Domain-Specific Languages

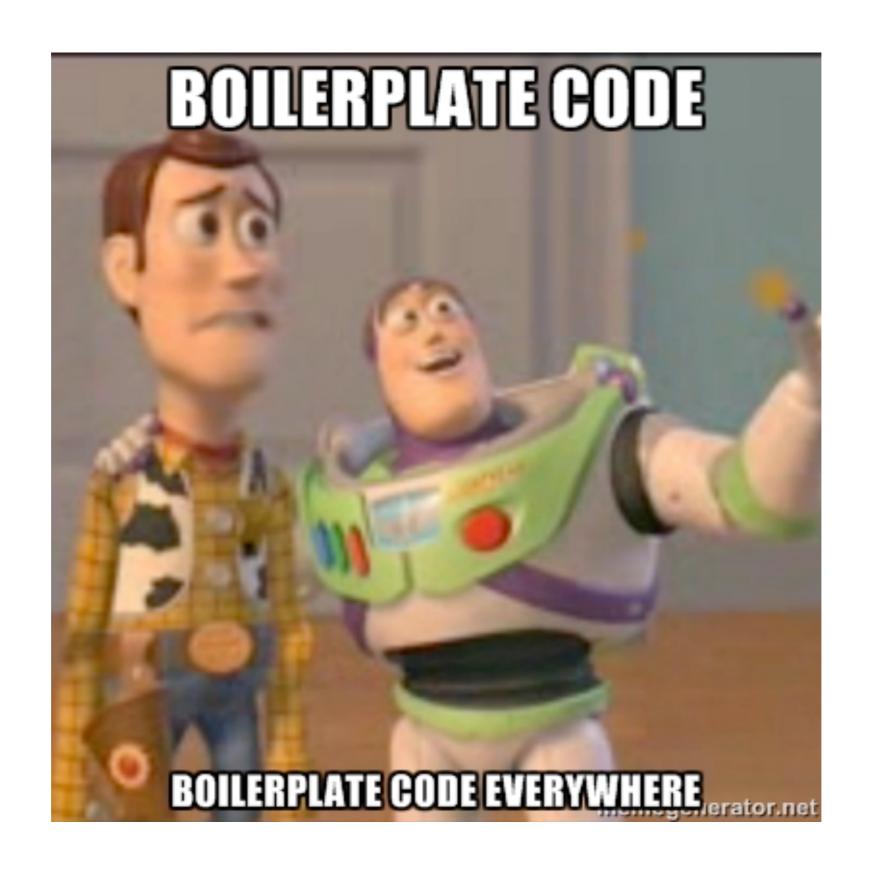
Tijs van der Storm (storm@cwi.nl / @tvdstorm)





Programming?





A programming language is low level when its programs require attention to the irrelevant

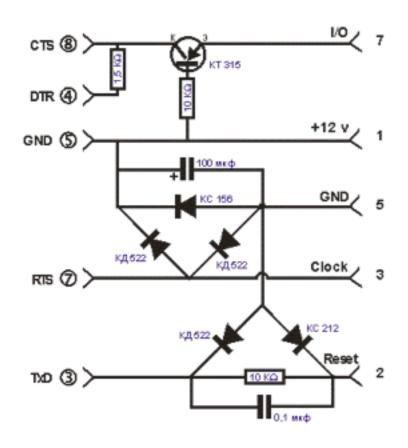
Some facts

- Fact 41. Maintenance typically consumes 40 to 80 percent of software costs.

 It is probably the most important life cycle phase of software.
- Foct 44. Understanding the existing product is the most difficult task of maintenance.
- Fact 21. For every 25 percent increase in problem complexity, there is a 100 percent increase in solution complexity.

Domain Specific Languages!

Domain specific languages



$$c_{t} = S_{t}N(h) - Xe^{-rt}N(h - \sigma\sqrt{\tau})$$

$$h = \left\{\ln(\frac{S}{X}) + r\tau + \frac{\sigma^{2}\tau}{2}\right\}$$

$$CH_{3}$$

$$CH_{3}$$

$$CH_{3}$$

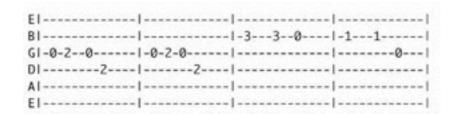
$$CH_{3}$$

$$CH_{3}$$

$$CH_{4}$$

$$CH_{2}$$

Domain specific languages

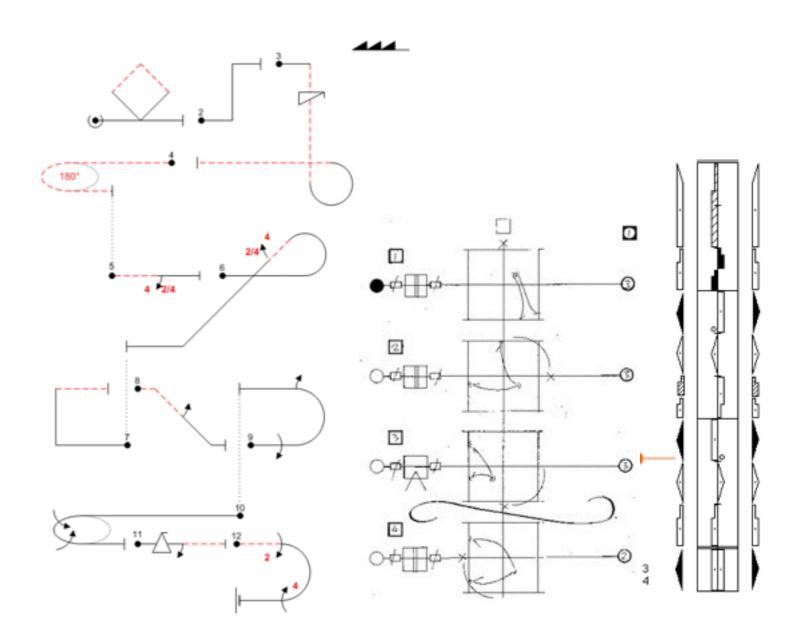


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Observations

- Special purpose
- Restricted
- Concise
- Expert usage
- Formalized
- Textual or graphic or combination

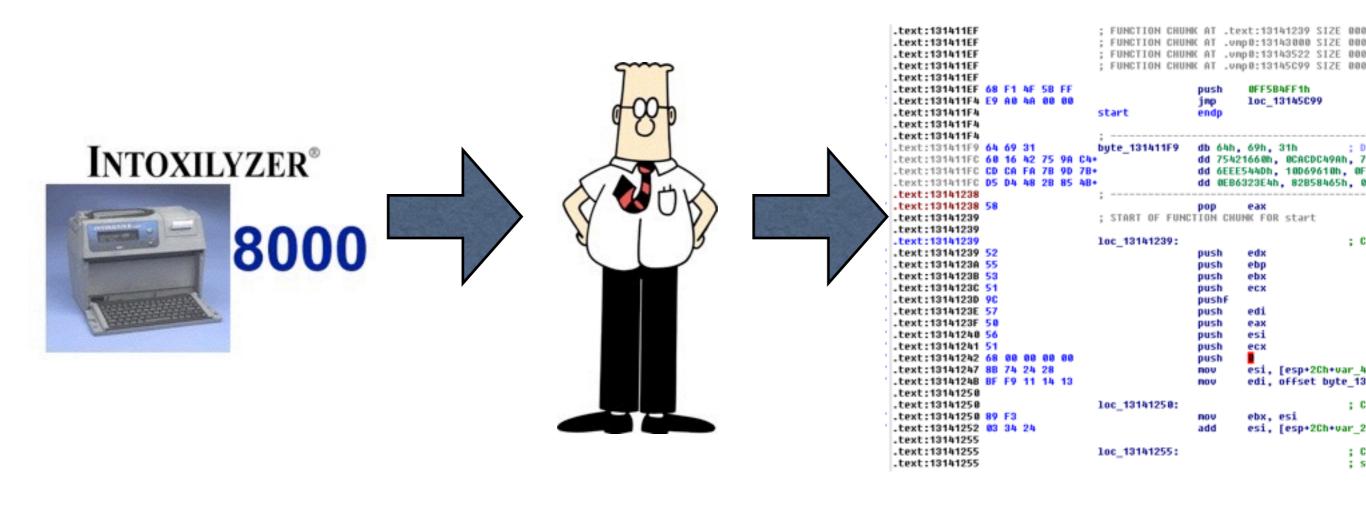
General purpose languages (GPLs)



DSLs



Programming



Domain

Programmer

Code

Programming



Domain

Programmer

Code

Programming is "lossy"

- encoding
- obfuscating
- encrypting
- dispersing
- tangling
- distorting

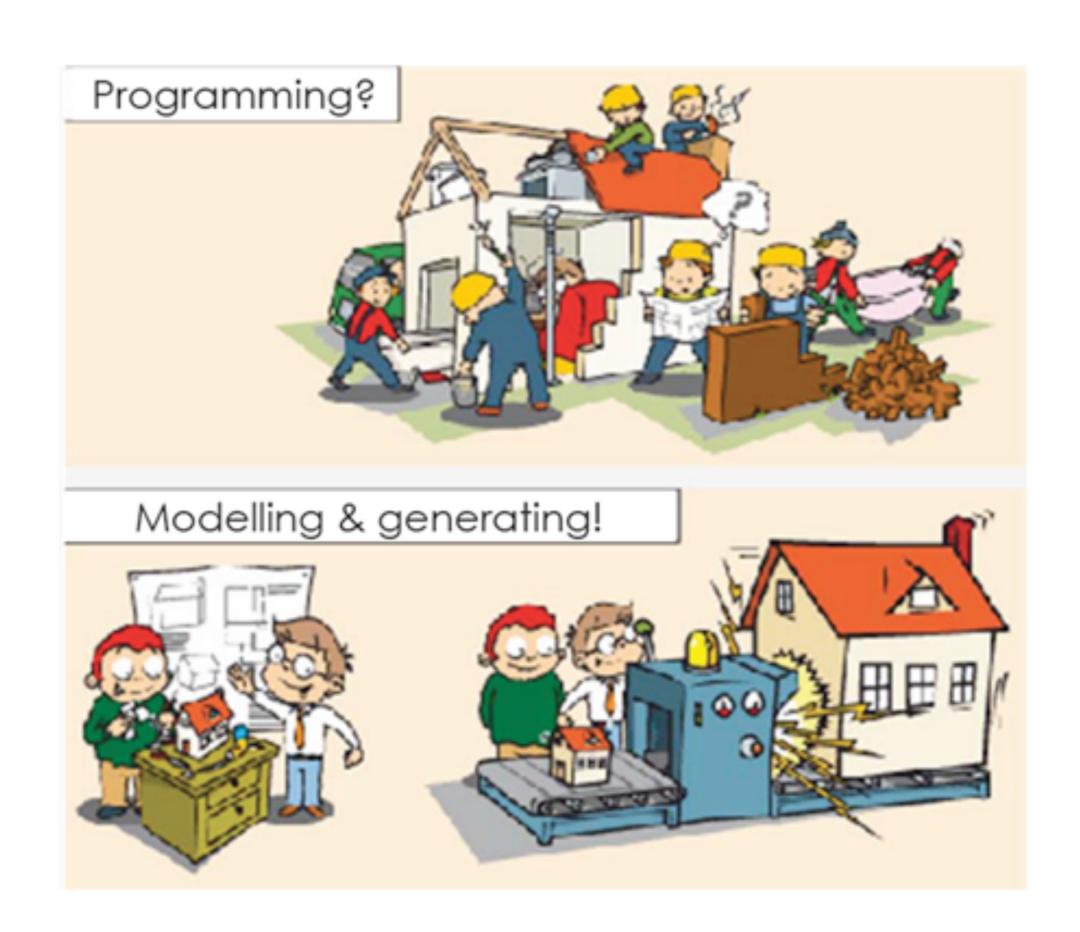


Cognitive distance



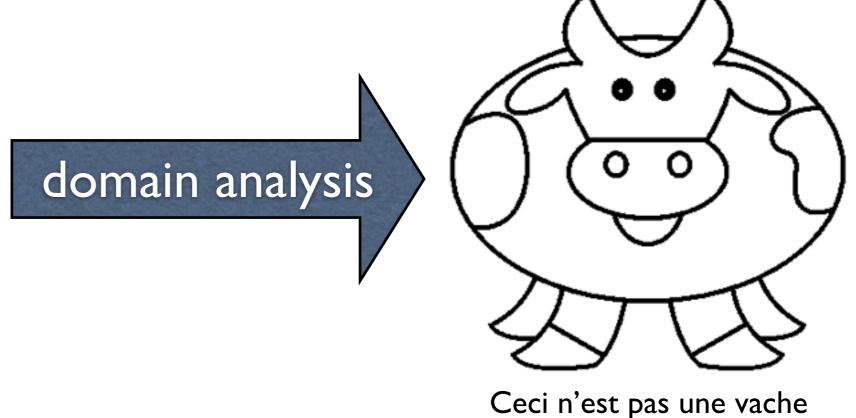
The problem

- a lot of code,
- low level code,
- characterized by lack of abstraction
- encoding domain knowledge
- and encoding design knowledge



Modeling the domain

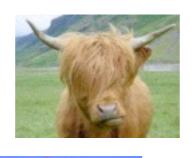




System families











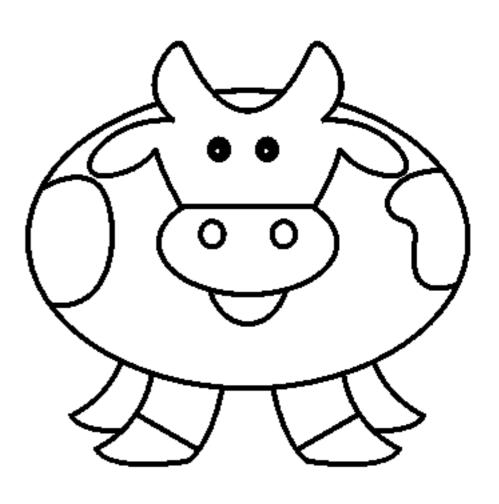




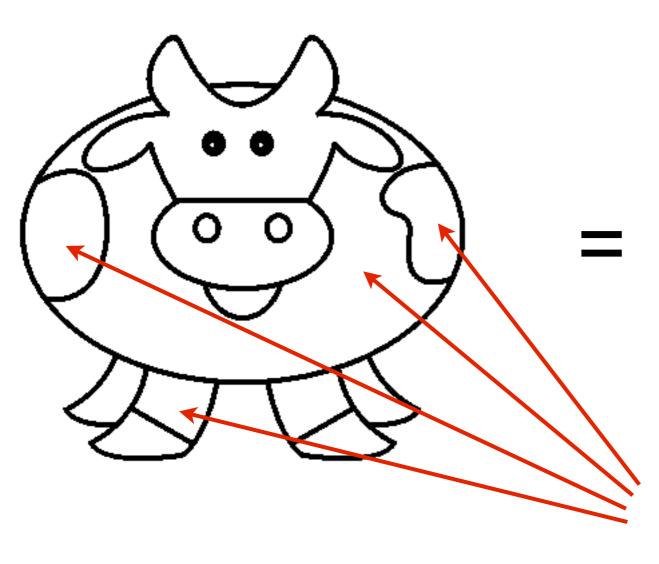








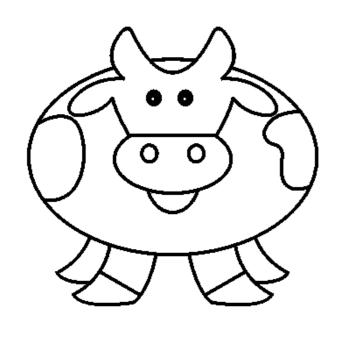
Domain Specific Language



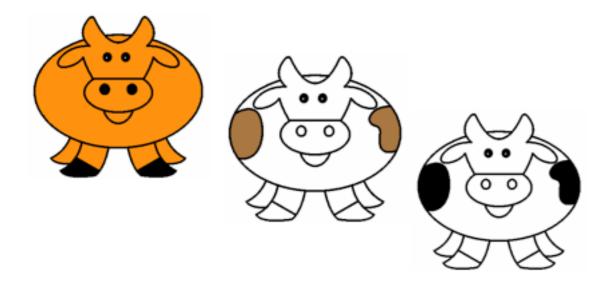
formalized notation capturing "Cows"

variation points

Domain Specific Languages

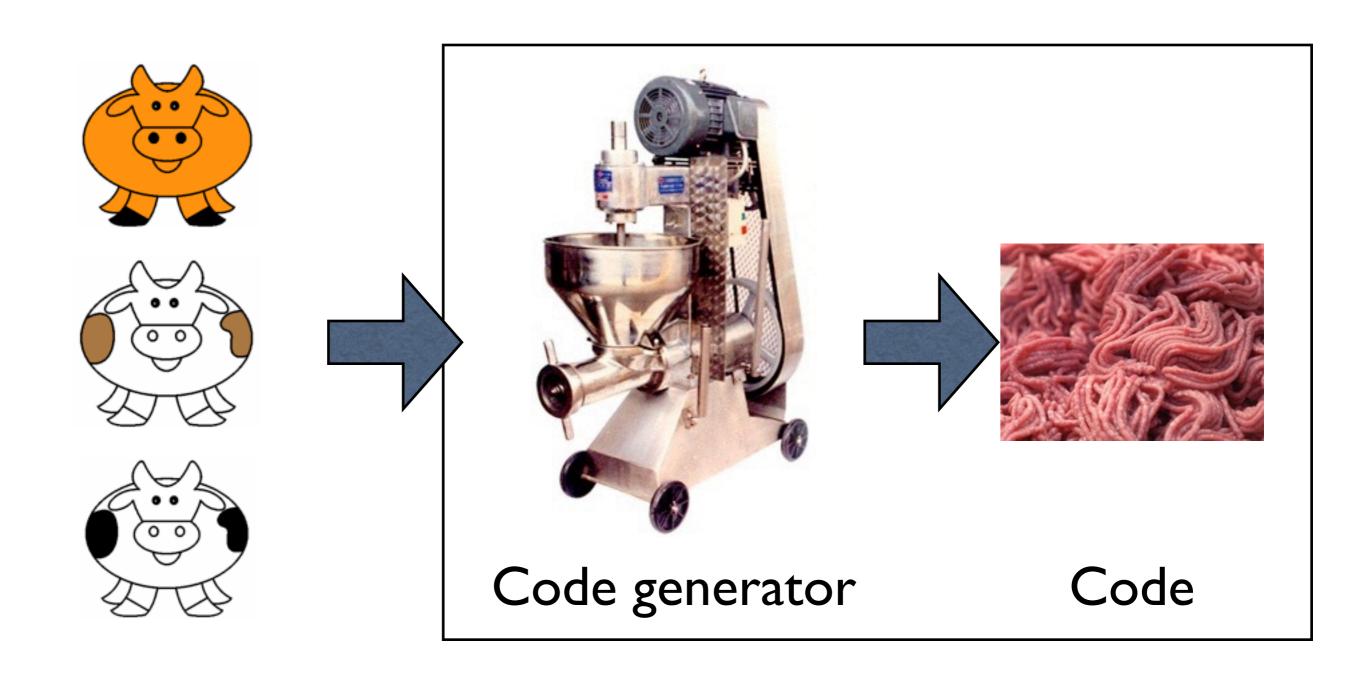


grammar,template,metamodel

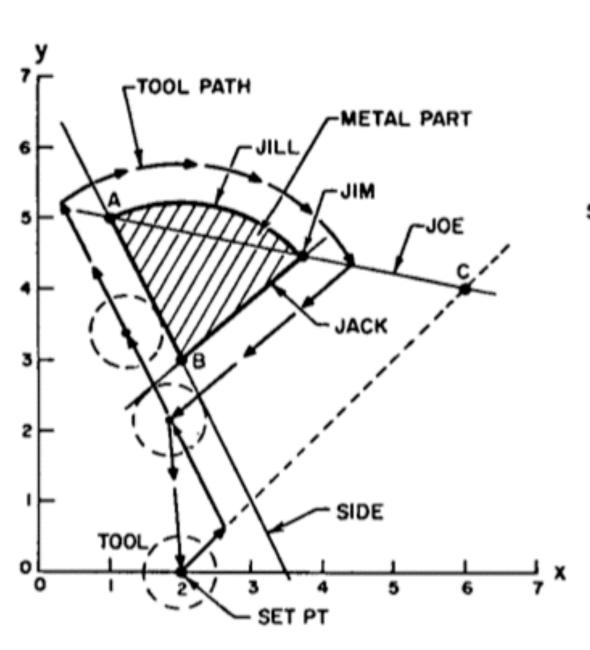


sentence,
= instance,
model

Code generation



APT: numerical control



A . POINT / I, 5

B = POINT/2, 3

C * POINT / 6, 4

TL DIA / +1.0, INCH

FEDRAT/ 30, IPM

SET PT = FROM, POINT/2, O

IN DIR, POINT/ C

SIDE . GO TO, LINE / THRU, A, AND, B

WITH, TL LFT, GO LFT, ALONG / SIDE

JILL = GO RGT, ALONG, CIRCLE / WITH, CTR AT, B, THRU, A

JOE = LINE / THRU, A, AND, C

JIM . POINT / X LARGE, INT OF, JOE, WITH, JILL

JACK = LINE / THRU, JIM, AND, B

GO RGT, ALONG/ JACK, UNTIL, TOOL, PAST, SIDE

GO TO/ SET PT

STOP, END, FINI

from the '50s (!)

La TeX: document preparation

```
\subsection{Application à un exemple: la coévolution proies-prédateurs}
            \subsubsection{Étape 1: Modèle écologique et stationnarité}
Nous nous intéresserons dans ce cas au modèle simple de Lotka-Volterra,
énoncé par le système~\ref{eq:lotka volterra}.
Dans ce modèle de base, il faut introduire une dépendance au trait sujet
à évolution qui nous intéresse. Ici, nous considérons la taille
 corporelle $x$ comme trait d'intérêt et supposons que la compétition
intraspécifique, $\alpha$, et la prédation, $\beta$, en dépendent ainsi:
\begin{eqnarray}
            \alpha(x 1)&=&\alpha 0+\alpha 2(x 1-x {1 {0}})^{2}\\
            \beta(x 1,x 2)&=&\beta 0
\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ens
2\beta 3\left(\frac{x 1}{\beta 1}\right)\left(\frac{x 2}{\beta 2}\right)
 \left(\frac{x 2}{\beta 2}\right)^{2}\right)
\end{eqnarray}
\begin{figure}[p]
            \begin{center}
                        \includegraphics[width=0.45\textwidth]{figures/func alp}
                        \includegraphics[width=0.45\textwidth]{figures/func bet}
                        \caption{Les fonctions choisies pour $\alpha$ et $\beta$}
            \end{center}
 \end{figure}
```

VHDL: hardware description

```
library IEEE;
use IEEE STD LOGIC 1164 ALL:
use IEEE STD LOGIC ARITH ALL:
use IEEE.STD LOGIC UNSIGNED.ALL;
entity ClkDiv is
    RegSel: in STD LOGIC VECTOR(1 downto 0); --<-- Seq CPLD
            RegStrb : in STD LOGIC;
                                                            --<-- Seq CPLD
            MClk : in STD LOGIC;
                                                            --<-- OSC
            SeqReset : in STD LOGIC;
                                                      --<-- Power Monitor
            ADC Clk : out STD LOGIC);
                                                           -->-- ADC
end ClkDiv;
architecture Behavioral of ClkDiv is
  signal ADC_div : STD_LOGIC_VECTOR(5 downto 0) := "001111";
signal ADCClk : STD_LOGIC := '0';
signal ClkSel : STD_LOGIC_VECTOR(2 downto 0) := "100";
begin
```

Risla: financial products

product LOAN

declaration contract data

PAMOUNT: amount

STARTDATE : date

MATURDATE : date

INTRATE: int-rate

RDMLIST := [] : cashflow-list %% List of redemptions.

Time to market went down from 3 months to 3 weeks.



%% Principal Amount

%% Starting date

%% Maturity data

%% Interest rate

information

PAF : cashflow-list

IAF : cashflow-list

%% Principal Amount Flow

%% Interest Amount Flow

registration

%% Register one redemption.

RDM (AMOUNT : amount, DATE : date)

QL

```
form Box1HouseOwning {
   "Did you sell a house in 2010?" hasSoldHouse: boolean
   "Did you by a house in 2010?" hasBoughtHouse: boolean
   "Did you enter a loan for maintenance?" hasMaintLoan: boolean
   if (hasSoldHouse) {
        "Private debts for the sold house:" privateDebt: money
        "Price the house was sold for:" sellingPrice: money
        "Value residue:" valueResidue = sellingPrice - privateDebt
   }
}
```

Other examples

- Make: software building
- Dot: graph visualization
- SQL: relational querying
- SWUL: Swing GUIs
- HTML: hypertext
- CLOPS: commandline options

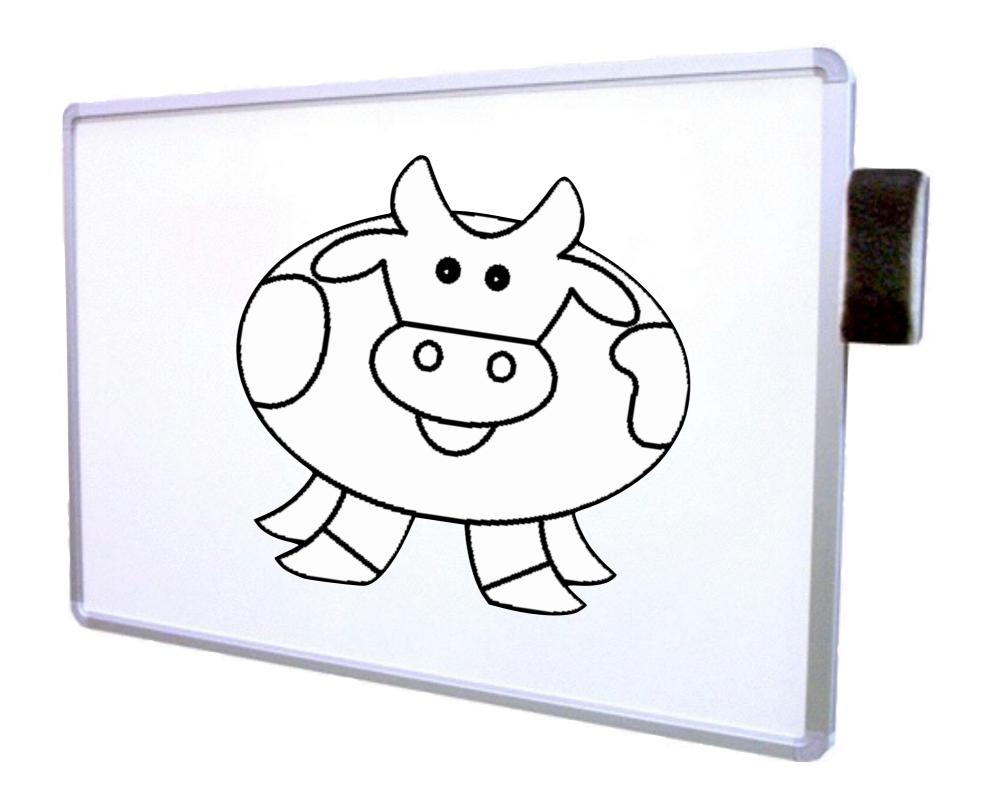
- GNUPlot: plotting
- R: statistics
- CML: kernel config
- Lex: lexical scanning
- Excel: spreadheets
- Rascal: meta-programming

• ...

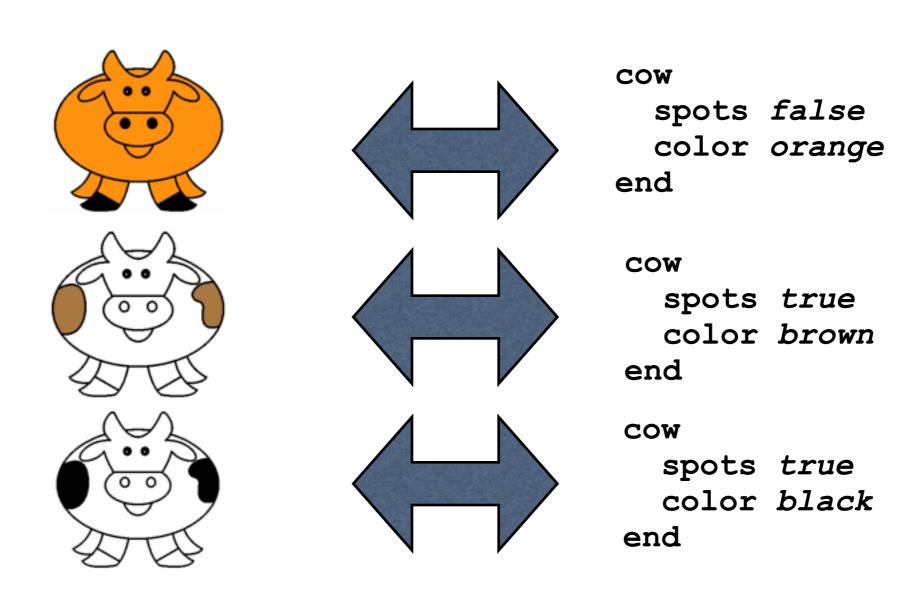
Domain-specific languages

- Better languages, for specific domains
- Capture families of systems
- Higher level of abstraction
- Focus on "what" vs "how"
- Reuse designs, not just code
- Language workbenches (e.g., Rascal)

DSL Implementation



DSL Code



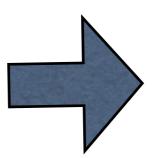
Internal DSLs



```
new Cow()
    .spots(false)
    .color("orange")
.end();
```



spots false
 color orange
end



(cow
 spots #t
 color 'orange)





cow do
 spots false
 color :orange
end



Advantages

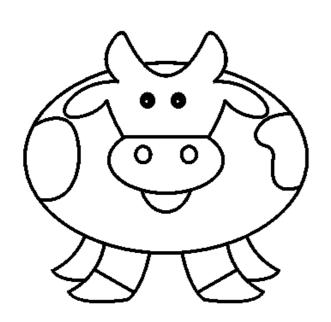
- No need to write/maintain parser
- Host language available if needed
- Use of existing tools (IDE) etc.

Drawbacks

- Restricted to host language
- Less static checking
- Fewer opportunities for optimization

External DSLs

Repetitio



alternative

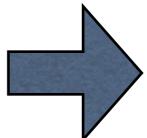
Parser generation

Grammar

yacc bison lemon

Parser



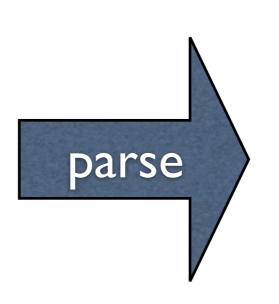


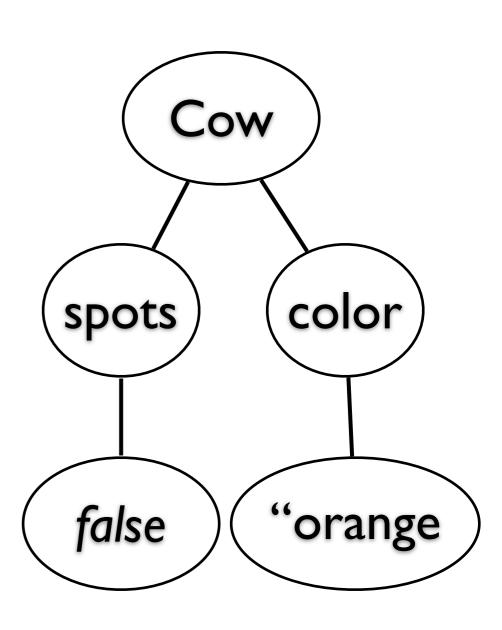
parse.exe

javacup antlr Rats!

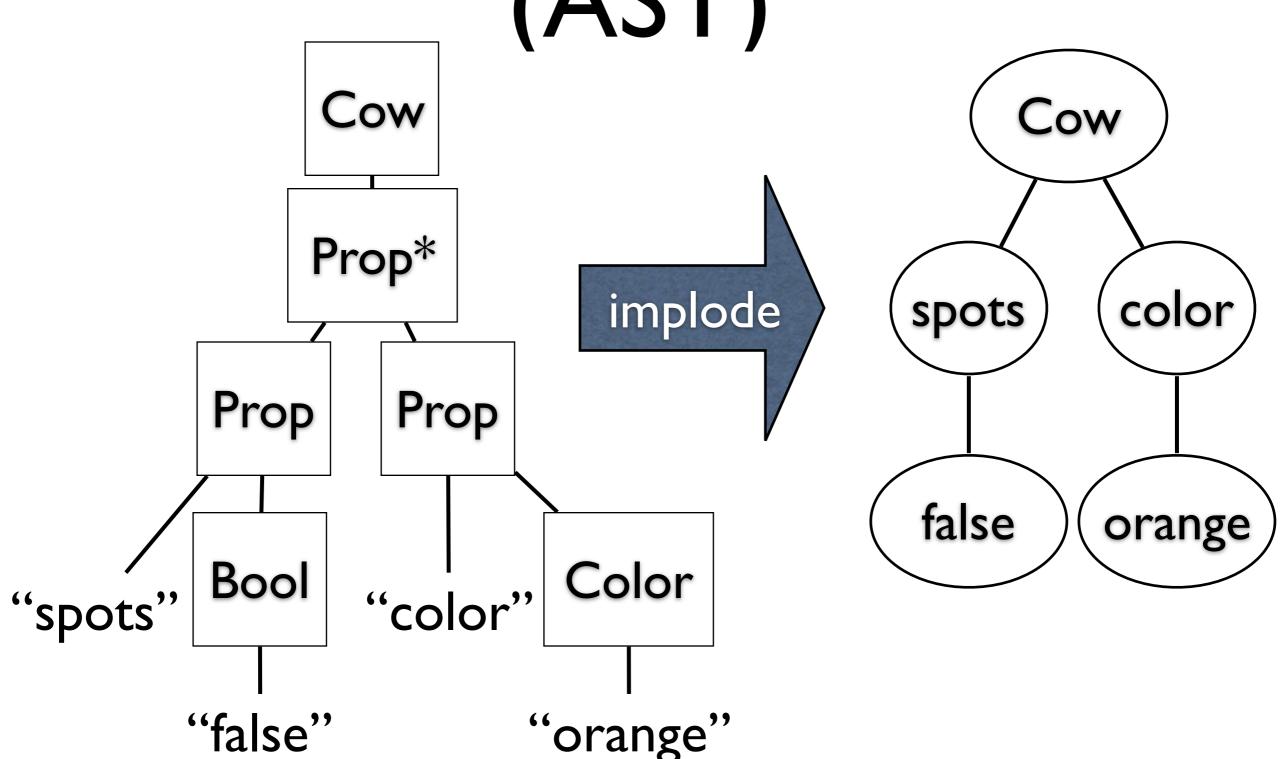
Parsing

cow
spots false
color orange
end



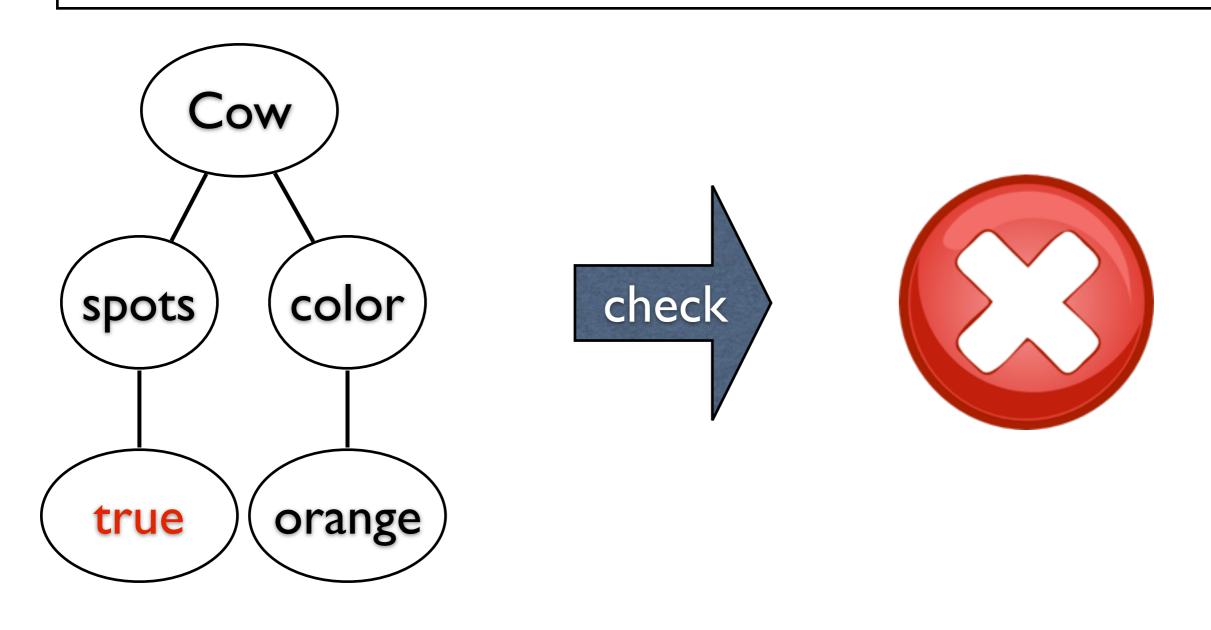


Abstract syntax tree (AST)

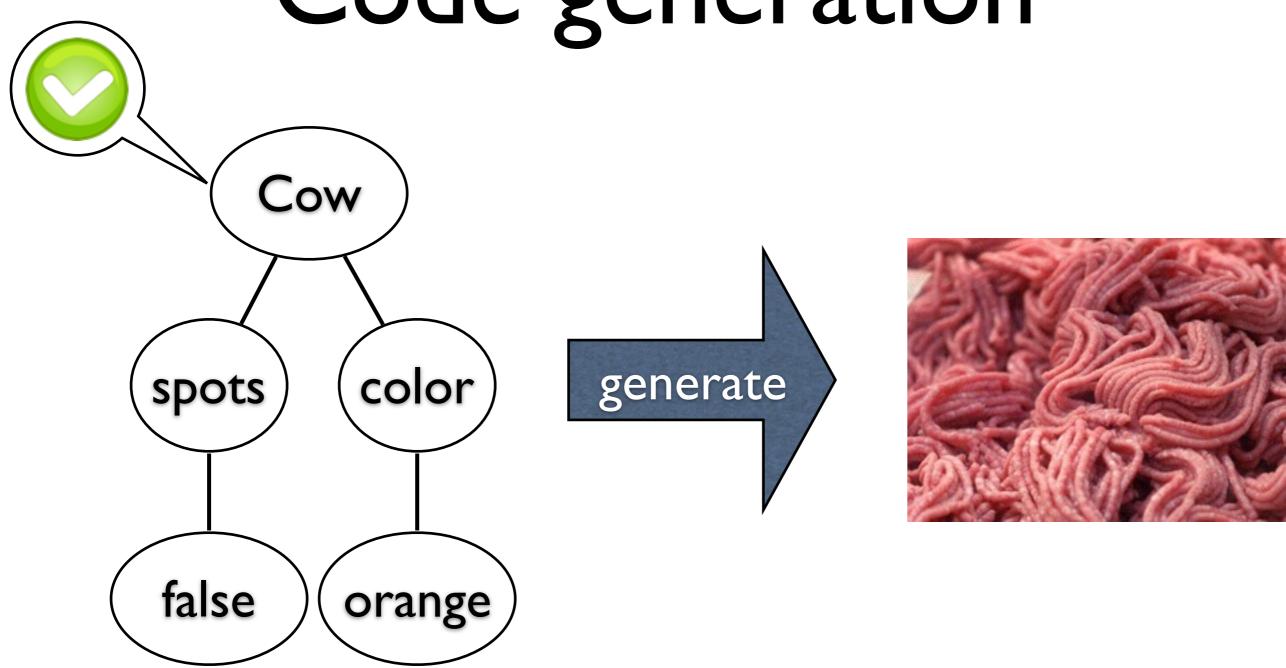


Semantic analysis

Constraint: $spots = true \Rightarrow color \neq orange$



Code generation



Tools

ANTLR
THE SETURN STATE

- Parser generators
- Attribute grammar systems
- Transformation systems
- Language workbenches









Xte≍t

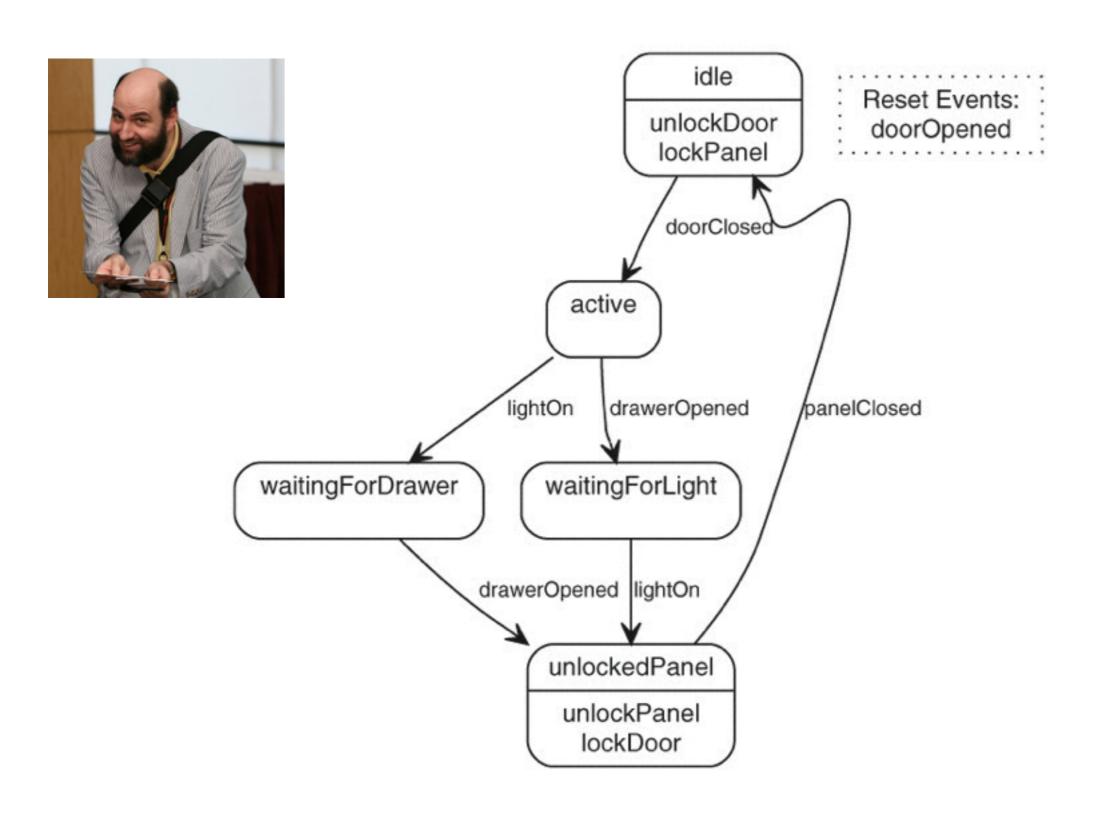








State machines

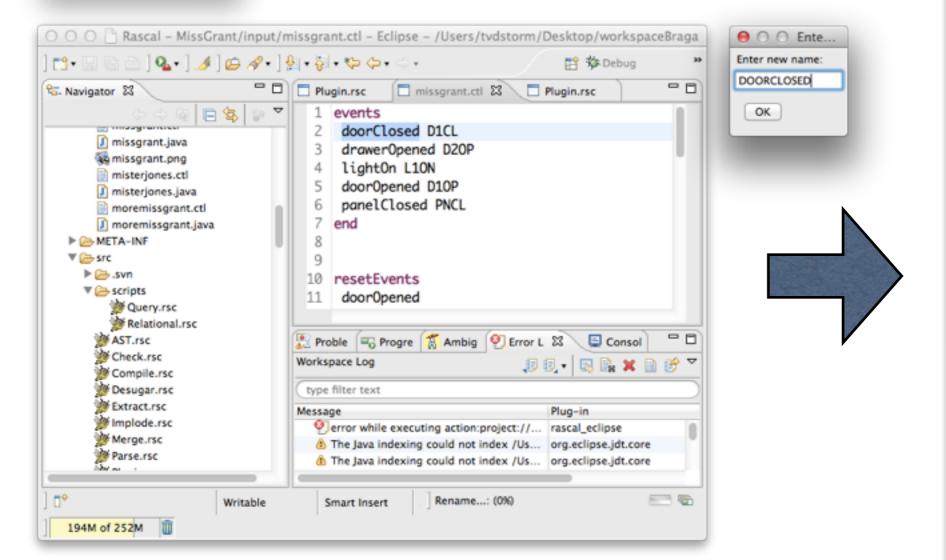


Textual notation

```
events
 doorClosed D1CL
 drawerOpened D2OP
 lightOn L10N
 doorOpened D10P
 panelClosed PNCL
end
resetEvents
door0pened
end
commands
 unlockPanel PNUL
 lockPanel PNLK
 lockDoor D1LK
 unlockDoor D1UL
end
state idle
 actions {unlockDoor lockPanel}
 doorClosed => active
end
```

```
state active
 drawerOpened => waitingForLight
lightOn => waitingForDrawer
end
state waitingForLight
lightOn => unlockedPanel
end
state waitingForDrawer
 drawerOpened => unlockedPanel
end
state unlockedPanel
 actions {unlockPanel lockDoor}
 panelClosed => idle
end
```

Visualize Rename...



events

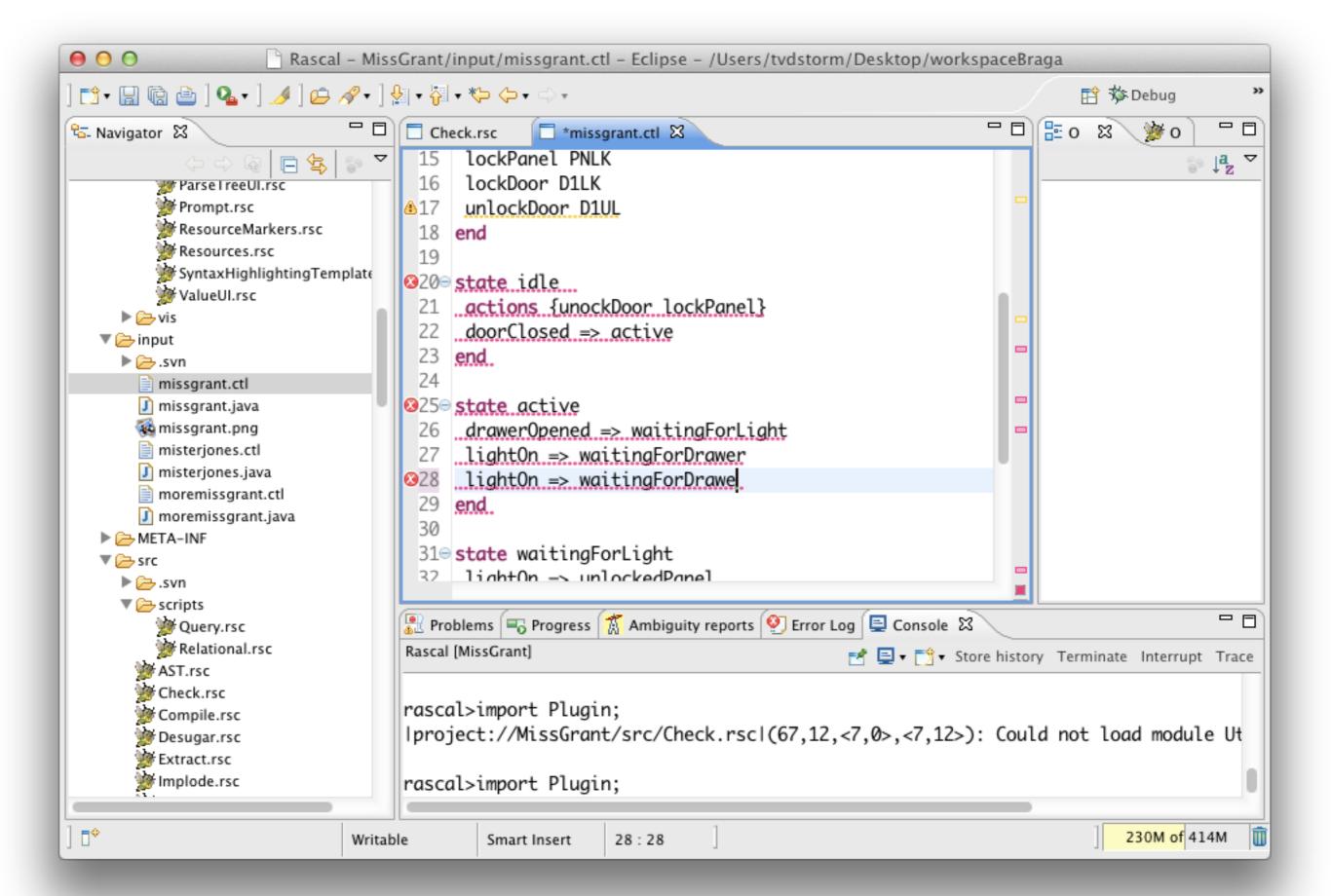
DOORCLOSED D1CL drawerOpened D2OP lightOn L1ON doorOpened D1OP panelClosed PNCL end

resetEvents doorOpened end

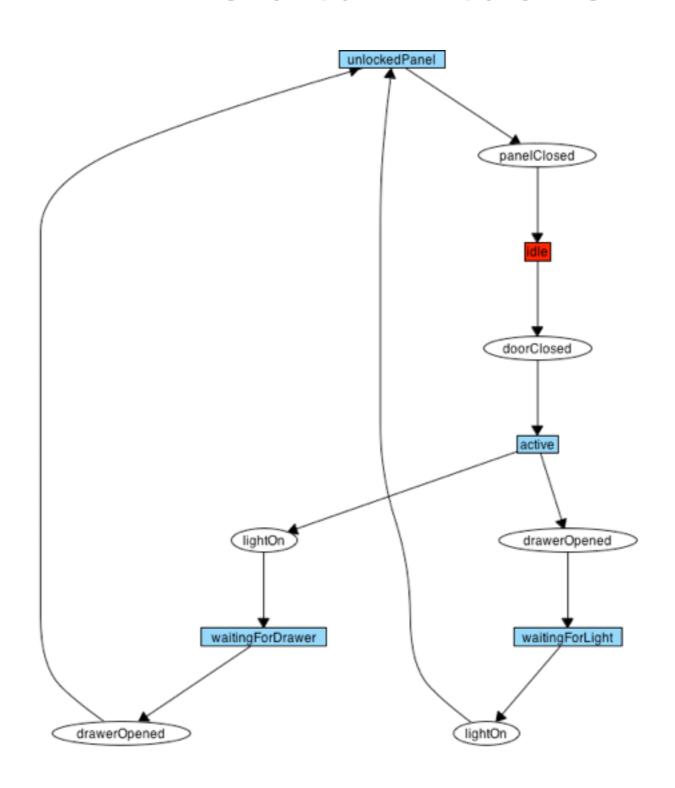
commands

unlockPanel PNUL lockPanel PNLK lockDoor D1LK unlockDoor D1UL end

state idle
 actions {unlockDoor lockPanel}
 DOORCLOSED => active
end

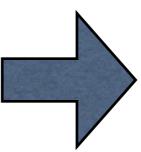


Visualization



Code generation

```
events
 doorClosed D1CL
 drawerOpened D20P
 lightOn L10N
 doorOpened D10P
 panelClosed PNCL
end
resetEvents
door0pened
end
commands
 unlockPanel PNUL
 lockPanel PNLK
 lockDoor D1LK
 unlockDoor D1UL
end
state idle
actions {unlockDoor lockPanel}
 doorClosed => active
end
state active
drawerOpened => waitingForLight
lightOn => waitingForDrawer
end
state waitingForLight
liahtOn => unlockedPanel
end
state waitingForDrawer
drawerOpened => unlockedPanel
end
state unlockedPanel
 actions {unlockPanel lockDoor}
 panelClosed => idle
end
```



```
public class missgrant {
      public static void main(String args[]) throws java.io.IOException {
            new missgrant().run(new java.util.Scanner(System.in),
                        new java.io.PrintWriter(System.out));
     private static final int state$idle = 0;
     private static final int state$active = 1;
     private static final int state$waitingForLight = 2;
     private static final int state$waitingForDrawer = 3;
     private static final int state$unlockedPanel = 4;
     public void run(java.util.Scanner input, java.io.Writer output)
                  throws java.io.IOException {
            int state = state$idle;
            while (true) {
                  String token = input.nextLine();
                  switch (state) {
                  case state$idle: {
                        unlockDoor(output);
                        lockPanel(output);
                        if (doorClosed(token)) {
                              state = state$active;
                        if (doorOpened(token)) {
                              state = state$idle;
                        break;
                 }
                  case state$active: {
                        if (drawerOpened(token)) {
                              state = state$waitingForLight;
                        if (lightOn(token)) {
                              state = state$waitingForDrawer;
                        if (doorOpened(token)) {
                              state = state$idle;
                        break;
                  }
                  case state$waitingForLight: {
```

Domain-specific languages

- Better languages, for specific domains
- Capture families of systems
- Higher level of abstraction
- Focus on "what" vs "how"
- Reuse designs, not just code
- Language workbenches (e.g., Rascal)

Take home

- How much code I write is actually relevant for the problem I'm solving?
- What are recurring patterns in the code I'm writing?
- How would I want to describe the solution?
- Could I formalize the relevant bits, ... in a DSL?

Thank you

- http://www.rascal-mpl.org
- http://www.cwi.nl/~storm
- storm@cwi.nl
- @tvdstorm

