## Trees review

Delicition A tree is a graph in there is one, and only one, path between any pair of nodes.

- Trees do not have loops
- Moralization of a directed tree
results in an undirected tree

HExact Inference is possible on trees using an algorithm much like that presented for chairs. Sum-product or B. P. Graphs with loops are now complicated - loopy I.P.

Before sum-product?

- Directed & andirected graphs express a global function as the product of factors over subsots of those vars

- Factors walce explicit modes for the factors

- Factors welce explicit modes for the factors in the product (in addition to the variables modes)

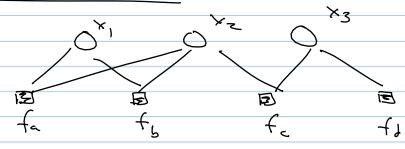
Joint can be written es

where is a subset of the variables

Individual varis are xi, xi way be complicated is metrix, vector, etc.

Where did the normalizing co-st. go?
- factor over an empty set of vars.

## Factor graph example

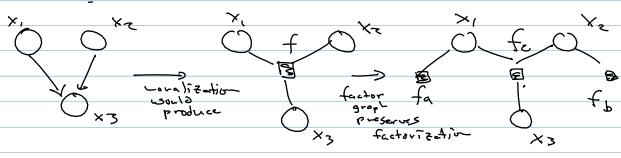


P(x) = fa(x, x2) fb(x,x2) fc(x2,x3) fd(x3)

A factor graph consists of

I circle for each variable
I squere for each factor

A factor graph is a more explicit characterization of the factorization of the joint distribution, keeping separate terms that might be lumped together in an undirected g.m. for instance



P(x)= P(x1) P(x2) P(x3 | x1, x2) -= P(x1) P(x2) P(x3 | x1, x2)

- can be reasons to

for 
$$p(x_1)$$

combine factors in this

 $f_b = p(x_2)$ 
 $f_c = p(x_1, x_2, x_3)$ 

The Sur-product algorith (Belief prop.) Powerful algorithm for inference off tree-structured factor graphs. Assurption all modes discrete

(could be linear-Gaussian lumerar)

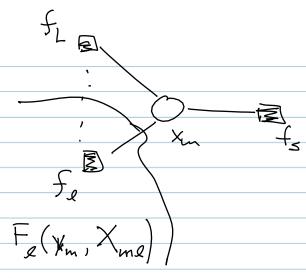
Recipe

- Convert graph to factor graph.

- Pass messages

o two kinds of messages in sum-prod. algorithm Start as before with fi-di-s a single warsi-al dist:  $p(x) = \sum_{X \mid X} p(\overline{X}) \qquad \stackrel{?}{X} \mid X \text{ is}$  co-sideri s' $F_s(x, X_s)$   $f_s = f_s$ One can directly see that  $P(\overset{>}{\times}) = \prod_{s \in X} T_s(x, X_s)$ where nc(x) are the factor node neighbors
of x and  $X_s$  denotes the set of all
vars in the subtree connected to x through F<sub>5</sub>(x, X<sub>5</sub>) = f<sub>5</sub>(x, x, .... x<sub>m</sub>) G(x, X<sub>5</sub>) .... G<sub>m</sub>(x<sub>m</sub>, X<sub>5</sub>m) is the prod of all factors in group' associated will for

We exploit this factorization to co-pute  $M_{f_s \to x}(x) = \sum_{X_s} F_s(x, X_s)$  $= \sum_{X_s} f_s(x_1, x_1, \dots, x_m) G_1(x_1, X_{s_2}) \cdots G_M(x_m, X_{s_M})$  $= \underbrace{\sum_{x_1} \underbrace{\sum_{x_m} f_s(x, x_1, ..., x_m)}_{m \in Ne(f_m)} \underbrace{\sum_{x_m} G_m(x_m, \underbrace{X_{s_m}})}_{m \in Ne(f_m)}$ = \( \sim \) \( \sim \ m & ne (fs) x all variables in Factor Fs Implicitely defined a new message Mxm > fs (xm) = \( \int \text{Gm}(xm, \overline{\text{Xsm}}) \)
which is a ressage from a var node to a factor node. How do we co-pute Mxm >f (xm)?



From this we can see that  $G_m(x_m, \overline{X_{sm}})$  can be written as another product of factors

$$G_m(x_m, X_m) = \prod_{l \in ne(x_m)} f_l(x_m, X_me)$$

So we can write

- Collecting of lits yields message passing

Message Types

 $\longrightarrow 0 \qquad \underset{f_s \to x}{\mu_{f_s \to x}} (x) = \underbrace{\sum_{x_1, \dots, x_M} f_s(x_1, x_1, \dots, x_M)}_{x_M} \underbrace{\prod_{x_m \to f_s(x_m)} \mu_{x_m \to f_s(x_m)}}_{\text{mene}(f_s) \to \infty}$ 

Mxm=fs(xm)= ||
lene(xm) fs Mfe=xm(xm)

To start, messages can be sent from leaf modes

 $\begin{array}{c}
M_{x \to f}(x) = I \\
X \longrightarrow f \\
X \longrightarrow f
\end{array}$ 

- Pecap

   Massage passing for evaluating marginal

  p(x)
  - Ptck x as root messages from leaves
    - Once root receives all messeges, pass all messages back to root
    - Easy to see that this yields valid algorithment with enough we ssages always available jit.

$$\frac{E_{xa-p}k}{\sum_{\alpha} f_{\alpha}}$$

$$\frac{\sum_{\alpha} f_{\alpha}}{\sum_{\alpha} f_{\alpha}}$$

$$\frac{\sum_{\alpha$$

Designate  $x_3$  as root and process wesseges from  $\mu_{x_1} \Rightarrow f_{c_1}(x_1) = 1$ lewes

$$M_{f_{a}} \rightarrow x_{z} (x_{z}) = \sum_{x_{1}} f_{a}(x_{1}, x_{z})$$

$$M_{x_{4}} \rightarrow f_{c}(x_{4}) = 1$$

$$M_{f_{c}} \rightarrow x_{z} (x_{z}) = \sum_{x_{4}} f_{c}(x_{z}, x_{4})$$

$$M_{x_{2}} \rightarrow f_{b}(x_{z}) = M_{f_{a} \rightarrow x_{z}}(x_{z}) M_{f_{c} \rightarrow x_{z}}(x_{z})$$

$$M_{f_{b} \rightarrow x_{3}}(x_{3}) = \sum_{x_{2}} f_{b}(x_{2}, x_{3}) M_{x_{2} \rightarrow f_{b}}$$

Now other way

$$M_{\chi_{3} \to f_{b}}(\chi_{3}) = 1$$

$$M_{f_{b} \to \chi_{z}}(\chi_{z}) = \sum_{\chi_{3}} f_{(\chi_{z} \times \chi_{3})}$$

$$M_{\chi_{z} \to f_{a}}(\chi_{z}) = M_{f_{b} \to \chi_{z}}(\chi_{z}) M_{f_{c} \to \chi_{z}}(\chi_{z})$$

$$M_{\chi_{z} \to f_{a}}(\chi_{z}) = M_{f_{b} \to \chi_{z}}(\chi_{z}) M_{\chi_{z} \to f_{a}}(\chi_{z})$$

$$M_{\chi_{z} \to f_{c}}(\chi_{z}) = M_{f_{c} \to \chi_{z}}(\chi_{z}) M_{\chi_{z} \to f_{c}}(\chi_{z})$$

$$M_{\chi_{z} \to f_{c}}(\chi_{z}) = M_{f_{c} \to \chi_{z}}(\chi_{z}) M_{\chi_{z} \to f_{c}}(\chi_{z})$$

$$M_{f_{c} \to \chi_{y}}(\chi_{y}) = \sum_{\chi_{z}} f_{c}(\chi_{z}, \chi_{y}) M_{\chi_{z} \to f_{c}}(\chi_{z})$$

$$M_{f_{c} \to \chi_{y}}(\chi_{y}) = \sum_{\chi_{z}} f_{c}(\chi_{z}, \chi_{y}) M_{\chi_{z} \to f_{c}}(\chi_{z})$$

To evaluate  $\vec{p}(x_2)$  for instance we use

$$P(x) = \prod_{s \in Ne(x)} M_{fs \to x}(x)$$

$$= \left[ \sum_{x_1} f_a(x_1, x_2) \right] \left[ \sum_{x_3} f_b(x_2, x_3) \right] \left[ \sum_{x_4} f_c(x_2, x_4) \right]$$

$$= \underbrace{\lesssim}_{x_1} \underbrace{\lesssim}_{x_2} \underbrace{f_{\alpha}(x_1,x_2)} \underbrace{f_{\beta}(x_2,x_3)} \underbrace{f_{\alpha}(x_2,x_3)} \underbrace{f_$$

Conditioning on observed values yields the same kind of indicator function approach.

- Discrete variable assumption not strictly necessary.

Technique extends to continuous real valled vary, etc.

Loopy B.P.
What abouts sraphs with loops?

loops until "convergence."