# **Deplump for Streaming Data**

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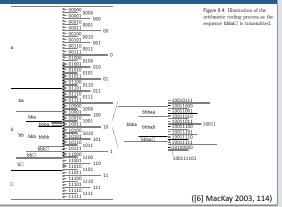


# Arithmetic Encoding [6]

An arithmetic encoder uses of the output from a predictive model to create a 1-1 correspondence between input streams and sub-intervals of the unit interval [0,1). The input is then encoded using the sub-interval to which it corresponds. Input streams which are likely under the predictive model correspond to larger sub-intervals and thus require fewer bits to

### For example:

Context (sequence thus far)	Probability of next symbol						
	P(a) = 0.425	P(b) = 0.425	$P(\Box) = 0.15$				
ъ	P(a   b) = 0.28	P(b   b) = 0.57	$P(\Box     \mathbf{b}) = 0.15$				
bb	P(a     bb) = 0.21	$P(\mathbf{b} \mathbf{bb}){=}0.64$	$P(\Box \mathrm{bb}){=}0.15$				
bbb	$P(\mathtt{a} \mathtt{bbb}){=}0.17$	$P(\mathbf{b} \mathbf{bbb}){=}0.68$	$P(\Box \mathtt{bbb}){=}0.15$				
bbba	$P(\mathtt{a} \mathtt{bbba}){=}0.28$	$P(\mathbf{b} \mathbf{bbba}){=}0.57$	$P(\Box \mathtt{bbba}){=}0.15$				



# Batch Deplump [4] and the Sequence Memoizer [11]

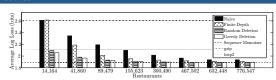
Batch deplump is an arithmetic compressor powered by a nonparametric Bayesian model called the sequence memoizer. The performance of deplump has been demonstrated on benchmark corpora, including the Calgary Corpus [2]. The aggregate performance of deplump is equal or better than that of comparable state of the art, general purpose, lossless compressors [4]. The spatial complexity of the sequence memoizer model grows unboundedly making batch deplump unrealistic for streaming data.

The performance of batch deplump on the Calgary Corpus is compared to PPM [3] and CTW [10]. Performance is measured in average bits per byte (lower is better). Bold text indicates best performance.

DEPLUMP PPM CTW								
		DEPLUMP						
File	Size	1PF	UKN	PPM*	PPMZ	CTW		
bib	111261	1.73	1.72	1.91	1.74	1.83		
book1	768771	2.17	2.20	2.40	2.21	2.18		
book2	610856	1.83	1.84	2.02	1.87	1.89		
geo	102400	4.40	4.40	4.83	4.64	4.53		
news	377109	2.20	2.20	2.42	2.24	2.35		
obj1	21504	3.64	3.65	4.00	3.66	3.72		
obj2	246814	2.21	2.19	2.43	2.23	2.40		
paper1	53161	2.21	2.20	2.37	2.22	2.29		
paper2	82199	2.18	2.18	2.36	2.21	2.23		
pic	513216	0.77	0.82	0.85	0.76	0.80		
progc	39611	2.23	2.21	2.40	2.25	2.33		
progl	71646	1.44	1.43	1.67	1.46	1.65		
progp	49379	1.44	1.42	1.62	1.47	1.68		
trans	93695	1.21	1.20	1.45	1.23	1.44		
avg.		2.12	2.12	2.34	2.16	2.24		
w. avg.		1.89	1.91	2.09	1.93	1.99		

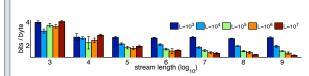
([4] Gasthaus, J.; Wood, F. and Teh, Y. W.)

### Results 1



The sequence memoizer model makes use of a suffix tree data structure. To create a streaming deplump compressor it is necessary to approximate the model using a data structure which does not grow with the length of the input sequence. Forgetting (pruning) of the suffix tree is used to achieve this and was demonstrated to have excellent empirical performance [1]. Results here are measured in bits per byte and shown versus an upper limit on the number of nodes in the suffix tree. Comparisons are made to naïve and other simple strategies to obtain constant spatial complexity.

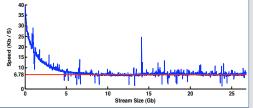
Streaming deplump was evaluated on a complete Wikipedia .xml dump [9]. For this result 10 100MB chunks of text were sampled with replacement and compressed using models limited to depths varying from 2 to 32. Performance is shown in bits per byte and each group contains results for models with a different upper limit on the node count of the data structure. As expected, larger models generally perform better. Using a larger depth appears advantageous up to ≈ 16.



The performance of streaming deplump as a function of stream length was also evaluated using the Wikipedia .xml dump. For this result stream lengths ranging from 10<sup>3</sup> to 10<sup>9</sup> bytes were compressed using models of varying size. Performance is again measured in bits per byte. The performance clearly improves as the length of the sequence increases.

## **Linear Time Verification**

Streaming deplump has asymptotic properties appropriate for streaming data. Shown here is the speed of the compressor on the entire Wikipedia corpus. The speed of the compressor is plotted with the size of the input stream. After an initial period the speed remains constant as the stream length increases. Streaming deplump compresses the 26.8Gb corpus to 4.0Gb, compared to 7.8Gb with gzip and 3.8Gb with paq9a.



# **Additional Approximations**

Representation

- · Each node can be represented in a constant amount of space [5].
- · Each node is associated with a context observed in the input sequence.
- The number of times a given context has been observed in the input sequence grows unboundedly with the length of the input sequence.
- . Operations on the suffix tree necessary for incremental construction and estimation of the model require space and time which grows with the number of times contexts have been observed in the input sequence [5].
- · Node representations are approximated by placing an upper bound on the recorded
- number of occurrences in the input sequence of any given context.

### Suffix Tree

- · Suffix tree data structures label edges using pointers into the input sequence.
- . The input sequence grows linearly.
- The removal of nodes through the pruning mechanism does not guarantee that the edges will be labeled using only a small subset of the original input sequence.
- The suffix tree data structure is approximated by only allowing edges to be labeled by a fixed length suffix of the input sequence.
- · Edges which cannot be labeled are removed from the model along with the descending sub tree.
- . To minimize the impact of this approximation edges are updated incrementally to point to later sections of the input sequence.

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[3] Cleary, J. G. and Teahan, W. J. Unbounded length contexts for PPM. The Computer Journal, 1997, 40:67–75.

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[10] Willems, F. M. J. , 2009. CTW website. URL: http:// www.ele.tue.nl/ctw/.

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## Hierarchical Pitman-Yor Processes (HPYP) [7,8]

The PYP is a distribution over distributions and is a generalization of the Dirichlet process [7]. Hierarchically composing PY processes is a way to smooth distribution estimates. The parameters of the process are known as discounts and concentrations and control the amount of smoothing in the model.

$$\begin{array}{cccc} \mathcal{G}_1|d_1,c_1,\mathcal{G}_0 & \sim & \mathcal{PY}(d_1,c_1,\mathcal{G}_0) \\ \\ \mathcal{G}_2|d_2,c_2,\mathcal{G}_1 & \sim & \mathcal{PY}(d_2,c_2,\mathcal{G}_1) \\ \\ \theta_i|\mathcal{G}_2 & \sim & \mathcal{G}_2 & i=1,\ldots,N. \end{array}$$

