# AN ABSTRACT OF THE DISSERTATION OF

<u>John Smith</u> for the degree of <u>Doctor of Philosophy</u> in <u>Computer Science</u> presented on
September 23, 2011.
Title: The Meaning of Life
Abstract approved:
Joan Smythe
This is an abstract statement.

©Copyright by John Smith September 23, 2011 All Rights Reserved

# The Meaning of Life

by

John Smith

### A DISSERTATION

submitted to

Oregon State University

in partial fulfillment of the requirements for the degree of

Doctor of Philosophy

Presented September 23, 2011 Commencement June 2012

<u>Doctor of Philosophy</u> dissertation of <u>John Smith</u> presented on <u>September 23, 2011</u> .
APPROVED:
Major Professor, representing Computer Science
Director of the School of Electrical Engineering and Computer Science
Dean of the Graduate School
I understand that my dissertation will become part of the permanent collection of Oregon State University libraries. My signature below authorizes release of my dissertation to any reader upon request.
John Smith, Author

# ACKNOWLEDGEMENTS

I would like to acknowledge the Starting State and the Transition Function.

## TABLE OF CONTENTS

				Page
1	Int	roduction		1
	1.1	Introduct	tion to the Variant Annotation	1
	1.2	Definition	n	1
		1.2.1	Gene Structure	1
		1.2.2	Potential Variant Effects	1
2	Me	ethods		4
	2.1	Input		4
		2.1.1	Genome data	4
		2.1.2	Gene structure data	4
		2.1.3	Variant data	4
	2.2	Output .		5
	2.3	How it w	vorks	6
		2.3.1	Outline of the algorithm	6
		2.3.2	Searching algorithm	6
		2.3.3	Array based searching	6
		2.3.4	Tree based searching	7
	2.4		sues	8
		2.4.1 2.4.2	how to get amino acid	8 10
		2.4.2	strand	10
3	Re	sult		11
	3.1	Running	time	11
		3.1.1	Sorted variants	11
	3.2	Memory	cost	11
	3.3	Success 1	rate	11
В	ibliog	graphy		11
A	ppen	dices		12
	A		ncv	13

## LIST OF FIGURES

Figure		Page
1.1	potential variant functional regions	1
2.1	Genome data example	4
2.2	GTF file example	5
2.3	VCF file example	5
2.4	VCF file example	5

## LIST OF ALGORITHMS

Algorith	<u>m</u>	Page
1	VARIANT ANNOTATION	7
2	ARRAY BASED SEARCHING	8
3	TREE BASED SEARCHING	9

## Chapter 1: Introduction

- 1.1 Introduction to the Variant Annotation
- 1.2 Definition
- 1.2.1 Gene Structure
- 1.2.2 Potential Variant Effects

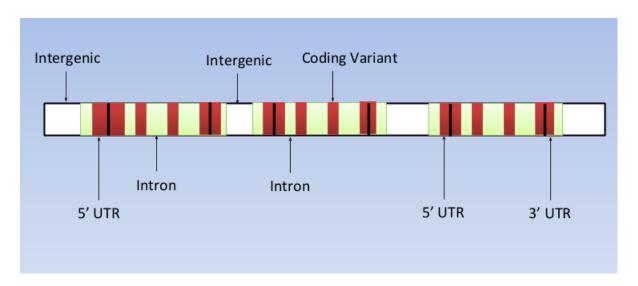


Figure 1.1: potential variant functional regions

In this project, we annotate 7 potential variant effects. The corresponding functional regions are show in Figure 1.1

• Intergenic variant
Intergenic variant does not reside in any genes, resulting in that it does not alter any amino acid sequences. So its putative impact is low.

#### • Intron variant

Intron variant resides in a gene, but not in an exon. Thus, it does not alter any amino acid either. Its putative impact is still low but a little higher than Intergenic variant.

#### • 5' UTR variant

5' UTR variant resides in an exon, but is prior to the start of coding sequence. Although it still does not alter any amino acid sequences, it is relatively more damaging than intergenic and intron variant.

### • 3' UTR variant

3' UTR variant resides in an exon, but after the stop codon of coding sequences. The putative impact is similar to 5' UTR variant.

### • Non-coding exon variant

Non-coding exon variant resides in an exon. But the corresponding gene has no start or stop codon, resulting in no coding sequences. Thus, no amino acid sequences would be altered.

## • Synonymous variant

Synonymous variant resides in an exon and coding sequences, leading to altering a coding sequence. However, since several DNA codons correspond to one amino acid, the resulting amino acid sequence does not change.

### • Non-synonymous variant

Non-synonymous variant resides in coding sequences and alters an amino acid sequence. Its putative impact is high.

The following table shows the notation and putative impact for each variant effect.

Putative impact
MODIFIER
LOW
MODERATE

For non-synonymous variant, we annotate the reference amino acid, alternate amino acid and the position of the amino acid instead of the term.

## Chapter 2: Methods

### 2.1 Input

#### 2.1.1 Genome data

The genome data contains the whole nucleotide sequences of each chromosome. It is often stored in a faidx-indexed reference file in the FASTA format. We will use this file to get the amino acid codon of a given variant. Example:

Figure 2.1: Genome data example

### 2.1.2 Gene structure data

The gene structure data contains information about gene structure, such as start codon, stop codon, exon, strand and gene id. This is the key data to determine the functional region of a given variant. Example:

### 2.1.3 Variant data

This is the data to be annotated. The key information includes chromosome, position, reference base and alternate variant alleles of the variant. The data is typically stored in VCF file. We will annotate the variant in the INFO field. Example:

chr1	canFam3_ensGene	exon	8350657	8351219	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	exon	8359233	8359447	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	start_co	don	8365671	8365673 0	.000000	+		gei
chr1	canFam3_ensGene	CDS	8365671	8365716	0.000000	+	0	gene_id	"El
chr1	canFam3_ensGene	exon	8365645	8365716	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	CDS	8366024	8366362	0.000000	+	2	gene_id	"El
chr1	canFam3_ensGene	exon	8366024	8366362	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	CDS	8410380	8410727	0.000000	+	2	gene_id	"El
chr1	canFam3_ensGene	exon	8410380	8410727	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	CDS	8427874	8427976	0.000000	+	2	gene_id	"El
chr1	canFam3_ensGene	exon	8427874	8427976	0.000000	+		gene_id	"El
chr1	canFam3_ensGene	CDS	8453827	8453842	0.000000	+	1	gene_id	"El

Figure 2.2: GTF file example

```
#CHROM
        POS
                 ΙD
                          REF
                                   ALT
                                            QUAL
                                                                      FORMAT
                                   T
        366459
                          G
                                                             SC=0; KS=NOVEL; CV=COVERED; PW=0.9999
                                                    PASS
        sum=40;n_alt_sum=KEEP
                                      RC:AC:RS:AS
                                                          954141:91:4:3398
                                                                                  0:GG:23.515183:
        724795
                                   Т
                                                             SC=0; KS=NOVEL; CV=COVERED; PW=0.9992
                          G
                                                    PASS
                   _sum=KEEP RC:AC:RS:AS
                                               2.1856:76:
                                                                         0:GG:24.594482:82
        1244743
                          C
                                                    PASS
                                                             SC=0; KS=NOVEL; CV=COVERED; PW=0.9999
=3163;n_ref_sum=0;n_alt_sum=KEEP
                                      RC:AC:RS:AS
                                                        2.027738:94:4:3575
                                                                                  0:CC:24.982382:
        1244775
                                                    PASS
                                                             SC=0; KS=NOVEL; CV=COVERED; PW=0.9998
```

Figure 2.3: VCF file example

### 2.2 Output

Genes may overlap. Hence, one variant may have several effects. For each variant, the output is a list of effects. The potential effects are listed in section 1.1. In the project, the output is stored in a VCF file in the INFO field. To be specific, a new subfield, named ANN, will be appended to INFO field. The most important information is the effect of variants and the putative impact of the effect.

Besides the effect of variants, the corresponding Ensemble transcript id and common gene name are also annotated. All the annotation formats follow the standard variant annotation in VCF format. Example:

ANN=A|intron\_variant|MODIFIER|ENSCAFT00000046825|ENSCAFG00000000012|ATP9B,

Figure 2.4: VCF file example

### 2.3 How it works

There are 3 datasets to be kept track of. If we search on the whole datasets for each variants, the total running time complexity is O(m\*n\*k), where m, n and k is the size of each file. In addition, the size of the datasets may be too large to load them all into RAM. For instance, the size of a typical human genome file is about 3.0 GB. To solve the difficulties in this problem, we develop several algorithms to annotate variants.

The whole problem can be split into 2 smaller subproblems. The first one is to obtain a list of genes in which a variant resides. The second one is to annotate each functional region. And if the variant resides in a coding sequence, get the corresponding DNA sequence around the variant and translate it into amino acids.

The most challenging part is to get the gene list since genes may overlap. The naive solution is to search on the whole genome for each gene, resulting in a time complexity of O(N\*K), where N is the number of genes and K is the number of variants. To speed up the annotation, 2 algorithms were proposed. One algorithm is array based and efficient for sorted variants while the other is tree based and efficient for unsorted variants.

## 2.3.1 Outline of the algorithm

## 2.3.2 Searching algorithm

2 algorithms are proposed for searching gene list. One is array based and most efficient for sorted variants while the other is tree based and most efficient for unsorted variants.

## 2.3.3 Array based searching

Array based search algorithm is most efficient for sorted variants. Genes are stored in sorted order in an array. While iterate though the variants, the algorithm keep track of the last visited gene. Each variant start searching from the last visited gene, not from the beginning. At the end, all the variants and genes are visited only once, resulting in a complexity of O(N+K), where N is the number of genes and K is the number of variants.

### **Algorithm 1** VARIANT ANNOTATION

```
input: gene structure data G. reference genome data R. variant data V.
output: The annotated variant data V
for each variant v in V at chromosome chr and position pos do
   gene\_list = get\_gene\_list(G, chr, pos)
   A = \emptyset
   if gene_list is empty then
       A = A + "intergenic_region"
   end
   for each gene in gene_list do
       if gene has neither start codon nor stop codon then
          A = A + \text{"non\_coding\_exon\_variant"}
       else if v does not reside in any exon of gene then
       A = A + "intron_variant"
       else if v resides prior to start codon of gene then
       A = A + "5" prime_UTR_variant"
       end
       else if v resides after stop codon of gene then
        A = A + "3" prime_UTR_variant"
       end
           get reference amino acid ra and alternative amino acid aa
            if ra == aa then
            A = A + "synonymous_variant"
           end
           A = A + "ra + pos + aa"
           end
       end
   end
   Annotate v with A
end
```

**return** The annotated variants V

## 2.3.4 Tree based searching

Tree based algorithm is most efficient for unsorted variants. In this algorithm, genes are stored in an interval tree. Although each variant searching on the whole gene, it takes only O(logN)

#### Algorithm 2 ARRAY BASED SEARCHING

```
input: gene structure data G, chromosome chr and position pos output: a list of genes which reside in chr and overlap pos get gene array gene\_arr and last visited index i from G on chromosome chr while i < length of gene\_arr and <math>pos > the \ end of gene\_arr[i] do | i = i + 1 end update last visited index i to G gene\_list = \varnothing j = i while j < length of gene\_arr and <math>gene\_arr[j] overlap pos do | add \ gene\_arr[j] to gene\_list | j = j + 1 end return gene\_list
```

time to get the gene list. Thus, the total time complexity is O(K\*(logN)). If sort the variants and apply array based algorithm, the total time complexity would be O(K\*(logK)+N), which is larger than O(K\*(logN)).

Interval tree is very similar to binary search tree. Standard binary search tree usually maintains a single value for each node while interval tree maintains an interval for each node. When searching genes, it makes a preorder traversal on the interval tree until the position is completely outside of the current node.

The key to keep binary search tree efficient is maintaining it balanced. Since genes rarely change, we don't need a balanced binary search tree algorithm here, such as Red Black tree. We could sort the genes by their starting position first and then build a balanced binary search tree with it.

### 2.4 Other issues

## 2.4.1 how to get amino acid

The whole genome sequence data is very large. For instance, the size of a typical human genome file is about 3.0 GB. However, since coding sequences amount to only 1.5% of whole genome,

### Algorithm 3 TREE BASED SEARCHING

```
input: gene structure data G, chromosome chr and position pos output: a list of genes which reside in chr and overlap pos
```

```
Function get\_gene\_list(G, chr, pos)
get gene tree root gene\_root from G on chromosome chr
gene\_list = \varnothing
binary\_search(gene\_root, pos, gene\_list)
return gene\_list
```

**Function** binary\_search(root, pos, gene\_list)

```
if root is empty then
| return
end

if pos resides outside region of root then
| return
end

g = the gene in root
if pos resides in g then
| add g to gene_list
end
```

binary\_search(root.leftchild, pos, gene\_list)
binary\_search(root.rightchild, pos, gene\_list)

we could preload the coding sequence to the gene structure data. In this project, the final gene structure data is about 100 MB, which is smaller enough to load in most machines' memory.

## 2.4.2 strand

Another annoying problem is that genes have different strand. Different from positive strand, negative strand starts from a bigger position and ends at a smaller position. In addition, the nucleotide sequence must be reversed according to the format A-T, C-G.

## Chapter 3: Result

# 3.1 Running time

## 3.1.1 Sorted variants

The following table shows the running time for sorted variants. There are 8777 variants in this experiment. The variants come from a real variant project. The time unit is second.

naive	array	tree	
12.64	0.75	0.93	

- 3.2 Memory cost
- 3.3 Success rate

## **APPENDICES**

# Appendix A: Redundancy

This appendix is inoperable.