

CS301 A1

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1 Problem 1

(a) $\Theta(n^3)$

(b) $\Theta(n^{\log_2 7})$

(c) $\Theta(\sqrt{n} \log n)$

(d) $\Theta(n^2)$

2 Problem 2

(a)

(i) Best asymptotic worst case means there is no common sub-sequence between these 2 words. That means the algorithm will try to find common sub-sequence and will fail. In that way, it will be searching from the beginning to the end. In figure 1, which shows the naive algorithm, since length of lcs is 0, else part will be called every time until the end. Else part is calling the max of $\text{lcs}(X, Y, i, j-1)$ and $\text{lcs}(X, Y, i-1, j)$ recursively. We want to find $T(m+n)$.

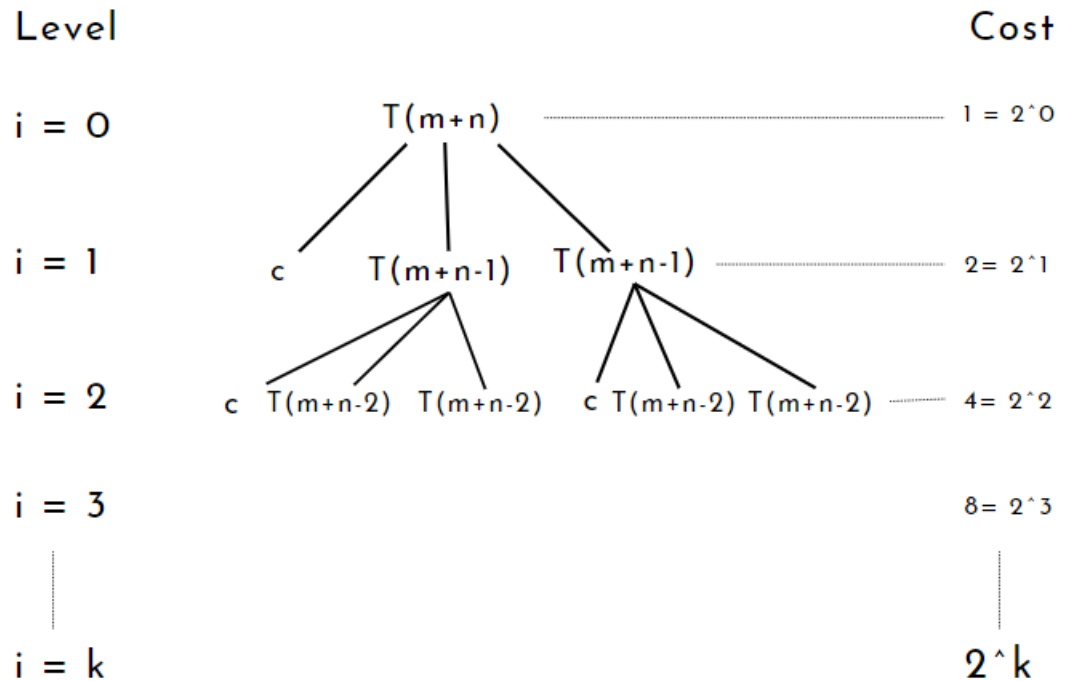
$$T(m+n) = T(m + (n-1)) + T((m-1) + n) + O(1)$$

it is equal to:

$$T(m+n) = 2T(m+n-1) + O(1), m+n \text{ is greater than } 0$$

$$T(m+n) = 1, m+n = 0$$

Let's solve this equation with recursion tree. We put c instead of $O(1)$



When we sum up the costs the total cost is $1 + 2 + 2^2 + 2^3 + \dots + 2^k$

According to the math rules it is equal to $2^{k+1} - 1$

We assume $m + n - k = 0$ then $k = m + n$

Then we get $2^{m+n} - 1$ as a result

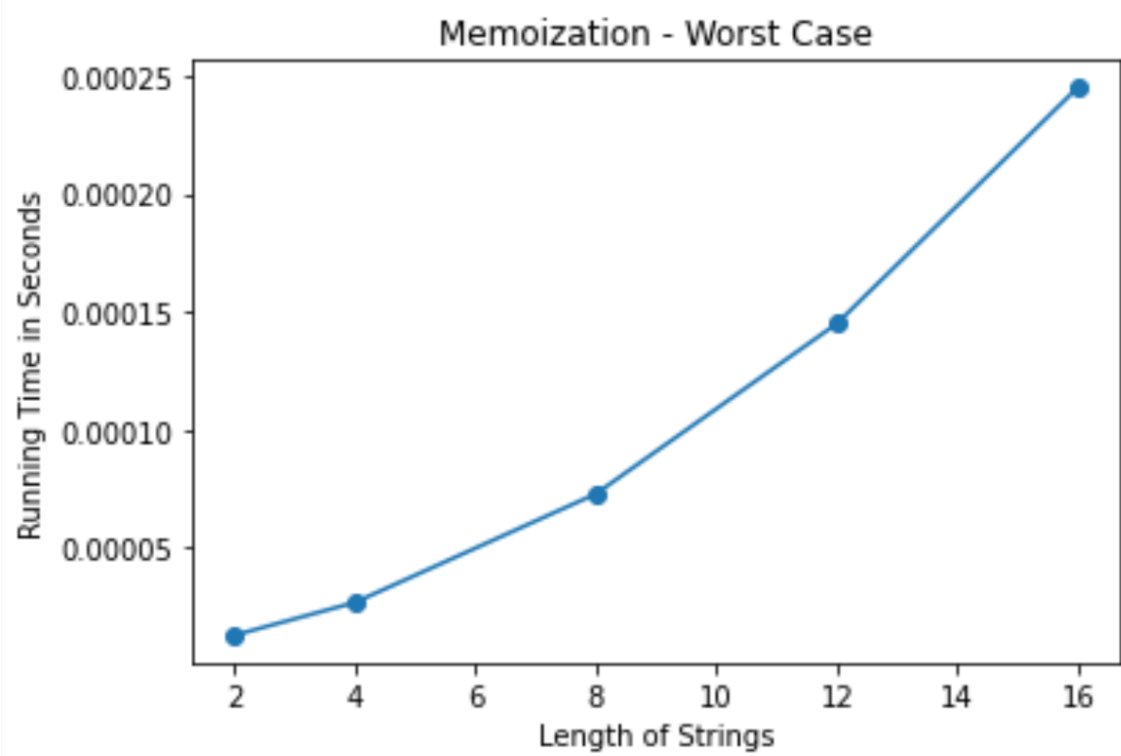
That means our time complexity is $O(2^{m+n})$

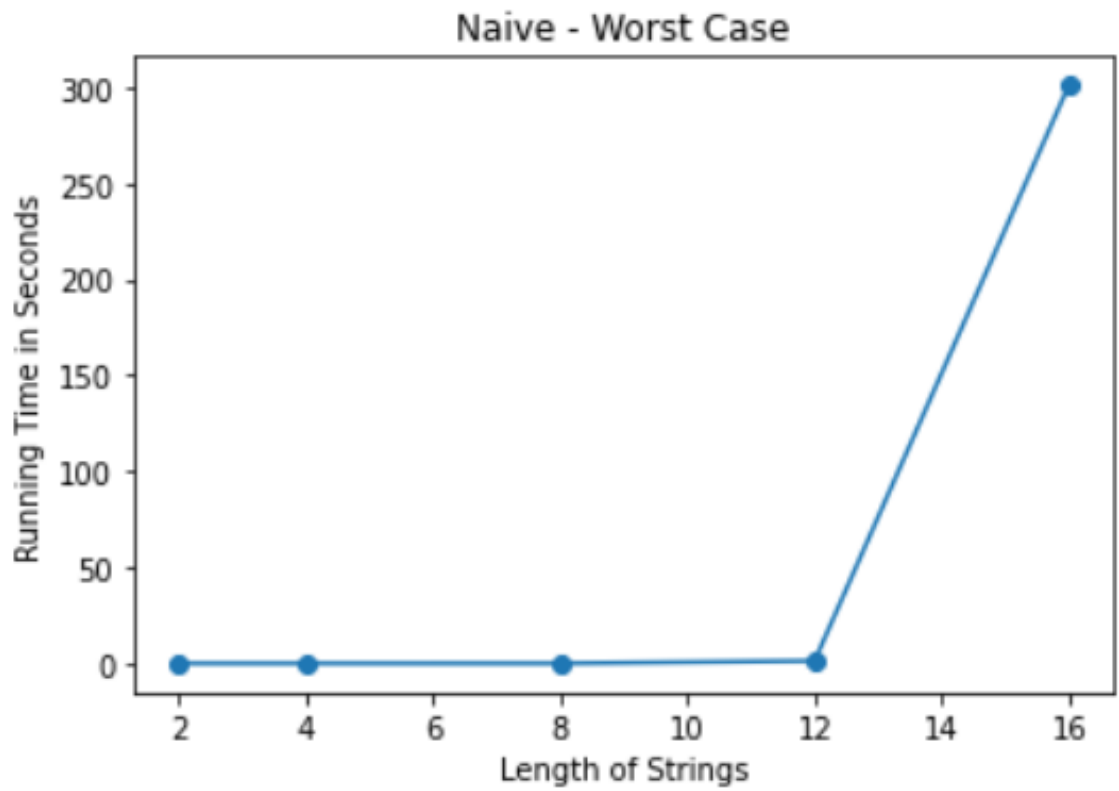
(b)

(i) macOS Monterey, 2,6 GHz 6-Core Intel Core i7, 16 GB RAM

Algorithm	$m = n = 2$	$m = n = 4$	$m = n = 8$	$m = n = 12$	$m = n = 16$
Naive	0.0000081	0.000052	0.007712	1.315021	300.984532
Memoization	0.0000132	0.000027	0.0000731	0.000148	0.000246

(ii)



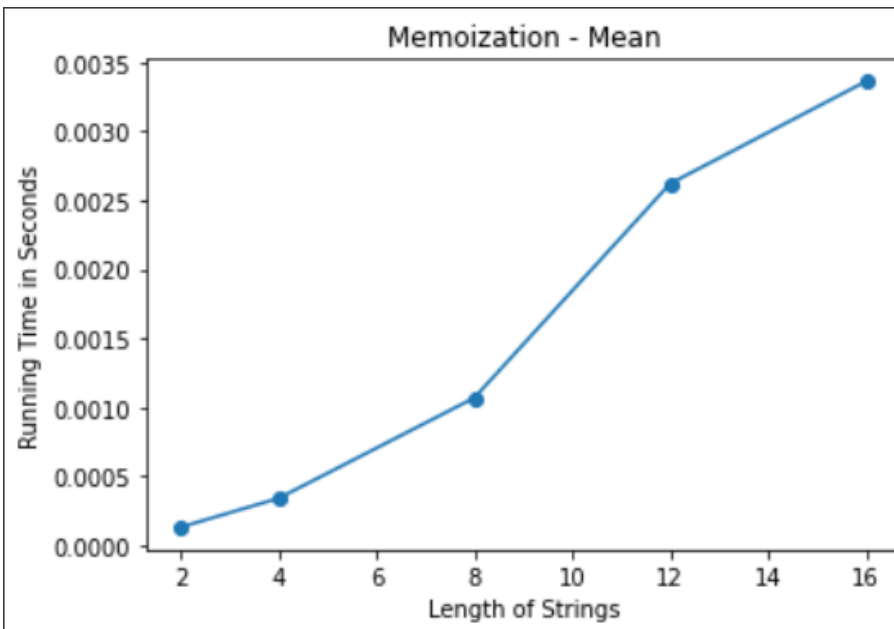
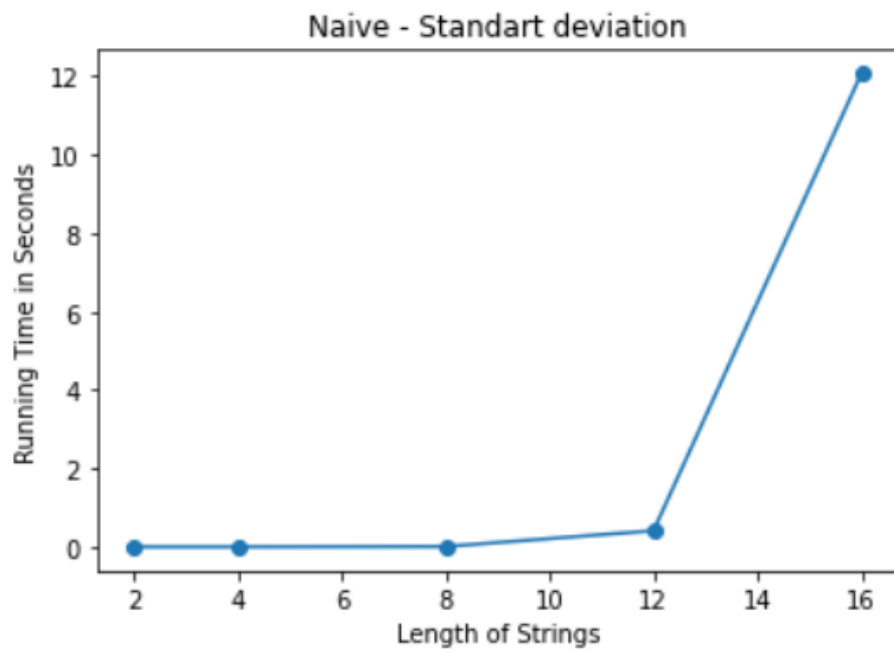


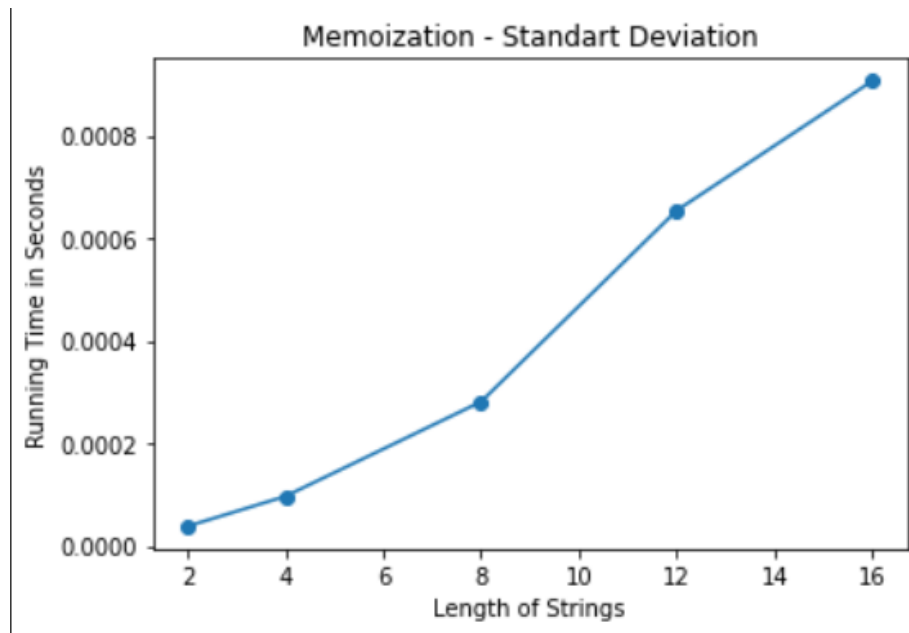
(c)
(i)

Algorithm	m = n = 2		m = n = 4		m = n = 8	
	μ	σ	μ	σ	μ	σ
Naive	0.0000078	0.0000062	0.000043	0.00018	0.00200	0.00829
Memoization	0.000134	0.000038	0.0003439	0.000096	0.001073	0.0002810

Algorithm	m = n = 12		m = n = 16	
	μ	σ	μ	σ
Naive	0.09420	0.410899	3.365674	12.0742688
Memoization	0.002621	0.000654	0.003368	0.000906

(ii)





For this homework I have studied with Ebrar Guler's laptop since my laptop is in service.