

# PyExaFMM: Designing a high-performance point fast multipole solver in Python with Numba

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**Abstract**—PyExaFMM is a Python implementation of the kernel-independent point fast multipole method. Building on the success of the ExaFMM project, we aimed to answer the question of whether it was possible to develop a highly-performant scientific code without resorting to a lower level language. The computational experience of domain specialists is often restricted to interpreted languages, such as Python or Matlab. Therefore the ability to develop high-performance software in a high-level language is appealing, reducing the barrier to entry for potential contributors, as well as making deployment to new architectures or platforms easier. The FMM is a good case study for understanding the maturity of Python for developing high-performance software, due its reliance on a complex hierarchical tree data structure. In this article we explore the software engineering and mathematical techniques used to extract performance for PyExaFMM, and we report that we achieve runtimes within  $\mathcal{O}(10)$  of the state of the art C++ implementation, with comparable accuracy, for three dimensional electrostatic problems in single precision.

■ **THE FAST MULTIPOLE METHOD** [FMM], [1], approximates the solution of the so called  $N$  body problem, in which one seeks to calculate

pairwise interactions between  $N$  objects. This problem arises in numerous contexts in science and engineering, for example in the calculation of electrostatic or gravitational potentials due to a set of charged or massive source objects. This is an important calculation, as with an estimate of potential one is able to describe the corresponding electric, or gravitational field due to these objects, and hence the forces they exert. In this article we use the electrostatic potential generated by point charges in three dimensions as our model to explore the FMM. Here, the potential,  $\phi(x_j)$ , for a given point at position  $x_j$ , or ‘target’, due to  $N$  points at positions  $x_i$ , or ‘sources’, where  $i \in [1, \dots, N]$ , each with a charge  $q_i$ , can be written as,

$$\phi(x_j) = \sum_{i=1}^N K(x_i, x_j) q_i, \quad (1)$$

here  $K(\cdot, \cdot)$  is called the Green’s function, or kernel function. For electrostatic problems in three dimensions it is,

$$K(x, y) = \frac{1}{4\epsilon_0\pi|x - y|}, \quad (2)$$

where  $\epsilon_0$  is the permittivity of free space, and is referred to as the Laplace kernel. Without loss of generality, we consider the set of sources and set of targets to correspond to the same set of points.

Intuitively, (2) can be understood as encoding a ‘decay’ relationship between two given points. The closer they are the higher their proportionate effect on each other’s potential. Attempting to evaluate the sum in (1) directly for  $N$  points, results in algorithm of  $\mathcal{O}(N^2)$  runtime complexity, as the sum will have to be repeated  $N$  times. However the FMM is able to approximate (1) with just  $\mathcal{O}(N)$  runtime complexity, with prescribed error bounds. This makes the simulation of granular physically realistic systems tractable.

### The FMM

The key idea behind the FMM is to encode the potential in the far field due to a cluster of points with a representative analytic series expansion centered on the cluster, referred to as a *multipole expansion*. This expansion can be

truncated to tune for a desired accuracy, which allows one to approximate the sum in (1) with fewer calculations. Alternatively, we can center the expansion far away from the cluster, the encoded potentials are then referred to as a *local expansion*. Translations between multipole and local expansions can be done analytically, and are critical in the development of the  $\mathcal{O}(N)$  algorithm. However, this approximation is only valid where the kernel function of a problem exhibits appropriate ‘decay’ characteristics, allowing us to assume that potential due to clusters of points far away from a region of interest can be approximated with a truncated representative expansion.

In the FMM, the problem domain is described by a cubic box enclosing all targets and sources, which is hierarchically partitioned into a structure known as a quadtree in two dimensions, and an octree in three dimensions. The partitioning is done recursively. Initially it is partitioned into four equal parts in two dimensions (eight in three dimensions), known as its children. Correspondingly the unrefined box is referred to as their parent. The children are each subsequently refined in a similar fashion. The final tree consisting of a series of hierarchically nested boxes.

The size of the box is described by its ‘level’, with level zero denoting the unrefined box, level one denoting the level of its children, and so forth. One can define the maximum level of refinement through a parameter dictating the maximum allowable number of points contained within the finest box covering a portion of the domain. These finest boxes that remain after refinement are known as leaves. Note that the set of all leaves cover the entire domain, without overlap. One can choose to refine *adaptively*, such that the leaves may be of different sizes, reflective of the underlying point distribution, or *non-adaptively*, such that the leaves are all of a uniform size.

The FMM then consists of two sequential traversals of this tree, where boxes are considered level by level. Initially, during the *upward pass*, multipole expansions for points contained in boxes at the leaf level are formed. This is referred to as the *point-to-multipole* [P2M] operation. In order to obtain the multipole expansions for their parents, the expansion centers of a given box’s children are shifted to the center of their shared

parent box, and their coefficients are summed. This is referred to as the *multipole-to-multipole* [M2M] operation. This is repeated bottom-up, level by level, until one is left with the multipole expansions of all boxes in the tree.

Subsequently, during the *downward pass*, the multipole expansions of non-adjacent boxes of a given box are defined as being in its far field. In order to evaluate their contribution towards the potential for points within a given box, one translates their multipole expansion to a local expansion centered at the given box. This is referred to as the *multipole-to-local* [M2L] operation. The local expansion is then transferred to the children of the given box by shifting the expansion center to the center of the child boxes, and updating the coefficients of the child boxes' local expansions. Proceeding top-down, considering boxes level by level, one is left with the local expansions for each leaf box. These are referred to as *local-to-local* [L2L] operations.

Crucially the extent of the far field not already encapsulated in the local expansion shrinks as we descend down the levels of the tree. This is because the far field of a box's parent is already captured in the child's local expansion. These local expansions compress all of the contribution from the far field towards the potential for targets within a given leaf. Therefore the far field's effect on the potential of targets within a leaf box is entirely described by the local expansion.

The evaluation of the local expansion at the target points in a given leaf box is known as the *local-to-point* [L2P] operation. Finally, the near fields of a given leaf box are evaluated directly using (1), this is referred to as the *point-to-point* [P2P] operation. As there are a fixed number of points per leaf node, the complexity of this direct evaluation is bounded. The operations in this scheme are illustrated for a two dimensional problem, with a uniform quadtree, in figure (1).

The  $\mathcal{O}(N)$  algorithmic runtime complexity can be roughly seen to be the result of the ability to exactly translate between the multipole and local expansion representations for all boxes deemed to be in the far field of a given box, as well as the recursive procedure of the FMM. As a result of this, each box must only consider its interaction with a constant number of other boxes. As there are  $\mathcal{O}(N)$  boxes in a given tree, the

entire algorithm can be seen to be bounded by a runtime complexity of  $\mathcal{O}(N)$ .

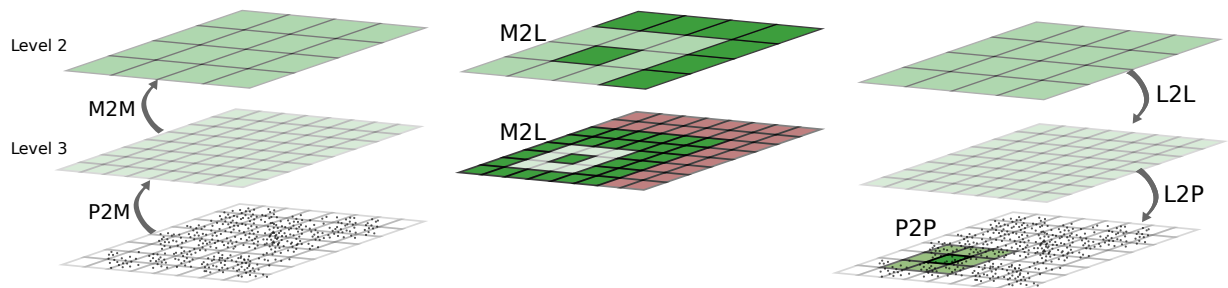
### Kernel Independent FMM

The coefficients of a given multipole expansion are kernel-dependent in the sense that they will depend on the form of (2), software implementations have to be rewritten for each specific physical model. Fast multipole methods for generic kernel functions have been developed [2], however there are fewer methods that rely only on numerical kernel evaluations, rather than analytic series expansions, examples include [3], [4]. The latter approaches are often referred to as kernel-independent fast multipole methods [KIFMM].

PyExaFMM utilizes the approach first presented by Ying et. al [4]. This KIFMM represents the multipole expansions due to a cluster of charges as set of equivalent charge densities supported on a surface enclosing the cluster. The fields generated by the charges are matched with an equivalent field generated by these equivalent densities via a least-squares fitting of the potential they generate at another surface in the far field. This method generalizes the FMM to non-oscillatory second-order elliptic partial differential equations with constant coefficients, for example the Laplace or Stokes equations. This includes many common problems, such as our model electrostatic problem.

### The Software Context

At present, numerous high-quality open source implementations exist to solve FMM problems, both using analytic expansions for various kernels [5], as well as using the KIFMM of Fong et. al [6], [7], and Ying et. al [8], [9]. Some of these implementations are able to scale on supercomputing clusters, and solve problems involving millions of degrees of freedom [9]. Others are designed to be run on a single workstation [6], [8]. These softwares are developed in compiled languages, such as Fortran or C++, often making use of numerous non-standard libraries and language features. Compiled languages are often accompanied by complex build processes, making deployment to different platforms or architectures non-trivial. Many domain specialists who may want to use the FMM to aid their research usually have limited computational skill sets, often re-



**Figure 1.** The FMM illustrated for a two dimensional problem with a uniform quadtree. The left column illustrates the first two steps of the upward pass, beginning with the P2M operation to find the multipole expansions at the leaf level, followed by an M2M operation. The central and right columns illustrate the downward pass. In the central column, the far-field of a given box is highlighted, we observe that at level three the far field highlighted in red is already encapsulated in the local expansion of the given box due to the L2L translation from its parent. In the right column we see how the algorithm concludes with a given box evaluating its near field of adjacent boxes, directly via the P2P operation.

stricted to subsets of Python or Matlab. Therefore it is not feasible for many users of to contribute to, or extend, many of these existing libraries without a significant investment in acquiring new skills. This problem is not unique to FMM implementations, but is general across high-performance computing [HPC] software.

#### Python for High Performance

In recent years, HPC libraries that aid in bypassing the limitations of the Python interpreter have emerged. The most pertinent example is Numba, a ‘just in time’ [JIT] compiler for CPython, a popular C implementation of the Python interpreter. JIT refers to the fact that the compiler waits until a Numba enhanced object, either a function or a class, is instantiated at runtime, before running optimizations to convert the Python source code into efficient machine code. Numba achieves performance by focussing solely on optimizing operations on ndarrays, the multi-dimensional array objects provided by Numpy . These objects have become the basis for numerical libraries written in Python as Numpy binds to widely used HPC libraries for numerical computing written in C or Fortran, such as LAPACK and BLAS.

Numba extends Numpy by removing inefficient layers of indirection for indexing into Numpy arrays, and instead translating array index operations into direct load and store operations from pointers [10]. Furthermore, it of-

fers simple multithreading support through a ‘prange’ statement that allows you to loop over processes, mirroring OpenMP’s parallel for loops. Numba even supports automatic single-instruction-multiple-data [SIMD] vectorization. Adapting to the instruction set architecture of your CPU, whether it supports Advanced Vector Extensions [AVX], AVX-512, or Streaming SIMD Extensions [SSE]. Therefore with Numba it’s easy to heavily optimize methods that loop over arrays. Finally, Numba allows one to decorate functions to be run on GPUs, allowing the development of heterogenous applications with a single Python source.

Importantly, Numba is a ‘drop-in’ library. It’s simply installed via package manager, and objects are marked to be optimized by a single decorator. If Numba is unable to perform an optimization, it fails silently, deferring to the ordinary Python runtime. Therefore users are able to seamlessly incorporate Numba into new or existing projects. Previously, achieving similar behavior in Python would have required the usage of a hybrid language such as Cython, or by writing custom C extensions on top of CPython.

Numba has been extended with optimized versions of many common array operations, as well as common linear algebra routines available through the Numpy library. These are called automatically when used in a marked object. For common array operations, Numba is able to achieve comparable performance with compiled

**Table 1. Laplace kernel evaluations over 100,000 randomly distributed particles. Coordinate components, and source charge densities chosen in interval  $[0, 1)$ . Repeated seven times for statistics.**

| Runtime (s)      |                    |
|------------------|--------------------|
| Numba (PyExaFMM) | $4.71 \pm 0.29$ s  |
| C++ (ExaFMM-T)   | $1.99s \pm 0.01$ s |

languages. See table (1), for a comparison of direct Laplace kernel evaluations, using (1), for 100,000 points between a Numba modified PyExaFMM, and ExaFMM-t, both of which are multi-threaded and SIMD vectorized, or [10] for further examples.<sup>1</sup>

### The Motivation for Building PyExaFMM

The question arises of whether it is feasible to develop a complex HPC application entirely within Python. The vision being that a domain specialist may be able to iterate from a prototype, to a high-performance application, using a language they are comfortable with. Python has the additional benefit of simple cross platform build tools, such as Conda, making it easy for developers to code once and deploy everywhere. An FMM implementation is an excellent case study to understand the maturity of Python for HPC development. As Numba is built to optimize operations on arrays, it represents a challenge to develop efficient representations and operations for the tree on which the algorithm is based.

PyExaFMM was developed to answer the above question. It is currently a single node implementation designed for three-dimensional problems, implementing parallel strategies on multicore architectures, similar to [6], [8]. We demonstrate the feasibility of building a high-performance implementation of a non-trivial algorithm solely in Python. As with all Python applications, performance is degraded by portions of code that are unavoidably interpreted. However, this represents a compromise between usability and performance. Where by usability, we mean the ease by which non-software specialists can understand, alter, or extend the codebase to their needs. Using lines-of-code [LOC] as a rough metric of usability, PyExaFMM consists of approximately 5,800 LOC, compared with

approximately 20,000 LOC and 50,000 LOC for the comparable C++ implementations, ExaFMM-T [8] and TBFMM [6] respectively.

In this article we begin by providing an overview of the mathematical formulation behind the KIFMM in [4], before discussing how Numba is used by PyExaFMM to accelerate computations. We continue by examining other mathematical and software based optimizations used to extract performance. Specifically, the efficient pre-computation, and application of FMM operations, and an overview of the software’s design. We conclude with a comparison of performance, in terms of accuracy, runtimes, and memory footprint, with the comparable state-of-the-art C++ implementation from the ExaFMM project, ExaFMM-T [8].

### Formulating the Kernel Independent FMM

Consider the P2M operation illustrated in figure (2a). A cubic *equivalent* surface, sharing the same center,  $S_e$  is used to enclose a leaf from the tree containing set of points. It is discretized with evenly spaced quadrature points placed its faces, each of which supports an equivalent charge density. In order to find these equivalent densities, we match the potential generated by the points directly,  $\phi_c$ , calculated using (1) with that generated by the equivalent densities at the *check* surface,  $S_c$ , in the far field, also centered on the box, which encloses both the box and the corresponding equivalent surface. Using (1) we find,

$$\sum_{x_i \in S_e, x_j \in S_c} K(x_i, x_j) q_i = \phi_c \quad (3)$$

where  $x_i$  and  $x_j$  are the quadrature points for the equivalent and check surfaces respectively,  $q_i$  is the equivalent charge density supported at  $x_i$ . Written in matrix form, we can write the equivalent densities,  $q_e$ , as the solution of the following linear system,

$$K q_e = \phi_c \quad (4)$$

$$q_e = K^{-1} \phi_c \quad (5)$$

<sup>1</sup>Note that all benchmarks are reported on a six core Intel i7-9750H CPU



where  $K$  denotes a matrix with elements,

$$K_{i,j} = K(x_i, x_j), \quad x_j \in S_c, \quad x_i \in S_e \quad (6)$$

We identify these equivalent densities as a multipole expansion. The order of this expansion,  $p$ , is related to the number of the quadrature points,  $n$ , by,

$$n = 6(p-1)^2 + 2 \quad (7)$$

Equation (5) is in general ill-conditioned, therefore we use a truncated singular value decomposition [SVD] of  $K$ , to calculate a pseudo-inverse. We cut off components corresponding to singular values smaller than machine epsilon in single precision.

As shown in figure (2a), the equivalent surface is enclosed by the check surface for the P2M operation. When this is the case we refer to the surfaces as the *upward* equivalent and check surface respectively. We find that the ratio of the side-length of the upward equivalent and check surfaces with respect to the tree node as 1.05 and 1.90 work well in practice.

We perform a similar calculation for the M2M operation, illustrated in figure (2b). Here the upward equivalent density of a child box, is matched to that of its parent,

$$K_c q_e^{child} = K_p q_e^{parent} \quad (8)$$

where  $K_c$  and  $K_p$  are matrices with elements as defined by (6), where  $S_c$  and  $S_e$  denote their relevant check and equivalent surfaces. To find the child's contribution to the parent's upward equivalent density, and hence multipole expansion, we solve,

$$q_e^{parent} = K_p^{-1} K_c q_e^{child} \quad (9)$$

we refer to  $K_p^{-1} K_c$  as the *M2M operator*, with four unique operators in two dimensions (eight in three dimensions) corresponding to all the children of a given box. To find the corresponding *L2L operators* and *M2L operators*, we simply repeat the procedure with the appropriate check and equivalent surfaces. For the L2L operation, the check surface is placed inside the equivalent surface. This reflects the region in which we expect the local expansion to valid.

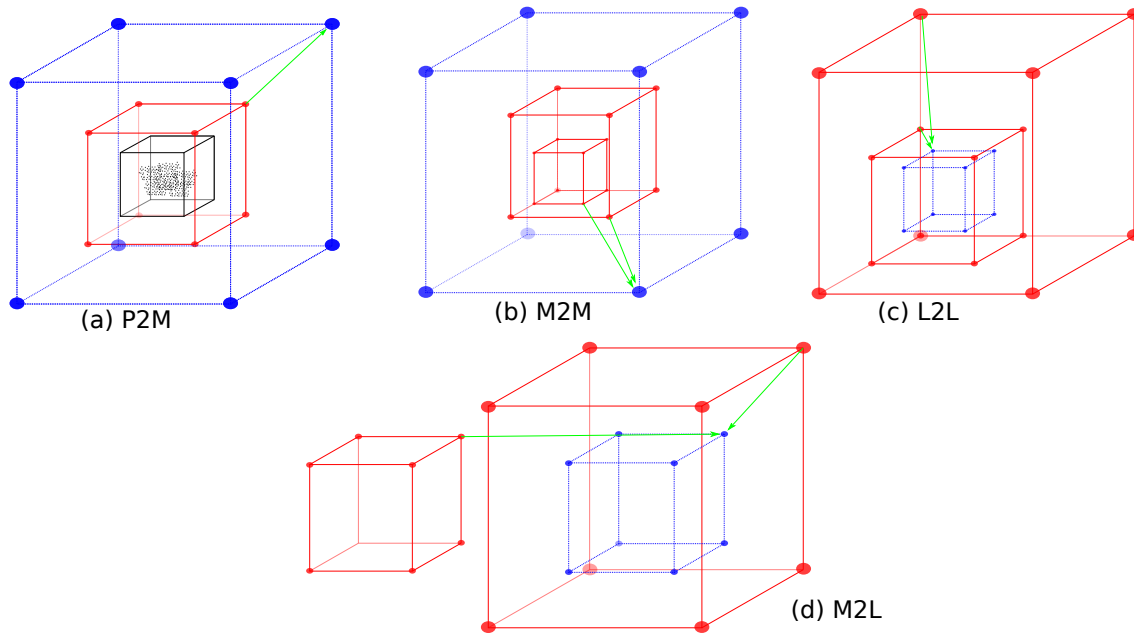
Here they are referred to as the *downward* check and equivalent surfaces respectively. For the M2L operation (fig. 2d) the upward equivalent surface of the source box, is matched to the potential generated by the downward equivalent surface at the target box.

We choose the upward equivalent surface to correspond to the downward check surface, and the downward equivalent surface to correspond to the upward check surface, for a given level. We then notice that the quantity  $K^{-1}$  in (5) can be re-used in the calculation of the M2M operator, as well as the M2L operator. Additionally, the M2M and L2L operators remain the same for all boxes in the tree.

As PyExaFMM uses adaptive octrees, there are further permutations in the composition of the near and far field we must describe. Consider a target box  $T$ , all source boxes  $S$  directly neighboring  $T$  (sharing an edge, face, or corner) are in  $T$ 's *U list*. Denoting a neighbor of a box at the same level as its *colleague*. The children of  $T$ 's colleagues which are not adjacent to  $T$  are in its *W list*, and the children of the colleagues of  $T$ 's parent which are not adjacent to  $T$  are in its *V list*. Finally, the *X list* of  $T$  consists of boxes for which  $T$  is in their *W list*. The *X* and *W* list are only calculated for leaf boxes. Collectively these lists are known as the *interaction lists* of  $T$ , and are illustrated for a two dimensional problem in figure (3).

Greater specification of the interaction lists allows for a more efficient algorithm implementation. Considering a source box in  $T$ 's *W list*, the multipole expansion of the source box is valid despite not being in  $T$ 's far field, therefore reducing the cost of calculating the near field contribution. For source boxes in  $T$ 's *X list*, the multipole expansion is not valid, as  $T$  is not in the far-field of the source box. However, the source charges can be evaluated directly at the quadrature points on  $T$ 's downward equivalent surface.

Implementing efficient KIFMMs relies on efficiently caching and loading the matrices that represent the P2M, M2M, L2L and M2L operators. As the algorithm consists of a series of matrix-vector products, one has to find a way to represent the coefficients of the multipole and local expansions, as well as tree nodes, to take



**Figure 2.** The operators of the KIFMM. Equivalent surfaces are denoted in red, check surfaces in blue, and the charged points are denoted in black. Surfaces are plotted with 8 quadrature points, for clarity.

|   |   |   |   |   |   |   |  |
|---|---|---|---|---|---|---|--|
| X | X |   |   |   | V | V |  |
|   |   |   |   |   | V | V |  |
| X | W | W | W | W | U | V |  |
|   | W | U | U | U |   |   |  |
| X | U | T | U | U | U | V |  |
|   | V | V | V | V | V | V |  |
|   |   |   |   |   |   |   |  |

**Figure 3.** Interaction Lists for a target leaf box,  $T$ . Unmarked nodes are not considered to interact with  $T$ .

advantage of the cache-locality and vectorization offered by modern CPUs, or GPUs. The P2M and M2L steps are the largest bottlenecks in implementing fast KIFMMs [9]. The M2L operator has to be applied for each box in a

given target's V list, and in three dimensions the maximum possible occupancy of a box's V list is 189. The P2M operator involves calculating the potential due to all particles in all leaves, at their check surfaces. However, the P2M operator is embarrassingly parallel over leaves, as is the M2L operator over a given V list.

## TECHNIQUES FOR ACHIEVING PERFORMANCE

### Tree Construction

- Morton encoding
- Linear octree representation
- Numba for morton operations
- Numba for parallel tree construction
- Numba barriers: less pythonic, cant really rely on dicts, sets, other higher level python operations. Difficulty with python's object model, specifically with throwing functions around like objects. End up writing C like code.
- Construction time benchmark table vs ex-afmm.

### Precomputing Operators

- HDF5, loading into memory at run time.
- Transfer vectors for M2L

- hashing, and why Numba is sometimes hard to use.

#### Compressing the M2L operator

- Introduce RSVD algorithm [11].
- Demonstrate that SVD error dominated by FMM error.
- Demonstrate impact on runtime and storage.

#### Software Design

The key is separating routines to be accelerated, and organising data ready to run. The data organisation part (and it's slowness) should be demonstrated as a bottleneck using experiment.

### PERFORMANCE

Description of main experiments, and how they are conducted. What are the main results? Are they expected from theory? What conclusions can be drawn about developing HPC codes entirely in Python, is it worth it?

### CONCLUSION

What have we learned, what will we be working on next?

### ACKNOWLEDGMENT

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