

1 Definitions

Definition 1.0.1. A *pre-log structure* on X is a sheaf of commutative monoids \mathcal{M} on $X_{\text{ét}}$ along with a morphism of sheaves of monoids,

$$\exp : \mathcal{M} \rightarrow (\mathcal{O}_X, \times)$$

A morphism of pre-log structures $\alpha : (\mathcal{M}_1, \exp_1) \rightarrow (\mathcal{M}_2, \exp_2)$ is a morphism of sheaves of monoids $\alpha : \mathcal{M}_1 \rightarrow \mathcal{M}_2$ such that the diagram,

$$\begin{array}{ccc} \mathcal{M}_1 & \xrightarrow{\alpha} & \mathcal{M}_2 \\ \exp_1 \searrow & & \swarrow \exp_2 \\ & \mathcal{O}_X & \end{array}$$

commutes.

Remark. From now on, when we write \mathcal{O}_X in the category of monoids we mean (\mathcal{O}_X, \times) .

Example 1.0.2. Some examples,

- (a) $\mathcal{O}_X^\times \hookrightarrow \mathcal{O}_X$
- (b) $\mathcal{O}_X \rightarrow \mathcal{O}_X$

Example 1.0.3. Suppose that P is a commutative monoid and $\exp : P \rightarrow \Gamma(W, \mathcal{O}_W)$ is a map then we get a pre-log structure $\exp : \underline{P} \rightarrow \mathcal{O}_W$ by adjunction.

For a monoid P let $X_P = \text{Spec}(k[P])$ then we see,

$$\text{Hom}(W, X_P) = \{\text{pre-log structures } \underline{P} \rightarrow \mathcal{O}_W\}$$

Definition 1.0.4. A pre-log structure $\exp : \mathcal{M} \rightarrow \mathcal{O}_X$ is called a *log structure* if $\exp^{-1}(\mathcal{O}_X^\times) \rightarrow \mathcal{O}_X^\times$ is an isomorphism.

Remark. In particular, there is a unique morphism $\alpha : \mathcal{O}_X^\times \rightarrow \mathcal{M}$ making the diagram,

$$\begin{array}{ccc} \mathcal{O}_X^\times & \xrightarrow{\alpha} & \mathcal{M} \\ & \searrow & \downarrow \exp \\ & & \mathcal{O}_X \end{array}$$

commute so all log-structures “lie between” \mathcal{O}_X^\times and \mathcal{O}_X in the sense that \mathcal{O}_X^\times is the initial object and \mathcal{O}_X the terminal object of the categories of log structures.

Definition 1.0.5. A *log scheme* (X, \mathcal{M}) is a scheme X equipped with a log structure \mathcal{M} .

Proposition 1.0.6. Let $Z \subset X$ be a closed subset. Let $\mathcal{M}_Z \subset \mathcal{O}_X$ be the subsheaf of functions invertible on $U = X \setminus Z$. Then \mathcal{M}_Z is a log structure.

Proof. This just says that if $f \in \mathcal{O}_X(V)$ then it is invertible on U which is obvious by restriction. \square

1.1 Logification

Proposition 1.1.1. There is a left-adjoint $\mathcal{M} \mapsto \mathcal{M}^{\log}$ to the forgetful functor $\{\log_X\} \rightarrow \{\text{pre-log}_X\}$ given by,

$$\begin{array}{ccc}
 \exp^{-1}(\mathcal{O}_X^\times) & \hookrightarrow & \mathcal{M} \\
 \downarrow & \lrcorner & \downarrow \\
 \mathcal{O}_X^\times & \longrightarrow & \mathcal{M}^{\log} \\
 & \searrow & \swarrow \text{exp}' \\
 & & \mathcal{O}_X
 \end{array}
 \quad \begin{array}{c}
 \text{exp} \\
 \text{exp}'
 \end{array}$$

Proof. (DO THIS!!!) □

Definition 1.1.2. Given a morphism $\pi : X \rightarrow Y$ and a pre-log structure $\text{exp} : \mathcal{M} \rightarrow \mathcal{O}_Y$ on Y . Then the pullback is $\pi^{-1}\mathcal{M} \rightarrow \pi^{-1}\mathcal{O}_Y \rightarrow \mathcal{O}_X$. If \mathcal{M} is a log structure on Y then the pullback is,

$$\pi^b\mathcal{M} = (\pi^{-1}\mathcal{M})^{\log}$$

By the universal property there is a map,

$$\begin{array}{ccc}
 \pi^{-1}\mathcal{M} & \longrightarrow & \mathcal{O}_X \\
 \searrow & & \nearrow \\
 & \pi^b\mathcal{M} &
 \end{array}$$

Remark. There is a tautological log structure $\underline{P}^{\log} \rightarrow \mathcal{M}_{X_P}$. In the toric case $P = \sigma^\vee \cap M$ we can define $Z = X_\sigma \setminus T$ where $T = \text{Spec}(k[M])$ is the torus then I claim that $\underline{P}^{\log} = \mathcal{M}_Z$. (PROVE THIS!!)

Remark. Consider a morphism $W \rightarrow X_P$ then the log structure $\underline{P}^{\log} \rightarrow \mathcal{O}_W$ is $\pi^b\mathcal{M}_P$. (PROVE TIS!!)

Proposition 1.1.3. Let $\pi : X \rightarrow Y$ be a morphism. Then the following diagram commutes,

$$\begin{array}{ccc}
 \{\text{pre-log}_Y\} & \xrightarrow{\pi^{-1}} & \{\text{pre-log}_X\} \\
 \downarrow \log & & \downarrow \log \\
 \{\log_Y\} & \xrightarrow{\pi^b} & \{\log_X\}
 \end{array}$$

Proof. IS THIS ACTUALLY TRUE?? □

Definition 1.1.4. A morphism $f : (X, \mathcal{M}) \rightarrow (Y, \mathcal{N})$ of *log schemes* is a morphism of schemes $f : X \rightarrow Y$ and a morphism of sheaves of monoids α such that,

$$\begin{array}{ccc}
 \mathcal{N} & \xrightarrow{\alpha} & \pi_*\mathcal{M} \\
 \downarrow & & \downarrow \\
 \mathcal{O}_Y & \longrightarrow & \pi_*\mathcal{O}_X
 \end{array}$$

commutes. Equivalently this is a morphism of log structures $\alpha : \pi^b\mathcal{N} \rightarrow \mathcal{M}$ because then automatically the diagram,

$$\begin{array}{ccccc}
\pi^{-1}\mathcal{N} & \longrightarrow & \pi^{\flat}\mathcal{N} & \longrightarrow & \mathcal{M} \\
\downarrow & & & & \downarrow \\
\pi^{-1}\mathcal{O}_Y & \longrightarrow & & \longrightarrow & \mathcal{O}_X
\end{array}$$

commutes.