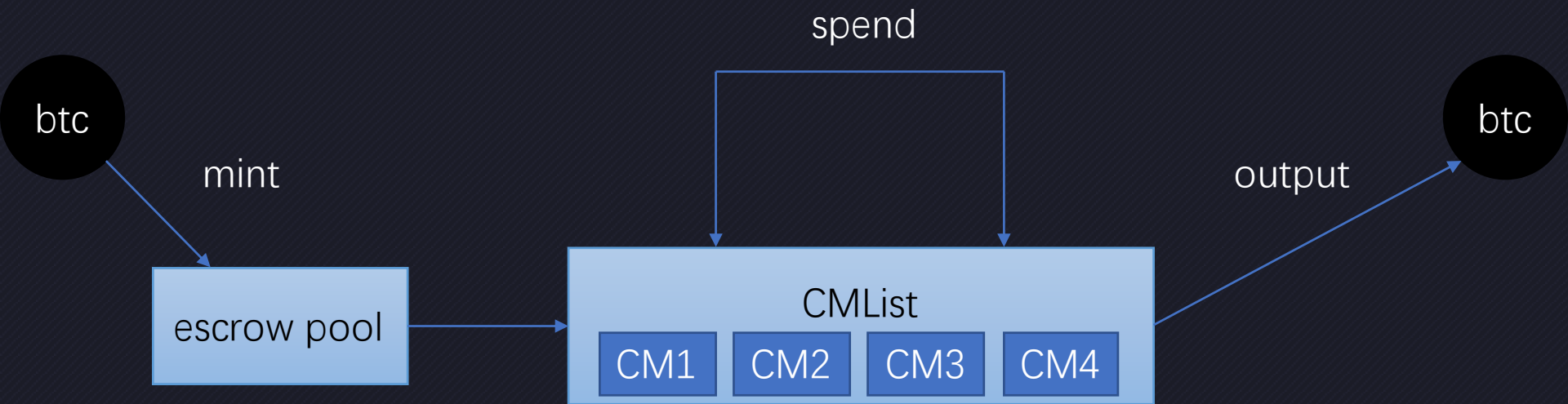




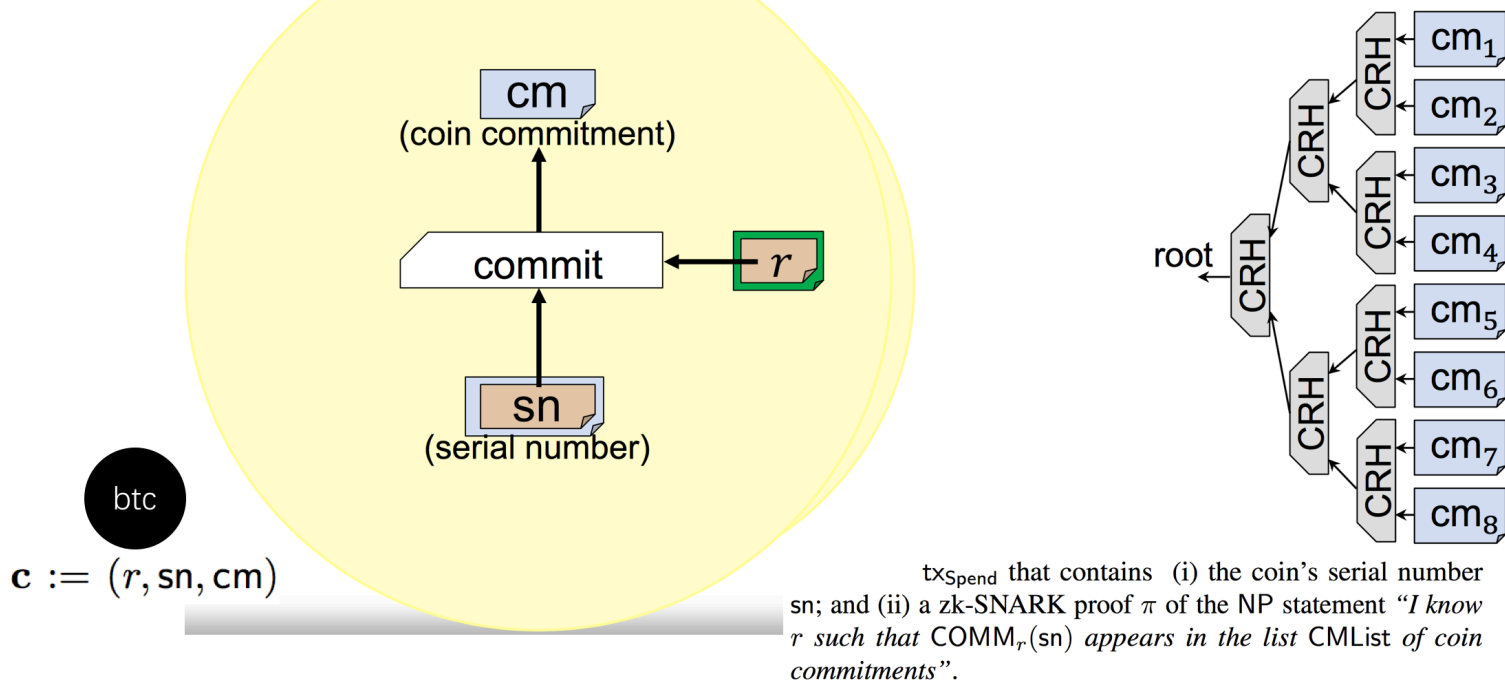
ZCash





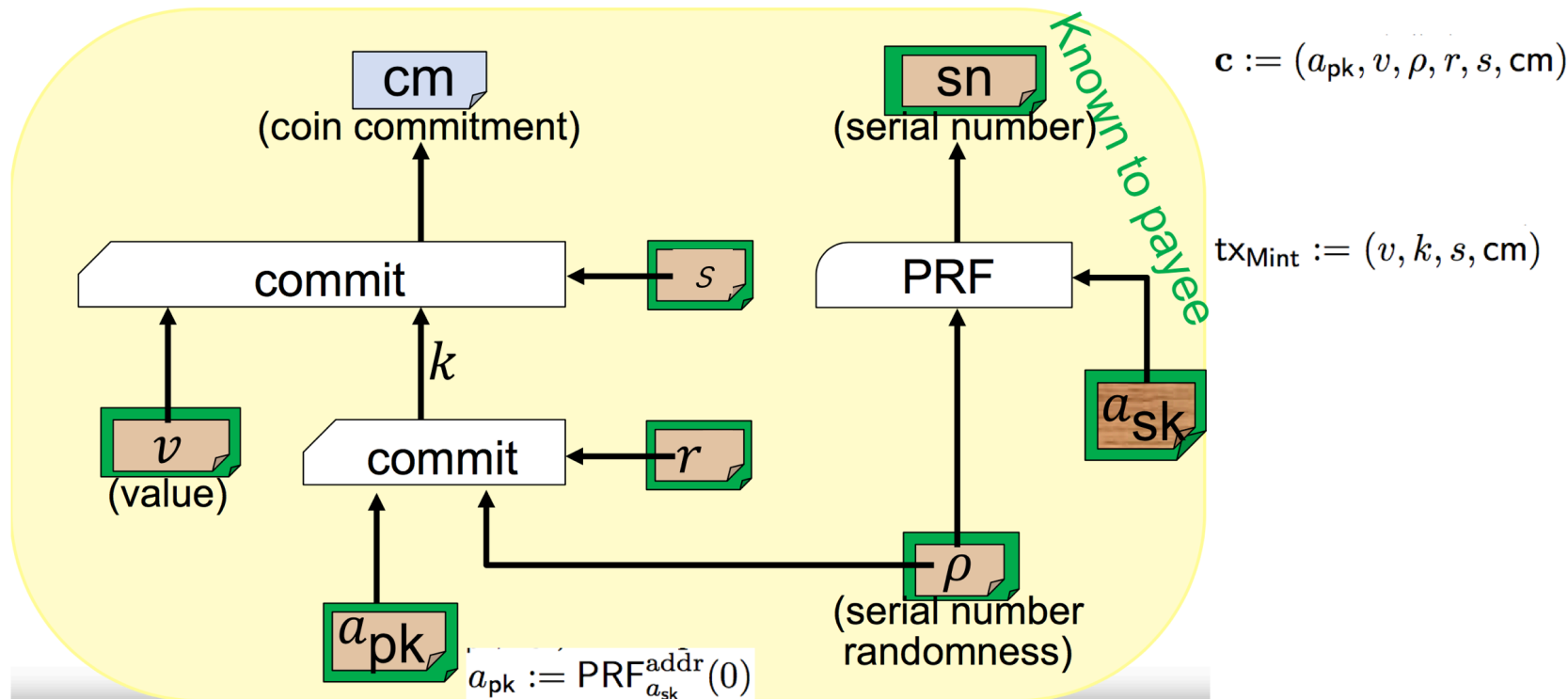
ZCash

$$cm := \text{COMM}_r(sn)$$





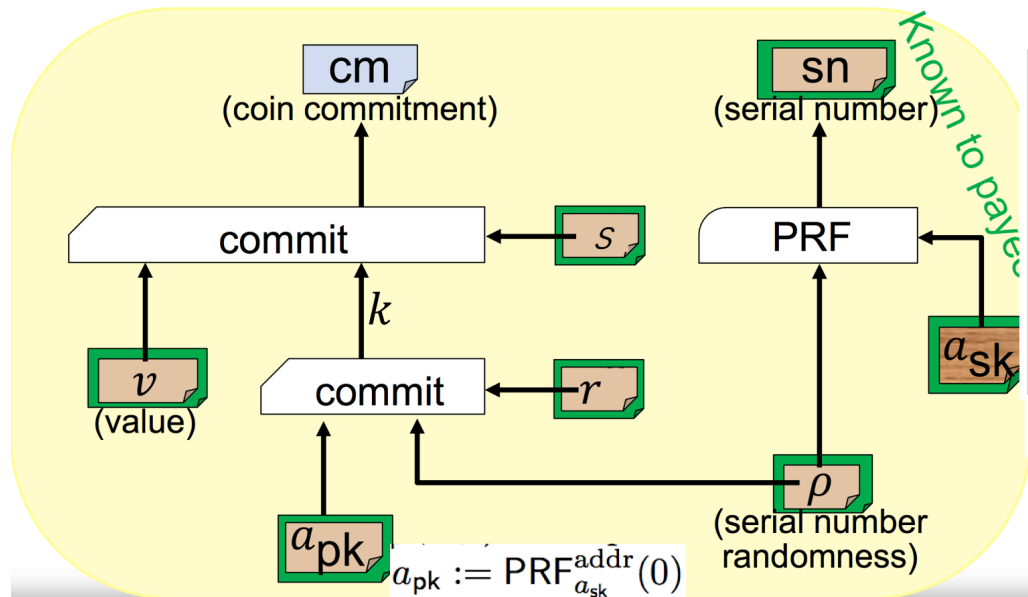
ZCash





ZCash

$$\mathbf{c} := (a_{\text{pk}}, v, \rho, r, s, \text{cm})$$



$$\mathbf{c}^{\text{old}} = (a_{\text{pk}}^{\text{old}}, v^{\text{old}}, \rho^{\text{old}}, r^{\text{old}}, s^{\text{old}}, \text{cm}^{\text{old}})$$

This yields the coins $\mathbf{c}_1^{\text{new}} := (a_{\text{pk},1}^{\text{new}}, v_1^{\text{new}}, \rho_1^{\text{new}}, r_1^{\text{new}}, s_1^{\text{new}}, \text{cm}_1^{\text{new}})$ and $\mathbf{c}_2^{\text{new}} := (a_{\text{pk},2}^{\text{new}}, v_2^{\text{new}}, \rho_2^{\text{new}}, r_2^{\text{new}}, s_2^{\text{new}}, \text{cm}_2^{\text{new}})$. Next, u produces a zk-SNARK proof π_{POUR} for the following NP statement, which we call POUR :

“Given the Merkle-tree root rt , serial number sn^{old} , and coin commitments $cm_1^{\text{new}}, cm_2^{\text{new}}$, I know coins $c^{\text{old}}, c_1^{\text{new}}, c_2^{\text{new}}$, and address secret key a_{sk}^{old} such that:

- *The coins are well-formed: for c^{old} it holds that $k^{\text{old}} = \text{COMM}_{r,\text{old}}(a_{\text{pk}}^{\text{old}} \parallel \rho^{\text{old}})$ and $\text{cm}^{\text{old}} = \text{COMM}_{s,\text{old}}(v^{\text{old}} \parallel k^{\text{old}})$; and similarly for c_1^{new} and c_2^{new} .*
- *The address secret key matches the public key: $a_{\text{pk}}^{\text{old}} = \text{PRF}_{a_{\text{sk}}^{\text{old}}}^{\text{addr}}(0)$.*
- *The serial number is computed correctly: $\text{sn}^{\text{old}} := \text{PRF}_{a_{\text{sk}}^{\text{old}}}^{\text{sn}}(\rho^{\text{old}})$.*
- *The coin commitment cm^{old} appears as a leaf of a Merkle-tree with root rt .*
- *The values add up: $v_1^{\text{new}} + v_2^{\text{new}} = v^{\text{old}}$.*

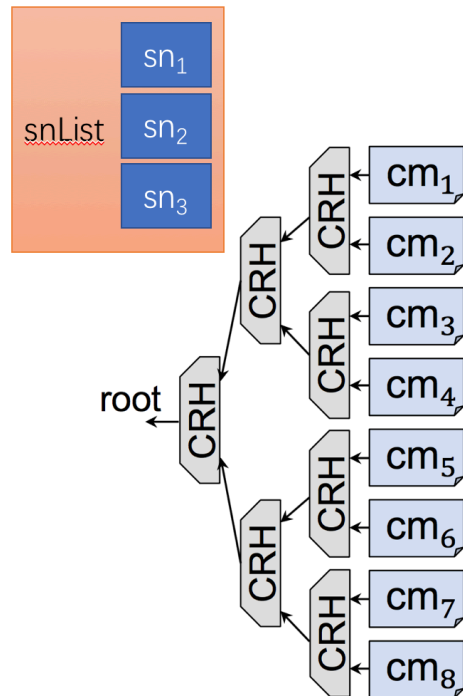
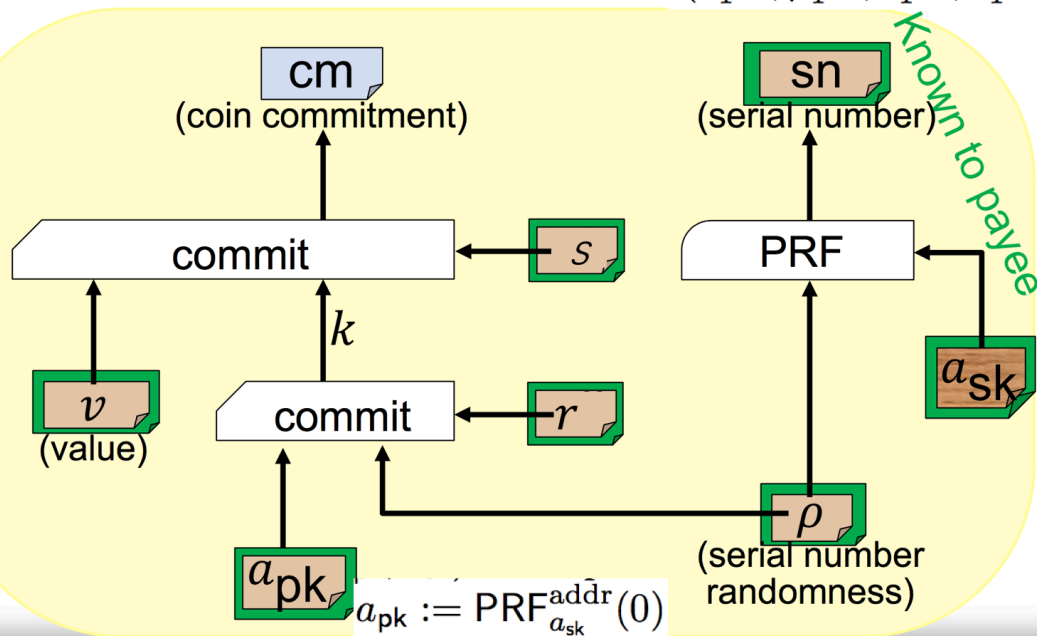


ZCash

加密传输给接收者

$(v_1^{\text{new}}, \rho_1^{\text{new}}, r_1^{\text{new}}, s_1^{\text{new}})$

$c := (a_{\text{pk}}, v, \rho, r, s, \text{cm})$





ZCash

