C++ Competitive Programming Library

DO NOT DISCLOSE OR DISTRIBUTE

bfs.07 - Bernardo Flores Salmeron

1. Template

```
#include <bits/stdc++.h>
   using namespace std;
   #define INF (111 << 62)
   #define pb push_back
   #define ii pair<int,int>
   #define OK cerr <<"OK"<< endl
   #define debug(x) cerr << \#x " = " << (x) << endl
10 #define ff first
11
   #define ss second
   #define int long long
   #define tt tuple<int, int, int>
13
  #define endl '\n'
15
16
   signed main () {
17
18
     ios_base::sync_with_stdio(false);
19
     cin.tie(NULL);
20
21
```

2. Data*structures*

2.1. Arvore Binaria

```
// C program to demonstrate delete operation in binary search tree
   #include < stdio.h>
3
   #include < stdlib.h>
   struct node {
     int kev;
     struct node *left, *right;
8
10 // A utility function to create a new BST node
11 struct node *newNode(int item) {
     struct node *temp = (struct node *)malloc(sizeof(struct node));
12
13
     temp->kev = item:
     temp->left = temp->right = NULL;
15
     return temp;
16
17
18 // A utility function to do inorder traversal of BST
19 void inorder(struct node *root) {
   if (root != NULL)
      inorder(root->left);
```

```
22
       printf("%d ", root->key);
23
       inorder(root->right);
24
25
26
   /* A utility function to insert a new node with given key in BST */
   struct node* insert(struct node* node, int key)
     /* If the tree is empty, return a new node */
29
     if (node == NULL) return newNode(key);
30
31
32
     /★ Otherwise, recur down the tree ★/
3.3
     if (key < node->key)
34
       node->left = insert(node->left, key);
35
       node->right = insert(node->right, key);
37
38
     /* return the (unchanged) node pointer */
39
     return node:
40
41
42 /* Given a non-empty binary search tree, return the node with minimum
    key value found in that tree. Note that the entire tree does not
     need to be searched. */
   struct node * minValueNode(struct node* node) {
45
46
     struct node* current = node;
47
     /* loop down to find the leftmost leaf */
49
     while (current->left != NULL)
50
       current = current->left;
51
52
     return current;
53
54
   /\star Given a binary search tree and a key, this function deletes the key
    and returns the new root */
   struct node* deleteNode(struct node* root, int key) {
    // base case
59
     if (root == NULL) return root;
60
61
     // If the key to be deleted is smaller than the root's key,
62
     // then it lies in left subtree
63
     if (key < root->key)
64
       root->left = deleteNode(root->left, key);
65
66
     // If the key to be deleted is greater than the root's key,
67
     // then it lies in right subtree
68
     else if (key > root->key)
69
       root->right = deleteNode(root->right, key);
70
     // if key is same as root's key, then This is the node
```

```
// to be deleted
 73
      else {
        // node with only one child or no child
 74
 75
        if (root->left == NULL) {
 76
          struct node *temp = root->right;
 77
          free (root);
 78
          return temp;
        } else if (root->right == NULL) {
 79
          struct node *temp = root->left;
 80
 81
          free (root.):
 82
          return temp;
 83
 84
 85
        // node with two children: Get the inorder successor (smallest
        // in the right subtree)
 86
 87
        struct node* temp = minValueNode(root->right);
 88
 89
        // Copy the inorder successor's content to this node
 90
        root->kev = temp->kev;
 91
 92
        // Delete the inorder successor
 93
        root->right = deleteNode(root->right, temp->key);
 94
 95
     return root:
 96
 97
 98
    // Driver Program to test above functions
 99
    int main() {
100
     /★ Let us create following BST
101
                50
102
103
            30
                    70
         20 40 60 80 */
104
105
      struct node *root = NULL;
106
      root = insert(root, 50);
107
      root = insert(root, 30);
108
109
      root = insert(root, 20);
      root = insert(root, 40);
110
111
      root = insert(root, 70);
112
      root = insert(root, 60);
      root = insert(root, 80);
113
114
      printf("Inorder traversal of the given tree \n");
115
      inorder(root);
117
      printf("\nDelete 20\n");
      root = deleteNode(root, 20);
119
      printf("Inorder traversal of the modified tree \n");
120
121
      inorder(root);
122
      printf("\nDelete 30\n");
123
      root = deleteNode(root, 30);
124
125
      printf("Inorder traversal of the modified tree \n");
126
      inorder(root):
127
128
      printf("\nDelete 50\n");
129
      root = deleteNode(root, 50);
      printf("Inorder traversal of the modified tree \n");
130
131
      inorder(root);
132
133
      return 0;
134
```

2.2. Bit2D

```
// INDEX BY ONE ALWAYS!!!
2 class BIT_2D {
3
    private:
    // row, column
    int n, m;
     vector<vector<int>> tree;
    // Returns an integer which constains only the least significant bit.
     int low(int i) {
11
       return i & (-i);
12
13
     void bit_update(const int x, const int y, const int delta) {
14
       for (int i = x; i < n; i + low(i))
15
16
         for(int j = y; j < m; j += low(j))</pre>
17
           this->tree[i][j] += delta;
1.8
19
2.0
     int bit query(const int x, const int y) {
21
       int ans = 0:
22
       for(int i = x; i > 0; i -= low(i))
         for(int j = y; j > 0; j -= low(j))
23
           ans += this->tree[i][j];
24
25
26
       return ans;
27
28
29
30
    // put the size of the array without 1 indexing.
     /// Time Complexity: O(n * m)
31
     BIT 2D(int n, int m) {
33
       this -> n = n + 1;
       this->m = m + 1;
34
35
36
       this->tree.resize(n, vector<int>(m, 0));
37
3.8
39
     /// Time Complexity: O(n * m * (log(n) + log(m)))
     BIT 2D(const vector<vector<int>> &mat) {
40
       // Check if it is 1 index.
41
42
       assert (mat[0][0] == 0);
43
       this->n = mat.size();
44
       this->m = mat.front().size();
45
       this->tree.resize(n, vector<int>(m, 0));
46
47
       for(int i = 1; i < n; i++)
         for(int j = 1; j < m; j++)
48
           update(i, j, mat[i][j]);
49
50
51
52
     /// Query from (1, 1) to (x, y).
53
54
     /// Time Complexity: O(log(n) + log(m))
     int prefix_query(const int x, const int y) {
56
       assert (0 < x); assert (x < this -> n);
57
       assert(0 < v); assert(v < this->m);
58
59
       return bit_query(x, y);
60
61
62
     /// Query from (x1, y1) to (x2, y2).
63
     ///
```

```
/// Time Complexity: O(log(n) + log(m))
     int query (const int x1, const int y1, const int x2, const int y2) {
66
       assert (0 < x1); assert (x1 <= x2); assert (x2 < this -> n);
67
       assert(0 < y1); assert(y1 <= y2); assert(y2 < this->m);
68
69
       return bit_query(x2, y2) - bit_query(x1 - 1, y2) - bit_query(x2, y1 - 1)
       + bit_query(x1 - 1, y1 - 1);
70
71
72
     /// Updates point (x, y).
73
74
     /// Time Complexity: O(log(n) + log(m))
75
     void update(const int x, const int y, const int delta) {
76
       assert (0 < x); assert (x < this -> n);
77
       assert(0 < y); assert(y < this->m);
78
79
       bit_update(x, y, delta);
80
81
   };
```

2.3. Merge Sort Tree (K-Esimo Maior Elemento Num Intervalo, Valores Maiores Que K Num Intervalo,

```
1 // retornar a qtd de números maiores q um numero k numa array de i...j
2
   struct Tree {
3
    vector<int> vet:
4
   Tree tree [4*(int)3e4];
   int arr[(int)5e4];
8
   int query(int 1,int r, int i, int j, int k, int pos) {
9
    if(1 > j || r < i)
10
       return 0;
11
12
     if(i <= 1 && r <= j) {
13
       auto it = upper_bound(tree[pos].vet.begin(), tree[pos].vet.end(), k);
       return tree[pos].vet.end()-it;
14
15
16
17
     int mid = (1+r) >> 1;
     return query(1, mid, i, j, k, 2*pos+1) + query(mid+1,r,i,j,k,2*pos+2);
19
20
   void build(int 1, int r, int pos) {
21
22
23
     if(1 == r) {
24
       tree[pos].vet.pb(arr[1]);
25
       return;
26
2.7
     int mid = (1+r) >> 1;
28
29
     build(1, mid, 2*pos+1);
30
     build (mid + 1, r, 2*pos+2);
31
32
     merge(tree[2*pos+1].vet.begin(), tree[2*pos+1].vet.end(),
       tree[2*pos+2].vet.begin(), tree[2*pos+2].vet.end(),
       back_inserter(tree[pos].vet));
33
```

2.4. Mos Algorithm

```
struct Tree {
int 1, r, ind;
```

```
4 Tree query[311111];
5 | int arr[311111];
6 | int freq[1111111];
7 | int ans[311111];
8 | int block = sqrt(n), cont = 0;
10 | bool cmp(Tree a, Tree b)
    if(a.l/block == b.l/block)
11
12
       return a.r < b.r;</pre>
     return a.l/block < b.l/block;
13
14
15
16
   void add(int pos) {
17
     freg[arr[pos]]++;
18
     if(freq[arr[pos]] == 1) {
19
       cont++;
20
21
22
   void del(int pos) {
     freg[arr[pos]]--;
     if(freq[arr[pos]] == 0)
25
       cont--:
26
27
   int main () {
     int n; cin >> n;
29
     block = sqrt(n);
30
     for(int i = 0; i < n; i++) {</pre>
31
32
        cin >> arr[i];
3.3
        freq[arr[i]] = 0;
34
35
36
     int m; cin >> m;
37
38
     for(int i = 0; i < m; i++) {
       cin >> query[i].l >> query[i].r;
39
40
       query[i].l--, query[i].r--;
41
       query[i].ind = i;
42
43
     sort (query, query + m, cmp);
45
     int s,e;
     s = e = querv[0].1:
47
     add(s);
      for(int i = 0; i < m; i++) {</pre>
49
        while(s > querv[i].l)
50
          add(--s);
51
        while(s < query[i].l)</pre>
52
          del(s++);
53
       while(e < query[i].r)</pre>
54
          add (++e);
       while(e > query[i].r)
55
56
          del(e--);
57
       ans[query[i].ind] = cont;
58
59
     for(int i = 0; i < m; i++)
60
61
       cout << ans[i] << endl;</pre>
62
```

2.5. Ordenacao De Estruturas (Pq, Etc)

```
1 struct cmp {
```

```
bool operator(ii a, ii b) {
   //ordena primeiro pelo first(decrescente), dps pelo second(crescente)
   if(a.first == b.first)
        return a.second < b.second;
   return a.first > b.first;
}

Ex: pq<ii,vector<ii>,cmp> fila;
```

2.6. Ordered Set (Policy Based Data Structures)

```
#include <bits/stdc++.h>
   #include <ext/pb_ds/assoc_container.hpp>
   #include <ext/pb_ds/trie_policy.hpp>
   using namespace std;
   using namespace __gnu_pbds;
   typedef tree<
   int,
10
   null_type,
11
   less<int>,
   rb_tree_tag,
12
13 tree_order_statistics_node_update>
14 ordered_set;
15
16 ordered set X:
17 | X.insert(1); X.insert(2);
18 X.insert(4); X.insert(8);
19 | X.insert(16);
20 // 1, 2, 4, 8, 16
21 // retorna o k-ésimo maior elemento a partir de 0
22 | cout << *X.find_by_order(1) << endl; // 2
23 cout << *X.find_by_order(2) << endl; // 4
24 | cout << *X.find_by_order(4) << endl; // 16
25 | cout << (end (X) == X.find_by_order (6)) << endl; // true</pre>
26
   // retorna o número de itens estritamente menores que o número
   cout << X.order of key(-5) << endl; // 0
   cout << X. order_of_key(1) << endl; // 0
30
   cout << X.order_of_key(3) << endl;
   cout << X.order_of_key(4) << endl;</pre>
   cout << X.order_of_key(400) << endl; // 5
```

2.7. Sqrt Decomposition

```
// Problem: Sum from 1 to r
   // Ver MO'S ALGORITHM
   // -----
4
   int getId(int indx,int blockSZ) {
       return indx/blockSZ;
7
   void init(int sz) {
8
      for(int i=0; i<=sz; i++)
    BLOCK[i]=inf;
10
11 | int query(int left, int right) {
12 int startBlockIndex=left/sqrt;
13 int endIBlockIndex = right / sqrt;
14 | int sum = 0;
15 for (int i = startBlockIndex + 1; i < endIBlockIndex; i++) {
16
           sum += blockSums[i];
17
```

2.8. Bit

```
1 /// INDEX THE ARRAY BY 1!!!
   class BIT {
   private:
     vector<int> bit;
5
     int n;
   private:
     int low(const int i) { return (i & (-i)); }
a
10
     // point update
     void bit_update(int i, const int delta) {
11
       while (i \le this -> n) {
12
1.3
          this->bit[i] += delta;
          i += this->low(i);
14
15
16
17
18
     // point query
19
     int bit_query(int i) {
       int sum = 0;
20
       while (i > 0)
21
22
         sum += bit[i];
23
         i -= this->low(i);
24
25
       return sum;
26
27
28
29
     BIT(const vector<int> &arr) { this->build(arr); }
30
31
     BIT(const int n) {
32
       // OBS: BIT IS INDEXED FROM 1
33
        // THE USE OF 1-BASED ARRAY IS RECOMMENDED
34
       this->n = n:
35
       this->bit.resize(n + 1, 0);
36
37
38
      // build the bit
     void build(const vector<int> &arr) {
39
       // OBS: BIT IS INDEXED FROM 1
40
       // THE USE OF 1-BASED ARRAY IS RECOMMENDED
41
       assert(arr.front() == 0);
42
43
       this->n = (int)arr.size() - 1;
44
       this->bit.resize(arr.size(), 0);
45
46
        for (int i = 1; i <= this->n; i++)
47
          this->bit_update(i, arr[i]);
48
49
50
     // point update
     void update(const int i, const int delta) {
51
52
       assert(1 <= i), assert(i <= this->n);
53
       this->bit update(i, delta);
54
55
56
     // point query
```

```
int query(const int i) {
58
       assert(1 <= i), assert(i <= this->n);
59
       return this->bit_query(i);
60
61
62
     // range query
63
     int query(const int 1, const int r) {
64
       assert (1 \le 1), assert (1 \le r), assert (r \le this - n);
        return this->bit_query(r) - this->bit_query(l - 1);
65
66
   };
```

2.9. Bit (Range Update)

```
1 /// INDEX THE ARRAY BY 1!!!
   class BIT {
3
    private:
     vector<int> bit1;
     vector<int> bit2;
     int n;
    private:
     int low(int i) {
10
       return (i & (-i));
11
12
13
     // point update
     void update(int i, const int delta, vector<int> &bit) {
14
15
       while(i <= this->n) {
16
         bit[i] += delta;
17
         i += this->low(i);
18
19
20
21
     // point query
22
     int query(int i, const vector<int> &bit) {
23
       int sum = 0;
24
       while (i > 0)
25
        sum += bit[i];
26
         i -= this->low(i);
27
28
       return sum;
29
30
31
     // build the bit
32
     void build(const vector<int> &arr) {
       // OBS: BIT IS INDEXED FROM 1
33
34
       // THE USE OF 1-BASED ARRAY IS MANDATORY
35
       assert(arr.front() == 0);
36
       this->n = (int)arr.size() - 1;
       this->bit1.resize(arr.size(), 0);
37
38
       this->bit2.resize(arr.size(), 0);
39
40
       for (int i = 1; i <= this->n; i++)
41
         this->update(i, arr[i]);
42
43
44
     BIT(const vector<int> &arr) {
46
       this->build(arr);
47
48
49
     BIT (const int n) {
       // OBS: BIT IS INDEXED FROM 1
```

```
// THE USE OF 1-BASED ARRAY IS MANDATORY
52
        this->n = n;
        this->bit1.resize(n + 1, 0);
53
        this->bit2.resize(n + 1, 0);
54
55
56
57
      // range update
      void update(const int 1, const int r, const int delta) {
5.8
59
        assert(1 <= 1); assert(1 <= r); assert(r <= this->n);
60
        this->update(l, delta, this->bit1);
        this->update(r + 1, -delta, this->bit1);
61
62
        this->update(1, delta * (1 - 1), this->bit2);
63
        this->update(r + 1, -delta * r, this->bit2);
64
65
66
      // point update
67
      void update(const int i, const int delta) {
68
        assert(1 <= i); assert(i <= this->n);
69
        this->update(i, i, delta);
70
71
72
      // range query
      int querv(const int 1, const int r) {
74
        assert(1 <= 1); assert(1 <= r); assert(r <= this->n);
75
        return this->query(r) - this->query(l - 1);
76
77
78
      // point query
79
      int query(const int i) {
80
        assert(1 <= i); assert(i <= this->n);
81
        return (this->query(i, this->bit1) * i) - this->query(i, this->bit2);
82
83 };
84
    // TESTS
85
    // signed main()
   // {
87
8.8
   // vector<int> input = {0,1,2,3,4,5,6,7};
89
91 // BIT ft(input);
92
93 // assert (1 == ft.query(1));
94 // assert (3 == ft.querv(2));
95 // assert (6 == ft.querv(3));
96 // assert (10 == ft.query(4));
97 // assert (15 == ft.guery(5));
98 // assert (21 == ft.query(6));
99 // assert (28 == ft.query(7));
100 // assert (12 == ft.query(3,5));
101 // assert (21 == ft.query(1,6));
102 // assert (28 == ft.query(1,7));
103 // }
```

2.10. Counting Inversions (Minimum Number Of Adjacent Swaps To Sort Array)

```
// REQUIRES bit.cpp!!
// REQUIRES point_compresion.cpp!!
int count_inversions(vector<int> &arr) {
    arr = compress(arr);
    int ans = 0;
    BIT bit(arr.size());
    for (int i = arr.size() - 1; i > 0; --i) {
```

57

58

59

60

61

62

63

64

65 66

67

68

69

70

71

72

73

74

75

76

77

78

79

80

81

82

83

84

85

86

87

88

89

90

91

92

93

94

95

96

97

98

99

100

102

103

104

105

106

107

108

109

110

111

112

113

114

115

116

117 118

2.11. Persistent Segment Tree

```
class Persistent Seg Tree {
     struct Node {
3
       int val;
4
       Node *left, *right;
5
       Node() {}
6
       Node (int v, Node \star1, Node \starr) : val(v), left(1), right(r) {}
8
   #define NEUTRAL_NODE Node(0, nullptr, nullptr);
9
    Node NEUTRAL NODE = Node(0, nullptr, nullptr);
10
  public:
11
12
     int merge_nodes(const int x, const int y) { return x + y; }
13
14
  private:
15
     int n;
16
     vector<Node *> version;
17
18
     Persistent_Seq_Tree() { this->n = -1; }
19
20
     /// Builds version[0] with the values in the array.
21
     /// Time complexity: O(n)
22
     Node *pst_build(Node *node, const int 1, const int r,
23
                      const vector<int> &arr) {
24
25
       node = new NEUTRAL_NODE;
       if (1 == r) {
26
27
         node->val = arr[l]:
28
         return node;
29
30
       int mid = (1 + r) / 2;
31
32
       node->left = pst_build(node->left, 1, mid, arr);
33
       node->right = pst_build(node->right, mid + 1, r, arr);
34
       node->val = merge_nodes(node->left->val, node->right->val);
35
       return node;
36
37
38
     /// Builds version[0] with 0.
39
     ///
40
     /// Time complexity: O(n)
41
     Node *pst_build_empty(Node *node, const int 1, const int r) {
42
       node = new NEUTRAL_NODE;
43
       if (1 == r)
44
         return node;
45
46
       int mid = (1 + r) / 2;
47
       node->left = pst build empty(node->left, 1, mid);
48
       node->right = pst_build_empty(node->right, mid + 1, r);
49
       node->val = merge_nodes(node->left->val, node->right->val);
50
       return node;
51
52
53
     Node *pst_update(Node *cur_tree, Node *prev_tree, const int 1, const int r,
54
                       const int idx, const int delta) {
55
       if (1 > idx || r < idx) {
56
         if (cur_tree != nullptr)
```

```
return cur tree:
     return prev tree;
   if (cur tree == nullptr)
     cur_tree = new Node(prev_tree->val, prev_tree->left, prev_tree->right);
     cur_tree = new Node(cur_tree->val, cur_tree->left, cur_tree->right);
   if (1 == r) {
     cur tree->val += delta;
     return cur_tree;
   int mid = (1 + r) / 2;
   cur tree->left =
       pst_update(cur_tree->left, prev_tree->left, 1, mid, idx, delta);
   cur_tree->right =
       pst_update(cur_tree->right, prev_tree->right, mid + 1, r, idx,
   cur_tree->val = merge_nodes(cur_tree->left->val, cur_tree->right->val);
   return cur tree;
 int pst_query(Node *node, const int 1, const int r, const int i,
               const int j) {
   if (1 > j || r < i)
     return _NEUTRAL_NODE.val;
   if (i <= 1 && r <= i)
     return node->val:
   int mid = (1 + r) / 2;
   return merge_nodes(pst_query(node->left, 1, mid, i, j),
                      pst_query(node->right, mid + 1, r, i, j));
public:
 Persistent_Seq_Tree(const int n, const int number_of_versions) {
   this->n = n;
   version.resize(number_of_versions);
   this->version[0] = this->pst_build_empty(this->version[0], 0, this->n -
   1);
 /// Constructor that allows to pass initial values to the leafs.
 Persistent_Seq_Tree(const vector<int> &arr, const int number_of_versions) {
   this->n = arr.size();
   version.resize(number of versions);
   this->version[0] = this->pst_build(this->version[0], 0, this->n - 1,
 /// Links the root of a version to a previous version.
 /// Time Complexity: O(1)
 void link(const int version, const int prev_version) {
   assert (this->n > -1);
   assert(0 <= prev_version);</pre>
   assert (prev version <= version);
   assert(version < this->version.size());
   this->version[version] = this->version[prev_version];
```

/// Updates an index in cur_tree based on prev_tree with a delta.

```
119
120
      /// Time Complexity: O(log(n))
121
      void update (const int cur_version, const int prev_version, const int idx,
122
                   const int delta) {
123
        assert (this->n > -1);
124
        assert(0 <= prev_version);</pre>
125
        assert (prev_version <= cur_version);</pre>
126
        assert(cur version < this->version.size());
        this->version[cur_version] = this->pst_update(this->version[cur_version],
127
128
        this->version[prev version],
129
                                                         0, this->n - 1, idx,
        delta);
130
131
132
      /// Query from 1 to r.
133
134
      /// Time Complexity: O(log(n))
135
      int query(const int version, const int 1, const int r) {
136
        assert (this->n > -1):
        assert(this->version[version] != nullptr);
138
        assert(0 <= 1);
139
        assert(1 <= r);
140
        assert(r < this->n);
        return this->pst_query(this->version[version], 0, this->n - 1, 1, r);
141
142
143
    };
```

2.12. Segment Tree

```
class Seq_Tree {
   public:
     struct Node {
       int val, lazy;
6
7
       Node (const int val, const int lazy) : val(val), lazy(lazy) {}
8
     };
9
   private:
10
11
    // // range sum
     // Node NEUTRAL_NODE = Node(0, 0);
12
     // Node merge_nodes(const Node &x, const Node &y) {
13
14
     // return Node(x.val + v.val, 0);
15
     // }
     // void apply_lazy(const int 1, const int r, const int pos) {
16
17
     // tree[pos].val += (r - l + 1) * tree[pos].lazy;
18
     // }
19
20
     // // RMQ max
     // Node NEUTRAL_NODE = Node(-INF, 0);
21
22
     // Node merge_nodes(const Node &x, const Node &y) {
23
     // return Node (max (x.val, y.val), 0);
24
     // }
2.5
     // void apply_lazy(const int 1, const int r, const int pos) {
26
     // tree[pos].val += tree[pos].lazy;
27
     // }
28
29
     // // RMO min
30
     // Node NEUTRAL_NODE = Node(INF, 0);
31
     // Node merge_nodes(const Node &x, const Node &y) {
32
     // return Node (min (x.val, y.val), 0);
33
     // }
     // void apply_lazy(const int 1, const int r, const int pos) {
```

```
// tree[pos].val += tree[pos].lazy;
36
     // }
37
38
     // XOR
39
     // Only works with point updates
40
     // Node NEUTRAL_NODE = Node(0, 0);
41
     // Node merge_nodes (const Node &x, const Node &y) {
42
     // return Node(x.val ^ y.val, 0);
43
44
     // void apply_lazy(const int 1, const int r, const int pos) {}
45
46
   private:
47
     int n;
48
49
   public:
50
     vector<Node> tree;
51
52 private:
53
     void st_propagate(const int 1, const int r, const int pos) {
54
       if (tree[pos].lazy != 0) {
55
         apply_lazy(l, r, pos);
56
         if (1 != r) {
57
           tree[2 * pos + 1].lazv += tree[pos].lazv;
58
           tree[2 * pos + 2].lazy += tree[pos].lazy;
59
60
         tree[pos].lazy = 0;
61
62
63
64
     Node st_build(const int 1, const int r, const vector<int> &arr,
6.5
                    const int pos) {
66
       if (1 == r)
67
         return tree[pos] = Node(arr[1], 0);
68
69
       int mid = (1 + r) / 2;
70
       return tree[pos] = merge nodes(st build(1, mid, arr, 2 * pos + 1),
71
                                       st_build(mid + 1, r, arr, 2 * pos + 2));
72
73
74
     int st_get_first(const int 1, const int r, const int v, const int pos) {
75
       st_propagate(1, r, pos);
76
77
       // Needs RMO MAX
78
       // Replace to <= for greater or equal or (with RMQ MIN) > for smaller or
79
       // equal or >= for smaller
80
       if (tree[pos].val < v)</pre>
81
         return -1;
82
83
       if (1 == r)
84
         return 1;
85
86
       int mid = (l + r) / 2;
87
       int aux = st_get_first(1, mid, v, 2 * pos + 1);
88
       if (aux != -1)
89
         return aux:
90
       return st_get_first(mid + 1, r, v, 2 * pos + 2);
91
92
93
     Node st_query(const int 1, const int r, const int i, const int j,
94
                    const int pos) {
95
       st_propagate(1, r, pos);
96
97
       if (1 > r || 1 > j || r < i)
98
         return NEUTRAL_NODE;
99
```

```
if (i <= 1 && r <= j)
101
          return tree[pos];
102
103
        int mid = (1 + r) / 2;
104
        return merge_nodes(st_query(1, mid, i, j, 2 * pos + 1),
105
                            st_query(mid + 1, r, i, j, 2 * pos + 2));
106
107
108
      // it adds a number delta to the range from i to j
109
      Node st_update(const int 1, const int r, const int i, const int j,
                      const int delta, const int pos) {
110
111
        st_propagate(1, r, pos);
112
113
        if (1 > r | | 1 > j | | r < i)
114
          return tree[pos];
115
116
        if (i <= 1 && r <= j) {
117
          tree[pos].lazy = delta;
118
          st_propagate(1, r, pos);
119
          return tree[pos];
120
121
122
        int mid = (1 + r) / 2;
123
        return tree[pos] =
124
                    merge_nodes(st_update(1, mid, i, j, delta, 2 * pos + 1),
125
                                st\_update(mid + 1, r, i, j, delta, 2 * pos + 2));
126
127
128
      void build(const vector<int> &arr) {
129
        this->n = arr.size();
130
        this->tree.resize(4 * this->n);
131
        this->st_build(0, this->n - 1, arr, 0);
132
133
134
    public:
135
      /// N equals to -1 means the Segment Tree hasn't been created yet.
136
      Seg\_Tree() : n(-1) {}
137
138
      /// Constructor responsible initialize a tree with 0.
139
      /// Time Complexity O(n)
140
141
      Seq_Tree(const int n): n(n) { this->tree.resize(4 * this->n, Node(0, 0));
142
143
      /// Constructor responsible for building the initial tree based on a
        vector.
144
145
      /// Time Complexity O(n)
      Seq Tree(const vector<int> &arr) { this->build(arr); }
147
148
      /// Returns the first index from left to right.
149
      /// Uncomment the line in the original funtion to get the proper element
150
      /// may be: GREATER OR EQUAL, GREATER, SMALLER OR EQUAL, SMALLER.
151
      /// Time Complexity O(log n)
152
153
      int get first(const int v) {
154
        assert (this->n >= 0);
155
        return this->st_get_first(0, this->n - 1, v, 0);
156
157
158
      /// Update at a single index.
159
      /// Time Complexity O(log n)
160
      void update(const int idx, const int delta) {
```

```
assert (this->n >= 0);
163
        assert(0 <= idx), assert(idx < this->n);
164
        this->st_update(0, this->n - 1, idx, idx, delta, 0);
165
166
      /// Range update from 1 to r.
167
168
169
      /// Time Complexity O(log n)
170
      void update(const int l, const int r, const int delta) {
        assert (this->n >= 0);
171
172
         assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
173
        this->st_update(0, this->n - 1, 1, r, delta, 0);
174
175
176
      /// Query at a single index.
177
      /// Time Complexity O(log n)
178
179
      int query(const int idx) {
180
        assert (this->n >= 0);
181
        assert(0 <= idx), assert(idx < this->n);
182
        return this->st_query(0, this->n - 1, idx, idx, 0).val;
183
184
185
      /// Range query from 1 to r.
186
      /// Time Complexity O(log n)
187
188
      int query(const int 1, const int r) {
189
        assert (this->n >= 0);
190
        assert(0 \le 1), assert(1 \le r), assert(r < this->n);
191
         return this->st_query(0, this->n - 1, 1, r, 0).val;
192
193 | };
```

2.13. Segment Tree 2D

```
1 // REQUIRES segment_tree.cpp!!
2 class Seg_Tree_2d {
    private:
3
     // // range sum
     // int NEUTRAL_VALUE = 0;
     // int merge_nodes(const int &x, const int &y) {
6
     // return x + y;
7
8
     // }
10
     // // RMQ max
     // int NEUTRAL_VALUE = -INF;
11
12
     // int merge_nodes(const int &x, const int &y) {
13
     // return max(x, y);
     // }
14
15
     // // RMQ min
16
17
     // int NEUTRAL_VALUE = INF;
18
     // int merge_nodes(const int &x, const int &y) {
19
     // return min(x, y);
     // }
20
21
    private:
    int n, m;
24
25
    public:
26
     vector<Seq_Tree> tree;
27
28
    private:
```

```
void st_build(const int 1, const int r, const int pos, const
       vector<vector<int>> &mat) {
30
       if(1 == r)
31
         tree[pos] = Seq_Tree(mat[1]);
32
       else {
33
         int mid = (l + r) / 2;
34
         st_build(1, mid, 2*pos + 1, mat);
         st_build(mid + 1, r, 2*pos + 2, mat);
35
         for(int i = 0; i < tree[2*pos + 1].tree.size(); i++)</pre>
36
37
           tree[pos].tree[i].val = merge_nodes(tree[2*pos + 1].tree[i].val,
38
                                                 tree[2*pos + 2].tree[i].val);
39
40
41
     int st_query(const int 1, const int r, const int x1, const int y1, const
42
       int x2, const int y2, const int pos) {
43
       if(1 > x2 | | r < x1)
         return NEUTRAL VALUE;
44
45
46
       if(x1 \le 1 \&\& r \le x2)
47
         return tree[pos].query(y1, y2);
48
49
       int mid = (1 + r) / 2;
50
       return merge_nodes(st_query(1, mid, x1, y1, x2, y2, 2*pos + 1),
51
                           st_query(mid + 1, r, x1, y1, x2, y2, 2*pos + 2));
52
53
     void st_update(const int 1, const int r, const int x, const int y, const
54
       int delta, const int pos) {
55
       if(1 > x | | r < x)
56
         return:
57
58
       // Only supports point updates.
       if(1 == r) {
59
60
         tree[pos].update(y, delta);
61
         return:
62
63
64
       int mid = (1 + r) / 2;
65
       st_update(1, mid, x, y, delta, 2*pos + 1);
66
       st update(mid + 1, r, x, v, delta, 2*pos + 2);
67
       tree[pos].update(y, delta);
68
69
70
    public:
71
     Seq_Tree_2d() {
72
       this->n = -1;
73
       this->m = -1;
74
75
76
     Seq_Tree_2d(const int n, const int m) {
77
       this->n = n:
78
       this->m = m;
79
       // MAY TLE IN BUILD, TEST IT OR UPDATE EACH NODE MANUALLY!
80
       assert (m < 10000);
81
       tree.resize(4 * n, Seg_Tree(m));
82
83
84
     Seg Tree 2d(const int n, const int m, const vector<vector<int>> &mat) {
85
       this -> n = n;
86
       this->m = m;
87
       // MAY TLE IN BUILD, TEST IT OR UPDATE EACH NODE MANUALLY!
88
       assert (m < 10000);
       tree.resize(4 * n, Seg_Tree(m));
89
       st build(0, n - 1, 0, mat);
90
```

```
92
93
      // Query from (x1, y1) to (x2, y2).
94
95
      // Time complexity: O((log n) * (log m))
96
      int query(const int x1, const int y1, const int x2, const int y2) {
97
        assert (this->n > -1);
98
        assert (0 \leq x1); assert (x1 \leq x2); assert (x2 \leq this->n);
99
        assert(0 \le y1); assert(y1 \le y2); assert(y2 < this -> n);
100
        return st_query(0, this->n - 1, x1, y1, x2, y2, 0);
101
102
103
      // Point updates on position (x, y).
104
      // Time complexity: O((log n) * (log m))
105
      void update(const int x, const int v, const int delta) {
106
107
        assert (0 <= x); assert (x < this->n);
108
        assert(0 <= y); assert(y < this->n);
109
        st_update(0, this->n - 1, x, y, delta, 0);
110
111 };
```

2.14. Segment Tree Polynomial

```
/// Works for the polynomial f(x) = z1*x + z0
   class Seg_Tree {
   public:
3
     struct Node {
       int val, z1, z0;
6
7
       Node() {}
8
       Node (const int val, const int z1, const int z0)
9
           : val(val), z1(z1), z0(z0) {}
10
11
12 private:
13
     // range sum
     Node NEUTRAL_NODE = Node (0, 0, 0);
14
15
     Node merge_nodes(const Node &x, const Node &y) {
16
       return Node(x.val + y.val, 0, 0);
17
18
     void apply_lazy(const int 1, const int r, const int pos) {
19
       tree[pos].val += (r - l + 1) * tree[pos].z0;
       tree[pos].val += (r - 1) * (r - 1 + 1) / 2 * tree[pos].z1;
20
21
22
23 private:
24
     int n;
25
26 public:
2.7
     vector<Node> tree;
28
     void st_propagate(const int 1, const int r, const int pos) {
31
       if (tree[pos].z0 != 0 || tree[pos].z1 != 0) {
32
         apply_lazy(l, r, pos);
33
         int mid = (1 + r) / 2;
34
         int sz left = mid - 1 + 1;
35
         if (1 != r) {
36
           tree[2 * pos + 1].z0 += tree[pos].z0;
37
           tree [2 * pos + 1].z1 += tree[pos].z1;
38
39
           tree[2 * pos + 2].z0 += tree[pos].z0 + sz_left * tree[pos].z1;
           tree[2 * pos + 2].z1 += tree[pos].z1;
40
```

```
42
          tree[pos].z0 = 0;
 43
          tree[posl.z1 = 0:
 44
 45
 46
 47
      Node st build(const int 1, const int r, const vector<int> &arr,
 48
                    const int pos) {
        if (1 == r)
 49
 50
          return tree[pos] = Node(arr[1], 0, 0);
 51
 52
        int mid = (1 + r) / 2;
        return tree[pos] = merge_nodes(st_build(1, mid, arr, 2 * pos + 1),
 53
                                         st_build(mid + 1, r, arr, 2 * pos + 2));
 54
 55
 56
 57
      Node st_query(const int 1, const int r, const int i, const int j,
 58
                     const int pos) {
 59
        st_propagate(l, r, pos);
 60
 61
        if (1 > r || 1 > j || r < i)
          return NEUTRAL NODE;
 62
 63
        if (i <= 1 && r <= i)
 64
 65
          return tree[pos];
 66
 67
        int mid = (1 + r) / 2;
 68
        return merge_nodes(st_query(l, mid, i, j, 2 * pos + 1),
                            st_query(mid + 1, r, i, j, 2 * pos + 2));
 69
 70
 71
 72
      // it adds a number delta to the range from i to j
 73
      Node st_update(const int 1, const int r, const int i, const int j,
                      const int z1, const int z0, const int pos) {
 74
 75
        st propagate(1, r, pos);
 76
 77
        if (1 > r || 1 > j || r < i)
          return tree[pos];
 78
 79
 80
        if (i <= 1 && r <= j) {
          tree[pos].z0 = (1 - i + 1) * z0;
 81
 82
          tree[pos].z1 = z1;
 83
          st_propagate(1, r, pos);
 84
          return tree[pos];
 85
 86
        int mid = (1 + r) / 2;
 87
        return tree[pos] =
 89
                    merge_nodes(st_update(1, mid, i, j, z1, z0, 2 * pos + 1),
 90
                                st\_update(mid + 1, r, i, j, z1, z0, 2 * pos + 2));
 91
 92
 93
    public:
      Seq\_Tree() : n(-1) {}
 95
 96
      Seg_Tree(const int n) : n(n) { this->tree.resize(4 * this->n, Node(0, 0));
 97
 98
      Seq Tree(const vector<int> &arr) { this->build(arr); }
 99
100
      void build(const vector<int> &arr) {
101
        this->n = arr.size();
102
        this->tree.resize(4 * this->n);
        this->st_build(0, this->n - 1, arr, 0);
103
104
```

```
106
      /// Index update of a polynomial f(x) = z1*x + z0
107
108
      /// Time Complexity O(log n)
109
      void update(const int i, const int z1, const int z0) {
110
        assert (this->n >= 0);
111
        assert(0 <= i), assert(i < this->n);
        this->st_update(0, this->n - 1, i, i, z1, z0, 0);
112
113
114
      /// Range update of a polynomial f(x) = z1 \star x + z0 from 1 to r
115
116
      /// Time Complexity O(log n)
117
      void update(const int 1, const int r, const int z1, const int z0) {
118
        assert (this->n >= 0);
119
        assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
120
121
        this->st_update(0, this->n - 1, 1, r, z1, z0, 0);
122
123
124
      /// Range sum query from 1 to r
125
126
      /// Time Complexity O(log n)
127
      int querv(const int 1, const int r) {
128
        assert(this->n >= 0);
129
        assert (0 \le 1), assert (1 \le r), assert (r < this -> n);
        return this->st_query(0, this->n - 1, 1, r, 0).val;
130
131
132 };
```

2.15. Sparse Table

```
1 // RMO min implementation
   class Sparse Table {
    private:
     int n;
5
     vector<vector<int> > table;
     vector<int> lq;
8
     /// lg[i] represents the log2(i)
9
     void build_log_array() {
       lg.resize(this->n + 1);
10
11
        for(int i = 2; i <= this->n; i++)
12
13
         lq[i] = lq[i/2] + 1;
14
15
      /// Time Complexity: O(n*log(n))
16
      /// Space Complexity: O(n*log(n))
17
     void build_sparse_table(const vector<int> &arr) {
18
19
20
       table.resize(lg[this->n] + 1, vector<int>(this->n));
21
22
       table[0] = arr;
       int pow2 = 1:
23
        for(int i = 1; i < table.size(); i++) {</pre>
         int lastsz = this->n - pow2 + 1;
26
27
          for(int j = 0; j + pow2 < lastsz; <math>j++) {
28
           table[i][j] = merge(table[i-1][j], table[i-1][j+pow2]);
29
30
         pow2 <<= 1;
31
32
33
```

```
35
     int merge(const int &1, const int &r) {
36
       return min(l, r);
37
38
39
    public:
40
     Sparse_Table(const vector<int> &arr) {
       this->n = arr.size();
41
42
43
       this->build_log_array();
       this->build_sparse_table(arr);
44
45
46
47
     void print() {
48
       int pow2 = 1;
       for(int i = 0; i < table.size(); i++) {</pre>
49
50
         int sz = (int) (table.front().size()) - pow2 + 1;
51
          for (int j = 0; j < sz; j++) {
52
            cout << table[i][j] << " \n"[(j+1) == sz];</pre>
53
          pow2 <<= 1;
54
55
56
57
58
     /// Query of a range from 1 to r.
59
60
     /// Time Complexity: O(1)
     /// Space Complexity: O(1)
61
62
     int query(int 1, int r) {
       assert(1 <= r);
63
       assert(0 <= 1 && r <= this->n - 1);
64
65
66
       int lgg = lg[(r - l + 1)];
67
       return merge(table[lgg][1], table[lgg][r - (1 << lgg) + 1]);</pre>
68
69
   };
```

3. Dp

3.1. Achar Major Palindromo

Fazer LCS da string com o reverso

3.2. Catalan

```
// The first few Catalan numbers for n = 0, 1, 2, 3, \ldots
   // are 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, ...
   // Formula Recursiva:
   // cat(0) = 0
   // cat(n+1) = somatorio(i from 0 to n)(cat(i)*cat(n-i))
   // Using Binomial Coefficient
   // We can also use the below formula to find nth catalan number in O(n) time.
   // Formula acima
11 // Returns value of Binomial Coefficient C(n, k)
13 | int binomialCoeff(int n, int k) {
14
    int res = 1;
15
16
     // Since C(n, k) = C(n, n-k)
    if (k > n - k)
17
     k = n - k;
```

```
20
     // Calculate value of [n*(n-1)*--*(n-k+1)] / [k*(k-1)*--*1]
21
     for (int i = 0; i < k; ++i) {</pre>
22
         res \star= (n - i);
23
          res /= (i + 1);
24
25
26
       return res;
27
28 // A Binomial coefficient based function to find nth catalan
29 // number in O(n) time
30 | int catalan(int n)
31
        // Calculate value of 2nCn
32
        int c = binomialCoeff(2*n, n);
33
34
        // return 2nCn/(n+1)
35
       return c/(n+1);
36 }
```

3.3. Digit Dp

```
1 /// How many numbers x are there in the range a to b, where the digit d
        occurs exactly k times in x?
2 vector<int> num;
3 int a, b, d, k;
4 int DP[12][12][2];
5 /// DP[p][c][f] = Number of valid numbers <= b from this state
6 /// p = current position from left side (zero based)
7 | /// c = number of times we have placed the digit d so far
8 / / / f = the number we are building has already become smaller than b? [0 =
       no, 1 = yes
10 | int call(int pos, int cnt, int f) {
     if(cnt > k) return 0;
12
13
     if(pos == num.size()){
14
       if(cnt == k) return 1;
15
       return 0;
16
17
18
     if(DP[pos][cnt][f] != -1) return DP[pos][cnt][f];
19
     int res = 0;
     int lim = (f ? 9 : num[pos]);
20
21
22
     /// Try to place all the valid digits such that the number doesn't exceed b
     for (int dgt = 0; dgt<=LMT; dgt++) {
23
24
       int nf = f;
25
       int ncnt = cnt;
26
       if (f == 0 && dgt < LMT) nf = 1; /// The number is getting smaller at
        this position
27
       if(dqt == d) ncnt++;
28
       if (ncnt <= k) res += call(pos+1, ncnt, nf);
2.9
30
31
     return DP[pos][cnt][f] = res;
32 }
33
34 int solve(int b) {
35
     num.clear();
36
     while(b>0){
37
       num.push_back(b%10);
38
       b/=10;
39
40
     reverse(num.begin(), num.end());
```

```
/// Stored all the digits of b in num for simplicity
42
43
     memset (DP, -1, sizeof (DP));
44
     int res = call(0, 0, 0);
45
     return res;
46
47
48
   int main () {
49
50
     cin >> a >> b >> d >> k;
     int res = solve(b) - solve(a-1);
51
52
     cout << res << endl;
53
54
     return 0;
55
```

3.4. Longest Common Subsequence

```
string lcs(string &s, string &t) {
3
     int n = s.size(), m = t.size();
4
     s.insert(s.begin(), '#');
     t.insert(t.begin(), '$');
7
8
     vector<vector<int>> mat(n + 1, vector<int>(m + 1, 0));
9
10
     for(int i = 1; i <= n; i++) {</pre>
       for (int j = 1; j <= m; j++) {</pre>
11
12
         if(s[i] == t[j])
13
           mat[i][j] = mat[i - 1][j - 1] + 1;
14
15
           mat[i][j] = max(mat[i - 1][j], mat[i][j - 1]);
16
17
18
19
     string ans;
20
     int i = n, j = m;
21
     while (i > 0 \&\& j > 0) {
22
       if(s[i] == t[j])
23
         ans += s[i], i--, j--;
24
        else if (mat[i][j - 1] > mat[i - 1][j])
25
        j--;
26
       else
27
         i--;
28
29
     reverse(ans.begin(), ans.end());
30
31
     return ans;
```

3.5. Longest Common Substring

```
int LCSubStr(char *X, char *Y, int m, int n) {
    // Create a table to store lengths of longest common suffixes of
    // substrings. Notethat LCSuff[i][j] contains length of longest
    // common suffix of X[0..i-1] and Y[0..j-1]. The first row and
    // first column entries have no logical meaning, they are used only
    // for simplicity of program
    int LCSuff[m+1][n+1];
    int result = 0; // To store length of the longest common substring
    /* Following steps build LCSuff[m+1][n+1] in bottom up fashion. */
```

```
for (int i=0; i<=m; i++) {</pre>
12
       for (int j=0; j<=n; j++) {
13
         if (i == 0 || j == 0)
14
           LCSuff[i][j] = 0;
15
16
          else if (X[i-1] == Y[j-1]) {
17
           LCSuff[i][j] = LCSuff[i-1][j-1] + 1;
18
            result = max(result, LCSuff[i][j]);
19
20
         else LCSuff[i][j] = 0;
21
22
23
     return result:
24
```

3.6. Longest Increasing Subsequence 2D (Not Sorted)

```
set<ii>> s[(int)2e6];
   bool check (ii par, int ind) {
3
     auto it = s[ind].lower_bound(ii(par.ff, -INF));
5
     if(it == s[ind].begin())
       return false;
7
8
     it--;
9
     if(it->ss < par.ss)</pre>
10
       return true;
11
12
     return false;
13
14
   int lis2d(vector<ii> &arr) {
15
16
     int n = arr.size();
18
     s[1].insert(arr[0]);
19
20
     int maior = 1;
21
     for(int i = 1; i < n; i++) {</pre>
22
23
       ii x = arr[i];
24
25
       int 1 = 1, r = maior;
26
        int ansbb = 0:
27
        while(1 <= r) {
28
          int mid = (1+r)/2;
29
          if(check(x, mid)) {
30
           l = mid + 1;
31
            ansbb = mid;
32
          } else {
33
            r = mid - 1;
34
35
36
37
        // inserting in list
38
        auto it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
39
        while (it != s[ansbb+1].end() && it->ss >= x.ss)
40
          it = s[ansbb+1].erase(it);
41
        it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
42
       if(s[ansbb+1].size() > 0 && it != s[ansbb+1].end() && it->ff == x.ff &&
43
        it->ss <= x.ss)
44
          continue;
45
        s[ansbb+1].insert(arr[i]);
46
```

3.7. Longest Increasing Subsequence 2D (Sorted)

```
set<ii>> s[(int)2e6];
   bool check(ii par, int ind) {
     auto it = s[ind].lower_bound(ii(par.ff, -INF));
     if(it == s[ind].begin())
       return false;
     it--;
10
     if(it->ss < par.ss)</pre>
11
       return true;
12
     return false;
13
14
   int lis2d(vector<ii> &arr) {
17
     int n = arr.size();
18
     s[1].insert(arr[0]);
19
20
     int maior = 1;
21
     for(int i = 1; i < n; i++) {</pre>
22
23
       ii x = arr[i];
24
25
       int 1 = 1, r = maior;
26
       int ansbb = 0;
27
       while(1 <= r) {
28
         int mid = (1+r)/2;
29
         if(check(x, mid)) {
30
           l = mid + 1;
31
           ansbb = mid;
32
         } else {
33
           r = mid - 1;
34
35
36
37
       // inserting in list
        auto it = s[ansbb+1].lower bound(ii(x.ff, -INF));
38
39
        while(it != s[ansbb+1].end() && it->ss >= x.ss)
40
         it = s[ansbb+1].erase(it);
41
42
        it = s[ansbb+1].lower_bound(ii(x.ff, -INF));
       if(s[ansbb+1].size() > 0 && it != s[ansbb+1].end() && it->ff == x.ff &&
43
        it->ss <= x.ss)
44
45
       s[ansbb+1].insert(arr[i]);
46
47
       maior = max(maior, ansbb + 1);
48
49
50
     return maior:
51
52
```

3.8. Longest Increasing Subsequence

```
1 int lis(vector<int> &arr){
     int n = arr.size();
     vector<int> lis:
     for(int i = 0; i < n; i++) {</pre>
       int l = 0, r = (int) lis.size() - 1;
        int ans i = -1;
        while(1 <= r) {
          int mid = (1+r)/2;
          // OBS: PARA >= TROCAR SINAL EMBAIXO POR <=
9
          if(arr[i] < lis[mid]) {
10
11
           r = mid - 1;
12
            ansj = mid;
1.3
14
          else 1 = mid + 1:
15
16
        if (ansi == -1) {
          // se arr[i] e maior que todos
17
18
          lis.push back(arr[i]);
19
20
        else {
21
          lis[ansj] = arr[i];
22
23
24
25
     return lis.size();
```

3.9. Subset Sum Com Bitset

```
bitset<312345> bit;
int arr[112345];

void subsetSum(int n) {
   bit.reset();
   bit.set(0);
   for(int i = 0; i < n; i++) {
   bit | (bit << arr[i]);
   }
}</pre>
```

3.10. Coin Change Problem

```
1 // função que recebe o valor de troco N. o número de moedas disponíveis M.
   // e um vetor com as moedas disponíveis arr
   // essa função deve retornar o número mínimo de moedas,
   // de acordo com a solução com Programação Dinamica.
   int num_moedas(int N, int M, int arr[]) {
     int dp[N+1];
     // caso base
     dp[0] = 0;
     // sub-problemas
     for(int i=1; i<=N; i++) {</pre>
10
11
     // é comum atribuir um valor alto, que concerteza
     // é maior que qualquer uma das próximas possibilidades,
12
     // sendo assim substituido
13
     dp[i] = 1000000;
     for(int j=0; j<M; j++) {</pre>
15
16
       if(i-arr[j] >= 0) {
17
         dp[i] = min(dp[i], dp[i-arr[j]]+1);
18
19
20
```

```
// solução
22
     return dp[N];
23
```

8

9

10

11

1.3

15

17

18

19

20

22

23

24

25

26

27

28

29

31

32

33

34

35

36

37

38

39

40 41

42 43

44

45

46

47

48

51

54

55

56

57

58

59

60

61

62

63

64

65

66

67

68

 \star xi = x2 + rx;

3.11. Knapsack

```
int dp[2001][2001];
   int moc(int q, int p, vector<ii> vec) {
     for(int i = 1; i <= q; i++)
4
       for (int j = 1; j \le p; j++) {
5
         if(j \ge vec[i-1].ff)
           dp[i][j] = max(dp[i-1][j], vec[i-1].ss + dp[i-1][j-vec[i-1].ff]);
          else
            dp[i][j] = dp[i-1][j];
10
11
12
     return dp[q][p];
13
   int main(int argc, char *argv[])
14
15
16
     int p,q;
17
     vector<ii> vec:
     cin >> p >> q;
18
19
     int x,y;
     for(int i = 0; i < q; i++) {
20
21
          cin >> x >> y;
22
          vec.push_back(make_pair(x,y));
23
2.4
     for(int i = 0; i <= p; i++)
25
       dp[0][i] = 0;
26
     for(int i = 1; i <= q; i++)
27
       dp[i][0] = 0;
28
     sort(vec.begin(), vec.end());
29
     cout << moc(q,p,vec) << endl;</pre>
30
```

4. Geometry

4.1. Centro De Massa De Um Poligono

```
double area = 0;
2
   pto c:
3
4
   c.x = c.y = 0;
   for (int i = 0; i < n; i++) {
     double aux = (arr[i].x * arr[i+1].y) - (arr[i].y * arr[i+1].x); // shoelace
     area += aux;
8
     c.x += aux*(arr[i].x + arr[i+1].x);
     c.y += aux*(arr[i].y + arr[i+1].y);
10
11
12
   c.x /= (3.0*area);
   c.y /= (3.0*area);
13
14
15 cout << c.x << ' ' << c.y << endl;
```

4.2. Circle-Circle Intersection

```
/* circle_circle_intersection() *
2
   * Determine the points where 2 circles in a common plane intersect.
3
```

```
* int circle_circle_intersection(
5
                                      // center and radius of 1st circle
                                      double x0, double y0, double r0,
6
7
                                      // center and radius of 2nd circle
                                      double x1, double v1, double r1,
                                      // 1st intersection point
                                      double *xi, double *yi,
                                      // 2nd intersection point
                                      double *xi_prime, double *yi_prime)
12
    * This is a public domain work. 3/26/2005 Tim Voght
14
16
    */
   int circle_circle_intersection(double x0, double y0, double r0, double x1,
                                  double v1, double r1, double *xi, double *vi,
                                   double *xi_prime, double *yi_prime) {
21
     double a, dx, dy, d, h, rx, ry;
     double x2, y2;
     /* dx and dv are the vertical and horizontal distances between
      * the circle centers.
      */
     dx = x1 - x0;
     dy = y1 - y0;
30
     /* Determine the straight-line distance between the centers. */
     // d = sqrt((dy*dy) + (dx*dx));
     d = hypot(dx, dy); // Suggested by Keith Briggs
     /★ Check for solvability. ★/
     if (d > (r0 + r1)) {
       /* no solution. circles do not intersect. */
       return 0;
     if (d < fabs(r0 - r1)) {
       /★ no solution. one circle is contained in the other ★/
       return 0:
     /\star 'point 2' is the point where the line through the circle
      * intersection points crosses the line between the circle
      * centers.
      */
     /* Determine the distance from point 0 to point 2. */
     a = ((r0 * r0) - (r1 * r1) + (d * d)) / (2.0 * d);
     /* Determine the coordinates of point 2. */
52
53
     x2 = x0 + (dx * a / d);
     y2 = y0 + (dy * a / d);
     /★ Determine the distance from point 2 to either of the
      * intersection points.
     h = sqrt((r0 * r0) - (a * a));
     /★ Now determine the offsets of the intersection points from
      * point 2.
      */
     rx = -dy * (h / d);
     ry = dx * (h / d);
     /* Determine the absolute intersection points. ★/
```

4.3. Closest Pair Of Points

```
1 | struct Point {
2
    int x, y;
3
   int compareX(const void *a,const void *b) {
    Point *p1 = (Point *)a, *p2 = (Point *)b;
     return (p1->x - p2->x);
7
8 int compareY(const void *a, const void *b) {
9
    Point *p1 = (Point *)a, *p2 = (Point *)b;
10
    return (p1->y - p2->y);
11
12 float dist(Point p1, Point p2) {
13
     return sqrt ((p1.x - p2.x) * (p1.x - p2.x) + (p1.y - p2.y) * (p1.y - p2.y));
14
15 | float bruteForce(Point P[], int n) {
16
     float min = FLT_MAX;
     for (int i = 0; i < n; ++i)
17
18
       for (int j = i+1; j < n; ++j)
         if (dist(P[i], P[j]) < min)</pre>
19
20
           min = dist(P[i], P[j]);
21
     return min;
22
23 | float min(float x, float y) {
    return (x < y)? x : y;
25
26 | float stripClosest(Point strip[], int size, float d) {
27
     float min = d;
28
     for (int i = 0; i < size; ++i)
29
       for (int j = i+1; j < size && (strip[j].y - strip[i].y) < min; ++j)
30
         if (dist(strip[i],strip[j]) < min)</pre>
31
           min = dist(strip[i], strip[j]);
32
    return min;
33
   float closestUtil(Point Px[], Point Py[], int n) {
34
35
    if (n <= 3)
36
       return bruteForce(Px, n);
     int mid = n/2;
37
     Point midPoint = Px[mid];
38
39
     Point Pyl[mid+1];
40
     Point Pyr[n-mid-1];
41
     int 1i = 0, ri = 0;
42
     for (int i = 0; i < n; i++)
43
       if (Py[i].x <= midPoint.x)</pre>
44
         Pyl[li++] = Py[i];
45
46
         Pyr[ri++] = Py[i];
47
     float dl = closestUtil(Px, Pyl, mid);
     float dr = closestUtil(Px + mid, Pyr, n-mid);
     float d = min(dl, dr);
51
     Point strip[n];
52
     int j = 0;
53
     for (int i = 0; i < n; i++)
54
       if (abs(Py[i].x - midPoint.x) < d)</pre>
         strip[j] = Py[i], j++;
```

```
return min(d, stripClosest(strip, j, d));
57 }
58
59 | float closest (Point P[], int n) {
60 | Point Px[n];
     Point Py[n];
61
62
     for (int i = 0; i < n; i++) {
       Px[i] = P[i];
63
64
       Py[i] = P[i];
6.5
     qsort(Px, n, sizeof(Point), compareX);
66
     gsort(Py, n, sizeof(Point), compareY);
     return closestUtil(Px, Pv, n);
68
69 }
```

4.4. Condicao De Existencia De Um Triangulo

```
1
2
| | b - c | < a < b + c
3 | a - c | < b < a + c
4 | a - b | < c < a + b
5
6
Para a < b < c, basta checar
7 a + b > c

OBS: Para um conjunto n >= 100 sempre exite um triângulo válido, pois a sequência de triângulos não válidos seguem a sequência de Fibonacci e Fib(100) > 2^64
```

4.5. Convex Hull

```
1 // Asymptotic complexity: O(n log n).
2 | struct pto {
     double x, y;
3
     bool operator <(const pto &p) const {</pre>
       return x < p.x | | (x == p.x && y < p.y);
6
       /* a impressao será em prioridade por mais a esquerda, mais
7
           abaixo, e antihorário pelo cross abaixo */
8
9
   };
1.0
   double cross(const pto &O, const pto &A, const pto &B) {
12
     return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
13
14
   vector<pto> convex hull(vector<pto> P) {
16
     int n = P.size(), k = 0;
17
     vector<pto> H(2 * n);
18
     // Sort points lexicographically
19
     sort(P.begin(), P.end());
20
      // Build lower hull
21
     for (int i = 0; i < n; ++i) {</pre>
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
2.3
        // trocar por >= 0
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
24
25
26
       H[k++] = P[i];
27
28
     // Build upper hull
29
      for (int i = n - 2, t = k + 1; i >= 0; i--) {
30
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
31
        // trocar por >= 0
       while (k \ge t \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
32
```

```
33 | k--;
34 | H[k++] = P[i];
35 | H.resize(k);
37 /* o último ponto do vetor é igual ao primeiro, atente para isso
38 as vezes é necessário mudar */
39 | return H;
40 |}
```

4.6. Cross Product

```
1 // Outra forma de produto vetorial
   // reta ab, ac se for zero e colinear
   // se for < 0 entao antiHorario, > 0 horario
  bool ehcol(pto a,pto b,pto c) {
5
    return ((b.y-a.y) * (c.x-a.x) - (b.x-a.x) * (c.y-a.y));
   //Produto vetorial AB x AC, se for zero e colinear
9 int cross(pto A, pto B, pto C) {
   pto AB, AC;
1.0
    AB.x = B.x-A.x;
11
    AB.y = B.y-A.y;
12
    AC.x = C.x-A.x;
13
14
    AC.y = C.y-A.y;
    int cross = AB.x*AC.y-AB.y * AC.x;
    return cross:
17
18
19 // OBS: DEFINE ÁREA DE QUADRILÁTERO FORMADO PELAS RETAS, A ÁREA DO TRIÂNGULO
       É A METADE
```

4.7. Distance Point Segment

```
// use struct point and line
double dist_point_segment(const Point p, const Point s, const Point t) {
   if(sgn(dot(p-s, t-s)) < 0)
      return (p-s).norm();
   if(sgn(dot(p-t, s-t)) < 0)
      return (p-t).norm();
   return abs(det(s-p, t-p) / dist(s, t));
}</pre>
```

4.8. Line-Line Intersection

```
1 // Intersecção de retas Ax + By = C dados pontos (x1,y1) e (x2,y2)
2 | A = y2-y1
3 \mid B = x1-x2
4 \mid C = A \star x 1 + B \star y 1
5 //Retas definidas pelas equações:
6 \mid A1x + B1y = C1
7 A2x + B2v = C2
8 //Encontrar x e y resolvendo o sistema
9 double det = A1*B2 - A2*B1;
10 | if (det == 0) {
    //Lines are parallel
11
12 | }else{
13
    double x = (B2*C1 - B1*C2)/det:
14
    double y = (A1*C2 - A2*C1)/det;
15
```

4.9. Line-Point Distance

```
1 double ptoReta(double x1, double y1, double x2, double y2, double pointX,
        double pointY, double *ptox, double *ptoy) {
     double diffX = x2 - x1;
     double diffY = y2 - y1;
     if ((diffX == 0) && (diffY == 0)) {
       diffX = pointX - x1;
       diffY = pointY - y1;
       //se os dois sao pontos
8
       return hypot (pointX - x1, pointY - y1);
9
10
     double t = ((pointX - x1) * diffX + (pointY - y1) * diffY) /
11
                      (diffX * diffX + diffY * diffY);
12
     if (t < 0) {
13
       //point is nearest to the first point i.e x1 and y1
14
        // Ex:
15
        // cord do pto na reta = pto inicial(x1,y1);
16
        \starptox = x1, \starptoy = y1;
        diffX = pointX - x1;
17
        diffY = pointY - y1;
18
       else if (t > 1) {
19
20
        //point is nearest to the end point i.e x2 and y2
21
        // Ex : .
22
        // cord do pto na reta = pto final(x2,y2);
23
        *ptox = x2, *ptoy = y2;
24
        diffX = pointX - x2;
25
        diffY = pointY - y2;
       else {
26
27
          //if perpendicular line intersect the line segment.
28
          // pto nao esta mais proximo de uma das bordas do segmento
29
          // Ex:
30
          //
31
          //
                               |(Ângulo Reto)
32
          // cord x do pto na reta = (x1 + t * diffX)
33
34
         // cord y do pto na reta = (y1 + t * diffY)
35
        \star ptox = (x1 + t \star diffX), \star ptoy = (y1 + t \star diffY);
36
        diffX = pointX - (x1 + t * diffX);
37
       diffY = pointY - (y1 + t * diffY);
3.8
39
     //returning shortest distance
40
     return sqrt(diffX * diffX + diffY * diffY);
```

4.10. Point Inside Convex Polygon - Log(N)

```
#include <bits/stdc++.h>
3 using namespace std;
5 #define INF 1e18
6 #define pb push back
7 #define ii pair<int,int>
8 #define OK cout<<"OK"<<endl
9 | #define debug(x) cout << #x " = " << (x) << endl
10 #define ff first
11 #define ss second
12 #define int long long
13
14 struct pto {
15 double x, y;
16
    bool operator <(const pto &p) const {</pre>
17
       return x < p.x || (x == p.x && y < p.y);
```

```
/★ a impressao será em prioridade por mais a esquerda, mais
19
          abaixo, e antihorário pelo cross abaixo */
20
21
22
   double cross(const pto &O, const pto &A, const pto &B) {
23
    return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
24
2.5
26
   vector<pto> lower, upper;
2.7
   vector<pto> convex hull(vector<pto> &P) {
29
     int n = P.size(), k = 0;
30
     vector<pto> H(2 * n);
31
     // Sort points lexicographically
32
     sort(P.begin(), P.end());
33
     // Build lower hull
     for (int i = 0; i < n; ++i) {
34
35
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
36
       // trocar por >= 0
37
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
38
         k--:
39
       H[k++] = P[i];
40
     // Build upper hull
41
42
     for (int i = n - 2, t = k + 1; i >= 0; i--) {
43
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
44
       // trocar por >= 0
45
       while (k \ge 1 \text{ & & cross}(H[k - 2], H[k - 1], P[i]) \le 0)
46
        k--;
47
       H[k++] = P[i];
48
49
     H.resize(k):
50
     /★ o último ponto do vetor é igual ao primeiro, atente para isso
51
     as vezes é necessário mudar */
52
53
     int j = 1;
54
     lower.pb(H.front());
55
     while (H[j].x >= H[j-1].x) {
56
      lower.pb(H[j++]);
57
58
59
     int 1 = H.size()-1:
     while(1 >= i) {
60
61
       upper.pb(H[1--]);
62
63
     upper.pb(H[1--]);
64
65
     return H;
66
67
68
   bool insidePolygon(pto p, vector<pto> &arr) {
69
     if (pair<double, double> (p.x, p.y) == pair<double, double> (lower[0].x,
70
       lower[0].v))
71
       return true;
72
73
     pto lo = {p.x, -(double)INF};
74
     pto hi = {p.x, (double) INF};
75
     auto itl = lower_bound(lower.begin(), lower.end(), lo);
76
     auto itu = lower_bound(upper.begin(), upper.end(), lo);
77
78
     if(itl == lower.begin() || itu == upper.begin()) {
       auto it = lower bound(arr.begin(), arr.end(), lo);
79
80
       auto it2 = lower_bound(arr.begin(), arr.end(), hi);
81
       it2--:
```

```
if(it2 >= it && p.x == it-> x && it->x == it2->x && it->y <= p.y && p.y
        \leq it2->v)
          return true;
83
84
        return false:
85
      if(itl == lower.end() || itu == upper.end()) {
86
87
        return false:
8.8
89
90
      auto ol = itl, ou = itu;
91
      ol--, ou--;
92
      if(cross(*ol, *itl, p) >= 0 && cross(*ou, *itu, p) <= 0)</pre>
93
        return true;
94
95
      auto it = lower bound(arr.begin(), arr.end(), lo);
      auto it2 = lower_bound(arr.begin(), arr.end(), hi);
96
97
      if(it2 >= it && p.x == it-> x && it->x == it2->x && it->y <= p.y && p.y <=
98
        it2->v)
99
        return true;
100
      return false:
101
102
103
104
    signed main () {
106
      ios_base::sync_with_stdio(false);
107
108
      cin.tie(NULL);
1 / 9
110
      double n, m, k;
111
112
      cin >> n >> m >> k;
113
114
      vector<pto> arr(n);
115
116
      for(pto &x: arr) {
117
        cin >> x.x >> x.y;
118
119
      convex_hull(arr);
120
121
122
      pto p;
123
124
      int c = 0;
125
      while (m--)
126
        cin >> p.x >> p.y;
127
        cout << (insidePolygon(p, arr) ? "dentro" : "fora") << endl;</pre>
128
129
130
```

4.11. Point Inside Polygon

```
/* Traça-se uma reta do ponto até um outro ponto qualquer fora do triangulo e checa o número de interseção com a borda do polígono se este for ímpar então está dentro se não está fora */

// Define Infinite (Using INT_MAX caused overflow problems)

#define INF 10000

struct pto {
   int x, y;
```

```
10
       pto(int x, int y) : x(x), y(y) {}
11
12
13 // Given three colinear ptos p, q, r, the function checks if
14 // pto q lies on line segment 'pr'
15 bool on Segment (pto p, pto q, pto r) {
    if (q.x \le max(p.x, r.x) \& q.x >= min(p.x, r.x) \& \&
16
17
         q.y \le max(p.y, r.y) && q.y >= min(p.y, r.y)
18
       return true:
19
     return false;
20
21
   // To find orientation of ordered triplet (p, q, r).
23
   // The function returns following values
   // 0 --> p, q and r are colinear
   // 1 --> Clockwise
   // 2 --> Counterclockwise
27
  int orientation(pto p, pto q, pto r) {
    int val = (q.y - p.y) * (r.x - q.x) -
               (q.x - p.x) * (r.y - q.y);
29
30
    if (val == 0) return 0; // colinear
    return (val > 0)? 1: 2; // clock or counterclock wise
33
34
35
   // The function that returns true if line segment 'plq1'
36 // and 'p2q2' intersect.
37 | bool doIntersect(pto p1, pto q1, pto p2, pto q2) {
     // Find the four orientations needed for general and
39
     // special cases
     int o1 = orientation(p1, q1, p2);
40
     int o2 = orientation(p1, q1, q2);
41
     int o3 = orientation(p2, q2, p1);
43
     int o4 = orientation(p2, q2, q1);
44
45
     // General case
     if (o1 != o2 && o3 != o4)
46
47
       return true;
48
49
     // Special Cases
     // p1, q1 and p2 are colinear and p2 lies on segment p1q1
50
     if (o1 == 0 && onSegment(p1, p2, q1)) return true;
51
53
     // p1, q1 and p2 are colinear and q2 lies on segment p1q1
     if (o2 == 0 && onSegment(p1, q2, q1)) return true;
56
     // p2, q2 and p1 are colinear and p1 lies on segment p2q2
57
     if (o3 == 0 && onSegment(p2, p1, q2)) return true;
58
59
      // p2, q2 and q1 are colinear and q1 lies on segment p2q2
60
     if (o4 == 0 && onSegment(p2, q1, q2)) return true;
61
62
     return false; // Doesn't fall in any of the above cases
63
64
   // Returns true if the pto p lies inside the polygon[] with n vertices
   bool isInside(pto polygon[], int n, pto p) {
     // There must be at least 3 vertices in polygon[]
68
     if (n < 3) return false;
69
70
     // Create a pto for line segment from p to infinite
71
     pto extreme = pto(INF, p.y);
72
     // Count intersections of the above line with sides of polygon
```

```
int count = 0, i = 0;
75
76
       int next = (i+1)%n;
77
78
       // Check if the line segment from 'p' to 'extreme' intersects
79
       // with the line segment from 'polygon[i]' to 'polygon[next]'
80
       if (doIntersect(polygon[i], polygon[next], p, extreme)) {
         // If the pto 'p' is colinear with line segment 'i-next',
81
         // then check if it lies on segment. If it lies, return true,
82
         // otherwise false
8.3
84
         if (orientation(polygon[i], p, polygon[next]) == 0)
85
           return onSegment(polygon[i], p, polygon[next]);
86
87
         count++;
88
89
       i = next;
    } while (i != 0);
90
91
     // Return true if count is odd, false otherwise
    return count&1; // Same as (count%2 == 1)
```

4.12. Points Inside And In Boundary Polygon

```
1 int cross(pto a, pto b) {
    return a.x * b.y - b.x * a.y;
3
5 int boundaryCount(pto a, pto b) {
    if(a.x == b.x)
       return abs(a.y-b.y)-1;
8
     if(a.y == b.y)
       return abs(a.x-b.x)-1;
10
     return _gcd(abs(a.x-b.x), abs(a.y-b.y))-1;
11
12
13 | int totalBoundaryPolygon(vector<pto> &arr, int n) {
15
     int boundPoint = n;
16
     for(int i = 0; i < n; i++) {
17
       boundPoint += boundaryCount(arr[i], arr[(i+1)%n]);
18
19
     return boundPoint;
20 }
21
22 int polygonArea2(vector<pto> &arr, int n) {
23
    int area = 0;
24
     // N = quantidade de pontos no polígono e armazenados em p;
2.5
     // OBS: VALE PARA CONVEXO E NÃO CONVEXO
26
     for(int i = 0; i<n; i++) {
       area += cross(arr[i], arr[(i+1)%n]);
27
28
29
     return abs(area);
30
31
32 | int internalCount(vector<pto> &arr, int n) {
33
34
     int area_2 = polygonArea2(arr, n);
     int boundPoints = totalBoundaryPolygon(arr,n);
    return (area_2 - boundPoints + 2)/2;
36
37 }
```

4.13. Polygon Area (3D)

```
#include <bits/stdc++.h>
3
   using namespace std;
4
   struct point {
    double x, y, z;
     void operator=(const point & b) {
       x = b.x;
9
       y = b.y;
10
       z = b.z;
11
12
13
14 point cross (point a, point b) {
1.5
    point ret;
16
    ret.x = a.y*b.z - b.y*a.z;
17
     ret.y = a.z*b.x - a.x*b.z;
18
     ret.z = a.x*b.y - a.y*b.x;
19
     return ret;
20
21
22
   int main() {
23
     int num;
24
     cin >> num;
     point v[num];
26
     for(int i=0; i<num; i++) cin >> v[i].x >> v[i].y >> v[i].z;
27
28
     point cur;
     cur.x = 0, cur.y = 0, cur.z = 0;
30
     for(int i=0; i<num; i++) {</pre>
32
       point res = cross(v[i], v[(i+1)%num]);
33
       cur.x += res.x;
34
       cur.y += res.y;
35
       cur.z += res.z;
36
37
38
     double ans = sqrt(cur.x*cur.x + cur.y*cur.y + cur.z*cur.z);
39
40
     double area = abs(ans);
41
42
     cout << fixed << setprecision(9) << area/2. << endl;
43
```

4.14. Polygon Area

```
double polygonArea(vector<pto> &arr, int n) {
   int area = 0;
   // N = quantidade de pontos no polígono e armazenados em p;
   // OBS: VALE PARA CONVEXO E NÃO CONVEXO
   for(int i = 0; i<n; i++) {
      area += cross(arr[i], arr[(i+1)%n]);
   }
   return (double)abs(area/2.0);
}</pre>
```

4.15. Segment-Segment Intersection

```
1 // Given three colinear points p, q, r, the function checks if
2 // point q lies on line segment 'pr'
3 int onSegment (Point p, Point q, Point r) {
```

```
if (q.x \le max(p.x, r.x) \& \& q.x \ge min(p.x, r.x) \& \& q.y \le max(p.y, r.y)
        && q.v >= min(p.v, r.v))
       return true;
     return false:
7
8 /★ PODE SER RETIRADO
9 int onSegmentNotBorda (Point p, Point q, Point r) {
       if (q.x < max(p.x, r.x) & q.x > min(p.x, r.x) & q.y <= max(p.y, r.y)
       && q.y >= min(p.y, r.y))
11
            return true:
       if (q.x \le max(p.x, r.x) \&  q.x \ge min(p.x, r.x) \&  q.y < max(p.y, r.y)
       && q.y > min(p.y, r.y)
            return true;
       return false;
14
15 }
16 */
17 // To find orientation of ordered triplet (p, q, r).
18 // The function returns following values
19 // 0 --> p, q and r are colinear
20 // 1 --> Clockwise
21 // 2 --> Counterclockwise
22 int orientation (Point p, Point q, Point r) {
   int val = (q.y - p.y) * (r.x - q.x) -
24
           (q.x - p.x) * (r.y - q.y);
25 | if (val == 0) return 0; // colinear
return (val > 0)? 1: 2; // clock or counterclock wise
27
28 // The main function that returns true if line segment 'p1p2'
29 // and 'q1q2' intersect.
30 | int doIntersect (Point p1, Point p2, Point q1, Point q2) {
31
    // Find the four orientations needed for general and
32
     // special cases
     int o1 = orientation(p1, p2, q1);
33
34
     int o2 = orientation(p1, p2, q2);
35
     int o3 = orientation(q1, q2, p1);
36
     int o4 = orientation(q1, q2, p2);
37
     // General case
3.8
     if (o1 != o2 && o3 != o4) return 2;
39
40
41 /* PODE SER RETIRADO
42 if(01 == 02 && 02 == 03 && 03 == 04 && 04 == 0) {
43
       //INTERCEPTAM EM RETA
44
       if (onSegmentNotBorda(p1,q1,p2) || onSegmentNotBorda(p1,q2,p2)) return 1;
45
       if(onSegmentNotBorda(q1,p1,q2) || onSegmentNotBorda(q1,p2,q2)) return 1;
46
47 */
    // Special Cases (INTERCEPTAM EM PONTO)
48
    // p1, p2 and q1 are colinear and q1 lies on segment p1p2
    if (o1 == 0 && onSegment(p1, q1, p2)) return 2;
     // p1, p2 and q1 are colinear and q2 lies on segment p1p2
52
     if (o2 == 0 && onSegment(p1, q2, p2)) return 2;
53
     // q1, q2 and p1 are colinear and p1 lies on segment q1q2
     if (o3 == 0 && onSegment(q1, p1, q2)) return 2;
55
     // q1, q2 and p2 are colinear and p2 lies on segment q1q2
56
     if (o4 == 0 && onSegment(q1, p2, q2)) return 2;
57
     return false; // Doesn't fall in any of the above cases
58
59 // OBS: SE (C2/A2 == C1/A1) SÃO COLINEARES
```

4.16. Upper And Lower Hull

```
struct pto {
double x, y;
```

```
bool operator < (const pto &p) const {
4
       return x < p.x || (x == p.x && y < p.y);
5
       /★ a impressao será em prioridade por mais a esquerda, mais
          abaixo, e antihorário pelo cross abaixo */
7
8
9
   double cross(const pto &O, const pto &A, const pto &B) {
    return (A.x - 0.x) \star (B.y - 0.y) - (A.y - 0.y) \star (B.x - 0.x);
10
11
12
   vector<pto> lower, upper;
13
14
   vector<pto> convex hull(vector<pto> &P) {
15
16
     int n = P.size(), k = 0;
17
     vector<pto> H(2 * n);
18
     // Sort points lexicographically
     sort(P.begin(), P.end());
19
     // Build lower hull
20
21
     for (int i = 0; i < n; ++i) {
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
22
23
       // trocar por >= 0
       while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
24
2.5
         k--;
       H[k++] = P[i];
26
27
28
     // Build upper hull
     for (int i = n - 2, t = k + 1; i >= 0; i--) {
29
30
       // esse <= 0 representa sentido anti-horario, caso deseje mudar
31
       // trocar por >= 0
32
       while (k \ge t \&\& cross(H[k - 2], H[k - 1], P[i]) \le 0)
        k--;
33
34
       H[k++] = P[i];
35
36
     H.resize(k);
37
     /★ o último ponto do vetor é igual ao primeiro, atente para isso
38
     as vezes é necessário mudar */
39
40
     int j = 1;
     lower.pb(H.front());
41
42
     while (H[j].x >= H[j-1].x) {
43
      lower.pb(H[j++]);
44
45
46
     int l = H.size()-1;
47
     while(1 >= i) {
48
       upper.pb(H[1--]);
49
50
     upper.pb(H[1--]);
51
52
     return H:
```

4.17. Struct Point And Line

```
int sgn(double x) {
   if(abs(x) < 1e-8)     return 0;
   return x > 0 ? 1 : -1;
}
inline double sqr(double x) {     return x * x; }

struct Point {
   double x, y, z;
   Point() {);
   Point(double a, double b): x(a), y(b) {};
```

```
Point (double x, double y, double z): x(x), y(y), z(z) {}
11
12
13
        void input() { scanf(" %lf %lf", &x, &y); };
14
        friend Point operator+(const Point &a, const Point &b) {
15
            return Point(a.x + b.x, a.v + b.v);
16
17
        friend Point operator-(const Point &a, const Point &b) {
18
           return Point(a.x - b.x, a.y - b.y);
19
2.0
21
       bool operator !=(const Point& a) const {
22
            return (x != a.x || y != a.y);
23
24
25
       bool operator <(const Point &a) const{</pre>
26
         if(x == a.x)
27
           return y < a.y;</pre>
28
          return x < a.x;</pre>
29
30
31
        double norm() {
32
            return sqrt(sqr(x) + sqr(y));
3.3
34
35
   double det (const Point &a, const Point &b) {
        return a.x * b.v - a.v * b.x;
37
   double dot(const Point &a, const Point &b) {
38
39
       return a.x * b.x + a.y * b.y;
40
41
   double dist (const Point &a, const Point &b) {
42
       return (a-b).norm();
43
44
45
   struct Line {
47
        Point a, b;
48
        Line() {}
49
        Line (Point x, Point y): a(x), b(y) {};
50
51
52 double dis_point_segment(const Point p, const Point s, const Point t) {
53
       if(sqn(dot(p-s, t-s)) < 0)
54
         return (p-s).norm();
        if(sqn(dot(p-t, s-t)) < 0)
55
56
         return (p-t).norm();
57
        return abs(det(s-p, t-p) / dist(s, t));
58 }
```

5. Graphs

5.1. Checa Grafo Bipartido

```
bool isBipartite(int src, int V) {

int colorArr[V + 1];
   memset(colorArr, -1, sizeof(colorArr));
   colorArr[src] = 1;

queue <int> q; q.push(src);

while (!q.empty()) {
   int u = q.front(); q.pop();
}
```

```
// Find all non-colored adjacent vertices
       for (auto it = adj[u].begin(); it != adj[u].end(); it++) {
13
       //Return false if there is a self-loop
14
15
           if (u == *it)
16
               return false;
17
         // An edge from u to v exists and destination v is not colored
18
19
         if (colorArr[*it] == -1) {
20
           // Assign alternate color to this adjacent v of u
21
           colorArr[*it] = 1 - colorArr[u];
22
           q.push(*it);
23
24
         // An edge from u to v exists and destination v is colored with same
       color as u
25
         else if (colorArr[*it] == colorArr[u])
26
           return false;
27
28
29
     // If we reach here, then all adjacent vertices can be colored with
30
     // alternate color
31
    return true;
32
```

5.2. Ciclo Grafo

```
int n;
2
   vector<vector<int>> adj;
3
   vector<char> color;
   vector<int> parent;
   int cycle_start, cycle_end;
   bool dfs(int v) {
    color[v] = 1;
    for (int u : adi[v]) {
10
       if (color[u] == 0) {
11
         parent[u] = v;
12
         if (dfs(u))
13
           return true;
       } else if (color[u] == 1) {
14
15
         cycle_end = v;
16
         cycle_start = u;
17
         return true;
18
19
20
     color[v] = 2:
21
     return false;
22
23
   void find_cycle() {
24
25
     color.assign(n, 0);
26
     parent.assign(n, -1);
27
     cycle_start = -1;
28
     for (int v = 0; v < n; v++) {
29
30
       if (dfs(v))
         break:
31
32
33
     if (cycle_start == -1) {
34
       cout << "Acyclic" << endl;
35
36
     } else {
37
       vector<int> cycle;
38
       cycle.push_back(cycle_start);
39
       for (int v = cycle_end; v != cycle_start; v = parent[v])
```

```
cycle.push_back(v);
40
41
        cvcle.push back(cvcle start);
       reverse(cycle.begin(), cycle.end());
42
43
44
       cout << "Cycle found: ";
45
       for (int v : cycle)
46
         cout << v << " ";
47
       cout << endl;
48
49
```

5.3. Diametro Em Arvore

1 Calcula qual o vértice a mais distante de um qualquer vértice X e **do** vértice A calcula-se o vértice B mais distante dele.

5.4. Floyd Warshall

```
// OBS: ZERAR adj[i][i] sempre
for(int i = 0; i < n; i++)
    adj[i][i] = 0;

for(int k = 0; k < n; k++) {
    for(int i = 0; i < n; i++) {
        for(int j = 0; j < n; j++) {
            adj[i][j] = min(adj[i][j],adj[i][k] + adj[k][j]);
        }
        }
    }
}</pre>
```

5.5. Ford Fulkersson (Maximum Flow)

```
1 | int rGraph[2000][2000];
   int graph[2000][2000];
3
4 int V;
5 | bool bfs(int s, int t, int parent[]) {
    bool visited[V];
6
     memset(visited, 0, sizeof(visited));
     // Create a queue, enqueue source vertex and mark source vertex
9
10
     // as visited
11
     queue <int> q:
12
     q.push(s);
     visited[s] = true;
13
     parent[s] = -1;
14
15
16
      // Standard BFS Loop
17
      while (!q.empty()) {
18
       int u = q.front();
19
       q.pop();
20
2.1
        for (int v=0; v<V; v++) {</pre>
22
         if (visited[v] == false && rGraph[u][v] > 0) {
23
           q.push(v);
24
           parent[v] = u;
25
           visited[v] = true;
26
27
28
29
     // If we reached sink in BFS starting from source, then return true, else
```

```
return (visited[t] == true);
31
32
33 // Returns the maximum flow from s to t in the given graph
34 int fordFulkerson(int s, int t) {
     int u, v;
     // Create a residual graph and fill the residual graph with given
       capacities in the original graph as residual capacities in residual
       graph residual capacity of edge from i to j (if there is an edge. If
       rGraph[i][i] is 0, then there is not)
     for (u = 0; u < V; u++)
38
       for (v = 0; v < V; v++)
39
         rGraph[u][v] = graph[u][v];
40
     int parent[V]; // This array is filled by BFS and to store path
41
42
43
     int max flow = 0;// There is no flow initially
44
45
     // Augment the flow while tere is path from source to sink
46
     while (bfs(s, t, parent)) {
       // Find minimum residual capacity of the edges along the path filled by
       BFS. Or we can say find the maximum flow through the path found.
       int path flow = INT MAX;
49
       for (v=t; v!=s; v=parent[v]) {
50
         u = parent[v];
51
         path_flow = min(path_flow, rGraph[u][v]);
52
53
54
       // update residual capacities of the edges and reverse edges
55
       // along the path
56
       for (v=t; v != s; v=parent[v]) {
57
         u = parent[v];
58
         rGraph[u][v] -= path flow;
59
         rGraph[v][u] += path flow;
60
61
62
       // Add path flow to overall flow
63
       max_flow += path_flow;
64
65
66
     // Return the overall flow
67
     return max flow:
68
   // PRINT THE FLOW AFTER RUNNING THE ALGORITHM
   void print(int n) {
     for(int i = 1; i <= m; i++) {
73
       for (int j = m+1; j \le m*2; j++) {
74
         cout << "flow from i(left) to j(right) is " << graph[i][j] -
       rGraph[i][i] << endl;
75
76
77
78
79
   void addEdge(int 1, int r, int n, int x) {
     graph[l][r+n] = x;
80
81
82
83
   void addEdgeSource(int 1, int x) {
84
    graph[0][1] = x;
85
86
   void addEdgeSink(int r, int n, int x) {
87
88
    graph[r+n][V-1] = x;
```

5.6. Pontes Num Grafo

```
1 //SE TIRA-LAS O GRAFO FICA DESCONEXO
2 // OBS: PRESTAR ATENCAO EM SELF-LOOPS, É MELHOR NÃO ADICIONA-LOS
3 // SO FUNCIONA EM GRAFO NÃO DIRECIONADO
   vector<int> T((int)2e6,0); //Tempo necessário para chegar naquele vértice na
   vector<int> adj[(int)2e6];
   vector<int> Low((int)2e6); // Tempo "mínimo" para chegar naquele vértice na
   vector<int> ciclo((int)2e6, false);
   vector<ii> bridges:
10 void dfs(int u, int p) {
    Low[u] = T[u] = t;
11
12
13
     for(auto v : adi[u]){
       if (v==p) {
14
15
          //checa arestas paralelas
16
          p = -1;
17
          continue;
18
19
        //se ele ainda não foi visited
2.0
        else if(T[v]==0){
21
          dfs(v,u);
          Low[u]=min(Low[u], Low[v]);
22
23
          if(Low[v]>T[u]) {
24
            bridges.pb(ii(min(u,v), (max(u,v))));
25
          // ponte de u para v
26
27
28
        else
29
         Low[u]=min(Low[u], T[v]);
30
        ciclo[u] \mid = (T[u] > = Low[v]);
31
        //checa se o vértice u faz parte de um ciclo
32
33
34
35
   void clear() {
36
37
     for(int i = 0; i <= n; i++) {
38
       T[i] = 0, Low[i] = 0, adj[i].clear(), ciclo[i] = false;
39
40
     bridges.clear();
41
42
43
    signed main () {
45
46
     for(int i = 0; i < n; i++)</pre>
47
       if(T[i] == 0)
48
          dfs(i, -1);
49
     sort(bridges.begin(), bridges.end());
51
52
     cout << (int)bridges.size() << endl;</pre>
53
     for(int i = 0; i < bridges.size(); i++) {</pre>
       cout << bridges[i].ff << " - " << bridges[i].ss << endl;</pre>
54
55
56
     cout << endl:
57
58
     clear();
```

```
59
60 }
```

5.7. Pontos De Articulação

```
// SE TIRAR TAIS VERTICES O GRAFO FICA DESCONEXO
2
   vector<bool> ap(100000, false);
   vector<int> low(100000,0), T(100000,0);
   int tempo = 1:
   list<int> adj[100000];
   void artPoint(int u, int p) {
9
10
     low[u] = T[u] = tempo++;
11
     int children = 0;
12
     for(int v: adj[u]) {
13
14
15
     // cuidado com arestas paralelas
16
     // se tiver nao podemos fazer assim
17
18
       if(T[v] == 0) {
19
20
          children++;
21
          artPoint(v,u);
22
         low[u] = min(low[v], low[u]);
23
       if(p == -1 && children > 1) {
24
25
          ap[u] = true;
26
27
28
       if(p != -1 && low[v] > T[u])
29
         ap[u] = true;
        } else if (v != p)
30
31
         low[u] = min(low[u], T[v]);
32
33
34
35
36
   int main() {
37
38
     for(int i = 0; i < n; i++)</pre>
39
       if(T[i] == 0)
40
          artPoint(i,-1);
41
```

5.8. Scc (Kosaraju)

```
class SCC {
    private:
     // number of vertices
     int n;
     // indicates whether it is indexed from 0 or 1
     int indexed from;
     // reversed graph
     vector<vector<int>> trans;
1.0
    private:
11
     void dfs_trans(int u, int id) {
       comp[u] = id;
12
13
       scc[id].push_back(u);
14
```

```
for (int v: trans[u])
15
16
         if (comp[v] == -1)
           dfs trans(v, id);
17
18
19
20
     void get_transpose(vector<vector<int>>& adj) {
21
        for (int u = indexed from; u < this->n + indexed from; u++)
          for(int v: adj[u])
22
23
           trans[v].push_back(u);
24
25
     void dfs_fill_order(int u, stack<int> &s, vector<vector<int>>& adj) {
26
27
       comp[u] = true;
28
        for(int v: adj[u])
29
30
         if(!comp[v])
31
           dfs_fill_order(v, s, adj);
32
33
       s.push(u);
34
35
36
     // The main function that finds all SCCs
37
     void compute SCC(vector<vector<int>>& adi) {
38
39
       // Fill vertices in stack according to their finishing times
40
        for(int i = indexed from; i < this->n + indexed from; i++)
41
42
         if(!comp[i])
43
           dfs_fill_order(i, s, adj);
44
4.5
        // Create a reversed graph
46
       get_transpose(adj);
47
48
       fill(comp.begin(), comp.end(), -1);
49
50
        // Now process all vertices in order defined by stack
51
       while(s.empty() == false) {
52
         int v = s.top();
53
          s.pop();
54
55
         if(comp[v] == -1)
56
           dfs_trans(v, this->number_of_comp++);
57
58
59
    public:
    // number of the component of the i-th vertex
62
     // it's always indexed from 0
     vector<int> comp;
63
     /\!/ the i-th vector contains the vertices that belong to the i-th scc
64
     // it's always indexed from 0
65
66
     vector<vector<int>> scc;
     int number of comp = 0;
67
68
69
     SCC(int n, int indexed_from, vector<vector<int>>& adj) {
70
       this->n = n;
71
       this->indexed from = indexed from:
72
       comp.resize(n + 1);
73
       trans.resize(n + 1);
74
       scc.resize(n + 1);
75
76
       this->compute_SCC(adj);
77
78 };
```

60

61

62

63

64

65

66

67 68

70

71

72

73

74

75

76

77

78

79

80

81

82

8.3

84

85

86

87

88

91

92

93

95

96

97

98

99

100

101

102

103

104

105

106

107

5.9. All Eulerian Path Or Tour

```
struct edge {
     int v, id;
3
     edge() {}
     edge(int v, int id) : v(v), id(id) {}
4
5
   // The undirected + path and directed + tour wasn't tested in a problem.
8 // TEST AGAIN BEFORE SUBMITTING IT!
9 namespace graph {
     // Namespace which auxiliary funcions are defined.
10
11
     namespace detail {
       pair<bool, pair<int, int>> check_both_directed(const
12
       vector<vector<edge>> &adj, const vector<int> &in_degree) {
         // source and destination
1.3
         int src = -1, dest = -1;
14
         // adj[i].size() represents the out degree of an vertex
15
         for (int i = 0; i < adj.size(); i++)
16
           if((int)adj[i].size() - in_degree[i] == 1) {
17
18
             if(src != -1)
19
               return make_pair(false, pair<int, int>());
20
              src = i;
21
           } else if((int)adj[i].size() - in_degree[i] == -1) {
22
              if(dest != -1)
               return make_pair(false, pair<int, int>());
23
24
             dest = i;
25
           } else if(abs((int)adj[i].size() - in_degree[i]) > 1)
26
              return make_pair(false, pair<int, int>());
27
28
         if (src == -1 && dest == -1)
29
           return make_pair(true, pair<int, int>(src, dest));
30
31
         else if (src !=-1 && dest !=-1)
32
           return make_pair(true, pair<int, int>(src, dest));
33
34
         return make_pair(false, pair<int, int>());
35
36
37
       /// Builds the path/tour for directed graphs.
38
       void build(const int u, vector<int> &tour, vector<vector<edge>> &adj,
       vector<bool> &used) {
39
         while(!adj[u].empty()) {
40
           const edge e = adi[u].back();
41
           if(!used[e.id]) {
42
             used[e.id] = true;
43
              adi[u].pop back();
             build (e.v, tour, adj, used);
44
45
           } else
46
              adj[u].pop_back();
47
48
49
         tour.push_back(u);
50
51
       /// Auxiliary function to build the eulerian tour/path.
52
53
       vector<int> set_build(vector<vector<edge>> &adj, const int E, const int
54
         vector<int> path;
                                                                                     108
55
         vector<bool> used(E + 3);
                                                                                     109
56
                                                                                     110
57
         build(first, path, adj, used);
                                                                                     111
58
                                                                                     112
59
         for (int i = 0; i < adj.size(); i++)
```

```
// if there are some remaining edges, it's not possible to build the
  tour.
      if(adj[i].size())
        return vector<int>();
    reverse(path.begin(), path.end());
    return path:
/// All vertices v should have in_degree[v] == out_degree[v]. It must not
  contain a specific
/// start and end vertices.
/// Time complexity: O(V * (log V) + E)
bool has euler tour directed (const vector < vector < edge >> & adj, const
  vector<int> &in degree) {
  const pair<bool, pair<int, int>> aux = detail::check_both_directed(adj,
  in degree);
  const bool valid = aux.first;
  const int src = aux.second.first;
  const int dest = aux.second.second;
  return (valid && src == -1 && dest == -1);
/// A directed graph has an eulerian path/tour if has:
/// - One vertex v such that out_degree[v] - in_degree[v] == 1
/// - One vertex v such that in_degree[v] - out_degree[v] == 1
/// - The remaining vertices v such that in_degree[v] == out_degree[v]
/// or
/// - All vertices v such that in_degree[v] - out_degree[v] == 0 -> TOUR
/// Returns a boolean value that indicates whether there's a path or not.
/// If there's a valid path it also returns two numbers: the source and
  the destination.
/// If the source and destination can be an arbitrary vertex it will
  return the pair (-1, -1)
/// for the source and destination (it means the contains an eulerian
  tour).
111
/// Time complexity: O(V + E)
pair<bool, pair<int, int>> has_euler_path_directed(const
  vector<vector<edge>> &adj, const vector<int> &in_degree) {
  return detail::check_both_directed(adj, in_degree);
/// Returns the euler path. If the graph doesn't have an euler path it
  returns an empty vector.
/// Time Complexity: O(V + E) for directed, O(V * log(V) + E) for
  undirected.
/// Time Complexity: O(adj.size() + sum(adj[i].size()))
vector<int> get_euler_path_directed(const int E, vector<vector<edge>>
  &adi, const vector<int> &in degree) {
  const pair<bool, pair<int, int>> aux = has_euler_path_directed(adj,
  in degree);
  const bool valid = aux.first;
  const int src = aux.second.first;
 const int dest = aux.second.second;
  if(!valid)
   return vector<int>();
  int first;
  if (src !=-1)
```

```
113
          first = src;
114
        else {
115
           first = 0:
          while (adj[first].empty())
116
117
            first++;
118
119
120
        return detail::set_build(adj, E, first);
121
122
      /// Returns the euler tour. If the graph doesn't have an euler tour it
123
        returns an empty vector.
124
      /// Time Complexity: O(V + E)
125
126
      /// Time Complexity: O(adj.size() + sum(adj[i].size()))
127
      vector<int> get_euler_tour_directed(const int E, vector<vector<edge>>
        &adj, const vector<int> &in_degree) {
128
        const bool valid = has_euler_tour_directed(adj, in_degree);
129
130
        if(!valid)
131
          return vector<int>();
132
133
        int first = 0:
134
        while(adj[first].empty())
135
          first++;
136
137
        return detail::set_build(adj, E, first);
138
139
140
      // The graph has a tour that passes to every edge exactly once and gets
141
      // back to the first edge on the tour.
142
143
      // A graph with an euler path has zero odd degree vertex.
144
145
      // Time Complexity: O(V)
146
      bool has euler tour undirected(const vector<int> &degree) {
        for(int i = 0; i < degree.size(); i++)</pre>
147
          if(degree[i] & 1)
148
149
            return false;
150
        return true;
151
152
153
      // The graph has a path that passes to every edge exactly once.
154
      // It doesn't necessarely gets back to the beginning.
155
156
      // A graph with an euler path has two or zero (tour) odd degree vertices.
157
      // Returns a pair with the startpoint/endpoint of the path.
158
159
      // Time Complexity: O(V)
160
      pair<bool, pair<int, int>> has_euler_path_undirected(const vector<int>
161
        &degree) {
        vector<int> odd_degree;
162
163
        for (int i = 0; i < degree.size(); i++)
          if(degree[i] & 1)
164
165
            odd_degree.pb(i);
166
167
        if(odd_degree.size() == 0)
168
          return make_pair(true, make_pair(-1, -1));
169
        else if (odd_degree.size() == 2)
170
          return make_pair(true, make_pair(odd_degree.front(),
        odd_degree.back()));
171
        else
172
           return make_pair(false, pair<int, int>());
173
```

```
175
      vector<int> get euler tour undirected(const int E, const vector<int>
         &degree, vector<vector<edge>> &adj) {
176
        if (!has_euler_tour_undirected(degree))
177
          return vector<int>();
178
179
        int first = 0:
180
        while(adj[first].empty())
181
          first++;
182
183
        return detail::set_build(adj, E, first);
184
185
186
      /// Returns the euler tour. If the graph doesn't have an euler tour it
        returns an empty vector.
187
      /// Time Complexity: O(V + E)
188
189
      /// Time Complexity: O(adj.size() + sum(adj[i].size()))
190
      vector<int> get_euler_path_undirected(const int E, const vector<int>
        &degree, vector<vector<edge>> &adj)
191
        auto aux = has_euler_path_undirected(degree);
        const bool valid = aux.first;
192
193
        const int x = aux.second.first;
194
        const int y = aux.second.second;
195
196
        if(!valid)
197
          return vector<int>();
198
199
        int first;
200
        if (x != -1) {
201
          first = x;
202
          adj[x].emplace_back(y, E + 1);
203
          adj[y].emplace_back(x, E + 1);
204
         } else {
205
          first = 0;
206
          while(adj[first].emptv())
207
            first++:
208
209
210
        vector<int> ans = detail::set_build(adj, E, first);
211
        reverse(ans.begin(), ans.end());
212
        if (x != -1)
213
          ans.pop_back();
214
        return ans:
215
216 };
```

5.10. Bellman Ford

```
struct edge {
     int src, dest, weight;
3
     edge() {}
     edge (int src, int dest, int weight) : src(src), dest(dest), weight(weight)
5
6
     bool operator<(const edge &a) const {</pre>
7
        return weight < a.weight;</pre>
8
9
   };
1.0
11 /// Works to find the shortest path with negative edges.
12 /// Also detects cycles.
13 | ///
14 /// Time Complexity: O(n * e)
```

```
15 | /// Space Complexity: O(n)
16 | bool bellman ford(vector<edge> &edges, int src, int n) {
17
     // n = qtd of vertices, E = qtd de arestas
18
19
     // To calculate the shortest path uncomment the line below
20
     // vector<int> dist(n, INF);
21
22
     // To check cycles uncomment the line below
23
     // vector<int> dist(n, 0);
2.4
25
     vector<int> pai(n, -1);
26
     int E = edges.size();
27
28
     dist[src] = 0;
29
     // Relax all edges n - 1 times.
30
     // A simple shortest path from src to any other vertex can have at-most n
       - 1 edges.
31
     for (int i = 1; i <= n - 1; i++) {
32
       for (int j = 0; j < E; j++) {
33
         int u = edges[j].src;
34
         int v = edges[j].dest;
35
         int weight = edges[i].weight;
36
         if (dist[u] != INF && dist[u] + weight < dist[v]) {</pre>
37
           dist[v] = dist[u] + weight;
38
           pai[v] = u;
39
40
41
42
43
     // Check for NEGATIVE-WEIGHT CYCLES.
44
     // The above step guarantees shortest distances if graph doesn't contain
       negative weight cycle.
45
     // If we get a shorter path, then there is a cycle.
     bool is cycle = false;
46
47
     int vert_in_cycle;
48
     for (int i = 0; i < E; i++) {
       int u = edges[i].src;
49
       int v = edges[i].dest;
50
       int weight = edges[i].weight;
51
52
       if (dist[u] != INF && dist[u] + weight < dist[v]) {
53
        is_cycle = true;
         pai[v] = u;
54
55
         vert_in_cycle = v;
56
57
58
59
     if(is cvcle) {
       for (int i = 0; i < n; i++)
60
61
         vert_in_cycle = pai[vert_in_cycle];
62
63
       vector<int> cvcle;
64
       for(int v = vert_in_cycle; (v != vert_in_cycle || cycle.size() <= 1); v</pre>
       = pai[v])
65
         cycle.pb(v);
66
67
       reverse(cycle.begin(), cycle.end());
68
69
       for(int x: cycle) {
70
         cout << x + 1 << ' ';
71
72
       cout << cycle.front() + 1 << endl;</pre>
73
       return true;
74
       else
75
       return false;
76
```

5.11. De Bruijn Sequence

```
1 // We can solve this problem by constructing a directed graph with
 2 // k^(n-1) nodes with each node having k outgoing edges order. Each node
       corresponds
3 \mid // to a string of size n-1. Every edge corresponds to one of the k characters
4 // in A and adds that character to the starting string.
5 // For example, if n=3 and k=2, then we construct the following graph:
8 //
                - 1 -> (O1) - 1 ->
                       `^ |
9 1/
10
  // 0 -> (00)
                        1 0
                                      (11) <- 1
11 | //
                        | v
12 //
                <- 0 - (10) <- 0 -
13
14 \mid // The node 01 is connected to node 11 through edge 1, as adding 1 to 01
15 // (and removing the first character) gives us 11.
16 //
17 // We can observe that every node in this graph has equal in-degree and
       out-degree,
  // which means that a Eulerian circuit exists in this graph.
19
20 namespace graph {
     namespace detail {
21
22
       // Finding an valid eulerian path
23
       void dfs(const string &node, const string &alphabet, set<string> &vis,
       string &edges_order) {
24
         for(char c: alphabet) {
25
           string nxt = node + c;
26
           if (vis.count (nxt))
27
             continue;
28
29
           vis.insert(nxt);
30
           nxt.erase(nxt.begin());
31
           dfs(nxt, alphabet, vis, edges_order);
32
           edges_order += c;
33
34
35
     };
36
     appears in the
     // resulting string exactly once.
3.8
39
40
     // Time Complexity: O(alphabet.size() ^ n * log2(alphabet.size() ^ n))
41
     string de_bruijn(const int n, const string &alphabet) {
42
       set<string> vis;
43
       string edges_order;
44
45
       string starting_node = string(n - 1, alphabet.front());
46
       detail::dfs(starting_node, alphabet, vis, edges_order);
47
48
       return edges_order + starting_node;
49
50 };
```

5.12. Dijkstra + Dij Graph

```
1 class Dijkstra {
private:
```

```
int src, dest;
     int n;
5
6
    private:
     int calculate(vector<vector<ii>>> &adj) {
8
       dist.resize(this->n + 1, INF);
9
       parent.resize(this->n + 1);
       vector<int> vis(this->n + 1, 0);
1.0
11
12
       for(int i = 0; i <= this->n; i++)
13
         parent[i].pb(i);
14
15
       priority_queue<ii, vector<ii>, greater<ii>> pq;
        pq.push(make_pair(0, this->src));
16
       dist[this->src] = 0;
17
18
19
       while(!pq.empty()) {
20
         int u = pq.top().ss;
21
          pq.pop();
22
          if (vis[u])
23
           continue;
24
          vis[u] = true;
25
26
          for (ii x: adj[u]) {
27
           int v = x.ff;
            int w = x.ss;
28
29
30
            if(dist[u] + w < dist[v]) {</pre>
31
             parent[v].clear();
32
              parent[v].pb(u);
33
              dist[v] = dist[u] + w;
34
             pq.push(ii(dist[v], v));
35
36
            else if(dist[u] + w == dist[v]) {
37
              parent[v].pb(u);
38
39
40
41
42
       return dist[dest];
43
44
45
     // Use a vector vis in a DFS on the dijkstra graph
46
     vector<vector<int>> gen_dij_graph() {
47
       vector<vector<int>> dijkstra_graph(this->n + 1);
48
       vector<bool> vis(this->n + 1);
49
50
       queue<int> q;
51
       q.push(this->dest);
       while (!q.empty())
52
53
         int v = q.front();
54
         q.pop();
55
56
          for(int u: parent[v]) {
57
           if(u == v)
58
              continue;
59
            dijkstra_graph[u].pb(v);
60
            if(!vis[u]) {
61
              q.push(u);
62
              vis[u] = true;
63
64
65
66
       return dijkstra_graph;
```

```
public:
69
70
      int min_path;
71
      vector<int> dist;
72
      vector<vector<int>> parent;
73
      vector<vector<int>> dij_graph;
74
75
      Dijkstra(int n, int src, int dest, vector<vector<ii>>> &adj) {
76
        this -> n = n;
77
        this->src = src:
78
        this->dest = dest;
        this->min_path = this->calculate(adj);
79
80
        /// Generates the dijkstra graph with the parent vector
        this->dij_graph = this->gen_dij_graph();
81
        d_graph = this->dij_graph;
82
83
84
85
      // Returns a path with minimum costs and a minimum length.
      vector<int> get_min_path() {
86
87
        vector<int> path;
88
        vector<int> pai(this->n + 1, -1);
89
        vector<int> d(this->n + 1, INF);
90
91
        queue<int> q:
92
        q.push(this->dest);
93
        d[this->dest] = 0;
94
95
        while(!q.empty()) {
96
          int v = q.front();
97
          q.pop();
98
99
          for(int u: parent[v]) {
100
            if(u == v)
101
              continue;
            if(d[v] + 1 < d[u]) {
102
103
              d[u] = d[v] + 1;
              pai[u] = v;
104
105
              q.push(u);
106
107
108
109
        int cur = this->src;
110
111
        while(cur != -1) {
112
          path.pb(cur);
113
          cur = pai[cur];
114
115
116
        return path;
117
118 | };
```

5.13. Dinic (Max Flow)

```
// Created by Ubiratan Neto

struct Dinic {

struct FlowEdge {
   int v, rev, c, cap;
   bool is_rev;
   FlowEdge() {}
   FlowEdge(int v, int c, int cap, int rev, bool is_rev) : v(v), c(c), cap(cap), rev(rev), is_rev(is_rev) {}
}
```

```
10
     };
11
12
     vector<vector<FlowEdge>> adi:
13
     vector<int> level, used;
     int src, snk, V;
14
15
     int sz;
16
     int max flow:
     bool calculated;
17
18
     Dinic(){}
19
     Dinic(int n) {
20
       calculated = false;
21
       src = 0;
22
       snk = n+1;
23
       adj.resize(n+2, vector< FlowEdge >());
24
       level.resize(n+2);
25
       used.resize(n+2);
26
       sz = n+2;
27
       V = n+2;
28
       max_flow = 0;
29
30
31
     void add edge(int u, int v, int c){
32
       int id1 = adi[u].size();
33
       int id2 = adj[v].size();
34
       adj[u].pb(FlowEdge(v, c, c, id2, false));
35
       adj[v].pb(FlowEdge(u, 0, 0, idl, true));
36
37
38
     void add_to_src(int v, int c){
39
       adj[src].pb(FlowEdge(v, c, c, -1, false));
40
41
42
     void add to snk(int u, int c){
43
       adj[u].pb(FlowEdge(snk, c, c, -1, false));
44
45
46
     bool bfs() {
47
       for (int i=0; i<sz; i++) {</pre>
48
         level[i] = -1;
49
50
51
       level[src] = 0;
52
       queue<int> q; q.push(src);
53
       while(!q.empty()){
54
55
         int cur = q.front();
56
          q.pop();
57
          for(FlowEdge e : adj[cur]){
58
           if(level[e.v] == -1 && e.c > 0){
59
             level[e.v] = level[cur]+1;
60
              q.push(e.v);
61
62
63
64
       return (level[snk] == -1 ? false : true);
65
66
67
68
     int send_flow(int u, int flow) {
69
       if(u == snk) return flow;
70
71
       for(int &i = used[u]; i<adj[u].size(); i++){</pre>
72
         FlowEdge &e = adj[u][i];
73
74
          if(level[u]+1 != level[e.v] || e.c <= 0) continue;</pre>
```

```
75
 76
           int new flow = min(flow, e.c);
 77
           int adjusted_flow = send_flow(e.v, new_flow);
 78
 79
           if(adjusted flow > 0){
 80
             e.c -= adjusted_flow;
 81
             if(e.rev != -1) adj[e.v][e.rev].c += adjusted_flow;
 82
             return adjusted_flow;
 83
 84
 85
 86
         return 0;
 87
 88
 89
       int calculate() {
 90
         if(src == snk) {max_flow = -1; return - 1;} //not sure if needed
 91
 92
         max_flow = 0;
 93
 94
         while(bfs()){
 95
           for(int i=0; i<sz; i++) used[i] = 0;</pre>
 96
           while(int inc = send flow(src, INF)) max flow += inc;
 97
         calculated = true:
 98
 99
100
         return max_flow;
101
102
103
104
       vector<ii> mincut (vector<vector<int>> &mat_adj) {
105
         assert (calculated);
106
         int mat[sz][sz];
107
         memset(mat, 0, sizeof mat);
108
         for(int i = 0; i < V; i++)
109
           for(FlowEdge x: adj[i])
110
             mat[i][x.v] += x.c;
111
         vector<bool> vis(sz);
112
113
         queue<int> q;
114
         q.push(src);
115
         vis[src] = true;
116
         while(!q.empty()) {
117
           int u = q.front();
118
           q.pop();
119
           for (int v = 0; v < sz; v++) {
120
             if (mat[u][v] > 0 && !vis[v]) {
121
               q.push(v);
122
               vis[v] = true;
123
124
125
126
127
         vector<ii> cut;
         for(int i = 0; i < sz; i++)</pre>
128
129
           for(int j = 0; j < sz; j++)
             if(vis[i] && !vis[j])
130
131
               // if there's an edge from i to j.
132
               if(mat\_adj[i][j] > 0)
133
                 cut.emplace_back(i, j);
134
135
         return cut;
136
137
138
      vector<ii> min_edge_cover() {
139
        bool covered[sz];
```

```
for(int i=0; i<sz; i++) covered[i] = false;</pre>
140
141
         vector< ii > edge cover;
142
         for (int i=1; i < sz-1; i++) {
143
           for(FlowEdge e : adj[i]) {
144
            if (e.cap == 0 | l | e.v > sz-2) continue;
145
             if(e.c == 0){
               edge_cover.pb(ii(i, e.v));
146
147
               covered[i] = true;
148
               covered[e.v] = true;
149
               break:
150
151
152
         for (int i=1; i<sz-1; i++) {
153
154
           for(FlowEdge e : adj[i]){
155
             if (e.cap == 0 \mid \mid e.v > sz-2) continue;
             if (e.c == 0) continue;
156
157
             if(!covered[i] || !covered[e.v]){
158
               edge_cover.pb(ii(i, e.v));
159
               covered[i] = true;
               covered[e.v] = true;
160
161
162
163
164
        return edge_cover;
165
166
167
      vector<vector<int>> allFlow() {
168
         assert (calculated);
169
         vector<vector<int>> ret(V, vector<int>(V, 0));
170
         for (int i = 0; i < V; i++) {
171
172
           for(FlowEdge x: adj[i]) {
173
             if(x.is rev)
174
               continue;
175
             // flow from vertex i to x.v
176
             ret[i][x.v] += x.cap - x.c;
177
178
179
180
         // for (int i = 0; i < V; i++) {
             for (int j = 0; j < V; j++) {
181
182
                cout << ret[i][j] << ' ';
183
184
             cout << endl;
185
         // }
186
187
         return ret;
188
189
190
      void dfs_build_path(int u, vector<int> &path, vector<vector<int>>
         &mat_flow, vector<vector<int>> &ans, vector<vector<int>> &adj) {
191
         path.pb(u);
192
193
         if(u == this -> snk) {
194
           ans.pb(path);
195
           return;
196
197
198
        for(int v: adj[u]) {
199
          if(mat_flow[u][v]) {
200
             mat_flow[u][v]--;
             dfs_build_path(v, path, mat_flow, ans, adj);
201
202
             return;
203
```

```
204
205
206
207
      vector<vector<int>> get_all_paths(vector<vector<int>> &adj) {
208
        assert (calculated);
209
210
        vector<vector<int>> mat_flow = allFlow();
211
        vector<vector<int>> ans;
212
        ans.reserve(max_flow);
213
214
         for(int i = 0; i < max_flow; i++) {
215
          vector<int> path;
216
           path.reserve(V);
217
           dfs_build_path(this->src, path, mat_flow, ans, adj);
218
219
220
        return ans;
221
222 };
```

5.14. Functional Graph

```
1 // Based on:
       http://maratona.ic.unicamp.br/MaratonaVerao2020/lecture-b/20200122.pdf
   class Functional_Graph {
    // FOR DIRECTED GRAPH
    private:
     void compute_cycle(int u, vector<int> &nxt, vector<bool> &vis) {
       int id_cycle = cycle_cnt++;
8
       int cur id = 0;
9
       this->first[id_cycle] = u;
10
11
       while(!vis[u]) {
12
         vis[u] = true;
13
14
         this->cycle[id_cycle].push_back(u);
15
         this->in_cycle[u] = true;
16
17
         this->cycle_id[u] = id_cycle;
18
         this->id_in_cycle[u] = cur_id;
19
         this->near_in_cycle[u] = u;
20
         this->id_near_cycle[u] = id_cycle;
21
         this->cycle_dist[u] = 0;
22
23
         u = nxt[u];
24
         cur id++;
25
26
27
28
     // Time Complexity: O(V)
29
     void build(int n, int indexed_from, vector<int> &nxt, vector<int>
       &in_degree) {
30
       queue<int> q;
31
       vector<bool> vis(n + indexed_from);
32
       for(int i = indexed_from; i < n + indexed_from; i++) {</pre>
33
         if(in_degree[i] == 0) {
34
           q.push(i);
35
           vis[i] = true;
36
37
38
39
       vector<int> process_order;
40
       process_order.reserve(n + indexed_from);
```

```
while(!q.empty()) {
 42
          int u = q.front();
                                                                                        104
 43
          q.pop();
                                                                                        105
 44
                                                                                        106
 45
                                                                                        107
          process order.push back(u);
 46
                                                                                        108
 47
          if(--in_degree[nxt[u]] == 0) {
                                                                                        109
 48
            q.push(nxt[u]);
                                                                                        110
 49
            vis[nxt[u]] = true;
 50
                                                                                        112
 51
 52
                                                                                        114
 53
        int cycle cnt = 0;
 54
        for (int i = indexed_from; i < n + indexed_from; i++)</pre>
                                                                                        115
 55
          if(!vis[i])
                                                                                        116
 56
            compute_cycle(i, nxt, vis);
                                                                                        117
 57
                                                                                        118
58
        for(int i = (int)process_order.size() - 1; i >= 0; i--) {
                                                                                        119
 59
          int u = process_order[i];
                                                                                        120
 60
                                                                                        121
          this->near in cycle[u] = this->near in cycle[nxt[u]];
                                                                                        122
 61
 62
           this->id near cycle[u] = this->id near cycle[nxt[u]];
                                                                                        123
 63
           this->cycle_dist[u] = this->cycle_dist[nxt[u]] + 1;
                                                                                        124
                                                                                        125
 64
 65
                                                                                        126
      }
                                                                                        127
 66
 67
      void allocate(int n, int indexed_from) {
                                                                                        128
 68
                                                                                        129
        this->cycle.resize(n + indexed_from);
 69
        this->first.resize(n + indexed_from);
                                                                                        130
 70
                                                                                        131
 71
        this->in cycle.resize(n + indexed from, false);
                                                                                        132
 72
        this->cycle_id.resize(n + indexed_from, -1);
                                                                                        133
 73
        this->id in cycle.resize(n + indexed from, -1);
                                                                                        134
 74
        this->near in cycle.resize(n + indexed from);
                                                                                        135
 75
        this->id_near_cycle.resize(n + indexed_from);
                                                                                        136
 76
        this->cycle dist.resize(n + indexed from);
                                                                                        137
 77
                                                                                        138
 78
                                                                                        139
 79
     public:
                                                                                        140
 80
      Functional_Graph(int n, int indexed_from, vector<int> &nxt, vector<int>
                                                                                        141
                                                                                        142
81
        this->allocate(n, indexed_from);
                                                                                        143
 82
        this->build(n, indexed_from, nxt, in_degree);
                                                                                        144
 83
                                                                                        145
 84
                                                                                        146
 85
      // THE CYCLES ARE ALWAYS INDEXED BY ZERO!
                                                                                        147
 86
 87
      // number of cycles
                                                                                        148
 88
      int cycle_cnt = 0;
                                                                                        149
      // Vertices present in the i-th cycle.
 89
                                                                                        150
 90
      vector<vector<int>> cycle;
                                                                                        151
 91
      // first vertex of the i-th cycle
                                                                                        152
 92
      vector<int> first;
                                                                                        153
 93
                                                                                        154
      // The i-th vertex is present in any cycle?
                                                                                        155
 94
 95
      vector<bool> in cycle;
                                                                                        156
      // id of the cycle that the vertex belongs. -1 if it doesn't belong to any
                                                                                        157
        cycle.
                                                                                        158
97
      vector<int> cycle_id;
                                                                                        159
 98
      // Represents the id of the cycle of the i-th vertex. -1 if it doesn't
                                                                                        160
        belong to any cycle.
                                                                                        161
99
      vector<int> id_in_cycle;
                                                                                        162
      // Represents the id of the nearest vertex present in a cycle.
                                                                                        163
100
101
      vector<int> near_in_cycle;
                                                                                        164
      // Represents the id of the nearest cycle.
                                                                                        165
```

```
103 | vector<int> id_near_cycle;
      // Distance to the nearest cycle.
      vector<int> cvcle dist:
      // Represent the id of the component of the vertex.
      // Equal to id near cycle
     vector<int> &comp = id_near_cycle;
111 class Functional Graph {
     // FOR UNDIRECTED GRAPH
      void compute_cycle(int u, vector<int> &nxt, vector<bool> &vis,
        vector<vector<int>> &adj) {
        int id_cycle = cycle_cnt++;
        int cur_id = 0;
        this->first[id cycle] = u;
        while(!vis[u]) {
          vis[u] = true;
          this->cvcle[id cvcle].push back(u);
          nxt[u] = find nxt(u, vis, adj);
          if(nxt[u] == -1)
            nxt[u] = this->first[id_cycle];
          this->in_cycle[u] = true;
          this->cycle_id[u] = id_cycle;
          this->id_in_cycle[u] = cur_id;
          this->near_in_cycle[u] = u;
          this->id_near_cycle[u] = id_cycle;
          this->cycle_dist[u] = 0;
          u = nxt[u];
          cur id++;
      int find_nxt(int u, vector<bool> &vis, vector<vector<int>> &adj) {
        for(int v: adj[u])
          if(!vis[v])
            return v:
        return -1;
      // Time Complexity: O(V + E)
      void build(int n, int indexed_from, vector<int> &degree,
        vector<vector<int>> &adj) {
        queue<int> q;
        vector<bool> vis(n + indexed_from, false);
        vector<int> nxt(n + indexed from);
        for(int i = indexed_from; i < n + indexed_from; i++) {</pre>
          if(adj[i].size() == 1) {
            q.push(i);
            vis[i] = true;
        vector<int> process_order;
        process_order.reserve(n + indexed_from);
        while(!q.empty()) {
          int u = q.front();
          q.pop();
          process_order.push_back(u);
```

```
nxt[u] = find_nxt(u, vis, adj);
167
          if(--degree[nxt[u]] == 1) {
168
            q.push(nxt[u]);
169
            vis[nxt[u]] = true;
170
171
172
173
        int cycle_cnt = 0;
        for (int i = indexed_from; i < n + indexed_from; i++)</pre>
174
175
          if(!vis[i])
            compute_cycle(i, nxt, vis, adj);
176
177
178
        for (int i = (int) process order.size() - 1; i \ge 0; i--) {
179
          int u = process_order[i];
180
181
          this->near_in_cycle[u] = this->near_in_cycle[nxt[u]];
          this->id_near_cycle[u] = this->id_near_cycle[nxt[u]];
182
183
          this->cycle_dist[u] = this->cycle_dist[nxt[u]] + 1;
184
185
186
187
      void allocate(int n, int indexed from) {
188
        this->cvcle.resize(n + indexed from);
189
        this->first.resize(n + indexed_from);
190
191
        this->in_cycle.resize(n + indexed_from, false);
192
        this->cycle_id.resize(n + indexed_from, -1);
193
        this->id_in_cycle.resize(n + indexed_from, -1);
194
        this->near_in_cycle.resize(n + indexed_from);
195
        this->id_near_cycle.resize(n + indexed_from);
196
        this->cycle_dist.resize(n + indexed_from);
197
198
199
     public:
200
      Functional_Graph(int n, int indexed_from, vector<int> degree,
        vector<vector<int>> &adj)
201
        this->allocate(n, indexed_from);
202
        this->build(n, indexed_from, degree, adj);
203
204
      // THE CYCLES ARE ALWAYS INDEXED BY ZERO!
205
206
207
      // number of cycles
208
      int cvcle cnt = 0;
      // Vertices present in the i-th cycle.
210
      vector<vector<int>> cycle;
211
      // first vertex of the i-th cycle
      vector<int> first;
212
213
214
      // The i-th vertex is present in any cycle?
215
      vector<bool> in_cycle;
216
      // id of the cycle that the vertex belongs. -1 if it doesn't belong to any
      vector<int> cycle_id;
      // Represents the id of the cycle of the i-th vertex. -1 if it doesn't
        belong to any cycle.
      vector<int> id_in_cycle;
      // Represents the id of the nearest vertex present in a cycle.
220
221
      vector<int> near_in_cycle;
222
      // Represents the id of the nearest cycle.
223
      vector<int> id_near_cycle;
224
      // Distance to the nearest cycle.
225
      vector<int> cycle_dist;
226
      // Represent the id of the component of the vertex.
      // Equal to id_near_cycle
```

```
228 | vector<int> &comp = id_near_cycle;
229 | };
```

5.15. Hld

```
1 class HID {
    private:
    int n;
     // number of nodes below the i-th node
     vector<int> sz:
    private:
8
     int get_sz(const int u, const int p, const vector<vector<int>> &adj) {
9
       this->sz[u] = 1;
10
        for(const int v: adj[u]) {
         if(v == p)
11
12
           continue;
13
          this->sz[u] += this->get_sz(v, u, adj);
14
15
       return this->sz[u];
16
17
18
     void dfs(const int u, const int id, const int p, const vector<vector<int>>>
        &adj) {
19
       this->chain_id[u] = id;
       this->id_in_chain[u] = chain_size[id];
20
21
       this->parent[u] = p;
22
23
        if(this->chain head[id] == -1)
24
         this->chain_head[id] = u;
25
        this->chain size[id]++;
26
27
       int maxx = -1, idx = -1;
28
        for(const int v: adj[u]) {
29
         if(v == p)
30
           continue;
31
         if(sz[v] > maxx) {
32
           maxx = sz[v];
33
           idx = v:
34
35
36
37
       if(idx != -1)
38
         this->dfs(idx, id, u, adj);
39
40
        for(const int v: adj[u]) {
         if(v == idx \mid | v == p)
41
42
43
          this->dfs(v, this->number_of_chains++, u, adj);
44
45
46
    public:
48
    /// Builds the paths.
49
     ///
     /// Time Complexity: O(n)
50
     HLD(const vector<vector<int>> &adj) {
52
       this->n = adj.size();
53
       this->chain_head.resize(this->n + 1, -1);
54
       this->id_in_chain.resize(this->n + 1, -1);
55
       this->chain id.resize(this->n + 1, -1);
56
       this->sz.resize(this->n + 1);
57
       this->chain_size.resize(this->n + 1);
58
       this->parent.resize(this->n + 1, -1);
```

```
this->get_sz(1, -1, adj);
       this->dfs(1, 0, -1, adj);
60
61
62
63
     // the chains are indexed from 0
64
     int number_of_chains = 1;
65
     // topmost node of the chain
     vector<int> chain head;
66
     // id of the i-th node in his chain
68
     vector<int> id in chain:
     // id of the chain that the i-th node belongs
70
     vector<int> chain_id;
71
     // size of the i-th chain
    vector<int> chain size;
72
    // parent of the i-th node, -1 for root
73
    vector<int> parent;
75
```

5.16. Kruskal + Dsu

```
class DSU {
   public:
     vector<int> root:
     vector<int> sz;
     DSU(int n) {
       this->root.resize(n + 1);
       iota(this->root.begin(), this->root.begin() + n + 1, 0);
10
11
       this->sz.resize(n + 1, 1);
12
13
     int Find(int x) {
14
       if(this->root[x] == x)
15
16
         return x;
17
       return this->root[x] = this->Find(this->root[x]);
18
19
20
     bool Union(int p, int q) {
21
22
       p = this->Find(p), q = this->Find(q);
23
24
       if(p == q)
25
         return false;
26
27
       if(this->sz[p] > this->sz[q]) {
28
         this->root[a] = p:
29
         this->sz[p] += this->sz[q];
30
       } else {
31
         this->root[p] = q;
32
         this->sz[q] += this->sz[p];
33
34
35
       return true;
36
37
38
39
40 struct edge {
    int u, v, w;
42
     edge() {}
43
     edge(int u, int v, int w) : u(u), v(v), w(w) {}
44
```

```
bool operator<(const edge &a) const {</pre>
46
        return w < a.w;</pre>
47
48 };
49
50 int kruskal(int n, vector<edge>& edges) {
51
52
     DSU dsu(n);
53
54
     sort(edges.begin(), edges.end());
55
56
     int weight = 0;
     for(int i = 0; i < (int)edges.size(); i++) {</pre>
57
58
       if (dsu.Union(edges[i].u, edges[i].v)) {
59
          weight += edges[i].w;
60
61
62
63
     // returns weight of mst
    return weight:
```

5.17. Lca

```
1 // #define DIST
2 // #define COST
3 /// UNCOMMENT ALSO THE LINE BELOW FOR COST!
5 class LCA {
   private:
    int n:
    // INDEXED from 0 or 1??
     int indexed from;
    /// Store all log2 from 1 to n
10
     vector<int> lq;
11
     // level of the i-th node (height)
12
13
     vector<int> level;
14
     \ensuremath{/\!/} matrix to store the ancestors of each node in power of 2 levels
15
     vector<vector<int>> anc:
16
17
     #ifdef DIST
18
       vector<int> dist:
19
      #endif
20
     #ifdef COST
21
       // int NEUTRAL VALUE = -INF; // MAX COST
        // int combine(const int a, const int b) {return max(a, b);}
22
       // int NEUTRAL VALUE = INF; // MIN COST
23
24
       // int combine(const int a, const int b) {return min(a, b);}
       vector<vector<int>> cost;
25
     #endif
26
27
28
    private:
2.9
     void allocate() {
30
       // initializes a matrix [n][lq n] with -1
31
       this->build log array();
32
       this->anc.resize(n + 1, vector<int>(lg[n] + 1, -1));
       this->level.resize(n + 1, -1);
33
34
35
        #ifdef DIST
         this->dist.resize(n + 1, 0);
36
37
        #endif
38
        #ifdef COST
39
         this->cost.resize(n + 1, vector<int>(lq[n] + 1, NEUTRAL_VALUE));
40
```

```
103 l
 42
                                                                                         104
      void build_log_array() {
 43
                                                                                         105
                                                                                                // go up k levels from x
 44
        this->lq.resize(this->n + 1);
                                                                                         106
                                                                                                int lca_go_up(int x, int k) {
 45
                                                                                         107
                                                                                                  for (int i = 0; k > 0; i++, k >>= 1)
 46
        for (int i = 2; i <= this->n; i++)
                                                                                         108
                                                                                                    if(k & 1) {
 47
           this->lq[i] = this->lq[i/2] + 1;
                                                                                         109
                                                                                                      x = this - > anc[x][i]:
                                                                                                      if(x == -1)
 48
                                                                                         110
 49
                                                                                         111
                                                                                                        return -1:
      void build_anc() {
 50
                                                                                         112
        for(int j = 1; j < anc.front().size(); j++)
for(int i = 0; i < anc.size(); i++)</pre>
 51
                                                                                         113
 52
                                                                                         114
                                                                                                  return x;
 53
             if(this->anc[i][j - 1] != -1) {
                                                                                         115
               this->anc[i][j] = this->anc[this->anc[i][j - 1]][j - 1];
 54
                                                                                         116
 55
                                                                                                #ifdef COST
               #ifdef COST
                                                                                         117
 56
                 this->cost[i][j] = combine(this->cost[i][j - 1],
                                                                                         118
                                                                                                /// Ouery between the an ancestor of v (p) and v. It returns the
         this->cost[anc[i][j - 1]][j - 1]);
                                                                                         119
                                                                                                /// max/min edge between them.
 57
               #endif
                                                                                         120
                                                                                                int lca_query_cost_in_line(int v, int p) {
 58
                                                                                         121
                                                                                                  assert(this->level[v] >= this->level[p]);
 59
                                                                                         122
                                                                                         123
                                                                                                  int k = this->level[v] - this->level[p];
 60
      void build weighted(const vector<vector<pair<int, int>>> &adj) {
                                                                                         124
                                                                                                  int ans = NEUTRAL VALUE;
 61
 62
         this->dfs LCA weighted(this->indexed from, -1, 1, 0, adj);
                                                                                         125
                                                                                                  for(int i = 0; k > 0; i++, k >>= 1)
 63
                                                                                         126
         this->build anc();
                                                                                         127
 64
                                                                                                    if(k & 1) {
                                                                                                      ans = combine(ans, this->cost[v][i]);
 65
                                                                                         128
 66
                                                                                         129
                                                                                                      v = this->anc[v][i];
      void dfs_LCA_weighted(const int u, const int p, const int l, const int d,
 67
                                                                                         130
         const vector<vector<pair<int, int>>> &adj) {
                                                                                         131
 68
         this->level[u] = 1;
                                                                                         132
                                                                                                  return ans;
 69
         this->anc[u][0] = p;
                                                                                         133
 70
         #ifdef DIST
                                                                                         134
                                                                                                #endif
 71
           this->dist[u] = d;
                                                                                         135
 72
         #endif
                                                                                         136
                                                                                                int get lca(int a, int b) {
 73
                                                                                         137
                                                                                                  // a is below b
 74
         for(const pair<int, int> &x: adj[u]) {
                                                                                         138
                                                                                                  if(this->level[b] > this->level[a])
 75
           int v = x.first, w = x.second;
                                                                                         139
                                                                                                    swap(a,b);
 76
           if(v == p)
                                                                                         140
 77
             continue;
                                                                                                  const int logg = lg[this->level[a]];
                                                                                         141
 78
           #ifdef COST
                                                                                         142
 79
             this->cost[v][0] = w:
                                                                                         143
                                                                                                  // putting a and b in the same level
                                                                                                  for(int i = logg; i >= 0; i--)
 80
                                                                                         144
           this->dfs_LCA_weighted(v, u, l + 1, d + w, adj);
                                                                                                    if(this->level[a] - (1 << i) >= this->level[b])
 81
                                                                                         145
                                                                                                      a = this->anc[a][i];
 82
                                                                                         146
 83
                                                                                         147
 84
                                                                                         148
                                                                                                  if(a == b)
 85
      void build unweighted(const vector<vector<int>> &adj) 
                                                                                         149
                                                                                                    return a;
         this->dfs_LCA_unweighted(this->indexed_from, -1, 1, 0, adj);
 86
                                                                                         150
                                                                                                  for(int i = logg; i >= 0; i--)
 87
                                                                                         151
                                                                                                    if (this->anc[a][i] != -1 && this->anc[a][i] != this->anc[b][i]) {
 88
        this->build anc();
                                                                                         152
 89
                                                                                         153
                                                                                                      a = this->anc[a][i];
 90
                                                                                         154
                                                                                                      b = this->anc[b][i];
 91
      void dfs_LCA_unweighted(const int u, const int p, const int l, const int
                                                                                         155
         d, const vector<vector<int>> &adj) {
                                                                                         156
 92
         this->level[u] = 1;
                                                                                                  return anc[a][0]:
                                                                                         157
 93
         this->anc[u][0] = p;
                                                                                         158
 94
         #ifdef DIST
                                                                                         159
 95
          this->dist[u] = d:
                                                                                         160
                                                                                               public:
 96
         #endif
                                                                                         161
                                                                                               /// Builds an weighted graph.
 97
                                                                                         162
 98
         for(const int v: adj[u]) {
                                                                                         163
                                                                                                /// Time Complexity: O(n*log(n))
 99
          if(v == p)
                                                                                         164
                                                                                                explicit LCA(const vector<vector<pair<int, int>>> &adj, const int
100
                                                                                                  indexed_from) {
             continue;
101
           this->dfs_LCA_unweighted(v, u, 1 + 1, d + 1, adj);
                                                                                         165
                                                                                                  this->n = adj.size();
102
                                                                                         166
                                                                                                  this->indexed from = indexed from;
```

```
167
        this->allocate();
168
169
        this->build_weighted(adj);
170
171
172
      /// Builds an unweighted graph.
173
      /// Time Complexity: O(n*log(n))
174
      explicit LCA(const vector<vector<int>> &adj, const int indexed_from) {
175
176
        this->n = adj.size();
177
        this->indexed_from = indexed_from;
178
        this->allocate();
179
180
        this->build_unweighted(adj);
181
182
183
      /// Goes up k levels from v. If it passes the root, returns -1.
184
      ///
185
      /// Time Complexity: O(log(k))
186
      int go_up(const int v, const int k) {
        assert(indexed_from <= v); assert(v < this->n + indexed_from);
187
188
        return this->lca_go_up(v, k);
189
190
191
192
      /// Returns the parent of v in the LCA dfs from 1.
193
194
      /// Time Complexity: O(1)
195
      int parent(int v) {
196
        assert(indexed_from <= v); assert(v < this->n + indexed_from);
197
198
        return this->anc[v][0];
199
200
201
      /// Returns the LCA of a and b.
202
      /// Time Complexity: O(log(n))
203
204
      int query_lca(const int a, const int b) {
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
205
        indexed_from);
206
207
        return this->get_lca(a, b);
208
209
210
      /// Returns the distance from a to b. When the graph is unweighted, it is
      /// 1 as the weight of the edges.
213
      /// Time Complexity: O(log(n))
214
      int query_dist(const int a, const int b) {
215
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
216
        indexed from);
217
        return this->dist[a] + this->dist[b] - 2*this->dist[this->get_lca(a, b)];
218
219
220
      #endif
221
222
      #ifdef COST
223
      /// Returns the max/min weight edge from a to b.
224
225
      /// Time Complexity: O(log(n))
226
      int query_cost(const int a, const int b) {
227
        assert(indexed_from <= min(a, b)); assert(max(a, b) < this->n +
        indexed_from);
```

5.18. Maximum Path Unweighted Graph

```
1 /// Returns the maximum path between the vertices 0 and n - 1 in a
        unweighted graph.
3 /// Time Complexity: O(V + E)
4 int maximum_path(int n) {
     vector<int> top_order = topological_sort(n);
     vector<int> pai(n, -1);
     if(top_order.empty())
       return -1;
10
     vector<int> dp(n);
     dp[0] = 1;
12
     for(int u: top_order)
13
       for(int v: adj[u])
14
         if (dp[u] \&\& dp[u] + 1 > dp[v]) {
            dp[v] = dp[u] + 1;
15
16
           pai[v] = u;
17
18
19
      if(dp[n - 1] == 0)
20
       return -1;
21
22
     vector<int> path;
23
     int cur = n - 1;
     while (cur !=-1)
24
25
       path.pb(cur);
26
       cur = pai[cur];
27
     reverse(path.begin(), path.end());
     // cout << path.size() << endl;</pre>
31
     // for(int x: path) {
          cout << x + 1 << ' ';
32
33
     // }
34
     // cout << endl;
35
36
     return dp[n - 1];
37
```

5.19. Number Of Different Spanning Trees In A Complete Graph

```
Cayley's formula
n ^ (n - 2)
```

5.20. Number Of Ways To Make A Graph Connected

```
1 s_{1} s_{2} s_{3} () s_{k} (n ^ (k 2))
2 n = number of vertices
3 s_{i} = size of the i-th connected component
4 k = number of connected components
```

5.21. Pruffer Decode

```
// IT MUST BE INDEXED BY 0.
   /// Returns the adjacency matrix of the decoded tree.
3
   ///
   /// Time Complexity: O(V)
   vector<vector<int>> pruefer decode(const vector<int> &code) {
     int n = code.size() + 2;
     vector<vector<int>> adj = vector<vector<int>> (n, vector<int>());
9
     vector<int> degree(n, 1);
10
     for (int x : code)
11
       degree[x]++;
12
13
     int ptr = 0;
14
     while (degree[ptr] > 1)
15
       ++ptr;
16
17
     int nxt = ptr;
18
     for (int u : code) {
19
       adj[u].push_back(nxt);
20
       adj[nxt].push back(u);
21
22
       if (--degree[u] == 1 && u < ptr)</pre>
23
         nxt = u;
24
       else {
25
         while (degree[++ptr] > 1)
26
27
         nxt = ptr;
28
29
30
     adj[n - 1].push_back(nxt);
     adj[nxt].push_back(n - 1);
32
33
     return adj;
```

5.22. Pruffer Encode

```
void dfs(int v, const vector<vector<int>> &adj, vector<int> &parent) {
     for (int u : adj[v]) {
3
       if (u != parent[v]) {
4
         parent[u] = v;
5
         dfs(u, adj, parent);
6
   // IT MUST BE INDEXED BY 0.
11
   /// Returns prueffer code of the tree.
12 ///
13 /// Time Complexity: O(V)
14 vector<int> pruefer_code(const vector<vector<int>> &adj) {
15
     int n = adj.size();
16
     vector<int> parent(n);
     parent[n-1] = -1;
17
     dfs(n - 1, adj, parent);
19
     int ptr = -1;
21
     vector<int> degree(n);
     for (int i = 0; i < n; i++) {
       degree[i] = adj[i].size();
23
       if (degree[i] == 1 && ptr == -1)
24
         ptr = i;
```

```
26
27
28
     vector<int> code(n - 2);
29
     int leaf = ptr;
30
     for (int i = 0; i < n - 2; i++) {
31
       int next = parent[leaf];
32
       code[i] = next;
3.3
       if (--degree[next] == 1 && next < ptr)</pre>
34
         leaf = next;
3.5
       else {
36
         ptr++;
37
          while (degree[ptr] != 1)
38
           ptr++;
39
          leaf = ptr;
40
41
42
43
    return code;
```

5.23. Pruffer Properties

- 1 * After constructing the Prüfer code two vertices will remain. One of them is the highest vertex n1, but nothing **else** can be said about the other one.
- 2 * Each vertex appears in the Prüfer code exactly a fixed number of times its degree minus one. This can be easily checked, since the degree will get smaller every time we record its label in the code, and we remove it once the degree is 1. For the two remaining vertices this fact is also true.

5.24. Remove All Bridges From Graph

1 1. Start a DFS **and** store the leafs in an array.
2 2. Connect the first leaf vertex in the array with the one in the middle,
3 the second one **and** the middle + 1, **and** so on.

5.25. Shortest Cycle In A Graph

```
int bfs(int vt) {
2
3
     vector<int> dist(MAXN, INF);
     queue<pair<int, int>> q;
     g.emplace(vt, -1);
     dist[vt] = 0:
9
     int ans = INF;
10
      while (!q.empty()) {
11
       pair<int, int> aux = q.front();
12
       int u = aux.first, p = aux.second;
13
       q.pop();
14
15
        for (int v : adj[u]) {
16
          if (v == p)
17
            continue;
          if (dist[v] < INF)</pre>
18
19
            ans = min(ans, dist[u] + dist[v] + 1);
20
21
            dist[v] = dist[u] + 1;
22
            q.emplace(v, u);
23
```

```
26
27
     return ans;
28
29
30
   /// Returns the shortest cycle in the graph
31
32 /// Time Complexity: O(V^2)
33 int get_girth(int n) {
    int ans = INF;
     for (int u = 1; u <= n; u++)
       ans = min(ans, bfs(u));
37
     return ans;
38
```

5.26. Topological Sort

```
/// INDEXED BY ZERO
3
   /// Time Complexity: O(n)
   vector<int> topological_sort(int n) {
     vector<int> in_degree(n, 0);
6
     for (int u = 0; u < n; u++)
       for(int v: adj[u])
          in_degree[v]++;
10
     queue<int> q;
11
12
     for(int i = 0; i < n; i++)
13
       if(in_degree[i] == 0)
14
         q.push(i);
15
16
     int cnt = 0;
     vector<int> top_order;
17
18
     while(!q.empty()) {
19
       int u = q.front();
20
       q.pop();
21
22
       top_order.push_back(u);
23
       cnt++;
24
25
       for(int v: adj[u])
26
         if(--in_degree[v] == 0)
27
            q.push(v);
28
29
30
     if(cnt != n) {
       cerr << "There exists a cycle in the graph" << endl;</pre>
31
32
       return vector<int>();
33
34
35
     return top_order;
```

5.27. Tree Distance

```
vector<pair<int, int>> sub(MAXN, pair<int, int>(0, 0));

void subu(int u, int p) {
  for (const pair<int, int> x : adj[u]) {
    int v = x.first, w = x.second;
    if (v == p)
```

```
continue;
        subu(v, u);
       if (sub[v].first + w > sub[u].first) {
10
         swap(sub[u].first, sub[u].second);
         sub[u].first = sub[v].first + w;
11
12
       } else if (sub[v].first + w > sub[u].second) {
13
         sub[u].second = sub[v].first + w;
14
15
16 }
17
   /// Contains the maximum distance to the node i
   vector<int> ans(MAXN);
20
21 void dfs(int u, int d, int p) {
22
     ans[u] = max(d, sub[u].first);
23
     for (const pair<int, int> x : adj[u]) {
24
       int v = x.first, w = x.second;
       if (v == p)
25
26
         continue;
27
       if (sub[v].first + w == ans[u]) {
28
         dfs(v, max(d, sub[u].second) + w, u);
29
30
         dfs(v, ans[u] + w, u);
31
32
33
34
35 // Returns the maximum tree distance
36 int solve() {
37
    subu(0, -1);
38
     dfs(0, 0, -1);
39
    return *max_element(ans.begin(), ans.end());
40 }
```

6. Language $_S tuff$

6.1. Binary String To Int

```
int y = bitset<qtdDeBits>(stringVar).to_ulong();
Ex: x = 1010, qtdDeBits = 32;
y = bitset<32>(x).to_ulong(); // y = 10
```

6.2. Climits

```
1 LONG_MIN -> (-2^31+1) :: LONG_MAX -> (2^31-1)
2 ULONG_MAX -> (2^32-1) -> UNSIGNED
3 LLONG_MIN, LLONG_MAX, ULLONG_MAX
```

6.3. Checagem Brute Force Com Solucao

```
1 $ g++ -std=c++11 gen.cpp && ./a.out > gen.out && g++ -std=c++11 brute.cpp && (./a.out < gen.in) > brute.out && g++ -std=c++11 sol.cpp && (./a.out < gen.in) > sol.out && diff brute.out sol.out
```

6.4. Checagem De Bits

```
// OBS: SO FUNCIONA PARA INT (NAO FUNCIONA COM LONG LONG)

builtin_popcount(int) -> Número de bits ativos;

builtin_ctz(int) -> Número de zeros à direita

builtin_clz(int) -> Número de zeros à esquerda

builtin_parity(int) -> Retorna se a quantidade de uns é impar(1) ou par(0)
```

6.5. Checagem E Tranformacao De Caractere

```
#include <cctype>
isdigit(str[i]);//checa se str[i] é número
isalpha(str[i]);//checa se é uma letra
islower(str[i]);//checa minúsculo
isupper(str[i]);//checa maiúsculo
isalnum(str[i]);//checa letra ou número
tolower(str[i]);//converte para minusculo
toupper(str[i]);//converte para maiusculo
```

6.6. Conta Digitos 1 Ate N

```
int solve(int n) {

int maxx = 9, minn = 1, dig = 1, ret = 0;

for(int i = 1; i <= 17; i++) {
    int q = min(maxx, n);
    ret += max(011, (q - minn + 1) * dig);
    maxx = (maxx * 10 + 9), minn * = 10, dig++;
}

return ret;
}

return ret;
}</pre>
```

6.7. Escrita Em Arquivo

```
1 ofstream cout("output.txt");
```

6.8. Gcd

```
1 int _gcd(int a, int b) {
2    if(a == 0 || b == 0) return 0;
3    else return abs(__gcd(a,b));
4 }
```

6.9. Hipotenusa

```
1 cout << hypot(3,4); // output: 5
```

6.10. Int To Binary String

```
string s = bitset<qtdDeBits>(intVar).to_string();
Ex: x = 10, qtdDeBits = 32;
s = bitset<32>(x).to_string(); // s = 00...0001010
```

6.11. Int To String

```
int a; string b;
b = to_string(a);
```

6.12. Leitura De Arquivo

```
1 ifstream cin("input.txt");
```

6.13. Max E Min Element Num Vetor

```
int maior = *max_element(arr.begin(), arr.end());
int menor = *min_element(arr.begin(), arr.end());
// OBS: Retorna iterador
```

6.14. Permutação

```
int v[] = {1,2,3};
sort(v, v+3);
do {
   cout << v[0] << ' ' << v[1] ' ' << v[2];
} while (next_permutation(v, v+3));</pre>
```

6.15. Printf De Uma String

```
char buffer [50];
int n, a=5, b=3;
n=sprintf (buffer, "%d plus %d is %d", a, b, a+b);
printf ("[%s] is a string %d chars long\n",buffer,n);
// Output:
// [5 plus 3 is 8] is a string 13 chars long
```

6.16. Remove Repeticoes Continuas Num Vetor

```
1  // arr = {10,20,20,20,30,20,20,10}
2  it = unique(arr.begin(), arr.end());
3  // arr = {10,20,30,20,10, iterator aponta pra aqui, ...}
4  arr.resize(distance(arr.begin(), it));
5  // arr = {10,20,30,20,10}
```

6.17. Rotate (Left)

```
Passado o inicio o meio e o fim ele rotaciona de forma que o meio seja o novo inicio.

vector<int> arr(n); // 1 2 3 4 5 6 7 8 9
rotate(arr.begin(),arr.begin()+3,arr.end()); //4 5 6 7 8 9 1 2 3
```

6.18. Rotate (Right)

```
vector<int> arr(n); // 1 2 3 4 5 6 7 8 9
rotate(arr.begin(),arr.rbegin()+3,arr.rend()); //7 8 9 1 2 3 4 5 6
```

6.19. Scanf De Uma String

```
char sentence[]="Rudolph is 12 years old";
char str [20]; int i;
sscanf (sentence,"%s %** %d",str,&i);
printf ("%s -> %d\n",str,i);
// Output: Rudolph -> 12
```

6.20. Split Function

```
// SEPARA STRING POR UM DELIMITADOR
// EX: str=A-B-C split -> x = {A,B,C}

vector<string> split(const string &s, char delim) {
    stringstream ss(s);
    string item;
```

```
vector<string> tokens;
while (getline(ss, item, delim)) {
    tokens.push_back(item);
}

return tokens;

int main () {
    vector<string> x = split("cap-one-best-opinion-language", '-');
// x = {cap,one,best,opinion,language};
}
```

6.21. String To Long Long

```
string s = "0xFFFF"; int base = 16;
string::size_type sz = 0;
int ll = stoll(s,&sz,base); // ll = 65535, sz = 6;
OBS: Não precisa colocar o sz, pode colocar 0; // stoll(s,0,base);
```

6.22. Substring

```
string s = "abcdef";
s.substr(posição inicial, qtd de char(opcional));
string s2 = s.substr(3,2); // s2 = "de"
string s3 = s.substr(2); // s3 = "cdef"
```

6.23. Using The Gnu Compiler Collection (Gcc) Integer Overflow Builtins

```
<!DOCTYPE html PUBLIC "-//W3C//DTD HTML 4.01 Transitional//EN"</pre>
       "http://www.w3.org/TR/html4/loose.dtd">
   <!-- saved from
       url=(0065)https://gcc.gnu.org/onlinedocs/gcc/Integer-Overflow-Builtins.html 50 | pre.format {font-family: inherit}
   <html><!-- Copyright (C) 1988-2019 Free Software Foundation, Inc.
4
5
   Permission is granted to copy, distribute and/or modify this document
   under the terms of the GNU Free Documentation License, Version 1.3 or
   any later version published by the Free Software Foundation; with the
   Invariant Sections being "Funding Free Software", the Front-Cover
   Texts being (a) (see below), and with the Back-Cover Texts being (b)
   (see below). A copy of the license is included in the section entitled
10
11
   "GNU Free Documentation License".
12
   (a) The FSF's Front-Cover Text is:
13
14
15
   A GNU Manual
16
17
   (b) The FSF's Back-Cover Text is:
18
   You have freedom to copy and modify this GNU Manual, like GNU
19
20
        software. Copies published by the Free Software Foundation raise
        funds for GNU development. --><!-- Created by GNU Texinfo 6.3,
21
       http://www.gnu.org/software/texinfo/ --><head><meta
       http-equiv="Content-Type" content="text/html; charset=UTF-8">
   <title>Using the GNU Compiler Collection (GCC): Integer Overflow
       Builtins</title>
   <meta name="description" content="Using the GNU Compiler Collection (GCC):</pre>
       Integer Overflow Builtins">
   <meta name="keywords" content="Using the GNU Compiler Collection (GCC):</pre>
       Integer Overflow Builtins">
26 <meta name="resource-type" content="document">
```

```
27 | <meta name="distribution" content="global">
     <meta name="Generator" content="makeinfo">
     <link href="https://gcc.qnu.org/onlinedocs/gcc/index.html#Top" rel="start"</pre>
30
           title="Top">
31 <link
           href="https://gcc.gnu.org/onlinedocs/gcc/Option-Index.html#Option-Index"
           rel="index" title="Option Index">
32 | 32 | 32 | 32 | 32 | 32 | 33 | 34 | 35 | 36 | 36 | 36 | 37 | 38 | 38 | 39 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 | 30 
           rel="contents" title="Table of Contents">
           href="https://gcc.gnu.org/onlinedocs/gcc/C-Extensions.html#C-Extensions"
           rel="up" title="C Extensions">
34 <link
           href="https://gcc.gnu.org/onlinedocs/gcc/x86-specific-memory-model-extensions-for
           rel="next" title="x86 specific memory model extensions for transactional
           memory">
35 <link
           href="https://gcc.gnu.org/onlinedocs/gcc/_005f_005fatomic-Builtins.html#g_t_005f_
           rel="prev" title="__atomic Builtins">
36 <style type="text/css">
37 <!--
38 a.summary-letter {text-decoration: none}
39 blockquote.indentedblock {margin-right: 0em}
40 | blockquote.smallindentedblock {margin-right: 0em; font-size: smaller}
41 | blockquote.smallquotation {font-size: smaller}
42 div.display {margin-left: 3.2em}
43 div.example {margin-left: 3.2em}
44 div.lisp {margin-left: 3.2em}
45 div.smalldisplay {margin-left: 3.2em}
46 div.smallexample {margin-left: 3.2em}
47 div.smalllisp {margin-left: 3.2em}
48 kbd {font-style: oblique}
49 pre.display {font-family: inherit}
51 pre.menu-comment {font-family: serif}
52 pre.menu-preformatted {font-family: serif}
53 pre.smalldisplay {font-family: inherit; font-size: smaller}
54 pre.smallexample {font-size: smaller}
55 pre.smallformat {font-family: inherit; font-size: smaller}
56 pre.smalllisp {font-size: smaller}
57 | span.nolinebreak {white-space: nowrap}
58 span.roman {font-family: initial; font-weight: normal}
59 | span.sansserif {font-family: sans-serif; font-weight: normal}
60 | ul.no-bullet {list-style: none}
61 | -->
62 </style>
     <link rel="stylesheet" type="text/css" href="./Using the GNU Compiler</pre>
           Collection (GCC) _ Integer Overflow Builtins_files/qcc.css">
64
65
     </head>
66
67
     <body lang="en">
68
     <a name="Integer-Overflow-Builtins"></a>
70 <div class="header">
71
72 Next: <a
           href="https://gcc.gnu.org/onlinedocs/gcc/x86-specific-memory-model-extensions-for
           accesskey="n" rel="next">x86 specific memory model extensions for
           transactional memory</a>, Previous: <a
           href="https://gcc.gnu.org/onlinedocs/gcc/_005f_005fatomic-Builtins.html#g_t_005f_
           accesskey="p" rel="prev">__atomic Builtins</a>, Up: <a
           href="https://gcc.gnu.org/onlinedocs/gcc/C-Extensions.html#C-Extensions"
           accesskey="u" rel="up">C Extensions</a> &nbsp; [<a
```

| | <pre>href="https://gcc.gnu.org/onlinedocs/gcc/index.html#SEC_Contents" title="Table of contents" rel="contents">Contents][Index]</pre> | 109 | <pre><dt>Built-in Function: bool _builtin_sub_overflow (<var>type1</var> a, <var>type2</var> b, <var>type3</var> *res)</dt></pre> |
|------------|---|------------|---|
| 73 74 | <hr/> | 110 | <pre><dt>Built-in Function: bool _builtin_ssub_overflow</dt></pre> |
| 75 | name="Built-in-Functions-to-Perform-Arithmetic-with-Overflow-Checking"> <td>>111</td> <td><pre>(int a, int b, int *res) <dt>Built-in</dt></pre></td> | >111 | <pre>(int a, int b, int *res) <dt>Built-in</dt></pre> |
| 76 | <h3 class="section">6.56 Built-in Functions to Perform Arithmetic with Overflow Checking</h3> | | Function: bool _builtin_ssubl_overflow (long int a, long int b, long int *res) |
| 77 78 | The following built-in functions allow performing simple arithmetic | 112 | <pre><dt>Built-in Function: bool _builtin_ssubll_overflow</dt></pre> |
| 79 | operations together with checking whether the operations overflowed. | 113 | <pre>(long long int a, long long int b, long long int *res)</pre> |
| | <pre></pre> | 113 | <pre><dt>Built-in Function: bool builtin_usub_overflow</dt></pre> |
| 81 | <di>></di> | | <pre>(unsigned int a, unsigned int b, unsigned int *res)</pre> |
| 82 | <pre><dt>Built-in</dt></pre> | 114 | |
| | Function: bool builtin_add_overflow (<var>type1</var> a, <var>type2</var> b, <var>type3</var> | | Function: bool builtin_usubl_overflow (unsigned long int a, unsigned long int b, unsigned long int |
| | *res) | | *res\ |
| 83 | | 115 | |
| | Function: bool builtin_sadd_overflow | | Function: bool builtin_usubll_overflow |
| 84 | <pre>(int a, int b, int *res) <dt>Built-in</dt></pre> | | <pre>(unsigned long long int a, unsigned long long int b, unsigned long long int *res)</pre> |
| 04 | Function: bool builtin_saddl_overflow | 116 | |
| | <pre>(long int a, long int b, long int *res)</pre> | 117 | These built-in functions are similar to the add overflow checking built-in |
| 85 | | 118 | |
| | <pre>Function: bool builtin_saddll_overflow (long long int a, long long int b, long long int *res)</pre> | 119 | argument from the first one, instead of addition. |
| 86 | <pre><dt>Built-in</dt></pre> | 120 | /p> |
| | Function: bool builtin_uadd_overflow | 121 | |
| 0.7 | <pre>(unsigned int a, unsigned int b, unsigned int *res)</pre> | 122 | <dl></dl> |
| 87 | <pre><dt>Built-in Function: bool builtin_uaddl_overflow</dt></pre> | 123 124 | <pre><dr><dt>Built-in</dt></dr></pre> |
| | <pre>(unsigned long int a, unsigned long int b, unsigned long int</pre> | | Function: bool builtin_mul_overflow |
| | *res) | | (<var>type1</var> a, <var>type2</var> b, <var>type3</var> |
| 88 | <pre><dt>Built-in Function: bool builtin_uaddll_overflow</dt></pre> | 125 | <pre>*res) <dt>< name="index005f_005fbuiltin_005fsmul_005foverflow">Built-in</dt></pre> |
| | <pre>(unsigned long long int a, unsigned long long int b, unsigned long</pre> | 123 | Function: bool _builtin_smul_overflow |
| | long int *res) | | (int a, int b, int *res) |
| 89 | <dd></dd> | 126 | |
| 90 | These built-in functions promote the first two operands into infinite precision signed | | Function: bool _builtin_smull_overflow (long int a, long int b, long int *res) |
| 91 | type and perform addition on those promoted operands. The result is then | 127 | |
| | cast to the type the third pointer argument points to and stored there. | | Function: bool _builtin_smulll_overflow |
| | If the stored result is equal to the infinite precision result, the built-in functions return <code>false</code> , otherwise they return | 128 | <pre>(long long int a, long long int b, long long int *res) <dt>Built-in</dt></pre> |
| 24 | <pre><code>true</code></pre> /code>. As the addition is | 120 | Function: bool _builtin_umul_overflow |
| 95 | performed in infinite signed precision, these built-in functions have fully | | <pre>(unsigned int a, unsigned int b, unsigned int *res)</pre> |
| 96 | defined behavior for all argument values. | 129 | <pre><dt>Built-in Function: bool _builtin_umull_overflow</dt></pre> |
| 97 | | | <pre>(unsigned long int a, unsigned long int b, unsigned long int</pre> |
| 98 | The first built-in function allows arbitrary integral types for operands | | *res) |
| | and | 130 | <pre><dt>Built-in</dt></pre> |
| 99 | the result type must be pointer to some integral type other than enumerated or | | Function: bool _builtin_umulll_overflow (unsigned long long int a, unsigned long long int b, unsigned long |
| 100 | boolean type, the rest of the built-in functions have explicit integer types. | | long int *res) |
| 101 | | | <dd></dd> |
| 102 | The compiler will attempt to use hardware instructions to implement those built in functions where possible, like conditional jump on everflow | 132 | 1 1 |
| | these built-in functions where possible, like conditional jump on overflow after addition, conditional jump on carry etc. | 133 134 | |
| 105 | | 135 | |
| 106 | | 136 | CONTROL following built in functions allow the bins if simple swithout's |
| 107 108 | <dl></dl> | 137 | The following built-in functions allow checking if simple arithmetic operation |
| - | | 1 | · - |

```
138 | would overflow.
139 
140 <dl>
141 <dt><a
        name="index-005f 005fbuiltin 005fadd 005foverflow 005fp"></a>Built-in
        Function: <em>bool</em> <strong>__builtin_add_overflow_p</strong>
        <em>(<var>type1</var> a, <var>type2</var> b, <var>type3</var>
        c) </em></dt>
142
    <dt><a
        name="index-_005f_005fbuiltin_005fsub_005foverflow_005fp"></a>Built-in
        Function: <em>bool</em> <strong>__builtin_sub_overflow_p</strong>
        <em>(<var>type1</var> a, <var>type2</var> b, <var>type3</var>
        c) </em></dt>
    <dt><a
143
        name="index-_005f_005fbuiltin_005fmul_005foverflow_005fp"></a>Built-in
        Function: <em>bool</em> <strong>__builtin_mul_overflow_p</strong>
        <em>(<var>type1</var> a, <var>type2</var> b, <var>type3</var>
        c) </em></dt>
    <dd>
144
    These built-in functions are similar to
145
        <code>__builtin_add_overflow</code>,
    <code>__builtin_sub_overflow</code>, or <code>__builtin_mul_overflow</code>,
    they dont store the result of the arithmetic operation anywhere and the
    last argument is not a pointer, but some expression with integral type other
149 than enumerated or boolean type.
1.51
    The built-in functions promote the first two operands into infinite
        precision signed type
    and perform addition on those promoted operands. The result is then
    cast to the type of the third argument. If the cast result is equal to the
        infinite
    precision result, the built-in functions return <code>false</code>,
        otherwise they return <code>true</code>.
    The value of the third argument is ignored, just the side effects in the
        third argument
    are evaluated, and no integral argument promotions are performed on the last
        argument.
    If the third argument is a bit-field, the type used for the result cast has
    precision and signedness of the given bit-field, rather than precision and
        signedness
159
    of the underlying type.
160
    <q\>
    For example, the following macro can be used to portably check, at
162 compile-time, whether or not adding two constant integers will overflow,
163 and perform the addition only when it is known to be safe and not to trigger
164 | a <samp>-Woverflow</samp> warning.
165 
    <div class="smallexample">
    #define INT_ADD_OVERFLOW_P(a, b) \
167
168
       __builtin_add_overflow_p (a, b, (__typeof__ ((a) + (b))) 0)
169
170
   enum {
171
        A = INT_MAX, B = 3,
        C = INT\_ADD\_OVERFLOW\_P (A, B) ? 0 : A + B,
172
173
        D = __builtin_add_overflow_p (1, SCHAR_MAX, (signed char) 0)
174
175
    </div>
176
177
    The compiler will attempt to use hardware instructions to implement
    these built-in functions where possible, like conditional jump on overflow
   after addition, conditional jump on carry etc.
179
    </p>
180
    </dd></dl>
```

```
183
     <hr>
184 <div class="header">
185
     >
186 | Next: <a
          href="https://gcc.gnu.org/onlinedocs/gcc/x86-specific-memory-model-extensions-for
          accesskey="n" rel="next">x86 specific memory model extensions for
          transactional memory</a>, Previous: <a
          href="https://gcc.gnu.org/onlinedocs/gcc/_005f_005fatomic-Builtins.html#g_t_005f_
          accesskey="p" rel="prev">__atomic Builtins</a>, Up: <a
          href="https://gcc.gnu.org/onlinedocs/gcc/C-Extensions.html#C-Extensions"
          accesskey="u" rel="up">C Extensions</a> &nbsp; [<a
         href="https://gcc.gnu.org/onlinedocs/gcc/index.html#SEC_Contents" title="Table of contents" rel="contents">Contents</a>][<a href="https://gcc.gnu.org/onlinedocs/gcc/Option-Index.html#Option-Index"
          title="Index" rel="index">Index</a>]
187
     </div>
188
189
190
191
192
     </body></html>
```

6.24. Width

6.25. Readint

```
int readInt(){
   int a = 0;
   char c;
   while(!(c >= '0' && c <= '9')) c = getchar();
   while(c>= '0' && c <= '9') a = 10*a + (c - '0'), c = getchar();
   return a;
}</pre>
```

6.26. Using The Gnu Compiler Collection (Gcc) $_{I}ntegerOverflowBuiltins_{F}ile$

6.27. Gcc

```
/* CSS for the GCC web site.
3
      Gerald Pfeifer <gerald@pfeifer.com>
    */
             { background-color: white; color: black; }
   bodv
             { color: #0066bb; text-decoration: none;
   a:visited { color: #003399; text-decoration: none;
10 a:hover { color: darkorange; text-decoration: none; }
11
12 h1
             { color: darkslategray; text-align:center; }
13 h2
             { color: darkslategray; }
14
15 | .highlight{ color: darkslategray; font-weight:bold; }
```

```
.smaller { font-size: 80%; }
17
18
   .left
              { text-align:left; }
19 | .right
              { text-align:right; }
   .center
              { text-align:center; margin-left:auto; margin-right:auto; }
20
   .top
21
               vertical-align:top; }
22
   .middle
              { vertical-align:middle; }
2.3
24
   .width33 { width:33%; }
2.5
   .border0 { border-width:0; }
26
27
   .no-margin-top { margin-top:0;
28
   .twocolumns { column-count:2;
29
   .imgleft { margin: 5px 20px; float: left; }
30
31
   img.right { float: right; }
32
33
   td.news
                 { width: 50%; padding-right: 8px; vertical-align: top; }
34
   td.news h2
                  font-size: 1.2em; margin-top: 0; margin-bottom: 2%; }
   td.news dl
                  margin-top:0; }
35
                  color:darkslategrey; font-weight:bold; margin-top:0.3em; }
   td.news dt
   td.news dd { margin-left:3ex; margin-top:0.1em; margin-bottom:0.1em; }
   td.news .date { color:darkslategrey; font-size:90%; margin-left:0.1ex; }
39
                { width: 50%; padding-left: 12px; vertical-align: top;
41
                  border-left: #3366cc thin solid; }
   td.status h2 { font-size: 1.2em; margin-top:0; margin-bottom: 1%; }
43 td.status dl { margin-top:0; }
44 td.status .version { font-weight:bold; }
45
   td.status .regress { font-size: 80%; }
46
   td.status dd { margin-left:3ex; }
47
48
   table.nav {
49
     padding-left: 32px;
50
     border-spacing: Opt;
51
52
   table.nav td {
53
54
     border-width: 0;
55
56
57
   table.navitem {
58
     width: 100%;
59
     border-spacing: Opt:
60
61
   table.navitem tr:nth-child(1) {
     border-color: #3366cc;
64
     border-style: solid;
65
     border-width: thin;
66
     color: #f2f2f9;
67
     background-color: #0066dd;
68
     font-weight: bold;
69
   table.navitem tr:nth-child(2) {
70
71
     padding-top: 3px;
72
     padding-left: 8px;
73
     padding-bottom: 3px;
74
     background-color: #f2f2f9;
75
     font-size: smaller;
76
77
78
   div.copyright {
79
     clear: both;
     font-size: smaller;
```

```
background: #f2f2f9;
      border: 2px solid #3366cc;
 83
      border-style: solid;
8.4
      border-width: thin;
 85
      padding: 4px;
 86
 87
    div.copyright p:nth-child(3) { margin-bottom: 0; }
 88
 89
    .bold
                 { font-weight:bold; }
 90
    .boldcvan
                   font-weight:bold; color:cvan;
    .boldlime
                   font-weight:bold; color:lime;
 91
 92
    .boldmagenta { font-weight:bold; color:magenta; }
 93
    .boldred
                   font-weight:bold; color:red;
 94
    .boldgreen
                   font-weight:bold; color:green;
 95
    .boldblue
                   font-weight:bold; color:blue; }
 96
    .red
                   color:red; }
 97
                   color:green; }
    .green
    .blue
                   color:blue; }
98
99
    .blackbg
                 { color:white; background-color: #000000; }
100
    /* Ouote an e-mail. The first \langle \text{div} \rangle has the sender, the second the quote. */
    blockguote.mail div:nth-child(2) { border-left: solid blue; padding-left:
        4pt; }
103
104
    /* This comes close to , alas a bit less bordersome. */
    table.border th { border:2px solid; }
106
107
    /* C++ status tables. */
108 table.cxxstatus th, table.cxxstatus td { border: 1px solid gray; }
    table.cxxstatus td:nth-child(3) { text-align:center; }
110
    table.cxxstatus tr.separator { background: #f2f2f9; }
111
112 /* C++ Defect Report table. */
113 table.cxxdrstatus th, table.cxxdrstatus td { border: 1px solid gray; }
    table.cxxdrstatus td:nth-child(4) { text-align:center; }
    table.cxxdrstatus tr.separator { background: #f2f2f9; }
    table.cxxdrstatus { width: 65%; }
117
118
    /* Padded tables. */
    table.padding5 th, td { border: 1px solid gray; padding:5px; }
119
    .supported { background-color: lightgreen; }
121
122 | .unsupported { background-color: lightsalmon; }
123 | .other { background-color: lightgray; }
    .partial { background-color: lightyellow; }
125 | .open { color: #AAAAAA; }
    /* Online documentation. */
127
128
129 pre.smallexample {
130
     font-size: medium;
      background: #f2f2f9;
131
132
      padding: 4px;
133
      display: inline-block;
134
135
136 /* Classpath versus libgci merge status page. */
137
138
    .classpath-only
                               background-color: #FFFFAA;
139
    .libgcj-only
                               background-color: #FFFFAA;
140 .VM-specific
                               background-color: #CCCCFF;
141 .GCJ-specific
                               background-color: #CCCCFF;
142 | .needsmerge
                               background-color: #FF9090;
                               background-color: #22FF22;
143 merged
                               background-color: #22FF22;
144 .merged-expected-diff
```

```
.merged-unexpected-diff { background-color: #FF4444; }
```

7. Math

7.1. Bell Numbers

```
int bellNumber(int n) {
   int bell[n+1][n+1];
   bell[0][0] = 1;
   for (int i=1; i<=n; i++) {
      // Explicitly fill for j = 0
      bell[i][0] = bell[i-1][i-1];

      // Fill for remaining values of j
      for (int j=1; j<=i; j++)
            bell[i][j] = bell[i-1][j-1] + bell[i][j-1];
    }

return bell[n][0];
}</pre>
```

7.2. Checagem De Primalidade

```
bool isPrime(int n) {
    if (n <= 1) return false;
    if (n <= 3) return true;

    // This is checked so that we can skip
    // middle five numbers in below loop
    if (n%2 == 0 || n%3 == 0)
    return false;
    for (int i=5; i*i<=n; i += 6)
    if (n%i == 0 || n%(i+2) == 0)
    return false;
    return false;
    return true;
}</pre>
```

7.3. Combinação Nor Mod Primo

```
inv[1] = 1;
for(int i = 2; i < m; ++i)
    inv[i] = (m - (m/i) * inv[m%i] % m) % m;

factorial[0] = 1;
for (int i = 1; i <= MAXN; i++) {
    factorial[i] = factorial[i - 1] * i % m;
}

int binomial_coefficient(int n, int k) {
    return factorial[n] * inverse(factorial[k]) % m * inverse(factorial[n - k]) % m;
}</pre>
```

7.4. Combinação Ncr

```
// Returns value of Binomial Coefficient C(n, k)
int binomialCoeff(int n, int k) {
   int res = 1;
   // Since C(n, k) = C(n, n-k)
   if (k > n - k)
        k = n - k;
   // Calculate value of [n*(n-1)*---*(n-k+1)] / [k*(k-1)*---*1]
   for (int i = 0; i < k; ++i) {</pre>
```

```
9 res *= (n - i);
10 res /= (i + 1);
11 }
12 return res;
13 }
```

7.5. Compressao De Pontos

```
map<int, int> rev;
for(int x : arr) {
    sl.insert(x);

    vector<int> aux;
    for(int x : sl) aux.pb(x);
    for(int i=0; i<n; i++) {
        int id = lower_bound(aux.begin(), aux.end(), arr[i]) - aux.begin();
        rev[id] = arr[i];
        arr[i] = id;
}</pre>
```

7.6. Equacao Diofantina

```
int gcd(int a, int b, int &x, int &y) {
     if (a == 0) {
3
       x = 0; y = 1;
       return b:
    int x1, y1;
6
     int d = gcd(b%a, a, x1, y1);
     x = y1 - (b / a) * x1;
    v = x1;
10
    return d;
11
12
13
   | bool find_any_solution(int a, int b, int c, int &x0, int &y0, int &q) {
14
    g = gcd(abs(a), abs(b), x0, y0);
1.5
     if (c % q)
16
       return false:
17
     x0 *= c / g;
18
     v0 *= c / q;
19
     if (a < 0)^{T} x0 = -x0;
20
21
     if (b < 0) y0 = -y0;
22
     return true;
23 }
```

7.7. Euclides Estendido

```
1 int gcd, x, y;
   //Ax + By = gcd(A,B)
   void extendedEuclidian(int a,int b){
     if(b==0){
       qcd=a;
9
       x=1:
10
       y=0;
11
      } else{
12
       extendedEuclidian(b, a%b);
13
14
       int temp = x;
```

```
15 | x=y;

16 | y = temp - (a/b)*y;

17 | }
```

7.8. Euler Totient

```
int phi(int n) {
     int result = n;
     for (int i = 2; i * i <= n; i++) {
       if(n % i == 0) {
         while (n % i == 0)
             n /= i;
           result -= result / i;
7
8
9
1.0
11
     if(n > 1)
12
       result -= result / n;
13
     return result:
14
```

7.9. Fatoracao Multiplas Queries

```
//stor smallest prime factor for every num
   int spf[MAXN];
   // Calculating SPF (Smallest Prime Factor) for every number till MAXN.
   // Time Complexity : O(nloglogn)
   void sieve() {
     spf[1] = 1;
     for (int i=2; i<MAXN; i++)</pre>
       // marking smallest prime factor for every number to be itself.
10
     // separatelyMarking spf for every even
11
12
     // number as 2
13
     for (int i=4; i<MAXN; i+=2)</pre>
14
       spf[i] = 2;
15
     for (int i=3; i*i<MAXN; i++) {</pre>
16
17
       // checking if i is prime
18
       if (spf[i] == i) {
19
         // marking SPF for all numbers divisible by i
20
         for (int j=i*i; j<MAXN; j+=i)</pre>
           // marking spf[j] if it is not previously marked
21
22
           if (spf[j]==j)
23
              spf[j] = i;
24
25
26
   // A O(log n) function returning primefactorization
   // by dividing by smallest prime factor at every step
   vector<int> getFactorization(int x) {
30
     vector<int> ret;
31
     while (x != 1)
32
       ret.push_back(spf[x]);
33
       x = x / spf[x];
34
35
     return ret;
36
```

7.10. Fatoracao Simples

```
1 map<int, int> primeFactors(int n) {
     set<int> ret;
     while (n%2 == 0) {
       m[2]++;
       n = n/2;
8
     int sq = sqrt(n);
     for (int i = 3; i \le sq+2; i = i+2) {
10
       while (n%i == 0) {
11
         m[i]++;;
12
         n = n/i;
1.3
       /★ OBS1
14
15
     IF(N < 1E7)
       FATORE COM SPF
16
17
       */
18
19
     if (n > 2)
20
21
       m[n]++;
22
23
     return ret;
24 }
```

7.11. Inclusao-Exclusao

```
1 \mid // \mid A \cup B \cup C \mid = \mid A \mid + \mid B \mid + \mid C \mid - \mid A \cap B \mid - \mid A \cap C \mid - \mid B \cap C \mid + \mid A \cap B \cap C \mid
 3 // EXEMPLO: Quantos números de 1 a 10^9 são múltiplos de 42, 54, 137 ou 201?
    int f(vector<int> arr, int LIMIT) {
      int n = arr.size();
      int c = 0;
9
10
      for(int mask = 1; mask < (1<<n); mask++) {</pre>
11
         int lcm = 1;
12
         for(int i = 0; i < n; i++)</pre>
13
           if(mask&(1<<i))
14
             lcm=lcm*arr[i]/__qcd(lcm,arr[i]);
         // se o numero de conjutos a unir for impar entao soma
15
16
         if( builtin popcount(mask) %2 == 1)
          c += LIMIT/lcm;
17
         else // se nao subtrai
18
           c -= LIMIT/lcm;
19
20
21
22
      return LIMIT-c;
23
24
```

7.12. Numero De Fatores

```
int calcFat(vector<int> fatores) {
   int x = fatores[1];
   auto lo = lower_bound(fatores.begin(), fatores.end(), x);
   auto up = upper_bound(fatores.begin(), fatores.end(), x);
   int fat = 1;
   while(up != fatores.end()) {
     fat *= (up-lo+1);
}
```

7.13. Pollard Rho (Find A Divisor)

```
1 /* Function to calculate (base^exponent) % modulus */
   int modular pow(int base, int exponent,
3
               int modulus) {
4
     /* initialize result */
     int result = 1;
     while (exponent > 0) {
       /* if y is odd, multiply base with result */
7
8
       if (exponent & 1)
9
        result = (result * base) % modulus;
10
       /* exponent = exponent/2 */
       exponent = exponent >> 1;
11
       /* base = base * base */
12
13
       base = (base * base) % modulus;
14
15
    return result;
16
17
18
   /* method to return prime divisor for n */
19
  int PollardRho(int n) {
    /* initialize random seed */
20
    srand (time(NULL));
22
23
     /* no prime divisor for 1 */
     if (n==1) return n;
24
25
26
     /* even number means one of the divisors is 2 */
27
     if (n % 2 == 0) return 2;
28
29
     /* we will pick from the range [2, N) */
30
     int x = (rand()%(n-2))+2;
31
     int y = x;
32
33
     /* the constant in f(x).
34
     * Algorithm can be re-run with a different c
35
     * if it throws failure for a composite. */
36
     int c = (rand()%(n-1))+1;
37
38
     /★ Initialize candidate divisor (or result) ★/
39
     int d = 1;
40
41
     /* until the prime factor isn't obtained.
42
     If n is prime, return n */
43
     while (d==1) {
44
       /* Tortoise Move: x(i+1) = f(x(i)) */
45
       x = (modular_pow(x, 2, n) + c + n)%n;
46
47
       /* Hare Move: y(i+1) = f(f(y(i))) */
48
       y = (modular_pow(y, 2, n) + c + n) %n;
49
       y = (modular_pow(y, 2, n) + c + n) %n;
50
51
       /* check gcd of |x-y| and n */
52
       d = \underline{gcd(abs(x-y), n)};
53
54
       /\star retry if the algorithm fails to find prime factor
       * with chosen x and c */
```

7.14. Precomputar Combinacao Ncr

```
int C[1123][1123];
3 int mod(int n) {return n%((int)1e9+7);}
5 int nCr(int n, int k) {
    for(int i = 0; i <= n; i++) {</pre>
       for(int j = 0; j <= min(i,k); j++) {</pre>
         if(j == 0 || j == i) {
8
           C[i][j] = 1;
9
10
         } else {
11
            C[i][j] = mod(C[i-1][j-1]+C[i-1][j]);
12
13
14
15 }
```

7.15. Teorema Chines Do Resto

```
1 int inv(int a, int m) {
    int m0 = m, t, q;
    int x0 = 0, x1 = 1;
5
    if (m == 1)
6
       return 0;
8 // Apply extended Euclid Algorithm
9
    while (a > 1) {
10 // q is quotient
11
       if(m == 0)
         return INF;
12
13
       q = a / m;
14
       // m is remainder now, process same as euclid's algo
15
       m = a % m, a = t;
16
17
       t = x0;
18
       x0 = x1 - q \star x0;
19
       x1 = t;
2.0
21
    // Make x1 positive
23 if (x1 < 0)
24
      x1 += m0;
25
    // debug(x1);
26
27
    return x1;
28
29 // k is size of num[] and rem[]. Returns the smallest
```

```
30 | // number x such that:
31 // x \% num[0] = rem[0],
32 // x % num[1] = rem[1],
34 // x % num[k-2] = rem[k-1]
35 // Assumption: Numbers in num[] are pairwise coprimes
36 // (gcd for every pair is 1)
  int findMinX(int num[], int rem[], int k){
     // Compute product of all numbers
39
     int prod = 1;
40
     for (int i = 0; i < k; i++)
41
       prod *= num[i];
42
43
     // Initialize result
     int result = 0;
44
45
46
     // Apply above formula
47
     for (int i = 0; i < k; i++) {
48
       int pp = prod / num[i];
49
       // debug(pp);
50
       int iv = inv(pp, num[i]);
51
       if(iv == INF)
52
        return INF:
53
       result+=rem[i] *inv(pp,num[i]) *pp;
54
55
56
     // IF IS NOT VALID RETURN INF
     return (result % prod == 0 ? INF: result % prod);
57
58
```

7.16. Binary Exponentiation

```
int power(const int x, const int p, const int MOD = ((int)1e9 + 7)) {
    if(p == 0)
        return 1%MOD;
    if(p == 1)
        return x%MOD;
    int res = power(x, p/2, MOD);
    res = (long long)res*res%MOD;
    if(p&1)
        res = (long long)res*x%MOD;
    return res;
}
```

7.17. Divisors

```
1 | // OBS: EACH NUMBER HAS AT MOST \sqrt[3]{N} DIVISORS
   vector<int> divisors(int n) {
     vector<int> ans;
     for (int i=1; i * i <= n; i++) {</pre>
        if (n%i==0) {
          // If divisors are equal, print only one
          if (n/i == i)
            ans.pb(i);
          else // Otherwise print both
10
            ans.pb(i), ans.pb(n/i);
11
12
13
     return ans;
14
```

7.18. Matrix Exponentiation

```
1 namespace matrix {
   #define Matrix vector<vector<int>>
3 const int MOD = 1e9 + 7;
5 /// Creates an n x n identity matrix.
6 ///
7 /// Time Complexity: O(n*n)
   Matrix identity (const int n) {
     assert (n > 0);
10
11
     Matrix mat_identity(n, vector<int>(n, 0));
12
13
     for (int i = 0; i < n; i++)</pre>
14
       mat_identity[i][i] = 1;
15
     return mat_identity;
16
17
18
   /// Multiplies matrices a and b.
19
   ///
20
21
   /// Time Complexity: O(mat.size() ^ 3)
   Matrix mult (const Matrix &a, const Matrix &b) {
     assert(a.front().size() == b.size());
23
24
25
     Matrix ans(a.size(), vector<int>(b.front().size(), 0));
26
     for (int i = 0; i < ans.size(); i++)</pre>
27
        for (int j = 0; j < ans.front().size(); j++)</pre>
28
          for (int k = 0; k < a.front().size(); k++)</pre>
29
            ans[i][j] = (ans[i][j] + a[i][k] * b[k][j]) % MOD;
30
31
     return ans;
32
33
34 /// Exponentiates the matrix mat to the power of p.
35 ///
   /// Time Complexity: O((mat.size() ^ 3) * log2(p))
37
   Matrix expo(Matrix &mat, int p) {
38
     assert (p >= 0);
39
40
     Matrix ans = identity(mat.size());
41
     Matrix cur_power;
42
     cur_power.swap(mat);
43
     while (p) {
44
45
       if (p & 1)
46
         ans = mult(ans, cur_power);
47
48
       cur_power = mult(cur_power, cur_power);
49
       p >>= 1;
50
51
52
     return ans;
53
54 | }; // namespace matrix
```

7.19. Modular Inverse

```
int gcd(int a, int b) {
   if (a == 0)
    return b;
   return gcd(b%a, a);
}
```

```
int power(int x, int p, int MOD) {
     if(p == 0)
       return 1%MOD;
9
10
     if(p == 1)
11
       return x%MOD;
12
     int res = power(x, p/2, MOD);
     res = (long long) res*res%MOD;
1.3
14
     if(p&1)
1.5
       res = (long long) res*x%MOD;
16
     return res;
17
18
19
   // A*B = x (mod m)
   // B = x * A^(-1)
   // Function to find modular inverse of a under modulo m
   // Assumption: m is prime
23 int modInverse(int a, int mod) {
24
     int g = gcd(a, mod);
25
     if (q != 1)
26
         return -1;
27
     else
28
       // If a and m are relatively prime, then modulo inverse
29
       // is a^(m-2) mod m
30
       return power(a, mod-2, mod);
31
```

7.20. Sieve + Segmented Sieve

```
const int MAXN = 1e6;
2
3
   /// Contains all the primes in the segments
   vector<int> seqPrimes;
   bitset<MAXN+5> primesInSeq:
   /// smallest prime factor
   int spf[MAXN+5];
9
   vector<int> primes;
10
11
   bitset<MAXN+5> isPrime;
12
13
   void sieve(int n = MAXN + 2) {
14
15
     for(int i = 0; i <= n; i++)
16
       spf[i] = i;
17
18
     isPrime.set();
     for(int i = 2; i <= n; i++) {</pre>
19
20
       if(!isPrime[i])
21
         continue;
22
23
       for (int j = i*i; j \le n; j+=i) {
24
         isPrime[j] = false;
25
         spf[j] = min(i, spf[j]);
2.6
27
       primes.pb(i);
28
29
30
31
   vector<int> getFactorization(int x) {
     vector<int> ret;
32
     while (x != 1)
33
34
       ret.push_back(spf[x]);
       x = x / spf[x];
```

```
37
     return ret;
38 }
39
40 /// Gets all primes from 1 to r
41 void segSieve(int l, int r) {
   // primes from 1 to r
     // transferred to 0..(l-r)
4.3
     segPrimes.clear();
44
4.5
     primesInSeq.set();
     int sq = sqrt(r) + 5;
46
47
48
     for(int p: primes) {
       if(p > sq)
49
50
         break;
51
52
        for(int i = 1 - 1%p; i <= r; i += p) {
53
         if(i - 1 < 0)
54
           continue;
55
56
         // if i is less than 1e6, it could be checked in the
57
         // array of the sieve
58
         if(i >= (int)1e6 || !isPrime[i])
59
           primesInSeq[i-l] = false;
60
61
62
63
     for (int i = 0; i < r-l+1; i++) {
64
       if (primesInSeq[i])
65
         seqPrimes.pb(i+1);
66
67
```

8. Miscellaneous

8.1. 2-Sat

```
1 class SCC {
    private:
3
    // number of vertices
     int n;
     // indicates whether it is indexed from 0 or 1
     int indexed from:
7
     // reversed graph
     vector<vector<int>> trans:
10
    private:
11
     void dfs_trans(int u, int id) {
12
       comp[u] = id;
       scc[id].push_back(u);
13
14
15
       for (int v: trans[u])
16
         if (comp[v] == -1)
17
           dfs_trans(v, id);
1.8
19
     void get_transpose(vector<vector<int>>& adj) {
20
21
       for (int u = indexed from; u < this->n + indexed from; u++)
22
         for(int v: adj[u])
23
           trans[v].push_back(u);
24
25
26
     void dfs_fill_order(int u, stack<int> &s, vector<vector<int>>& adj) {
27
       comp[u] = true;
```

```
29
       for(int v: adi[u])
30
         if(!comp[v])
31
           dfs_fill_order(v, s, adj);
32
33
       s.push(u);
34
35
36
     // The main function that finds all SCCs
37
     void compute_SCC(vector<vector<int>>& adj) {
38
39
       stack<int> s;
40
       // Fill vertices in stack according to their finishing times
       for (int i = indexed_from; i < this->n + indexed_from; i++)
41
42
         if(!comp[i])
43
           dfs_fill_order(i, s, adj);
44
45
       // Create a reversed graph
46
       get_transpose(adj);
47
48
       fill(comp.begin(), comp.end(), -1);
49
50
       // Now process all vertices in order defined by stack
51
       int id = 0:
52
       while(s.empty() == false) {
53
         int v = s.top();
54
         s.pop();
55
56
         if(comp[v] == -1)
57
           dfs_trans(v, id++);
58
59
60
61
62
     // number of the component of the i-th vertex
63
     // it's always indexed from 0
     vector<int> comp;
64
65
     // the i-th vector contains the vertices that belong to the i-th scc
66
     // it's always indexed from 0
67
     vector<vector<int>> scc:
68
69
     SCC(int n, int indexed_from, vector<vector<int>>& adj) {
70
       this -> n = n:
71
       this->indexed from = indexed from:
72
       comp.resize(n + 1);
73
       trans.resize(n + 1);
74
       scc.resize(n + 1);
75
76
       this->compute_SCC(adj);
77
78
   };
79
80
   // OBS: INDEXED FROM 0
81
   class SAT {
82
83
    private:
     vector<vector<int>> adi;
84
85
     int n;
86
87
    public:
88
     vector<bool> ans;
89
90
     SAT(int n) {
91
       this->n = n;
92
       adj.resize(2*n);
```

```
93
        ans.resize(n);
94
95
      // (X v Y) = (X -> Y) (X -> Y)
96
97
      void add_or(int x, bool pos_x, int y, bool pos_y) {
98
        assert (0 \le x); assert (x \le n);
99
        assert (0 \leq y); assert (x \leq n);
100
        adj[(x << 1) ^pos_x].pb((y << 1) ^(pos_y ^1));
101
        adj[(y << 1) ^pos_y].pb((x << 1) ^(pos_x ^1));
102
103
104
      // (X xor Y) = (X v Y) (X v Y)
105
      // for this function the result is always 0 1 or 1 0
      void add_xor(int x, bool pos_x, int y, bool pos_y) {
106
107
        assert (0 \le x); assert (x \le n);
        assert(0 <= y); assert(x < n);
108
109
        add_or(x, y, pos_x, pos_y);
110
        add_or(x, y, pos_x^1, pos_y^1);
111
112
      bool check() {
113
        SCC scc(2*n, 0, adi);
114
115
        for(int i = 0; i < n; i++) {</pre>
116
117
          if(scc.comp[(i << 1) | 1] == scc.comp[(i << 1) | 0])
118
            return false;
119
          ans[i] = (scc.comp[(i << 1) | 1] < scc.comp[(i << 1) | 0]);
120
121
122
        return true;
123
124 };
```

8.2. 3Sum Problem

```
// vetor arr e valor x, a soma de três valores desse vetor deve ser igual a x
   |bool sum3(int arr[], int x, int n) {
     sort(arr, arr + n);
4
     for (int i = 0; i < n-2; i++) {
       int 1 = i+1, r = n-1;
     /★ 2SUM problem -> ponteiro que aponta para o primeiro e ultimo da
        sequencia e caso a soma for menor do que x adianta em uma casa o
        ponteiro da esquerda caso seja maior diminui em uma casa o ponteiro da
       direita */
8
       while(1 < r) {
9
         if(arr[i] + arr[l] + arr[r] == x) {
1.0
           return true;
11
         } else if(arr[i] + arr[l] + arr[r] < x)</pre>
12
           1++;
13
         else
14
           r--;
15
16
17
     return false;
18
```

8.3. Fibonacci Matrix Exponentiation

```
int fib (int n) {
  long long fib[2][2]= {{1,1},{1,0}};
  int ret[2][2]= {{1,0},{0,1}};
  int tmp[2][2]= {{0,0},{0,0}};
```

```
int i, j, k;
     while(n)
7
       if(n&1)
8
          memset(tmp,0,sizeof tmp);
9
          for(i=0; i<2; i++)
10
            for(j=0; j<2; j++)
11
              for(k=0; k<2; k++)
12
                tmp[i][j] = (tmp[i][j] + ret[i][k] * fib[k][j]);
13
          for(i=0; i<2; i++)
14
            for(j=0; j<2; j++)
              ret[i][j]=tmp[i][j];
15
16
17
       memset(tmp,0,sizeof tmp);
18
        for(i=0; i<2; i++)
19
          for(j=0; j<2; j++)
20
            for(k=0; k<2; k++)
21
             tmp[i][j] = (tmp[i][j] + fib[i][k] * fib[k][j]);
22
       for(i=0; i<2; i++)
23
          for(j=0; j<2; j++)
24
            fib[i][j]=tmp[i][j];
25
26
27
     return (ret[0][1]);
28
```

8.4. Infix To Prefix

```
int main() {
   map<char,int> prec;
   stack<char> op;
   string postfix;
   string infix;
   cin >> infix;
8
9
   prec['+'] = prec['-'] = 1;
   prec['*'] = prec['/'] = 2;
10
   prec['^'] = 3;
11
   for(int i = 0; i < infix.length(); i++) {</pre>
12
1.3
     char x = infix[i];
     if('0' <= x && x <= '9') {
14
15
       for (i; i < infix.length() && ('0' <= infix[i] && infix[i] <= '9'); i++)
16
         postfix += infix[i];
17
18
     \} else if(('a' <= x && x <= 'z') || ('A' <= x && x <= 'Z')) {
       postfix += x;
19
     } else if (x = '(')
20
21
       op.push('(');
     else if(x == ')')
22
23
       while(!op.empty() && op.top() != '(') {
24
         postfix += op.top();
25
         op.pop();
26
27
       op.pop();
28
     } else {
29
       while(!op.empty() && prec[op.top()] >= prec[x]) {
30
       postfix += op.top();
31
       op.pop();
32
33
     op.push(x);
34
35
36 while(!op.empty()) {
37 | postfix += op.top();
```

```
38 | op.pop();

39 | }

40 | cout << postfix << endl;

41 | }
```

8.5. Interval Scheduling

```
1 | 1 -> Ordena pelo final do evento, depois pelo inicio.
2 | 2 -> Vai iterando pelos eventos, se eles não tiverem horário em comum então adiciona o evento à lista.
```

8.6. Kadane (Maior Soma Num Vetor)

```
int kadane(int arr[], int 1) {
    int soma, total;
    soma = total = arr[0];

for(int i = 1; i < 1; i++) {
    soma = max(arr[i], arr[i] +soma);
    if(soma > total)
        total = soma;
}

return total;
}
```

8.7. Kadane 2D

```
2 // Program to find maximum sum subarray in a given 2D array
3 #include <stdio.h>
4 #include <string.h>
5 #include <limits.h>
6 | int mat[1001][1001]
   int ROW = 1000, COL = 1000;
   // Implementation of Kadane's algorithm for 1D array. The function
   // returns the maximum sum and stores starting and ending indexes of the
11 // maximum sum subarray at addresses pointed by start and finish pointers
12 // respectively.
   int kadane(int* arr, int* start, int* finish, int n) {
13
14
        // initialize sum, maxSum and
15
       int sum = 0, maxSum = INT MIN, i;
16
17
        // Just some initial value to check for all negative values case
        \starfinish = -1;
18
19
20
        // local variable
21
       int local_start = 0;
22
23
       for (i = 0; i < n; ++i) {
2.4
           sum += arr[i]:
25
           if (sum < 0) {
26
                sum = 0;
27
                local start = i+1;
28
29
           else if (sum > maxSum) {
30
                maxSum = sum;
31
                *start = local_start;
32
                \starfinish = i;
33
```

```
35
36
        // There is at-least one non-negative number
37
       if (\starfinish != -1)
38
           return maxSum;
39
40
       // Special Case: When all numbers in arr[] are negative
41
       maxSum = arr[0];
42
       \starstart = \starfinish = 0;
43
44
       // Find the maximum element in array
       for (i = 1; i < n; i++) {
45
           if (arr[i] > maxSum) {
46
47
               maxSum = arr[i];
               *start = *finish = i;
48
49
50
51
       return maxSum;
52
53
   // The main function that finds maximum sum rectangle in mat[][]
   int findMaxSum() {
       // Variables to store the final output
57
       int maxSum = INT_MIN, finalLeft, finalRight, finalTop, finalBottom;
58
59
       int left, right, i;
60
       int temp[ROW], sum, start, finish;
61
62
       // Set the left column
       for (left = 0; left < COL; ++left) {
63
64
            // Initialize all elements of temp as 0
65
           for (int i = 0; i < ROW; i++)
66
                temp[i] = 0;
67
68
           // Set the right column for the left column set by outer loop
           for (right = left; right < COL; ++right) {
69
70
               // Calculate sum between current left and right for every row 'i'
71
               for (i = 0; i < ROW; ++i)
72
                   temp[i] += mat[i][right];
73
74
                // Find the maximum sum subarray in temp[]. The kadane()
75
                // function also sets values of start and finish. So 'sum' is
76
                // sum of rectangle between (start, left) and (finish, right)
77
                // which is the maximum sum with boundary columns strictly as
78
                // left and right.
79
                sum = kadane(temp, &start, &finish, ROW);
                // Compare sum with maximum sum so far. If sum is more, then
81
                // update maxSum and other output values
82
               if (sum > maxSum) {
83
                   maxSum = sum;
84
85
                    finalLeft = left;
                    finalRight = right;
86
87
                    finalTop = start:
                    finalBottom = finish:
88
89
90
91
92
93
       return maxSum;
94
       // Print final values
95
       printf("(Top, Left) (%d, %d)\n", finalTop, finalLeft);
96
       printf("(Bottom, Right) (%d, %d)\n", finalBottom, finalRight);
97
       printf("Max sum is: %d\n", maxSum);
```

8.8. Oito Rainhas

```
#define N 4
   bool isSafe(int mat[N][N],int row,int col) {
     for(int i = row - 1; i >= 0; i--)
        if (mat[i][col])
5
          return false:
      for (int i = row - 1, j = col - 1; i >= 0 && j >= 0; i--, j--)
        if (mat[i][j])
8
          return false:
      for (int i = row - 1, j = col + 1; i >= 0 && j < N; i--, j++)
9
10
       if (mat[i][j])
11
          return false;
12
      return true;
13
   // inicialmente a matriz esta zerada
15 | int queen(int mat[N][N], int row = 0) {
     if(row >= N) {
        for(int i = 0; i < N; i++) {</pre>
17
          for (int j = 0; j < N; j++) {
18
19
            cout << mat[i][j] << ' ';
20
21
          cout << endl:
22
23
        cout << endl << endl;</pre>
24
        return false:
25
26
      for(int i = 0; i < N; i++) {</pre>
27
       if(isSafe(mat,row,i)) {
28
          mat[row][i] = 1;
29
          if (queen (mat, row+1))
30
            return true;
31
          mat[row][i] = 0;
32
33
34
     return false;
```

8.9. Sliding Window Minimum

```
1 // minimo num vetor arr de arr[0] ... arr[k-1], arr[1] ... arr[k], arr[2]
        ... arr[k+1]
   void swma(vector<int> arr, int k) {
     deque<ii>> window;
     for(int i = 0; i < arr.size(); i++) {</pre>
        while(!window.empty() && window.back().ff > arr[i])
7
          window.pop_back();
8
        window.pb(ii(arr[i],i));
       while(window.front().ss <= i - k)</pre>
9
10
          window.pop_front();
11
     if(i >= k)
12
       cout <<'' ';
13
     if(i - k + 1 >= 0)
15
       cout << window.front().ff;</pre>
16
17 }
```

8.10. Torre De Hanoi

```
#include <stdio.h>
2
   // C recursive function to solve tower of hanoi puzzle
   void towerOfHanoi(int n, char from rod, char to rod, char aux rod) {
       printf("\n Move disk 1 from rod %c to rod %c", from_rod, to_rod);
7
8
9
     towerOfHanoi(n-1, from_rod, aux_rod, to_rod);
10
     printf("\n Move disk %d from rod %c to rod %c", n, from_rod, to_rod);
11
     towerOfHanoi(n-1, aux_rod, to_rod, from_rod);
12
1.3
14 | int main() {
15
    int n = 4; // Number of disks
     towerOfHanoi(n, 'A', 'C', 'B'); // A, B and C are names of rods
17
18
```

8.11. Kadane (Segment Tree)

```
struct Node {
2
     int pref, suf, tot, best;
3
     Node () {}
     Node (int pref, int suf, int tot, int best) : pref(pref), suf(suf),
       tot(tot), best(best) {}
   const int MAXN = 2E5 + 10;
8
   Node tree[5*MAXN];
9
   int arr[MAXN];
10
11
   Node query (const int 1, const int r, const int i, const int j, const int
       pos) {
12
     if(l > r || l > j || r < i)
13
14
       return Node(-INF, -INF, -INF, -INF);
1.5
     if(i <= 1 && r <= j)
16
17
       return Node (tree [pos].pref, tree [pos].suf, tree [pos].tot,
       tree[posl.best);
18
19
     int mid = (1 + r) / 2;
20
     Node left = query(1, mid, i, j, 2*pos+1), right = query(mid+1, r, i, j, 2*pos+2);
21
     x.pref = max({left.pref, left.tot, left.tot + right.pref});
22
23
     x.suf = max({right.suf, right.tot, right.tot + left.suf});
     x.tot = left.tot + right.tot;
24
25
     x.best = max({left.best,right.best, left.suf + right.pref});
26
     return x;
27
28
29 // Update arr[idx] to v
30 // ITS NOT DELTA!!!
31 | void update(int 1, int r, const int idx, const int v, const int pos) {
     if(1 > r || 1 > idx || r < idx)
33
       return;
34
35
     if(l == idx && r == idx) {
36
       tree[pos] = Node(v, v, v, v);
37
       return;
38
```

```
39
40
     int mid = (1 + r)/2;
41
     update(1, mid, idx, v, 2*pos+1); update(mid+1, r, idx, v, 2*pos+2);
42
     1 = 2 * pos + 1, r = 2 * pos + 2;
     tree[pos].pref = max({tree[l].pref, tree[l].tot, tree[l].tot +
       tree[r].pref});
44
     tree[pos].suf = max({tree[r].suf, tree[r].tot, tree[r].tot + tree[l].suf});
45
     tree[pos].tot = tree[l].tot + tree[r].tot;
     tree[pos].best = max({tree[1].best,tree[r].best, tree[1].suf +
       tree[r].pref});
47
48
   void build(int 1, int r, const int pos) {
49
50
51
     if(1 == r) {
       tree[pos] = Node(arr[l], arr[l], arr[l], arr[l]);
52
53
       return:
54
55
56
     int mid = (1 + r)/2:
     build(1, mid, 2*pos+1); build(mid+1, r, 2*pos+2);
     1 = 2*pos+1, r = 2*pos+2;
58
     tree[pos].pref = max({tree[l].pref, tree[l].tot, tree[l].tot +
       tree[r].pref});
     tree[pos].suf = max({tree[r].suf, tree[r].tot, tree[r].tot + tree[l].suf});
61
     tree[pos].tot = tree[l].tot + tree[r].tot;
     tree[pos].best = max({tree[1].best,tree[r].best, tree[1].suf +
62
       tree[r].pref});
63 }
```

8.12. Point Compression

```
// map<int, int> rev;
3 /// Compress points in the array arr to the range [0..n-1].
4 ///
5 /// Time Complexity: O(n log n)
   | vector<int> compress(vector<int> &arr) {
     vector<int> aux = arr;
     sort(aux.begin(), aux.end());
     aux.erase(unique(aux.begin(), aux.end()), aux.end());
10
11
     for (size t i = 0; i < arr.size(); i++) {
       int id = lower_bound(aux.begin(), aux.end(), arr[i]) - aux.begin();
12
13
       // rev[id] = arr[i];
14
       arr[i] = id;
15
16
     return arr;
17 }
```

9. Strings

9.1. Kmp

```
vector<int> pi(const string &s) {
  int n = s.size();
  vector<int> pi(n);

int l = 0, r = 1;
  while (r < n) {
  if (s[l] == s[r]) {
    l++;
    pi[r] = 1;
}</pre>
```

```
10
          r++;
11
        } else {
12
         if (1 == 0) {
13
           pi[r] = 0;
14
           r++;
15
         } else
16
           1 = pi[1 - 1];
17
18
19
     return pi;
20
21
22
   // returns the index of first occurence of a pat in a txt
23
   int kmp(const string &txt, const string &pat) {
     int n = txt.size(), m = pat.size();
24
25
26
     int t = 0, p = 0;
     vector<int> pi_pat = pi(pat);
27
28
     // vector<int> occ;
29
30
     while (t < n) {
31
32
       if (txt[t] == pat[p]) {
33
         t++, p++;
34
         if (p == m) {
35
           return t - m;
36
            // if you want to continue searching
37
            // occ.pb(t - m);
38
            // p = pi_pat[p - 1];
39
40
       } else {
41
          if (p == 0) {
42
           t++;
43
          } else
           p = pi_pat[p - 1];
44
45
46
47
48
     return -1;
     // return occ;
49
```

9.2. Suffix Array

```
// Created by Ubiratan Neto
2
3
4
                                       Suf.
                                              lcp
   0 banana
                                     5 a
                                               0
   1 anana
               Sort the Suffixes
                                     3 ana
                                               1
  2 nana
               ---->
                                     1 anana
                                               3
8
  3 ana
                alphabetically
                                     0 banana 0
9
   4 na
                                     4 na
                                               0
  15 a
10
                                     2 nana
11
   lcp = number of characters equal prefi
12
13
14
   struct SuffixArray {
15
16
17
     vector<int> rnk,tmp,sa, sa_aux, lcp, pot, sp[22];
18
19
     int block, n;
20
```

```
string s;
22
23
      SuffixArrav() {}
24
25
      SuffixArray(string t){
26
       s = t;
27
       n = t.size():
28
        rnk.resize(n+1);
29
        for(int i=0; i<22; i++) sp[i].resize(n+1);</pre>
3.0
        pot.resize(n+1):
31
        tmp.resize(max((int)257, n+1));
        // sa stores index of first char of sufix
32
33
        sa.resize(n+1);
34
        sa aux.resize(n+1);
35
        // lcp stores value between the string and next string
36
       lcp.resize(n+1);
37
       block = 0;
38
39
40
      bool suffixcmp(int i, int i){
41
        if(rnk[i] != rnk[j]) return rnk[i] < rnk[j];</pre>
42
        i+=block, j+=block;
43
        i%=n;
44
        i%=n;
45
        return rnk[i] < rnk[j];</pre>
46
47
48
      void suffixSort(int MAX VAL){
        for (int i=0; i<=MAX_VAL; i++) tmp[i] = 0;</pre>
49
        for(int i=0; i<n; i++) tmp[rnk[i]]++;</pre>
50
51
        for(int i=1; i<=MAX_VAL; i++) tmp[i] += tmp[i-1];</pre>
52
        for(int i = n-1; i>=0; i--) {
53
            int aux = sa[i]-block;
54
            aux%=n;
55
            if (aux < 0) aux+=n;
56
            sa aux[--tmp[rnk[aux]]] = aux;
57
58
        for(int i=0; i<n; i++) sa[i] = sa_aux[i];</pre>
59
        tmp[0] = 0;
60
        for (int i=1; i < n; i++) tmp[i] = tmp[i-1] + suffixcmp(sa[i-1], sa[i]);
61
        for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
62
63
64
      void calculate() {
65
       s+='\0';
66
67
        for(int i=0; i<n; i++) {</pre>
68
          sa[i] = i;
69
          rnk[i] = s[i];
70
          tmp[i] = 0;
71
72
        suffixSort (256):
73
        block = 1;
74
        while (tmp[n-1] != n-1) {
75
          suffixSort(tmp[n-1]);
76
          block*=2;
77
78
        for(int i=0; i<n-1; i++) sa[i] = sa[i+1];</pre>
79
80
81
        for(int i=1; i<n; i++) tmp[i] = tmp[i-1] + suffixcmp(sa[i-1], sa[i]);</pre>
82
        for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
83
        s.pop_back();
84
        sa.pop_back();
85
```

```
87
       void calculate lcp() {
 88
         int last = 0;
 89
         for (int i=0; i<n; i++) {</pre>
 90
           if(rnk[i] == n-1) continue;
 91
           int x = rnk[i];
 92
           lcp[x] = max((int)0, last-1);
 93
           while (sa[x] + lcp[x] < n &  sa[x+1] + lcp[x] < n &  s[sa[x]+lcp[x]] ==
         s[sa[x+1]+lcp[x]])
 94
             lcp[x]++;
 95
 96
           last = lcp[x];
 97
 98
 99
      void build_lcp_table() {
100
        int k = 0;
101
102
         for (int j = 0; (1<<j) <= 2*n; j++) {
103
           for(; k \le n \&\& k < (1 << \dot{j}); k++) {
104
             pot[k] = j-1;
105
106
107
         for (int i=0; i<n; i++) {</pre>
           sp[0][i] = lcp[i];
108
109
         for(int i = 1; (1<<i) <= n; i++) {
110
111
           for (int j = 0; j+(1<<i) <= n; <math>j++) {
             sp[i][j] = min(sp[i-1][j], sp[i-1][j+(1<<(i-1))]);
112
113
114
115
116
117
       // to find lcp of two different sufixes starting at x and y
118
      int query_lcp(int x, int y) {
119
         if(x == y) return n - x;
120
         if(rnk[x] > rnk[v]) swap(x,v);
        int 1 = rnk[x], r = rnk[y]-1;
return min(sp[pot[r-l+1]][l], sp[pot[r-l+1]][r-(1LL<<pot[r-l+1])+1]);</pre>
121
122
123
124
125
       // needs calculate and calculate lcp first
       int number of substrings(){
126
127
         int ans = n - sa[0];
128
         for(int i=0; i<n-1; i++){
129
           int length = n - sa[i+1];
130
           ans += length - lcp[i];
131
132
         return ans;
133
134
135
       // needs calculate and calculate lcp first
136
       int lcs(string &x, string &y) {
137
         string s = x + "#" + y;
138
139
         int n = (int)s.size() - 1;
140
141
         int ans = 0;
         for (int i = 0; i < n - 1; i++) {
142
143
           int ida = sa[i];
144
           int idb = sa[i+1];
           if(ida < x.size() && idb > x.size() || ida > x.size() && idb <</pre>
145
         x.size()) {
146
             ans = max(ans, lcp[i]);
147
148
```

```
149
150
151
152 };
```

9.3. Trie - Maximum Xor Sum

```
// XOR(L,R) = XOR(1,L-1) ^ XOR(1,R)
ans= pre = 0
Trie.insert(0)
for i=1 to N:
    pre = pre XOR a[i]
    Trie.insert(pre)
    ans=max(ans, Trie.query(pre))
print ans
// a funcao query é a mesma da maximum xor between two elements
```

9.4. Trie - Maximum Xor Two Elements

1. Dada uma trie de números binários e um numero X, tente achar o número máximo que resultante da operação XOR

2

3 Ex: Para o número 10 (=(1010)2), o número que resulta no **xor** máximo é (0101)2, tente acha-lo na trie.

9.5. Z-Function

```
1 // What is 7 Array?
   // For a string str[0..n-1], Z array is of same length as string.
3 | // An element Z[i] of Z array stores length of the longest substring
4 // starting from str[i] which is also a prefix of str[0..n-1]. The
5 // first entry of Z array is meaning less as complete string is always
   // prefix of itself.
   // Example:
8 // Index
9 // 0 1 2 3 4
                               7 8 9 10 11
10 // Text
11 // a a b
               c a a b
                               x
12 // % values
13 // X 1 0 0 3 1 0 0 2 2 1 0
14 // More Examples:
15 // str = "aaaaaa"
16 // Z[] = {x, 5, 4, 3, 2, 1}
17
18 // str = "aabaacd"
19 // Z[] = {x, 1, 0, 2, 1, 0, 0}
20
21 // str = "abababab"
22 // Z[] = \{x, 0, 6, 0, 4, 0, 2, 0\}
23
24 | vector<int> z_function(const string &s) {
25
    vector<int> z(s.size());
     int 1 = -1, r = -1;
26
27
     for (int i = 1; i < s.size(); ++i) {</pre>
       z[i] = i >= r ? 0 : min(r - i, z[i - 1]);
28
29
       while (i + z[i] < s.size() & s[i + z[i]] == s[z[i]])
30
        z[i]++;
31
       if (i + z[i] > r)
32
        l = i, r = i + z[i];
33
     return z;
34
35 }
```

9.6. Aho Corasick

```
/// REQUIRES trie.cpp
2
3
   class Aho {
   private:
4
     // node of the output list
     struct Out Node {
       vector<int> str_idx;
8
       Out_Node *next = nullptr;
9
     };
10
     vector<Trie::Node *> fail;
11
12
     Trie trie;
     // list of nodes of output
13
     vector<Out_Node *> out_node;
14
     const vector<string> arr;
15
16
17
     /// Time Complexity: O(number of characters in arr)
     void build_trie() {
18
19
       const int n = arr.size();
2.0
       int node cnt = 1;
21
22
       for (int i = 0; i < n; ++i)
23
         node cnt += arr[i].size();
24
25
       out_node.reserve(node_cnt);
26
       for (int i = 0; i < node cnt; ++i)
27
         out_node.push_back(new Out_Node());
28
29
       fail.resize(node cnt);
30
       for (int i = 0; i < n; ++i) {</pre>
31
         const int id = trie.insert(arr[i]);
32
          out_node[id]->str_idx.push_back(i);
33
34
35
       this->build_failures();
36
37
38
     /// Returns the fail node of cur.
39
     Trie::Node *find_fail_node(Trie::Node *cur, char c) {
40
       while (cur != this->trie.root() && !cur->next.count(c))
         cur = fail[cur->id];
41
42
       // if cur is pointing to the root node and c is not a child
43
       if (!cur->next.count(c))
         return trie.root();
44
45
       return cur->next[c];
46
47
48
     /// Time Complexity: O(number of characters in arr)
49
     void build_failures() {
50
       queue < const Trie:: Node *> q;
51
52
       fail[trie.root()->id] = trie.root();
       for (const pair<char, Trie::Node *> v : trie.root()->next) {
53
54
         g.emplace(v.second);
55
         fail[v.second->id] = trie.root();
56
         out_node[v.second->id]->next = out_node[trie.root()->id];
57
58
59
       while (!q.empty()) {
60
         const Trie::Node *u = q.front();
61
         q.pop();
62
63
         for (const pair<char, Trie::Node *> x : u->next) {
```

```
64
            const char c = x.first;
65
            const Trie::Node *v = x.second;
66
            Trie::Node *fail_node = find_fail_node(fail[u->id], c);
67
            fail[v->id] = fail_node;
68
69
            if (!out_node[fail_node->id]->str_idx.empty())
70
              out_node[v->id]->next = out_node[fail_node->id];
71
            else
72
              out_node[v->id]->next = out_node[fail_node->id]->next;
7.3
74
            q.emplace(v);
75
76
77
78
79
      vector<vector<pair<int, int>>> aho_find_occurrences(const string &text) {
80
        vector<vector<pair<int, int>>> ans(arr.size());
81
        Trie::Node *cur = trie.root();
82
83
        for (int i = 0; i < text.size(); ++i) {</pre>
          cur = find_fail_node(cur, text[i]);
84
85
          for (Out Node *node = out node[cur->id]; node != nullptr;
86
               node = node->next)
87
            for (const int idx : node->str_idx)
88
              ans[idx].emplace_back(i - (int)arr[idx].size() + 1, i);
89
90
        return ans;
91
92
    public:
93
94
      /// Constructor that builds the trie and the failures.
95
96
      /// Time Complexity: O(number of characters in arr)
97
      Aho(const vector<string> &arr) : arr(arr) { this->build trie(); }
98
99
      /// Searches in text for all occurrences of all strings in array arr.
100
101
      /// Time Complexity: O(text.size() + number of characters in arr)
      vector<vector<pair<int, int>>> find_occurrences(const string &text) {
102
103
        return this->aho_find_occurrences(text);
104
105 };
```

9.7. Hashing

```
// OBS: CHOOSE THE OFFSET AND THE PRIMES BELOW!!
   class Hash {
     /// Prime numbers to be used in mod operations
3
     /// OBS: if you change m's size, please change the return type of both
     /// and _query methods.
5
     vector<int> m = {1000000007, 1000000009};
     // Case the alphabet goes from 'a' to 'z'.
8
     static constexpr int OFFSET = 'a';
     // Choose primes greater than the size of the alphabet.
10
     vector<int> prime = \{31, 37\};
13
     // Case the alphabet goes from 'A' to 'z'.
     // constexpr int OFFSET = 'A';
14
15
     // // Choose primes greater than the size of the alphabet.
16
     // vector<int> prime = {61, 67};
17
18
     vector<vector<int>> hash table;
```

```
vector<vector<int>> pot;
20
     // size of the string
21
     int n:
22
23 private:
     int mod(int n, int m) {
24
25
       n %= m;
2.6
       if (n < 0)
27
         n += m;
2.8
       return n;
29
30
31
     /// Time Complexity: O(1)
32
     pair<int, int> _query(const int 1, const int r) {
33
       vector<int> ans(m.size());
34
35
       if (1 == 0) {
         for (int i = 0; i < m.size(); i++)
36
37
           ans[i] = hash_table[i][r];
38
       } else {
39
         for (int i = 0; i < m.size(); i++)
40
           ans[i] =
41
               mod((hash\_table[i][r] - hash\_table[i][l - 1] * pot[i][r - l +
       1]),
42
                   m[i]);
43
44
45
       return {ans.front(), ans.back()};
46
47
48
     /// Builds the hash table and the pot table.
49
50
     /// Time Complexity: O(n)
51
     void build(string &s) {
52
       pot.resize(m.size(), vector<int>(this->n));
53
       hash table.resize(m.size(), vector<int>(this->n));
54
       // Remapping the string
55
       for (char &c : s)
         c -= OFFSET;
56
57
58
       for (int i = 0; i < m.size(); i++) {
         hash_table[i][0] = s[0];
59
60
         pot[i][0] = 1;
61
         for (int j = 1; j < this -> n; j++) {
62
           hash\_table[i][j] = (s[j] + hash\_table[i][j - 1] * prime[i]) % m[i];
63
           pot[i][j] = (pot[i][j-1] * prime[i]) % m[i];
64
65
66
67
68
   public:
     /// Constructor that is responsible for building the hash table and pot
       table.
70
     111
71
     /// Time Complexity: O(n)
72
     Hash(string s) {
73
       assert(m.size() == prime.size());
74
       this->n = s.size();
75
76
       build(s);
77
78
79
     /// Returns the hash from 1 to r.
80
     /// Time Complexity: O(1) -> Actually O(number_of_primes)
```

```
82    pair<int, int> query(const int l, const int r) {
83          assert(0 <= l), assert(l <= r), assert(r < this->n);
84          return _query(l, r);
85     }
86     };
```

9.8. Lcs K Strings

```
1 // Make the change below in SuffixArray code.
2 int MaximumNumberOfStrings;
   void build suffix array() {
     vector<pair<Rank, int>> ranks(this->n + 1);
     vector<int> arr;
8
     for (int i = 1, separators = 0; i <= n; i++)</pre>
9
       if(this->s[i] > 0) {
          ranks[i] = pair<Rank, int>(Rank((int)this->s[i] +
10
        MaximumNumberOfStrings, 0), i);
11
          this->s[i] += MaximumNumberOfStrings;
12
13
          ranks[i] = pair<Rank, int>(Rank(separators, 0), i);
14
          this->s[i] = separators;
15
          separators++;
16
17
18
     RadixSort::sort_pairs(ranks, 256 + MaximumNumberOfStrings);
19
20 }
21
22 /// Program to find the LCS between k different strings.
23 ///
24 /// Time Complexity: O(n*log(n))
25 /// Space Complexity: O(n*log(n))
26 int main() {
27
     int n;
28
29
     cin >> n;
30
31
     MaximumNumberOfStrings = n;
32
33
     vector<string> arr(n);
34
35
     int sum = 0;
36
     for(string &x: arr) {
37
       cin >> x;
38
       sum += x.size() + 1;
39
40
41
     string concat;
     vector<int> ind(sum + 1);
43
     int cnt = 0;
44
     for(string &x: arr) {
45
       if(concat.size())
46
         concat += (char)cnt;
47
       concat += x;
48
49
50
     cnt = 0;
51
     for(int i = 0; i < concat.size(); i++) {</pre>
52
       ind[i + 1] = cnt;
53
       if(concat[i] < MaximumNumberOfStrings)</pre>
54
         cnt++;
55
```

```
57
     Suffix Array say(concat);
58
     vector<int> sa = say.get_suffix_array();
59
     Sparse_Table spt(say.get_lcp());
60
61
     vector<int> freq(n);
62
     int cnt1 = 0:
63
     /// Ignore separators
64
65
     int i = n, j = n - 1;
66
     int ans = 0;
67
68
     while(true) {
69
70
       if(cnt1 == n) {
71
72
         ans = max(ans, spt.query(i, j - 1));
73
74
         int idx = ind[sa[i]];
75
          frea[idx]--:
76
          if(freq[idx] == 0)
77
           cnt1--;
78
79
        } else if(j == (int)sa.size() - 1)
80
         break;
81
        else {
82
          int idx = ind[sa[j]];
83
          freq[idx]++;
84
85
         if(freq[idx] == 1)
86
            cnt1++;
87
88
89
90
     cout << ans << endl;
91
```

9.9. Lexicographically Smallest Rotation

```
1 int booth (string &s) {
     s += s;
3
     int n = s.size();
4
     vector<int> f(n, -1);
     int k = 0;
     for (int j = 1; j < n; j++) {
       int sj = s[j];
9
        int i = f[j - k - 1];
        while (i !=-1 && sj != s[k + i + 1]) {
10
11
         if(sj < s[k + i + 1])
12
           k = j - i - 1;
13
          i = f[\tilde{i}];
14
15
        if(sj != s[k + i + 1]) {
16
         if(sj < s[k])
17
           k = j;
          f[j - \tilde{k}] = -1;
18
19
20
        else
          f[j - k] = i + 1;
21
22
23
     return k;
24
```

9.10. Manacher (Longest Palindrome)

```
1 //
       https://medium.com/hackernoon/manachers-algorithm-explained-longest-palindromic-s
3 /// Create a string containing '#' characters between any two characters.
4 string get_modified_string(string &s){
    string ret:
     for(int i = 0; i < s.size(); i++){
       ret.push back('#');
       ret.push back(s[i]);
9
10
    ret.push back('#');
11
     return ret:
12 }
13
14 /// Returns the first occurence of the longest palindrome based on the lps
       arrav.
15 ///
16 /// Time Complexity: O(n)
17 string get best (const int max len, const string &str, const vector<int>
        &lps) {
     for(int i = 0; i < lps.size(); i++) {</pre>
18
       if(lps[i] == max len) {
19
20
         string ans:
         int cnt = max_len / 2;
21
         int io = i - \overline{1};
22
23
         while (cnt)
           if(str[io] != '#') {
24
             ans += str[io];
25
26
              cnt--:
2.7
28
           io--;
29
30
         reverse(ans.begin(), ans.end());
31
         if(str[i] != '#')
32
           ans += str[i];
33
         cnt = max len / 2;
34
         io = i + 1;
35
         while (cnt) {
           if(str[io] != '#') {
36
37
              ans += str[io];
38
              cnt--;
39
40
            io++;
41
42
         return ans;
43
44
45
46
   /// Returns a pair containing the size of the longest palindrome and the
        first occurence of it.
48 ///
49 /// Time Complexity: O(n)
50 pair<int, string> manacher(string &s) {
51  int n = s.size();
     string str = get modified string(s);
     int len = (2 * n) + 1;
    //the i-th index contains the longest palindromic substring with the i-th
       char as the center
     vector<int> lps(len);
56
     int c = 0; //stores the center of the longest palindromic substring until
```

```
int r = 0; //stores the right boundary of the longest palindromic
       substring until now
     int max_len = 0;
59
     for(int i = 0; i < len; i++) {
       //get mirror index of i
60
61
       int mirror = (2 * c) - i;
62
63
       //see if the mirror of i is expanding beyond the left boundary of
       current longest palindrome at center c
64
       //if it is, then take r - i as lps[i]
65
       //else take lps[mirror] as lps[i]
       if(i < r)
66
67
         lps[i] = min(r - i, lps[mirror]);
68
69
       //expand at i
       int a = i + (1 + lps[i]);
70
       int b = i - (1 + lps[i]);
71
72
       while(a < len && b >= 0 && str[a] == str[b]) {
73
         lps[i]++;
74
         a++;
75
         b--;
76
77
78
       //check if the expanded palindrome at i is expanding beyond the right
       boundary of current longest palindrome at center c
       //if it is, the new center is i
79
80
       if(i + lps[i] > r) {
81
        c = i;
82
         r = i + lps[i];
83
84
         if(lps[i] > max_len) //update max_len
85
           max_len = lps[i];
86
87
88
89
     return make pair(max len, get best(max len, str, lps));
```

9.11. Suffix Array

```
namespace RadixSort {
2
   /// Sorts the array arr stably in ascending order.
3
   ///
4
   /// Time Complexity: O(n + max_element)
   /// Space Complexity: O(n + max_element)
   template <typename T>
   void sort(vector<T> &arr, const int max element, int (*get key) (T &),
8
             int begin = 0) {
9
     const int n = arr.size();
10
     vector<T> new_order(n);
11
     vector<int> count(max_element + 1, 0);
12
13
     for (int i = begin; i < n; i++)</pre>
14
       count[get_key(arr[i])]++;
1.5
     for (int i = 1; i <= max_element; i++)</pre>
16
17
       count[i] += count[i - 1];
18
19
     for (int i = n - 1; i >= begin; i--) {
20
       new_order[count[get_key(arr[i])] - (begin == 0)] = arr[i];
21
       count[get_key(arr[i])]--;
22
23
     arr.swap(new_order);
```

```
25 | }
26
27 /// Sorts an array by their pair of ranks stably in ascending order.
28 template <typename T> void sort_pairs(vector<T> &arr, const int rank_size) {
29
    // Sort by the second rank
     RadixSort::sort<T>(arr, rank_size, [](T &item) { return item.first.second;
3.0
       },
31
                         011):
32
3.3
     // Sort by the first rank
34
     RadixSort::sort<T>(arr, rank_size, [](T &item) { return item.first.first;
35
                         011);
36
37 | } // namespace RadixSort
38
39 /// It is indexed by 0.
40 /// Let the given string be "banana".
41 ///
42 /// 0 banana
                                          5 a
43 /// 1 anana
                   Sort the Suffixes
                                          3 ana
44 /// 2 nana
                   ----->
                                         1 anana
45 /// 3 ana
                    alphabetically
                                         0 banana
46 /// 4 na
                                         4 na
47 /// 5 a
48 /// So the suffix array for "banana" is {5, 3, 1, 0, 4, 2}
49 ///
50 /// LCP
51 ///
52 /// 1 a
53 /// 3 ana
54 /// 0 anana
55 /// 0 banana
56 /// 2 na
57 /// 0 nana (The last position will always be zero)
58 ///
59 /// So the LCP for "banana" is {1, 3, 0, 0, 2, 0}
60 ///
61 class Suffix_Array {
62 private:
63
     string s;
64
     int n;
65
     typedef pair<int, int> Rank;
67
68
   public:
69
     Suffix Array(string &s) {
70
       this->n = s.size();
71
       this->s = s;
72
       // little optimization, remove the line above
73
       // this->s.swap(s);
74
75
       this->sa = build_suffix_array();
76
       this->lcp = build_lcp();
77
78
79
   private:
80
     /// The vector containing the ranks will be present at ret
81
     void build_ranks(const vector<pair<Rank, int>> &ranks, vector<int> &ret) {
82
       ret[ranks[0].second] = 1;
83
       for (int i = 1; i < n; i++) {</pre>
84
         // If their rank are equal, than its position should be the same.
85
         if (ranks[i - 1].first == ranks[i].first)
86
           ret[ranks[i].second] = ret[ranks[i - 1].second];
87
         else
```

```
ret[ranks[i].second] = ret[ranks[i - 1].second] + 1;
 89
 90
 91
 92
      /// Builds the Suffix Array for the string s.
 93
 94
      /// Time Complexity: O(n*log(n))
      /// Space Complexity: O(n)
 95
      vector<int> build_suffix_array() {
 96
 97
        // This tuple below represents the rank and the index associated with it.
 98
        vector<pair<Rank, int>> ranks(this->n);
 99
        vector<int> arr(this->n);
100
101
        for (int i = 0; i < n; i++)
          ranks[i] = pair<Rank, int>(Rank(s[i], 0), i);
102
103
104
        RadixSort::sort_pairs(ranks, 256);
105
        build_ranks(ranks, arr);
106
107
          int jump = 1;
108
          int max rank = arr[ranks.back().second];
109
110
          // It will be compared intervals a pair of intervals (i, jump-1), (i +
           // jump, i + 2* jump - 1). The variable jump is always a power of 2.
111
112
           while (max_rank != this->n) {
            for (int i = 0; i < this->n; i++) {
113
114
               ranks[i].first.first = arr[i];
               ranks[i].first.second = (i + jump < this->n ? arr[i + jump] : 0);
115
116
              ranks[i].second = i;
117
118
            RadixSort::sort_pairs(ranks, n);
119
120
            build ranks (ranks, arr);
121
            max_rank = arr[ranks.back().second];
122
123
            jump \star = 2;
124
125
126
127
        vector<int> sa(this->n);
128
        for (int i = 0; i < this -> n; i++)
          sa[arr[i] - 1] = i;
129
130
        return sa:
131
132
133
      /// Builds the lcp (Longest Common Prefix) array for the string s.
134
      /// A value lcp[i] indicates length of the longest common prefix of the
      /// suffixes indexed by i and i + 1. Implementation of the Kasais
135
        Algorithm.
136
137
      /// Time Complexity: O(n)
138
      /// Space Complexity: O(n)
139
      vector<int> build_lcp() {
140
        lcp.resize(n, 0);
141
        vector<int> inverse suffix(this->n);
142
143
        for (int i = 0; i < this -> n; i++)
          inverse_suffix[sa[i]] = i;
144
145
146
        int k = 0;
147
148
        for (int i = 0; i < this->n; i++) {
149
          if (inverse_suffix[i] == this->n - 1) {
150
            k = 0;
151
            continue;
```

```
152
153
154
           int j = sa[inverse_suffix[i] + 1];
155
156
           while (i + k < this -> n && j + k < this -> n && s[i + k] == s[j + k])
157
            k++;
158
159
          lcp[inverse_suffix[i]] = k;
160
161
          if (k > 0)
            k--;
163
164
165
        return lcp;
166
167
    public:
168
169
      vector<int> sa;
170
      vector<int> lcp;
171
      /// LCS of two strings A and B.
173
174
      /// The string s must be initialized in the constructor as the string (A +
        '$'
175
      /// + B).
176
      ///
177
      /// The string A starts at index 1 and ends at index (separator - 1).
178
      /// The string B starts at index (separator + 1) and ends at the end of the
179
      /// string.
180
      ///
181
      /// Time Complexity: O(n)
182
      /// Space Complexity: O(1)
183
      int lcs(int separator) {
184
        assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
185
186
        int ans = 0;
187
         for (int i = 0; i + 1 < this->sa.size(); i++) {
188
          int left = this->sa[i];
189
190
           int right = this->sa[i + 1];
191
192
           if ((left < separator && right > separator) ||
193
               (left > separator && right < separator))</pre>
194
            ans = max(ans, lcp[i]);
195
196
197
        return ans:
198
199 };
```

9.12. Suffix Array Pessoa

```
1 // OBS: Suffix Array build code imported from:
       https://github.com/gabrielpessoal/Biblioteca-Maratona/blob/master/code/String/Suf
3 // Because it's faster.
5 /// It is indexed by 0.
6 /// Let the given string be "banana".
7 ///
8 /// 0 banana
                                       5 a
9 /// 1 anana
                  Sort the Suffixes
                                       3 ana
10 /// 2 nana
                  _____>
                                       1 anana
11 /// 3 ana
                   alphabetically
                                       0 banana
```

```
12 | /// 4 na
                                           4 na
13 /// 5 a
                                           2 nana
14 /// So the suffix array for "banana" is {5, 3, 1, 0, 4, 2}
15 ///
16 /// LCP
17 ///
18 /// 1 a
19 /// 3 ana
20 /// 0 anana
   /// 0 banana
2.1
   /// 2 na
22
23
   /// 0 nana (The last position will always be zero)
24
   /// So the LCP for "banana" is {1, 3, 0, 0, 2, 0}
25
26
   ///
27
   class Suffix_Array {
28
   private:
29
     string s;
30
     int n;
31
     typedef pair<int, int> Rank;
32
33
34
   public:
35
     Suffix_Array(string &s) {
36
       this->n = s.size();
37
       this->s = s;
38
       // little optimization, remove the line above
39
       // this->s.swap(s);
40
41
       this->sa = build_suffix_array();
42
       this->lcp = build_lcp();
43
44
45
   private:
46
     /// Builds the Suffix Array for the string s.
47
48
     /// Time Complexity: O(n*log(n))
49
     /// Space Complexity: O(n)
50
     vector<int> build_suffix_array() {
51
       int n = this -> s.size(), c = 0;
52
       vector<int> temp(n), posBucket(n), bucket(n), bpos(n), out(n);
53
       for (int i = 0; i < n; i++)
54
         out[i] = i;
       sort(out.begin(), out.end(),
55
56
             [&] (int a, int b) { return this->s[a] < this->s[b]; });
57
       for (int i = 0; i < n; i++) {
58
         bucket[i] = c;
59
         if (i + 1 == n \mid | this -> s[out[i]] != this -> s[out[i + 1]])
60
61
62
       for (int h = 1; h < n && c < n; h <<= 1) {
63
          for (int i = 0; i < n; i++)
           posBucket[out[i]] = bucket[i];
64
65
          for (int i = n - 1; i >= 0; i--)
           bpos[bucket[i]] = i;
66
67
          for (int i = 0; i < n; i++) {
68
           if (out[i] >= n - h)
69
             temp[bpos[bucket[i]]++] = out[i];
70
71
          for (int i = 0; i < n; i++) {
72
           if (out[i] >= h)
73
             temp[bpos[posBucket[out[i] - h]]++] = out[i] - h;
74
75
76
          for (int i = 0; i + 1 < n; i++) {
```

```
int a = (bucket[i] != bucket[i + 1]) || (temp[i] >= n - h) ||
 78
                      (posBucket[temp[i + 1] + h] != posBucket[temp[i] + h]);
 79
            bucket[i] = c;
 80
            c += a;
 81
 82
           bucket[n-1]=c++;
 83
           temp.swap(out);
 84
 85
        return out;
 86
 87
 88
      /// Builds the lcp (Longest Common Prefix) array for the string s.
       /// A value lcp[i] indicates length of the longest common prefix of the
 89
      /// suffixes indexed by i and i + 1. Implementation of the Kasais
 90
         Algorithm.
 91
      /// Time Complexity: O(n)
 92
 93
      /// Space Complexity: O(n)
 94
      vector<int> build_lcp() {
 95
        lcp.resize(n, 0);
         vector<int> inverse_suffix(this->n);
 96
 97
 98
         for (int i = 0; i < this -> n; i++)
 99
           inverse_suffix[sa[i]] = i;
100
101
         int k = 0;
102
103
         for (int i = 0; i < this -> n; i++) {
104
           if (inverse_suffix[i] == this->n - 1) {
105
            k = 0;
106
            continue;
107
108
109
           int j = sa[inverse suffix[i] + 1];
110
111
           while (i + k < this -> n \&\& j + k < this -> n \&\& s[i + k] == s[j + k])
112
            k++;
113
114
           lcp[inverse_suffix[i]] = k;
115
           if (k > 0)
116
117
            k--;
118
119
120
         return lcp;
121
122
123
    public:
124
      vector<int> sa;
125
      vector<int> lcp;
126
127
      /// LCS of two strings A and B.
128
      ///
      /// The string s must be initialized in the constructor as the string (A \pm
129
        '$'
      /// + B).
130
131
132
      /// The string A starts at index 1 and ends at index (separator - 1).
133
      /// The string B starts at index (separator + 1) and ends at the end of the
134
      /// string.
135
      ///
136
      /// Time Complexity: O(n)
137
      /// Space Complexity: O(1)
138
      int lcs(int separator) {
139
        assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
```

```
140
141
         int ans = 0;
142
143
         for (int i = 0; i + 1 < this->sa.size(); <math>i++) {
144
          int left = this->sa[i];
145
           int right = this->sa[i + 1];
146
           if ((left < separator && right > separator) ||
147
148
               (left > separator && right < separator))</pre>
149
             ans = max(ans, lcp[i]);
150
151
152
         return ans:
153
154
     };
```

57

61

63

67

79

81

82

83

84

85

87

89

91

101

102

103

104

105

107

9.13. Suffix Array With Additional Memory

```
namespace RadixSort {
   /// Sorts the array arr stably in ascending order.
   /// Time Complexity: O(n + max_element)
   /// Space Complexity: O(n + max_element)
   template <typename T>
   void sort(vector<T> &arr, const int max_element, int (*get_key)(T &),
             int begin = 0) {
8
     const int n = arr.size();
     vector<T> new_order(n);
10
11
     vector<int> count(max_element + 1, 0);
12
13
     for (int i = begin; i < n; i++)
14
       count[get_key(arr[i])]++;
15
     for (int i = 1; i <= max element; i++)
17
       count[i] += count[i - 1];
18
19
     for (int i = n - 1; i >= begin; i--) {
20
       new_order[count[get_key(arr[i])]] = arr[i];
21
       count[get_key(arr[i])]--;
22
23
24
     arr = new_order;
25
26
   /// Sorts an array by their pair of ranks stably in ascending order.
28
   template <typename T> void sort_pairs(vector<T> &arr, const int rank_size) {
29
     // Sort by the second rank
30
     RadixSort::sort<T>(
31
         arr, rank_size, [](T &item) { return item.first.second; }, 111);
32
33
     // Sort by the first rank
34
     RadixSort::sort<T>(
35
         arr, rank_size, [](T &item) { return item.first.first; }, 111);
36
37
   } // namespace RadixSort
39 /// It is indexed by 1.
40 class Suffix Array {
41 private:
42
     string s;
43
     int n;
44
45
     typedef pair<int, int> Rank;
     vector<int> suffix array;
```

```
vector<int> lcp;
      vector<vector<int>> rank table:
50
      vector<int> log_array;
51
52 public:
53
      Suffix Arrav(const string &s) {
54
        this->n = s.size();
55
        this->s = "#" + s;
56
        build log array();
58
        build_suffix_array();
59
        lcp = build lcp();
 60
 62
    private:
      vector<int> build_ranks(const vector<pair<Rank, int>> &ranks) {
        vector<int> arr(this->n + 1);
 65
 66
        arr[ranks[1].second] = 1:
        for (int i = 2; i <= n; i++) {
 68
          // If their rank are equal, than its position should be the same.
 69
          if (ranks[i - 1].first == ranks[i].first)
70
            arr[ranks[i].second] = arr[ranks[i - 1].second];
71
72
            arr[ranks[i].second] = arr[ranks[i - 1].second] + 1;
73
74
75
        return arr;
76
77
78
      /// Builds the Suffix Array for the string s.
80
      /// Time Complexity: O(n*log(n))
      /// Space Complexity: O(n*log(n))
      void build suffix array() {
        // This tuple below represents the rank and the index associated with it.
        vector<pair<Rank, int>> ranks(this->n + 1);
        vector<int> arr;
86
        int rank table size = 0;
88
        this->rank_table.resize(log_array[this->n] + 2);
90
        for (int i = 1; i <= this->n; i++)
          ranks[i] = pair<Rank, int>(Rank(s[i], 0), i);
92
93
        // Inserting only the ranks in the table.
 94
        transform(ranks.begin(), ranks.end(),
 95
                   back_inserter(rank_table[rank_table_size++]),
96
                   [] (pair<Rank, int> &pair) { return pair.first.first; });
 97
 98
        RadixSort::sort_pairs(ranks, 256);
99
        arr = build ranks(ranks);
100
          int jump = 1;
          int max rank = arr[ranks.back().second];
          // It will be compared intervals a pair of intervals (i, jump-1), (i +
106
          // jump, i + 2*jump - 1). The variable jump is always a power of 2.
          while (jump < n) {</pre>
108
            for (int i = 1; i <= this->n; i++) {
109
              ranks[i].first.first = arr[i];
110
              ranks[i].first.second = (i + jump <= this->n ? arr[i + jump] : 0);
111
              ranks[i].second = i;
```

```
113
114
             // Inserting only the ranks in the table.
115
             transform(ranks.begin(), ranks.end(),
116
                       back inserter(rank table [rank table size++]),
                       [] (pair<Rank, int> &pair) { return pair.first.first; });
117
118
119
             RadixSort::sort_pairs(ranks, n);
120
121
             arr = build ranks(ranks);
122
123
             max_rank = arr[ranks.back().second];
124
             jump *= 2;
125
126
127
           for (int i = 1; i <= n; i++) {
128
             ranks[i].first.first = arr[i];
129
             ranks[i].first.second = (i + jump <= this->n ? arr[i + jump] : 0);
130
             ranks[i].second = i;
131
132
           // Inserting only the ranks in the table.
133
           transform(ranks.begin(), ranks.end(),
134
135
                     back_inserter(rank_table[rank_table_size++]),
136
                     [](pair<Rank, int> &pair) { return pair.first.first; });
137
138
139
         this->suffix_array.resize(this->n + 1);
         for (int i = 1; i <= this->n; i++)
140
141
           this->suffix_array[arr[i]] = i;
142
143
144
       /// Builds the lcp (Longest Common Prefix) array for the string s.
145
      /// A value lcp[i] indicates length of the longest common prefix of the
146
      /// suffixes indexed by i and i + 1. Implementation of the Kasais
         Algorithm.
147
      /// Time Complexity: O(n)
148
149
      /// Space Complexity: O(n)
150
      vector<int> build lcp() {
151
         vector<int> lcp(this->n + 1, 0);
152
         vector<int> inverse_suffix(this->n + 1, 0);
153
154
         for (int i = 1; i <= n; i++)
155
          inverse_suffix[suffix_array[i]] = i;
156
157
         int k = 0;
158
         for (int i = 1; i <= n; i++) {</pre>
159
          if (inverse suffix[i] == n) {
160
            k = 0;
161
             continue;
162
163
164
           int j = suffix_array[inverse_suffix[i] + 1];
165
166
           while (i + k <= this->n && i + k <= this->n && s[i + k] == s[i + k])
167
168
            k++;
169
170
           lcp[inverse_suffix[i]] = k;
171
172
           if (k > 0)
173
            k--;
174
175
```

```
176
         return lcp;
177
178
179
      void build_log_array() {
180
        log array.resize(this->n + 1, 0);
181
182
         for (int i = 2; i <= this->n; i++)
183
           log_array[i] = log_array[i / 2] + 1;
184
185
186
     public:
187
      const vector<int> &get_suffix_array() { return suffix_array; }
188
189
      const vector<int> &get_lcp() { return lcp; }
190
191
       /// LCS of two strings A and B.
192
193
      /// The string s must be initialized in the constructor as the string (A +
         ' $'
194
      /// + B).
      ///
      /// The string A starts at index 1 and ends at index (separator - 1).
197
      /// The string B starts at index (separator + 1) and ends at the end of the
      /// string.
198
199
200
       /// Time Complexity: O(n)
201
       /// Space Complexity: O(1)
202
      int lcs(int separator) {
203
         separator++;
204
         assert(!isalpha(this->s[separator] && !isdigit(this->s[separator])));
205
206
         int ans = 0;
207
208
         for (int i = 1; i < this->n - 1; i++) {
209
           int left = this->suffix_array[i];
210
           int right = this->suffix array[i + 1];
211
212
           if ((left < separator && right > separator) ||
213
               (left > separator && right < separator))
214
             ans = max(ans, lcp[i]);
215
216
217
         return ans;
218
219
220
      /// Compares two substrings beginning at indexes i and j of a fixed length.
221
222
      /// OBS: Necessary build rank_table (uncomment build_suffix_array) and
        build
223
      /// log_array.
224
225
       /// Time Complexity: O(1)
226
       /// Space Complexity: O(1)
227
       int compare(const int i, const int j, const int length) {
228
         assert (1 <= i && i <= this->n && 1 <= j && j <= this->n);
229
         assert(!this->log_array.empty() && !this->rank_table.empty());
230
         assert(i + length - 1 <= this->n && j + length - 1 <= this->n);
231
232
         // Greatest k such that 2^k <= 1
233
         const int k = this->log array[length];
234
235
         const int jump = length - (1 << k);</pre>
236
237
         const pair<int, int> iRank = {
238
             this->rank_table[k][i],
```

```
(i + jump <= this->n ? this->rank_table[k][i + jump] : -1)};
239
240
        const pair<int, int> jRank = {
241
            this->rank_table[k][j],
242
             (j + jump <= this->n ? this->rank_table[k][j + jump] : -1)};
243
244
        return iRank == jRank ? 0 : iRank < jRank ? -1 : 1;
245
246
    };
```

56

57

59

60

61

62 6.3

64 65

66

67

68

69

70

71

72

73

74

75

76

77

78

79

80 81

82

83

85

86

87

88

89

90

91

92

93

94

95

96

97

98

100

101

106

107

108

109

110

111

112

113

114

115

116

117

9.14. Trie

```
1 class Trie {
   private:
     static const int INT LEN = 31;
     // static const int INT_LEN = 63;
  public:
6
     struct Node {
       map<char, Node *> next;
8
9
       int id:
10
       // cnt counts the number of words which pass in that node
11
       int cnt = 0:
12
       // word counts the number of words ending at that node
13
       int word cnt = 0;
14
15
       Node(const int x) : id(x) {}
16
     };
17
18
   private:
19
     int trie size = 0;
     // contains the next id to be used in a node
21
     int node cnt = 0;
     Node *trie_root = this->make_node();
22
23
24 private:
     Node *make_node() { return new Node(node_cnt++); }
26
27
     int trie_insert(const string &s) {
28
       Node *aux = this->root();
29
       for (const char c : s) {
30
         if (!aux->next.count(c))
31
           aux->next[c] = this->make_node();
32
         aux = aux->next[c];
33
         ++aux->cnt;
34
35
       ++aux->word cnt;
36
       ++this->trie size;
37
       return aux->id:
38
39
40
     void trie_erase(const string &s) {
41
       Node *aux = this->root():
42
       for (const char c : s) {
43
         Node *last = aux:
         aux = aux->next[c];
44
45
         --aux->cnt;
         if (aux->cnt == 0) {
46
47
           last->next.erase(c);
48
           aux = nullptr;
49
           break;
50
51
52
       if (aux != nullptr)
53
         --aux->word cnt;
```

```
54 l
        --this->trie size:
55
      int trie_count(const string &s) {
58
        Node *aux = this->root();
        for (const char c : s) {
          if (aux->next.count(c))
            aux = aux->next[c];
            return 0:
        return aux->word_cnt;
      int trie_query_xor_max(const string &s) {
        Node *aux = this->root();
        int ans = 0;
        for (const char c : s) {
          const char inv = (c == '0' ? '1' : '0');
          if (aux->next.count(inv)) {
            ans = (ans << 111) | (inv - '0');
            aux = aux->next[inv];
          } else {
            ans = (ans << 111) | (c - '0');
            aux = aux->next[c];
        return ans;
84
    public:
     Trie() {}
      Node *root() { return this->trie_root; }
      int size() { return this->trie size; }
      /// Returns the number of nodes present in the trie.
      int node count() { return this->node cnt; }
      /// Inserts s in the trie.
      /// Returns the id of the last character of the string in the trie.
      /// Time Complexity: O(s.size())
      int insert(const string &s) { return this->trie_insert(s); }
      /// Inserts the binary representation of x in the trie.
102
      ///
      /// Time Complexity: O(log x)
103
      int insert(const int x) {
104
        assert (x >= 0):
105
        // converting x to binary representation
        return this->trie_insert(bitset<INT_LEN>(x).to_string());
      /// Removes the string s from the trie.
      ///
      /// Time Complexity: O(s.size())
      void erase(const string &s) { this->trie erase(s); }
      /// Removes the binary representation of x from the trie.
      /// Time Complexity: O(log x)
      void erase(const int x) {
118
```

```
119
        assert (x >= 0);
        // converting x to binary representation
120
121
        this->trie_erase(bitset<INT_LEN>(x).to_string());
122
123
124
      /\!/\!/ Returns the number of maximum xor sum with x present in the trie.
125
126
      /// Time Complexity: O(log x)
127
      int query_xor_max(const int x) {
        assert(x \ge 0);
128
129
        // converting x to binary representation
        return this->trie_query_xor_max(bitset<INT_LEN>(x).to_string());
130
131
132
133
      /// Returns the number of strings equal to s present in the trie.
134
135
      /// Time Complexity: O(s.size())
     int count(const string &s) { return this->trie_count(s); }
136
137
    } ;
```