

# CPSC-354 Report

Brandon Foley  
Chapman University

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## Abstract

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## 1 Introduction

## 2 Week by Week

### 2.1 Week 1

#### 2.1.1 Homework

Here is my work on the MU Puzzle:

$$MIU \rightarrow MIUIU \rightarrow MIUIUIU \rightarrow \dots$$

$$MI \rightarrow MII \rightarrow MIII \rightarrow MUI \Rightarrow MUI$$

$$MUI \rightarrow MUIU$$

$MUI \rightarrow MUIU \rightarrow MUIUIU \rightarrow MUIUIUIU \rightarrow MUUIU \rightarrow MUUIUIU \rightarrow \dots$

$MUIIU \rightarrow MUIIUIU \rightarrow MUIU$

$MUIIU \rightarrow MUUIU \rightarrow MUIU$

$MUIU \rightarrow MUUIU \rightarrow MUIUIU \rightarrow MUIUIUIU \rightarrow \dots$

$MUIUIUIU \rightarrow MUIUIUIUIU \rightarrow MUIUIUIUIUIU \rightarrow \dots$

The MU Puzzle is unsolvable because it is impossible to reach the string MU starting from MI using the given rules. After a little research and thinking on my part it is the number of I's that make this impossible: rule applications can double the count, add one, or remove three at a time, however, none of these operations ever reduce the number of I's to exactly zero. Since MU has no I's, it can never be derived.

### 2.1.2 Questions

What rules could we add or remove in order to make this possible? I was thinking are there any ways to make this question solvable by adding only one more rule or maybe altering a rule. I think that this is an interesting question and concept.

## 2.2 Week 2

### 2.2.1 Homework

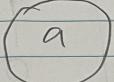
Here is my work on the Rewriting Homework:

## Rewriting

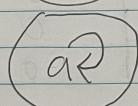
1.  $A = \{3\}$



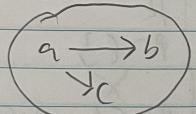
2.  $A = \{a\}$   $R = \{\}$



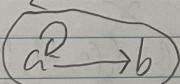
3.  $A = \{a\}$   $R = \{(a, a)\}$



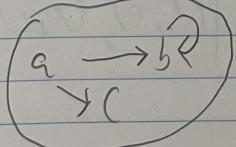
4.  $A = \{a, b, c\}$   $R = \{(a, b), (a, c)\}$



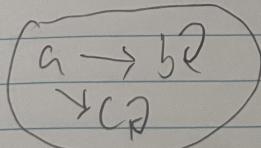
5.  $A = \{a, b\}$   $R = \{(a, a), (a, b)\}$



6.  $A = \{a, b, c\}$   $R = \{(a, b), (b, b), (a, c)\}$



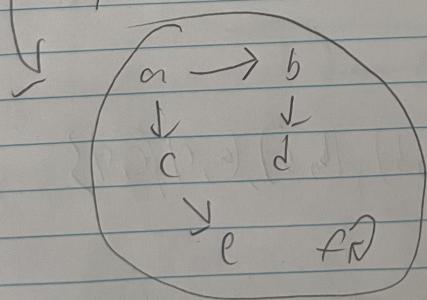
7.  $A = \{a, b, c\}$   $R = \{(a, b), (b, b), (a, c), (c, c)\}$



$0 = \text{false}$   
 $1 = \text{true}$   
 $C = \text{confluent}$   
 $T = \text{terminating}$   
 $U = \text{unique normal form}$

	Confluent	Terminating	Unique normal form
1.	Not		
2.	0	1	
3.	Not 0	1	
4.	0	1	0
5.	1	0	
6.	0	0	
7.	0	0	

	C	+	U	A	R
2.	1	1	1	$A = \{\alpha\}$	$R = \{\beta\}$
				X	X
5.	1	0	1	$A = \{\alpha, \beta\}$	$R = \{\beta_1, \gamma\}, (\alpha, \beta)\}$
				X	X
4.	0	1	1	X	X
6.	0	0	1	$A = \{\alpha, \beta, \gamma\}$	$R = \{(\alpha, \beta), (\beta, \gamma)\}$
				X	X
	0	0	0	$A = \{\alpha, \beta, \gamma, \delta, \epsilon, \zeta\}$	$R = \{(\alpha, \beta), (\beta, \gamma), (\gamma, \delta), (\delta, \epsilon), (\epsilon, \zeta)\}$



## 2.2.2 Questions

When relating to programming, if a system is confluent but not terminating, like in some programs with infinite loops, how does confluence still guarantee that the result of a program is unique when it exists?

**3 Essay**

**4 Evidence of Participation**

**5 Conclusion**

**References**

[BLA] Author, [Title](#), Publisher, Year.