C++ Competitive Programming Library ***DO NOT DISCLOSE OR DISTRIBUTE***

bfs.07 - Bernardo Flores Salmeron

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1. Template

```
#include <bits/stdc++.h>
3 using namespace std;
5 #define INF (111 << 62)
6 #define pb push_back
7 | #define ii pair<int,int>
8 #define OK cerr <<"OK"<< endl
9 | #define debug(x) cerr << #x " = " << (x) << endl
10 #define ff first
11 | #define ss second
12 #define int long long
13 | #define tt tuple<int, int, int>
14 #define endl '\n'
15
16 | signed main () {
17
18
     ios_base::sync_with_stdio(false);
19
     cin.tie(NULL);
20
21 }
```

2. Data Structures

2.1. Arvore Binaria

```
// C program to demonstrate delete operation in binary search tree
2 #include<stdio.h>
3 #include<stdlib.h>
5 struct node {
    int kev;
     struct node *left, *right;
8
10 // A utility function to create a new BST node
11 | struct node *newNode(int item) {
struct node *temp = (struct node *) malloc(sizeof(struct node));
13 temp->key = item;
   temp->left = temp->right = NULL;
14
15
    return temp;
16 }
17
18 // A utility function to do inorder traversal of BST
19 | void inorder(struct node *root) {
   if (root != NULL) {
20
21
       inorder(root->left);
       printf("%d ", root->key);
22
23
       inorder(root->right);
24
25 }
26
27 | /* A utility function to insert a new node with given key in BST */
28 struct node* insert(struct node* node, int key)
29 /* If the tree is empty, return a new node */
    if (node == NULL) return newNode(key);
31
     /★ Otherwise, recur down the tree ★/
32
33
     if (key < node->key)
34
      node->left = insert(node->left, key);
35
36
      node->right = insert(node->right, key);
```

```
38
      /* return the (unchanged) node pointer */
 39
      return node:
 40
 41
 42
    /* Given a non-empty binary search tree, return the node with minimum
 43
     key value found in that tree. Note that the entire tree does not
      need to be searched. */
 44
    struct node * minValueNode(struct node* node) {
 46
      struct node* current = node:
      /* loop down to find the leftmost leaf */
 48
      while (current->left != NULL)
 49
        current = current->left;
 50
51
52
     return current;
53
54
 55
    /* Given a binary search tree and a key, this function deletes the key
     and returns the new root */
    struct node* deleteNode(struct node* root, int key) {
     // base case
      if (root == NULL) return root;
 60
      // If the key to be deleted is smaller than the root's key,
      // then it lies in left subtree
 62
      if (key < root->key)
 63
        root->left = deleteNode(root->left, key);
 64
 65
 66
      // If the key to be deleted is greater than the root's key,
 67
      // then it lies in right subtree
 68
      else if (key > root->key)
 69
        root->right = deleteNode(root->right, key);
 70
71
      // if key is same as root's key, then This is the node
      // to be deleted
 72
 73
      else {
 74
        // node with only one child or no child
 75
        if (root->left == NULL) {
 76
          struct node *temp = root->right;
 77
          free (root);
 78
          return temp;
 79
        } else if (root->right == NULL) {
          struct node *temp = root->left:
          free (root);
 82
          return temp;
 83
 84
 85
        // node with two children: Get the inorder successor (smallest
        // in the right subtree)
 86
        struct node * temp = minValueNode(root->right);
 87
 88
 89
        // Copy the inorder successor's content to this node
 90
        root->key = temp->key;
 91
 92
        // Delete the inorder successor
 93
        root->right = deleteNode(root->right, temp->kev);
94
95
      return root;
96
97
98 // Driver Program to test above functions
99 int main() {
100
    /★ Let us create following BST
```

```
102
                    7.0
103
104
105
         20 40 60 80 */
106
      struct node *root = NULL;
      root = insert(root, 50);
107
      root = insert(root, 30);
      root = insert(root, 20);
109
      root = insert(root, 40);
110
111
      root = insert(root, 70);
      root = insert(root, 60);
      root = insert(root, 80);
113
114
      printf("Inorder traversal of the given tree \n");
115
      inorder(root);
116
117
118
      printf("\nDelete 20\n");
      root = deleteNode(root, 20);
119
120
      printf("Inorder traversal of the modified tree \n");
      inorder(root);
122
123
      printf("\nDelete 30\n");
124
      root = deleteNode(root, 30);
125
      printf("Inorder traversal of the modified tree \n");
126
      inorder(root);
127
      printf("\nDelete 50\n");
128
129
      root = deleteNode(root, 50);
      printf("Inorder traversal of the modified tree \n");
130
131
      inorder(root);
132
133
      return 0;
134 }
```

2.2. Bit.2D

```
1 // INDEX BY ONE ALWAYS!!!
2 class BIT 2D {
    private:
3
4
    // row, column
5
     int n, m;
6
     vector<vector<int>> tree;
9
     // Returns an integer which constains only the least significant bit.
10
     int low(int i) {
        return i & (-i);
11
12
13
      void bit_update(const int x, const int y, const int delta) {
14
15
        for(int i = x; i < n; i += low(i))</pre>
          for(int j = y; j < m; j += low(j))
    this->tree[i][j] += delta;
16
17
18
19
      int bit_query(const int x, const int y) {
20
        int ans = \bar{0};
21
22
        for (int i = x; i > 0; i -= low(i))
23
          for (int j = y; j > 0; j -= low(j))
            ans += this->tree[i][j];
24
25
26
        return ans:
27
28
```

```
public:
     // put the size of the array without 1 indexing.
31
     /// Time Complexity: O(n * m)
32
     BIT_2D(int n, int m) {
33
       this->n = n + 1;
34
       this->m = m + 1;
35
36
       this->tree.resize(n, vector<int>(m, 0));
37
38
39
     /// Time Complexity: O(n * m * (log(n) + log(m)))
40
     BIT_2D(const vector<vector<int>> &mat) {
41
       // Check if it is 1 index.
42
       assert(mat[0][0] == 0);
43
       this->n = mat.size();
44
       this->m = mat.front().size();
45
46
       this->tree.resize(n, vector<int>(m, 0));
47
       for(int i = 1; i < n; i++)</pre>
48
         for (int j = 1; j < m; j++)
49
           update(i, j, mat[i][j]);
50
51
52
     /// Query from (1, 1) to (x, y).
53
54
     /// Time Complexity: O(log(n) + log(m))
55
     int prefix_query(const int x, const int y) {
       assert (0 < x); assert (x < this -> n);
56
57
       assert(0 < y); assert(y < this->m);
58
59
       return bit_query(x, y);
60
61
62
     /// Query from (x1, y1) to (x2, y2).
63
64
     /// Time Complexity: O(log(n) + log(m))
65
     int query (const int x1, const int y1, const int x2, const int y2) {
66
       assert (0 < x1); assert (x1 <= x2); assert (x2 < this->n);
67
       assert(0 < y1); assert(y1 <= y2); assert(y2 < this->m);
68
69
       return bit_query(x2, y2) - bit_query(x1 - 1, y2) - bit_query(x2, y1 - 1)
       + bit_query(x1 - 1, y1 - 1);
70
71
72
     /// Updates point (x, y).
73
     ///
74
     /// Time Complexity: O(log(n) + log(m))
75
     void update(const int x, const int y, const int delta) {
76
       assert (0 < x); assert (x < this->n);
77
       assert(0 < y); assert(y < this->m);
78
79
       bit_update(x, y, delta);
80
81
   } ;
```

2.3. Merge Sort Tree (K-Esimo Maior Elemento Num Intervalo, Valores Maiores Que K Num Intervalo,

```
// retornar a qtd de números maiores q um numero k numa array de i...j
struct Tree {
    vector<int> vet;
};
Tree tree[4*(int)3e4];
int arr[(int)5e4];
```

```
8 | int query(int 1, int r, int i, int j, int k, int pos) {
     if(1 > j || r < i)
10
       return 0;
11
12
     if(i <= 1 && r <= j) {
13
       auto it = upper_bound(tree[pos].vet.begin(), tree[pos].vet.end(), k);
14
       return tree[pos].vet.end()-it;
15
16
17
     int mid = (1+r) >> 1:
18
     return query(1, mid, i, j, k, 2*pos+1) + query(mid+1,r,i,j,k,2*pos+2);
19
20
21
   void build(int 1, int r, int pos) {
22
23
     if(1 == r) {
24
       tree[pos].vet.pb(arr[1]);
25
       return;
26
2.7
28
     int mid = (1+r) >> 1:
2.9
     build(l, mid, 2*pos+1);
     build(mid + 1, r, 2*pos+2);
31
32
     merge(tree[2*pos+1].vet.begin(), tree[2*pos+1].vet.end(),
        tree[2*pos+2].vet.begin(), tree[2*pos+2].vet.end(),
        back_inserter(tree[pos].vet));
33 }
```

2.4. Mos Algorithm

```
1 struct Tree {
    int l, r, ind:
3 };
4 | Tree query[311111];
5 | int arr[311111];
6 | int freq[1111111];
   int ans[311111];
   int block = sqrt(n), cont = 0;
10 | bool cmp(Tree a, Tree b) {
    if(a.1/block == b.1/block)
11
12
       return a.r < b.r;</pre>
13
     return a.l/block < b.l/block;
14
15
16 void add(int pos) {
17
     freq[arr[pos]]++;
18
     if(freq[arr[pos]] == 1) {
19
       cont++;
20
21
22
   void del(int pos) {
     freg[arr[pos]]--;
     if(freq[arr[pos]] == 0)
25
        cont--;
26
27 | int main () {
28
    int n; cin >> n;
29
     block = sqrt(n);
30
31
     for(int i = 0; i < n; i++) {</pre>
32
       cin >> arr[i];
```

```
freq[arr[i]] = 0;
34
35
36
     int m; cin >> m;
37
38
     for(int i = 0; i < m; i++) {</pre>
39
        cin >> query[i].l >> query[i].r;
40
        query[i].l--, query[i].r--;
        query[i].ind = i;
41
42
43
     sort (query, query + m, cmp);
44
45
     int s,e;
46
     s = e = query[0].1;
47
     add(s);
     for(int i = 0; i < m; i++) {</pre>
48
49
        while(s > query[i].l)
50
          add (--s);
51
        while(s < query[i].l)</pre>
52
          del(s++);
53
        while(e < query[i].r)</pre>
54
          add(++e);
55
        while(e > query[i].r)
56
          del(e--);
57
        ans[query[i].ind] = cont;
58
59
60
     for(int i = 0; i < m; i++)</pre>
61
        cout << ans[i] << endl;</pre>
62
```

2.5. Ordenacao De Estruturas (Pq, Etc)

```
struct cmp {
  bool operator(ii a, ii b) {
  //ordena primeiro pelo first(decrescente), dps pelo second(crescente)
  if(a.first == b.first)
    return a.second < b.second;
  return a.first > b.first;
}

Ex: pq<ii,vector<ii>,cmp> fila;
```

2.6. Ordered Set (Policy Based Data Structures)

```
#include <bits/stdc++.h>
   #include <ext/pb_ds/assoc_container.hpp>
   #include <ext/pb_ds/trie_policy.hpp>
   using namespace std;
   using namespace __gnu_pbds;
   typedef tree<
9 int,
10 null_type,
11 less<int>,
12 rb tree tag,
13 tree_order_statistics_node_update>
14 ordered_set;
15
16 ordered_set X;
17 X.insert(1); X.insert(2);
18 X.insert(4); X.insert(8);
```

2.7. Sqrt Decomposition

```
1 // Problem: Sum from 1 to r
2 // Ver MO'S ALGORITHM
4 int getId(int indx, int blockSZ) {
       return indx/blockSZ;
7 | void init(int sz) {
8
      for(int i=0; i<=sz; i++)
    BLOCK[i]=inf;
10
11 | int query(int left, int right) {
12 | int startBlockIndex=left/sqrt;
13 | int endIBlockIndex = right / sqrt;
14 int sum = 0;
15 | for (int i = startBlockIndex + 1; i < endIBlockIndex; i++) {
16
           sum += blockSums[i];
17
18 for(i=left...(startBlockIndex*BLOCK_SIZE-1))
   sum += a[i];
19
20
   for(j = endIBlockIndex*BLOCK SIZE ... right)
    sum += a[i];
21
22 }
```

2.8. Bit

```
/// INDEX THE ARRAY BY 1!!!
   class BIT {
   private:
     vector<int> bit;
     int n:
     int low(const int i) { return (i & (-i)); }
     // point update
10
     void bit_update(int i, const int delta) {
11
12
       while (i <= this->n) {
13
         this->bit[i] += delta:
14
         i += this->low(i);
15
16
17
18
     // point query
19
     int bit_query(int i) {
20
       int sum = 0;
21
       while (i > 0) {
```

```
sum += bit[i];
23
         i -= this->low(i);
24
25
       return sum:
26
27
28
   public:
     BIT(const vector<int> &arr) { this->build(arr); }
31
     BIT(const int n)
       // OBS: BIT IS INDEXED FROM 1
32
33
       // THE USE OF 1-BASED ARRAY IS RECOMMENDED
34
       this->n = n;
35
       this->bit.resize(n + 1, 0);
36
37
     // build the bit
38
39
     void build(const vector<int> &arr) {
40
       // OBS: BIT IS INDEXED FROM 1
41
       // THE USE OF 1-BASED ARRAY IS RECOMMENDED
42
       assert(arr.front() == 0);
43
       this->n = (int)arr.size() - 1;
44
       this->bit.resize(arr.size(), 0);
45
46
       for (int i = 1; i <= this->n; i++)
         this->bit_update(i, arr[i]);
47
48
49
50
     // point update
51
     void update(const int i, const int delta) {
52
       assert(1 <= i), assert(i <= this->n);
53
       this->bit_update(i, delta);
54
55
56
     // point query
57
     int query(const int i) {
       assert(1 <= i), assert(i <= this->n);
58
59
       return this->bit_query(i);
60
61
62
     // range guery
     int query(const int 1, const int r) {
       assert(1 \leq 1), assert(1 \leq r), assert(r \leq this\rightarrown);
64
       return this->bit_query(r) - this->bit_query(l - 1);
66
67
   };
```

2.9. Bit (Range Update)

```
/// INDEX THE ARRAY BY 1!!!
2 class BIT {
    private:
     vector<int> bit1:
     vector<int> bit2;
     int n:
    private:
9
     int low(int i) {
10
       return (i & (-i));
11
12
13
     // point update
14
     void update(int i, const int delta, vector<int> &bit) {
       while(i <= this->n) {
```

```
bit[i] += delta;
17
         i += this->low(i);
18
19
20
21
     // point query
22
     int query(int i, const vector<int> &bit) {
       int sum = 0;
23
24
       while (i > 0)
         sum += bit[i];
2.5
26
          i -= this->low(i);
27
28
       return sum;
29
30
31
     // build the bit
32
     void build(const vector<int> &arr) {
33
       // OBS: BIT IS INDEXED FROM 1
34
       // THE USE OF 1-BASED ARRAY IS MANDATORY
35
       assert(arr.front() == 0);
36
       this->n = (int)arr.size() - 1;
37
       this->bit1.resize(arr.size(), 0);
38
       this->bit2.resize(arr.size(), 0);
39
40
       for(int i = 1; i <= this->n; i++)
41
         this->update(i, arr[i]);
42
43
44
    public:
45
     BIT(const vector<int> &arr) {
46
       this->build(arr);
47
48
49
     BIT(const int n) {
       // OBS: BIT IS INDEXED FROM 1
50
        // THE USE OF 1-BASED ARRAY IS MANDATORY
51
52
       this->n = n:
       this->bit1.resize(n + 1, 0);
53
       this->bit2.resize(n + 1, 0);
54
55
56
57
     // range update
     void update(const int 1, const int r, const int delta) {
59
       assert(1 \leq 1); assert(1 \leq r); assert(r \leq this->n);
60
       this->update(1, delta, this->bit1);
61
       this->update(r + 1, -delta, this->bit1);
62
       this->update(l, delta * (l - 1), this->bit2);
63
       this->update(r + 1, -delta * r, this->bit2);
64
65
66
     // point update
67
     void update(const int i, const int delta) {
68
       assert(1 <= i); assert(i <= this->n);
69
       this->update(i, i, delta);
70
71
72
     // range guery
73
     int query(const int 1, const int r) {
       assert(1 <= 1); assert(1 <= r); assert(r <= this->n);
74
75
       return this->query(r) - this->query(l - 1);
76
77
78
     // point query
     int query(const int i) {
79
       assert(1 <= i); assert(i <= this->n);
```

23

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79 80

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83

84 85

```
return (this->query(i, this->bit1) * i) - this->query(i, this->bit2);
 82
 83
    };
 84
   // TESTS
 85
 86 // signed main()
87
   // {
 88
89
   // vector<int> input = \{0,1,2,3,4,5,6,7\};
 90
 91
    // BIT ft(input);
 92
 93
    // assert (1 == ft.guerv(1));
    // assert (3 == ft.query(2));
    // assert (6 == ft.query(3));
    // assert (10 == ft.query(4));
   // assert (15 == ft.query(5));
98 // assert (21 == ft.query(6));
99 // assert (28 == ft.guerv(7));
100 // assert (12 == ft.querv(3,5));
101 // assert (21 == ft.query(1,6));
102 // assert (28 == ft.guerv(1,7));
103 // }
```

2.10. Counting Inversions (Minimum Number Of Adjacent Swaps To Sort Array)

```
// REQUIRES bit.cpp!!
// REQUIRES point_compresion.cpp!!
int count_inversions(vector<int> &arr) {
    arr = compress(arr);
    int ans = 0;
    BIT bit(arr.size());
    for (int i = arr.size() - 1; i > 0; --i) {
        ans += bit.query(arr[i] - 1);
        bit.update(arr[i], 1);
    }

return ans;
}
```

2.11. Persistent Segment Tree

```
class Persistent Seg Tree {
     struct Node {
3
       int val;
4
       Node *left, *right;
       Node() {}
       Node (int v, Node \star1, Node \starr) : val(v), left(l), right(r) {}
   #define NEUTRAL_NODE Node(0, nullptr, nullptr);
     Node _NEUTRAL_NODE = Node(0, nullptr, nullptr);
10
11
     int merge_nodes(const int x, const int y) { return x + y; }
12
13
14 private:
     vector<Node *> version;
17
18 public:
    Persistent_Seq_Tree() { this->n = -1; }
1 9
     /// Builds version[0] with the values in the array.
21
    ///
```

```
/// Time complexity: O(n)
Node *pst build (Node *node, const int 1, const int r,
                const vector<int> &arr) {
  node = new NEUTRAL_NODE;
 if (1 == r) {
   node->val = arr[l];
   return node:
 int mid = (1 + r) / 2;
 node->left = pst_build(node->left, 1, mid, arr);
 node->right = pst_build(node->right, mid + 1, r, arr);
 node->val = merge nodes(node->left->val, node->right->val);
  return node;
/// Builds version[0] with 0.
111
/// Time complexity: O(n)
Node *pst_build_empty(Node *node, const int 1, const int r) {
 node = new NEUTRAL NODE;
  if (l == r)
   return node:
  int mid = (1 + r) / 2;
 node->left = pst_build_empty(node->left, 1, mid);
 node->right = pst_build_empty(node->right, mid + 1, r);
 node->val = merge_nodes(node->left->val, node->right->val);
 return node;
Node *pst_update(Node *cur_tree, Node *prev_tree, const int 1, const int r,
                 const int idx, const int delta) {
  if (1 > idx || r < idx)
   if (cur tree != nullptr)
     return cur tree:
    return prev tree;
  if (cur tree == nullptr)
   cur_tree = new Node(prev_tree->val, prev_tree->left, prev_tree->right);
    cur_tree = new Node(cur_tree->val, cur_tree->left, cur_tree->right);
  if (l == r) {
   cur tree->val += delta;
    return cur tree;
 int mid = (1 + r) / 2;
 cur tree->left =
      pst_update(cur_tree->left, prev_tree->left, 1, mid, idx, delta);
 cur_tree->right =
     pst_update(cur_tree->right, prev_tree->right, mid + 1, r, idx,
  cur_tree->val = merge_nodes(cur_tree->left->val, cur_tree->right->val);
  return cur tree;
int pst_query(Node *node, const int 1, const int r, const int i,
              const int j) {
  if (1 > j || r < i)
   return _NEUTRAL_NODE.val;
  if (i <= 1 && r <= j)
```

```
86
          return node->val:
 87
 88
        int mid = (1 + r) / 2;
 89
        return merge_nodes(pst_query(node->left, 1, mid, i, j),
 90
                            pst query(node->right, mid + 1, r, i, j));
 91
 92
 93
    public:
      Persistent_Seq_Tree(const int n, const int number_of_versions) {
 94
 95
        this -> n = n:
 96
        version.resize(number of versions);
        this->version[0] = this->pst_build_empty(this->version[0], 0, this->n -
 97
        1);
 98
 99
100
      /// Constructor that allows to pass initial values to the leafs.
101
      Persistent_Seg_Tree(const vector<int> &arr, const int number_of_versions) {
102
        this->n = arr.size();
103
        version.resize(number of versions);
104
        this->version[0] = this->pst_build(this->version[0], 0, this->n - 1,
105
106
      /// Links the root of a version to a previous version.
107
108
109
      /// Time Complexity: O(1)
      void link(const int version, const int prev_version) {
110
        assert (this -> n > -1);
111
112
        assert(0 <= prev_version);</pre>
113
        assert (prev version <= version);
114
        assert(version < this->version.size());
115
        this->version[version] = this->version[prev_version];
116
117
      /// Updates an index in cur_tree based on prev_tree with a delta.
118
119
      /// Time Complexity: O(log(n))
120
121
      void update (const int cur_version, const int prev_version, const int idx,
122
                   const int delta) {
123
        assert (this->n > -1);
124
        assert (0 <= prev version);
125
        assert (prev version <= cur version);
        assert(cur version < this->version.size());
126
        this->version[cur_version] = this->pst_update(this->version[cur_version],
127
128
        this->version[prev_version],
129
                                                        0, this->n - 1, idx,
        delta);
130
1.31
132
      /// Ouery from 1 to r.
133
      ///
134
      /// Time Complexity: O(log(n))
135
      int query(const int version, const int 1, const int r) {
        assert (this->n > -1);
136
137
        assert(this->version[version] != nullptr);
138
        assert (0 \le 1):
139
        assert(1 <= r);
140
        assert(r < this->n);
141
        return this->pst_query(this->version[version], 0, this->n - 1, 1, r);
142
143
    };
```

```
class Seq Tree {
   public:
3
     struct Node {
       int val, lazy;
7
       Node (const int val, const int lazy) : val(val), lazy(lazy) {}
9
10
   private:
11
    // // range sum
12
     // Node NEUTRAL_NODE = Node(0, 0);
13
     // Node merge nodes (const Node &x, const Node &v) {
14
     // return Node(x.val + y.val, 0);
1.5
     // }
16
     // void apply_lazy(const int 1, const int r, const int pos) {
17
     // tree[pos].val += (r - l + 1) * tree[pos].lazy;
1.8
19
20
     // // RMO max
21
     // Node NEUTRAL NODE = Node(-INF, 0);
     // Node merge_nodes (const Node &x, const Node &y) {
22
23
     // return Node (max (x.val, y.val), 0);
24
     // }
25
     // void apply_lazy(const int l, const int r, const int pos) {
26
     // tree[pos].val += tree[pos].lazy;
2.7
     11 }
28
29
     // // RMO min
3.0
     // Node NEUTRAL NODE = Node(INF, 0);
31
     // Node merge nodes (const Node &x, const Node &v) {
32
          return Node (min (x.val, y.val), 0);
     // }
33
     // void apply_lazy(const int 1, const int r, const int pos) {
34
35
         tree[pos].val += tree[pos].lazy;
     // }
36
37
38
39
     // Only works with point updates
40
     // Node NEUTRAL_NODE = Node(0, 0);
41
     // Node merge nodes (const Node &x, const Node &v) {
42
         return Node(x.val ^ v.val, 0);
43
44
     // void apply lazy(const int l, const int r, const int pos) {}
45
   private:
46
47
     int n;
48
   public:
49
     vector<Node> tree;
50
51
52 private:
     void st propagate(const int 1, const int r, const int pos) {
       if (tree[pos].lazy != 0) {
55
         apply_lazy(l, r, pos);
         if (1 != r) {
56
57
           tree[2 * pos + 1].lazy += tree[pos].lazy;
58
           tree[2 * pos + 2].lazy += tree[pos].lazy;
59
60
         tree[pos].lazy = 0;
61
62
63
     Node st build(const int 1, const int r, const vector<int> &arr,
```

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188

189 190

```
const int pos) {
 66
         if (1 == r)
 67
          return tree[pos] = Node(arr[1], 0);
 68
 69
        int mid = (1 + r) / 2;
 70
        return tree[pos] = merge_nodes(st_build(1, mid, arr, 2 * pos + 1),
 71
                                         st_build(mid + 1, r, arr, 2 * pos + 2));
 72
 73
 74
      int st_get_first(const int 1, const int r, const int v, const int pos) {
 75
        st_propagate(l, r, pos);
 76
 77
        // Needs RMO MAX
        // Replace to <= for greater or equal or (with RMQ MIN) > for smaller or
 78
 79
         // equal or >= for smaller
 80
        if (tree[pos].val < v)</pre>
 81
          return -1;
 82
 83
        if (1 == r)
 84
          return 1:
 85
 86
        int mid = (1 + r) / 2;
 87
        int aux = st_get_first(1, mid, v, 2 * pos + 1);
 88
        if (aux !=-\overline{1})
 89
           return aux;
        return st_get_first(mid + 1, r, v, 2 * pos + 2);
 90
 91
 92
 93
      Node st_query(const int 1, const int r, const int i, const int j,
 94
                     const int pos) {
 9.5
        st_propagate(1, r, pos);
 96
 97
        if (1 > r || 1 > j || r < i)
 98
          return NEUTRAL NODE;
 99
100
        if (i <= 1 && r <= i)
101
          return tree[pos];
102
103
        int mid = (1 + r) / 2;
104
        return merge_nodes(st_query(l, mid, i, j, 2 * pos + 1),
105
                            st_query(mid + 1, r, i, j, 2 * pos + 2));
106
107
108
      // it adds a number delta to the range from i to j
109
      Node st_update(const int 1, const int r, const int i, const int j,
110
                      const int delta, const int pos) {
111
        st_propagate(l, r, pos);
112
113
        if (1 > r || 1 > i || r < i)
          return tree[pos];
114
115
116
        if (i <= 1 && r <= i) {
          tree[pos].lazy = delta;
117
           st_propagate(1, r, pos);
118
          return tree[pos];
119
120
121
122
        int mid = (1 + r) / 2;
123
        return tree[pos] =
                    merge_nodes(st_update(l, mid, i, j, delta, 2 * pos + 1),
124
125
                                 st\_update(mid + 1, r, i, j, delta, 2 * pos + 2));
126
127
128
      void build(const vector<int> &arr) {
129
        this->n = arr.size();
```

```
130
        this->tree.resize(4 * this->n);
        this->st build(0, this->n - 1, arr, 0);
134 public:
      /// N equals to -1 means the Segment Tree hasn't been created yet.
      Seg Tree(): n(-1) {}
      /// Constructor responsible initialize a tree with 0.
      /// Time Complexity O(n)
      Seg_Tree (const int n) : n(n) { this->tree.resize(4 * this->n, Node(0, 0));
      /// Constructor responsible for building the initial tree based on a
        vector.
      /// Time Complexity O(n)
      Seg_Tree(const vector<int> &arr) { this->build(arr); }
      /// Returns the first index from left to right.
      /// Uncomment the line in the original funtion to get the proper element
      /// may be: GREATER OR EQUAL, GREATER, SMALLER OR EQUAL, SMALLER.
      /// Time Complexity O(log n)
      int get first(const int v) {
        assert (this->n >= 0);
        return this->st_get_first(0, this->n - 1, v, 0);
      /// Update at a single index.
      /// Time Complexity O(log n)
      void update(const int idx, const int delta) {
        assert(this->n >= 0);
        assert(0 <= idx), assert(idx < this->n);
        this->st_update(0, this->n - 1, idx, idx, delta, 0);
      /// Range update from 1 to r.
      ///
      /// Time Complexity O(log n)
      void update(const int l, const int r, const int delta) {
        assert (this->n >= 0);
        assert(0 \le 1), assert(1 \le r), assert(r < this->n);
        this->st update(0, this->n - 1, 1, r, delta, 0);
      /// Query at a single index.
      /// Time Complexity O(log n)
      int query(const int idx) {
        assert(this->n >= 0);
        assert(0 <= idx), assert(idx < this->n);
        return this->st_query(0, this->n - 1, idx, idx, 0).val;
      /// Range query from 1 to r.
      /// Time Complexity O(log n)
      int query(const int 1, const int r) {
        assert (this->n >= 0);
        assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
191
        return this->st_query(0, this->n - 1, 1, r, 0).val;
```

```
192 };
```

2.13. Segment Tree 2D

```
1 // REOUIRES segment tree.cpp!!
  class Seg_Tree_2d {
    private:
     // // range sum
     // int NEUTRAL_VALUE = 0;
     // int merge_nodes(const int &x, const int &y) {
     // return x + y;
8
     // }
9
10
     // // RMO max
11
     // int NEUTRAL_VALUE = -INF;
     // int merge_nodes(const int &x, const int &y) {
12
13
     // return max(x, y);
14
15
     // // RMQ min
16
17
     // int NEUTRAL VALUE = INF;
     // int merge_nodes(const int &x, const int &y) {
18
19
     // return min(x, y);
20
     // }
21
22
    private:
23
     int n, m;
24
    public:
     vector<Seq_Tree> tree;
27
28
    private:
     void st_build(const int 1, const int r, const int pos, const
       vector<vector<int>> &mat) {
30
       if(1 == r)
31
         tree[pos] = Seq_Tree(mat[1]);
32
       else {
33
         int mid = (1 + r) / 2;
34
         st build(1, mid, 2*pos + 1, mat);
35
         st build(mid + 1, r, 2*pos + 2, mat);
         for(int i = 0; i < tree[2*pos + 1].tree.size(); i++)</pre>
36
37
           tree[pos].tree[i].val = merge_nodes(tree[2*pos + 1].tree[i].val,
38
                                                 tree[2*pos + 2].tree[i].val);
39
40
41
42
     int st_query(const int 1, const int r, const int x1, const int y1, const
       int x2, const int y2, const int pos) {
43
       if(1 > x2 | | r < x1)
44
         return NEUTRAL VALUE;
45
46
       if(x1 \le 1 \&\& r \le x2)
47
         return tree[pos].query(y1, y2);
48
49
       int mid = (1 + r) / 2;
50
       return merge_nodes(st_query(1, mid, x1, y1, x2, y2, 2*pos + 1),
51
                           st_query(mid + 1, r, x1, y1, x2, y2, 2*pos + 2));
52
53
     void st_update(const int 1, const int r, const int x, const int y, const
       int delta, const int pos) {
55
       if(1 > x | | r < x)
56
         return;
```

```
58
        // Only supports point updates.
59
        if(1 == r) {
60
          tree[pos].update(y, delta);
61
          return;
62
63
64
        int mid = (1 + r) / 2;
        st\_update(1, mid, x, y, delta, 2*pos + 1);
65
66
        st\_update(mid + 1, r, x, y, delta, 2*pos + 2);
67
        tree[pos].update(y, delta);
68
69
70
     public:
71
      Seq_Tree_2d() {
72
        this->n = -1;
73
        this->m = -1;
74
75
76
      Seq Tree_2d(const int n, const int m) {
77
        this \rightarrow n = n:
78
        this->m = m;
79
        // MAY TLE IN BUILD, TEST IT OR UPDATE EACH NODE MANUALLY!
80
        assert (m < 10000):
        tree.resize(4 * n, Seg_Tree(m));
81
82
83
84
      Seg_Tree_2d(const int n, const int m, const vector<vector<int>> &mat) {
85
        this->n = n:
        this->m = m;
86
87
        // MAY TIE IN BUILD. TEST IT OR UPDATE EACH NODE MANUALLY!
88
        assert (m < 10000):
89
        tree.resize(4 * n, Seg Tree(m));
90
        st_build(0, n - 1, 0, mat);
91
92
93
      // Query from (x1, y1) to (x2, y2).
94
95
      // Time complexity: O((log n) * (log m))
96
      int query (const int x1, const int y1, const int x2, const int y2) {
        assert (this->n > -1):
        assert(0 \le x1); assert(x1 \le x2); assert(x2 < this -> n);
        assert(0 \leq y1); assert(y1 \leq y2); assert(y2 \leq this->n);
100
        return st_query(0, this->n - 1, x1, y1, x2, y2, 0);
101
102
103
      // Point updates on position (x, y).
104
      // Time complexity: O((log n) * (log m))
105
106
      void update(const int x, const int y, const int delta) {
107
        assert (0 \leq x); assert (x \leq this\rightarrown);
108
        assert(0 <= y); assert(y < this->n);
109
        st update(0, this->n - 1, x, v, delta, 0);
110
111 };
```

2.14. Segment Tree Polynomial

```
/// Works for the polynomial f(x) = z1*x + z0
class Seg_Tree {
public:
    struct Node {
    int val, z1, z0;
}
```

```
8
       Node (const int val, const int z1, const int z0)
9
           : val(val), z1(z1), z0(z0) {}
10
11
12 private:
13
     // range sum
     Node NEUTRAL_NODE = Node(0, 0, 0);
14
     Node merge_nodes(const Node &x, const Node &y) {
15
16
       return Node (x.val + v.val, 0, 0);
17
     void apply_lazy(const int 1, const int r, const int pos) {
18
19
       tree[pos].val += (r - l + 1) * tree[pos].z0;
       tree[pos].val += (r - 1) * (r - 1 + 1) / 2 * tree[pos].z1;
20
21
22
23
  private:
    int n:
25
26
  public:
    vector<Node> tree:
28
    void st_propagate(const int 1, const int r, const int pos) {
       if (tree[pos].z0 != 0 || tree[pos].z1 != 0) {
32
         apply lazy(l, r, pos);
         int mid = (1 + r) / 2;
33
34
         int sz_left = mid - 1 + 1;
         if (l != r) {
35
36
           tree[2 * pos + 1].z0 += tree[pos].z0;
37
           tree[2 * pos + 1].z1 += tree[pos].z1;
38
39
           tree[2 * pos + 2].z0 += tree[pos].z0 + sz_left * tree[pos].z1;
40
           tree[2 * pos + 2].z1 += tree[pos].z1;
41
42
         tree[posl.z0 = 0;
43
         tree[pos].z1 = 0;
44
45
46
47
     Node st_build(const int 1, const int r, const vector<int> &arr,
48
                   const int pos) {
49
       if (1 == r)
         return tree[pos] = Node(arr[1], 0, 0);
50
51
52
       int mid = (1 + r) / 2;
53
       return tree[pos] = merge_nodes(st_build(1, mid, arr, 2 * pos + 1),
54
                                       st\_build(mid + 1, r, arr, 2 * pos + 2));
55
56
57
     Node st_query(const int 1, const int r, const int i, const int j,
58
                   const int pos) {
59
       st_propagate(1, r, pos);
60
       if (l > r || l > j || r < i)
61
62
         return NEUTRAL NODE;
63
64
       if (i <= 1 && r <= j)
65
         return tree[pos];
66
67
       int mid = (1 + r) / 2;
68
       return merge_nodes(st_query(l, mid, i, j, 2 * pos + 1),
69
                          st_query(mid + 1, r, i, j, 2 * pos + 2));
70
71
```

```
// it adds a number delta to the range from i to j
73
      Node st update (const int 1, const int r, const int i, const int j,
74
                     const int z1, const int z0, const int pos) {
75
        st_propagate(1, r, pos);
76
77
        if (1 > r || 1 > j || r < i)
78
          return tree[pos];
79
        if (i <= l && r <= j) {
80
81
          tree[pos].z0 = (1 - i + 1) * z0;
82
          tree[pos].z1 = z1;
          st_propagate(1, r, pos);
83
84
          return tree[posl:
85
86
87
        int mid = (1 + r) / 2;
88
        return tree[pos] =
89
                   merge_nodes(st_update(1, mid, i, j, z1, z0, 2 * pos + 1),
90
                                st\_update(mid + 1, r, i, j, z1, z0, 2 * pos + 2));
91
92
   public:
93
      Seg_Tree() : n(-1) {}
      Seg_Tree (const int n) : n(n) { this->tree.resize(4 * this->n, Node(0, 0));
97
      Seq_Tree(const vector<int> &arr) { this->build(arr); }
98
99
100
      void build(const vector<int> &arr) {
101
        this->n = arr.size();
102
        this->tree.resize(4 * this->n);
103
        this->st_build(0, this->n - 1, arr, 0);
104
106
      /// Index update of a polynomial f(x) = z1*x + z0
107
      /// Time Complexity O(log n)
108
      void update(const int i, const int z1, const int z0) {
109
        assert (this->n >= 0);
110
        assert(0 <= i), assert(i < this->n);
111
        this->st_update(0, this->n - 1, i, i, z1, z0, 0);
112
113
114
115
      /// Range update of a polynomial f(x) = z1 \star x + z0 from 1 to r
116
      ///
      /// Time Complexity O(log n)
117
      void update(const int 1, const int r, const int z1, const int z0) {
118
119
        assert (this->n >= 0);
        assert (0 \le 1), assert (1 \le r), assert (r < this->n);
120
        this->st_update(0, this->n - 1, 1, r, z1, z0, 0);
121
122
123
124
      /// Range sum query from 1 to r
125
      /// Time Complexity O(log n)
126
      int querv(const int 1, const int r) {
127
128
        assert (this->n >= 0);
129
        assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
130
        return this->st query(0, this->n - 1, 1, r, 0).val;
131
132 };
```

2.15. Sparse Table

```
// RMO min implementation
   class Sparse_Table {
    private:
     int n;
     vector<vector<int> > table;
     vector<int> lq;
     /// lg[i] represents the log2(i)
     void build_log_array() {
9
10
       lq.resize(this->n + 1);
11
12
       for(int i = 2; i <= this->n; i++)
13
         lq[i] = lq[i/2] + 1;
14
15
16
     /// Time Complexity: O(n*log(n))
     /// Space Complexity: O(n*log(n))
17
     void build_sparse_table(const vector<int> &arr) {
18
19
20
       table.resize(lq[this->n] + 1, vector<int>(this->n));
21
22
       table[0] = arr;
23
       int pow2 = 1;
24
25
       for(int i = 1; i < table.size(); i++) {</pre>
26
         int lastsz = this->n - pow2 + 1;
27
         for (int j = 0; j + pow2 < lastsz; <math>j++) {
28
           table[i][j] = merge(table[i-1][j], table[i-1][j+pow2]);
29
         pow2 <<= 1;
30
31
32
33
34
35
     int merge(const int &1, const int &r) {
36
       return min(l, r);
37
38
39
40
     Sparse_Table(const vector<int> &arr) {
41
       this->n = arr.size();
42
43
       this->build_log_array();
44
       this->build sparse table(arr);
45
46
47
     void print() {
48
       int pow2 = 1:
       for(int i = 0; i < table.size(); i++) {</pre>
49
50
         int sz = (int) (table.front().size()) - pow2 + 1;
          for(int j = 0; j < sz; j++) {</pre>
51
           cout << table[i][j] << " \n"[(j+1) == sz];</pre>
52
53
54
         pow2 <<= 1;
55
56
57
58
     /// Query of a range from 1 to r.
59
60
     /// Time Complexity: O(1)
61
     /// Space Complexity: O(1)
62
     int query(int 1, int r) {
63
       assert(1 <= r);
       assert(0 <= 1 && r <= this->n - 1);
```

```
int lgg = lg[(r - l + 1)];
return merge(table[lgg][1], table[lgg][r - (1 << lgg) + 1]);
};
</pre>
```

- 3. Dp
- 3.1. Achar Maior Palindromo

```
1 Fazer LCS da string com o reverso
```

3.2. Catalan

```
1 // The first few Catalan numbers for n = 0, 1, 2, 3, ...
2 // are 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, ...
3 // Formula Recursiva:
4 / / \text{cat}(0) = 0
5 / /  cat (n+1) =  somatorio (i from 0 to n) (cat <math>(i) * cat (n-i))
7 // Using Binomial Coefficient
   // We can also use the below formula to find nth catalan number in O(n) time.
   // Formula acima
10
11 // Returns value of Binomial Coefficient C(n, k)
12
13 | int binomialCoeff(int n, int k) {
14
     int res = 1;
15
16
     // Since C(n, k) = C(n, n-k)
17
     if (k > n - k)
       k = n - k:
18
19
     // Calculate value of [n*(n-1)*--*(n-k+1)] / [k*(k-1)*--*1]
20
21
     for (int i = 0; i < k; ++i) {
22
         res \star = (n - i);
         res /= (i + 1);
23
24
25
26
       return res:
2.7
28 // A Binomial coefficient based function to find nth catalan
29 // number in O(n) time
30 int catalan(int n) {
31
        // Calculate value of 2nCn
32
       int c = binomialCoeff(2*n, n);
33
34
        // return 2nCn/(n+1)
35
       return c/(n+1);
36 }
```

3.3. Digit Dp

```
8 \mid //// f = the number we are building has already become smaller than b? [0 =
       no, 1 = ves1
1.0
   int call(int pos, int cnt, int f) {
     if(cnt > k) return 0;
11
12
13
     if(pos == num.size()){
14
       if(cnt == k) return 1;
15
       return 0:
16
17
18
     if (DP[pos][cnt][f] != -1) return DP[pos][cnt][f];
19
     int res = 0;
     int lim = (f ? 9 : num[pos]);
20
21
     /// Try to place all the valid digits such that the number doesn't exceed b
22
23
     for (int dgt = 0; dgt <= LMT; dgt ++) {
24
       int nf = f:
25
       int ncnt = cnt;
26
       if (f == 0 && dqt < LMT) nf = 1; /// The number is getting smaller at
       this position
       if(dqt == d) ncnt++;
28
       if(ncnt <= k) res += call(pos+1, ncnt, nf);</pre>
29
30
31
     return DP[pos][cnt][f] = res;
32
33
34
   int solve(int b){
     num.clear();
35
36
     while(b>0){
37
       num.push_back(b%10);
38
       b/=10;
39
     reverse(num.begin(), num.end());
40
41
     /// Stored all the digits of b in num for simplicity
42
43
     memset (DP, -1, sizeof (DP));
44
     int res = call(0, 0, 0);
45
     return res:
46
47
48
   int main () {
    cin >> a >> b >> d >> k;
    int res = solve(b) - solve(a-1);
     cout << res << endl;
53
54
    return 0;
55
```

3.4. Longest Common Subsequence

```
string lcs(string &s, string &t) {
   int n = s.size(), m = t.size();
   s.insert(s.begin(), '#');
   t.insert(t.begin(), '$');
   vector<vector<int>> mat(n + 1, vector<int>(m + 1, 0));
   for(int i = 1; i <= n; i++) {
      for(int j = 1; j <= m; j++) {</pre>
```

```
if(s[i] == t[j])
13
           mat[i][j] = mat[i - 1][j - 1] + 1;
14
15
           mat[i][j] = max(mat[i - 1][j], mat[i][j - 1]);
16
17
18
19
     string ans;
     int i = n, j = m;
20
2.1
      while(i > 0 &   j > 0) {
       if(s[i] == t[j])
22
23
         ans += s[i], i--, j--;
       else if (mat[i][j - 1] > mat[i - 1][j])
24
25
       else
26
27
         i--;
28
29
30
     reverse(ans.begin(), ans.end());
     return ans:
```

3.5. Longest Common Substring

```
int LCSubStr(char *X, char *Y, int m, int n) {
    // Create a table to store lengths of longest common suffixes of
    // substrings. Notethat LCSuff[i][j] contains length of longest
     // common suffix of X[0..i-1] and Y[0..i-1]. The first row and
     // first column entries have no logical meaning, they are used only
     // for simplicity of program
     int LCSuff[m+1][n+1];
     int result = 0; // To store length of the longest common substring
     /\star Following steps build LCSuff[m+1][n+1] in bottom up fashion. \star/
1.0
11
     for (int i=0; i<=m; i++) {
12
        for (int j=0; j<=n; j++) {</pre>
         if (i == 0 | | j == 0)
13
14
           LCSuff[i][j] = 0;
15
16
         else if (X[i-1] == Y[i-1]) {
17
           LCSuff[i][j] = LCSuff[i-1][j-1] + 1;
18
            result = max(result, LCSuff[i][j]);
19
20
         else LCSuff[i][j] = 0;
21
22
23
     return result;
24
```

3.6. Longest Increasing Subsequence 2D (Not Sorted)

```
set<ii>set<ii>set(int)2e6];
bool check(ii par, int ind) {

auto it = s[ind].lower_bound(ii(par.ff, -INF));
if(it == s[ind].begin())
return false;

it--;

if(it->ss < par.ss)
return true;
return false;</pre>
```

```
13 | }
14
   int lis2d(vector<ii> &arr) {
16
17
     int n = arr.size();
18
     s[1].insert(arr[0]);
19
2.0
     int maior = 1;
     for(int i = 1; i < n; i++) {</pre>
21
22
23
       ii x = arr[i];
24
25
       int 1 = 1, r = maior;
26
       int ansbb = 0;
       while(1 <= r) {
27
         int mid = (1+r)/2;
28
29
         if(check(x, mid)) {
30
           1 = mid + 1:
31
           ansbb = mid;
32
          } else {
33
           r = mid - 1:
34
35
36
37
       // inserting in list
38
       auto it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
39
       while(it != s[ansbb+1].end() && it->ss >= x.ss)
         it = s[ansbb+1].erase(it);
40
41
42
       it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
43
       if(s[ansbb+1].size() > 0 && it != s[ansbb+1].end() && it->ff == x.ff &&
        it->ss <= x.ss)
44
45
        s[ansbb+1].insert(arr[i]);
46
47
       maior = max(maior, ansbb + 1);
48
49
50
     return maior;
51
52
```

3.7. Longest Increasing Subsequence 2D (Sorted)

```
set<ii>> s[(int)2e6];
   bool check(ii par, int ind) {
4
     auto it = s[ind].lower_bound(ii(par.ff, -INF));
     if(it == s[ind].begin())
       return false;
     it--;
10
     if(it->ss < par.ss)</pre>
11
       return true:
     return false:
12
13
14
   int lis2d(vector<ii> &arr) {
16
17
     int n = arr.size();
18
     s[1].insert(arr[0]);
19
     int maior = 1;
```

```
for(int i = 1; i < n; i++) {</pre>
22
23
       ii x = arr[i]:
24
25
       int 1 = 1, r = maior;
26
       int ansbb = 0;
27
       while(1 <= r) {
         int mid = (1+r)/2;
28
         if(check(x, mid)) {
29
30
           1 = mid + 1:
31
           ansbb = mid;
32
          } else {
33
           r = mid - 1;
34
35
36
37
        // inserting in list
38
       auto it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
39
       while(it != s[ansbb+1].end() && it->ss >= x.ss)
40
         it = s[ansbb+1].erase(it);
41
42
       it = s[ansbb+1].lower bound(ii(x.ff, -INF));
43
       if(s[ansbb+1].size() > 0 && it != s[ansbb+1].end() && it->ff == x.ff &&
        it->ss <= x.ss)
44
         continue;
45
       s[ansbb+1].insert(arr[i]);
46
47
       maior = max(maior, ansbb + 1);
48
49
50
     return maior;
51
52
```

3.8. Longest Increasing Subsequence

```
1 | int lis(vector<int> &arr) {
     int n = arr.size();
     vector<int> lis;
     for(int i = 0; i < n; i++) {</pre>
       int l = 0, r = (int) lis.size() - 1;
       int ans i = -1;
        while(1 <= r){
8
          int mid = (1+r)/2;
          // OBS: PARA >= TROCAR SINAL EMBAIXO POR <=
9
10
          if(arr[i] < lis[mid]) {
11
           r = mid - 1;
12
            ansj = mid;
13
14
          else l = mid + 1;
15
16
        if(ansj == -1){
17
          // se arr[i] e maior que todos
18
          lis.push_back(arr[i]);
19
20
        else H
21
          lis[ansj] = arr[i];
22
23
24
25
     return lis.size();
26
```

3.9. Subset Sum Com Bitset

```
bitset<312345> bit;
int arr[112345];

void subsetSum(int n) {
   bit.reset();
   bit.set(0);
   for(int i = 0; i < n; i++) {
      bit | = (bit << arr[i]);
   }
}</pre>
```

3.10. Coin Change Problem

```
1 // função que recebe o valor de troco N, o número de moedas disponíveis M,
   // e um vetor com as moedas disponíveis arr
   // essa função deve retornar o número mínimo de moedas,
   // de acordo com a solução com Programação Dinamica.
   int num_moedas(int N, int M, int arr[]) {
     int dp[N+1];
     // caso base
     dp[0] = 0;
     // sub-problemas
10
     for(int i=1; i<=N; i++) {</pre>
     // é comum atribuir um valor alto, que concerteza
11
     // é maior que qualquer uma das próximas possibilidades,
12
     // sendo assim substituido
13
14
     dp[i] = 1000000;
15
     for(int j=0; j<M; j++) {</pre>
16
       if(i-arr[i] >= 0) {
17
         dp[i] = min(dp[i], dp[i-arr[j]]+1);
18
19
20
     // solução
21
22
     return dp[N];
23
```

3.11. Knapsack

```
int dp[2001][2001];
   int moc(int q, int p, vector<ii> vec) {
3
     for(int i = 1; i <= q; i++)
4
5
       for(int j = 1; j <= p; j++) {
         if(j \ge vec[i-1].ff)
           dp[i][j] = max(dp[i-1][j], vec[i-1].ss + dp[i-1][j-vec[i-1].ff]);
          else
           dp[i][j] = dp[i-1][j];
10
11
12
     return dp[q][p];
13
   int main(int argc, char *argv[])
14
15
16
     int p,q;
17
     vector<ii> vec;
18
     cin >> p >> q;
19
     int x,y;
20
     for (int i = 0; i < q; i++) {
21
          cin >> x >> y;
22
          vec.push_back(make_pair(x,y));
23
```

4. Geometry

4.1. Centro De Massa De Um Poligono

```
double area = 0;
pto c;

c.x = c.y = 0;
for(int i = 0; i < n; i++) {
    double aux = (arr[i].x * arr[i+1].y) - (arr[i].y * arr[i+1].x); // shoelace
    area += aux;
    c.x += aux*(arr[i].x + arr[i+1].x);
    c.y += aux*(arr[i].y + arr[i+1].y);

c.y /= (3.0*area);
    c.y /= (3.0*area);
    cut << c.x << ' ' << c.y << endl;</pre>
```

4.2. Circle-Circle Intersection

```
1 /* circle_circle_intersection() *
    * Determine the points where 2 circles in a common plane intersect.
    * int circle_circle_intersection(
                                      // center and radius of 1st circle
5
                                      double x0, double y0, double r0,
 6
                                      // center and radius of 2nd circle
 8
                                      double x1, double y1, double r1,
 9
                                      // 1st intersection point
10
                                      double *xi, double *yi,
11
                                      // 2nd intersection point
12
                                      double *xi_prime, double *yi_prime)
13
    * This is a public domain work. 3/26/2005 Tim Voght
14
15
16
    */
17
    int circle_circle_intersection(double x0, double y0, double r0, double x1,
19
                                   double y1, double r1, double *xi, double *yi,
20
                                   double *xi_prime, double *yi_prime) {
21
     double a, dx, dy, d, h, rx, ry;
22
     double x2, y2;
23
2.4
     /* dx and dy are the vertical and horizontal distances between
25
      * the circle centers.
26
      */
27
     dx = x1 - x0;
28
     dy = y1 - y0;
29
30
     /* Determine the straight-line distance between the centers. */
31
     // d = sqrt((dy*dy) + (dx*dx));
32
     d = hypot(dx, dy); // Suggested by Keith Briggs
33
```

```
/* Check for solvability. */
35
     if (d > (r0 + r1)) {
36
       /* no solution. circles do not intersect. */
37
       return 0;
38
39
     if (d < fabs(r0 - r1)) {
40
       /* no solution. one circle is contained in the other */
41
42
43
     /* 'point 2' is the point where the line through the circle
44
      * intersection points crosses the line between the circle
46
      * centers.
47
      */
48
49
     /★ Determine the distance from point 0 to point 2. ★/
50
     a = ((r0 * r0) - (r1 * r1) + (d * d)) / (2.0 * d);
51
52
     /* Determine the coordinates of point 2. */
53
     x2 = x0 + (dx * a / d);
     y2 = y0 + (dy * a / d);
54
55
     /* Determine the distance from point 2 to either of the
57
      * intersection points.
58
59
     h = sqrt((r0 * r0) - (a * a));
60
61
     /★ Now determine the offsets of the intersection points from
62
      * point 2.
63
      */
64
     rx = -dv * (h / d);
65
     ry = dx * (h / d);
66
67
     /* Determine the absolute intersection points. */
68
     \star xi = x2 + rx:
69
     *xi prime = x2 - rx;
70
     \staryi = y2 + ry;
71
     *yi_prime = y2 - ry;
72
73
     return 1;
```

4.3. Closest Pair Of Points

```
struct Point {
2
    int x, y;
3
   int compareX(const void *a,const void *b) {
4
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->x - p2->x);
   int compareY(const void *a,const void *b) {
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->y - p2->y);
10
11
12 | float dist(Point p1, Point p2) {
    return sqrt((p1.x-p2.x)*(p1.x-p2.x) + (p1.y - p2.y)*(p1.y - p2.y));
14
15 | float bruteForce(Point P[], int n) {
16
    float min = FLT_MAX;
    for (int i = 0; i < n; ++i)
17
       for (int j = i+1; j < n; ++j)
1.8
19
        if (dist(P[i], P[j]) < min)
20
          min = dist(P[i], P[i]);
```

```
return min;
22
23 float min(float x, float y) {
24 return (x < y)? x : y;
25
26 | float stripClosest(Point strip[], int size, float d) {
27 float min = d:
     for (int i = 0; i < size; ++i)
2.8
        for (int j = i+1; j < size && (strip[j].y - strip[i].y) < min; ++j)
29
30
         if (dist(strip[i], strip[j]) < min)</pre>
31
           min = dist(strip[i], strip[j]);
32
     return min;
33
34 | float closestUtil(Point Px[], Point Py[], int n) {
35
    if (n <= 3)
       return bruteForce(Px, n);
36
     int mid = n/2;
37
38
    Point midPoint = Px[mid];
39
     Point Pvl[mid+1];
40
     Point Pyr[n-mid-1];
     int li = 0, ri = 0;
     for (int i = 0; i < n; i++)
       if (Pv[i].x <= midPoint.x)</pre>
44
         Pyl[li++] = Py[i];
45
46
         Pyr[ri++] = Py[i];
47
48
     float dl = closestUtil(Px, Pyl, mid);
49
     float dr = closestUtil(Px + mid, Pyr, n-mid);
     float d = min(dl, dr);
50
51
     Point strip[n];
52
     int j = 0;
53
     for (int i = 0; i < n; i++)
54
       if (abs(Py[i].x - midPoint.x) < d)</pre>
55
         strip[j] = Py[i], j++;
56
     return min(d, stripClosest(strip, j, d));
57
58
59 | float closest (Point P[], int n) {
60
    Point Px[n]:
61
     Point Py[n];
62
     for (int i = 0; i < n; i++) {
       Px[i] = P[i];
63
       Py[i] = P[i];
64
65
     gsort(Px, n, sizeof(Point), compareX);
     gsort(Pv, n, sizeof(Point), compareY);
68
     return closestUtil(Px, Py, n);
69 }
```

4.4. Condicao De Existencia De Um Triangulo

```
1
2
  | b - c | < a < b + c
3  | a - c | < b < a + c
4  | a - b | < c < a + b
5
6  Para a < b < c, basta checar
7  a + b > c

0BS: Para um conjunto n >= 100 sempre exite um triângulo válido, pois a sequência de triângulos não válidos seguem a sequência de Fibonacci e Fib(100) > 2^64
```

4.5. Convex Hull

```
1 // Asymptotic complexity: O(n log n).
   struct pto {
3
     double x, y;
4
     bool operator <(const pto &p) const {</pre>
       return x < p.x | | (x == p.x && y < p.y);
6
       /★ a impressao será em prioridade por mais a esquerda, mais
7
          abaixo, e antihorário pelo cross abaixo */
8
9
   };
10
11
   double cross(const pto &O, const pto &A, const pto &B) {
12
     return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
13
14
15
   vector<pto> convex_hull(vector<pto> P) {
     int n = P.size(), k = 0;
16
17
     vector<pto> H(2 * n);
     // Sort points lexicographically
18
19
     sort(P.begin(), P.end());
     // Build lower hull
20
     for (int i = 0; i < n; ++i) {
22
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
23
       // trocar por >= 0
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
24
25
        k--;
       H[k++] = P[i];
26
27
28
     // Build upper hull
29
     for (int i = n - 2, t = k + 1; i >= 0; i--) {
30
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
31
       // trocar por >= 0
32
       while (k \ge t \& cross(H[k - 2], H[k - 1], P[i]) \le 0)
33
         k--;
34
       H[k++] = P[i];
35
36
     H.resize(k);
37
     /★ o último ponto do vetor é iqual ao primeiro, atente para isso
38
     as vezes é necessário mudar */
39
     return H;
40
```

4.6. Cross Product

```
// Outra forma de produto vetorial
   // reta ab, ac se for zero e colinear
   // se for < 0 entao antiHorario, > 0 horario
  bool ehcol(pto a,pto b,pto c) {
    return ((b.y-a.y)*(c.x-a.x) - (b.x-a.x)*(c.y-a.y));
8 //Produto vetorial AB x AC, se for zero e colinear
9 int cross(pto A, pto B, pto C) {
    pto AB, AC;
    AB.x = B.x-A.x;
    AB.v = B.v-A.v;
13
    AC.x = C.x-A.x;
14
    AC.y = C.y-A.y;
15
    int cross = AB.x*AC.y-AB.y * AC.x;
16
    return cross;
17
18
```

19 | // OBS: DEFINE ÁREA DE QUADRILÁTERO FORMADO PELAS RETAS, A ÁREA DO TRIÂNGULO É A METADE

4.7. Distance Point Segment

```
// use struct point and line
double dist_point_segment(const Point p, const Point s, const Point t) {
   if (sgn(dot(p-s, t-s)) < 0)
      return (p-s).norm();
   if (sgn(dot(p-t, s-t)) < 0)
      return (p-t).norm();
   return abs(det(s-p, t-p) / dist(s, t));
}</pre>
```

4.8. Line-Line Intersection

```
1 // Intersecção de retas Ax + Bv = C
                                            dados pontos (x1,y1) e (x2,y2)
2 A = v2-v1
3 \mid B = x1-x2
4 \mid C = A \star x1 + B \star v1
5 //Retas definidas pelas equações:
6 \mid A1x + B1y = C1
7 | A2x + B2y = C2
8 //Encontrar x e y resolvendo o sistema
9 double det = A1*B2 - A2*B1;
10 | if (det == 0) {
   //Lines are parallel
11
12 }else{
double x = (B2*C1 - B1*C2)/det;
    double v = (A1*C2 - A2*C1)/det;
```

4.9. Line-Point Distance

```
1 | double ptoReta(double x1, double y1, double x2, double y2, double pointX,
        double pointY, double *ptox, double *ptoy) {
     double diffX = x^2 - x^1;
     double diffY = y2 - y1;
3
     if ((diffX == 0) && (diffY == 0)) {
5
       diffX = pointX - x1;
       diffY = pointY - y1;
6
7
       //se os dois sao pontos
8
       return hypot (pointX - x1, pointY - y1);
9
10
     double t = ((pointX - x1) * diffX + (pointY - y1) * diffY) /
                      (diffX * diffX + diffY * diffY);
11
12
     if (t < 0) {
13
       //point is nearest to the first point i.e x1 and y1
14
15
       // cord do pto na reta = pto inicial(x1,y1);
        \starptox = x1, \starptoy = y1;
16
       diffX = pointX - x1;
17
       diffY = pointY - y1;
1.8
       else if (t > 1) {
20
        //point is nearest to the end point i.e x2 and y2
21
        // Ex : .
       // cord do pto na reta = pto final(x2,y2);
22
23
        *ptox = x2, *ptoy = y2;
24
       diffX = pointX - x2;
25
       diffY = pointY - y2;
26
     } else {
27
         //if perpendicular line intersect the line segment.
```

```
// pto nao esta mais proximo de uma das bordas do segmento
29
         // Ex:
30
         //
31
         //
                              |(Ângulo Reto)
32
         //
33
         // cord x do pto na reta = (x1 + t * diffX)
34
         // cord v do pto na reta = (v1 + t * diffY)
35
       *ptox = (x1 + t * diffX), *ptoy = (y1 + t * diffY);
36
       diffX = pointX - (x1 + t * diffX);
       diffY = pointY - (v1 + t * diffY);
37
38
39
     //returning shortest distance
     return sqrt (diffX * diffX + diffY * diffY);
40
41
```

4.10. Point Inside Convex Polygon - Log(N)

```
#include <bits/stdc++.h>
2
   using namespace std;
3
   #define INF 1e18
   #define pb push_back
   #define ii pair<int,int>
   #define OK cout << "OK" << endl
   #define debug(x) cout << \#x " = " << (x) << endl
10 #define ff first
   #define ss second
11
   #define int long long
12
13
14 struct pto {
     double x, v;
16
    bool operator < (const pto &p) const {
17
       return x < p.x \mid | (x == p.x \&\& v < p.v);
18
       /★ a impressao será em prioridade por mais a esquerda, mais
19
          abaixo, e antihorário pelo cross abaixo */
20
21
   double cross(const pto &O, const pto &A, const pto &B) {
23
    return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
24
25
   vector<pto> lower, upper;
26
27
28
   vector<pto> convex hull(vector<pto> &P) {
     int n = P.size(), k = 0;
30
     vector<pto> H(2 * n);
31
     // Sort points lexicographically
32
     sort(P.begin(), P.end());
33
     // Build lower hull
34
     for (int i = 0; i < n; ++i) {
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
35
36
       // trocar por >= 0
37
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
38
        k--:
39
       H[k++] = P[i];
40
41
     // Build upper hull
42
     for (int i = n - 2, t = k + 1; i >= 0; i--) {
43
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
44
       // trocar por >= 0
       while (k \ge t \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
45
46
        k--;
47
       H[k++] = P[i];
```

```
49
     H.resize(k);
50
      /★ o último ponto do vetor é iqual ao primeiro, atente para isso
51
      as vezes é necessário mudar */
52
53
      int j = 1;
54
      lower.pb(H.front());
      while (\bar{H}[j].x >= H[j-1].x) {
5.5
56
        lower.pb(H[j++]);
57
58
59
      int l = H.size()-1;
 60
      while (1 >= i)
        upper.pb(H[1--]);
61
62
 63
      upper.pb(H[1--]);
 64
 65
      return H;
 66
 67
    bool insidePolygon(pto p, vector<pto> &arr) {
70
      if (pair<double, double>(p.x, p.y) == pair<double, double>(lower[0].x,
         lower[0].y))
71
        return true;
72
73
      pto lo = {p.x, -(double) INF};
      pto hi = {p.x, (double) INF};
74
75
      auto itl = lower_bound(lower.begin(), lower.end(), lo);
76
      auto itu = lower_bound(upper.begin(), upper.end(), lo);
77
78
      if(itl == lower.begin() || itu == upper.begin())
79
        auto it = lower_bound(arr.begin(), arr.end(), lo);
8.0
         auto it2 = lower_bound(arr.begin(), arr.end(), hi);
81
        it2--;
82
        if(it2 >= it && p.x == it-> x && it->x == it2->x && it->v <= p.v && p.v
         \langle = it.2-\rangle v
83
          return true;
84
        return false;
85
86
      if(itl == lower.end() || itu == upper.end()) {
87
        return false;
88
89
90
      auto ol = itl, ou = itu;
91
      ol--, ou--;
92
      if(cross(*ol, *itl, p) >= 0 && cross(*ou, *itu, p) <= 0)</pre>
93
        return true;
94
9.5
      auto it = lower bound(arr.begin(), arr.end(), lo);
      auto it2 = lower_bound(arr.begin(), arr.end(), hi);
 97
      it2--:
98
      if(it2 >= it && p.x == it-> x && it->x == it2->x && it->y <= p.y && p.y <=
        it2->y)
        return true;
99
100
101
      return false:
102
103
104
105
    signed main () {
106
      ios_base::sync_with_stdio(false);
107
108
      cin.tie(NULL);
109
```

```
double n, m, k;
111
112
       cin >> n >> m >> k;
113
114
       vector<pto> arr(n);
115
116
       for (pto &x: arr) {
117
         cin >> x.x >> x.y;
118
119
120
       convex hull(arr);
121
122
      pto p;
123
124
       int c = 0;
125
       while (m--) {
126
        cin >> p.x >> p.y;
         cout << (insidePolygon(p, arr) ? "dentro" : "fora") << endl;</pre>
127
128
129
130
```

4.11. Point Inside Polygon

```
2
   /* Traça-se uma reta do ponto até um outro ponto qualquer fora do triangulo
        e checa o número de interseção com a borda do polígono se este for ímpar
        então está dentro se não está fora */
4
   // Define Infinite (Using INT_MAX caused overflow problems)
5
   #define INF 10000
6
7
   struct pto {
8
       int x, y;
9
       pto() {}
10
       pto(int x, int y) : x(x), y(y) {}
11
12
13 // Given three colinear ptos p, q, r, the function checks if 14 // pto q lies on line segment 'pr'
15
   bool onSegment(pto p, pto q, pto r) {
16
     if (q.x \le max(p.x, r.x) \& q.x >= min(p.x, r.x) \& \&
17
          q.y \le max(p.y, r.y) && q.y >= min(p.y, r.y)
18
        return true;
19
     return false:
20
21
22
   // To find orientation of ordered triplet (p, q, r).
   // The function returns following values
24
   // 0 --> p, q and r are colinear
25
   /// 1 --> Clockwise
   // 2 --> Counterclockwise
26
27
   int orientation(pto p, pto q, pto r) {
28
     int val = (q.y - p.y) * (r.x - q.x) -
29
                (q.x - p.x) * (r.y - q.y);
30
31
     if (val == 0) return 0; // colinear
32
     return (val > 0)? 1: 2; // clock or counterclock wise
33
34
35 // The function that returns true if line segment 'plq1'
36 // and 'p2g2' intersect.
37 | bool doIntersect(pto p1, pto q1, pto p2, pto q2) {
    // Find the four orientations needed for general and
```

```
// special cases
     int o1 = orientation(p1, q1, p2);
     int o2 = orientation(p1, q1, q2);
     int o3 = orientation(p2, q2, p1);
     int o4 = orientation(p2, q2, q1);
43
11
45
     // General case
46
     if (01 != 02 && 03 != 04)
47
       return true;
48
49
     // Special Cases
50
     // pl, q1 and p2 are colinear and p2 lies on segment p1q1
51
     if (o1 == 0 && onSegment(p1, p2, q1)) return true;
52
53
     // pl, ql and p2 are colinear and q2 lies on segment plq1
54
     if (o2 == 0 && onSegment(p1, q2, q1)) return true;
55
     // p2, q2 and p1 are colinear and p1 lies on segment p2q2
56
57
     if (o3 == 0 && onSegment(p2, p1, q2)) return true;
5.8
59
      // p2, q2 and q1 are colinear and q1 lies on segment p2q2
     if (o4 == 0 && onSegment(p2, q1, q2)) return true;
60
61
     return false; // Doesn't fall in any of the above cases
62
63
64
   // Returns true if the pto p lies inside the polygon[] with n vertices
   bool isInside(pto polygon[], int n, pto p) {
     // There must be at least 3 vertices in polygon[]
67
68
     if (n < 3) return false;
69
70
     // Create a pto for line segment from p to infinite
71
     pto extreme = pto(INF, p.y);
72
73
     // Count intersections of the above line with sides of polygon
74
     int count = 0, i = 0;
75
76
       int next = (i+1)%n;
77
78
       // Check if the line segment from 'p' to 'extreme' intersects
       // with the line segment from 'polygon[i]' to 'polygon[next]'
79
8.0
       if (doIntersect(polygon[i], polygon[next], p, extreme)) {
         // If the pto 'p' is colinear with line segment 'i-next',
81
         // then check if it lies on segment. If it lies, return true,
82
83
         // otherwise false
84
         if (orientation(polygon[i], p, polygon[next]) == 0)
85
           return onSegment(polygon[i], p, polygon[next]);
86
87
         count++;
88
89
       i = next;
90
     } while (i != 0);
91
     // Return true if count is odd, false otherwise
     return count&1; // Same as (count%2 == 1)
93
94
```

4.12. Points Inside And In Boundary Polygon

```
int cross(pto a, pto b) {
   return a.x * b.y - b.x * a.y;
}
int boundaryCount(pto a, pto b) {
```

```
if(a.x == b.x)
       return abs(a.v-b.v)-1;
     if(a.y == b.y)
9
       return abs(a.x-b.x)-1;
10
     return _gcd(abs(a.x-b.x), abs(a.y-b.y))-1;
11
12
13
   int totalBoundaryPolygon(vector<pto> &arr, int n) {
14
1.5
     int boundPoint = n;
     for(int i = 0; i < n; i++) {</pre>
16
       boundPoint += boundaryCount(arr[i], arr[(i+1)%n]);
17
18
19
     return boundPoint;
20
21
22
   int polygonArea2(vector<pto> &arr, int n) {
23
     int area = 0;
24
     // N = quantidade de pontos no polígono e armazenados em p;
25
     // OBS: VALE PARA CONVEXO E NÃO CONVEXO
     for(int i = 0; i<n; i++) {
26
27
       area += cross(arr[i], arr[(i+1)%n]);
28
29
    return abs(area);
30
31
32
   int internalCount(vector<pto> &arr, int n) {
33
34
     int area_2 = polygonArea2(arr, n);
35
     int boundPoints = totalBoundaryPolygon(arr,n);
36
    return (area_2 - boundPoints + 2)/2;
37
```

4.13. Polygon Area (3D)

```
#include <bits/stdc++.h>
2
3
   using namespace std;
4
5
   struct point {
6
     double x, y, z;
     void operator=(const point & b) {
8
       x = b.x
9
       v = b.v;
10
       z = b.z;
11
12
13
14
   point cross(point a, point b) {
15
     point ret;
16
    ret.x = a.y*b.z - b.y*a.z;
17
     ret.y = a.z*b.x - a.x*b.z;
     ret.z = a.x*b.y - a.y*b.x;
18
19
     return ret;
2.0
21
22 | int main() {
     int num;
     cin >> num:
25
     point v[num];
26
     for(int i=0; i<num; i++) cin >> v[i].x >> v[i].y >> v[i].z;
27
28
     point cur;
     cur.x = 0, cur.y = 0, cur.z = 0;
```

```
31
     for(int i=0; i<num; i++) {</pre>
32
       point res = cross(v[i], v[(i+1)%num]);
33
       cur.x += res.x;
34
       cur.v += res.v;
35
       cur.z += res.z;
36
37
38
     double ans = sqrt(cur.x*cur.x + cur.y*cur.y + cur.z*cur.z);
39
     double area = abs(ans);
40
41
42
     cout << fixed << setprecision(9) << area/2. << endl;</pre>
43
```

4.14. Polygon Area

```
double polygonArea(vector<pto> &arr, int n) {
   int area = 0;
   // N = quantidade de pontos no polígono e armazenados em p;
   // OBS: VALE PARA CONVEXO E NÃO CONVEXO
   for(int i = 0; i<n; i++) {
      area += cross(arr[i], arr[(i+1)%n]);
   }
   return (double) abs(area/2.0);
}</pre>
```

4.15. Segment-Segment Intersection

```
1 | // Given three colinear points p, q, r, the function checks if
2 // point q lies on line segment 'pr
3 int onSegment(Point p, Point q, Point r) {
     if (q.x \le max(p.x, r.x)) && q.x \ge min(p.x, r.x) && q.y \le max(p.y, r.y)
       && q.y \Rightarrow min(p.y, r.y))
       return true;
     return false;
7
8 /★ PODE SER RETIRADO
9 int onSegmentNotBorda (Point p, Point q, Point r) {
10
       if (q.x < max(p.x, r.x) \& q.x > min(p.x, r.x) \& q.y <= max(p.y, r.y)
       && q.y >= min(p.y, r.y)
11
            return true;
       if (q.x \le max(p.x, r.x) \& q.x \ge min(p.x, r.x) \& q.y < max(p.y, r.y)
       && q.y > min(p.y, r.y)
13
            return true;
14
       return false:
15
16 | */
17 // To find orientation of ordered triplet (p, q, r).
18 // The function returns following values
19 // 0 --> p, q and r are colinear
20 // 1 --> Clockwise
21 // 2 --> Counterclockwise
22 int orientation (Point p, Point q, Point r) {
    int val = (q.y - p.y) * (r.x - q.x) -
                (q.x - p.x) * (r.y - q.y);
     if (val == 0) return 0; // colinear
     return (val > 0)? 1: 2; // clock or counterclock wise
26
27
28 // The main function that returns true if line segment 'p1p2'
29 // and 'q1q2' intersect.
30 | int doIntersect(Point p1, Point p2, Point q1, Point q2) {
```

```
// Find the four orientations needed for general and
     // special cases
     int o1 = orientation(p1, p2, q1);
     int o2 = orientation(\bar{p}1, \bar{p}2, q2);
     int o3 = orientation(q1, q2, p1);
36
     int o4 = orientation(q1, q2, p2);
37
38
     // General case
     if (01 != 02 && 03 != 04) return 2;
39
40
   /* PODE SER RETIRADO
41
    if (o1 == o2 && o2 == o3 && o3 == o4 && o4 == 0) {
42
43
       //INTERCEPTAM EM RETA
44
       if(onSegmentNotBorda(p1,q1,p2) || onSegmentNotBorda(p1,q2,p2)) return 1;
45
       if(onSegmentNotBorda(q1,p1,q2) || onSegmentNotBorda(q1,p2,q2)) return 1;
46
47
48
     // Special Cases (INTERCEPTAM EM PONTO)
     // p1, p2 and q1 are colinear and q1 lies on segment p1p2
49
     if (o1 == 0 && onSegment(p1, q1, p2)) return 2;
     // p1, p2 and q1 are colinear and q2 lies on segment p1p2
51
     if (o2 == 0 && onSegment(p1, q2, p2)) return 2;
     // q1, q2 and p1 are colinear and p1 lies on segment q1q2
     if (o3 == 0 && onSegment(q1, p1, q2)) return 2;
     // q1, q2 and p2 are colinear and p2 lies on segment q1q2
    if (o4 == 0 && onSegment(q1, p2, q2)) return 2;
    return false; // Doesn't fall in any of the above cases
58
59 // OBS: SE (C2/A2 == C1/A1) SÃO COLINEARES
```

4.16. Upper And Lower Hull

```
1 struct pto {
     double x, v;
     bool operator <(const pto &p) const {</pre>
4
       return x < p.x || (x == p.x && y < p.y);
5
       /★ a impressao será em prioridade por mais a esquerda, mais
6
          abaixo, e antihorário pelo cross abaixo */
7
8
9
   double cross(const pto &O, const pto &A, const pto &B) {
    return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
10
11
12
   vector<pto> lower, upper;
13
14
15
   vector<pto> convex hull(vector<pto> &P) {
16
     int n = P.size(), k = 0;
17
     vector<pto> H(2 * n);
18
     // Sort points lexicographically
19
     sort(P.begin(), P.end());
20
     // Build lower hull
21
     for (int i = 0; i < n; ++i) {
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
22
23
       // trocar por >= 0
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
24
25
26
       H[k++] = P[i];
27
     // Build upper hull
28
29
     for (int i = n - 2, t = k + 1; i >= 0; i--) {
30
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
       // trocar por >= 0
31
       while (k \ge t \& cross(H[k - 2], H[k - 1], P[i]) \le 0)
```

```
34
       H[k++] = P[i];
35
36
     H.resize(k);
37
     /★ o último ponto do vetor é igual ao primeiro, atente para isso
38
     as vezes é necessário mudar */
39
40
     int j = 1;
     lower.pb(H.front());
41
     while (\bar{H}[j].x >= H[j-1].x) {
42
43
       lower.pb(H[j++]);
44
45
46
     int l = H.size()-1;
47
     while (1 >= j) {
       upper.pb(H[1--]);
49
50
     upper.pb(H[1--]);
51
52
    return H;
```

4.17. Struct Point And Line

```
int sqn(double x) {
       if(abs(x) < 1e-8) return 0;
       return x > 0 ? 1 : -1;
3
   inline double sqr(double x) { return x * x; }
   struct Point {
        double x, y, z;
9
       Point() {};
10
       Point (double a, double b): x(a), y(b) {};
11
       Point (double x, double y, double z): x(x), y(y), z(z) {}
12
13
       void input() { scanf(" %lf %lf", &x, &y); };
14
        friend Point operator+(const Point &a, const Point &b) {
15
            return Point(a.x + b.x, a.y + b.y);
16
17
        friend Point operator-(const Point &a, const Point &b) {
18
            return Point(a.x - b.x, a.y - b.y);
19
20
21
       bool operator !=(const Point& a) const {
22
            return (x != a.x || y != a.y);
23
24
25
       bool operator <(const Point &a) const{</pre>
26
         if(x == a.x)
27
           return y < a.y;</pre>
28
          return x < a.x;</pre>
2.9
30
31
        double norm() {
32
           return sqrt(sqr(x) + sqr(y));
33
34 };
35 double det (const Point &a, const Point &b) {
36
       return a.x * b.y - a.y * b.x;
37
38 | double dot(const Point &a, const Point &b) {
39
       return a.x * b.x + a.y * b.y;
40 }
```

```
41 | double dist (const Point &a, const Point &b) {
       return (a-b).norm();
43
44
45
46
  struct Line {
47
       Point a, b;
48
       Line() {}
49
       Line (Point x, Point y): a(x), b(y) {};
50
51
   double dis_point_segment(const Point p, const Point s, const Point t) {
53
       if(sqn(dot(p-s, t-s)) < 0)
54
         return (p-s).norm();
55
       if(sqn(dot(p-t, s-t)) < 0)
56
        return (p-t).norm();
57
       return abs (det (s-p, t-p) / dist(s, t));
58
```

5. Graphs

5.1. Checa Grafo Bipartido

```
bool isBipartite(int src, int V) {
2
     int colorArr[V + 1];
     memset(colorArr, -1, sizeof(colorArr));
     colorArr[src] = 1;
     queue <int> q; q.push(src);
8
     while (!q.empty()) {
       int u = q.front(); q.pop();
10
11
12
       // Find all non-colored adjacent vertices
13
       for (auto it = adj[u].begin(); it != adj[u].end(); it++) {
14
       //Return false if there is a self-loop
15
           if (u == *it)
16
               return false;
17
         // An edge from u to v exists and destination v is not colored
18
19
         if (colorArr[*it] == -1) {
           // Assign alternate color to this adjacent v of u
2.0
21
           colorArr[*it] = 1 - colorArr[u];
22
           q.push(*it);
23
24
         // An edge from u to v exists and destination v is colored with same
         else if (colorArr[*it] == colorArr[u])
26
           return false;
27
28
29
     // If we reach here, then all adjacent vertices can be colored with
30
     // alternate color
31
     return true;
32
```

5.2. Ciclo Grafo

```
int n;
vector<vector<int>> adj;
vector<char> color;
vector<int> parent;
```

```
5 | int cycle_start, cycle_end;
7 bool dfs(int v) {
    color[v] = 1;
     for (int u : adj[v])
10
       if (color[u] == 0) {
11
         parent[u] = v;
12
         if (dfs(u))
13
           return true;
14
       } else if (color[u] == 1) {
15
         cvcle end = v;
16
         cycle_start = u;
17
          return true;
18
19
20
     color[v] = 2;
21
     return false;
22
23
24
   void find_cycle()
     color.assign(n. 0):
25
     parent.assign(n, -1);
26
     cvcle start = -1:
28
29
     for (int v = 0; v < n; v++) {
30
       if (dfs(v))
31
         break;
32
33
34
     if (cycle_start == -1) {
3.5
       cout << "Acyclic" << endl;</pre>
36
      } else {
37
       vector<int> cvcle;
38
        cycle.push back(cycle start);
39
        for (int v = cycle_end; v != cycle_start; v = parent[v])
40
         cvcle.push back(v);
        cycle.push_back(cycle_start);
41
42
       reverse (cycle.begin(), cycle.end());
43
44
       cout << "Cycle found: ";
       for (int v : cycle)
45
         cout << v << " ";
46
47
       cout << endl;
48
49 }
```

5.3. Diametro Em Arvore

1 Calcula qual o vértice a mais distante de um qualquer vértice X e **do** vértice A calcula-se o vértice B mais distante dele.

5.4. Floyd Warshall

```
// OBS: ZERAR adj[i][i] sempre
for(int i = 0; i < n; i++)
    adj[i][i] = 0;

for(int k = 0; k < n; k++) {
    for(int i = 0; i < n; i++) {
        for(int j = 0; j < n; j++) {
            adj[i][j] = min(adj[i][j],adj[i][k] + adj[k][j]);
      }
}</pre>
```

.1 }

5.5. Ford Fulkersson (Maximum Flow)

```
int rGraph[2000][2000];
2
   int graph[2000][2000];
3
4
   int V;
   bool bfs(int s, int t, int parent[]) {
     bool visited[V];
     memset(visited, 0, sizeof(visited));
9
     // Create a queue, enqueue source vertex and mark source vertex
10
     // as visited
     queue <int> q;
11
12
     q.push(s);
13
     visited[s] = true;
     parent[s] = -1;
14
15
     // Standard BFS Loop
16
17
     while (!q.empty()) {
18
       int u = q.front();
19
       q.pop();
20
21
       for (int v=0; v<V; v++) {</pre>
22
         if (visited[v] == false && rGraph[u][v] > 0) {
23
           q.push(v);
24
           parent[v] = u;
25
           visited[v] = true;
26
27
28
     // If we reached sink in BFS starting from source, then return true, else
29
     return (visited[t] == true);
31
32
33
   // Returns the maximum flow from s to t in the given graph
34 int fordFulkerson(int s, int t) {
35
    int u, v;
     // Create a residual graph and fill the residual graph with given
       capacities in the original graph as residual capacities in residual
       graph residual capacity of edge from i to j (if there is an edge. If
       rGraph[i][j] is 0, then there is not)
     for (u = 0; u < V; u++)
37
38
       for (v = 0; v < V; v++)
39
         rGraph[u][v] = graph[u][v];
40
41
     int parent[V]; // This array is filled by BFS and to store path
42
43
     int max_flow = 0;// There is no flow initially
44
45
     // Augment the flow while tere is path from source to sink
46
     while (bfs(s, t, parent)) {
47
       // Find minimum residual capacity of the edges along the path filled by
       BFS. Or we can say find the maximum flow through the path found.
48
       int path_flow = INT_MAX;
49
       for (v=t; v!=s; v=parent[v]) {
50
         u = parent[v];
51
         path_flow = min(path_flow, rGraph[u][v]);
52
53
54
       // update residual capacities of the edges and reverse edges
       // along the path
```

```
for (v=t; v != s; v=parent[v]) {
57
         u = parent[v];
58
         rGraph[u][v] -= path_flow;
59
         rGraph[v][u] += path_flow;
60
61
62
       // Add path flow to overall flow
6.3
       max_flow += path_flow;
64
6.5
     // Return the overall flow
66
67
     return max_flow;
68
69
   // PRINT THE FLOW AFTER RUNNING THE ALGORITHM
70
71
   void print(int n) {
    for(int i = 1; i <= m; i++) {
73
       for(int j = m+1; j <= m*2; j++) {</pre>
74
         cout << "flow from i(left) to j(right) is " << graph[i][j] -</pre>
       rGraph[i][j] << endl;
75
76
77
78
79
   void addEdge(int 1, int r, int n, int x) {
80
    graph[l][r+n] = x;
81 }
82
83 void addEdgeSource(int 1, int x) {
84
    graph[0][1] = x;
8.5
86
87
   void addEdgeSink(int r, int n, int x) {
    graph[r+n][V-1] = x;
89
```

5.6. Pontes Num Grafo

```
1 //SE TIRA-LAS O GRAFO FICA DESCONEXO
   // OBS: PRESTAR ATENCAO EM SELF-LOOPS, é MELHOR NÃO ADICIONA-LOS
3 // SO FUNCIONA EM GRAFO NÃO DIRECIONADO
   int t=1:
 5 vector int T((int)2e6,0); //Tempo necessário para chegar naquele vértice na
   vector<int> adi[(int)2e6];
   vector<int> Low((int)2e6); // Tempo "mínimo" para chegar naquele vértice na
   vector<int> ciclo((int)2e6, false);
   vector<ii> bridges;
10 | void dfs(int u, int p) {
11
    Low[u] = T[u] = t;
12
     t++;
1.3
     for(auto v : adj[u]) {
14
       if(v==p){
1.5
         //checa arestas paralelas
16
         p=-1:
17
         continue;
18
       //se ele ainda não foi visited
19
20
       else if(T[v]==0){
21
         dfs(v,u);
22
         Low[u]=min(Low[u], Low[v]);
23
         if(Low[v]>T[u]) {
           bridges.pb(ii(min(u,v), (max(u,v))));
24
```

```
// ponte de u para v
26
27
28
        else
29
         Low[u]=min(Low[u], T[v]);
30
        ciclo[u] \mid = (T[u] > = Low[v]);
31
        //checa se o vértice u faz parte de um ciclo
32
33
34
35
   void clear() {
36
37
     for(int i = 0; i <= n; i++) {</pre>
38
       T[i] = 0, Low[i] = 0, adj[i].clear(), ciclo[i] = false;
39
40
     bridges.clear();
41
42
43
44
   signed main () {
45
     for(int i = 0; i < n; i++)
46
        if(T[i] == 0)
47
48
          dfs(i, -1);
49
50
     sort(bridges.begin(), bridges.end());
51
52
     cout << (int)bridges.size() << endl;</pre>
     for(int i = 0; i < bridges.size(); i++) {</pre>
53
54
       cout << bridges[i].ff << " - " << bridges[i].ss << endl;</pre>
55
56
     cout << endl;
57
58
     clear();
59
```

5.7. Pontos De Articulação

```
// SE TIRAR TAIS VERTICES O GRAFO FICA DESCONEXO
2
3
   vector<bool> ap(100000, false);
   vector<int> low(100000,0), T(100000,0);
   int tempo = 1;
   list<int> adi[100000];
   void artPoint(int u, int p) {
10
     low[u] = T[u] = tempo++;
11
     int children = 0;
12
13
     for(int v: adj[u]) {
14
15
     // cuidado com arestas paralelas
     // se tiver nao podemos fazer assim
16
17
       if(T[v] == 0) {
18
19
         children++;
20
21
         artPoint(v,u);
22
         low[u] = min(low[v], low[u]);
23
24
       if(p == -1 \&\& children > 1) {
25
         ap[u] = true;
```

```
26
27
       if(p != -1 && low[v] > T[u])
28
29
         ap[u] = true;
30
       } else if(v != p)
31
         low[u] = min(low[u], T[v]);
32
33
34
3.5
36
   int main() {
37
38
     for(int i = 0; i < n; i++)
39
       if(T[i] == 0)
40
         artPoint(i,-1);
41 }
```

5.8. Scc (Kosaraju)

1 class SCC {

```
private:
     // number of vertices
    int n;
     // indicates whether it is indexed from 0 or 1
     int indexed from:
     // reversed graph
     vector<vector<int>> trans;
10
    private:
     void dfs_trans(int u, int id) {
11
12
       comp[u] = id;
       scc[id].push_back(u);
13
14
15
        for (int v: trans[u])
         if (comp[v] == -1)
16
17
           dfs_trans(v, id);
18
19
     void get_transpose(vector<vector<int>>& adj) {
20
21
       for (int u = indexed_from; u < this->n + indexed_from; u++)
22
          for(int v: adj[u])
23
           trans[v].push_back(u);
24
25
26
     void dfs_fill_order(int u, stack<int> &s, vector<vector<int>>& adj) {
27
       comp[u] = true;
28
        for(int v: adj[u])
29
30
         if(!comp[v])
31
           dfs_fill_order(v, s, adj);
32
33
       s.push(u);
34
35
36
     // The main function that finds all SCCs
37
     void compute_SCC(vector<vector<int>>& adj) {
38
39
40
        // Fill vertices in stack according to their finishing times
        for(int i = indexed_from; i < this->n + indexed_from; i++)
41
42
         if(!comp[i])
           dfs_fill_order(i, s, adj);
43
44
        // Create a reversed graph
45
```

28

29

30

31

32

33

34

3.5

36

37

38

39

40

41

42

43

44

45

46

47

48

49

50

51

52

5.3

54

55

56

57

58

59

60

61

62

63

64

65

66

67

70

71

72

7.3

74

75

76

77

78

79

80

81

82

8.3

84

85

```
get_transpose(adj);
47
48
       fill(comp.begin(), comp.end(), -1);
49
50
       // Now process all vertices in order defined by stack
51
       while(s.empty() == false) {
52
         int v = s.top();
53
         s.pop();
54
55
         if(comp[v] == -1)
56
           dfs trans(v, this->number of comp++);
57
58
59
    public:
60
61
     // number of the component of the i-th vertex
     // it's always indexed from 0
63
     vector<int> comp;
64
     // the i-th vector contains the vertices that belong to the i-th scc
65
     // it's always indexed from 0
     vector<vector<int>> scc:
     int number of comp = 0;
68
69
     SCC(int n, int indexed from, vector<vector<int>>& adi) {
       this->n = n;
70
71
       this->indexed_from = indexed_from;
72
       comp.resize(n + 1);
73
       trans.resize(n + 1);
74
       scc.resize(n + 1);
75
76
       this->compute SCC(adj);
77
78
   };
```

5.9. All Eulerian Path Or Tour

```
1 struct edge
    int v, id;
3
     edge() {}
4
     edge(int v, int id) : v(v), id(id) {}
5
   // The undirected + path and directed + tour wasn't tested in a problem.
   // TEST AGAIN BEFORE SUBMITTING IT!
   namespace graph {
     // Namespace which auxiliary funcions are defined.
10
     namespace detail {
11
12
       pair<bool, pair<int, int>> check_both_directed(const
       vector<vector<edge>> &adj, const vector<int> &in_degree) {
         // source and destination
13
14
         int src = -1, dest = -1;
15
         // adj[i].size() represents the out degree of an vertex
16
         for(int i = 0; i < adj.size(); i++) {
17
           if((int)adj[i].size() - in_degree[i] == 1) {
18
             if (src != -1)
19
               return make_pair(false, pair<int, int>());
20
21
           \} else if((int)adj[i].size() - in degree[i] == -1) {
              if(dest != -1)
22
23
               return make_pair(false, pair<int, int>());
24
             dest = i;
25
           } else if(abs((int)adj[i].size() - in_degree[i]) > 1)
26
             return make_pair(false, pair<int, int>());
27
```

```
if (src == -1 && dest == -1)
      return make_pair(true, pair<int, int>(src, dest));
    else if (src != -1 && dest != -1)
      return make pair(true, pair<int, int>(src, dest));
    return make pair(false, pair<int, int>());
  /// Builds the path/tour for directed graphs.
  void build (const int u, vector<int> &tour, vector<vector<edge>> &adj,
  vector<bool> &used) {
    while(!adi[u].emptv()) {
      const edge e = adj[u].back();
      if(!used[e.id]) {
        used[e.id] = true;
        adj[u].pop_back();
        build(e.v, tour, adj, used);
        adj[u].pop_back();
    tour.push back(u);
  /// Auxiliary function to build the eulerian tour/path.
  vector<int> set_build(vector<vector<edge>> &adj, const int E, const int
  first) {
    vector<int> path;
    vector<bool> used(E + 3);
    build(first, path, adj, used);
    for (int i = 0; i < adj.size(); i++)
      // if there are some remaining edges, it's not possible to build the
      if(adi[i].size())
        return vector<int>();
    reverse(path.begin(), path.end());
    return path:
/// All vertices v should have in degree[v] == out degree[v]. It must not
  contain a specific
/// start and end vertices.
/// Time complexity: O(V * (log V) + E)
bool has euler tour directed (const vector < vector < edge >> & adj, const
  vector<int> &in degree) {
  const pair<bool, pair<int, int>> aux = detail::check_both_directed(adj,
  in degree);
  const bool valid = aux.first;
  const int src = aux.second.first;
  const int dest = aux.second.second;
  return (valid && src == -1 && dest == -1);
/// A directed graph has an eulerian path/tour if has:
/// - One vertex v such that out_degree[v] - in_degree[v] == 1
/// - One vertex v such that in_degree[v] - out_degree[v] == 1
/// - The remaining vertices v such that in_degree[v] == out_degree[v]
/// or
/// - All vertices v such that in_degree[v] - out_degree[v] == 0 -> TOUR
```

```
/// Returns a boolean value that indicates whether there's a path or not.
                                                                                       143
      /// If there's a valid path it also returns two numbers: the source and
                                                                                       144
         the destination.
                                                                                       145
      /// If the source and destination can be an arbitrary vertex it will
                                                                                       146
         return the pair (-1, -1)
                                                                                        1/17
      /// for the source and destination (it means the contains an eulerian
                                                                                        148
                                                                                        149
      111
 92
                                                                                        150
                                                                                                return true;
 93
      /// Time complexity: O(V + E)
                                                                                        151
      pair<bool, pair<int, int>> has euler path directed(const
                                                                                        152
         vector<vector<edge>> &adj, const vector<int> &in_degree) {
                                                                                        153
 95
         return detail::check both directed(adj, in degree);
                                                                                        154
 96
                                                                                        155
 97
                                                                                       156
 98
      /// Returns the euler path. If the graph doesn't have an euler path it
                                                                                       157
        returns an empty vector.
                                                                                       158
 99
                                                                                       159
100
      /// Time Complexity: O(V + E) for directed, O(V * log(V) + E) for
                                                                                       160
        undirected.
                                                                                       161
      /// Time Complexity: O(adj.size() + sum(adj[i].size()))
                                                                                                &degree) {
      vector<int> get euler path directed(const int E, vector<vector<edge>>
                                                                                        162
         &adi, const vector<int> &in degree) {
                                                                                        163
103
         const pair<bool, pair<int, int>> aux = has_euler_path_directed(adj,
                                                                                        164
                                                                                       165
         in degree);
         const bool valid = aux.first;
104
                                                                                       166
105
         const int src = aux.second.first;
                                                                                       167
106
        const int dest = aux.second.second;
                                                                                       168
107
                                                                                       169
108
        if(!valid)
                                                                                       170
109
          return vector<int>();
                                                                                        171
110
111
        int first;
                                                                                        172
112
        if(src != -1)
                                                                                        173
           first = src:
113
                                                                                        174
114
         else {
                                                                                       175
115
           first = 0:
           while(adj[first].empty())
                                                                                       176
116
                                                                                       177
117
             first++;
118
                                                                                       178
119
                                                                                       179
120
        return detail::set_build(adj, E, first);
                                                                                       180
121
                                                                                       181
                                                                                                  first++:
122
                                                                                        182
      /// Returns the euler tour. If the graph doesn't have an euler tour it
                                                                                       183
         returns an empty vector.
                                                                                        184
                                                                                        185
      /// Time Complexity: O(V + E)
125
                                                                                        186
      /// Time Complexity: O(adj.size() + sum(adj[i].size()))
126
      vector<int> get_euler_tour_directed(const int E, vector<vector<edge>>
                                                                                        187
         &adj, const vector<int> &in_degree) {
                                                                                        188
128
         const bool valid = has_euler_tour_directed(adj, in_degree);
                                                                                        189
129
                                                                                       190
130
        if(!valid)
131
          return vector<int>();
                                                                                        191
132
                                                                                        192
133
        int first = 0:
                                                                                        193
        while(adj[first].empty())
134
                                                                                       194
135
          first++;
                                                                                       195
136
                                                                                       196
                                                                                                if(!valid)
137
        return detail::set_build(adj, E, first);
                                                                                       197
138
                                                                                       198
139
                                                                                       199
                                                                                                int first;
140
      // The graph has a tour that passes to every edge exactly once and gets
                                                                                       200
                                                                                                if (x != -1) {
      // back to the first edge on the tour.
                                                                                       201
                                                                                                  first = x;
```

```
// A graph with an euler path has zero odd degree vertex.
// Time Complexity: O(V)
bool has euler tour undirected (const vector < int > & degree) {
  for(int i = 0; i < degree.size(); i++)
  if(degree[i] & 1)</pre>
      return false:
// The graph has a path that passes to every edge exactly once.
// It doesn't necessarely gets back to the beginning.
// A graph with an euler path has two or zero (tour) odd degree vertices.
// Returns a pair with the startpoint/endpoint of the path.
// Time Complexity: O(V)
pair<bool, pair<int, int>> has_euler_path_undirected(const vector<int>
  vector<int> odd degree;
  for (int i = 0; i < degree.size(); i++)
    if(degree[i] & 1)
      odd_degree.pb(i);
  if(odd degree.size() == 0)
    return make_pair(true, make_pair(-1, -1));
  else if (odd_degree.size() == 2)
    return make_pair(true, make_pair(odd_degree.front(),
  odd degree.back()));
    return make_pair(false, pair<int, int>());
vector<int> get euler tour undirected(const int E, const vector<int>
  &degree, vector<vector<edge>> &adi) {
  if (!has_euler_tour_undirected(degree))
    return vector<int>();
  int first = 0:
  while(adj[first].empty())
  return detail::set_build(adj, E, first);
/// Returns the euler tour. If the graph doesn't have an euler tour it
  returns an empty vector.
/// Time Complexity: O(V + E)
/// Time Complexity: O(adj.size() + sum(adj[i].size()))
vector<int> get_euler_path_undirected(const int E, const vector<int>
  &degree, vector<vector<edge>> &adj)
  auto aux = has_euler_path_undirected(degree);
  const bool valid = aux.first;
  const int x = aux.second.first;
  const int y = aux.second.second;
    return vector<int>();
```

```
adj[x].emplace_back(y, E + 1);
203
           adi[v].emplace back(x, E + 1);
204
        } else {
205
           first = 0;
206
           while(adj[first].empty())
207
            first++;
208
209
210
        vector<int> ans = detail::set_build(adj, E, first);
211
         reverse(ans.begin(), ans.end());
        if (x != -1)
212
213
          ans.pop_back();
214
        return ans;
215
216
    };
```

5.10. Bellman Ford

```
struct edge {
2
     int src, dest, weight;
3
     edge() {}
     edge(int src, int dest, int weight) : src(src), dest(dest), weight(weight)
5
6
     bool operator<(const edge &a) const {</pre>
7
       return weight < a.weight;</pre>
8
9
   };
10
   /// Works to find the shortest path with negative edges.
11
   /// Also detects cycles.
12
13 ///
14 /// Time Complexity: O(n * e)
15 /// Space Complexity: O(n)
   | bool bellman_ford(vector<edge> &edges, int src, int n) {
17
     // n = qtd of vertices, E = qtd de arestas
18
19
     // To calculate the shortest path uncomment the line below
20
     // vector<int> dist(n, INF);
21
22
     // To check cycles uncomment the line below
23
     // vector<int> dist(n, 0);
24
25
     vector<int> pai(n, -1);
26
     int E = edges.size();
27
28
     dist[src] = 0;
29
     // Relax all edges n - 1 times.
30
     // A simple shortest path from src to any other vertex can have at-most n
       - 1 edges.
31
     for (int i = 1; i <= n - 1; i++) {
32
       for (int j = 0; j < E; j++) {
33
         int u = edges[j].src;
34
         int v = edges[j].dest;
35
         int weight = edges[j].weight;
36
         if (dist[u] != INF && dist[u] + weight < dist[v]) {
37
           dist[v] = dist[u] + weight;
38
           pai[v] = u;
39
40
41
42
     // Check for NEGATIVE-WEIGHT CYCLES.
```

```
44 // The above step quarantees shortest distances if graph doesn't contain
       negative weight cycle.
     // If we get a shorter path, then there is a cycle.
46
     bool is_cycle = false;
47
     int vert_in_cycle;
48
     for (int i = \bar{0}; i < E; i++) {
49
       int u = edges[i].src;
       int v = edges[i].dest;
50
51
       int weight = edges[i].weight;
52
       if (dist[u] != INF && dist[u] + weight < dist[v]) {
53
         is cvcle = true;
54
         pai[v] = u;
55
         vert_in_cycle = v;
56
57
58
59
     if(is_cycle) {
       for (int i = 0; i < n; i++)
60
61
         vert_in_cycle = pai[vert_in_cycle];
62
63
       vector<int> cvcle:
64
       for(int v = vert_in_cycle; (v != vert_in_cycle || cycle.size() <= 1) ; v</pre>
       = pai[v])
65
         cycle.pb(v);
66
67
       reverse(cycle.begin(), cycle.end());
68
69
       for(int x: cycle) {
70
         cout << x + 1 << ' ';
71
72
       cout << cycle.front() + 1 << endl;</pre>
73
       return true;
74
     } else
75
       return false;
```

5.11. De Bruijn Sequence

```
1 // We can solve this problem by constructing a directed graph with
2 // k^(n-1) nodes with each node having k outgoing edges_order. Each node
3 // corresponds to a string of size n-1. Every edge corresponds to one of the
 4 // characters in A and adds that character to the starting string. For
5 \mid // \text{ if } n=3 \text{ and } k=2, then we construct the following graph:
6
7 1/
                 - 1 -> (01) - 1 ->
                        ^ |
8 | //
                         1 0
9 // 0 -> (00)
                                      (11) <- 1
                        l v
10 //
11 //
                <- 0 - (10) <- 0 -
12
13 // The node '01' is connected to node '11' through edge '1', as adding '1' to
14 // '01' (and removing the first character) gives us '11'.
15 //
16 // We can observe that every node in this graph has equal in-degree and
17 // out-degree, which means that a Eulerian circuit exists in this graph.
19 namespace graph {
20 namespace detail {
21 // Finding an valid eulerian path
22 void dfs (const string &node, const string &alphabet, set<string> &vis,
23
           string &edges_order) {
24
    for (char c : alphabet) {
```

```
string nxt = node + c;
       if (vis.count(nxt))
26
27
         continue;
28
29
       vis.insert(nxt);
30
       nxt.erase(nxt.begin());
31
       dfs(nxt, alphabet, vis, edges_order);
32
       edges_order += c;
33
34
35
   }; // namespace detail
36
   // Returns a string in which every string of the alphabet of size n appears
   // the resulting string exactly once.
38
39
40
   // Time Complexity: O(alphabet.size() ^ n * log2(alphabet.size() ^ n))
   string de_bruijn(const int n, const string &alphabet) {
42
     set <string> vis;
43
     string edges_order;
44
45
     string starting node = string(n - 1, alphabet.front());
46
     detail::dfs(starting_node, alphabet, vis, edges_order);
47
48
     return edges_order + starting_node;
49
   }; // namespace graph
```

5.12. Dijkstra + Dij Graph

```
class Dijkstra {
    private:
     int src, dest;
     int n;
5
     int calculate(vector<vector<ii>>> &adj) {
8
       dist.resize(this->n + 1, INF);
9
       parent.resize(this->n + 1);
1.0
       vector<int> vis(this->n + 1, 0);
11
12
       for (int i = 0; i <= this->n; i++)
13
         parent[i].pb(i);
14
15
       priority_queue<ii, vector<ii>, greater<ii>>> pg;
16
       pg.push(make pair(0, this->src));
17
       dist[this->src] = 0;
18
19
       while(!pq.empty()) {
20
         int u = pq.top().ss;
21
          pq.pop();
22
         if(vis[u])
23
           continue;
24
          vis[u] = true;
2.5
          for (ii x: adj[u]) {
26
27
           int v = x.ff;
28
            int w = x.ss;
29
            if(dist[u] + w < dist[v]) {</pre>
30
31
             parent[v].clear();
32
              parent[v].pb(u);
33
              dist[v] = dist[u] + w;
34
              pq.push(ii(dist[v], v));
```

```
36
            else if(dist[u] + w == dist[v]) {
37
              parent[v].pb(u);
38
39
40
41
42
       return dist[dest];
43
44
45
     // Use a vector vis in a DFS on the dijkstra graph
46
     vector<vector<int>> gen_dij_graph() {
47
       vector<vector<int>> dijkstra_graph(this->n + 1);
48
       vector<bool> vis(this->n + 1);
49
50
        queue<int> q;
51
        q.push (this->dest);
52
        while(!q.empty()) {
53
         int v = q.front();
54
          q.pop();
55
56
          for(int u: parent[v]) {
57
           if(u == v)
58
              continue;
59
            dijkstra_graph[u].pb(v);
60
           if(!vis[u]) {
61
              q.push(u);
62
              vis[u] = true;
63
64
6.5
66
       return dijkstra_graph;
67
68
    public:
69
70
     int min path;
71
     vector<int> dist:
72
     vector<vector<int>> parent;
73
     vector<vector<int>> dij_graph;
74
75
     Dijkstra(int n, int src, int dest, vector<vector<ii>>> &adj) {
       this -> n = n:
76
77
       this->src = src;
78
       this->dest = dest;
79
       this->min path = this->calculate(adj);
80
        /// Generates the dijkstra graph with the parent vector
81
       this->dij_graph = this->gen_dij_graph();
82
       d_graph = this->dij_graph;
83
84
85
     // Returns a path with minimum costs and a minimum length.
86
     vector<int> get_min_path() {
87
       vector<int> path;
88
       vector<int> pai(this->n + 1, -1);
89
       vector<int> d(this->n + 1, INF);
90
91
       queue<int> q:
92
        q.push (this->dest);
93
       d[this->dest] = 0;
94
95
       while(!q.empty()) {
96
         int v = q.front();
97
          q.pop();
98
99
          for(int u: parent[v]) {
```

43

44

45

46

47

48

49

50

51

52

53

54

55

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60

61 62

63 64 65

66 67

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82

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84

85

86

87

88

89

90

91

92

93

94

95

96

97

98

99

100

101

102

103

104

105

106

int mat[sz][sz];

```
if(u == v)
100
101
               continue;
102
             if(d[v] + 1 < d[u]) {
103
               d[u] = d[v] + 1;
104
               pai[u] = v;
105
               q.push(u);
106
107
108
109
110
         int cur = this->src;
         while (cur !=-1) {
111
112
          path.pb(cur);
113
           cur = pai[cur];
114
115
116
         return path;
117
118
    };
```

5.13. Dinic (Max Flow)

```
// Created by Ubiratan Neto
3
   struct Dinic {
4
5
     struct FlowEdge {
       int v, rev, c, cap;
       bool is_rev;
8
       FlowEdge() {}
9
       FlowEdge (int v, int c, int cap, int rev, bool is_rev) : v(v), c(c),
       cap(cap), rev(rev), is_rev(is_rev) {}
10
     };
11
12
     vector<vector<FlowEdge>> adj;
13
     vector<int> level, used;
14
     int src, snk, V;
15
     int sz;
     int max_flow;
16
17
     bool calculated:
18
     Dinic(){}
     Dinic(int n) {
19
20
       calculated = false:
21
       src = 0;
22
       snk = n+1;
23
       adj.resize(n+2, vector< FlowEdge >());
       level.resize(n+2);
24
25
       used.resize(n+2):
26
       sz = n+2;
27
       V = n+2;
       max_flow = 0;
28
29
30
     void add_edge(int u, int v, int c){
31
32
       int id1 = adj[u].size();
33
       int id2 = adj[v].size();
       adj[u].pb(FlowEdge(v, c, c, id2, false));
34
35
       adj[v].pb(FlowEdge(u, 0, 0, id1, true));
36
37
38
     void add_to_src(int v, int c){
39
       adj[src].pb(FlowEdge(v, c, c, -1, false));
40
41
```

```
void add_to_snk(int u, int c){
 adj[u].pb(FlowEdge(snk, c, c, -1, false));
bool bfs() {
  for(int i=0; i<sz; i++) {</pre>
    level[i] = -1;
 level[src] = 0;
  queue<int> q; q.push(src);
  while(!q.emptv()){
    int cur = q.front();
    q.pop();
    for(FlowEdge e : adj[cur]) {
      if(level[e.v] == -1 && e.c > 0){
        level[e.v] = level[cur]+1;
        q.push(e.v);
  return (level[snk] == -1 ? false : true);
int send_flow(int u, int flow) {
  if(u == snk) return flow;
  for(int &i = used[u]; i<adj[u].size(); i++){</pre>
    FlowEdge &e = adj[u][i];
    if(level[u]+1 != level[e.v] || e.c <= 0) continue;</pre>
    int new_flow = min(flow, e.c);
    int adjusted flow = send flow(e.v, new flow);
    if(adjusted_flow > 0) {
      e.c -= adjusted_flow;
      if (e.rev != -1) adj[e.v][e.rev].c += adjusted_flow;
      return adjusted_flow;
 }
  return 0;
int calculate()
  if(src == snk) {max flow = -1; return - 1;} //not sure if needed
  max_flow = 0;
  while(bfs()){
    for(int i=0; i<sz; i++) used[i] = 0;</pre>
    while(int inc = send_flow(src, INF)) max_flow += inc;
  calculated = true;
  return max_flow;
vector<ii> mincut (vector<vector<int>> &mat_adj) {
  assert (calculated);
```

```
memset(mat, 0, sizeof mat);
107
         for(int i = 0; i < V; i++)</pre>
108
109
           for(FlowEdge x: adj[i])
110
            mat[i][x.v] += x.c;
111
112
        vector<bool> vis(sz);
113
         queue<int> q;
114
         q.push(src);
         vis[src] = true;
115
         while(!q.empty()){
116
117
          int u = q.front();
118
           q.pop();
119
           for(int v = 0; v < sz; v++) {
             if(mat[u][v] > 0 && !vis[v]) {
120
121
               q.push(v);
122
               vis[v] = true;
123
124
125
126
         vector<ii> cut:
127
128
         for (int i = 0; i < sz; i++)
129
           for(int i = 0; i < sz; i++)
             if(vis[i] && !vis[j])
130
131
               // if there's an edge from i to j.
132
               if (mat_adj[i][j] > 0)
133
                 cut.emplace_back(i, j);
134
135
        return cut;
136
137
138
      vector<ii> min_edge_cover() {
139
        bool covered[sz];
140
         for(int i=0; i<sz; i++) covered[i] = false;</pre>
         vector< ii > edge_cover;
141
         for (int i=1; i<sz-1; i++)
142
143
           for(FlowEdge e : adj[i]){
             if (e.cap == 0 \mid \mid e.v > sz-2) continue;
144
145
             if(e.c == 0){
146
               edge_cover.pb(ii(i, e.v));
147
               covered[i] = true;
148
               covered[e.v] = true;
149
               break;
150
151
152
153
         for (int i=1: i < sz-1: i++) {
154
           for(FlowEdge e : adj[i]) {
155
             if (e.cap == 0 \mid \mid e.v > sz-2) continue;
156
             if(e.c == 0) continue;
             if(!covered[i] || !covered[e.v]){
157
158
               edge cover.pb(ii(i, e.v));
159
               covered[i] = true;
160
               covered[e.v] = true;
161
162
163
164
        return edge_cover;
165
166
167
      vector<vector<int>> allFlow() {
168
         assert (calculated):
169
         vector<vector<int>> ret(V, vector<int>(V, 0));
170
171
         for (int i = 0; i < V; i++) {
```

```
for(FlowEdge x: adj[i]) {
172
173
            if(x.is rev)
174
              continue:
175
            // flow from vertex i to x.v
176
            ret[i][x.v] += x.cap - x.c;
177
178
179
180
        // for (int i = 0; i < V; i++) {
        181
182
183
        //
184
             cout << endl;
185
186
187
        return ret;
188
189
190
      void dfs_build_path(int u, vector<int> &path, vector<vector<int>>
        &mat_flow, vector<vector<int>> &ans, vector<vector<int>> &adj) {
191
        path.pb(u);
192
193
        if(u == this -> snk) {
194
          ans.pb(path);
195
          return;
196
197
198
        for(int v: adj[u]) {
199
          if(mat_flow[u][v]) {
200
            mat_flow[u][v]--;
201
            dfs_build_path(v, path, mat_flow, ans, adj);
202
            return:
203
204
205
206
207
      vector<vector<int>> get_all_paths(vector<vector<int>> &adj) {
        assert (calculated);
208
209
        vector<vector<int>> mat_flow = allFlow();
210
211
        vector<vector<int>> ans:
212
        ans.reserve(max flow);
213
214
        for (int i = 0; i < max flow; i++) {
215
          vector<int> path;
216
          path.reserve(V);
217
          dfs_build_path(this->src, path, mat_flow, ans, adj);
218
219
220
        return ans;
221
222 };
```

5.14. Functional Graph

```
// Based on:
   http://maratona.ic.unicamp.br/MaratonaVerao2020/lecture-b/20200122.pdf

class Functional_Graph {
   // FOR DIRECTED GRAPH
   private:
   void compute_cycle(int u, vector<int> &nxt, vector<bool> &vis) {
      int id_cycle = cycle_cnt++;
      int cur_id = 0;
}
```

this->first[id_cycle] = u;

9

```
10
11
        while(!vis[u]) {
12
          vis[u] = true;
13
14
          this->cycle[id_cycle].push_back(u);
15
16
          this->in_cycle[u] = true;
          this->cycle_id[u] = id_cycle;
17
18
          this->id in cycle[u] = cur id;
19
          this->near in cycle[u] = u;
          this->id_near_cycle[u] = id_cycle;
20
21
          this->cvcle dist[u] = 0;
22
23
          u = nxt[u];
24
          cur id++;
25
26
27
28
     // Time Complexity: O(V)
     void build(int n, int indexed from, vector<int> &nxt, vector<int>
        &in degree) {
        queue<int> q;
30
31
        vector<bool> vis(n + indexed_from);
32
        for(int i = indexed_from; i < n + indexed_from; i++) {</pre>
         if(in_degree[i] == 0) {
33
34
            q.push(i);
35
            vis[i] = true;
36
37
38
39
        vector<int> process_order;
40
        process order.reserve(n + indexed from);
41
        while(!q.emptv()) {
42
          int u = q.front();
43
          ; () qoq.p
44
45
          process_order.push_back(u);
46
47
          if(--in_degree[nxt[u]] == 0) {
48
            a.push(nxt[u]);
49
            vis[nxt[u]] = true;
50
51
52
53
        int cycle_cnt = 0;
54
        for(int i = indexed from; i < n + indexed from; i++)</pre>
55
         if(!vis[i])
56
            compute_cycle(i, nxt, vis);
57
58
        for(int i = (int)process_order.size() - 1; i >= 0; i--) {
59
          int u = process_order[i];
60
          this->near_in_cycle[u] = this->near_in_cycle[nxt[u]];
this->id_near_cycle[u] = this->id_near_cycle[nxt[u]];
61
62
63
          this->cycle_dist[u] = this->cycle_dist[nxt[u]] + 1;
64
65
66
67
     void allocate(int n, int indexed_from) {
68
        this->cycle.resize(n + indexed_from);
69
       this->first.resize(n + indexed_from);
70
        this->in_cycle.resize(n + indexed_from, false);
71
72
        this->cycle_id.resize(n + indexed_from, -1);
```

```
this->id_in_cycle.resize(n + indexed_from, -1);
        this->near in cycle.resize(n + indexed from);
 74
 75
        this->id near cycle.resize(n + indexed from);
 76
        this->cycle_dist.resize(n + indexed_from);
 77
 78
 79
     public:
      Functional_Graph(int n, int indexed_from, vector<int> &nxt, vector<int>
 8.0
         &in degree) {
 81
         this->allocate(n, indexed from);
 82
        this->build(n, indexed_from, nxt, in_degree);
 83
 84
 85
      // THE CYCLES ARE ALWAYS INDEXED BY ZERO!
 86
 87
      // number of cycles
      int cycle_cnt = 0;
      // Vertices present in the i-th cycle.
 90
      vector<vector<int>> cvcle;
      // first vertex of the i-th cycle
      vector<int> first:
 9.3
 94
      // The i-th vertex is present in any cycle?
 95
      vector<bool> in cvcle:
 96
      // id of the cycle that the vertex belongs. -1 if it doesn't belong to any
         cycle.
 97
      vector<int> cvcle id;
      // Represents the id of the cycle of the i-th vertex. -1 if it doesn't
         belong to any cycle.
      vector<int> id_in_cycle;
 99
      // Represents the id of the nearest vertex present in a cycle.
100
101
      vector<int> near_in_cycle;
102
      // Represents the id of the nearest cycle.
103
      vector<int> id near cycle;
      // Distance to the nearest cycle.
104
      vector<int> cycle_dist;
105
      // Represent the id of the component of the vertex.
106
107
      // Equal to id_near_cycle
      vector<int> &comp = id_near_cycle;
108
109
    };
110
111 class Functional Graph {
112 // FOR UNDIRECTED GRAPH
     private:
114
      void compute_cycle(int u, vector<int> &nxt, vector<bool> &vis,
         vector<vector<int>> &adj) {
115
        int id cycle = cycle cnt++;
        int cur_id = 0;
116
117
        this->first[id_cycle] = u;
118
119
         while(!vis[u]) {
120
          vis[u] = true;
121
122
           this->cycle[id_cycle].push_back(u);
123
           nxt[u] = find_nxt(u, vis, adj);
124
           if(nxt[u] == -1)
125
            nxt[u] = this->first[id cvcle];
126
127
           this->in_cycle[u] = true;
          this->cycle_id[u] = id_cycle;
this->id_in_cycle[u] = cur_id;
128
129
130
           this->near_in_cycle[u] = u;
131
           this->id_near_cycle[u] = id_cycle;
132
           this->cycle_dist[u] = 0;
133
```

u = nxt[u]:

```
135
          cur id++;
136
137
138
139
      int find_nxt(int u, vector<bool> &vis, vector<vector<int>> &adj) {
140
        for(int v: adj[u])
          if(!vis[v])
141
142
            return v:
143
        return -1:
144
145
146
      // Time Complexity: O(V + E)
      void build(int n, int indexed_from, vector<int> &degree,
147
        vector<vector<int>> &adj) {
        queue<int> q;
148
149
        vector<bool> vis(n + indexed_from, false);
150
        vector<int> nxt(n + indexed_from);
151
        for(int i = indexed from; i < n + indexed from; i++) {</pre>
152
          if(adj[i].size() == 1) {
153
            q.push(i);
            vis[i] = true;
154
155
156
157
158
        vector<int> process_order;
159
        process order.reserve(n + indexed from);
        while(!q.empty()) {
160
161
          int u = q.front();
162
          q.pop();
163
164
          process_order.push_back(u);
165
166
          nxt[u] = find_nxt(u, vis, adj);
167
          if(--degree[nxt[u]] == 1) {
168
            q.push(nxt[u]);
169
            vis[nxt[u]] = true;
170
171
172
173
        int cvcle cnt = 0;
        for(int i = indexed_from; i < n + indexed_from; i++)</pre>
174
175
          if(!vis[i])
176
            compute_cycle(i, nxt, vis, adj);
177
178
        for(int i = (int)process_order.size() - 1; i >= 0; i--) {
179
          int u = process order[i];
180
          this->near_in_cycle[u] = this->near_in_cycle[nxt[u]];
181
182
          this->id_near_cycle[u] = this->id_near_cycle[nxt[u]];
183
          this->cycle_dist[u] = this->cycle_dist[nxt[u]] + 1;
184
185
186
187
      void allocate(int n, int indexed from) {
188
        this->cvcle.resize(n + indexed from);
189
        this->first.resize(n + indexed from);
190
191
        this->in_cycle.resize(n + indexed_from, false);
192
        this->cycle_id.resize(n + indexed_from, -1);
        this->id_in_cycle.resize(n + indexed_from, -1);
193
194
        this->near_in_cycle.resize(n + indexed_from);
195
        this->id_near_cycle.resize(n + indexed_from);
        this->cycle_dist.resize(n + indexed_from);
196
197
```

```
199
     public:
200
     Functional_Graph(int n, int indexed_from, vector<int> degree,
        vector<vector<int>> &adj) {
201
        this->allocate(n, indexed from);
202
        this->build(n, indexed_from, degree, adj);
203
204
      // THE CYCLES ARE ALWAYS INDEXED BY ZERO!
205
206
207
      // number of cycles
      int cycle_cnt = 0;
208
209
      // Vertices present in the i-th cycle.
      vector<vector<int>> cycle;
210
      // first vertex of the i-th cycle
211
212
      vector<int> first;
213
214
      // The i-th vertex is present in any cycle?
215
      vector<bool> in cycle;
      // id of the cycle that the vertex belongs. -1 if it doesn't belong to any
        cvcle.
217
      vector<int> cvcle id;
      // Represents the id of the cycle of the i-th vertex. -1 if it doesn't
        belong to any cycle.
219
      vector<int> id_in_cycle;
      // Represents the id of the nearest vertex present in a cycle.
220
      vector<int> near_in_cycle;
2.2.1
222
      // Represents the id of the nearest cycle.
      vector<int> id_near_cycle;
223
224
      // Distance to the nearest cycle.
      vector<int> cycle dist;
      // Represent the id of the component of the vertex.
226
227
      // Equal to id near cycle
228
     vector<int> &comp = id near cycle;
229 };
```

5.15. Hld

```
1 class HLD {
    private:
3
    int n;
     // number of nodes below the i-th node
     vector<int> sz:
7
    private:
     int get_sz(const int u, const int p, const vector<vector<int>> &adj) {
9
       this -> sz[u] = 1;
10
        for(const int v: adj[u]) {
         if(v == p)
11
12
           continue:
         this->sz[u] += this->get_sz(v, u, adj);
13
14
15
       return this->sz[u];
16
17
     void dfs(const int u, const int id, const int p, const vector<vector<int>>
        &adi) {
19
       this->chain id[u] = id;
        this->id_in_chain[u] = chain_size[id];
20
21
       this->parent[u] = p;
22
23
       if(this->chain_head[id] == -1)
24
         this->chain_head[id] = u;
25
        this->chain size[id]++;
```

```
27
       int maxx = -1, idx = -1;
       for(const int v: adi[u]) {
28
29
         if(v == p)
30
           continue;
31
         if(sz[v] > maxx) {
32
          maxx = sz[v];
33
           idx = v:
34
35
36
37
       if(idx != -1)
38
         this->dfs(idx, id, u, adj);
39
40
       for(const int v: adj[u]) {
         if(v == idx \mid | v == p)
41
42
           continue;
43
         this->dfs(v, this->number of chains++, u, adj);
44
45
46
    public:
    /// Builds the paths.
     /// Time Complexity: O(n)
50
     HLD(const vector<vector<int>> &adj) {
       this->n = adj.size();
52
53
       this->chain_head.resize(this->n + 1, -1);
54
       this->id_in_chain.resize(this->n + 1, -1);
       this->chain_id.resize(this->n + 1, -1);
55
56
       this->sz.resize(this->n + 1);
57
       this->chain size.resize(this->n + 1);
58
       this->parent.resize(this->n + 1, -1);
59
       this->get_sz(1, -1, adj);
       this->dfs(1, 0, -1, adj);
60
61
62
63
     // the chains are indexed from 0
     int number_of_chains = 1;
64
     // topmost node of the chain
65
     vector<int> chain head;
     // id of the i-th node in his chain
     vector<int> id in chain;
     // id of the chain that the i-th node belongs
     vector<int> chain id:
     // size of the i-th chain
     vector<int> chain size;
     // parent of the i-th node, -1 for root
74
    vector<int> parent;
75
```

5.16. Kruskal + Dsu

```
class DSU {
   public:
     vector<int> root;
     vector<int> sz;

     DSU(int n) {
        this->root.resize(n + 1);
        iota(this->root.begin(), this->root.begin() + n + 1, 0);
        this->sz.resize(n + 1, 1);
}
```

```
12 l
13
14
     int Find(int x) {
15
       if(this->root[x] == x)
16
          return x;
17
       return this->root[x] = this->Find(this->root[x]);
18
19
     bool Union(int p, int q) {
20
2.1
22
       p = this->Find(p), q = this->Find(q);
23
24
       if(p == q)
25
         return false;
26
27
        if(this->sz[p] > this->sz[q]) {
28
          this->root[q] = p;
29
          this->sz[p] += this->sz[q];
30
        } else {
31
          this->root[p] = q;
32
          this->sz[q] += this->sz[p];
33
34
35
       return true;
36
37
38
39
40
   struct edge {
41
     int u, v, w;
42
     edge() {}
43
     edge(int u, int v, int w) : u(u), v(v), w(w) {}
44
45
     bool operator<(const edge &a) const {</pre>
46
       return w < a.w;</pre>
47
48
    };
49
   int kruskal(int n, vector<edge>& edges) {
50
51
52
     DSU dsu(n);
5.3
     sort(edges.begin(), edges.end());
     int weight = 0;
     for(int i = 0; i < (int)edges.size(); i++) {</pre>
58
       if (dsu.Union(edges[i].u, edges[i].v)) {
59
          weight += edges[i].w;
60
61
62
63
     // returns weight of mst
     return weight;
```

5.17. Lca

```
// #define DIST
// #define COST
// UNCOMMENT ALSO THE LINE BELOW FOR COST!

class LCA {
  private:
  int n;
```

77

78 79

8.0

81

82

83

84

85

86

87

8.8

89

90

92

97

100

101

103

104

105

106

107

109

111

113

122

123

124

125

126

127

128

129

130

131

133

```
// INDEXED from 0 or 1??
9
     int indexed from;
1.0
     /// Store all log2 from 1 to n
11
     vector<int> lq;
     // level of the i-th node (height)
12
13
     vector<int> level;
14
     // matrix to store the ancestors of each node in power of 2 levels
1.5
     vector<vector<int>> anc;
16
17
     #ifdef DIST
18
       vector<int> dist;
19
     #endif
20
     #ifdef COST
21
       // int NEUTRAL_VALUE = -INF; // MAX COST
22
       // int combine(const int a, const int b) {return max(a, b);}
       // int NEUTRAL_VALUE = INF; // MIN COST
23
24
       // int combine(const int a, const int b) {return min(a, b);}
25
       vector<vector<int>> cost;
26
     #endif
27
    private:
28
     void allocate() {
30
       // initializes a matrix [n][lq n] with -1
31
       this->build_log_array();
32
       this->anc.resize(n + 1, vector<int>(lg[n] + 1, -1));
       this->level.resize(n + 1, -1);
33
34
35
       #ifdef DIST
36
         this->dist.resize(n + 1, 0);
37
       #endif
38
       #ifdef COST
39
         this->cost.resize(n + 1, vector<int>(lq[n] + 1, NEUTRAL_VALUE));
40
41
42
43
     void build log array() {
44
       this->lq.resize(this->n + 1);
45
46
       for(int i = 2; i <= this->n; i++)
47
         this->lq[i] = this->lq[i/2] + 1;
48
49
50
     void build anc()
       for(int j = 1; j < anc.front().size(); j++)
52
         for(int i = 0; i < anc.size(); i++)</pre>
53
           if(this->anc[i][j - 1] != -1) {
54
             this->anc[i][j] = this->anc[this->anc[i][j - 1]][j - 1];
55
56
               this->cost[i][j] = combine(this->cost[i][j - 1],
       this->cost[anc[i][j - 1]][j - 1]);
57
             #endif
5.8
59
60
     void build weighted(const vector<vector<pair<int, int>>> &adj) {
61
62
       this->dfs_LCA_weighted(this->indexed_from, -1, 1, 0, adj);
63
       this->build anc():
64
65
66
67
     void dfs_LCA_weighted(const int u, const int p, const int l, const int d,
       const vector<vector<pair<int, int>>> &adj) {
68
       this->level[u] = 1;
69
       this->anc[u][0] = p;
70
       #ifdef DIST
```

```
this->dist[u] = d;
 71
 72
         #endif
 73
 74
         for(const pair<int, int> &x: adj[u]) {
 75
          int v = x.first, w = x.second;
 76
          if(v == p)
            continue;
          #ifdef COST
            this->cost[v][0] = w;
           this->dfs_LCA_weighted(v, u, l + 1, d + w, adj);
      void build_unweighted(const vector<vector<int>> &adj) {
        this->dfs_LCA_unweighted(this->indexed_from, -1, 1, 0, adj);
        this->build anc();
      void dfs_LCA_unweighted(const int u, const int p, const int l, const int
         d, const vector<vector<int>> &adj) {
        this->level[u] = 1;
 93
        this->anc[u][0] = p;
 94
        #ifdef DIST
 95
          this->dist[u] = d;
 96
         #endif
 98
         for(const int v: adj[u]) {
 99
          if(v == p)
           this->dfs_LCA_unweighted(v, u, l + 1, d + 1, adj);
102
      // go up k levels from x
      int lca_go_up(int x, int k) {
        for (int i = 0; k > 0; i++, k >>= 1)
108
          if(k & 1) {
            x = this \rightarrow anc[x][i];
110
            if(x == -1)
               return -1:
112
114
        return x;
115
116
      #ifdef COST
117
      /// Query between the an ancestor of v (p) and v. It returns the
118
      /// max/min edge between them.
119
120
      int lca_query_cost_in_line(int v, int p) {
121
        assert(this->level[v] >= this->level[p]);
        int k = this->level[v] - this->level[p];
        int ans = NEUTRAL VALUE;
         for(int i = 0; k > 0; i++, k >>= 1)
          if(k & 1) {
            ans = combine(ans, this->cost[v][i]);
            v = this->anc[v][i];
132
        return ans;
134
      #endif
```

```
135
136
      int get lca(int a, int b) {
137
        // a is below b
138
        if(this->level[b] > this->level[a])
139
          swap(a,b);
140
141
        const int logg = lg[this->level[a]];
142
        // putting a and b in the same level
143
144
        for(int i = logg; i >= 0; i--)
          if(this->level[a] - (1 << i) >= this->level[b])
145
            a = this->anc[a][i];
146
147
148
        if(a == b)
149
          return a;
150
151
        for(int i = logg; i >= 0; i--)
152
          if(this->anc[a][i] != -1 && this->anc[a][i] != this->anc[b][i]) {
153
            a = this->anc[a][i];
154
            b = this->anc[b][i];
155
156
157
        return anc[a][0];
158
159
160
     public:
      /// Builds an weighted graph.
161
162
163
      /// Time Complexity: O(n*log(n))
164
      explicit LCA(const vector<vector<pair<int, int>>> &adj, const int
        indexed from) {
165
        this->n = adj.size();
166
        this->indexed from = indexed from;
167
        this->allocate();
168
169
        this->build weighted(adj);
170
171
172
      /// Builds an unweighted graph.
173
174
      /// Time Complexity: O(n*log(n))
      explicit LCA(const vector<vector<int>> &adj, const int indexed_from) {
175
176
        this->n = adi.size();
177
        this->indexed from = indexed from:
178
        this->allocate();
179
180
        this->build_unweighted(adj);
181
182
183
      /// Goes up k levels from v. If it passes the root, returns -1.
184
185
      /// Time Complexity: O(log(k))
186
      int go_up(const int v, const int k) {
        assert(indexed from <= v); assert(v < this->n + indexed_from);
187
188
189
        return this->lca_go_up(v, k);
190
191
192
      /// Returns the parent of v in the LCA dfs from 1.
193
194
      /// Time Complexity: O(1)
195
      int parent(int v) {
        assert(indexed_from <= v); assert(v < this->n + indexed_from);
196
197
198
        return this->anc[v][0];
```

```
199
200
201
      /// Returns the LCA of a and b.
202
203
      /// Time Complexity: O(log(n))
204
      int query_lca(const int a, const int b) {
205
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
        indexed from);
206
207
        return this->get_lca(a, b);
208
209
210
      /// Returns the distance from a to b. When the graph is unweighted, it is
211
        considered
212
      /// 1 as the weight of the edges.
213
214
      /// Time Complexity: O(log(n))
215
      int query dist(const int a, const int b) {
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
         indexed from):
217
218
        return this->dist[a] + this->dist[b] - 2*this->dist[this->get lca(a, b)];
219
220
      #endif
221
222
       #ifdef COST
223
      /// Returns the max/min weight edge from a to b.
224
      /// Time Complexity: O(log(n))
225
226
      int query_cost(const int a, const int b) {
227
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
         indexed from);
228
        const int l = this->query_lca(a, b);
229
230
        return combine (this->lca query cost in line (a, 1),
        this->lca_query_cost_in_line(b, 1));
231
232
      #endif
233 };
```

5.18. Maximum Path Unweighted Graph

```
/// Returns the maximum path between the vertices 0 and n-1 in a
       unweighted graph.
   /// Time Complexity: O(V + E)
  int maximum_path(int n) {
4
5
     vector<int> top_order = topological_sort(n);
     vector<int> pai(n, -1);
     if(top_order.empty())
8
       return -1;
10
     vector<int> dp(n);
11
     dp[0] = 1:
     for(int u: top order)
12
13
       for(int v: adj[u])
14
         if (dp[u] & dp[u] + 1 > dp[v]) {
15
           dp[v] = dp[u] + 1;
16
           pai[v] = u;
17
1.8
19
     if (dp[n-1] == 0)
20
       return -1;
```

```
22
     vector<int> path;
23
     int cur = n - 1:
     while (cur != -1) {
24
25
       path.pb(cur);
26
        cur = pai[cur];
27
28
     reverse(path.begin(), path.end());
29
30
     // cout << path.size() << endl;</pre>
31
     // for(int x: path) {
     // cout << x + 1 << ' ';
32
33
34
     // cout << endl;
35
36
     return dp[n - 1];
```

5.19. Number Of Different Spanning Trees In A Complete Graph

```
1 Cayley's formula
2 n ^ (n - 2)
```

5.20. Number Of Ways To Make A Graph Connected

```
1 s_{1} \star s_{2} \star s_{3} \star \ldots \star s_{k} \star (n \hat{k} - 2)

2 n = \text{number of vertices}

3 s_{i} = \text{size of the } i-\text{th connected component}

4 k = \text{number of connected components}
```

5.21. Pruffer Decode

```
// IT MUST BE INDEXED BY 0.
   /// Returns the adjacency matrix of the decoded tree.
3 ///
   /// Time Complexity: O(V)
   vector<vector<int>> pruefer_decode(const vector<int> &code) {
     int n = code.size() + 2;
     vector<vector<int>> adj = vector<vector<int>> (n, vector<int>());
8
     vector<int> degree(n, 1);
10
     for (int x : code)
11
       degree[x]++;
12
13
     int ptr = 0:
14
     while (degree[ptr] > 1)
15
       ++ptr;
16
17
     int nxt = ptr;
18
     for (int u : code) {
19
       adj[u].push_back(nxt);
2.0
       adj[nxt].push_back(u);
21
       if (--degree[u] == 1 && u < ptr)</pre>
22
23
         nxt = u;
24
25
         while (degree[++ptr] > 1)
26
27
         nxt = ptr;
28
29
```

```
30    adj[n - 1].push_back(nxt);
31    adj[nxt].push_back(n - 1);
32
33    return adj;
34  }
```

5.22. Pruffer Encode

```
| void dfs(int v, const vector<vector<int>> &adi, vector<int> &parent) {
     for (int u : adj[v]) {
       if (u != parent[v]) {
          parent[u] = v;
5
          dfs(u, adj, parent);
6
7
8 }
9
10 // IT MUST BE INDEXED BY 0.
11 /// Returns prueffer code of the tree.
12 ///
13 /// Time Complexity: O(V)
   vector<int> pruefer code(const vector<vector<int>> &adi) {
     int n = adj.size();
     vector<int> parent(n);
17
     parent[n-1] = -1;
     dfs(n - 1, adj, parent);
18
19
20
     int ptr = -1;
     vector<int> degree(n);
21
2.2
     for (int i = 0; i < n; i++) {
23
       degree[i] = adj[i].size();
24
       if (degree[i] == 1 && ptr == -1)
25
         ptr = i;
26
27
     vector<int> code(n - 2);
28
29
     int leaf = ptr;
     for (int i = 0; i < n - 2; i++) {
30
31
       int next = parent[leaf];
32
       code[i] = next;
3.3
       if (--degree[next] == 1 && next < ptr)</pre>
34
         leaf = next;
35
36
          ptr++;
37
          while (degree[ptr] != 1)
38
           ptr++;
39
         leaf = ptr;
40
41
42
43
     return code;
```

5.23. Pruffer Properties

- 1 * After constructing the Prüfer code two vertices will remain. One of them is the highest vertex n-1, but nothing **else** can be said about the other one.
- 2 * Each vertex appears in the Prüfer code exactly a fixed number of times its degree minus one. This can be easily checked, since the degree will get smaller every time we record its label in the code, and we remove it once the degree is 1. For the two remaining vertices this fact is also true.

5.24. Remove All Bridges From Graph

5.25. Shortest Cycle In A Graph

```
int bfs(int vt) {
2
     vector<int> dist(MAXN, INF);
     queue<pair<int, int>> q;
5
     q.emplace(vt, -1);
     dist[vt] = 0;
8
9
     int ans = INF;
10
     while (!q.empty()) {
11
       pair<int, int> aux = q.front();
       int u = aux.first, p = aux.second;
12
13
       q.pop();
14
15
        for (int v : adj[u]) {
16
         if (v == p)
17
           continue;
18
         if (dist[v] < INF)</pre>
19
           ans = min(ans, dist[u] + dist[v] + 1);
20
          else {
21
           dist[v] = dist[u] + 1;
22
            q.emplace(v, u);
23
24
25
26
27
     return ans;
28
29
30
   /// Returns the shortest cycle in the graph
31
   ///
32
   /// Time Complexity: O(V^2)
33 int get_girth(int n) {
     int ans = INF;
34
35
     for (int u = 1; u <= n; u++)
36
       ans = min(ans, bfs(u));
37
     return ans;
38
```

5.26. Topological Sort

```
/// INDEXED BY ZERO
2
   111
   /// Time Complexity: O(n)
   vector<int> topological_sort(int n) {
     vector<int> in_degree(n, 0);
     for (int u = 0; u < n; u++)
8
       for(int v: adi[u])
9
         in_degree[v]++;
1.0
11
     queue<int> q;
     for(int i = 0; i < n; i++)
12
13
       if(in_degree[i] == 0)
14
         q.push(i);
```

```
15
16
      int cnt = 0;
17
      vector<int> top_order;
18
      while(!q.empty()) {
19
       int u = q.front();
20
       q.pop();
21
22
        top_order.push_back(u);
23
        cnt++;
24
25
        for(int v: adj[u])
26
          if(--in\_degree[v] == 0)
27
            q.push(v);
28
29
30
      if(cnt != n) {
        cerr << "There exists a cycle in the graph" << endl;</pre>
31
32
        return vector<int>();
33
34
     return top_order;
35
36 }
```

5.27. Tree Distance

```
1 | vector<pair<int, int>> sub(MAXN, pair<int, int>(0, 0));
3 void subu(int u, int p) {
     for (const pair<int, int> x : adj[u]) {
       int v = x.first, w = x.second;
       if (v == p)
7
         continue;
8
       subu(v, u);
       if (sub[v].first + w > sub[u].first) {
         swap(sub[u].first, sub[u].second);
10
11
         sub[u].first = sub[v].first + w;
12
       } else if (sub[v].first + w > sub[u].second) {
13
         sub[u].second = sub[v].first + w;
14
15
16
17
   /// Contains the maximum distance to the node i
19
   vector<int> ans(MAXN);
20
21
   void dfs(int u, int d, int p) {
22
     ans[u] = max(d, sub[u].first);
     for (const pair<int, int> x : adj[u]) {
23
24
       int v = x.first, w = x.second;
25
       if (v == p)
26
         continue;
27
       if (sub[v].first + w == ans[u]) {
28
         dfs(v, max(d, sub[u].second) + w, u);
29
       } else {
3.0
         dfs(v, ans[u] + w, u);
31
32
33 }
35 // Returns the maximum tree distance
36 | int solve() {
37
    subu(0, -1);
38
     dfs(0, 0, -1);
     return *max_element(ans.begin(), ans.end());
```

40 }

- 6. Language Stuff
- 6.1. Binary String To Int

```
int y = bitset<qtdDeBits>(stringVar).to_ulong();
Ex: x = 1010, qtdDeBits = 32;
y = bitset<32>(x).to_ulong(); // y = 10
```

6.2. Climits

```
LONG_MIN -> (-2^31+1) :: LONG_MAX -> (2^31-1)
ULONG_MAX -> (2^32-1) -> UNSIGNED
LLONG_MIN, LLONG_MAX, ULLONG_MAX
```

6.3. Checagem Brute Force Com Solucao

```
$ g++ -std=c++11 gen.cpp && ./a.out > gen.out && g++ -std=c++11 brute.cpp && (./a.out < gen.in) > brute.out && g++ -std=c++11 sol.cpp && (./a.out < gen.in) > sol.out && diff brute.out sol.out
```

6.4. Checagem De Bits

```
// OBS: SO FUNCIONA PARA INT (NAO FUNCIONA COM LONG LONG)

__builtin_popcount(int) -> Número de bits ativos;

__builtin_ctz(int) -> Número de zeros à direita

__builtin_clz(int) -> Número de zeros à esquerda

__builtin_parity(int) -> Retorna se a quantidade de uns é impar(1) ou par(0)
```

6.5. Checagem E Tranformacao De Caractere

```
#include <cctype>
isdigit(str[i]);//checa se str[i] é número
isalpha(str[i]);//checa se é uma letra
islower(str[i]);//checa minúsculo
isupper(str[i]);//checa maiúsculo
isalnum(str[i]);//checa letra ou número
tolower(str[i]);//converte para minusculo
toupper(str[i]);//converte para maiusculo
```

6.6. Conta Digitos 1 Ate N

```
int solve(int n) {
   int maxx = 9, minn = 1, dig = 1, ret = 0;

   for(int i = 1; i <= 17; i++) {
      int q = min(maxx, n);
      ret += max(011, (q - minn + 1) * dig);
      maxx = (maxx * 10 + 9), minn *= 10, dig++;
   }

   return ret;
}</pre>
```

6.7. Escrita Em Arquivo

```
1 ofstream cout("output.txt");
```

6.8. Gcd

```
1 int _gcd(int a, int b) {
2    if(a == 0 || b == 0) return 0;
3    else return abs(__gcd(a,b));
4  }
```

6.9. Hipotenusa

```
1 cout << hypot(3,4); // output: 5
```

6.10. Int To Binary String

```
string s = bitset<qtdDeBits>(intVar).to_string();
Ex: x = 10, qtdDeBits = 32;
s = bitset<32>(x).to_string(); // s = 00...0001010
```

6.11. Int To String

```
int a; string b;
b = to_string(a);
```

6.12. Leitura De Arquivo

```
1 ifstream cin("input.txt");
```

6.13. Max E Min Element Num Vetor

```
int maior = *max_element(arr.begin(), arr.end());
int menor = *min_element(arr.begin(), arr.end());
// OBS: Retorna iterador
```

6.14. Permutacao

```
int v[] = {1,2,3};
sort(v, v+3);
do {
   cout << v[0] << ' ' << v[1] ' ' << v[2];
} while(next_permutation(v, v+3));</pre>
```

6.15. Printf De Uma String

```
char buffer [50];
int n, a=5, b=3;
n=sprintf (buffer, "%d plus %d is %d", a, b, a+b);
printf ("[%s] is a string %d chars long\n",buffer,n);
// Output:
// [5 plus 3 is 8] is a string 13 chars long
```

6.16. Remove Repeticoes Continuas Num Vetor

```
1  // arr = {10,20,20,20,30,20,10}
2  it = unique(arr.begin(), arr.end());
3  // arr = {10,20,30,20,10, iterator aponta pra aqui, ...}
4  arr.resize(distance(arr.begin(), it));
5  // arr = {10,20,30,20,10}
```

6.17. Rotate (Left)

```
Passado o inicio o meio e o fim ele rotaciona de forma que o meio seja o novo inicio.

vector<int> arr(n); // 1 2 3 4 5 6 7 8 9
rotate(arr.begin(),arr.begin()+3,arr.end()); //4 5 6 7 8 9 1 2 3
```

6.18. Rotate (Right)

```
vector<int> arr(n); // 1 2 3 4 5 6 7 8 9
rotate(arr.begin(),arr.rbegin()+3,arr.rend()); //7 8 9 1 2 3 4 5 6
```

6.19. Scanf De Uma String

```
char sentence[]="Rudolph is 12 years old";
char str [20]; int i;
sscanf (sentence, "%s %** %d", str, &i);
printf ("%s -> %d\n", str, i);
// Output: Rudolph -> 12
```

6.20. Split Function

```
1 // SEPARA STRING POR UM DELIMITADOR
   // EX: str=A-B-C split \rightarrow x = {A,B,C}
   vector<string> split(const string &s, char delim) {
       stringstream ss(s);
5
       string item;
6
       vector<string> tokens;
       while (getline(ss, item, delim)) {
8
           tokens.push_back(item);
9
10
       return tokens:
11
12
   int main () {
    vector<string> x = split("cap-one-best-opinion-language", '-');
   // x = {cap, one, best, opinion, language};
14
15
```

6.21. String To Long Long

```
string s = "0xFFFF"; int base = 16;
string::size_type sz = 0;
int l1 = stoll(s,&sz,base); // l1 = 65535, sz = 6;
0BS: Não precisa colocar o sz, pode colocar 0; // stoll(s,0,base);
```

6.22. Substring

```
string s = "abcdef";
s.substr(posição inicial, qtd de char(opcional));
string s2 = s.substr(3,2); // s2 = "de"
string s3 = s.substr(2); // s3 = "cdef"
```

6.23. Width

6.24. Check Overflow

```
bool __builtin_add_overflow (type1 a, type2 b, type3 *res)
  | bool __builtin_sadd_overflow (int a, int b, int *res)
   | bool __builtin_saddl_overflow (long int a, long int b, long int *res)
   bool __builtin_saddll_overflow (long long int a, long long int b, long long
       int *res)
   bool __builtin_uadd_overflow (unsigned int a, unsigned int b, unsigned int
   bool builtin uaddl overflow (unsigned long int a, unsigned long int b,
       unsigned long int *res)
7 bool __builtin_uaddll_overflow (unsigned long long int a, unsigned long long
       int b, unsigned long long int *res)
   bool __builtin_sub_overflow (type1 a, type2 b, type3 *res)
10 bool builtin ssub overflow (int a, int b, int *res)
11 | bool __builtin_ssubl_overflow (long int a, long int b, long int *res)
12 bool __builtin_ssubll_overflow (long long int a, long long int b, long long
13 bool builtin usub overflow (unsigned int a, unsigned int b, unsigned int
14 | bool | builtin usubl overflow (unsigned long int a, unsigned long int b,
       unsigned long int *res)
15 bool __builtin_usubll_overflow (unsigned long long int a, unsigned long long
       int b, unsigned long long int *res)
17 | bool __builtin_mul_overflow (type1 a, type2 b, type3 *res)
18 | bool __builtin_smul_overflow (int a, int b, int *res)
19 | bool __builtin_smull_overflow (long int a, long int b, long int *res)
20 | bool | builtin smull overflow (long long int a, long long int b, long long
       int *res)
21 | bool __builtin_umul_overflow (unsigned int a, unsigned int b, unsigned int
  bool builtin umull overflow (unsigned long int a, unsigned long int b,
       unsigned long int *res)
23 bool __builtin_umulll_overflow (unsigned long long int a, unsigned long long
       int b, unsigned long long int *res)
```

6.25. Readint

```
int readInt(){
  int a = 0;
  char c;
  while(!(c >= '0' && c <= '9') ) c = getchar();
  while(c>= '0' && c <= '9') a = 10*a + (c - '0') , c = getchar();
  return a;
}</pre>
```

7. Math

7.1. Bell Numbers

```
int bellNumber(int n) {
    int bell[n+1][n+1];
    bell[0][0] = 1;
    for (int i=1; i<=n; i++) {
        // Explicitly fill for j = 0
        bell[i][0] = bell[i-1][i-1];
        // Fill for remaining values of j
        for (int j=1; j<=i; j++)
        bell[i][j] = bell[i-1][j-1] + bell[i][j-1];
    }
    return bell[n][0];
}</pre>
```

7.2. Checagem De Primalidade

```
bool isPrime(int n) {
    if (n <= 1) return false;
    if (n <= 3) return true;

// This is checked so that we can skip
// middle five numbers in below loop

if (n%2 == 0 || n%3 == 0)

return false;

for (int i=5; i*i<=n; i += 6)
    if (n%i == 0 || n% (i+2) == 0)
    return false;

return true;

return true;
}</pre>
```

7.3. Combinação Nor Mod Primo

```
inv[1] = 1;
for(int i = 2; i < m; ++i)
    inv[i] = (m - (m/i) * inv[m%i] % m) % m;

factorial[0] = 1;
for (int i = 1; i <= MAXN; i++) {
    factorial[i] = factorial[i - 1] * i % m;
}

int binomial_coefficient(int n, int k) {
    return factorial[n] * inverse(factorial[k]) % m * inverse(factorial[n - k]) % m;
}
</pre>
```

7.4. Combinação Ncr

```
// Returns value of Binomial Coefficient C(n, k)
int binomialCoeff(int n, int k) {
   int res = 1;
   // Since C(n, k) = C(n, n-k)
   if (k > n - k)
        k = n - k;
   // Calculate value of [n*(n-1)*---*(n-k+1)] / [k*(k-1)*---*1]
   for (int i = 0; i < k; ++i) {
      res *= (n - i);
      res /= (i + 1);
   }
   return res;
}</pre>
```

7.5. Compressao De Pontos

```
map<int, int> rev;
for(int x : arr) {
    sl.insert(x);

vector<int> aux;
for(int x : sl) aux.pb(x);
for(int i=0; i<n; i++) {
    int id = lower_bound(aux.begin(), aux.end(), arr[i]) - aux.begin();
    rev[id] = arr[i];
    arr[i] = id;
}</pre>
```

7.6. Equacao Diofantina

```
int gcd(int a, int b, int &x, int &y) {
2
     if (a == 0) {
       x = 0; v = 1;
       return b;
5
    int x1, y1;
     int d = gcd(b%a, a, x1, y1);
     x = y1 - (b / a) * x1;
9
    y = x1;
10
    return d;
11
12
1.3
   bool find_any_solution(int a, int b, int c, int &x0, int &y0, int &g) {
14
    g = gcd(abs(a), abs(b), x0, y0);
15
     if (c % q)
16
       return false;
17
18
     x0 \star = c / q;
     y0 *= c / g;
19
     if (a < 0) x0 = -x0;
20
21
     if (b < 0) y0 = -y0;
22
     return true;
23 }
```

7.7. Euclides Estendido

```
1 int gcd, x, y;
  //Ax + By = gcd(A, B)
   void extendedEuclidian(int a, int b) {
     if(b==0){
8
       qcd=a;
9
       x=1;
10
       v=0;
11
      } else{
12
       extendedEuclidian(b, a%b);
13
14
       int temp = x;
15
16
       y = temp - (a/b) *y;
17
18 }
```

7.8. Euler Totient

```
int phi(int n) {
     int result = n;
     for (int i = 2; i * i <= n; i++) {
       if(n % i == 0) {
         while (n \% i == 0)
             n /= i;
           result -= result / i;
7
8
9
10
11
     if(n > 1)
12
       result -= result / n;
13
     return result;
14
```

7.9. Fatoracao Multiplas Queries

```
//stor smallest prime factor for every num
   int spf[MAXN];
   // Calculating SPF (Smallest Prime Factor) for every number till MAXN.
   // Time Complexity : O(nloglogn)
   void sieve() {
     spf[1] = 1;
     for (int i=2; i<MAXN; i++)</pre>
       // marking smallest prime factor for every number to be itself.
       spf[i] = i;
10
     // separatelyMarking spf for every even
11
12
     // number as 2
13
     for (int i=4; i<MAXN; i+=2)
14
       spf[i] = 2;
15
     for (int i=3; i*i<MAXN; i++) {</pre>
16
17
       // checking if i is prime
       if (spf[i] == i) {
18
         // marking SPF for all numbers divisible by i
19
         for (int j=i*i; j<MAXN; j+=i)
20
21
           // marking spf[j] if it is not previously marked
22
           if (spf[j]==j)
23
             spf[j] = i;
24
25
26
27
   // A O(log n) function returning primefactorization
   // by dividing by smallest prime factor at every step
   vector<int> getFactorization(int x) {
30
     vector<int> ret;
31
     while (x != 1)
32
       ret.push_back(spf[x]);
33
       x = x / spf[x];
34
35
     return ret;
```

7.10. Fatoracao Simples

```
map<int, int> primeFactors(int n) {
    set<int> ret;
    while (n%2 == 0) {
        m[2]++;
        n = n/2;
    }
}
```

```
8
     int sq = sqrt(n);
     for (int i = 3; i <= sq+2; i = i+2) {
       while (n%i == 0) {
10
11
         m[i]++;;
12
          n = n/i;
1.3
        /★ OBS1
14
1.5
     TF(N < 1E7)
16
       FATORE COM SPF
17
        */
18
19
20
     if (n > 2)
21
       m[n]++;
22
23
     return ret;
24 }
```

7.11. Inclusao-Exclusao

```
| // | A ∪ B ∪ C|=|A|+|B|+|C|-|A ∩ B|-|A ∩ C|-|B ∩ C|+|A ∩ B ∩ C|
   // EXEMPLO: Quantos números de 1 a 10^9 são múltiplos de 42, 54, 137 ou 201?
   int f(vector<int> arr, int LIMIT) {
7
     int n = arr.size();
     int c = 0;
10
      for(int mask = 1; mask < (1<<n); mask++) {</pre>
11
       int lcm = 1;
12
        for(int i = 0; i < n; i++)</pre>
13
          if (mask& (1<<i))
            lcm=lcm*arr[i]/__gcd(lcm,arr[i]);
14
15
        // se o numero de conjutos a unir for impar entao soma
16
       if(__builtin_popcount(mask) %2 == 1)
17
         c += LIMIT/lcm;
18
        else // se nao subtrai
19
          c -= LIMIT/1cm;
20
21
22
     return LIMIT-c;
23
24 }
```

7.12. Numero De Fatores

```
int calcFat(vector<int> fatores) {
   int x = fatores[1];
   auto lo = lower_bound(fatores.begin(), fatores.end(), x);
   auto up = upper_bound(fatores.begin(), fatores.end(), x);
   int fat = 1;
   while(up != fatores.end()) {
      fat *= (up-lo+l);
      lo = lower_bound(fatores.begin(), fatores.end(), *up);
      up = upper_bound(fatores.begin(), fatores.end(), *up);
   }
   fat *= (up-lo+l);
   return fat;
}
```

7.13. Pollard Rho (Find A Divisor)

```
/* Function to calculate (base^exponent) % modulus */
   int modular_pow(int base, int exponent,
3
               int modulus) {
     /* initialize result */
     int result = 1;
     while (exponent > 0) {
       /* if y is odd, multiply base with result */
       if (exponent & 1)
9
        result = (result * base) % modulus;
10
       /* exponent = exponent/2 */
       exponent = exponent >> 1;
11
       /* base = base * base */
12
13
       base = (base * base) % modulus;
14
15
    return result:
16
17
   /* method to return prime divisor for n */
   int PollardRho(int n) {
19
     /* initialize random seed */
     srand (time(NULL));
22
23
     /* no prime divisor for 1 */
24
     if (n==1) return n;
25
26
     /★ even number means one of the divisors is 2 ★/
27
     if (n % 2 == 0) return 2;
28
     /* we will pick from the range [2, N) */
29
30
     int x = (rand()%(n-2))+2;
31
     int v = x;
32
33
     /* the constant in f(x).
     * Algorithm can be re-run with a different c
34
35
     * if it throws failure for a composite. */
36
     int c = (rand() % (n-1)) +1;
37
38
     /* Initialize candidate divisor (or result) */
39
     int d = 1:
40
41
     /* until the prime factor isn't obtained.
42
     If n is prime, return n */
     while (d==1) {
43
44
       /* Tortoise Move: x(i+1) = f(x(i)) */
45
       x = (modular_pow(x, 2, n) + c + n) %n;
46
       /* Hare Move: y(i+1) = f(f(y(i))) */
47
48
       y = (modular_pow(y, 2, n) + c + n)%n;
49
       y = (modular_pow(y, 2, n) + c + n)%n;
50
51
       /* check gcd of |x-y| and n */
52
       d = \underline{gcd(abs(x-y), n)};
53
54
       /★ retry if the algorithm fails to find prime factor
       * with chosen x and c */
56
       if (d==n) return PollardRho(n);
57
58
59
     return d;
60 }
62 /* driver function */
63 | signed main() {
```

```
int n = 12;
printf("One of the divisors for %lld is %lld.",
n, PollardRho(n));
return 0;
}
```

7.14. Precomputar Combinacao Ncr

```
1 int C[1123][1123];
3 int mod(int n) {return n%((int)1e9+7);}
5 int nCr(int n, int k) {
    for(int i = 0; i <= n; i++) {</pre>
       for(int j = 0; j <= min(i,k); j++) {</pre>
8
         if(j == 0 || j == i) {
           C[i][j] = 1;
10
         } else {
11
           C[i][j] = mod(C[i-1][j-1]+C[i-1][j]);
12
13
14
15 }
```

7.15. Teorema Chines Do Resto

```
1 int inv(int a, int m) {
     int m0 = m, t, q;
    int x0 = 0, x1 = 1;
    if (m == 1)
       return 0;
8 // Apply extended Euclid Algorithm
   while (a > 1) {
10 // g is quotient
11
       if(m == 0)
12
        return INF;
13
       q = a / m;
14
      t = m;
15
      // m is remainder now, process same as euclid's algo
16
      m = a % m, a = t;
17
      t = x0;
18
       x0 = x1 - q \star x0;
19
      x1 = t;
20
21
22
     // Make x1 positive
23
    if (x1 < 0)
24
      x1 += m0;
25
    // debug(x1);
2.6
27
    return x1;
28 }
29 // k is size of num[] and rem[]. Returns the smallest
30 // number x such that:
31 // x % num[0] = rem[0],
32 // x \% num[1] = rem[1],
33 // .......
34 // x \% num[k-2] = rem[k-1]
35 // Assumption: Numbers in num[] are pairwise coprimes
36 // (gcd for every pair is 1)
37 | int findMinX(int num[], int rem[], int k){
```

```
// Compute product of all numbers
39
     int prod = 1;
     for (int i = 0; i < k; i++)
41
       prod *= num[i];
42
43
     // Initialize result
44
     int result = 0:
45
46
     // Apply above formula
     for (int i = 0; i < k; i++) {
47
       int pp = prod / num[i];
48
49
       // debug(pp);
50
       int iv = inv(pp, num[i]);
51
       if(iv == INF)
52
         return INF;
53
       result+=rem[i]*inv(pp,num[i])*pp;
54
55
56
     // IF IS NOT VALID RETURN INF
57
     return (result % prod == 0 ? INF: result % prod);
```

7.16. Binary Exponentiation

```
int power(const int x, const int p, const int MOD = ((int)1e9 + 7)) {
    if(p == 0)
    return 1%MOD;
    if(p == 1)
        return x%MOD;
    int res = power(x, p/2, MOD);
    res = (long long)res*res%MOD;
    if(p&1)
        return res;
}
```

7.17. Divisors

```
1 I/I OBS: EACH NUMBER HAS AT MOST \sqrt[3]{N} DIVISORS
   vector<int> divisors(int n) {
3
     vector<int> ans;
     for (int i=1; i * i <= n; i++) {
       if (n%i==0) {
          // If divisors are equal, print only one
         if (n/i == i)
           ans.pb(i);
          else // Otherwise print both
10
            ans.pb(i), ans.pb(n/i);
11
12
13
     return ans;
14
```

7.18. Matrix Exponentiation

```
namespace matrix {
#define Matrix vector<vector<int>>
const int MOD = 1e9 + 7;

/// Creates an n x n identity matrix.
/// Time Complexity: O(n*n)
```

```
8 | Matrix identity(const int n) {
9
     assert (n > 0);
10
11
     Matrix mat_identity(n, vector<int>(n, 0));
12
13
     for (int i = 0; i < n; i++)</pre>
14
       mat_identity[i][i] = 1;
1.5
16
     return mat_identity;
17 }
18
19 /// Multiplies matrices a and b.
20
   /// Time Complexity: O(mat.size() ^ 3)
21
   Matrix mult (const Matrix &a, const Matrix &b) {
     assert(a.front().size() == b.size());
23
24
25
     Matrix ans(a.size(), vector<int>(b.front().size(), 0));
26
     for (int i = 0; i < ans.size(); i++)
27
        for (int j = 0; j < ans.front().size(); j++)</pre>
28
          for (int k = 0; k < a.front().size(); k++)</pre>
29
            ans[i][j] = (ans[i][j] + a[i][k] * b[k][j]) % MOD;
30
31
     return ans:
32
33
34 /// Exponentiates the matrix mat to the power of p.
35 ///
36 /// Time Complexity: O((\text{mat.size}() ^ 3) * \log_2(p))
37 Matrix expo(Matrix &mat, int p) {
38
     assert(p >= 0);
39
40
     Matrix ans = identity(mat.size());
41
     Matrix cur power;
42
     cur_power.swap(mat);
43
44
     while (p)
45
       if (p & 1)
46
         ans = mult(ans, cur_power);
47
48
       cur_power = mult(cur_power, cur_power);
49
       p >>= 1;
50
51
52
     return ans;
54 }; // namespace matrix
```

7.19. Modular Inverse

```
1 int gcd(int a, int b) {
    if (a == 0)
       return b;
     return gcd(b%a, a);
5
7 int power(int x, int p, int MOD) {
    if(p == 0)
       return 1%MOD;
10
    if(p == 1)
11
       return x%MOD;
12
     int res = power(x, p/2, MOD);
13
     res = (long long) res*res%MOD;
14
     if(p&1)
```

```
res = (long long) res*x%MOD;
16
     return res;
17
18
19 // A*B = x (mod m)
20 // B = x * A^(-1)
21 // Function to find modular inverse of a under modulo m
22 // Assumption: m is prime
  int modInverse(int a, int mod) {
2.4
     int g = gcd(a, mod);
25
     if (g != 1)
26
         return -1;
27
     else
28
       // If a and m are relatively prime, then modulo inverse
29
       // is a^(m-2) mod m
30
       return power(a, mod-2, mod);
31
```

7.20. Sieve + Segmented Sieve

```
const int MAXN = 1e6;
   /// Contains all the primes in the segments
   vector<int> seqPrimes;
   bitset < MAXN+5> primesInSeq;
   /// smallest prime factor
   int spf[MAXN+5];
   vector<int> primes;
10
   bitset<MAXN+5> isPrime;
11
12
   void sieve(int n = MAXN + 2) {
13
14
15
     for(int i = 0; i <= n; i++)
16
       spf[i] = i;
17
18
     isPrime.set();
19
     for(int i = 2; i <= n; i++) {</pre>
20
       if(!isPrime[i])
21
         continue;
22
23
       for (int j = i*i; j <= n; j+=i) {
24
         isPrime[j] = false;
25
         spf[j] = min(i, spf[j]);
26
27
       primes.pb(i);
28
29
30
31
   vector<int> getFactorization(int x) {
32
     vector<int> ret;
33
     while (x != 1)
34
       ret.push_back(spf[x]);
35
       x = x / spf[x];
36
37
     return ret;
38
39
40 /// Gets all primes from 1 to r
41 void segSieve(int l, int r) {
    // primes from 1 to r
43
    // transferred to 0..(l-r)
     segPrimes.clear();
```

```
primesInSeq.set();
     int sq = sqrt(r) + 5;
47
48
     for(int p: primes)
49
       if(p > sq)
50
         break:
51
52
        for(int i = 1 - 1%p; i <= r; i += p) {
53
         if(i - 1 < 0)
54
           continue:
55
         // if i is less than 1e6, it could be checked in the
56
57
         // array of the sieve
         if(i >= (int)1e6 || !isPrime[i])
58
59
           primesInSeq[i-l] = false;
60
61
62
63
     for (int i = 0; i < r-l+1; i++) {
64
       if(primesInSeg[i])
65
         segPrimes.pb(i+l);
66
67 }
```

8. Miscellaneous

8.1. 2-Sat

```
1 class SCC {
2 private:
    // number of vertices
    int n;
    // indicates whether it is indexed from 0 or 1
    int indexed from:
     // reversed graph
     vector<vector<int>> trans;
10 private:
11
     void dfs_trans(int u, int id) {
12
       comp[u] = id;
13
       scc[id].push_back(u);
14
15
        for (int v : trans[u])
16
         if (comp[v] == -1)
17
           dfs_trans(v, id);
18
19
20
     void get_transpose(vector<vector<int>> &adj) {
21
       for (int u = indexed_from; u < this->n + indexed_from; u++)
22
         for (int v : adj[u])
23
           trans[v].push_back(u);
24
25
     void dfs_fill_order(int u, stack<int> &s, vector<vector<int>> &adj) {
26
2.7
       comp[u] = true;
28
        for (int v : adj[u])
29
30
         if (!comp[v])
31
           dfs_fill_order(v, s, adj);
32
33
       s.push(u);
34
35
36
     // The main function that finds all SCCs
```

```
void compute_SCC(vector<vector<int>> &adj) {
 38
 39
         stack<int> s:
 40
        // Fill vertices in stack according to their finishing times
 41
        for (int i = indexed from; i < this->n + indexed from; i++)
 42
          if (!comp[i])
 43
            dfs_fill_order(i, s, adj);
 44
        // Create a reversed graph
 45
 46
        get_transpose(adj);
 47
 48
        fill(comp.begin(), comp.end(), -1);
 49
 50
        // Now process all vertices in order defined by stack
 51
        int id = 0;
 52
        while (s.empty() == false) {
 53
          int v = s.top();
 54
           s.pop();
 55
 56
          if (comp[v] == -1)
 57
            dfs_trans(v, id++);
 58
 59
      }
 60
    public:
     // number of the component of the i-th vertex
      // it's always indexed from 0
      vector<int> comp;
 64
 65
      // the i-th vector contains the vertices that belong to the i-th scc
 66
      // it's always indexed from 0
 67
      vector<vector<int>> scc;
 68
 69
      SCC(int n, int indexed from, vector<vector<int>> &adj) {
 70
        this->n = n;
 71
        this->indexed_from = indexed_from;
 72
        comp.resize(n + 1);
 73
        trans.resize(n + 1):
 74
        scc.resize(n + 1);
 75
 76
        this->compute_SCC(adj);
 77
 78
    };
 79
    // OBS: INDEXED FROM 0
    class SAT {
    private:
      vector<vector<int>> adj;
 85
      int n;
 86
    public:
 87
      vector<bool> ans:
 88
 89
 90
      SAT(int n) {
 91
        this->n = n;
 92
        adj.resize(2 \star n);
 93
        ans.resize(n);
 94
 95
 96
      // (X V Y) = (X -> ~Y) & (~X -> Y)
 97
      void add_or(int x, bool pos_x, int y, bool pos_y) {
 98
        assert(0 \le x);
 99
        assert(x < n);
100
        assert(0 \le v);
101
        assert (x < n);
```

```
ad_{1}[(x << 1) ^pos_{x}].pb((y << 1) ^(pos_{y} ^1));
103
        adi[(v << 1) ^pos v].pb((x << 1) ^(pos x ^ 1));
104
105
106
       // (X \times Y) = (X \vee Y) & (\sim X \vee \sim Y)
107
       // for this function the result is always 0 1 or 1 0
108
       void add_xor(int x, bool pos_x, int y, bool pos_y) {
109
         assert(0 \le x);
110
         assert (x < n):
         assert (0 \le v);
111
112
         assert (x < n);
         add_or(x, y, pos_x, pos_y);
113
114
         add_or(x, y, pos_x ^ 1, pos_y ^ 1);
115
116
117
      bool check() {
118
        SCC scc(2 * n, 0, adj);
119
120
         for (int i = 0; i < n; i++) {</pre>
121
           if (scc.comp[(i << 1) | 1] == scc.comp[(i << 1) | 0])</pre>
122
             return false;
123
           ans[i] = (scc.comp[(i << 1) | 1] < scc.comp[(i << 1) | 0]);
124
125
126
         return true;
127
128 };
```

8.2. 3Sum Problem

```
1 // vetor arr e valor x, a soma de três valores desse vetor deve ser iqual a x
3 bool sum3(int arr[], int x, int n) {
    sort(arr,arr + n);
4
     for (int i = 0; i < n-2; i++) {
       int 1 = i+1, r = n-1;
     /★ 2SUM problem -> ponteiro que aponta para o primeiro e ultimo da
        sequencia e caso a soma for menor do que x adianta em uma casa o
        ponteiro da esquerda caso seja maior diminui em uma casa o ponteiro da
        direita */
8
       while(1 < r)  {
         if(arr[i] + arr[l] + arr[r] == x) {
9
10
           return true;
         } else if(arr[i] + arr[l] + arr[r] < x)</pre>
11
12
           1++:
13
         else
14
           r--:
15
16
17
     return false;
18 }
```

8.3. Fibonacci Matrix Exponentiation

```
int fib (int n) {
  long long fib[2][2]= {{1,1},{1,0}};
  int ret[2][2]= {{1,0},{0,1}};
  int tmp[2][2]= {{0,0},{0,0}};
  int i,j,k;
  while(n) {
  if(n&1) {
    memset(tmp,0,sizeof tmp);
  for(i=0; i<2; i++)</pre>
```

```
for(j=0; j<2; j++)
11
              for(k=0; k<2; k++)
12
                tmp[i][j] = (tmp[i][j] + ret[i][k] * fib[k][j]);
13
          for(i=0; i<2; i++)
            for(j=0; j<2; j++)
14
15
              ret[i][j]=tmp[i][j];
16
17
        memset(tmp, 0, sizeof tmp);
        for(i=0; i<2; i++)
18
19
          for(j=0; j<2; j++)
20
            for (k=0; k<2; k++)
              tmp[i][j] = (tmp[i][j] + fib[i][k] * fib[k][j]);
21
22
        for(i=0; i<2; i++)
23
          for(j=0; j<2; j++)
            fib[i][j]=tmp[i][j];
24
25
26
27
     return (ret[0][1]);
28
```

8.4. Infix To Prefix

```
int main() {
   map<char.int> prec:
   stack<char> op;
   string postfix;
   string infix;
   cin >> infix;
   prec['+'] = prec['-'] = 1;
   prec['*'] = prec['/'] = 2;
11
   prec['^1] = 3;
  for(int i = 0; i < infix.length(); i++) {</pre>
12
13
     char x = infix[i];
14
     if('0' <= x && x <= '9') {
15
       for(i;i < infix.length() && ('0' <= infix[i] && infix[i] <= '9');i++)</pre>
16
         postfix += infix[i];
17
18
     \} else if(('a' <= x && x <= 'z') || ('A' <= x && x <= 'Z')) {
19
       postfix += x;
     } else if (x == '(')
20
       op.push('(');
21
     else if(x == ')') {
22
23
       while(!op.empty() && op.top() != '(') {
24
         postfix += op.top();
25
         op.pop();
26
27
       op.pop();
28
       else {
29
       while(!op.empty() && prec[op.top()] >= prec[x]) {
30
       postfix += op.top();
31
       op.pop();
32
33
     op.push(x);
34
35
   while(!op.empty()) {
  postfix += op.top();
38 op.pop();
39
   cout << postfix << endl;
41
```

8.5. Interval Scheduling

```
1 | 1 -> Ordena pelo final do evento, depois pelo inicio.
2 | 2 -> Vai iterando pelos eventos, se eles não tiverem horário em comum então adiciona o evento à lista.
```

8.6. Kadane (Maior Soma Num Vetor)

```
int kadane(int arr[], int 1) {
    int soma, total;
    soma = total = arr[0];

for(int i = 1; i < 1; i++) {
    soma = max(arr[i], arr[i] +soma);
    if(soma > total)
        total = soma;
}

return total;
}
```

8.7. Kadane 2D

```
2 /// Program to find maximum sum subarray in a given 2D array
3 #include <stdio.h>
4 #include <string.h>
5 #include <limits.h>
6 int mat[1001][1001]
7 int ROW = 1000, COL = 1000;
   // Implementation of Kadane's algorithm for 1D array. The function
10 // returns the maximum sum and stores starting and ending indexes of the
11 // maximum sum subarray at addresses pointed by start and finish pointers
12 // respectively.
13 | int kadane(int* arr, int* start, int* finish, int n) {
14
        // initialize sum, maxSum and
15
        int sum = 0, maxSum = INT_MIN, i;
16
17
        // Just some initial value to check for all negative values case
18
        \starfinish = -1:
19
20
        // local variable
21
       int local start = 0;
22
23
        for (i = 0; i < n; ++i) {
           sum += arr[i];
24
           if (sum < 0) {
25
26
                sum = 0:
2.7
                local_start = i+1;
2.8
            else if (sum > maxSum) {
29
30
                maxSum = sum:
31
                *start = local_start;
32
                \starfinish = i;
33
        }
34
35
36
         // There is at-least one non-negative number
37
       if (\starfinish != -1)
38
           return maxSum:
39
```

```
// Special Case: When all numbers in arr[] are negative
41
       maxSum = arr[0];
42
       *start = *finish = 0;
43
44
       // Find the maximum element in array
45
       for (i = 1; i < n; i++) {
46
           if (arr[i] > maxSum) {
               maxSum = arr[i];
47
               *start = *finish = i;
48
49
50
51
       return maxSum;
52
53
54
   // The main function that finds maximum sum rectangle in mat[][]
55
   int findMaxSum() {
56
       // Variables to store the final output
57
       int maxSum = INT_MIN, finalLeft, finalRight, finalTop, finalBottom;
58
59
       int left, right, i;
60
       int temp[ROW], sum, start, finish;
61
62
       // Set the left column
63
       for (left = 0; left < COL; ++left) {
64
           // Initialize all elements of temp as 0
65
           for (int i = 0; i < ROW; i++)
66
               temp[i] = 0;
67
68
           // Set the right column for the left column set by outer loop
69
           for (right = left; right < COL; ++right) {</pre>
70
               // Calculate sum between current left and right for every row 'i'
71
               for (i = 0; i < ROW; ++i)
72
                    temp[i] += mat[i][right];
73
74
               // Find the maximum sum subarray in temp[]. The kadane()
75
               // function also sets values of start and finish. So 'sum' is
76
               // sum of rectangle between (start, left) and (finish, right)
77
               // which is the maximum sum with boundary columns strictly as
78
               // left and right.
79
               sum = kadane(temp, &start, &finish, ROW);
80
81
               // Compare sum with maximum sum so far. If sum is more, then
               // update maxSum and other output values
82
               if (sum > maxSum) {
84
                   maxSum = sum;
                    finalLeft = left;
                    finalRight = right;
86
87
                    finalTop = start;
88
                    finalBottom = finish;
89
90
91
92
93
       return maxSum:
94
       // Print final values
95
       printf("(Top, Left) (%d, %d)\n", finalTop, finalLeft);
96
       printf("(Bottom, Right) (%d, %d)\n", finalBottom, finalRight);
97
       printf("Max sum is: %d\n", maxSum);
98
```

8.8. Oito Rainhas

```
#define N 4
bool isSafe(int mat[N][N],int row,int col) {
```

```
for(int i = row - 1; i >= 0; i--)
        if (mat[i][col])
5
          return false:
      for (int i = row - 1, j = col - 1; i >= 0 && <math>j >= 0; i--, j--)
6
7
        if(mat[i][i])
8
          return false:
9
      for (int i = row - 1, j = col + 1; i >= 0 && j < N; i --, j++)
1.0
        if (mat[i][j])
11
          return false:
12
      return true:
13
   // inicialmente a matriz esta zerada
14
15 | int queen (int mat[N][N], int row = 0) {
     if(row >= N) {
16
        for(int i = 0; i < N; i++) {
17
          for(int j = 0; j < N; j++) {
   cout << mat[i][j] << ' ';</pre>
18
19
20
21
          cout << endl;
22
23
        cout << endl << endl:
24
        return false;
25
      for(int i = 0; i < N; i++) {
26
27
        if(isSafe(mat,row,i)) {
28
          mat[row][i] = 1;
29
          if (queen (mat, row+1))
3.0
            return true;
31
          mat[row][i] = 0;
32
33
34
     return false:
35
```

8.9. Sliding Window Minimum

```
// mínimo num vetor arr de arr[0] ... arr[k-1], arr[1] ... arr[k], arr[2]
        ... arr[k+1]
   void swma(vector<int> arr, int k) {
     deque<ii>> window:
     for(int i = 0; i < arr.size(); i++) {</pre>
        while(!window.empty() && window.back().ff > arr[i])
          window.pop_back();
8
        window.pb(ii(arr[i],i));
        while(window.front().ss <= i - k)</pre>
9
10
          window.pop_front();
11
     if(i >= k)
12
       cout << ' ';
13
     if(i - k + 1 >= 0)
15
        cout << window.front().ff;</pre>
16
17 }
```

8.10. Torre De Hanoi

```
#include <stdio.h>

// C recursive function to solve tower of hanoi puzzle

void towerOfHanoi(int n, char from_rod, char to_rod, char aux_rod) {
   if (n == 1) {
      printf("\n Move disk 1 from rod %c to rod %c", from_rod, to_rod);
}
```

```
7
       return;
8
9
     towerOfHanoi(n-1, from rod, aux rod, to rod);
10
     printf("\n Move disk %d from rod %c to rod %c", n, from_rod, to_rod);
11
     towerOfHanoi(n-1, aux rod, to rod, from rod);
12
13
14 | int main() {
    int n = 4; // Number of disks
15
16
     towerOfHanoi(n, 'A', 'C', 'B'); // A, B and C are names of rods
17
18
```

8.11. Kadane (Segment Tree)

```
struct Node {
     int pref, suf, tot, best;
3
     Node () {}
     Node (int pref, int suf, int tot, int best) : pref(pref), suf(suf),
       tot(tot), best(best) {}
   const int MAXN = 2E5 + 10;
8
   Node tree[5*MAXN];
   int arr[MAXN];
10
   Node query (const int 1, const int r, const int i, const int j, const int
11
       pos) {
12
     if(1 > r || 1 > i || r < i)
13
14
       return Node(-INF, -INF, -INF, -INF);
15
     if(i <= 1 && r <= i)
16
       return Node (tree [pos].pref, tree [pos].suf, tree [pos].tot,
17
       tree[posl.best);
18
19
     int mid = (l + r) / 2;
20
     Node left = query(1, mid, i, j, 2*pos+1), right = query(mid+1, r, i, j, 2*pos+2);
21
     Node x:
     x.pref = max({left.pref, left.tot, left.tot + right.pref});
     x.suf = max({right.suf, right.tot, right.tot + left.suf});
24
     x.tot = left.tot + right.tot;
25
     x.best = max({left.best, right.best, left.suf + right.pref});
     return x;
26
27
28
29
   // Update arr[idx] to v
   // ITS NOT DELTA!!!
30
31
   void update(int 1, int r, const int idx, const int v, const int pos) {
32
     if(l > r || l > idx || r < idx)
33
       return;
34
35
     if(1 == idx && r == idx) {
36
       tree[pos] = Node(v, v, v, v);
37
       return:
38
39
     int mid = (1 + r)/2;
     update (1, mid, idx, v, 2*pos+1); update (mid+1, r, idx, v, 2*pos+2);
42
     1 = 2 * pos + 1, r = 2 * pos + 2;
     tree[pos].pref = max({tree[1].pref, tree[1].tot, tree[1].tot +
       tree[r].pref});
     tree[pos].suf = max({tree[r].suf, tree[r].tot, tree[r].tot + tree[l].suf});
     tree[pos].tot = tree[l].tot + tree[r].tot;
```

```
tree[pos].best = max({tree[1].best,tree[r].best, tree[1].suf +
       tree[r].pref});
47 }
48
   void build(int 1, int r, const int pos) {
49
5.0
51
     if(1 == r) {
52
       tree[pos] = Node(arr[l], arr[l], arr[l]);
53
       return:
54
55
56
     int mid = (1 + r)/2;
57
     build(1, mid, 2*pos+1); build(mid+1, r, 2*pos+2);
     1 = 2*pos+1, r = 2*pos+2;
58
     tree[pos].pref = max({tree[l].pref, tree[l].tot, tree[l].tot +
       tree[r].pref});
60
     tree[pos].suf = max({tree[r].suf, tree[r].tot, tree[r].tot + tree[l].suf});
61
     tree[pos].tot = tree[l].tot + tree[r].tot;
62
     tree[pos].best = max({tree[l].best, tree[r].best, tree[l].suf +
       tree[r].pref});
```

8.12. Point Compression

```
// map<int, int> rev;
3 /// Compress points in the array arr to the range [0..n-1].
4 ///
   /// Time Complexity: O(n log n)
   vector<int> compress(vector<int> &arr) {
     vector<int> aux = arr;
     sort(aux.begin(), aux.end());
     aux.erase(unique(aux.begin(), aux.end()), aux.end());
10
     for (size_t i = 0; i < arr.size(); i++) {</pre>
11
12
       int id = lower_bound(aux.begin(), aux.end(), arr[i]) - aux.begin();
13
       // rev[id] = arr[i];
14
       arr[i] = id;
15
16
     return arr;
17 }
```

9. Strings

9.1. Kmp

```
vector<int> pi(const string &s) {
     int n = s.size();
.3
     vector<int> pi(n);
     int 1 = 0, r = 1;
5
     while (r < n) {
       if (s[1] == s[r]) {
8
         1++:
         pi[r] = 1;
10
         r++;
11
        } else {
12
         if (1 == 0)
13
           pi[r] = 0;
14
           r++;
15
         } else
16
           1 = pi[1 - 1];
17
```

```
19
     return pi;
20
21
   // returns the index of first occurence of a pat in a txt
   int kmp(const string &txt, const string &pat) {
24
     int n = txt.size(), m = pat.size();
2.5
26
     int t = 0, p = 0;
     vector<int> pi_pat = pi(pat);
2.7
28
     // vector<int> occ;
29
30
     while (t < n) {
31
32
       if (txt[t] == pat[p]) {
33
         t++, p++;
34
         if (p == m) {
35
           return t - m;
36
            // if you want to continue searching
37
            // occ.pb(t - m);
38
           // p = pi_pat[p - 1];
39
40
       } else {
41
          if (p == 0) {
42
           t++;
43
          } else {
44
           p = pi_pat[p - 1];
45
46
47
48
     return -1:
49
     // return occ;
50
```

9.2. Suffix Array

```
1 // Created by Ubiratan Neto
2
3
4
                                       Suf,
                                              lcp
5
   0 banana
                                     5 a
                                               0
   1 anana
               Sort the Suffixes
                                     3 ana
               ---->
   2 nana
                                     1 anana
8
   3 ana
                alphabetically
                                     0 banana 0
9
   4 na
                                     4 na
                                               0
   5 a
10
11
   lcp = number of characters equal prefi
12
13
   */
14
15
   struct SuffixArray {
16
17
     vector<int> rnk,tmp,sa, sa_aux, lcp, pot, sp[22];
18
19
     int block, n;
20
21
     string s;
22
23
     SuffixArray() {}
24
25
     SuffixArray(string t){
26
       s = t;
27
       n = t.size();
28
       rnk.resize(n+1);
```

```
for(int i=0; i<22; i++) sp[i].resize(n+1);</pre>
29
30
        pot.resize(n+1);
31
        tmp.resize(max((int)257, n+1));
32
        // sa stores index of first char of sufix
33
        sa.resize(n+1);
34
        sa aux.resize(n+1);
35
        // lcp stores value between the string and next string
36
       lcp.resize(n+1);
37
       block = 0;
38
39
40
      bool suffixcmp(int i, int j){
        if(rnk[i] != rnk[j]) return rnk[i] < rnk[j];</pre>
41
42
        i+=block, j+=block;
43
        i%=n;
44
        j%=n;
45
        return rnk[i] < rnk[j];</pre>
46
47
48
      void suffixSort(int MAX VAL){
49
        for(int i=0; i<=MAX_VAL; i++) tmp[i] = 0;</pre>
        for(int i=0; i<n; i++) tmp[rnk[i]]++;</pre>
50
        for (int i=1; i<=MAX_VAL; i++) tmp[i] += tmp[i-1];</pre>
51
        for(int i = n-1; i>=0; i--){
52
53
            int aux = sa[i]-block;
            aux%=n;
54
55
            if (aux < 0) aux+=n;
            sa_aux[--tmp[rnk[aux]]] = aux;
56
57
58
        for(int i=0; i<n; i++) sa[i] = sa_aux[i];</pre>
59
        t.mp[0] = 0:
60
        for (int i=1; i < n; i++) tmp[i] = tmp[i-1] + suffix cmp(sa[i-1], sa[i]);
61
        for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
62
63
64
      void calculate() {
        s+='\0';
65
       n++;
66
67
        for(int i=0; i<n; i++) {</pre>
          sa[i] = i;
68
          rnk[i] = s[i];
69
70
          tmp[i] = 0;
71
72
        suffixSort (256);
73
        block = 1;
74
        while (tmp[n-1] != n-1) {
75
          suffixSort(tmp[n-1]);
76
          block*=2;
77
78
        for(int i=0; i<n-1; i++) sa[i] = sa[i+1];</pre>
79
8.0
        tmp[0] = 0;
81
        for (int i=1; i < n; i++) tmp[i] = tmp[i-1] + suffix cmp(sa[i-1], sa[i]);
        for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
82
83
        s.pop_back();
84
        sa.pop_back();
85
86
87
      void calculate_lcp() {
88
        int last = 0;
89
        for(int i=0; i<n; i++) {</pre>
90
          if(rnk[i] == n-1) continue;
91
          int x = rnk[i];
92
          lcp[x] = max((int)0, last-1);
```

```
while (sa[x] + lcp[x] < n \&  sa[x+1] + lcp[x] < n \&  s[sa[x]+lcp[x]] ==
         s[sa[x+1]+lcp[x]]){
 94
             lcp[x]++;
 95
 96
           last = lcp[x];
 97
 98
 99
      void build_lcp_table() {
100
101
         int k = 0;
102
         for (int j = 0; (1<<j) <= 2*n; j++) {
103
           for (; k \le n \&\& k < (1 << \dot{\uparrow}); k++) {
104
             pot[k] = j-1;
105
106
107
         for (int i=0; i<n; i++) {</pre>
108
           sp[0][i] = lcp[i];
109
110
         for (int i = 1; (1<<ii) <= n; i++) {
111
           for (int j = 0; j+(1 << i) <= n; <math>j++) {
             sp[i][j] = min(sp[i-1][j], sp[i-1][j+(1<<(i-1))]);
112
113
114
115
116
117
      // to find lcp of two different sufixes starting at x and y
118
      int query_lcp(int x, int y) {
         if (x = y) return (x = y)
119
120
         if(rnk[x] > rnk[y]) swap(x,y);
         int l = rnk[x], r = rnk[y]-1;
121
         return min(sp[pot[r-l+1]][1], sp[pot[r-l+1]][r-(1LL<<pot[r-l+1])+1]);</pre>
122
123
124
125
      // needs calculate and calculate lcp first
      int number_of_substrings() {
126
127
         int ans = n - sa[0];
         for(int i=0; i<n-1; i++){
128
129
           int length = n - sa[i+1];
           ans += length - lcp[i];
130
131
132
         return ans:
133
134
135
      // needs calculate and calculate lcp first
136
      int lcs(string &x, string &y) {
137
138
         string s = x + "#" + y;
139
140
         int n = (int) s.size() - 1;
141
         int ans = 0;
142
         for (int i = 0; i < n - 1; i++) {
143
          int ida = sa[i];
144
           int idb = sa[i+1];
145
           if(ida < x.size() && idb > x.size() || ida > x.size() && idb <</pre>
         x.size()) {
146
             ans = max(ans, lcp[i]);
147
148
149
150
151
152
```

```
17
20
26
27
28
29
30
31
32
33
34
```

```
// XOR(L,R) = XOR(1,L-1) ^ XOR(1,R)
ans= pre = 0
Trie.insert(0)
for i=1 to N:
    pre = pre XOR a[i]
    Trie.insert(pre)
    ans=max(ans, Trie.query(pre))
print ans
// a funcao query é a mesma da maximum xor between two elements
```

9.4. Trie - Maximum Xor Two Elements

```
1. Dada uma trie de números binários e um numero X, tente achar o número máximo que resultante da operação XOR

2

3 Ex: Para o número 10 (=(1010)2), o número que resulta no xor máximo é (0101)2, tente acha-lo na trie.
```

9.5. Z-Function

```
1 // What is Z Array?
2 / / For a string str[0..n-1], Z array is of same length as string.
3 // An element Z[i] of Z array stores length of the longest substring
4 // starting from str[i] which is also a prefix of str[0..n-1]. The
5 // first entry of Z array is meaning less as complete string is always
6 // prefix of itself.
7 // Example:
8 // Index
9 // 0 1
            2 3 4 5 6 7 8 9 10 11
10 // Text
11 //aabcaabxaaaz
12 // Z values
13 // X 1 0 0 3 1 0 0 2 2 1 0
14 // More Examples:
15 // str = "aaaaaa"
16 // Z[] = {x, 5, 4, 3, 2, 1}
18 // str = "aabaacd"
19 // Z[] = {x, 1, 0, 2, 1, 0, 0}
21 // str = "abababab"
22 // Z[] = {x, 0, 6, 0, 4, 0, 2, 0}
   vector<int> z function(const string &s) {
    vector<int> z(s.size());
     int 1 = -1, r = -1;
     for (int i = 1; i < s.size(); ++i) {</pre>
      z[i] = i >= r ? 0 : min(r - i, z[i - 1]);
      while (i + z[i] < s.size() && s[i + z[i]] == s[z[i]])
        z[i]++;
      if (i + z[i] > r)
        1 = i, r = i + z[i];
     return z;
35 }
```

9.6. Aho Corasick

```
1 /// REQUIRES trie.cpp
```

```
3 | class Aho {
   private:
     // node of the output list
     struct Out_Node {
7
       vector<int> str_idx;
8
       Out_Node *next = nullptr;
9
1.0
     vector<Trie::Node *> fail:
11
12
     Trie trie;
     // list of nodes of output
13
     vector<Out_Node *> out_node;
14
15
     const vector<string> arr;
16
17
     /// Time Complexity: O(number of characters in arr)
18
     void build trie() {
19
       const int n = arr.size();
20
       int node_cnt = 1;
21
22
       for (int i = 0; i < n; ++i)
23
         node_cnt += arr[i].size();
24
25
       out node.reserve(node cnt);
26
       for (int i = 0; i < node_cnt; ++i)</pre>
27
         out_node.push_back(new Out_Node());
28
29
       fail.resize(node_cnt);
30
       for (int i = 0; \overline{i} < n; ++i) {
31
         const int id = trie.insert(arr[i]);
32
         out_node[id]->str_idx.push_back(i);
33
34
35
       this->build failures();
36
37
38
     /// Returns the fail node of cur.
39
     Trie::Node *find_fail_node(Trie::Node *cur, char c) {
40
       while (cur != this->trie.root() && !cur->next.count(c))
41
         cur = fail[cur->id];
42
       // if cur is pointing to the root node and c is not a child
43
       if (!cur->next.count(c))
44
         return trie.root():
45
       return cur->next[c];
46
47
48
     /// Time Complexity: O(number of characters in arr)
49
     void build failures() {
50
       queue < const Trie:: Node *> q;
51
52
       fail[trie.root()->id] = trie.root();
53
       for (const pair<char, Trie::Node *> v : trie.root()->next) {
54
         q.emplace(v.second);
         fail[v.second->id] = trie.root();
55
56
         out_node[v.second->id]->next = out_node[trie.root()->id];
57
58
59
       while (!a.emptv()) {
60
         const Trie::Node *u = q.front();
61
62
63
         for (const pair<char, Trie::Node *> x : u->next) {
64
           const char c = x.first;
65
           const Trie::Node *v = x.second;
66
           Trie::Node *fail_node = find_fail_node(fail[u->id], c);
67
           fail[v->id] = fail_node;
```

```
69
            if (!out node[fail node->id]->str idx.empty())
70
              out_node[v->id]->next = out_node[fail_node->id];
71
            else
72
              out node[v->id]->next = out node[fail node->id]->next;
73
74
            q.emplace(v);
7.5
76
77
78
      vector<vector<pair<int, int>>> aho_find_occurrences(const string &text) {
79
80
        vector<vector<pair<int, int>>> ans(arr.size());
        Trie::Node *cur = trie.root();
81
82
83
        for (int i = 0; i < text.size(); ++i) {</pre>
84
          cur = find_fail_node(cur, text[i]);
85
          for (Out_Node *node = out_node[cur->id]; node != nullptr;
86
               node = node->next)
87
            for (const int idx : node->str idx)
88
              ans[idx].emplace_back(i - (int)arr[idx].size() + 1, i);
89
90
        return ans:
91
92
93
    public:
94
      /// Constructor that builds the trie and the failures.
9.5
      /// Time Complexity: O(number of characters in arr)
96
97
      Aho(const vector<string> &arr) : arr(arr) { this->build_trie(); }
98
99
      /// Searches in text for all occurrences of all strings in array arr.
100
101
      /// Time Complexity: O(text.size() + number of characters in arr)
      vector<vector<pair<int, int>>> find_occurrences(const string &text) {
102
103
        return this->aho find occurrences(text);
104
105
    };
```

9.7. Hashing

```
// OBS: CHOOSE THE OFFSET AND THE PRIMES BELOW!!
   class Hash {
3
     /// Prime numbers to be used in mod operations
     /// OBS: if you change m's size, please change the return type of both
     /// and _query methods.
5
     vector<int> m = {1000000007, 1000000009};
6
8
     // Case the alphabet goes from 'a' to 'z'.
     static constexpr int OFFSET = 'a';
10
     // Choose primes greater than the size of the alphabet.
     vector<int> prime = {31, 37};
11
12
1.3
     // Case the alphabet goes from 'A' to 'z'.
     // constexpr int OFFSET = 'A';
14
15
     // // Choose primes greater than the size of the alphabet.
16
     // vector<int> prime = {61, 67};
17
     vector<vector<int>> hash_table;
1.8
19
     vector<vector<int>> pot;
20
     // size of the string
21
     int n:
22
```

```
23 | private:
     int mod(int n, int m) {
25
       n %= m;
26
       if (n < 0)
27
         n += m;
28
       return n;
29
30
31
     /// Time Complexity: O(1)
     pair<int, int> _query(const int 1, const int r) {
32
33
       vector<int> ans(m.size());
34
35
       if (1 == 0) {
36
         for (int i = 0; i < m.size(); i++)
37
           ans[i] = hash_table[i][r];
38
39
         for (int i = 0; i < m.size(); i++)
40
           ans[i] =
41
                mod((hash\_table[i][r] - hash\_table[i][l - 1] * pot[i][r - l +
       1]),
                   m[i]);
42
43
44
45
       return {ans.front(), ans.back()};
46
47
48
     /// Builds the hash table and the pot table.
49
     /// Time Complexity: O(n)
50
51
     void build(string &s) {
52
       pot.resize(m.size(), vector<int>(this->n));
53
       hash_table.resize(m.size(), vector<int>(this->n));
54
       // Remapping the string
55
       for (char &c : s)
56
         c -= OFFSET;
57
58
       for (int i = 0; i < m.size(); i++) {
59
         hash_table[i][0] = s[0];
60
         pot[i][0] = 1;
61
         for (int j = 1; j < this->n; j++) {
           hash_table[i][j] = (s[j] + hash_table[i][j - 1] * prime[i]) % m[i];
62
63
           pot[i][j] = (pot[i][j-1] * prime[i]) % m[i];
64
65
66
     /// Constructor that is responsible for building the hash table and pot
       table.
70
71
     /// Time Complexity: O(n)
72
     Hash(string s) {
73
       assert(m.size() == prime.size());
74
       this->n = s.size();
75
76
       build(s);
77
78
79
     /// Returns the hash from 1 to r.
80
81
     /// Time Complexity: O(1) -> Actually O(number_of_primes)
     pair<int, int> query(const int 1, const int r) {
       assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
83
84
       return _query(1, r);
85
```

86 | };

9.8. Lcs K Strings

```
// Make the change below in SuffixArray code.
   int MaximumNumberOfStrings;
   void build_suffix_array() {
     vector<pair<Rank, int>> ranks(this->n + 1);
     vector<int> arr;
     for (int i = 1, separators = 0; i <= n; i++)</pre>
8
       if(this->s[i] > 0) {
10
         ranks[i] = pair<Rank, int>(Rank((int)this->s[i] +
        MaximumNumberOfStrings, 0), i);
11
         this->s[i] += MaximumNumberOfStrings;
12
13
         ranks[i] = pair<Rank, int>(Rank(separators, 0), i);
         this->s[i] = separators;
14
15
         separators++;
16
17
18
     RadixSort::sort_pairs(ranks, 256 + MaximumNumberOfStrings);
19
20
21
22 /// Program to find the LCS between k different strings.
23 ///
24 /// Time Complexity: O(n*log(n))
   /// Space Complexity: O(n*log(n))
26 int main() {
     int n;
2.7
28
     cin >> n:
3.0
31
     MaximumNumberOfStrings = n;
32
33
     vector<string> arr(n);
34
35
     int sum = 0;
36
     for(string &x: arr) {
       cin >> x;
37
38
       sum += x.size() + 1;
39
40
41
     string concat;
     vector<int> ind(sum + 1);
43
     int cnt = 0:
44
     for(string &x: arr) {
45
       if(concat.size())
46
         concat += (char)cnt;
47
       concat += x;
48
49
50
     cnt = 0:
     for(int i = 0; i < concat.size(); i++) {</pre>
52
       ind[i + 1] = cnt;
53
       if(concat[i] < MaximumNumberOfStrings)</pre>
54
55
56
57
     Suffix_Array say(concat);
58
     vector<int> sa = say.get_suffix_array();
     Sparse_Table spt(say.get_lcp());
```

```
61
     vector<int> freq(n);
62
     int cnt1 = 0:
63
64
     /// Ignore separators
65
     int i = n, j = n - 1;
66
     int ans = 0;
67
68
     while(true) {
69
70
       if(cnt1 == n) {
71
72
          ans = max(ans, spt.query(i, j - 1));
73
74
          int idx = ind[sa[i]];
75
          freg[idx]--;
          if(freq[idx] == 0)
76
77
           cnt1--;
78
          i++;
79
        } else if(j == (int)sa.size() - 1)
         break:
81
        else {
82
83
          int idx = ind[sa[j]];
84
          freq[idx]++;
85
          if(freq[idx] == 1)
86
            cnt1++;
87
88
89
90
     cout << ans << endl;
```

9.9. Lexicographically Smallest Rotation

```
int booth (string &s) {
     s += s;
     int n = s.size();
      vector<int> f(n, -1);
     int k = 0;
      for (int j = 1; j < n; j++) {
        int sj = s[j];
int i = f[j - k - 1];
9
10
        while(i != -1 \&\& sj != s[k + i + 1]) {
          if(sj < s[k + i + 1])
11
            k = j - i - 1;
12
13
          i = f[i];
14
15
        if(sj != s[k + i + 1]) {
16
          if(si < s[k])
17
          k = j;
f[j - k] = -1;
18
19
20
        else
          f[j - k] = i + 1;
21
22
23
      return k;
```

9.10. Manacher (Longest Palindrome)

```
https://medium.com/hackernoon/manachers-algorithm-explained-longest-palindromic-s
3 /// Create a string containing '#' characters between any two characters.
4 string get modified string(string &s) {
5 string ret;
    for(int i = 0; i < s.size(); i++) {</pre>
       ret.push_back('#');
       ret.push_back(s[i]);
8
9
10
    ret.push back('#');
11
     return ret;
12
13
14 /// Returns the first occurence of the longest palindrome based on the lps
15
16 /// Time Complexity: O(n)
17 string get_best(const int max_len, const string &str, const vector<int>
     for(int i = 0; i < lps.size(); i++) {</pre>
19
       if(lps[i] == max len) {
20
         string ans:
         int cnt = max_len / 2;
21
22
         int io = i - 1;
23
         while (cnt)
24
           if(str[io] != '#') {
25
             ans += str[io];
26
             cnt--;
27
2.8
           io--:
29
30
         reverse(ans.begin(), ans.end());
31
         if(str[i] != '#')
32
           ans += str[i];
         cnt = max len / 2;
33
34
         io = i + 1;
35
         while (cnt)
36
           if(str[io] != '#') {
             ans += str[io];
37
38
             cnt--:
39
40
           io++;
41
42
         return ans;
43
44
45 }
46
47 /// Returns a pair containing the size of the longest palindrome and the
        first occurence of it.
48 ///
49 /// Time Complexity: O(n)
50 pair<int, string> manacher(string &s) {
    int n = s.size();
     string str = get_modified_string(s);
52
     int len = (2 * n) + 1;
53
     //the i-th index contains the longest palindromic substring with the i-th
       char as the center
55
     vector<int> lps(len);
56
     int c = 0; //stores the center of the longest palindromic substring until
57
     int r = 0; //stores the right boundary of the longest palindromic
       substring until now
     int max len = 0;
```

```
for(int i = 0; i < len; i++) {</pre>
60
       //get mirror index of i
61
       int mirror = (2 * c) - i;
62
63
       //see if the mirror of i is expanding beyond the left boundary of
       current longest palindrome at center c
       //if it is, then take r - i as lps[i]
64
65
       //else take lps[mirror] as lps[i]
66
       if(i < r)
67
         lps[i] = min(r - i, lps[mirror]);
68
69
       //expand at i
70
       int a = i + (1 + lps[i]);
71
       int b = i - (1 + lps[i]);
72
       while(a < len && b >= 0 && str[a] == str[b]) {
73
         lps[i]++;
74
         a++;
75
         b--;
76
77
78
       //check if the expanded palindrome at i is expanding beyond the right
       boundary of current longest palindrome at center c
79
       //if it is, the new center is i
80
       if(i + lps[i] > r) {
81
         c = i;
82
         r = i + lps[i];
83
84
         if(lps[i] > max_len) //update max_len
85
           max_len = lps[i];
86
87
88
89
     return make_pair(max_len, get_best(max_len, str, lps));
```

9.11. Suffix Array

```
namespace RadixSort {
2
   /// Sorts the array arr stably in ascending order.
3
   ///
4
   /// Time Complexity: O(n + max_element)
5
   /// Space Complexity: O(n + max_element)
6
   template <typename T>
7
   void sort(vector<T> &arr, const int max_element, int (*get_key)(T &),
8
             int begin = 0) {
     const int n = arr.size();
9
10
     vector<T> new order(n);
11
     vector<int> count(max_element + 1, 0);
12
13
     for (int i = begin; i < n; i++)
14
       count[get_key(arr[i])]++;
15
16
     for (int i = 1; i <= max_element; i++)</pre>
17
       count[i] += count[i - 1];
18
19
     for (int i = n - 1; i >= begin; i--) {
20
       new_order[count[get_key(arr[i])] - (begin == 0)] = arr[i];
21
       count[get_key(arr[i])]--;
22
23
24
     arr.swap(new_order);
25
26
27 /// Sorts an array by their pair of ranks stably in ascending order.
```

```
28 | template <typename T> void sort_pairs(vector<T> &arr, const int rank_size) {
29
     // Sort by the second rank
30
     RadixSort::sort<T>(
31
         arr, rank_size, [](T &item) { return item.first.second; }, 011);
32
33
     // Sort by the first rank
34
     RadixSort::sort<T>(
35
         arr, rank_size, [](T &item) { return item.first.first; }, 011);
36
37 | // namespace RadixSort
38
39 /// It is indexed by 0.
40 /// Let the given string be "banana".
41 ///
42 /// 0 banana
                                          5 a
43 /// 1 anana
                   Sort the Suffixes
                                          3 ana
44 /// 2 nana
                   ----->
                                         1 anana
45 /// 3 ana
                    alphabetically
                                         0 banana
46 /// 4 na
                                         4 na
47 /// 5 a
                                         2 nana
48 /// So the suffix array for "banana" is {5, 3, 1, 0, 4, 2}
49 | ///
50 /// LCP
51 ///
52 | /// 1 a
53 /// 3 ana
54 /// 0 anana
55 |/// 0 banana
56 | /// 2 na
57 /// O nana (The last position will always be zero)
58 ///
59 /// So the LCP for "banana" is {1, 3, 0, 0, 2, 0}
60 ///
61 | class Suffix_Array {
62
   private:
63
     string s;
64
     int n;
65
66
     typedef pair<int, int> Rank;
67
68
   public:
     Suffix_Array(string &s) {
69
70
       this->n = s.size();
71
       this -> s = s;
72
       // little optimization, remove the line above
73
       // this->s.swap(s);
74
75
       this->sa = build_suffix_array();
76
       this->lcp = build_lcp();
77
78
79
   private:
80
     /// The vector containing the ranks will be present at ret
81
     void build_ranks(const vector<pair<Rank, int>> &ranks, vector<int> &ret) {
82
       ret[ranks[0].second] = 1;
83
       for (int i = 1; i < n; i++)
84
         // If their rank are equal, than its position should be the same.
85
         if (ranks[i - 1].first == ranks[i].first)
86
           ret[ranks[i].second] = ret[ranks[i - 1].second];
87
88
           ret[ranks[i].second] = ret[ranks[i - 1].second] + 1;
89
90
91
92
     /// Builds the Suffix Array for the string s.
```

```
94
      /// Time Complexity: O(n*log(n))
 95
      /// Space Complexity: O(n)
 96
      vector<int> build_suffix_array() {
 97
        // This tuple below represents the rank and the index associated with it.
 98
        vector<pair<Rank, int>> ranks(this->n);
 99
        vector<int> arr(this->n);
100
        for (int i = 0; i < n; i++)</pre>
101
102
          ranks[i] = pair<Rank, int>(Rank(s[i], 0), i);
103
104
        RadixSort::sort_pairs(ranks, 256);
105
        build ranks (ranks, arr);
106
107
108
          int jump = 1;
109
          int max_rank = arr[ranks.back().second];
110
          // It will be compared intervals a pair of intervals (i, jump-1), (i +
           // jump, i + 2*jump - 1). The variable jump is always a power of 2.
111
112
           while (max_rank != this->n) {
            for (int i = 0; i < this->n; i++) {
113
               ranks[i].first.first = arr[i];
114
115
               ranks[i].first.second = (i + jump < this->n ? arr[i + jump] : 0);
116
               ranks[i].second = i;
117
118
119
            RadixSort::sort_pairs(ranks, n);
120
            build_ranks(ranks, arr);
121
122
            max_rank = arr[ranks.back().second];
123
             jump \star = 2;
124
125
126
127
        vector<int> sa(this->n);
128
        for (int i = 0; i < this->n; i++)
129
          sa[arr[i] - 1] = i;
130
        return sa;
131
132
133
      /// Builds the lcp (Longest Common Prefix) array for the string s.
      /// A value lcp[i] indicates length of the longest common prefix of the
134
      /// suffixes indexed by i and i + 1. Implementation of the Kasai's
        Algorithm.
136
      /// Time Complexity: O(n)
138
      /// Space Complexity: O(n)
      vector<int> build_lcp() {
139
140
        lcp.resize(n, 0);
141
        vector<int> inverse suffix(this->n);
142
143
        for (int i = 0; i < this -> n; i++)
144
         inverse suffix[sa[i]] = i;
145
146
        int k = 0:
147
148
        for (int i = 0; i < this->n; i++) {
          if (inverse_suffix[i] == this->n - 1) {
149
150
            k = 0;
151
            continue;
152
153
          int j = sa[inverse_suffix[i] + 1];
154
155
156
          while (i + k < this - n \& j + k < this - n \& s[i + k] == s[j + k])
```

```
157
            k++;
158
159
           lcp[inverse_suffix[i]] = k;
160
161
          if (k > 0)
162
            k--;
163
164
165
        return lcp;
166
167
168 public:
169
      vector<int> sa;
170
      vector<int> lcp;
171
172
      /// LCS of two strings A and B.
173
174
      /// The string s must be initialized in the constructor as the string (A +
         '$'
175
      /// + B).
176
      ///
      /// The string A starts at index 1 and ends at index (separator - 1).
177
      /// The string B starts at index (separator + 1) and ends at the end of the
179
      /// string.
180
181
      /// Time Complexity: O(n)
182
      /// Space Complexity: O(1)
183
      int lcs(int separator) {
184
        assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
185
186
        int ans = 0;
187
188
         for (int i = 0; i + 1 < this->sa.size(); <math>i++) {
189
           int left = this->sa[i];
           int right = this->sa[i + 1];
190
191
192
           if ((left < separator && right > separator) ||
193
               (left > separator && right < separator))
194
             ans = max(ans, lcp[i]);
195
196
197
        return ans;
198
199 };
```

9.12. Suffix Array Pessoa

```
1 // OBS: Suffix Array build code imported from:
 2 //
       https://github.com/gabrielpessoal/Biblioteca-Maratona/blob/master/code/String/Suf
 3 // Because it's faster.
5 /// It is indexed by 0.
6 /// Let the given string be "banana".
7 1//
8 /// 0 banana
                                        5 a
9 /// 1 anana
                   Sort the Suffixes
                   ----->
10 /// 2 nana
                                        1 anana
11 /// 3 ana
                   alphabetically
                                        0 banana
12 /// 4 na
                                        4 na
13 /// 5 a
                                        2 nana
14 /// So the suffix array for "banana" is \{5, 3, 1, 0, 4, 2\}
15 ///
16 /// LCP
```

17 | ///

```
18 /// 1 a
19 /// 3 ana
20 /// 0 anana
21 /// 0 banana
22 /// 2 na
23 /// 0 nana (The last position will always be zero)
24 ///
25 /// So the LCP for "banana" is {1, 3, 0, 0, 2, 0}
26 ///
   class Suffix Array {
27
28
   private:
29
     string s;
30
     int n;
31
32
     typedef pair<int, int> Rank;
33
34
   public:
35
     Suffix Array(string &s) {
36
       this->n = s.size();
37
       this->s = s:
38
       // little optimization, remove the line above
39
       // this->s.swap(s):
40
41
       this->sa = build_suffix_array();
42
       this->lcp = build_lcp();
43
44
45
   private:
46
     /// Builds the Suffix Array for the string s.
47
48
     /// Time Complexity: O(n*log(n))
49
     /// Space Complexity: O(n)
50
     vector<int> build suffix array() {
51
       int n = this -> s.size(), c = 0;
52
       vector<int> temp(n), posBucket(n), bucket(n), bpos(n), out(n);
53
       for (int i = 0; i < n; i++)
54
         out[i] = i;
55
       sort(out.begin(), out.end(),
56
             [&] (int a, int b) { return this->s[a] < this->s[b]; });
       for (int i = 0; i < n; i++) {
57
58
         bucket[i] = c;
59
         if (i + 1 == n || this->s[out[i]] != this->s[out[i + 1]])
60
           c++:
61
62
       for (int h = 1; h < n && c < n; h <<= 1) {
         for (int i = 0; i < n; i++)
63
64
           posBucket[out[i]] = bucket[i];
65
         for (int i = n - 1; i >= 0; i--)
           bpos[bucket[i]] = i;
66
         for (int i = 0; i < n; i++) {
67
68
           if (out[i] >= n - h)
69
             temp[bpos[bucket[i]]++] = out[i];
70
71
         for (int i = 0; i < n; i++) {
72
           if (out[i] >= h)
73
             temp[bpos[posBucket[out[i] - h]]++] = out[i] - h;
74
75
76
         for (int i = 0; i + 1 < n; i++) {
77
           int a = (bucket[i] != bucket[i + 1]) || (temp[i] >= n - h) ||
78
                    (posBucket[temp[i + 1] + h] != posBucket[temp[i] + h]);
79
           bucket[i] = c:
80
           c += a;
81
```

```
bucket[n - 1] = c++;
82
83
          temp.swap(out);
84
85
        return out;
86
27
88
      /// Builds the lcp (Longest Common Prefix) array for the string s.
      /// A value lcp[i] indicates length of the longest common prefix of the
89
      /// suffixes indexed by i and i + 1. Implementation of the Kasai's
         Algorithm.
 91
      /// Time Complexity: O(n)
92
93
      /// Space Complexity: O(n)
      vector<int> build_lcp() {
94
95
        lcp.resize(n, 0);
96
        vector<int> inverse suffix(this->n);
97
98
         for (int i = 0; i < this->n; i++)
99
          inverse suffix[sa[i]] = i;
100
        int k = 0:
101
102
103
         for (int i = 0; i < this->n; i++) {
          if (inverse_suffix[i] == this->n - 1) {
104
105
            k = 0;
106
            continue;
107
108
109
          int j = sa[inverse_suffix[i] + 1];
110
111
          while (i + k < this -> n && j + k < this -> n && s[i + k] == s[j + k])
112
            k++;
113
114
          lcp[inverse suffix[i]] = k;
115
116
          if (k > 0)
117
            k--:
118
119
120
         return lcp:
121
122
123 public:
      vector<int> sa:
      vector<int> lcp;
126
127
      /// LCS of two strings A and B.
128
      ///
129
      /// The string s must be initialized in the constructor as the string (A +
        '$'
130
      /// + B).
131
      ///
132
      /// The string A starts at index 1 and ends at index (separator - 1).
133
      /// The string B starts at index (separator + 1) and ends at the end of the
134
      /// string.
135
      ///
136
      /// Time Complexity: O(n)
      /// Space Complexity: O(1)
137
138
      int lcs(int separator) {
139
        assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
140
141
        int ans = 0;
142
143
         for (int i = 0; i + 1 < this->sa.size(); i++) {
144
          int left = this->sa[i];
```

```
145
          int right = this->sa[i + 1];
146
147
           if ((left < separator && right > separator) ||
148
               (left > separator && right < separator))
            ans = max(ans, lcp[i]);
149
150
151
152
        return ans;
153
154
    };
```

53

54

55

56

57

58

59

60

61

62 63

64 65

66

67

68

69

7.0

71

72

73

74

75

76

77

78

79

80

81

82

83

84

85

86

87

88

89

90

91

92

93

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95

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100

101

102

103

104

105

106

107

108

109

110

111

112

113

114

115

116

9.13. Suffix Array With Additional Memory

```
1 namespace RadixSort {
   /// Sorts the array arr stably in ascending order.
  ///
4 /// Time Complexity: O(n + max_element)
  /// Space Complexity: O(n + max_element)
   template <typename T>
   void sort(vector<T> &arr, const int max_element, int (*get_key)(T &),
8
             int begin = 0) {
     const int n = arr.size();
10
     vector<T> new_order(n);
11
     vector<int> count(max element + 1, 0);
12
     for (int i = begin; i < n; i++)</pre>
13
14
       count[get_key(arr[i])]++;
15
     for (int i = 1; i <= max_element; i++)</pre>
16
17
       count[i] += count[i - 1];
18
19
     for (int i = n - 1; i >= begin; i--) {
       new_order[count[get_key(arr[i])]] = arr[i];
20
21
       count[get_key(arr[i])]--;
22
23
24
     arr = new_order;
25
26
27
   /// Sorts an array by their pair of ranks stably in ascending order.
   template <typename T> void sort_pairs(vector<T> &arr, const int rank_size) {
29
     // Sort by the second rank
     RadixSort::sort<T>(
30
31
         arr, rank size, [](T &item) { return item.first.second; }, 111);
32
33
     // Sort by the first rank
34
     RadixSort::sort<T>(
35
         arr, rank_size, [](T &item) { return item.first.first; }, 111);
36
37
   } // namespace RadixSort
38
39
   /// It is indexed by 1.
   class Suffix_Array {
40
41
   private:
42
     string s;
43
     int n;
44
     typedef pair<int, int> Rank;
     vector<int> suffix_array;
47
     vector<int> lcp;
48
49
     vector<vector<int>> rank_table;
50
     vector<int> log_array;
51
```

```
52 | public:
     Suffix_Array(const string &s) {
       this->n = s.size();
       this->s = "#" + s;
       build_log_array();
       build_suffix_array();
       lcp = build_lcp();
     vector<int> build_ranks(const vector<pair<Rank, int>> &ranks) {
       vector<int> arr(this->n + 1);
       arr[ranks[1].second] = 1;
       for (int i = 2; i <= n; i++) {
         // If their rank are equal, than its position should be the same.
         if (ranks[i - 1].first == ranks[i].first)
           arr[ranks[i].second] = arr[ranks[i - 1].second];
         else
           arr[ranks[i].second] = arr[ranks[i - 1].second] + 1;
       return arr:
     /// Builds the Suffix Array for the string s.
     /// Time Complexity: O(n*log(n))
     /// Space Complexity: O(n*log(n))
     void build suffix array() {
       // This tuple below represents the rank and the index associated with it.
       vector<pair<Rank, int>> ranks(this->n + 1);
       vector<int> arr;
       int rank table size = 0;
       this->rank_table.resize(log_array[this->n] + 2);
       for (int i = 1; i <= this->n; i++)
         ranks[i] = pair<Rank, int>(Rank(s[i], 0), i);
       // Inserting only the ranks in the table.
       transform(ranks.begin(), ranks.end(),
                  back_inserter(rank_table[rank_table_size++]),
                  [](pair<Rank, int> &pair) { return pair.first.first; });
       RadixSort::sort pairs(ranks, 256);
       arr = build ranks(ranks);
         int jump = 1;
         int max_rank = arr[ranks.back().second];
         // It will be compared intervals a pair of intervals (i, jump-1), (i +
         // jump, i + 2*jump - 1). The variable jump is always a power of 2.
         while (jump < n) {</pre>
           for (int i = 1; i <= this->n; i++) {
             ranks[i].first.first = arr[i];
             ranks[i].first.second = (i + jump <= this->n ? arr[i + jump] : 0);
             ranks[i].second = i;
           // Inserting only the ranks in the table.
           transform(ranks.begin(), ranks.end(),
                     back_inserter(rank_table[rank_table_size++]),
```

```
117
                        [](pair<Rank, int> &pair) { return pair.first.first; });
118
                                                                                         182
                                                                                                  for (int i = 2; i \le this -> n; i++)
119
             RadixSort::sort_pairs(ranks, n);
                                                                                         183
                                                                                                    log_array[i] = log_array[i / 2] + 1;
120
                                                                                         184
121
                                                                                         185
             arr = build ranks(ranks);
122
                                                                                         186
                                                                                             public:
123
             max rank = arr[ranks.back().second];
                                                                                         187
                                                                                               const vector<int> &get_suffix_array() { return suffix_array; }
124
             jump \star = 2;
                                                                                         188
125
                                                                                         189
                                                                                               const vector<int> &get_lcp() { return lcp; }
126
                                                                                         190
                                                                                               /// LCS of two strings A and B.
           for (int i = 1; i <= n; i++) {
127
                                                                                         191
128
             ranks[i].first.first = arr[i];
                                                                                         192
                                                                                               ///
129
             ranks[i].first.second = (i + jump <= this->n ? arr[i + jump] : 0);
                                                                                         193
                                                                                               /// The string s must be initialized in the constructor as the string (A +
130
             ranks[i].second = i;
                                                                                               /// + B).
                                                                                         194
131
132
                                                                                         195
                                                                                               ///
133
           // Inserting only the ranks in the table.
                                                                                         196
                                                                                               /// The string A starts at index 1 and ends at index (separator - 1).
134
           transform(ranks.begin(), ranks.end(),
                                                                                         197
                                                                                               /// The string B starts at index (separator + 1) and ends at the end of the
135
                     back_inserter(rank_table[rank_table_size++]),
                                                                                         198
                                                                                               /// string.
136
                     [](pair<Rank, int> &pair) { return pair.first.first; });
                                                                                         199
                                                                                               111
                                                                                               /// Time Complexity: O(n)
137
                                                                                         200
                                                                                               /// Space Complexity: 0(1)
138
139
         this->suffix arrav.resize(this->n + 1);
                                                                                         202
                                                                                               int lcs(int separator) {
140
         for (int i = 1; i <= this->n; i++)
                                                                                         203
                                                                                                 separator++;
                                                                                         204
141
           this->suffix_array[arr[i]] = i;
                                                                                                 assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
142
                                                                                         205
143
                                                                                         206
                                                                                                 int ans = 0;
      /// Builds the lcp (Longest Common Prefix) array for the string s.
                                                                                         207
144
145
      /// A value lcp[i] indicates length of the longest common prefix of the
                                                                                         208
                                                                                                  for (int i = 1; i < this -> n - 1; i++) {
146
      /// suffixes indexed by i and i + 1. Implementation of the Kasai's
                                                                                         209
                                                                                                    int left = this->suffix array[i];
         Algorithm.
                                                                                         210
                                                                                                    int right = this->suffix_array[i + 1];
147
                                                                                         211
148
      /// Time Complexity: O(n)
                                                                                         212
                                                                                                    if ((left < separator && right > separator) ||
149
      /// Space Complexity: O(n)
                                                                                         213
                                                                                                        (left > separator && right < separator))</pre>
150
      vector<int> build_lcp() {
                                                                                         214
                                                                                                      ans = max(ans, lcp[i]);
151
         vector<int> lcp(this->n + 1, 0);
                                                                                         215
152
         vector<int> inverse_suffix(this->n + 1, 0);
                                                                                         216
153
                                                                                         217
                                                                                                 return ans;
154
         for (int i = 1; i \le n; i++)
                                                                                         218
155
           inverse_suffix[suffix_array[i]] = i;
                                                                                         219
156
                                                                                         220
                                                                                               /// Compares two substrings beginning at indexes i and j of a fixed length.
157
        int k = 0;
                                                                                         221
                                                                                               ///
158
                                                                                         222
                                                                                               /// OBS: Necessary build rank_table (uncomment build_suffix_array) and
159
         for (int i = 1; i \le n; i++) {
                                                                                                 build
160
           if (inverse suffix[i] == n) {
                                                                                         223
                                                                                               /// log_array.
             k = 0;
                                                                                         224
                                                                                               ///
162
             continue;
                                                                                         225
                                                                                               /// Time Complexity: O(1)
163
                                                                                         226
                                                                                               /// Space Complexity: O(1)
164
                                                                                         227
                                                                                               int compare (const int i, const int j, const int length) {
165
           int j = suffix array[inverse suffix[i] + 1];
                                                                                         228
                                                                                                 assert(1 <= i \& \& i <= this -> n \& \& 1 <= j \& \& j <= this -> n);
166
                                                                                         229
                                                                                                 assert(!this->log_array.empty() && !this->rank_table.empty());
           while (i + k <= this->n && j + k <= this->n && s[i + k] == s[j + k])
                                                                                                 assert(i + length - 1 \stackrel{\leftarrow}{=} this->n && j + length - 1 \stackrel{\leftarrow}{=} this->n);
167
                                                                                         230
168
            k++;
                                                                                         231
169
                                                                                         232
                                                                                                  // Greatest k such that 2^k <= 1
           lcp[inverse_suffix[i]] = k;
170
                                                                                         233
                                                                                                 const int k = this->log_array[length];
171
                                                                                         234
172
           if (k > 0)
                                                                                         235
                                                                                                 const int jump = length - (1 << k);</pre>
173
             k--;
                                                                                         236
174
                                                                                         237
                                                                                                 const pair<int, int> iRank = {
175
                                                                                         238
                                                                                                      this->rank table[k][i],
176
         return lcp;
                                                                                         239
                                                                                                      (i + jump \le this - n ? this - rank_table[k][i + jump] : -1);
177
                                                                                                 const pair<int, int> jRank = {
                                                                                        240
178
                                                                                        241
                                                                                                      this->rank_table[k][j],
179
      void build_log_array() {
                                                                                        242
                                                                                                      (j + jump \le this - n ? this - rank_table[k][j + jump] : -1);
180
        log_array.resize(this->n + 1, 0);
                                                                                        243
```

```
return iRank == jRank ? 0 : iRank < jRank ? -1 : 1;
244
245
246
    };
    9.14. Trie
    class Trie {
    private:
      static const int INT LEN = 31;
```

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64 65

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123

```
// static const int INT LEN = 63;
     struct Node {
8
       map<char, Node *> next;
9
       int id;
       // cnt counts the number of words which pass in that node
10
       int cnt = 0;
11
12
       /\!/ word counts the number of words ending at that node
13
       int word_cnt = 0;
14
1.5
       Node(const int x) : id(x) {}
16
     };
17
18
   private:
     int trie_size = 0;
19
     // contains the next id to be used in a node
20
21
     int node cnt = 0;
     Node *trie_root = this->make_node();
22
23
24
25
     Node *make_node() { return new Node(node_cnt++); }
26
27
     int trie insert(const string &s) {
       Node *aux = this->root();
28
       for (const char c : s) {
29
30
         if (!aux->next.count(c))
           aux->next[c] = this->make_node();
31
32
         aux = aux->next[c];
33
         ++aux->cnt;
34
35
       ++aux->word cnt;
36
       ++this->trie size:
37
       return aux->id:
38
39
40
     void trie erase(const string &s) {
41
       Node *aux = this->root();
       for (const char c : s) {
42
43
         Node *last = aux;
44
         aux = aux->next[c];
45
         --aux->cnt;
46
         if (aux -> cnt == 0) {
47
           last->next.erase(c);
48
           aux = nullptr;
49
           break;
50
51
52
       if (aux != nullptr)
53
         --aux->word cnt;
54
       --this->trie_size;
55
56
57
     int trie_count(const string &s) {
       Node *aux = this->root();
```

```
59
       for (const char c : s) {
         if (aux->next.count(c))
           aux = aux->next[c];
         else
           return 0;
       return aux->word cnt;
     int trie_query_xor_max(const string &s) {
       Node *aux = this->root();
       int ans = 0:
       for (const char c : s) {
         const char inv = (c == '0' ? '1' : '0');
         if (aux->next.count(inv)) {
           ans = (ans << 111) | (inv - '0');
           aux = aux->next[inv];
         } else {
           ans = (ans << 111) | (c - '0');
           aux = aux->next[c];
       return ans;
84 public:
     Trie() {}
     Node *root() { return this->trie_root; }
     int size() { return this->trie_size; }
     /// Returns the number of nodes present in the trie.
     int node_count() { return this->node_cnt; }
     /// Inserts s in the trie.
     /// Returns the id of the last character of the string in the trie.
     /// Time Complexity: O(s.size())
     int insert(const string &s) { return this->trie insert(s); }
     /// Inserts the binary representation of x in the trie.
     /// Time Complexity: O(log x)
     int insert(const int x) {
       assert (x >= 0);
       // converting x to binary representation
       return this->trie insert(bitset<INT LEN>(x).to string());
     /// Removes the string s from the trie.
     /// Time Complexity: O(s.size())
     void erase(const string &s) { this->trie_erase(s); }
     /// Removes the binary representation of x from the trie.
     /// Time Complexity: O(log x)
     void erase(const int x) {
       assert (x >= 0);
       // converting x to binary representation
       this->trie_erase(bitset<INT_LEN>(x).to_string());
```

```
/// Returns the number of maximum xor sum with x present in the trie.
125
      /// Time Complexity: O(log x)
int query_xor_max(const int x) {
126
127
128
        assert(x >= 0);
129
         // converting x to binary representation
        return this->trie_query_xor_max(bitset<INT_LEN>(x).to_string());
130
131
132
      /// Returns the number of strings equal to s present in the trie. ///
133
134
      ///
/// Time Complexity: O(s.size())
135
      int count(const string &s) { return this->trie_count(s); }
136
137
```