Benchmarking packet reception with AF_XDP, DPDK and io_uring

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The performance of networking stacks

Why use a different networking stack?

- A need for processing packets on commodity hardware with low overhead.^[1]
- Portability and integration with the operating system versus performance.

io_uring

Keeps evolving \rightarrow need for new benchmarks.

^[1]G. Wan, F. Gong, T. Barbette, and Z. Durumeric, "Retina: analyzing 100GbE traffic on commodity hardware," in *Proceedings of the ACM SIGCOMM 2022 Conference*, in SIGCOMM '22. Association for Computing Machinery, 2022, pp. 530–544. doi: 10.1145/3544216.3544227.

Comparisons between networking stacks

Networking stacks support different layers.

Ethernet

- AF_XDP
- DPDK

TCP

- epoll
- io_uring



DPDK

Description

- Kernel bypass → no overhead due to transition between kernel space and user space.
- Networking code is written to be as fast as possible → faster than Linux networking stack.
- Need to reserve an entire NIC for the application → no sharing of resources.
- Linux NIC drivers cannot be reused.

Express Data Path (XDP)

Description

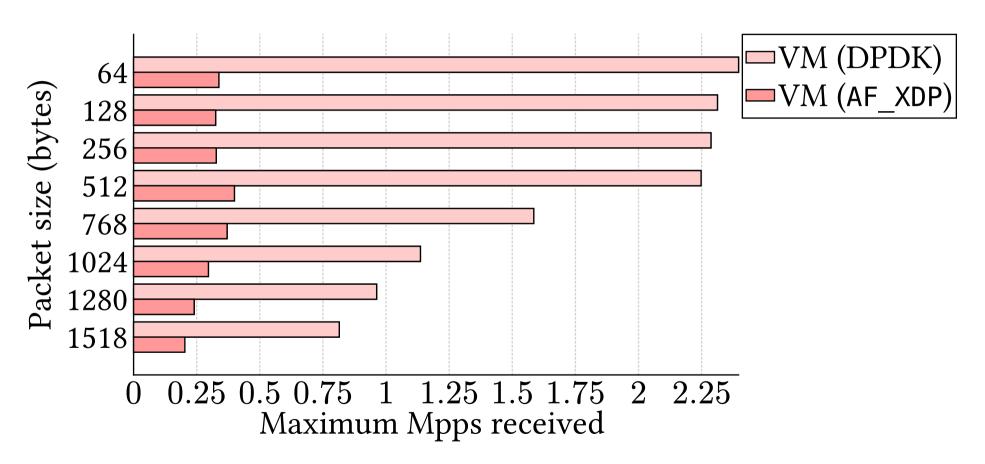
- XDP programs are BPF programs that are called for every incoming packet just after reception but before allocating memory for a socket buffer.
- They can drop packets, modify them, and chose to pass them to the networking stack, redirect them to a port or to userspace for further processing.
- Packets redirected to userspace are received on a AF_XDP socket.
- Need to reserve a NIC queue for the application.
- Uses Linux NIC drivers.

Experiment

- *Sender* machine sends packets as fast as possible to *receiver* machine using Pktgen-DPDK.
- Measure the maximum number of packets received per second.
- Vary the number of cores and packet size.



Results



Conclusions

- AF_XDP performs poorly compared to DPDK and has more variability.
- Packet size is less important with AF_XDP than DPDK → the overhead is not in memory copying.
- It's possible that I made a mistake in the setup.
- I was not able to test on bare-metal, but I think these results can still be useful for cloud providers that run everything in VMs.

TCP

BSD socket API

- Create a socket: socket, accept
- Receiving data: recv, recvfrom, read
- Sending data: send, sendto, write

Characteristics

- One system call per operation
- **Blocking**: no compute or IO can be done in the meanwhile

BSD socket API

Example

```
import random
import socket
s = socket.socket(socket.AF INET, socket.SOCK DGRAM)
s.bind(('0.0.0.0', 53))
while True:
    datagram, address = s.recvfrom(2048)
    print(f'Received {datagram} from {address}.')
```

select and friends

```
import socket, select
s = socket.socket(socket.AF INET, socket.SOCK STREAM)
s.bind(('0.0.0.0', 80))
s.listen(8)
sockets = [s]
while True:
    ready, _, _ = select.select(sockets, [], [])
for r in ready:
    if r == s:
        c, addr = s.accept()
        sockets.append(c)
    else:
        data = sock.recv(2048)
        if data:
            print(f'Received {data} from {r}')
        else:
            l.remove(r)
```

- Readiness-based
- One system call per operation
- Non-blocking
- Variants: poll, epoll, kqueue

Readiness-based interface

```
/* s is ready... */
for (;;) {
    char buf[2048];
    ssize t n = recv(s, buf, sizeof(buf), 0);
    if (n == 0)
        break;
    process(buf, n);
```

epoll

- Also readiness-based
- Modern version of select with less copying
- FreeBSD equivalent: kqueue

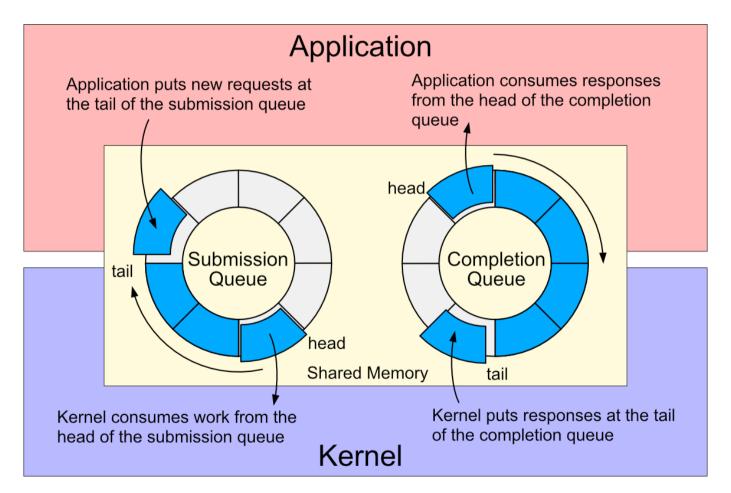
- Completion-based
- Start many operations with a single system call

Completion-based interface

```
char *buf = malloc(2048);
start_read_operation(s, buf, sizeof(buf), 1234);

/* Later... */
int id = process_completed_operation();
printf("%d\n", id); // 1234
```

Presentation



Performance improvements

- Fixed files and buffers (May 2019)
- Buffer ring (July 2022)
- Zero-copy reception (not merged yet!)

Fixed files

Pre-register FDs with the ring to reduce the overhead of:

- reference counting, and
- descriptor table lookup.

```
int io_uring_register_files(
    struct io_uring *ring,
    const int *files,
    unsigned nr_files,
);
```

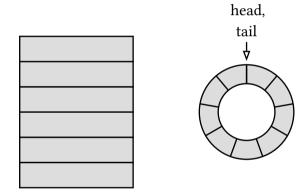
Buffer ring

Instead of specifying a buffer to use in read operations, let the kernel pick a buffer from another ring buffer.

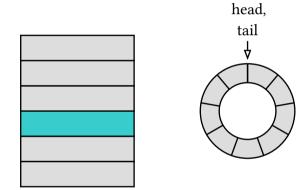
```
struct io_uring_buf_ring *io_uring_setup_buf_ring(
    struct io_uring *ring,
    unsigned int nentries,
    int bgid,
    unsigned int flags,
    int *ret,
);
```

Zero-copy reception

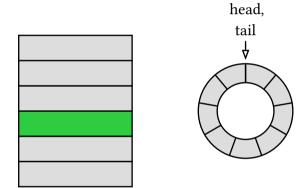
A new receive operation where the NIC directly writes to application memory.



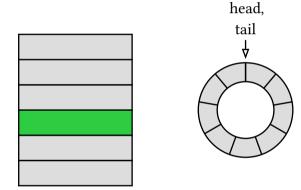
1. App submits RECV_ZC operation.



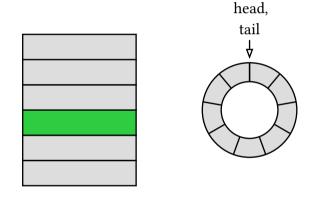
- 1. App submits RECV_ZC operation.
- 2. Kernel picks free buffer in area.



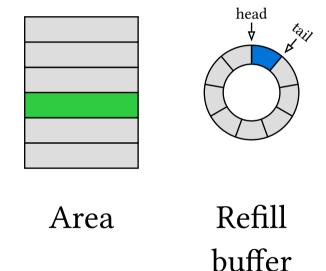
- 1. App submits RECV_ZC operation.
- 2. Kernel picks free buffer in area.
- 3. NIC writes to buffer.



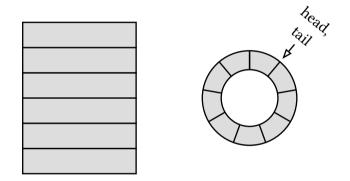
- 1. App submits RECV_ZC operation.
- 2. Kernel picks free buffer in area.
- 3. NIC writes to buffer.
- 4. Kernel enqueues completion.



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- 5. App processes data in buffer.
- 6. App enqueues buffer ready to be reused.



Area Refill buffer

- 1. App submits RECV_ZC operation.
- 2. Kernel picks free buffer in area.
- 3. NIC writes to buffer.
- 4. Kernel enqueues completion.
- 5. App processes data in buffer.
- 6. App enqueues buffer ready to be reused.
- 7. Kernel marks buffer as available.

Support

This work is not merged yet into Linux and liburing (a library to use io_uring). The patch series have some example code, but not much other than that.

I created a Rust wrapper to setup the area and refill buffer needed to use the new RECV_ZC operation: https://github.com/beviu/io-uring-zcrx.

I modified the io-uring Rust crate to add support for this new operation: https://github.com/beviu/io-uring/compare/master...zcrx.

Experiment

- *Sender* machine sends data on TCP socket as fast as possible to *receiver* machine.
- Two TCP server implementations (see server-io-uring and server-epoll directories).
- Measure the bandwidth.



Results

Unfortunately, I was not able to run the benchmarks.



Virtual machine setup

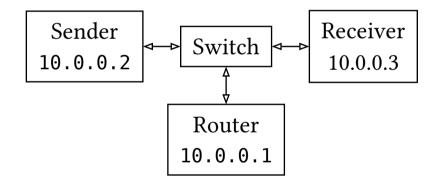
I prepared the benchmark setup beforehand, and tested it in virtual machines on my PC first to make sure I was ready before testing on actual Grid5000 hardware. Unfortunately, in the end, I was not able to test on Grid5000.

I installed Arch Linux in two VMs and configured a bridge with a tap device for each VM in the host. I used virtio-net NICs.

Virtual machine setup

Host network interfaces setup

```
ip tuntap add mode tap tap0
ip tuntap add mode tap tap1
ip link add name br0 type bridge
ip link set dev br0 up
ip link set tap0 up
ip link set tap1 up
ip link set tap0 master br0
ip link set tap1 master br0
ip addr add 10.0.0.1/24 dev br0
```



Appendix

Configure a NAT to give the tap devices access to the Internet:

```
sysctl net.ipv4.ip_forward=1
nft add table nat
nft 'add chain nat postrouting { type nat hook postrouting
priority 100 ; }'
nft add rule nat postrouting ip saddr 10.0.0.0/24 masquerade
```

Virtual machine setup

QEMU command line

Make sure to replace the tap interface name and MAC address for the receiver VM. The arguments for specifying the drives were omitted.

```
qemu-system-x86_64 \
  -machine type=q35,accel=kvm,kernel-irqchip=split \
  -cpu host -smp 4 -m 2G \
  -device intel-iommu,intremap=on,caching-mode=on \
  -vga none \
  -serial mon:stdio \
  -monitor none \
  -nographic \
  -netdev tap,id=tap,ifname=tap0,vhost=on,script=no,downscript=no \
  -device virtio-net-pci,mq=on,vectors=2,netdev=tap,mac=52:54:00:f8:e2:e3
```

Kernel configuration

I followed the instructions from the DPDK documentation.

I compiled the kernel at commit
76544811c850a1f4c055aa182b513b7a843868ea with a custom
configuration and added console=ttyS0,11520 intel_iommu=on

The first option is for using the terminal QEMU is running on as the Linux console.

vfio.enable unsafe noiommu mode=1 options to the command line.

Virtual machine setup

Guest network interfaces setup

On the sender VM:

ip addr add 10.0.0.2/24 dev enp0s2

On the receiver VM:

ip addr add 10.0.0.3/24 dev enp0s2

Compiling DPDK

```
curl https://fast.dpdk.org/rel/dpdk-24.11.1.tar.xz -0
tar -xf dpdk-24.11.1.tar.xz
cd dpdk-stable-24.11.1/
meson setup build -Dplatform=native
cd build
ninja
sudo ninja install
sudo sh -c 'echo /usr/local/lib > /etc/ld.so.conf.d/
local.conf' # Needed on Arch Linux.
sudo ldconfig
```

Compiling Pktgen-DPDK

```
git clone https://github.com/pktgen/Pktgen-DPDK.git --branch
pktgen-24.10.3
cd Pktgen-DPDK
PKG_CONFIG_PATH=/usr/local/lib/pkgconfig meson setup build
cd build
ninja
sudo ninja install
```

Running Pktgen-DPDK

Make sure to replace the PCI slot with your NIC's.

```
sudo dpdk-devbind.py --bind vfio-pci 0000:00:02.0 --force
sudo sh -c 'echo 512 > /sys/kernel/mm/hugepages/
hugepages-2048kB/nr_hugepages'
sudo pktgen -l 0,1 -n 1 -a 0000:00:02.0 -- -P -T -m 1.0
```

Running the AF_XDP server

I used an existing server implementation. It receives packets and drops them right away.

```
git clone https://github.com/xdp-project/bpf-examples.git --
recurse-submodules
cd bpf-examples
cd AF_XDP-example
make
```

Running the AF_XDP server

Run the *rxdrop* example:

```
sudo taskset -c 0 ./xdpsock \
    --interface enp0s2 \
    --busy-poll \
    --shared-umem \
    --rxdrop
```