

The Alethe Proof Format

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Abstract

Specifying more rules for alethe in the Cutting Planes reasoning

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1 Introduction

In this article we describe how the cutting planes reasoning can be expressed in the alethe proof format, in order to make use of checkers (carcara).

2 Pseudo Boolean Inequality in Alethe

The Pseudo Boolean format consists of:

$$\sum_i a_i l_i \geq A \quad (1)$$

where,

$$\begin{aligned} A, a_i &\in \mathbb{N} \\ l_i &\in \{x_i, \bar{x}_i\}, \quad x_i + \bar{x}_i = 1 \end{aligned} \quad (2)$$

In order to express it in alethe, we define an expression:

$$(>= (+ <TERMS> 0) <A>)$$

where `<TERMS>` is a list of either:

1. `(* <a_i> <l_i>)` in a plain literal $a_i x_i$.
 2. `(* (- 1 <a_i>) <l_i>)` in a negated literal $a_i \bar{x}_i$
- and `<A>` is the natural constant A .

3 Cutting Planes Rules in Alethe

3.1 Rule 1: cp_addition

We can apply the rules, using the ‘step’ syntax:

```
(assume c1 (>= (+ (* 2 x1) 0) 1))
(assume c2 (>= (+ (* 1 (- 1 x1)) 0) 1))
(step t1 (c1 (>= (+ (* 1 x1) 0) 1))
  :rule cp\_addition :premises (c1 c2)
)
```

3.2 Rule 2: cp_multiplication

```
(assume c1 (>= (+ (* 1 x1) 0) 1))
(step t1 (c1 (>= (+ (* 2 x1) 0) 2))
  :rule cp\_multiplication :premises (c1) :args (2)
)
```

3.3 Rule 3: cp_division

```
(assume c1 (>= (+ (* 2 x1) 0) 2))
(step t1 (c1 (>= (+ (* 1 x1) 0) 1))
  :rule cp\_division :premises (c1) :args (2)
)
```

3.4 Rule 4: cp_saturation

```
(assume c1 (>= (+ (* 2 x1) 0) 1))
(step t1 (c1 (>= (+ (* 1 x1) 0) 1))
  :rule cp\_saturation :premises (c1)
)
```

References