

# Example: Minimal DFA $NFA \rightarrow DFA$ $\epsilon$ closure / composite states

“The DFA uses its state to keep track of all possible states the NFA can be in after reading each input symbol.”

**Algorithm** Subset construction: Constructing a DFA from an NFA

*Input* An NFA  $N$

*Output:* A DFA  $D$  accepting the same language

*Goal:* “Each DFA state is a set of NFA states and we construct  $D$ -tran so that  $D$  will simulate in parallel all possible moves  $N$  can make on a given input string.”

*Method.* Construct a transition table  $D$ -tran for  $D$ . Apply  $\epsilon$  closure and move methods,

Table 1: default

|                            |  |
|----------------------------|--|
| $\epsilon$ -closure( $s$ ) | Set of NFA states reachable from NFA state $s$ on $\epsilon$ -transitions alone                    |
| $\epsilon$ -closure( $T$ ) | Set of NFA states reachable from some NFA state $s \in T$ on $\epsilon$ -transitions alone         |
| $move(T, a)$               | Set of NFA states to which there is a transition on input symbol $a$ from some NFA state $s \in T$ |