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Controller Area Network (CAN)

Distributed Embedded Systems

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October 8, 2014

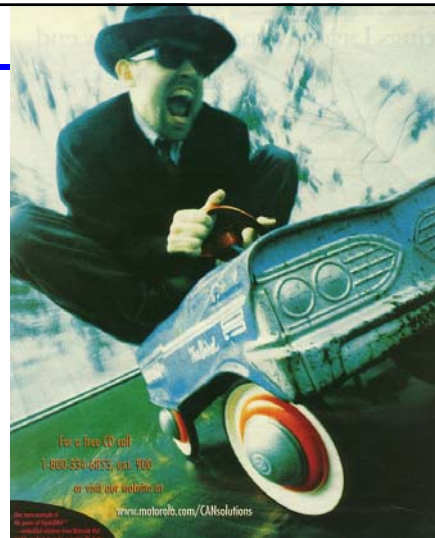
Significant material (CAN pictures) drawn from
a presentation by Siemens Corp. "CANPRES 2.0, Oct 1998"

**Carnegie
Mellon**

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Where Are We Now?

- ◆ **Where we've been:**
 - Protocol Overview
- ◆ **Where we're going today:**
 - CAN -- an important embedded protocol
 - Primarily automotive, but used in many places
- ◆ **Where we're going next:**
 - CAN performance
 - Other protocols
- ◆ **REMINDER – look at lessons learned slides for ideas on how to do better on second half of project!**



Preview

◆ CAN – important automotive protocol

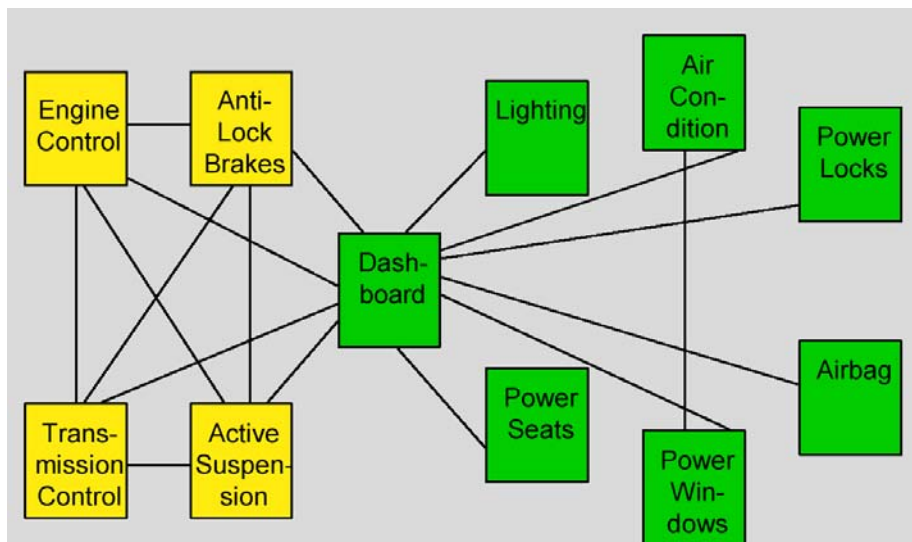
- Physical layer – built on bit dominance
- Protocol layer – binary countdown
- Message filtering layer (with add-on protocols)

◆ Keep an eye out for:

- Message prioritization
- How “small” nodes can be kept from overloading with received messages
- Tradeoffs

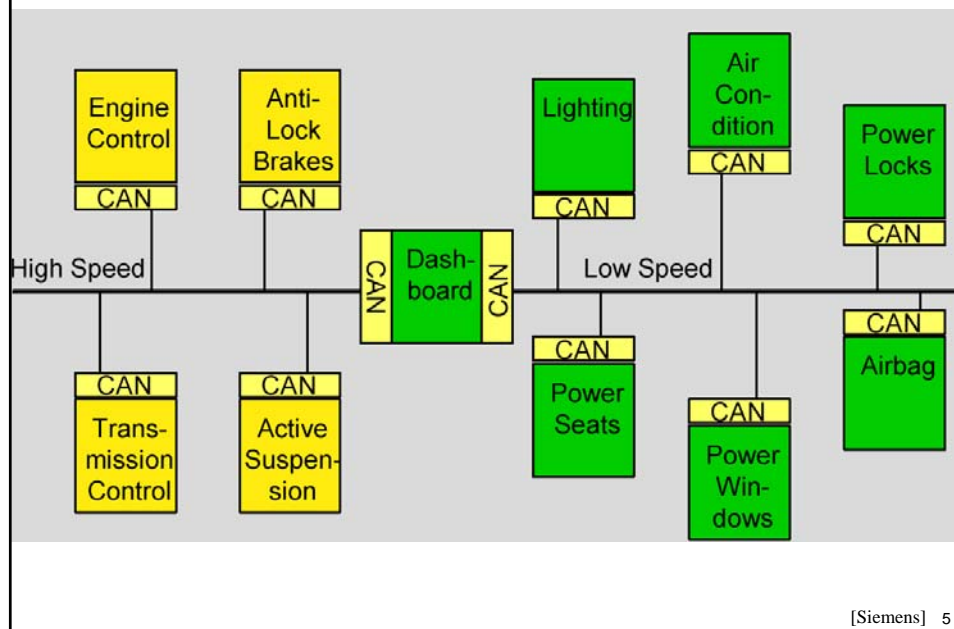
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Before CAN

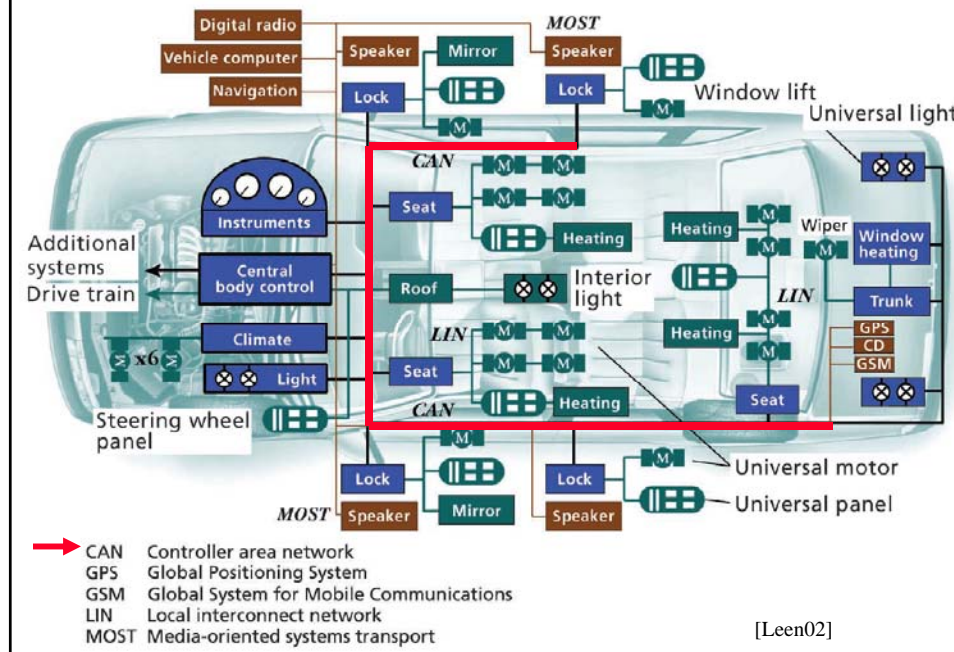


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With CAN



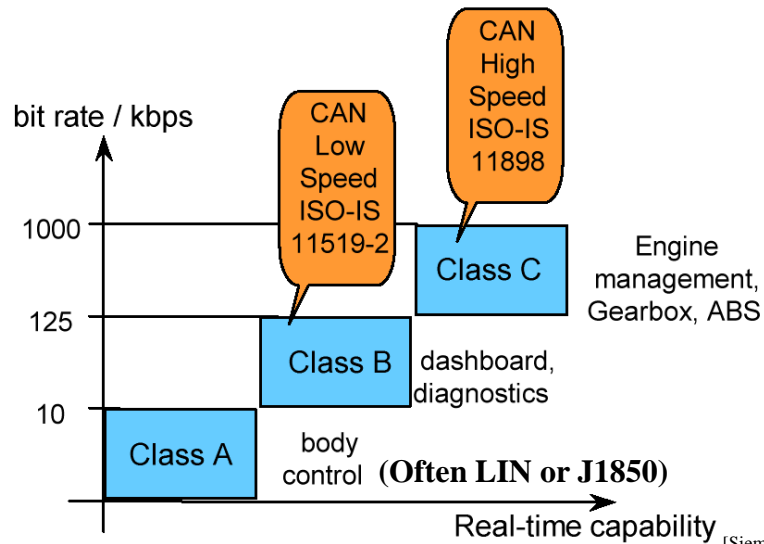
CAN Is Central To Automotive Networks



SAE Message Classes

◆ Fast tends to correlate with critical control

- But, this is not always true; just often true



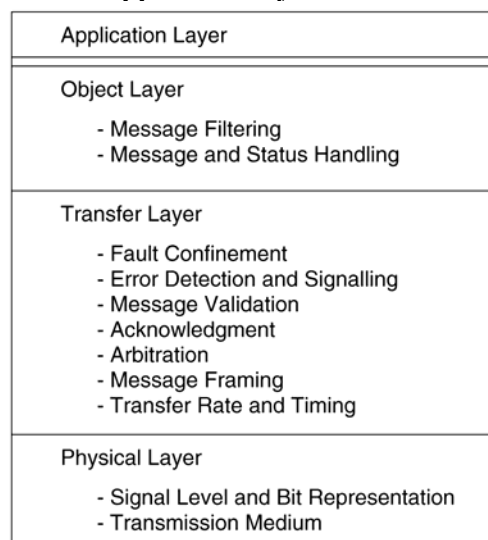
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CAN & the Protocol Layers

◆ CAN only standardizes the lower layers

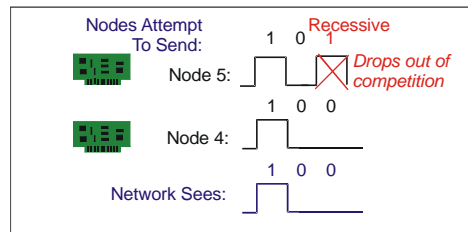
◆ Other high-level protocols are used for application layer

- User defined
- Other standards
- We'll see one possibility at the end of this lecture



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Remember This? Binary Countdown



◆ Operation

- Each node is assigned a unique identification number
- All nodes wishing to transmit compete for the channel by transmitting a binary signal based on their identification value
- A node drops out the competition if it detects a dominant state while transmitting a passive state
- Thus, the node with the **lowest** identification value wins

◆ Examples

- CAN – 500 Kbps or 1 Mbps
- SAE J1850 – pretty much same as CAN, except slower (around 10 Kbps)

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CAN – Bit Dominance In More Detail

◆ CAN uses the idea of recessive and dominant bits

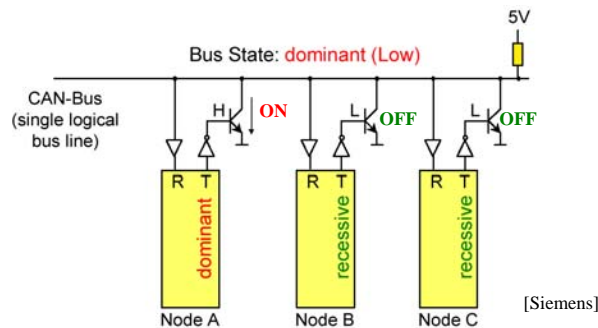
- Wired “OR” design
- Bus floats high unless a transmitter pulls it down (dominant)
- (Other bus wire in differential transmission floats low and transmitter pulls up)

◆ High is “recessive” value

- Sending a “1” can’t override the value seen on the bus

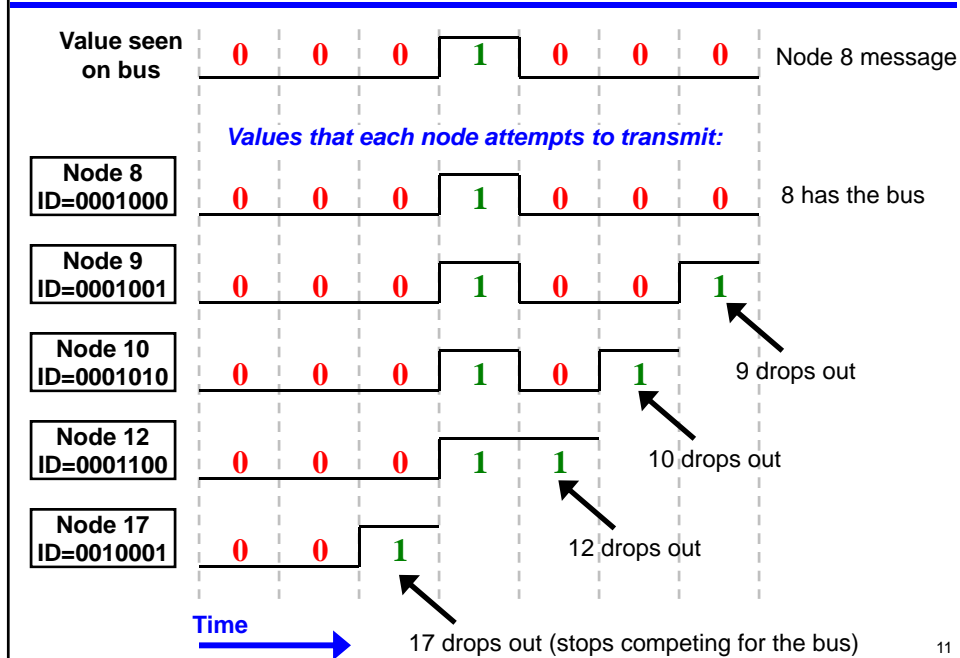
◆ Low is “dominant” value

- Sending a “0” forces the bus low no matter what another node is sending



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Example: Binary Countdown (highest bit first)



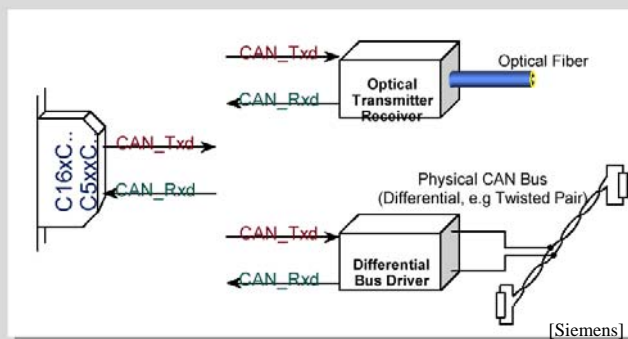
Physical Layer Possibilities

◆ MUST support bit dominance

- Specifically rules out transformer coupling for high-noise applications
- Differential driver used
 - Voltage across wires is dominant; high impedance (0V differential) is recessive
 - Opto-isolators are commonly used as well

□ Usual ISO Physical Layer :-

- Bus wires twisted pair, 120R Termination at each end
- 2 wires driven with differential signal (CAN_H, CAN_L)

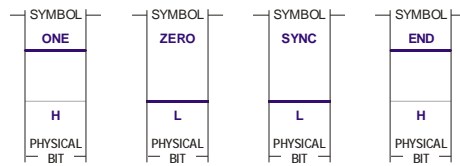


Non-Return to Zero (NRZ) Encoding

◆ Send a Zero as LO; send One as HI

- Worst case can have all zero or all one in a message – no edges in data
- Simplest solution is to limit data length to perhaps 8 bits
 - SYNC and END are opposite values, guaranteeing two edges per message
 - This is the technique commonly used on computer serial ports / UARTs
- Bandwidth is one edge per bit
 - Same bandwidth as Miller encoding, but no guarantee of frequent edges

Simple NRZ Bit Encoding



Simple NRZ Encoding Example: 11010001



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Bit Stuffing To Add Edges To NRZ Encoding

◆ Long NRZ messages cause problems in receivers

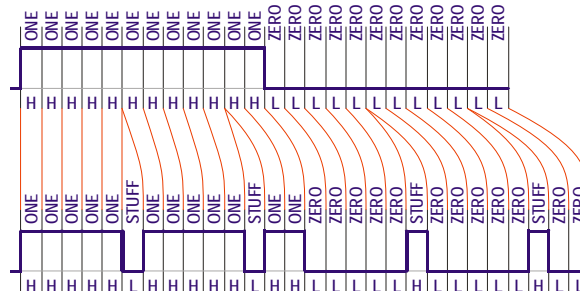
- Clock drift means that if there are no edges, receivers lose track of bits
- Periodic edges allow receiver to resynchronize to sender clock

◆ Solution: add “stuff bits”

- Stuff bits are extra bits added to force transitions regardless of data
- Typical approach: add an opposite-valued stuff bit after every 5 identical bits
- In best case you don't need stuff bits – they only are needed for runs of values

BIT STUFF IDEA:

SIMPLE NRZ ENCODING OF: 1111 1111 1111 0000 0000 0000:



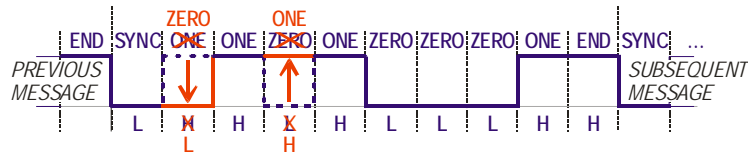
BIT-STUFFED NRZ ENCODING OF: 1111 1111 1111 0000 0000 0000:

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NRZ Encoding Error Susceptibility

◆ A single inverted physical bit is undetectable with Simple NRZ

- High efficiency comes at price of poor error detection



- (Can be detected via CRC sometimes; but CRCs have limitations)

◆ Bit stuffing error detection in general case:

- Improves error detection if stuffing rule is violated
- Any six identical data bits in a row is an stuffing error
- But, there is a subtle problem with bit stuffing...

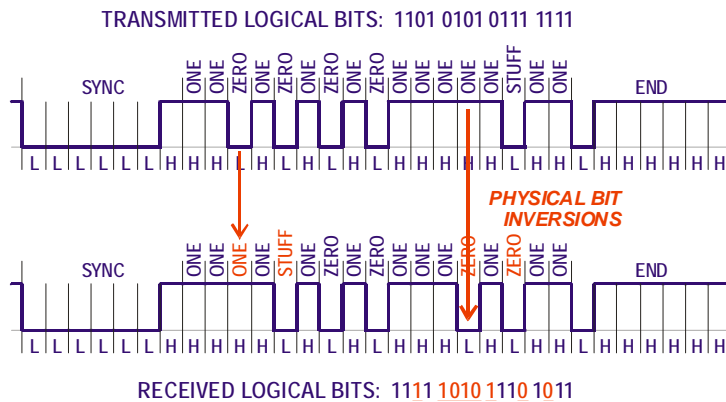
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Cascaded Bit Stuffing Errors

◆ Bit inversions in just the wrong place can confuse bit stuffing logic

- Worst errors occur in pairs that create and then break runs of bits
- Data bit is converted to stuff bit; stuff bit to data bit
- Net effect is same message length BUT, it shifts intervening data bits
- CAN has this problem; can cause 2-bit error to escape CRC detection!

Cascaded bit stuff error example:



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General CAN Message Format

SYNC	HEADER	DATA	ERROR DETECTION	END
------	--------	------	-----------------	-----

◆ Header

- Application can set any desired value in 11- or 29-bit header
- Global priority information (which message gets on bus first?)
- Header often contains source, destination, and message ID

◆ Data

- Application- or high-level-standard defined data fields
- 0 to 8 bytes of data for CAN

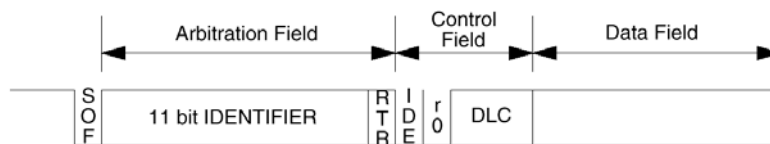
◆ Error detection

- Detects corrupted data (uses a 15-bit CRC):
 - All 15-bit or shorter burst errors (groups of flipped bits clumped together)
 - All 5-bit errors regardless of where they occur ...
 - ... except bit stuffing problem reduces this to all 1-bit errors

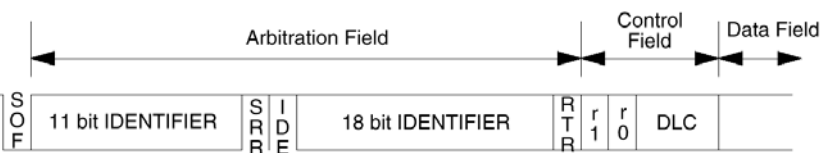
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Two Sizes of CAN Arbitration Fields

Standard Format



Extended Format



[Bosch] 18

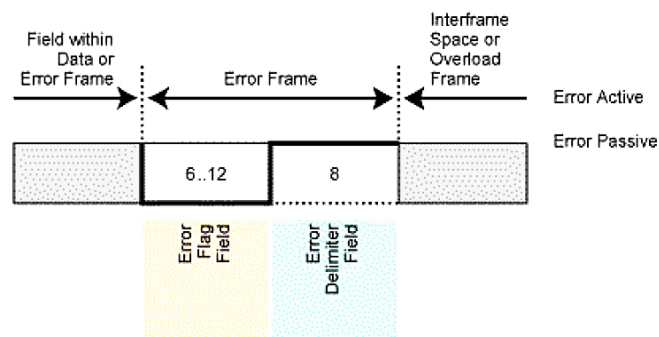
CAN Message Fields

- ◆ **SOF – Start of frame (SYNC symbol)**
 - Single dominant bit
- ◆ **Arbitration field – binary countdown priority value; set by application**
 - Also an RTR (remote transmission) field for atomic transactions; seldom used
 - SRR is a dummy bit to let standard format RTR messages win arbitration
- ◆ **Control field**
 - 4-bit data length (number of bytes in data field); valid values: 0 .. 8
 - 1 bit specifies standard or extended format; 1 bit unused
- ◆ **Data field**
 - 0 to 8 bytes
- ◆ **CRC field**
 - 15-bit CRC, followed by one recessive delimiter bit
- ◆ **Ack field**
 - If message received OK, assert as dominant bit (**at least one node** received)
- ◆ **END of frame delimiter**
 - Seven recessive bits mark end of frame (phase violation for bit stuff pattern)

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Error Frame Messages

- ◆ **Error frame alerts transmitter if message garbled at some receivers**
 - Sent if bit stuff violation detected or CRC error detected
 - Error flag is six dominant bits in a row – guaranteed to violate bit stuffing rules
 - (Unless the Error Frame itself suffers a bit error)
 - If transmitter sees that it has been pre-empted by an error frame, it attempts retransmission
 - Note – this is a source of nondeterminism in protocol – timing varies depending on errors encountered!

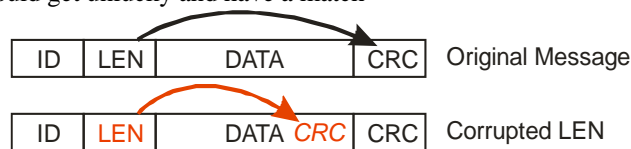


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CAN vs. FlexRay Length Field Corruptions

◆ CAN does not protect length field against ONE-BIT errors

- Corrupted length field will point to wrong location for CRC!
 - One bit error in length field circumvents HD=6 CRC
 - Could get unlucky and have a match



◆ FlexRay solves this with a header CRC to protect Length

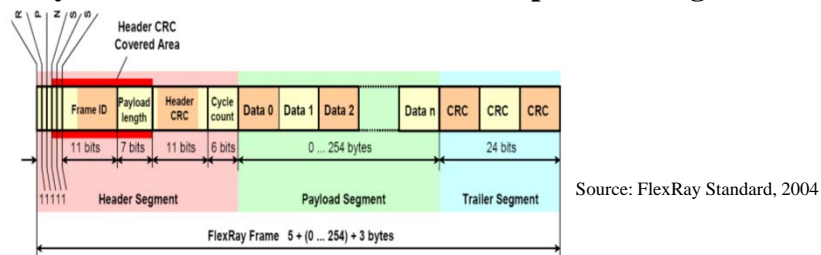


Figure 4-1: FlexRay frame format.

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Other CAN Issues

<http://betterembsw.blogspot.com/2012/02/can-protocol-vulnerabilities.html>

◆ CAN advertises “exactly once” delivery semantics

- A message corrupted only in some places is retransmitted...
... and received more than once by some nodes

[Rufino 1998]

◆ “Stuck at zero” (dominant) transmitter output locks up network

◆ CAN retry:

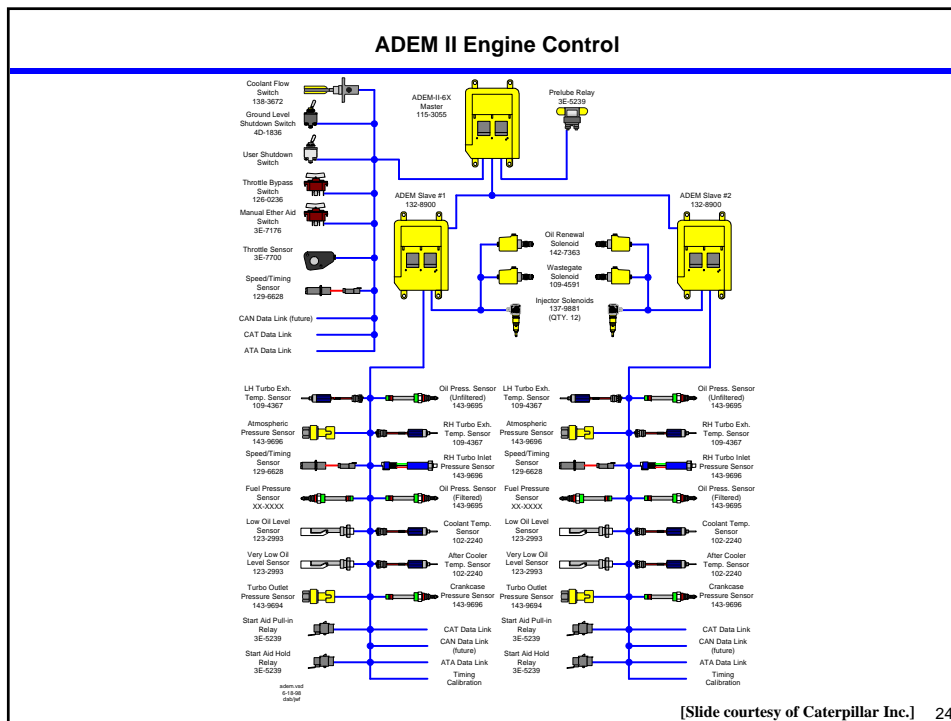
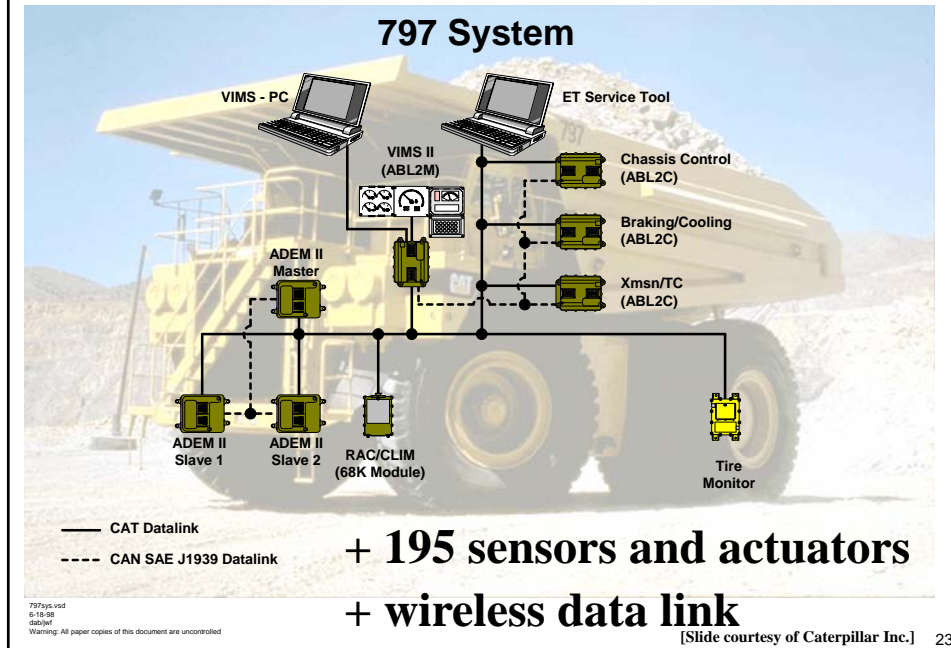
- “Error frame” scheme causes re-transmit
and ALSO,
- Node monitors network and looks for data sent == data received
 - If no match, assumes corruption and tries again
- But what if the transmitter is what is broken?
 - CAN node can lock up the network with retries [Perez 2003]

◆ In general, CAN unsuitable for highly critical applications

- That’s one reason we have FlexRay (and TTP) protocols
- This is news to some embedded folks (e.g., ARINC 825 aviation standard)

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CAN (SAE J1939) Example: Caterpillar 797



More Than Just Vehicles

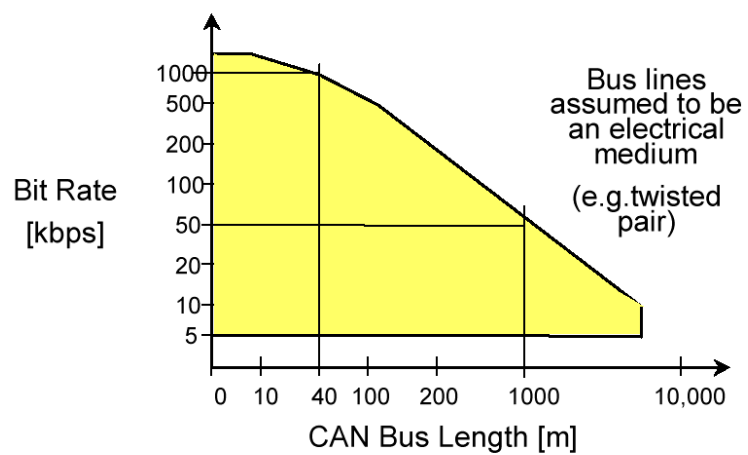
Liebert Modular UPS
(uses CAN)



Arbitration Limits Network Size

◆ Need $2 \cdot t_{pd}$ per bit maximum speed

□ Up to 1Mbit / sec @40m bus length (130 feet)



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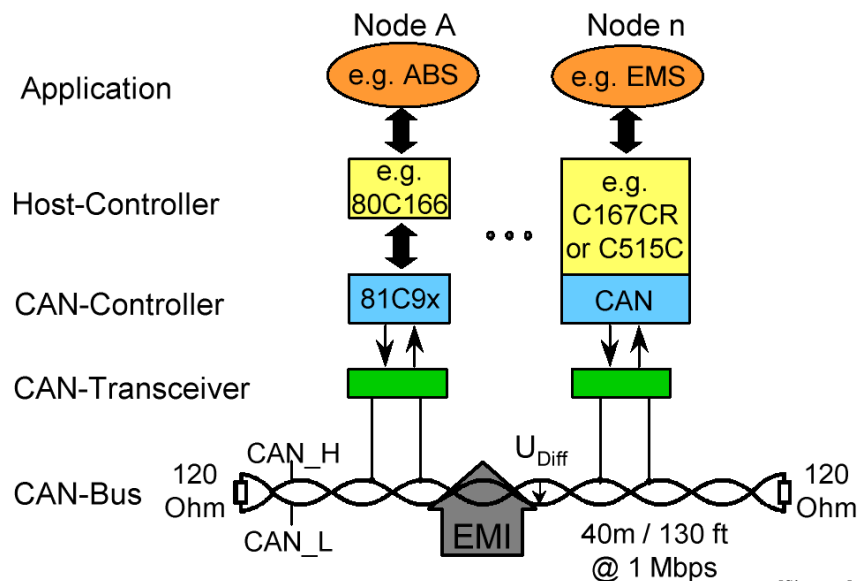
“Big” & “Small” Nodes

- ◆ **Some nodes can handle a lot of messages**
 - Many message mailboxes/filters
 - Fast processor
- ◆ **Some small nodes have limited capacity**
 - One or two mailboxes/filters
 - Slow processor
- ◆ **System designer has to prevent message over-run via one of:**
 - Dedicated mailbox per message (hardware ensures no data lost)
 - If mailbox shared, ensure messages to slow processors are spaced apart
 - Must be infrequent
 - Must *ALSO* not be clumped closer than receiver response time
 - This ends up being a constraint for real time scheduling (a later lecture)

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Generic CAN Network Implementation

- ◆ **Signals usually sent differentially – CAN_H and CAN_L**



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Example CAN Microcontrollers

◆ Motorola 68HC05 Family

- 11-bit headers; 1 Tx buffers; 2 Rx message buffers; 8-bit accept mask
- 8-bit CPU; up to 32 KB on-chip ROM; 28- or 64-pin housing
- (Also 68HC08 with 29-bit support and more buffers)

◆ Motorola 68HC912 Family

- 11- & 29-bit headers; 3 Tx buffers; 2 Rx message buffers; 2 accept masks
- 16-bit CPU; up to 128 KB on-chip Flash; 80- or 112-pin housing

◆ Motorola 6837X Family

- 11- & 29-bit headers; 16 Tx/Rx buffers; 16 accept masks
- 32-bit CPU; 256 KB on-chip Flash

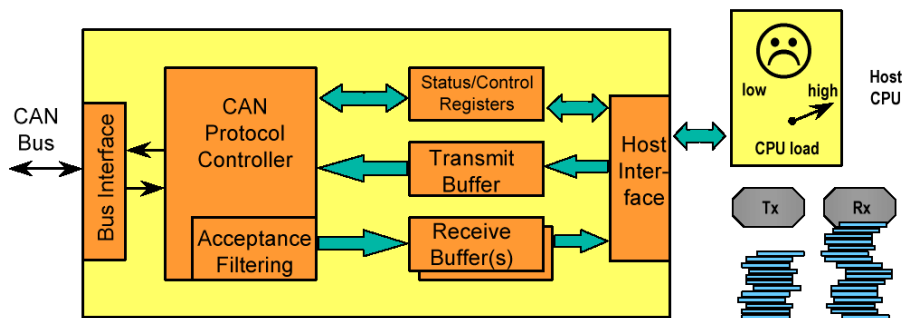
◆ Many other companies support CAN of course – these are just examples

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Basic CAN Controller (avoid this one if possible)

◆ “Cheap” node

- Could get over-run with messages even if it didn't need them

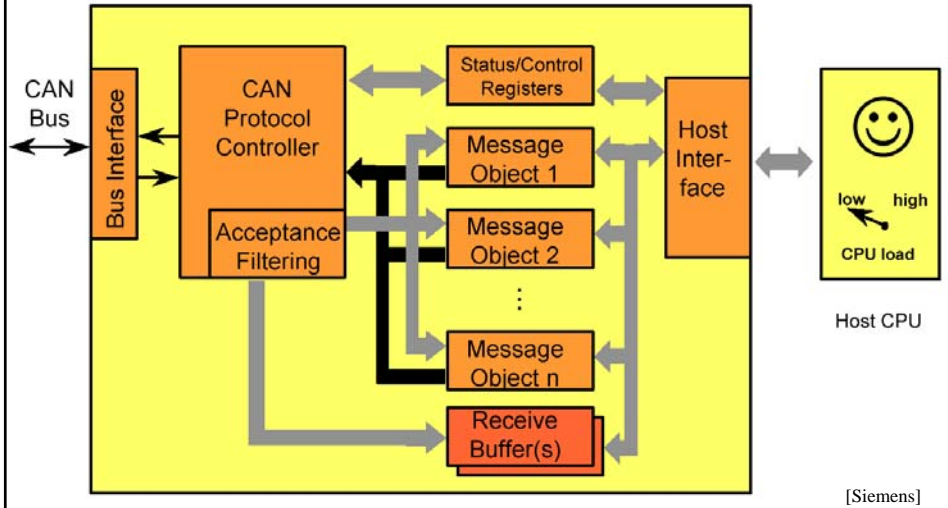


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Full CAN Controller

◆ Hardware message filters sort & filter messages without interrupting CPU

- Message object holds most recent message for that type – not a queue!



Mask Registers

◆ Used to set up message filters

- Mask register selects bits to examine
- Object Arbitration register selects bits that must match to be accepted
- Map multiple messages into each message object “mailbox”

Mask Register (std ID)	10 1 1 1 1 1 1 1 0 1 1 0 0	(= 0x7f6)
Message Object Arbitration Register	1 0 0 1 0 0 1 0 0 0 0	(= 0x490)
Resulting ID matches (X = don't care)	1 0 0 1 0 0 1 X 0 0 X	
ID's received:	1 0 0 1 0 0 1 0 0 0 0	(= 0x490)
	1 0 0 1 0 0 1 0 0 0 1	(= 0x491)
	1 0 0 1 0 0 1 1 0 0 0	(= 0x498)
	1 0 0 1 0 0 1 1 0 0 1	(= 0x499)

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Mask Register Example

- ◆ **Mask Register:** 1 1 0 1 1 1 0 1 0 1 1
- ◆ **Message Object Arbitration:** 1 0 0 0 1 0 0 0 0 0 1
- ◆ **Effective Match Value:** 1 0 * 0 1 0 * 0 * 0 1

- ◆ **Matches these message IDs:** 1 0 0 0 1 0 0 0 0 0 1
- 1 0 0 0 1 0 0 0 1 0 1
- 1 0 0 0 1 0 1 0 0 0 1
- 1 0 0 0 1 0 1 0 1 0 1
- 1 0 1 0 1 0 0 0 0 0 1
- 1 0 1 0 1 0 0 0 1 0 1
- 1 0 1 0 1 0 1 0 0 0 1
- 1 0 1 0 1 0 1 0 1 0 1

- ◆ **More likely, you mask a few bits next to each other**
 - See DeviceNet later in lecture

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DeviceNet

- ◆ **One of several higher-level protocols**
 - Based on top of CAN
 - Used for industrial control (valves, motor starters, display panels, ...)
 - Caterpillar is a member of ODVA as well (Open DeviceNet Vendors Assn.), but for factory automation.
- ◆ **Basic ideas:**
 - CAN is used in high volumes = cheaper network chips than competitors
 - Use structured approach to message formats to standardize operation
- ◆ **Does *NOT* standardize specific message contents**
 - But it does specify a hierarchy of message ID formats

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DeviceNet Message ID Scheme

- ◆ Each node on network “owns” a source node or message ID (or both)

Message Identifier Bits

10	9	8	7	6	5	4	3	2	1	0	Hex Range	Identity Usage
0	Message ID				Source Node #						000 - 3ff	Group 1
1	0	Source Node #						Msg ID			400 - 5ff	Group 2
1	1	Msg ID (0..6)			Source Node #						600 - 7bf	Group 3
1	1	1	1	1	Message ID (0..2f)						7c0 - 7ef	Group 4
1	1	1	1	1	1	1	X	X	X	X	7f0 - 7ff	Invalid

- ◆ Use message filters to only listen to messages you care about
 - E.g., Use message object arbitration to subscribe to a particular message ID
 - E.g., Use mask object to accept that message ID from any source node #
 - Elevator example: message ID is button press; source node # tells which button
 - Single receiver mailbox then holds most recently received button press message
 - Message must be processed before next such message is received!

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DeviceNet Group Strategy

- ◆ **Group 1**
 - Prioritized by Message ID / Node number
 - High priority messages with fairness to nodes
- ◆ **Group 2**
 - Prioritized by Node number / Message ID
 - Gives nodes priority
- ◆ **Group 3**
 - Essentially same as Group 1, but allows Group 2 to have higher priority
- ◆ **Group 4**
 - Global housekeeping messages / must be unique in system (no node number)

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Other Approaches Are Possible

- ◆ And, you can invent your own too...
- ◆ Variations include:
 - Automatic assignment of node numbers (include hot-swap)
 - Automatic assignment of message numbers (include hot-swap)
 - Mixes of node-based vs. message-ID based headers
- ◆ Can you have two transmitters using the same exact header field?
 - No – that would produce a bus conflict
 - Unless you have middleware that ensures only one node can transmit at a time
 - For example use a low priority message as a token to emulate token-passing
- ◆ Higher level protocols define message types
 - For example, J1939 defines message ID meanings, mostly for trucks and buses

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CAN Workloads – Spreadsheets

- ◆ “SAE Standard Workload” (53 messages) V/C = Vehicle Controller [Tindell]

Signal Number	Signal Description	Size /bits	J /ms	T /ms	Periodic /Sporadic	D /ms	From	To
1	Traction Battery Voltage	8	0.6	100.0	P	100.0	Battery	V/C
2	Traction Battery Current	8	0.7	100.0	P	100.0	Battery	V/C
3	Traction Battery Temp, Average	8	1.0	1000.0	P	1000.0	Battery	V/C
4	Auxiliary Battery Voltage	8	0.8	100.0	P	100.0	Battery	V/C
5	Traction Battery Temp, Max.	8	1.1	1000.0	P	1000.0	Battery	V/C
6	Auxiliary Battery Current	8	0.9	100.0	P	100.0	Battery	V/C
7	Accelerator Position	8	0.1	5.0	P	5.0	Driver	V/C
8	Brake Pressure, Master Cylinder	8	0.1	5.0	P	5.0	Brakes	V/C
9	Brake Pressure, Line	8	0.2	5.0	P	5.0	Brakes	V/C
10	Transaxle Lubrication Pressure	8	0.2	100.0	P	100.0	Trans	V/C
11	Transaxle Clutch Line Pressure	8	0.1	5.0	P	5.0	Trans	V/C
12	Vehicle Speed	8	0.4	100.0	P	100.0	Brakes	V/C
13	Traction Battery Ground Fault	1	1.2	1000.0	P	1000.0	Battery	V/C
14	Hi&Lo Contactor Open/Close	4	0.1	50.0	S	5.0	Battery	V/C
15	Key Switch Run	1	0.2	50.0	S	20.0	Driver	V/C
16	Key Switch Start	1	0.3	50.0	S	20.0	Driver	V/C
17	Accelerator Switch	2	0.4	50.0	S	20.0	Driver	V/C
18	Brake Switch	1	0.3	20.0	S	20.0	Brakes	V/C
19	Emergency Brake	1	0.5	50.0	S	20.0	Driver	V/C
20	Shift Lever (PRNDL)	3	0.6	50.0	S	20.0	Driver	V/C
21	Motor/Trans Over Temperature	2	0.3	1000.0	P	1000.0	Trans	V/C
22	Speed Control	3	0.7	50.0	S	20.0	Driver	V/C
23	12V Power Ack Vehicle Control	1	0.2	50.0	S	20.0	Battery	V/C
24	12V Power Ack Inverter	1	0.3	50.0	S	20.0	Battery	V/C
25	12V Power Ack I/M Contr.	1	0.4	50.0	S	20.0	Battery	V/C
26	Brake Mode (Parallel/Split)	1	0.8	50.0	S	20.0	Driver	V/C

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CAN Tradeoffs

◆ Advantages

- High throughput under light loads
- Local and global prioritization possible
- Arbitration is part of the message - low overhead

◆ Disadvantages

- Requires bit dominance (can't be used with transformer coupling)
- Propagation delay limits bus length ($2 t_{pd}$ bit length)
- Unfair access - node with a high priority can "hog" the network
 - Can be reduced in severity with Message + Node # prioritization
 - Can, in principle, use a bus guardian to limit duty cycle of each node
- Poor latency for low priority nodes
 - Starvation is possible

◆ Optimized for:

- Moderately large number of message types
- Arbitration overhead is constant
- Global prioritization (*but* limited mechanisms for fairness)

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A COMPARISON OF VARIOUS AUTOMOBILE NETWORKING STANDARDS				
Name	Protocol specification	Interface	Type	Speed (kbits/s)
J1850	Yes	1 wire	Control	41.6
CAN	Yes	No	Control	Variable
CAN-A	CAN	2 wire	Control	33, 83
CAN-B	CAN	2 wire	Control	250
CAN-C	CAN	2 wire	Control	1000
SAE J2284	CAN	2 wire	Control	500
SAE J1939	CAN	2 wire	Control	125
SAE J2411	CAN	1 wire	Control	24
LIN	Yes	1 wire ISO 9141	Control	20
TTP/A	Yes	1 wire ISO 9141	Control	20
TTP/C	Yes	2 wire	X-by-wire	2000
Flexray	Yes	2 x 2 wire	X-by-wire	5000
IDB-C	CAN	2 wire	Multimedia	250
IEEE 1394	Yes	2 wire	Multimedia	300,000
Smartwire	Yes	2 wire	Multimedia	22,000
MOST	Yes	Fiber	Multimedia	25,000
				Comment
				Proprietary implementations.
				General protocol.
				Used in U.S.
				Used in U.S.
				Used in U.S.
				Used for power-train control.
				Recommended Practice for Serial Control and Communications Vehicle Network Class C by the SAE Truck & Bus Control and Communications Network Subcommittee of the Truck & Bus Electrical Committee.
				Unique to General Motors.
				Master/slave. Doesn't require crystal.
				Master/slave. Supports hot plug-and-play. Also supports higher speeds and fiber.
				Higher speeds possible.
				Developed by Philips.
				Developed by the IDB Forum.
				Being adapted to the automotive environment.
				Ring topology.
				Multiple master. Up to 64 devices.

[Electronic Design; Jan 8, 2001, pg. 66]

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Review

◆ Controller Area Network

- Binary-countdown arbitration
- Standard used in automotive & industrial control

◆ CAN Tradeoffs

- Good at global priority (but difficult to be “fair”)
- Efficient use of bandwidth
- Requires bit-dominance in physical layer
- Message filters are required to keep small nodes from being overloaded
 - Only works if small node can read data before next data in that mailbox arrives