



# RISC-V S-mode Physical Memory Protection (SPMP)

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Version 0.9.10, 5/2025: This document is in development. Assume everything can change. See <http://riscv.org/spec-state> for details.

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## Preamble



*This document is in the [Development state](#)*

*Assume everything can change. This draft specification will change before being accepted as standard, so implementations made to this draft specification will likely not conform to the future standard.*

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## Chapter 1. Introduction

This document describes RISC-V S-mode Physical Memory Protection (SPMP) proposal to provide isolation when MMU is unavailable or disabled. RISC-V based processors recently stimulated great interest in the emerging internet of things (IoT) and automotive devices. However, page-based virtual memory (MMU) is usually undesired in order to meet resource and latency constraints. It is hard to isolate the S-mode OSes (e.g., RTOS) and user-mode applications for such devices. To support secure processing and isolate faults of U-mode software, the SPMP is desirable to enable S-mode OS to limit the physical addresses accessible by U-mode software on a hart.

## Chapter 2. S-mode Physical Memory Protection (SPMP)

An optional RISC-V S-mode Physical Memory Protection (SPMP) provides per-hart supervisor-mode control registers to allow physical memory access privileges (read, write, execute) to be specified for each physical memory region.

If PMP/ePMP is implemented, accesses succeed only if both PMP/ePMP and SPMP permission checks pass. The implementation can perform SPMP checks in parallel with PMA and PMP. The SPMP exception reports have higher priority than PMP or PMA exceptions (i.e., if the access violates both SPMP and PMP/PMA, the SPMP exception will be reported).

SPMP checks will be applied to all accesses whose effective privilege mode is S or U, including instruction fetches and data accesses in S and U mode, and data accesses in M-mode when the MPRV bit in mstatus is set and the MPP field in mstatus contains S or U.

SPMP registers can always be modified by M-mode and S-mode software.

SPMP can grant permissions to U-mode, which has none by default. SPMP can also revoke permissions from S-mode.

### 2.1. Requirements

1. S mode should be implemented
2. The `sstatus.SUM` (permit Supervisor User Memory access) bit must be **writeable**.

*The Privileged Architecture specification states the following*



*SUM has no effect when page-based virtual memory is not in effect, nor when executing in U-mode. SUM is read-only 0 if satp.MODE is read-only 0.*

*— Supervisor-Level ISA, Version 1.13 >> "Memory Privilege in `sstatus` Register"*

*In SPMP, this bit modifies the privilege with which S-mode loads and stores access to physical memory, hence the need to make it writeable.*

3. The `sstatus.MXR` (Make eXecutable Readable) bit must be **writeable**.

*The Privileged Architecture specification states that*



*MXR has no effect when page-based virtual memory is not in effect.*

*— Machine-Level ISA, Version 1.13 >> "Memory Privilege in `mstatus` Register"*

*In SPMP, the MXR bit modifies the privilege with which loads access physical memory. Its semantics are consistent with those of the Machine-Level ISA.*

*In SPMP, this bit is made writeable to support M-mode emulation handlers where instructions are read with MXR=1 and MPRV=1.*

## 2.2. S-mode Physical Memory Protection CSRs

Like PMP, SPMP entries are described by an 8-bit configuration register and one XLEN-bit address register. Some SPMP settings additionally use the address register associated with the preceding SPMP entry. Up to 64 SPMP entries are supported.

The SPMP configuration registers are packed into CSRs the same way as PMP. For RV32, 16 CSRs, `spmpcfg0`-`spmpcfg15`, hold the configurations `spmp0cfg`-`spmp63cfg` for the 64 SPMP entries. For RV64, even numbered CSRs (i.e., `spmpcfg0`, `spmpcfg2`, ..., `spmpcfg14`) hold the configurations for the 64 SPMP entries; odd numbered CSRs (e.g., `spmpcfg1`) are illegal. Figure 1 and Figure 2 demonstrate the first 16 entries of SPMP. The layout of the rest of the entries is identical.



*The terms, entry and rule, are similar to ePMP.*

*The implementation must decode all SPMP CSRs, and it can modify the number of **writable SPMP entries** while the remaining SPMP CSRs are read-only zero.*

*The lowest-numbered SPMP entries must be implemented first.*

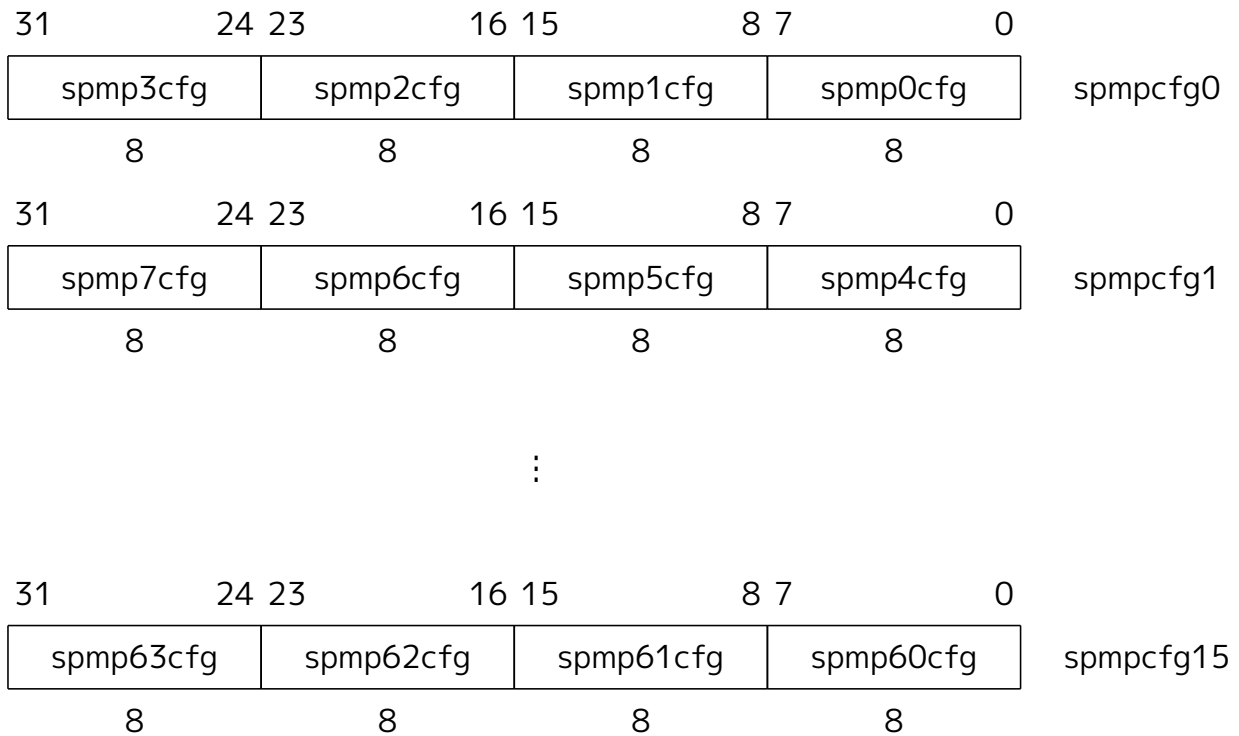


Figure 1. RV32 SPMP configuration CSR layout.

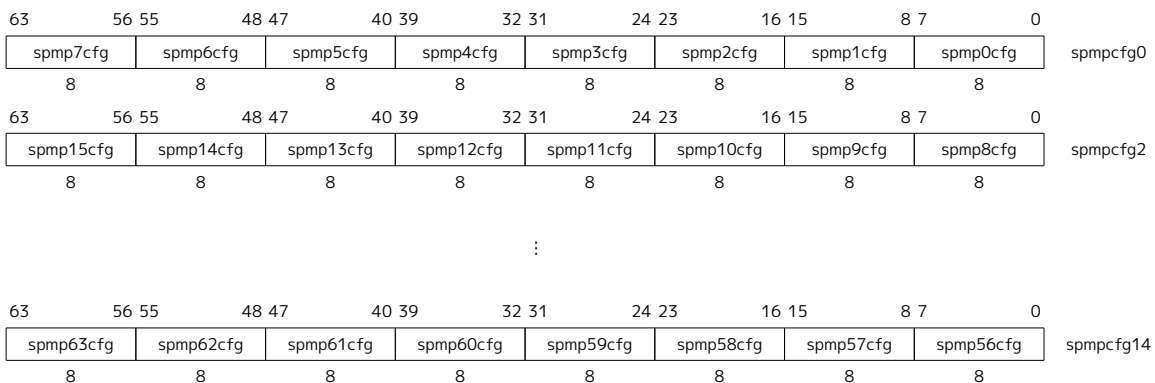


Figure 2. RV64 SPMP configuration CSR layout.



The SPMP address registers are CSRs named `spmpaddr0`–`spmpaddr63`. Each SPMP address register encodes bits 33–2 of 34-bit physical address for RV32, as shown in Figure 3. For RV64, each SPMP address encodes bits 55–2 of a 56-bit physical address, as shown in Figure 4. Fewer address bits may be implemented for specific reasons, e.g., systems with smaller physical address space. The number of address bits should be the same for all **writable SPMP entries**. Implemented address bits must extend to the LSB format, except as otherwise permitted by granularity rules. See the Privileged Architecture specification, Section 3.7: Physical Memory Protection, Address Matching.

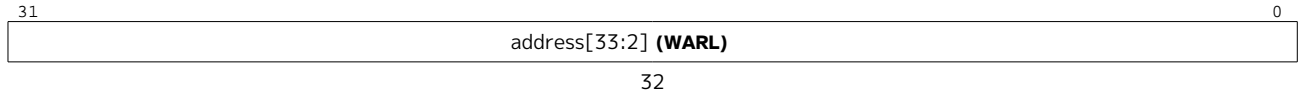


Figure 3. SPMP address register format, RV32.

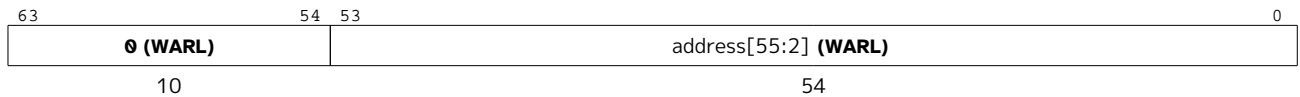


Figure 4. SPMP address register format, RV64.

The layout of SPMP configuration registers is the same as PMP configuration registers, as shown in Figure 5. The register is WARL. The rules and encodings for permission are explained in section 2.4, which resembles the encoding of ePMP (except SPMP does not use locked rules).

1. The S bit marks a rule as **S-mode-only** when set and **U-mode-only** when unset.
2. Bit 5 and 6 are reserved for future use.
3. The A field will be described in the following sections (2.3).
4. The R/W/X bits control read, write, and instruction execution permissions.

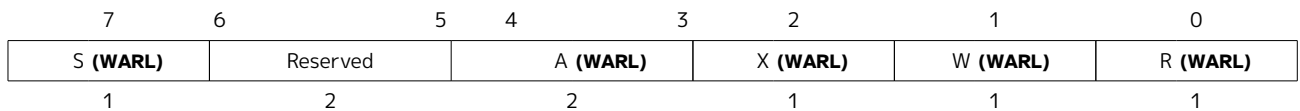


Figure 5. SPMP configuration register format.

## 2.3. Encoding of Permissions

SPMP has three kinds of rules: **S-mode-only**, **U-mode-only** and **Shared-Region** rules.

1. An **S-mode-only** rule is **enforced** on Supervisor mode and **denied** on User mode.
2. A **U-mode-only** rule is **enforced** on User modes and is either **denied** or **enforced** on Supervisor mode depending on the value of `sstatus.SUM` bit:
  - If `sstatus.SUM` is set, a U-mode-only rule is enforced on Supervisor mode, yet not be executable. This ensures the Supervisor Memory Execution Prevention (SMEP).
  - If `sstatus.SUM` is unset, a U-mode-only rule is denied on Supervisor mode. This ensures the Supervisor Memory Access Prevention (SMAP).
3. A **Shared-Region** rule is enforced on both Supervisor and User modes, with restrictions depending on the `spmpcfg.S` and `spmpcfg.X` bits:
  - If `spmpcfg.S` is not set, the region can be used for sharing data between S-mode and U-mode, yet not be executable. S-mode has RW permission to that region. U-mode has RW permission if `spmpcfg.X` is set, and it is restricted to read-only if `spmpcfg.X` is cleared.

- If `spmpcfg.S` is set, the region can be used for sharing code between S-mode and U-mode, yet not be writeable. S-mode has RX permission to that region. U-mode has RX permission if `spmpcfg.X` is set, and it is restricted to execute-only if `spmpcfg.X` is cleared.

4. The encoding `spmpcfg.SRWX=1000` is reserved for future standard use.

The encoding and results are shown in the [Figure 6](#):

	spmpcfg.S=0			spmpcfg.S=1		
spmpcfg	S mode Access	S mode Access	U mode Access	S mode Access	S mode Access	U mode Access
RWX	sstatus.SUM =0	sstatus.SUM =1	sstatus.SUM =x	sstatus.SUM =0	sstatus.SUM =1	sstatus.SUM =x
R - -	Deny	EnforceNoX	Enforce	Enforce	Enforce	Deny
R - X	Deny	EnforceNoX	Enforce	Enforce	Enforce	Deny
- - X	Deny	EnforceNoX	Enforce	Enforce	Enforce	Deny
- - -	Deny	EnforceNoX	Enforce	RSVD		
RW -	Deny	EnforceNoX	Enforce	Enforce	Enforce	Deny
RWX	Deny	EnforceNoX	Enforce	Enforce	Enforce	Deny
- WX	SHR RW			SHR RX		
- W -	SHR RW		SHR RO	SHR RX		SHR X

Figure 6. SPMP Encoding Table

**Deny:** Access fails.

**Enforce:** The R/W/X permissions are enforced on accesses.

**EnforceNoX:** The R/W permissions are enforced on accesses, while the X bit is forced to be zero.

**SHR:** It is shared between S/U modes with X, RX, RW, or ReadOnly privileges.

**RSVD:** It is reserved for future use.

**SUM bit:** The SPMP uses the `sstatus.SUM` (permit Supervisor User Memory access) bit to modify the privilege with which S-mode loads and stores access to physical memory. The semantics of `sstatus.SUM` in SPMP are consistent with those of the Machine-Level ISA (Please refer to the "Memory Privilege in mstatus Register" subsection in the riscv-privileged spec for detailed information).

## 2.4. Address Matching

The A field in an SPMP entry's configuration register encodes the address-matching mode of the associated SPMP address register. It is the same as PMP/ePMP.

Please refer to the "Address Matching" subsection of PMP in the riscv-privileged spec for detailed information.



*Software may determine the SPMP granularity by writing zero to `spmp0cfg`, then writing all ones to `spmpaddr0`, then reading back `spmpaddr0`. If  $G$  is the index of the least-significant bit set, the SPMP granularity is  $2^{G+2}$ .*

## 2.5. Supervisor Security Configuration (sseccfg) CSR

**Supervisor Security Configuration (sseccfg)** is a new Supervisor mode CSR used for configuring SPMP features. All sseccfg fields defined on this proposal are WARL, and the remaining bits are reserved for future standard use and should always read zero. This CSR has one field:

- Bit 0. sseccfg.SMAA (Supervisor Memory Access Allowlist Policy): When set, this bit changes the default SPMP policy for S-Mode when accessing memory regions that don't have a matching rule to denied instead of ignored. This bit must reset to 0.

For RV64 sseccfg is 64 bits wide, while for RV32 sseccfg is divided into sseccfg (lower 32 bits) and sseccfgh (upper 32 bits).

The sseccfg register must be cleared on reset. The presence of SPMP can be determined by probing the sseccfg CSR.

## 2.6. Matching Logic

- SPMP entries are statically prioritized, similar to PMP entries
- The lowest-numbered SPMP entry that matches any byte of access (indicated by an address and the accessed length) determines whether that access is allowed or denied
- The SPMP entry must match **all** bytes of access, or the access fails
- This matching is done irrespective of the S, R, W, and X bits

On some implementations, misaligned loads, stores, and instruction fetches may also be decomposed into multiple accesses, some of which may succeed before an exception occurs. In particular, a portion of a misaligned store that passes the SPMP check may become visible, even if another portion fails the SPMP check. The same behavior may manifest for stores wider than XLEN bits (e.g., the FSD instruction in RV32D), even when the store address is naturally aligned.

1. If the effective privilege mode of the access is M, the access is allowed;
2. If the effective privilege mode of the access is S and no SPMP entry matches, if sseccfg.SMAA is clear the access is allowed, otherwise if sseccfg.SMAA is set, the access is denied;
3. If the effective privilege mode of the access is U and no SPMP entry matches, but at least one SPMP entry is implemented, the access is denied;
4. Otherwise, each access is checked according to the permission bits in the matching SPMP entry. That access is allowed if it satisfies the permission checking with the SRWX encoding corresponding to the access type.

## 2.7. SPMP and Paging

The table below shows which mechanism to use. (Assume both paged virtual memory and SPMP are implemented.)

satp	Isolation mechanism
satp.mode == Bare	SPMP only
satp.mode != Bare	Paged Virtual Memory only

SPMP and paged virtual memory cannot be active simultaneously for two reasons:

1. An additional permission check layer would be introduced for each memory access.
2. Sufficient protection is provided by paged virtual memory.

That means SPMP is enabled when `satp.mode==Bare` and SPMP is implemented.



*Please refer to Table "Encoding of satp MODE field" in the riscv-privileged spec for detailed information on the satp.MODE field.*

## 2.8. Exceptions

When an access fails, SPMP generates an exception based on the access type (i.e., load accesses, store/AMO accesses, and instruction fetches). Each exception has a different code.

The SPMP reuses page fault exception codes for SPMP faults since page faults are typically delegated to S-mode. S-mode software (i.e., OS) can distinguish between SPMP and page faults by checking `satp.mode`, since SPMP and paged virtual memory cannot be active simultaneously (as described in [Section 2.7](#)). **SPMP proposes to rename page fault to SPMP/page fault for clarity.**

Note that a single instruction may generate multiple accesses, which may not be mutually atomic.

Table of renamed exception codes:

Interrupt	Exception Code	Description
0	12	Instruction SPMP/page fault
0	13	Load SPMP/page fault
0	15	Store/AMO SPMP/page fault



*Please refer to Table "Supervisor cause register (scause) values after trap" in the riscv-privileged spec for detailed information on exception codes.*

**Delegation:** Unlike PMP, which uses access faults for violations, SPMP uses SPMP/page faults for violations. The benefit of using SPMP/page faults is that the violations caused by SPMP can be delegated to S-mode, while the access violations caused by PMP can still be handled by machine mode.

## 2.9. Context Switching Optimization

Context switching with SPMP requires storing 64 address and 8 configuration registers (RV64), creating significant overhead. To optimize this:

- In RV32: two XLEN-bit read/write CSRs called `spmpswitch` and `spmpswitchh` are added, as depicted in [Figure 7](#).
- In RV64: one XLEN-bit read/write CSR called `spmpswitch` is added, as depicted in [Figure 8](#).

Each bit controls the activation of its corresponding SPMP entry. An entry is active only when both its `spmpswitch[i]` bit and `spmp[i]cfg.A` field are set, i.e., `spmpswitch[i] & spmp[i]cfg.A != 0`.



Figure 7. SPMP domain switch registers (spmpswitch and spmpswitchh), RV32.



Figure 8. SPMP domain switch register (spmpswitch), RV64.



When `spmpswitch` is implemented and `spmpcfg[i].A == TOR`, an entry matches any address  $y$  where:

- $\text{spmpaddr}[i-1] \leq y < \text{spmpaddr}[i]$
- This matching occurs regardless of `spmpcfg[i-1]` and `spmpswitch[i-1]` values

The `spmpswitch` registers must be cleared on reset.

## 2.10. Access Methods of SPMP CSRs

How SPMP CSRs are accessed depends on whether the `Sscsrind` extension is implemented or not.

**Direct CSR access:** SPMP CSRs can be accessed directly with corresponding CSR numbers if the `Sscsrind` extension is not implemented.

**Indirect CSR access:** The SPMP supports indirect CSR access if the `Sscsrind` extension is implemented. Each combination of `siselect` and `sireg` represents an access to the corresponding SPMP CSR.

siselect number	indirect CSR access of <code>sireg</code>
siselect#1	<code>sireg</code> → <code>spmpaddr[0]</code> , <code>sireg2</code> → <code>spmp[0]cfg</code>
siselect#2	<code>sireg</code> → <code>spmpaddr[1]</code> , <code>sireg2</code> → <code>spmp[1]cfg</code>
...	...
siselect#64	<code>sireg</code> → <code>spmpaddr[63]</code> , <code>sireg2</code> → <code>spmp[63]cfg</code>



The specific value of `siselect#1-64` will be allocated after review by the Arch Review Committee.

The rationale for SPMP only assign one rule per `siselect` value is due to performance consideration. If multiple SPMP rules are assigned to each `siselect`, a jump table or additional calculations would be needed to determine which `sireg` to assess.

Please refer to the `Sscsrind` extension specification for details on indirect CSR accesses: [github.com/riscv/riscv-indirect-csr-access](https://github.com/riscv/riscv-indirect-csr-access)

## Chapter 3. Extension of Resource Sharing

Given that PMP and SPMP have similar layout of address/config registers and the same address matching logic. Reusing registers and comparators between PMP and SPMP may be beneficial (in some cases) to save hardware resources. This section introduces the resource sharing extension that can support dynamic reallocation of hardware resource between PMP and SPMP.

Implementations can consider PMP/SPMP entries as a resource pool (called PMP\_Resource). Specifically, each PMP\_Resource consists of an address CSR, a configuration CSR, and associated micro-architecture state. A new M-mode CSR called `mpmpdeleg` is introduced to control the sharing of PMP\_Resource between PMP and SPMP.

In the following description, we will refer to the PMP/SPMP from the hardware perspective as PMP\_Resource, and the PMP/SPMP from the software perspective as entry.

The 16-bit CSR shown in [Figure 9](#) has one `pmpnum` field:

1. `pmpnum` is 7-bit, allowing a value of 0–64 to specify the number of PMP entries.
2. Any PMP\_Resource greater than or equal to the `pmpnum` is delegated to S-mode (SPMP). The lower numbered PMP\_Resource are left for M-mode (PMP). Delegating higher indexed PMP\_Resource to SPMP enables implementations with a single priority tree.
3. M-mode could set `pmpnum` ≥ 64 (the number of implemented PMP\_Resource), to reserve all resources for PMP.
4. M-mode could set `pmpnum` = 0 to delegate all PMP\_Resource to SPMP.
5. The reset value of `pmpnum` is 0b100\_0000.

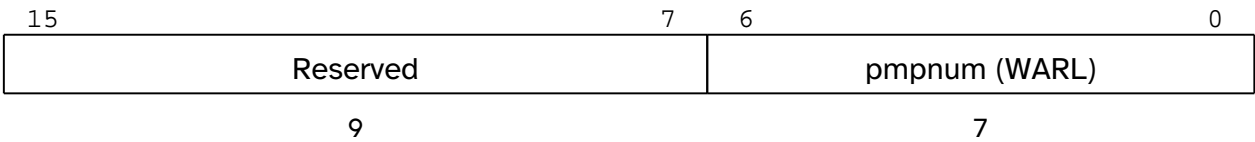


Figure 9. `mpmpdeleg` CSR format.

### Constraints:

1. With RV32, the values of `pmpnum`, can only be a multiple of 4; with RV64, it can only be a multiple of 8. This design avoids sharing the same configuration CSR between S-mode and M-mode.
2. The `pmpnum` is a WARL field. Illegal writes (e.g., values that are not multiples of 4 (RV32) or 8 (RV64)) should be ignored.
3. If the SPMP entry with lowest CSR number is configured with TOR address-matching mode, zero is used for the lower bound.



*In case a single pool of PMP\_Resource should be statically assigned to PMP/SPMP respectively, an implementation shall hardwire the value in `mpmpdeleg` accordingly for PMP and SPMP.*

### Addressing:

Both PMP and SPMP entries will be supported contiguously. The PMP entries begin with the lowest CSR number, while the SPMP entries begin with `mpmpdeleg.pmpnum`. For instance, given an implementation with a total of 64 PMP Resource entries, if `mpmpdeleg.pmpnum` is set to 16 during

runtime, `PMPResource[0]` to `PMPResource[15]` would map to `PMP[0]` to `PMP[15]`. The remaining entries, `PMPResource[16]` to `PMPResource[63]`, would be mapped as `SPMP[0]` to `SPMP[47]`.

From a software perspective, the SPMP entries start from `SPMP[0]`. The available number of SPMP entries can be discovered by writing to and reading from the SPMP CSRs. Illegal writes to SPMP CSRs will be ignored.

**Re-configuration:**

M-mode software can re-configure the entries for PMP and SPMP by modifying the `mpmpdeleg` CSR. A re-configuration involving locked PMP entry will leave `mpmpdeleg` unchanged.

## Chapter 4. Summary of Hardware Changes

Item	Changes
CSR for SPMP control	1 new CSR
CSRs for SPMP address	64 new CSRs
CSRs for SPMP configuration	16 new CSRs for RV32 and 8 for RV64
CSR for domain switch	2 new CSRs for RV32 and 1 for RV64
CSR for resource sharing	1 new CSR
Indirect access to SPMP CSRs	80 new siselect values
Renamed exception code	<b><i>Instruction page fault</i></b> renamed to <b><i>Instruction SPMP/page fault</i></b> <b><i>Load page fault</i></b> renamed to <b><i>Load SPMP/page fault</i></b> <b><i>Store/AMO page fault</i></b> renamed to <b><i>Store/AMO SPMP/page fault</i></b>



## Chapter 5. Recommended Programming Guidelines

When configuring SPMP to isolate user-mode tasks from each other and from the operating system (OS) executing in supervisor mode, two primary usage models arise, depending on whether the available SPMP entries can simultaneously accommodate all required memory ranges for both the user tasks and the OS:

- **Static Configuration:** all SPMP entries are programmed during system initialization. This model assumes that the number of available entries is sufficient to cover the complete set of memory regions assigned to user tasks and the OS without further modification.
- **Dynamic Configuration:** SPMP entries are reprogrammed on each context switch. This model is employed when the number of available SPMP entries is insufficient to simultaneously represent all relevant memory regions, requiring dynamic updates to enforce memory isolation between tasks.

### 5.1. Static Configuration

In the static configuration model, the number of available SPMP entries is sufficient to accommodate all required memory ranges for user-mode tasks and the OS. In this case, SPMP entries are programmed once during system initialization and remain unchanged at runtime. Only the `spmpswitch` register(s) need to be updated during context switches between user-mode tasks.

The OS begins by allocating SPMP entries and populating the `spmpaddr[i]` and `spmpcfg[i]` CSRs with the appropriate address ranges and permissions for each user-mode task. If the OS itself is to be protected using SPMP, additional entries must be allocated and configured with the OS memory ranges, with the `S` bit (Supervisor bit) set in `spmpcfg[i]`.

After initializing the entries, the OS must also set the corresponding bits in the `spmpswitch` register to activate these entries, and enable the allowlist policy for supervisor mode by setting the `SMAA` bit in the `ssecfg` CSR.

Prior to launching the first user task, the OS sets the bits in `spmpswitch` corresponding to the SPMP entries assigned to that task. During a context switch, the OS clears the current task's entry bits and sets those of the newly scheduled task using the `csrc` and `csrs` instructions, respectively. On RV32 systems, both `spmpswitch` and `spmpswitchh` must be managed as a continuous pair.

### 5.2. Dynamic Reconfiguration

In this configuration model, the available SPMP entries are insufficient to simultaneously represent the memory ranges required for all user-mode tasks and the supervisor. As a result, the OS must dynamically reconfigure SPMP entries for user tasks on every context switch. Notably, for any given hart, the number of SPMP entries must still be sufficient to hold both the supervisor entries and the entries for the currently executing user-mode task.

Where sufficient SPMP entries exist to cover all tasks and the OS, an implementation may simply update `spmpswitch` on context switches. Otherwise, the following sequence is recommended for dynamic reconfiguration:

1. **Disable Entries for the Outgoing Task.** Use the `csrc` instruction to clear the `spmpswitch` bits corresponding to the SPMP entries of the outgoing task. A bitmask representing active entries (typically stored in the task's control block) should be passed as an argument.

2. **Update SPMP Address Registers.** Write the `spmpaddr[i]` CSRs with the address ranges corresponding to the memory regions of the incoming task.
3. **Update SPMP Configuration Registers.** For each corresponding `spmpcfg[i]` field:
  - Clear the existing configuration bits using the `csrc` instruction with an appropriate mask;
  - Set the desired configuration using the `csrs` instruction.
4. **Enable Entries for the Incoming Task.** Use the `csrs` instruction to write to `spmpswitch`, passing a bitmask that enables the SPMP entries allocated to the incoming task.



*It is recommended that SPMP entries configured to protect the supervisor (i.e., entries with `spmpXcfg[i].S = 1`) remain resident and are not reprogrammed during the context-switch process. Maintaining these entries as persistent minimizes reconfiguration overhead and ensures consistent enforcement of memory protection for the supervisor across task switches.*

### 5.3. Entry Configuration Recommendations

When programming SPMP entries, a trade-off exists between using exclusively Naturally Aligned Power-Of-Two (NAPOT) or Top-Of-Range (TOR) address-matching modes.

While NAPOT allows compact encoding for power-of-two-aligned regions using a single entry, it may lead to internal fragmentation if the region size exceeds the actual requirement, resulting in memory waste. On the other hand, TOR mode (particularly in its generic form) requires two SPMP entries per protected region (base and top), which may exhaust available entries more quickly. However, for regions that are naturally power-of-two aligned, TOR may still be encoded with fewer entries.

This trade-off becomes especially relevant in MCU environments, where local memories are often sparsely mapped to fixed address ranges associated with specific core functions. In such cases, exclusive use of NAPOT or naïvely pairing TOR entries may be inefficient or lead to undesirable gaps in protection. Additionally, while it is possible to define multiple contiguous regions with different access permissions using overlapping or consecutive TOR entries, this technique can introduce subtle dependencies. Sharing a top address between two address spaces (e.g., supervisor and user) can lead to unintended interactions: reducing the supervisor region may inadvertently expand the adjacent user region. Similarly, replacing a shared top address during context switches could expose previously protected memory.

Given these risks, the following configuration model is recommended:

- Use TOR mode exclusively, treating each pair of `spmpaddr[i]` registers (even/odd indexed) as a base/top pair defining a single memory region;
- Conceptually, view SPMP entries as organized into pairs: (0/1), (2/3), ..., (62/63). Only the odd-indexed entries are enabled via the corresponding bits in `spmpswitch`, as each odd entry finalizes the definition of a region;
- Allocate and populate SPMP entries in descending index order (i.e., from lower priority to higher priority), starting from the highest index downward. This allocation strategy allows the OS to define temporary subregions by configuring unused lower-index entries without needing to reconfigure existing higher-index (priority) ones.

This disciplined use of TOR-mode SPMP entries ensures clearer isolation boundaries, reduces the likelihood of configuration errors, and improves runtime flexibility for memory protection schemes.

## 5.4. Re-configuration Non-preemption and Synchronization

To preserve the integrity of SPMP state, the reconfiguration process during a context switch must be executed as a non-preemptible critical section. This requirement stems from the need to update multiple control and configuration CSRs, and any interruption or concurrent modification during this process can result in transient inconsistencies or unintended access permissions.

In the **dynamic reconfiguration model**, the critical section must be enforced across updates to the following CSRs: `smpaddr[i]`, `smpcfg[i]`, and `smpswitch`. In the **static configuration model**, this concern is relevant only for RV32 systems with more than 32 SPMP entries, where both `smpswitch` and `smpswitchh` must be updated in coordination.

To prevent asynchronous supervisor-level interrupts during reconfiguration, the supervisor must clear the SIE (Supervisor Interrupt Enable) bit in the `sstatus` CSR. Furthermore, synchronisation mechanisms (e.g., mutexes or spinlocks) must be employed to serialise access to SPMP CSRs in multi-threaded or multi-core systems, ensuring that concurrent modifications do not result in conflicting or corrupted configurations.

By enforcing non-preemption and proper synchronisation, software ensures that SPMP protections remain deterministic, secure, and verifiable across context switches.

## Chapter 6. Interaction with other proposals

This section discusses how SPMP interacts with other proposals.

**J-extension pointer masking proposal:** When both PM and SPMP are used, SPMP checking should be performed using the actual addresses generated by PM (pointer masking).

**Hypervisor extension:** Virtualization support for the SPMP will be an independent proposal.

**Smstateen extension:** SPMP adds readable and writable states in S-mode, which can be abused as a covert channel if the OS/hypervisor is not aware of SPMP (thus the states won't be context-switched). It is desired that SPMP occupies a bit in mstateen register of Smstateen extension, which can control supervisor access to SPMP states.