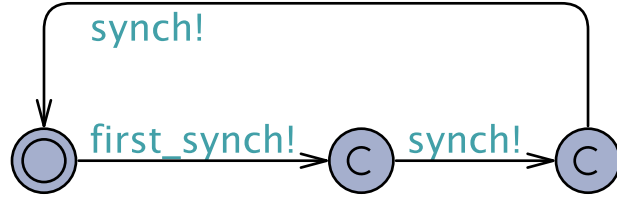


### (a) SYNCH



### (b) TASK

$x \geq s\_min()$  &&  
 $x < s\_max()$  &&  
 $t \leq deadline()$   
 $synch?$

Start  
 $t \leq offset()$   
 $t == offset()$   
 $synch?$   
 $t = 0,$   
 $x = 0,$   
 $seg\_idx = 0$

Suspended

$t > deadline()$

$x \leq s\_max()$

$x \geq s\_min()$  &&  
 $t \leq deadline()$   
 $synch?$   
 $enqueue()$

Ready

$t > deadline()$

$run[id]?$   
 $x = 0$

Running

$t > deadline()$

$x \leq c\_max()$

$!is\_last\_segment()$  &&  
 $x \geq c\_min()$  &&  
 $t \leq deadline()$   
 $first\_synch?$   
 $x = 0,$   
 $seg\_idx++,$   
 $avail\_processors++$

$x \geq c\_min()$  &&  
 $x < c\_max()$  &&  
 $t \leq deadline()$   
 $first\_synch?$

### (c) SCHED

$job\_ready() \&\&$   
 $processor\_avail()$   
 $run[front()]!$   
 $dequeue(),$   
 $avail\_processors--$

Scheduling

### (d) DECLARATIONS

```
int[0, M] avail_processors = M;  
urgent chan run[N];  
broadcast chan synch, first_synch;  
chan priority first_synch < run;  
chan priority synch < run;
```

```
bool is_last_segment() {  
    return seg_idx ==  
        Tasks[id].k - 1;  
}
```

```
bool job_ready() {  
    return queue_len > 0;  
}
```

```
bool processor_avail() {  
    return avail_processors > 0;  
}
```

$t \geq period()$   
 $synch?$

Completed

$t \geq period()$   
 $synch?$   
 $t = 0,$   
 $x = 0$

$is\_last\_segment() \&\&$   
 $x \geq c\_min()$  &&  
 $t \leq deadline()$   
 $first\_synch?$   
 $seg\_idx = 0,$   
 $avail\_processors++$