

July 21, **2007** 

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Null @ Root #44u61l5f

# Agenda



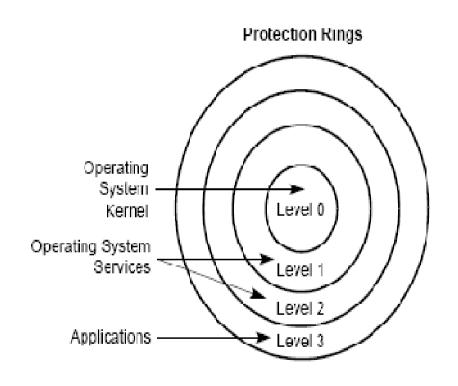
- 1) What is a kernel hook used for?
- 2) Get into the Ring0!
- 3) Review
- 4) Hook the planet!
- 5) Hook me, If you can!
- 6) ?

# What is a kernel hook used for?



- 1) Kernel hooks are global (relatively speaking).
- 2) Rootkit / Protection / Detection software are both in Ring Zero.
- 3) By using kernel hook, We can study the behavior of the system.
- 4) By using kernel hook, we can get performance data for specific tasks and generating statics.





### How to make code run on Ring0?

- Call gate
- : CPL change 3 from to 0
- Software Interrupt
- : API Call, Exception, ..
- Device Driver
- : It doesn't mean Usermode driver but Kernelmode drvier, because Kernelmode driver run on Level0, however, Usermode driver run on Level3.



☐ Code of exe file is located into Paged Memory, In contrast, Code of sys file is located into NonPaged Memory.

☐ Both exe file and sys file has PE(Portable Executable) file format.

#### **Header of 1stDriver.sys:**

		J		
EntryPoint:	00001000	Subsystem:	0001	
ImageBase:	00010000	NumberOfSections:	0006	
SizeOfImage:	00007000	TimeDateStamp:	466BAB7E	
BaseOfCode:	00001000	SizeOfHeaders:	00000400	? +
BaseOfData:	00002000	Characteristics:	010E	
SectionAlignment:	00001000	Checksum:	0000685F	?
FileAlignment:	00000200	SizeOfOptionalHeader:	00E0	
Magic:	010B	NumOfRvaAndSizes:	00000010	+ -

### **Header of cmd.exe:**

EntryPoint:	00005056	Subsystem:	0003
ImageBase:	4AD00000	NumberOfSections:	0003
SizeOflmage:	00078000	TimeDateStamp:	41107EBE
BaseOfCode:	00001000	SizeOfHeaders:	00000400 ? +
BaseOfData:	0001F000	Characteristics:	010F
SectionAlignment:	00001000	Checksum:	0007DCF6 ?
FileAlignment:	00000200	SizeOfOptionalHeader:	00E0
Magic:	010B	NumOfRvaAndSizes:	00000010 + -

☐ DPL of CS in sys file has value 0, however, DPL of CS in exe file has value 3. It means CS of sys file can access all of instructions and memorys, In contrast, CS of exe file has limited access right.

### www.CodeEnan.com

# Get into the Ringo!

) {

### 1. Compile C code by *DDK*

IN PDRIVER\_OBJECT pDriverObject, IN PUNICODE STRING pRegistryPath

NTSTATUS DriverEntru (

```
NTSTATUS status:
   int i:
   PDEVICE OBJECT Device Object = NULL:
   UNICODE STRING Device Name;
   UNICODE STRING Win32NameString;
   UNICODE STRING uFileName.uFileName2;
   ANSI STRING logNameString.logNameString2:
   OBJECT ATTRIBUTES obj attrib.obj attrib2:
   IO STATUS BLOCK file status file status2;
   IO STATUS BLOCK io status, io status2;
   LARGE INTEGER timeout:
   RtlInitUnicodeString(&Device Name, WIN DEVICE NA
   status = IoCreateDevice(pDriverObject,
                          66,878 KFCL .zip
883-82-84
            10:23
                            267 MAKEFILE
007-01-31
            88:22
                    (DIR)
                                obj
887-81-31
             88:22
                                objohk
997-01-31
            BB:48
                    (RID)
                                objfre
                             67 SOURCES
                            115.496 HOE
           5개 디렉터리 5,756,487,888 바이트 남음
G:Wsaruengang>build -Gez
BUILD: Object rest set to: ==> objchk
BUILD: Adding /Y to COPYCHD so xcopy eps won't hang.
BUILD: /i switch ignored
BUILD: Compile and Link for 1386
BUILD: Examining c:Warmengang directory for files to compile.
BUILD: Compiling c: Warmengang directory
ompiling - kfci.c for 1386
BUILD: Linking c: Maruengang directory
inking Executable - objohk#1386Wkfci.sys for 1386
BUILD: Done
  2 files compiled
  1 executable built
```

### 2. Assemble Asm code by KmdKit

```
DriverEntry proc pDriverObject:PDRIVER_OBJECT, pusRegi
  invoke DbgPrint, $CTAO("\nFindShadowTable: Entering
  mov eax, KeServiceDescriptorTable
  mov eax, [eax]
  invoke DbgPrint, $CTAO("FindShadowTable: ServiceDesc
  invoke GetServiceDescriptorTableShadowAddress
  .if eax != NULL
   invoke DbgPrint, $CTAO("FindShadowTable: ServiceDe
  .endif
  invoke DbgPrint, $CTAO("FindShadowTable: Leaving Dri
```

```
C:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#KndKit#examples#simple#GetKernelBase2set drv-GetKernelBase
G:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#KndKit#examples#simple#GetKernelBase2set drv-GetKernelBase
G:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#KndKit#examples#simple#GetKernelBase2#simple#GetKernelBase.bat
Assenbling: GetKernelBase.bat
G:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#kndKit#examples#simple#GetKernelBase.by
G:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#KndKit#examples#simple#GetKernelBase.sys /subsystem:native GetKernelBase.obj
C:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#XndKit#examples#simple#GetKernelBase2del GetKernelBase.obj
C:#Bocuments and Settings#구사무엘#바탕 화면#KndKit#XndKit#examples#simple#GetKernelBase2del GetKernelBase.obj
```



You can link DDK to Visual Studio.

If you use Visual Studio.net as IDE, You can use this way.

- 1. Download ddkbuild.bat from <a href="http://www.hollistech.com/Resources/ddkbuild/ddkbuild.htm">http://www.hollistech.com/Resources/ddkbuild/ddkbuild.htm</a>
- 2. Copied it at c:₩progra~₩microsoft visu~₩vc7₩bin₩
- 3. Edit ddkbuild.bat set WNETBASE= C:₩WinDDK (Directory that DDK installed)
- 4. Making Project
  - 1. Select makefile project.
  - 2. Build cmdline: ddkbuild -WNET checked.
  - 3. output / arrangement cmd skip
  - 4. Build cmdline: ddkbuild -WNET chekced. -cZ



It`s

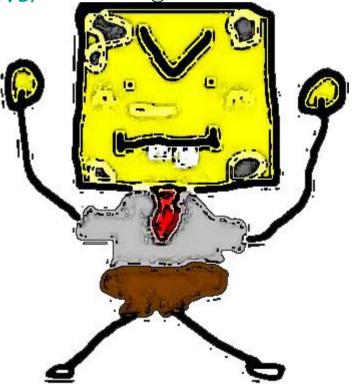
If you use Visual Studio 6.0 as IDE, You can use this way

So easy!!

1) Download easysys from this url: http://sourceforge.net/projects/easysvs/

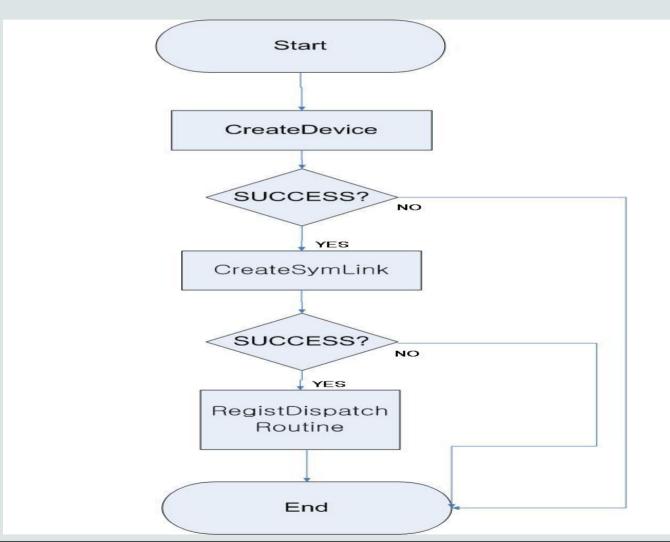
2) Easysys







### **Simple Device Driver:**



## www.

# Get into the Ringo!

### 1. Using SCM API

- 1) When a driver is loaded using the SCM, it is non-pageable. This means your callback functions, IRP-handling functions, and other Important code will not vanish from Memory.
- 2) You can select start mode of Driver.

### SERVICE\_BOOT\_START(0x0)

: Driver will be loaded by system loader.

### SERVICE\_SYSTEM\_START(0x1)

: Load Driver when IoInitSystem is called.

### SERVICE\_AUTO\_START(0x2)

: Load by SCM

#### **SERVICE DEMAND START(0x3)**

: Load by calling StartService() API.

### SERVICE\_DISABLED(0x4)

: Makes driver can not be loaded.

### 2. Using Undocumented API

#### Pro:

By using this way, you can load a Driver into the kernel without having to create registry key.

#### Con:

- 1) The problem with this approach is That the driver will be pageable. Sometimes when memory is paged out, it cannot be accessed; It will occur **BSOD**(Blue Screen Of Death) with system crash.
- 2) Once it is loaded, it cannot be Unloaded until reboot.





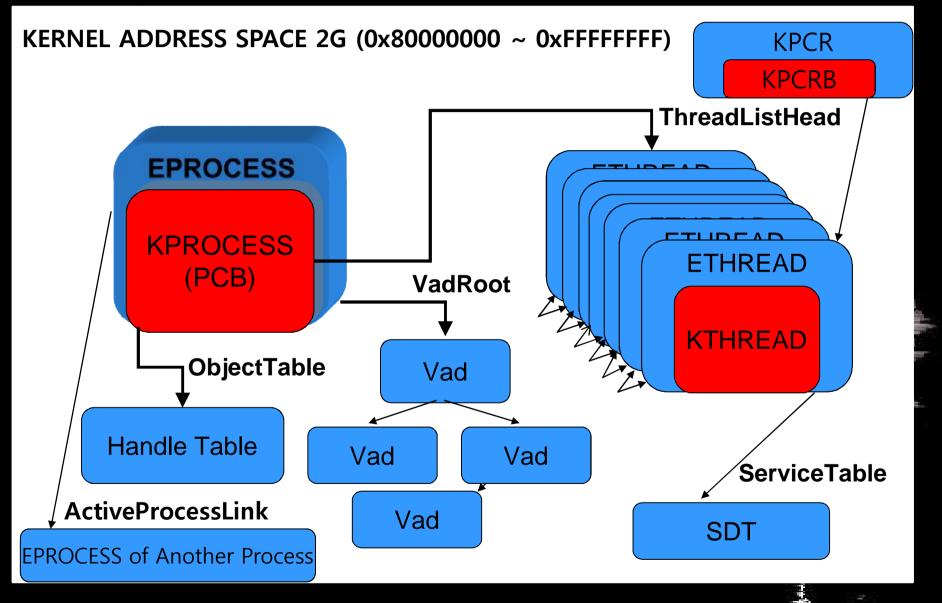






- 1) KPCR (Kernel's Processor Control Region)
- 2) EPROCESS (Executable Process)
- 3) TEB (Thread Environment Block)
- 4) ETHREAD (Executable Thread)





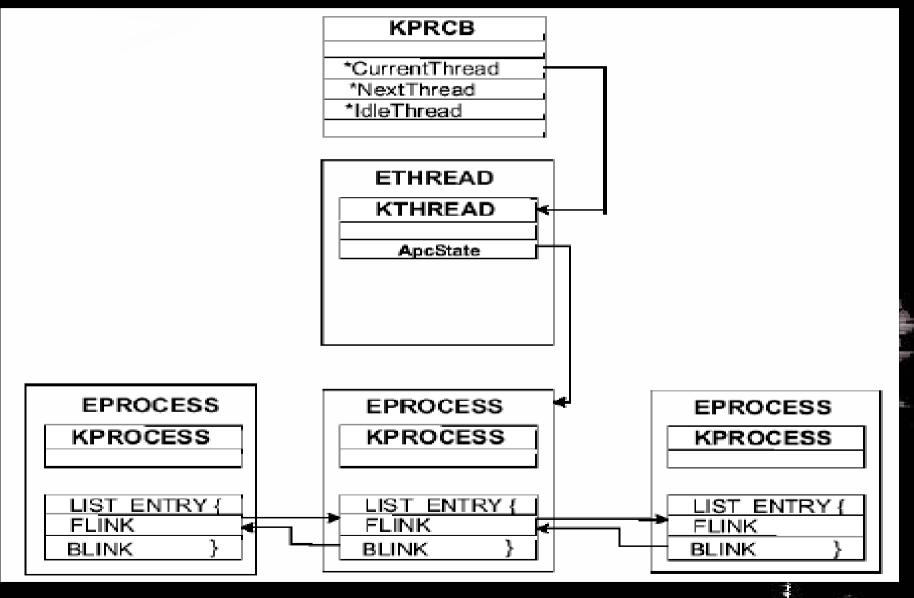


<b>KPCR</b> ( Kernel's Processor	Control Region )
----------------------------------	------------------

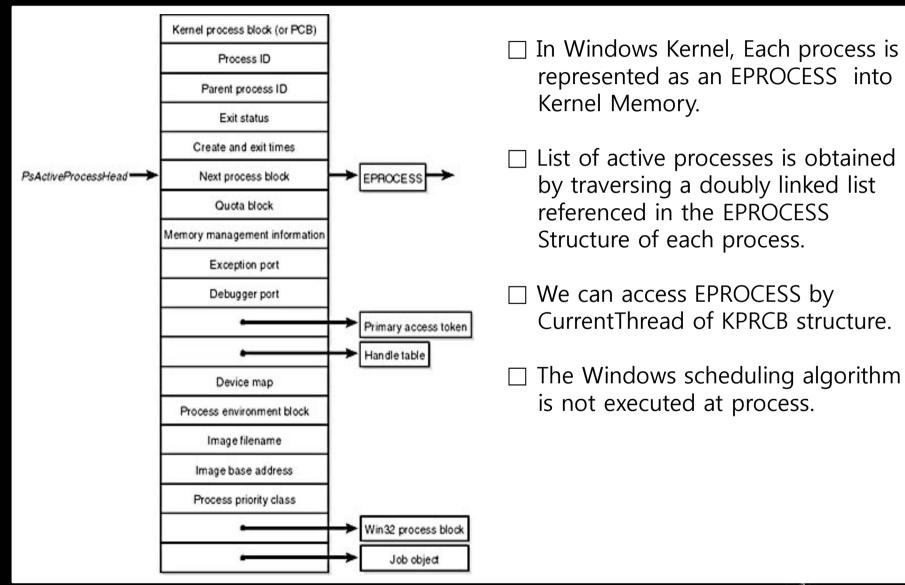
In Kernel Level, FS:0 = KPCR = 0xFFDFF000

b nter	
nter l	
Arbitrary UserPointer	
SelfTib	
1	

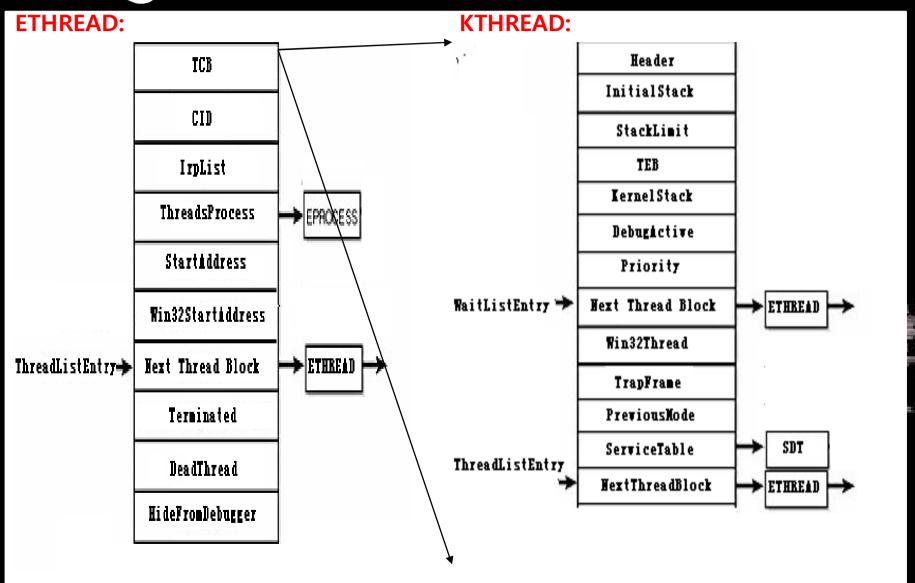






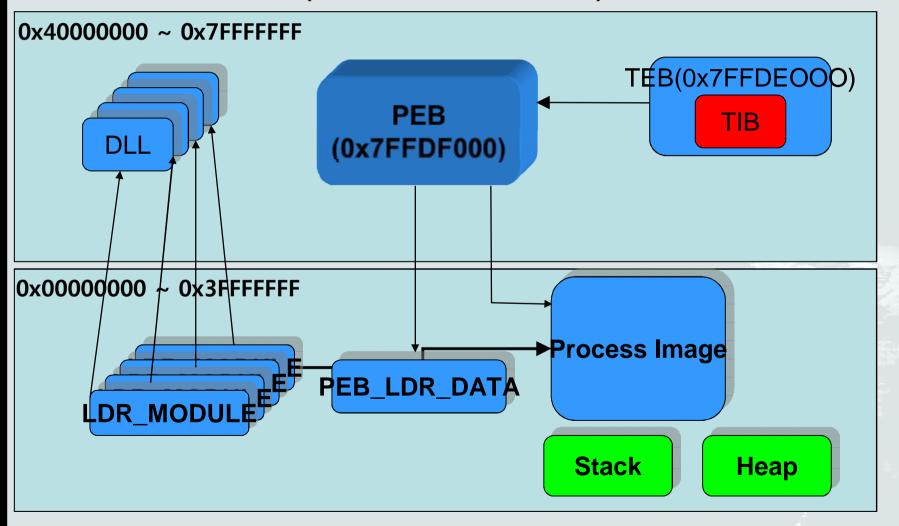








### **USER ADDRESS SPACE 2G (0x00000000 ~ 0x7FFFFFFF)**





TIB

CID

### **TEB** (Thread Environment Block)

ExceptionList

Stack Base

StackLimit.

Gub System Tib

Arbitrary UserPo Inter

SelfTib

**EnvironmentPointer** 

**UniqueProcess** 

Unique Thread

**RpcHandle** 

ThreadLocalStorage

\*PEB

LastErrorValue

In UserLevel, FS:0 = TEB = 0x7ffde000 (1st) Process)

- ☐ UserMode Application can access this structure directly.
- ☐ GetCurrentProcessId() function use this structure for getting PID.
- ☐ We can get PEB address from this structure.



### **PEB** ( Process Environment Block )

BeingDebugged ImageBaseAddress kernel32, dll-→ Next Module Block Ldr → **ProcessParameters** ProcessHeap GdiSharedHandleTable OSMajorVersion **OSMinorVersion OSBuildNumber** GdiHandleBuffer

UserMode Application can access this structure directly.
 This structure have loaded modules list.

☐ This structure have ProcessParameters information.

☐ IsDebuggerPresent() function use this structure.



## www.C

# Hook the planet!

- 1) Win32 UserLevel API Global Hooking
- 2) SSDT(System Service Distpach Table) hooking
- 3) IDT(Interrupt Descriptor Table) hooking
- 4) One byte hooking
- 5) Blind hooking by using DRx



Win32 UserLevel API Global Hooking Motivation

- ☐ Is there no way to Dll Injection without read/write process memory, and CreateRemoteThread() API?
- $\square$  Is there no way to hook Win32 API globaly?

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# Hook the planet!

- ☐ In Windows, DLL memory and mapped memory are shared.
- ☐ If we can modify memory of DLL without PageFault(Copy-On-Write), It will be applied to all processes.
- ☐ If we have chance to execute our code by victim process self, We don't need CreateRemoteThread() anymore.
- ☐ In Windows, There is a region that will get mapped into every process address space, The name of this area is KUSER\_SHARED\_DATA.

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# Hook the planet!

#### **Process**

- 1. Get the EPROCESS of explorer.exe
- 2. Attach to explorer.exe
- 3. Get the PEB from EPROCESS
- 4. Find Kernel32.dll address from Ldr in PEB.
- Find IsDebuggerPresent() address from IMAGE\_EXPORT\_DIRECTORY of kernel32.dll
- Modify IsDebuggerPresent() code to jmp KUSER\_SHARED\_DATA+0x8
   with CR0 trick.
- 7. Now just waiting for victim process call IsDebuggerPresent()



### Original IsDebuggerPresent() code:

MOV EAX,DWORD PTR FS:[18]
MOV EAX,DWORD PTR DS:[EAX+30]
MOVZX EAX,BYTE PTR DS:[EAX+2]

Modifed IsDebuggerPresent() code:

MOV EAX , KUSER\_SHARED\_DATA+0x8 JMP EAX



### PAYLOAD:

### <u>LoadLibrary address</u>

Return address

DLL

loading

code

Full DLL path

### DLL loading code:

pushad

pushfd

push KSD + 0x8 + sizeof(DLL load

code)

call [KUSER\_SHARED\_DATA]

popfd

popad

jmp [KUSER\_SHARED\_DATA+0x4]



### Demonstration





SSDT(System Service Distpach Table) hooking Motivation

- ☐ Is there no way to hook Native API as easier as IAT hook?
- ☐ Is there no way to hook Native API globaly?



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- $\square$  Native system services's addresses are listed in SSDT.
- ☐ KeServiceDescriptorTable is exported by the ntoskrnl.exe
- Windows XP later versions of the OS make the SSDT read-only, but we can bypass this protection with CR0 trick or MDL.
- ☐ We can get index number of Native API that we want to hook from ntdll.dll

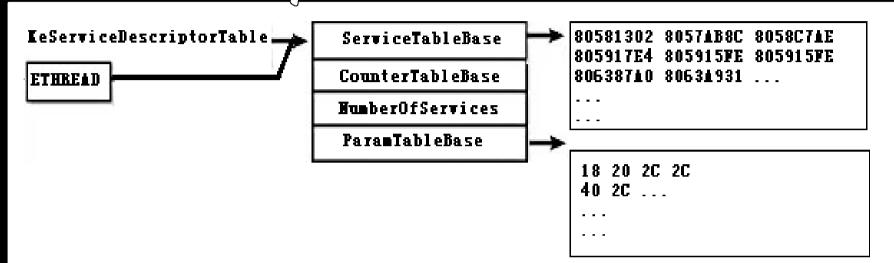


#### **Process**

- 1. Build the function that has same prototype with Native API to hook.
- 2. Get the index number of Native API.
- 3. Get address of SSDT by referencing KeServiceDescriptorTable.
- 4. Make SSDT memory area writeable.
- 5. KeServiceDescriptorTable->ServiceTableBase[index] = HookFunction;

### www.CodeEngn.com

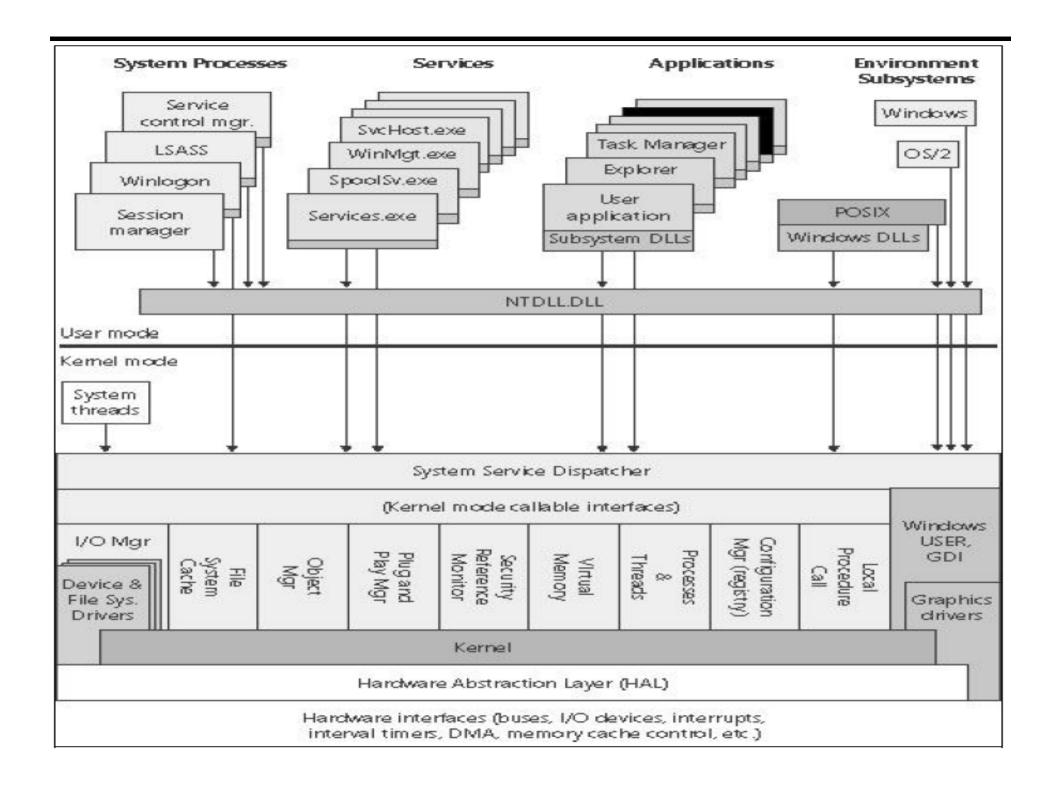
## Hook the planet!



- ☐ The KeServiceDescriptorTable is a table exported by the kernel,
  This table contains the core system services implemented in **ntoskrnl.exe**.
- ☐ There is another table in Windows Kernel, called KeServiceDescriptorTableShadow, that contains the address of USER & GDI services implemented in win32k.sys.

This table is not exported.

☐ ServiceTable pointer of ETHREAD is not always have same value with KSDT.



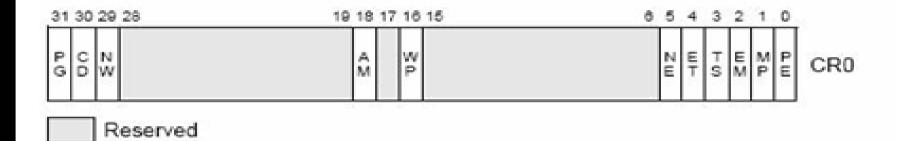
### www.CodeEngn.com

# Hook the planet!

- ☐ A system service dispatch is triggered when an INT 2E or SYSENTER instruction is called.
- □ Lower version than Windows 2000, Windows use INT 2E for service dispatch, Higher than Windows XP, Windows use SYSENTER for service dispatch.
- □ Normally, System call from UserLevel through ntdll.dll, but direct system call is possible.
- ☐ In Windows XP, we can use INT 2E instruction.
- ☐ System call number is contained in the EAX register.



What is the CR0 trick?



- ☐ The WP bit controls whether the processor will allow writes to memory pages marked as read-only.
- ☐ Setting WP(Write Protection) to zero disables memory protection.



What is the MDL?

#### MDL (Memory Descriptor List):

Next				
Size				
MdlFlags				
Process				
MappedSystemVa				
StartVa				
ByteCount				
ByteOffset				

- 1) Create the memory into our domain.
  - MmCreateMdI()
- 2) MDL build for NonPage
  - MmBuildMdlForNonPagedPool()
- 3) Change the flags of the MDL
  - MdlFlags |=MDL\_MAPPED\_TO\_SYSTEM\_VA
- 4) Lock that page
  - MmMapLockedPages()



How to get index number?

#### ZwWriteFile in Ntdll.dll:

7093E9F3 7093E9F8		MOV EAX, <mark>(12</mark> ) MOV EDX,7FFE0300
7C93E9FD 7C93E9FF	· · <del></del>	CALL DWORD PTR DS:[EDX] RETN 24

#### InvalidateRect in user32.dll:

77CFB5F5	B8 C2110000	MOV EAX, 11C2
77CFB5FA	BA 0003FE7F	MOV EDX,7FFE0300
77CFB5FF	FF12	CALL DWORD PTR DS:[EDX]
77CFB601	C2 0C00	RETN 0C



#### Demonstration



IDT(Interrupt Descriptor Table) hooking Motivation

- ☐ IDT is used to handle interrupts, so there are many tasty things.
- ☐ IDT hooking is more powerful than SDT hooking.





Concept

☐ We can use sidt, lidt instruction for getting and saving IDT information.



#### **IDTENTRY**

LowOffset

selector

unused\_lo

Type

Always0

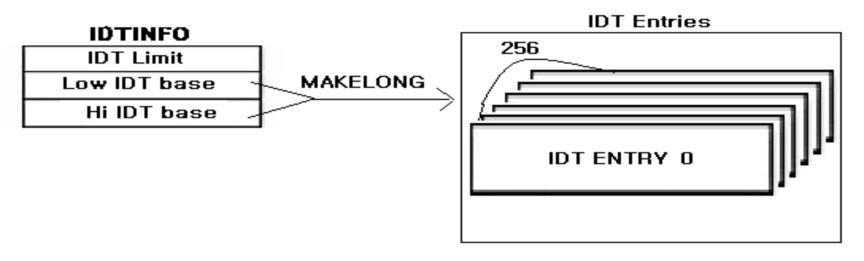
DPL

present bit

HiOffset

- ☐ Handler Address =MAKELONG(LowOffset, HiOffset);
- ☐ If DPL value of gate is 3, It can be called from both User and Kernel level.
- ☐ There are 3 types of gate.
  - Interrupt Gate (X 1 1 0)
  - Trap Gate (X 1 1 1)
  - Task Gate (0 1 0 1)





- ☐ The SIDT instruction is used to find the IDT in memory, It returns the address of the IDTINFO structure. The LIDT instruction is used for saving information into IDT.
- □ The IDT specifies how to process interrupts such as those fired when a key pressed, when a page fault occurs.
- $\square$  The Total number of IDT gates is 256.

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### Hook the planet!

#### Selector:

```
#define KGDT_NULL
#define KGDT_R0_CODE
#define KGDT R0 DATA
#define KGDT R3 CODE
                        24
#define KGDT R3 DATA
                       32
#define KGDT TSS
                    40
#define KGDT R0 PCR
                     48
                      56
#define KGDT_R3_TEB
#define KGDT VDM TILE 64
#define KGDT LDT
                     72
#define KGDT DF TSS
                      80
#define KGDT_NMI_TSS
                      88
```



#### There are so many interrupts:

Int.	DPL	p	Descriptor Type	Descriptic	)n
0000	0	P	Interrupt Gate	Fault	Divide Krror
0001			Interrupt Gate	Fault/Trap	Debug
0002	0	P	Task Gate	Interrupt	NMI Interrupt
0003	3	P	Interrupt Gate	Trap	Breakpoint
0004	3	P	Interrupt Gate	Trap	Overflow
0005	0	P	Interrupt Gate	Fault	BOUND Range Exceeded
0006	0	P	Interrupt Gate	Fault	Invalid Opcode (Undefined Opcode). Was intr
0007	0	P	Interrupt Gate	Fault	Device Not Available (No Math Coprocessor)
0008	0	P	Task Gate	Abort	Double Fault



#### Demonstration





One byte hooking Motivation

- ☐ I`m so lazy person to restore original codes of hooked function **J**
- ☐ Inline hooking is so static.



Concept

- □ Some functions in Windows XP later versions of the OS have MOV EDI, EDI instruction.
- ☐ MOV EDI, EDI doesn`t take the effect to code flow.
- $\square$  MOV EDI,EDI (=0x8B,0xFF) -> INT 0xFF (=0xCD,0xFF)
- $\square$  We can set handler at *INT 0xFF manually.*



#### **Process**

- 1. Get address of function.
- 2. Hook IDT (INT 0xFF) for all processors.
- 3. Make memory region of function to hook writeable.
- 4. Overwrite MOV EDI, EDI with 0xCD, the 'int' opcode making INT 0xFF.



#### Demonstration





Blind hooking by using DRx Motivation

- $\square$  I want to hook function without memory or table patching.
- ☐ Are there no way to hook dinamically?



Concept

- $\square$  We can use DR0  $\sim$  DR3 for set addresses to hook.
- $\square$  We can set handler at *INT 0x01* (Debug Exception).



#### Process

- 1. Get address of function.
- 2. Hook IDT (INT 0x01) for all processors.
- 3. Set addresses at DR0 ~ DR3.
- 4. Set hook type(E/W/RW) at DR7.



```
typedef struct tagDebugReg7
                                             DR0:
   unsigned L0:1; //
   unsigned G0:1; //
                                             ADDRESS 1
   unsigned L1:1; //
   unsigned G1:1; //
   unsigned L2:1; //
                                             DR1:
   unsigned G2:1; //
   unsigned L3:1; //
   unsigned G3:1; //
                                             ADDRESS 2
   unsigned GL:1; //
   unsigned GE:1; //
   unsigned undefined1 :3; ///001
                                            DR2:
   unsigned GD:1; //
   unsigned undefined2 ;2; // 00/
   unsigned RW0:2;
                                             ADDRESS 3
   unsigned LEN0 :2;
   unsigned RW1 :2;
   unsigned LEN1:2;
                                             DR3:
   unsigned RW2:2;
   unsigned LEN2:2;
   unsigned RW3:2;
                                             ADDRESS 4
   unsigned LEN3:2;
   } DebugReg7;
```





Lx?

Gx?

GD?

RWx? 00(execute), 01(write), 11(read & write).

LENx? 00(byte), 01(word), 11(dword).



```
typedef struct
  DebugReg6
unsigned B0:1;
unsigned B1:1;
unsigned B2:1;
unsigned B3:1;
unsigned undefined1:9;
unsigned BD :1;
unsigned BS:1;
unsigned BT:1;
unsigned undefined2:16;
} DebugReg6;
```

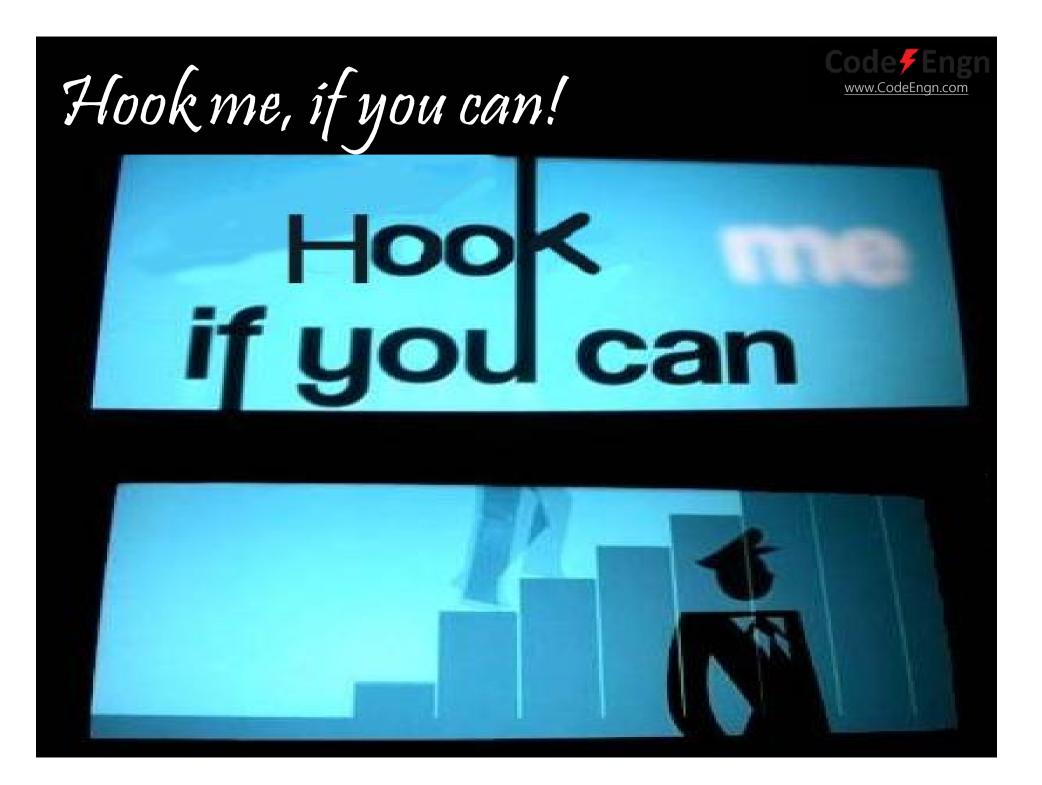
When Interrupt1 was occurred by DR0 ~ DR3.

When Interrupt1 was occurred by Gd bit.

When Interrupt1 was occurred by TF bit.



#### Demonstration



- 1) SDT Restore
- 2) SDT Relocation
- 3) KiFastCallEntry imitation
- 4) Bypassing Sysenter hook



**SDT Restore Motivation** 

- ☐ I want to be free from SDT hooking!
- ☐ Is there no way to restore SDT?



Concept

- We can get original copy of the ServiceTable by loading ntoskrnl.exe into memory.



#### **Process**

- 1. Use NtOpenSection to get a handle ₩device₩Physicalmemory with SECTION\_MAP\_READ | SECTION\_MAP\_WRITE access.
- 2. Load ntoskrnl.exe into memory.
- 3. Use NtMapViewOfSection to map in the physical memory page.
- 4. Get the address of ServiceTable from the page.
- 5. Use the address of ServiceTable to offset into the loaded ntoskrnl.exe
- 6. comparing the copy in the kernel memory with the copy in the loaded ntoskrnl.exe



#### Demonstration



**SDT Relocation Motivation** 

- ☐ How about change SDT as new thing?
- ☐ Is there no way to make SDT hook localy?



Concept

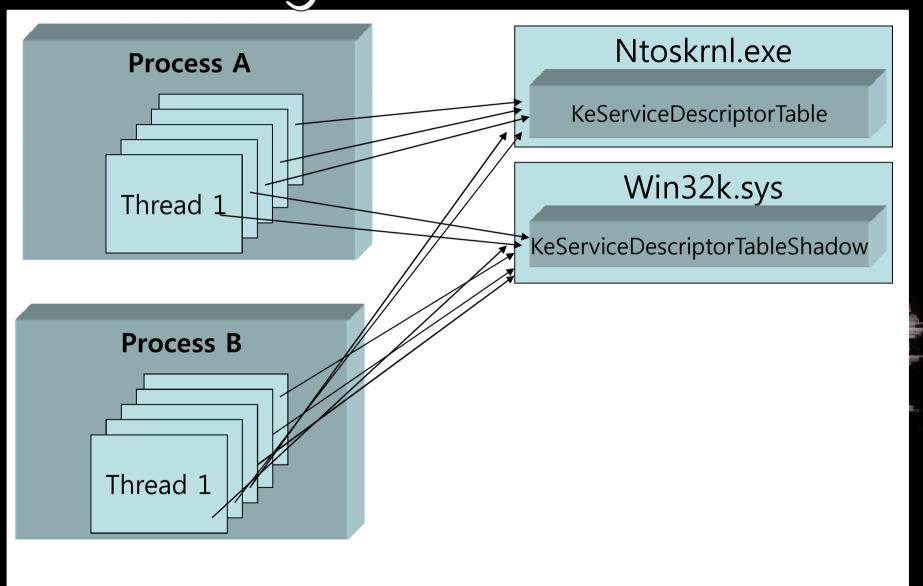
- □ Windows doesn`t use KeServiceDescriptorTable for getting address of ServiceTable. Actually, Windows use ETHREAD->ServiceTable to get address of ServiceTable.
- ☐ When new thread is created, we can know this by using PsSetCreateThreadNotifyRoutine()



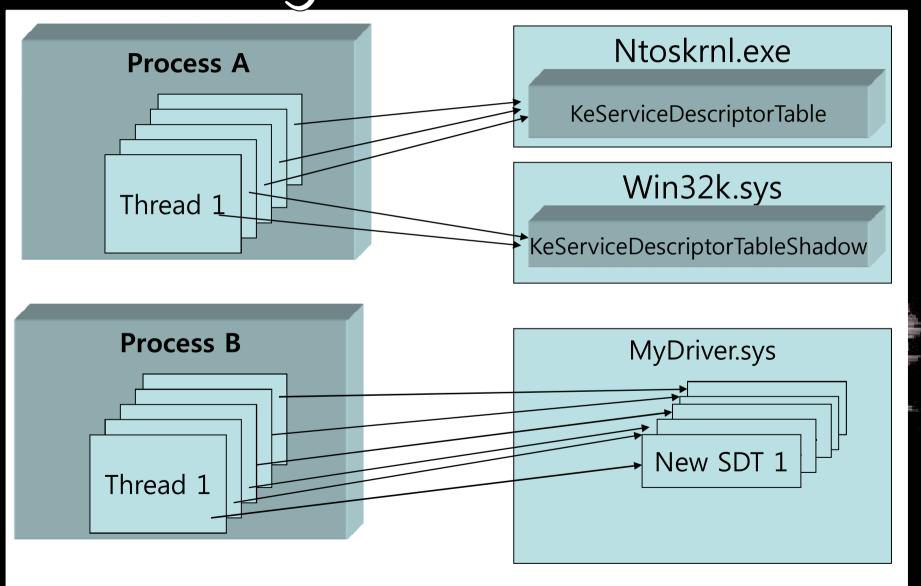
#### **Process**

- 1. Copy the SerivceTable from ntoskrnl.exe
- 2. Tracing threads of seleted process, and make NewSDT at each thread.
- 3. Regist PsSetCreateThreadNotifyRoutine()











#### What is the problem?

☐ When UserThread is maded, ETHREAD->ServiceTable of that UserThread doesn`t have KeServiceDescriptorTableShadow.

Where does problem come from?

☐ This problem is caused by PsConvertToGuiThread() that called by KiFastCallEntry().

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## Hookme, if you can!

```
PsConvertToGuiThread():
//Check Kernel Mode
cmp byte ptr[esi+0x140],bl //KTHREAD PREVIOUS MODE,kernelmode
 jz InvalidParam
 cmp dword ptr[0xBF820ECE],ebx //PspWin32ProcessCallback?
 jz AccessDenied
 cmp dword ptr[esi+0xE0],KeServiceDescriptorTable
                                                    //KTHREAD_SERUICE_TABLE
 inz AlreadyWin32
```



#### Demonstration



#### **KiFastCallEntry Imitation Motivation**

- $\square$  Is there no way to bypass SDT hooking by hooking?
- ☐ Is there no way to handle syscall by myself?



#### Concept

□ The SYSENTER instruction passes control to the address specified in one of the Model-Specific Registers(MSRs).
 □ The Name of this register is IA32\_SYSENTER\_EIP,
 □ We can read and write to this register, By using RDMSR , WRMSR instruction.



#### **Process**

- Copy the KeServiceDescriptorTable from ntoskrnl.exe,
   Copy the KeServiceDescriptorTableShadow from Win32k.sys
- 2. Change IA32\_SYSENTER\_EIP for all processors.
- 3. Hadling system call **J**

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### Hookme, if you can!

```
How we can do it?
KiFastCallEntry():
 mov ebx,dword ptr[edi+0xC] //SERVICE DESCRIPTOR NUMBER
 xor ecx,ecx
 mov cl,byte ptr [eax+ebx]
 //Get the function point
 mov edi,dword ptr[edi] //Service Descriptor Base
```



What is the thing that I did?

```
//Get the function point
mov edi,dword ptr[edi] //Service_Descriptor_Base

//Blocking SSDT Hooking
cmp edi,Orginal_ServiceTable
jne Shadow

mov edi,NewServiceTable
jmp NotShadow
Shadow:
mov edi,NewShadowTable
NotShadow:
mov ebx,dword ptr[edi+eax*4]
```



### Demonstration



Bypassing sysenter hook Motivation

- ☐ Is there no way to bypass sysenter hook?
- ☐ Is there more way to bypass SDT hook?



#### Concept

- ☐ In Windows XP, Defaultly, Windows use sysenter to handle system call, but using int 0x2e is possible.
- ☐ We can modify the code of KiFastSystemCall() in the ntdll.dll by using Win32 UserLevel API global hook.



#### **Process**

- 1. Hook IDT (INT 0x2E) for all processors.
- 2. Modify the code of KiFastSystemCall()
- 3. Hadling system call **J**



### Before:

7C93EB8B 8BD4 MOV EDX,ESP 7C93EB8D 0F34 SYSENTER

### After:

7C93EB8B	8BD4	MOV EDX,ESP
7C93EB8D	CD 2E	INT 2E



How we can do it?

```
__declspec(naked) MyKiSystemService()
{
    __asm
    {
        ......

        ExitService:
        sti
        jmp SyscallEntry
        .....
```



```
SysCallEntry:
       mov edi,eax
       shr edi,0x8 //SERUICE TABLE SHIFT
       and edi, 0x30 //SERVICE_TABLE_MASK
       mov ecx,edi
       mov cl,byte ptr [eax+ebx]
       //Get the function point
       mov edi,dword ptr[edi] //Service Descriptor Base
```



### What is the thing that I did?

```
//Get the function point
mov edi,dword ptr[edi] //Service_Descriptor_Base

//Blocking SSDT Hooking
cmp edi,Orginal_ServiceTable
jne Shadow

mov edi,NewServiceTable
jmp NotShadow
Shadow:
mov edi,NewShadowTable
NotShadow:
mov ebx,dword ptr[edi+eax*4]
```





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