

# Assignment 1 Part 1

## Analysis and Report

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## 1 Proof of the Heuristic Function

The *input1.txt* example is shown below :

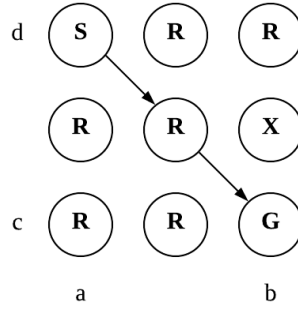


Figure 1: Minimum Heuristic from Start to Goal for *input1.txt*

Using the above example, we will prove that the heuristic is valid. The heuristic used in this case is the maximum between the difference from the current state to the goal state. The derivation of  $h(n)$  is stated below:

$$dx = |b - a| \quad (1)$$

$$dy = |d - c| \quad (2)$$

$$h(n) = \max(dx, dy) \quad (3)$$

Below is the python implementation for the heuristic:

```

1  def heuristic(self,x,y):
2      '''
3      method used to calculate the heuristic value given the
4      current x and y coordinates
5      @param x: current x coordinate
6      @param y: current y coordinate
7      @return: returns the heuristic value
8      '''
9      # Calculate the difference in both x and y directions
10     dx = abs(x - self.GOAL_COORD[0])
11     dy = abs(y - self.GOAL_COORD[1])
12
13     # Returns the max of either the x or y direction
14     return max(dx,dy)
15

```

When determining the heuristic we assume that there are no ridges in the map. Therefore the rules are relaxed compared to the actual rules of the environment. In order for the heuristic to be valid, it needs to be admissible and monotonic. Now let us prove that the above heuristic is Admissible and Monotonic.

## 1.1 Admissibility

In order for a heuristic to be admissible it needs to satisfy :

$$\forall n \quad h(n) \leq h^*(n) \quad (4)$$

Since we get the maximum value between dx and dy as the heuristic, it will always be the minimum amount of diagonal moves between the start and goal state. Therefore we can deduce that all cost will be greater than the heuristic and will never be less than it. So for the best case the heuristic will be equal to the actual cost, but for the worst case the heuristic will be underestimating since the cost of non-diagonal moves are greater than one.

Therefore the above heuristic is admissible as it satisfy Equation 4 and it does not overestimate the actual cost.

## 1.2 Monotonicity

In order for a heuristic to be monotonic, it should satisfy the following condition:

$$\forall n \quad h(n) \leq c(m,n) + h(m) \quad (5)$$

where m is a child of n

So rearranging the terms we can derive:

$$\forall n \quad h(n) - h(m) \leq c(m,n) \quad (6)$$

Since the rules of the environment states that:

$$c(m,n) \geq 1 \quad (7)$$

and our heuristic will resolve:

$$h(n) - h(m) == 1 \quad (8)$$

thus satisfying equation 5.

### 1.3 Is the resulting algorithm A or A\*?

Since the actual cost is always used for the  $g(n)$  value and not an estimate/ heuristic, we can state that the resulting algorithm is A\*.

## 2 Tie Breaker Implementation

As show below, to implement the tie breaker, we override the Node instance's less than operator:

```
1  def __lt__(self, other):
2      '''
3      This method is used to override the less than operator in
4      python to use the f cost for comparison
5      @param other: the other node to be compared
6      @return: boolean value stating whether its less than or not
7      '''
8      if (self.f < other.f):
9          return True
10     elif(self.f == other.f ):
11         if self.operator in self.best_operators:
12             return True
13         elif other.operator in other.best_operators:
14             return False
15         else:
16             return True
17     else:
18         return False
19
```

As per line 10, if the cost of the nodes are equal, we prioritize the node which was generated using a diagonal operator such as “LU, LD, RU, RD”. So the tie breaker implementation will always prioritize paths with diagonals.

## 3 Test Cases

### 3.1 Output for all the test cases

Below are the test cases and the resulting path from each algorithm:

#### 3.1.1 input1.txt

```
1      3
2      SRR
3      RRX
4      RRG
5
6      DLS : S-RD-D-R-G 5
7      A*  : S-RD-D-R-G 5
8
```

### 3.1.2 input2.txt

```
1      5
2      SRRXG
3      RXXRR
4      RRRXR
5      XXRRR
6      RRRRX
7
8      DLS : NO-PATH
9      A* : S-D-D-R-D-D-R-R-U-R-U-U-U-G 24
10
```

### 3.1.3 input3.txt

```
1      5
2      SRXXX
3      RRRXG
4      XRRRR
5      XRRRR
6      RXXRX
7
8      DLS : S-RD-RD-RD-RU-U-G 6
9      A* : S-RD-RD-RD-RU-U-G 6
10
```

### 3.1.4 input4.txt

```
1      7
2      RRRXRRR
3      RXRRRXR
4      RXXXXXR
5      RRXSXXX
6      XRXRXRX
7      XRXRXRX
8      XRRRXGR
9
10     DLS : NO-PATH
11     A* : S-D-D-D-L-L-U-U-U-L-U-U-U-R-R-D-R-R-U-R-R-D-D-D-D-D-L-G 54
12
```

### 3.1.5 input5.txt

```
1      5
2      SRRRG
3      RRRRR
4      XXXXX
5      RRRRR
6      RRRRR
7
8      DLS : S-RD-R-R-RU-G 6
9      A* : S-RD-RU-RD-RU-G 4
10
```

### 3.1.6 input6.txt

```
1      10
2      SRRRRRRRRR
3      RRRRRRRRRR
4      RRRRRRRRRR
5      RRRRRRRRRR
6      RRRRRRRRRR
7      RRRRRRRRRR
8      RRRRRRRRRR
9      RRRRRRRRRR
10     RRRRRRRRRR
11     GRRRRRRRRR
12
13     DLS : NO-PATH
14     A* : S-RD-LD-RD-RD-LD-LD-RD-LD-D-G 10
15
```

### 3.1.7 input7.txt

```
1      4
2      XRGR
3      SXRR
4      RRRX
5      RRRX
6
7
8     DLS : NO-PATH
9     A* : NO-PATH
10
```

### 3.1.8 input8.txt

```
1      6
2      SRRRRR
3      RRRXXR
4      RXRRRR
5      RRRXXR
6      XRRRRR
7      GRRRXR
8
9     DLS : NO-PATH
10     A* : S-D-D-D-R-D-D-L-G 14
11
```

### 3.1.9 input9.txt

```
1      6
2      GRRRRR
3      RRRRRR
4      RRRRRR
5      RRRRRR
6      RRRRRR
7      RRRRRS
8
9     DLS : S-LU-LU-LU-LU-LU-G 5
10     A* : S-LU-LU-LU-LU-LU-G 5
11
```

## 3.2 Node expansion outputs

Below are the five node expansions for *input1.txt*, *input2.txt* and *input3.txt* when using **A\*** algorithm:

### 3.2.1 input1.txt

```
1  3
2  SRR
3  RRX
4  RRG
5
6  A* : S-RD-D-R-G 5
7
8  N0:S      1 0 0 0
9  Children:      {N1:S-R,N2:S-D,N3:S-RD}
10 OPEN:      {(N3:S-RD 1 2 3)(N2:S-D 2 2 4)(N1:S-R 2 3 5)}
11 CLOSED: {(N0:S 1 0 0 0)}
12
13 N3:S-RD 2 1 2 3
14 Children:      {N1:S-R,N2:S-D,N4:S-RD-LD,N5:S-RD-D}
15 OPEN:      {(N4:S-RD-LD 2 2 4)(N2:S-D 2 2 4)(N5:S-RD-D 3 1 4)(N1:S-R 2 3 5)}
16 CLOSED: {(N0:S 1 0 0 0)(N3:S-RD 2 1 2 3)}
17
18 N4:S-RD-LD      3 2 2 4
19 Children:      {N2:S-D,N5:S-RD-D}
20 OPEN:      {(N5:S-RD-D 3 1 4)(N2:S-D 2 2 4)(N1:S-R 2 3 5)}
21 CLOSED: {(N0:S 1 0 0 0)(N3:S-RD 2 1 2 3)(N4:S-RD-LD 3 2 2 4)}
22
23 N5:S-RD-D      4 3 1 4
24 Children:      {N2:S-D,N4:S-RD-LD,N6:S-RD-D-R-G}
25 OPEN:      {(N2:S-D 2 2 4)(N1:S-R 2 3 5)(N6:S-RD-D-R-G 5 1 6)}
26 CLOSED: {(N0:S 1 0 0 0)(N3:S-RD 2 1 2 3)(N4:S-RD-LD 3 2 2 4)(N5:S-RD-D 4 3 1 4)}
27
28 N2:S-D 5 2 2 4
29 Children:      {N1:S-R,N3:S-RD,N4:S-RD-LD,N5:S-RD-D}
30 OPEN:      {(N1:S-R 2 3 5)(N6:S-RD-D-R-G 5 1 6)}
31 CLOSED: {(N0:S 1 0 0 0)(N3:S-RD 2 1 2 3)(N4:S-RD-LD 3 2 2 4)(N5:S-RD-D 4 3 1 4)(N2:
      S-D 5 2 2 4)}
32
33
```

### 3.2.2 input2.txt

```
1 5
2 SRRXG
3 RXXRX
4 RRRXR
5 XXRRR
6 RRRRX
7
8 A* : S-D-D-R-D-D-R-R-U-R-U-U-U-G 24
9
10 N0:S      1 0 0 0
11 Children: {N1:S-R,N2:S-D}
12 OPEN:     {(N1:S-R 2 3 5)(N2:S-D 2 4 6)}
13 CLOSED:   {(N0:S 1 0 0 0)}
14
15 N1:S-R    2 2 3 5
16 Children: {N3:S-R-R}
17 OPEN:     {(N2:S-D 2 4 6)(N3:S-R-R 4 2 6)}
18 CLOSED:   {(N0:S 1 0 0 0)(N1:S-R 2 2 3 5)}
19
20 N2:S-D    3 2 4 6
21 Children: {N4:S-D-D}
22 OPEN:     {(N3:S-R-R 4 2 6)(N4:S-D-D 4 4 8)}
23 CLOSED:   {(N0:S 1 0 0 0)(N1:S-R 2 2 3 5)(N2:S-D 3 2 4 6)}
24
25 N3:S-R-R      4 4 2 6
26 Children:     {N5:S-R-R-D}
27 OPEN:        {(N4:S-D-D 4 4 8)(N5:S-R-R-D 6 2 8)}
28 CLOSED:      {(N0:S 1 0 0 0)(N1:S-R 2 2 3 5)(N2:S-D 3 2 4 6)(N3:S-R-R 4 4 2 6)}
29
30 N4:S-D-D      5 4 4 8
31 Children:     {N6:S-D-D-R}
32 OPEN:        {(N5:S-R-R-D 6 2 8)(N6:S-D-D-R 6 3 9)}
33 CLOSED:      {(N0:S 1 0 0 0)(N1:S-R 2 2 3 5)(N2:S-D 3 2 4 6)(N3:S-R-R 4 4 2 6)(N4:S-D-D
34 5 4 4 8)}
35
36
```

### 3.2.3 input3.txt

```
1 5
2 SRXXX
3 RRRXG
4 XRRRR
5 XRRRR
6 RXXRX
7
8 A* : S-RD-RD-RD-RU-U-G 6
9
10 NO:S      1 0 0 0
11 Children: {N1:S-R,N2:S-D,N3:S-RD}
12 OPEN:      {(N3:S-RD 1 3 4)(N1:S-R 2 3 5)(N2:S-D 2 4 6)}
13 CLOSED:    {(NO:S 1 0 0 0)}
14
15 N3:S-RD 2 1 3 4
16 Children:   {N1:S-R,N2:S-D,N4:S-RD-R,N5:S-RD-D,N6:S-RD-RD}
17 OPEN:       {(N6:S-RD-RD 2 2 4)(N1:S-R 2 3 5)(N4:S-RD-R 3 2 5)(N2:S-D 2 4 6)(N5:S-RD-D
18 3 3 6)}
19 CLOSED:     {(NO:S 1 0 0 0)(N3:S-RD 2 1 3 4)}
20
21 N6:S-RD-RD      3 2 2 4
22 Children:        {N4:S-RD-R,N5:S-RD-D,N7:S-RD-RD-R,N8:S-RD-RD-LD,N9:S-RD-RD-D,N10:S-
23 RD-RD-RD}
24 OPEN:            {(N10:S-RD-RD-RD 3 1 4)(N4:S-RD-R 3 2 5)(N1:S-R 2 3 5)(N7:S-RD-RD-R 4 1 5)(
25 N8:S-RD-RD-LD 3 3 6)(N5:S-RD-D 3 3 6)(N2:S-D 2 4 6)(N9:S-RD-RD-D 4 2 6)}
26 CLOSED:          {(NO:S 1 0 0 0)(N3:S-RD 2 1 3 4)(N6:S-RD-RD 3 2 2 4)}
27
28 N10:S-RD-RD-RD 4 3 1 4
29 Children:       {N7:S-RD-RD-R,N11:S-RD-RD-RD-RU,N9:S-RD-RD-D,N12:S-RD-RD-RD-R,N13:S-
30 RD-RD-RD-D}
31 OPEN:           {(N11:S-RD-RD-RD-RU 4 0 4)(N7:S-RD-RD-R 4 1 5)(N1:S-R 2 3 5)(N4:S-RD-R 3 2
32 5)(N8:S-RD-RD-LD 3 3 6)(N9:S-RD-RD-D 4 2 6)(N2:S-D 2 4 6)(N5:S-RD-D 3 3 6)(N12:S-
33 RD-RD-RD-R 5 1 6)(N13:S-RD-RD-RD-D 5 2 7)}
34 CLOSED:         {(NO:S 1 0 0 0)(N3:S-RD 2 1 3 4)(N6:S-RD-RD 3 2 2 4)(N10:S-RD-RD-RD 4 3 1
35 4)(N11:S-RD-RD-RD-RU 5 4 0 4)}
36
37 N11:S-RD-RD-RD-RU      5 4 0 4
38 Children:              {N14:S-RD-RD-RD-RU-U-G,N7:S-RD-RD-R,N12:S-RD-RD-RD-R}
39 OPEN:                  {(N4:S-RD-R 3 2 5)(N1:S-R 2 3 5)(N7:S-RD-RD-R 4 1 5)(N8:S-RD-RD-LD 3 3 6)(
40 N12:S-RD-RD-RD-R 5 1 6)(N5:S-RD-D 3 3 6)(N2:S-D 2 4 6)(N9:S-RD-RD-D 4 2 6)(N13:S-
41 RD-RD-RD-D 5 2 7)(N14:S-RD-RD-RD-RU-U-G 6 1 7)}
42 CLOSED:                {(NO:S 1 0 0 0)(N3:S-RD 2 1 3 4)(N6:S-RD-RD 3 2 2 4)(N10:S-RD-RD-RD 4 3 1
43 4)(N11:S-RD-RD-RD-RU 5 4 0 4)}
```



## 4 Analysis for test maps

In order to perform an analysis, multiple test cases were generated and tested on the two algorithms. Below is a table with the time taken for each input by the two algorithms. Please do note that each time is an **average of five runs** and the **bound of the DLS is set to 5**.

Input File	Time Taken											
	DLS						A*					
	1	2	3	4	5	Average	1	2	3	4	5	Average
input1.txt	0.00024	0.00017	0.00023	0.00017	0.00017	<b>0.00020</b>	0.00028	0.00028	0.00028	0.00033	0.00028	<b>0.00029</b>
input2.txt	0.00037	0.00036	0.00037	0.00038	0.00037	<b>0.00037</b>	0.00132	0.00090	0.00092	0.00083	0.00083	<b>0.00096</b>
input3.txt	0.00035	0.00034	0.00035	0.00035	0.00035	<b>0.00035</b>	0.00081	0.00068	0.00058	0.00057	0.00057	<b>0.00064</b>
input4.txt	0.00023	0.00022	0.00023	0.00023	0.00024	<b>0.00023</b>	0.00152	0.00152	0.00150	0.00152	0.00150	<b>0.00151</b>
input5.txt	0.00035	0.00027	0.00027	0.00027	0.00027	<b>0.00029</b>	0.00029	0.00029	0.00028	0.00027	0.00029	<b>0.00028</b>
input6.txt	0.00201	0.00197	0.00202	0.00198	0.00197	<b>0.00199</b>	0.00336	0.00373	0.00369	0.00446	0.00431	<b>0.00385</b>
input7.txt	0.00026	0.00021	0.00023	0.00022	0.00022	<b>0.00023</b>	0.00022	0.00026	0.00023	0.00023	0.00023	<b>0.00023</b>
input8.txt	0.00061	0.00059	0.00049	0.00049	0.00050	<b>0.00054</b>	0.00081	0.00082	0.00087	0.00083	0.00083	<b>0.00083</b>
input9.txt	0.00349	0.00347	0.00349	0.00349	0.00347	<b>0.00348</b>	0.00077	0.00078	0.00094	0.00078	0.00077	<b>0.00080</b>

Table 1: Time Taken for each Algorithm on each input file

Now let us analyze the results for each test map. For all the analysis, we considered a bound of 5 for DLS. The reason for this is to test the boundary cases where the only path for DLS is the optimal path. So for some of the cases, there will be no solution as DLS is not complete.

### 4.1 input1.txt

1	3
2	SRR
3	RRX
4	RRG
5	
6	DLS : S-RD-D-R-G 5
7	A* : S-RD-D-R-G 5
8	

In this test case, both the DLS and A\* perform equally. The reason for this is the graph search updates the cost of the nodes in the open list to the nodes with the minimum cost to a certain location. So if the DLS get an expensive path to a node, it is not updated. Therefore, in DLS, if the correct actions are picked first, both of them can perform equally and can lead to an optimal path. But A\* will guarantee an optimal path but will take more time. So from Table 1 we can see that for *input1.txt* on average A\* was slower.

## 4.2 input2.txt

```
1      5
2      SRRXG
3      RXXRX
4      RRRXR
5      XRXRR
6      RRRRX
7
8      DLS : NO-PATH
9      A* : S-D-D-R-D-D-R-R-U-R-U-U-U-G 24
10
```

In this test case, DLS returns NO-PATH since the depth of the optimal solution is more than 6. So there is no path for DLS to reach the goal in less than 5 steps. This shows that DLS is not complete.

## 4.3 input3.txt

```
1      5
2      SRXXX
3      RRRXG
4      XRRRR
5      XRRRR
6      RXXRX
7
8      DLS : S-RD-RD-RD-RU-U-G 6
9      A* : S-RD-RD-RD-RU-U-G 6
10
```

Similar to *input1.txt*, DLS also returns the optimal path and A\* takes more time to run. This map was used to test how the agent react to when it finds a shorter path at the last step. So if the agent stopped when the goal node was generated, we would get an error as there is a better path than the initial path.

## 4.4 input4.txt

```
1      7
2      RRRXRRR
3      RXXRRXR
4      RXXXXXR
5      RRXSXXR
6      XRXRXRX
7      XRXRXRX
8      XRRRXGR
9
10     DLS : NO-PATH
11     A* : S-D-D-D-L-L-U-U-U-L-U-U-U-R-R-D-R-R-U-R-R-D-D-D-D-D-D-L-G 54
12
```

Similar to *input2.txt*, DLS returns NO-PATH and A\* takes more time to run. This map is used to test whether the agent follows the rules of the environment and not choose paths out of bounds.

## 4.5 input5.txt

```
1      5
2      SRRRG
3      RRRRR
4      XXXXX
5      RRRRR
6      RRRRR
7
8      DLS : S-RD-R-R-RU-G 6
9      A* : S-RD-RU-RD-RU-G 4
10
```

The ridges blocks the bottom half of the map. So this map is useful in testing whether the agent takes the shorter path or the longer path to get to the goal. This also test whether the rules of the environment are stable. The run time of both algorithms are similar since the optimal path is similar to a DLS, where picking the diagonals is the best option. But as the results indicate, DLS takes a longer, non optimal path compared to A\*. This is due to the order of actions taken by the agent.

## 4.6 input6.txt

```
1      10
2      SRRRRRRRRR
3      RRRRRRRRRR
4      RRRRRRRRRR
5      RRRRRRRRRR
6      RRRRRRRRRR
7      RRRRRRRRRR
8      RRRRRRRRRR
9      RRRRRRRRRR
10     RRRRRRRRRR
11     GRRRRRRRRR
12
13     DLS : NO-PATH
14     A* : S-RD-LD-RD-RD-LD-LD-RD-LD-D-G 10
15
```

This test can be used to see if the agent chooses the most optimal path, as taking diagonal actions only doesn't lead to the goal node. The run time of A\* is more because DLS is bounded to 5.

## 4.7 input7.txt

```
1      4
2      XRGR
3      SXRR
4      RRXR
5      RRRX
6
7
8      DLS : NO-PATH
9      A* : NO-PATH
10
```

This test map was used to test how the agent handles, when there is no path to the goal. So both agents returned NO-PATH, with equal run time as they had to exhaust all the paths.

## 4.8 input8.txt

```
1      6
2      SRRRRR
3      RRRXXR
4      RXRRRR
5      RRXXRR
6      XRRRRR
7      GRRRXR
8
9      DLS : NO-PATH
10     A* : S-D-D-D-R-D-D-L-G 14
11
```

This test map was used to test how the agent handles, when there is a shorter path with more horizontal moves than diagonal moves. So the DLS did not find a path while A\* found the path while taking more run time.

## 4.9 input9.txt

```
1      6
2      GRRRRR
3      RRRRRR
4      RRRRRR
5      RRRRRR
6      RRRRRR
7      RRRRRS
8
9      DLS : S-LU-LU-LU-LU-LU-G 5
10     A* : S-LU-LU-LU-LU-LU-G 5
11
```

This test map was used to test how the agent handles, when the Start and Goal nodes are swapped with no obstacles. Another constraint is the path for DLS is the same length as the bound. So the DLS should return the most optimal path. So looking at the results, we see that the DLS takes more time compared to A\*. This is because the A\* algorithm only used 5 moves and gets to the goal node, where as DLS has to go through all the possible paths upto the bound until it finds the single optimal path.

## 5 Conclusion

Therefore in conclusion we can see that the A\* algorithm finds the optimal path compared to DLS, while DLS is faster at finding paths. But if the path is outside the bound value, DLS doesn't find the path. Therefore DLS is not complete while on the other hand, A\*, given that there exist a path will always find it.