sqrt Discriminant}) / (2.0*A)
 9
           X2 = (~B - {Float.sqrt Discriminant}) / (2.0*A)
10
        else
11
           RealSol = false
12
           X1 = X2 = noSol
13
        end
14
    end
```

The procedure {QuadraticEquation A B C ?RealSol ?X1 ?X2} takes three parameters A, B, and C, which correspond to the coefficients in the quadratic equation $Ax^2 + Bx + C = 0$. It calculates the discriminant $B^2 - 4AC$ to determine the nature of the solutions. If the discriminant is non-negative, it binds RealSol to true and computes real roots X1 and X2 using the quadratic formula. Otherwise, it binds RealSol to false and both X1 and X2 to the atom noSol, indicating that no real solutions exist.

```
RealSol = true
X1 = 0.5
X2 = ~0.5
When A = 2, B = 2, and C = 2, then:
```

b) When A = 2, B = 2, and $C = ^1$, then:

- RealSol = false
- X1 = noSol
- X2 = noSol
- c) Procedural abstractions are useful because they allow you to encapsulate a sequence of related statements into a reusable procedure. By abstracting these operations into a single unit, the procedure can be invoked with different arguments, simplifying code, enhancing clarity, and facilitating code reuse.
- d) A procedure is a syntactic construct that executes a series of statements, performing operations without producing a return value. A function, on the other hand, executes statements (if applicable) and returns the value of an expression based on its inputs.

```
fun {Sum List}
case List

of nil then 0

[] Head|Tail then Head + {Sum Tail}
end
end
```

Task 3

a)

```
fun {RightFold List Op U}
case List
of nil then U
{
    [] Head|Tail then {Op Head {RightFold Tail Op U}}
end
end
```

b) Line 1 defines the function RightFold with the three parameters List, Op, and U. Line 2 starts a case expression to pattern match on the input List. It checks the structure of the List to decide the next steps. Line 3 handles the base case: If List is empty, it returns U, which is the neutral element. Line 4 handles the recursive case: If List has Head element followed by Tail list, it applies Op to Head and the result of recursively folding the Tail using RightFold. Lines 5 and 6 indicate the end of the case expression and the RightFold function, respectively.

c)

```
fun {Length List}
{RightFold List fun {$ _ Y} Y + 1 end 0}
end
```

```
fun {Sum List}
{RightFold List fun {$ X Y} X + Y end 0}
end
```

- d) LeftFold and RightFold would not give different results for the Sum and Length operations, as they are both associative and commutative.
- e) For implementing the product of list elements using RightFold, the appropriate value for U would be 1. This is because the neutral element for multiplication is 1, as any number multiplied by 1 remains unchanged.

```
fun {Quadratic A B C}
fun {$ X}
A*X*X + B*X + C
end
end
```

The function {Quadratic A B C} takes three parameters corresponding to the coefficients of the quadratic polynomial $Ax^2 + Bx + C$. It returns an anonymous function that accepts one parameter X and calculates the polynomial's value, using the coefficients that are already bound from the outer Quadratic function.

Task 5

a)

```
fun {LazyNumberGenerator StartValue}
[StartValue {LazyNumberGenerator StartValue + 1}]
end
```

b) The {LazyNumberGenerator StartValue} function implements a lazy infinite sequence of integers. It takes a starting value StartValue as an argument and returns a list with two elements: The head, which is StartValue (the current value), and the tail, which is a lazy, unevaluated call to generate the next value in the sequence. This lazy call increments the current value by 1 for the subsequent iteration.

The solution has several limitations worth noting. Firstly, while it is lazy, accessing many elements could create a long chain of unevaluated functions, potentially consuming significant memory. Depending on the Oz environment and system resources, deep recursion may lead to stack overflow. Additionally, there is no way to terminate the sequence once the function is executed, which could pose issues in certain scenarios.

a)

```
fun {Sum List}
1
        fun {DoSum List Acc}
2
           case List
3
           of nil then Acc
4
           [] Head | Tail then {DoSum Tail Acc}
        end
7
8
     in
        {DoSum List 0}
9
     end
10
```

To make the Sum function tail recursive, we can introduce a nested helper function DoSum that takes an additional accumulator Acc parameter to store the running sum of the List. We can then modify the base case to return the accumulator instead of 0, and change the recursive call to be the last operation in the function, passing the updated accumulator. The main Sum function now calls DoSum with the initial list and an accumulator value of 0.

- b) Tail recursion in Oz offers significant memory efficiency, as the compiler can optimize tail recursive functions to use constant stack space, regardless of input size. This optimization replaces function calls with jump instructions, improving performance by eliminating the overhead of stack frame allocations. Furthermore, it encourages a functional programming style, promoting immutability and a more declarative approach to writing functions.
- c) Not all programming languages that support recursion benefit equally from tail recursion. Languages like Clojure, Elixir, Erlang, and Haskell, which have built-in tail call optimization¹, gain significant advantages from tail-recursive functions. In contrast, languages such as C, Go, and Python, which lack guaranteed tail call optimization, do not experience the same benefits.

¹https://en.wikipedia.org/wiki/Tail_call#Language_support