

# CS3205: Introduction to Computer Networks

## Simulation of the TCP Congestion Control Algorithm

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### Introduction:

TCP is a protocol that is used to transmit information from one computer on the internet to another. For every segment of data that is sent from sender to receiver, receiver sends an “acknowledgement” segment back to sender indicating that it got that message. It Functions using congestion window which limits the flow of data using parameters  $K_i, K_m, K_n, K_f, P_s$ .

**Factors influence the CW change over the duration of the session.**

**$K_i$  : initial congestion window(CW)**

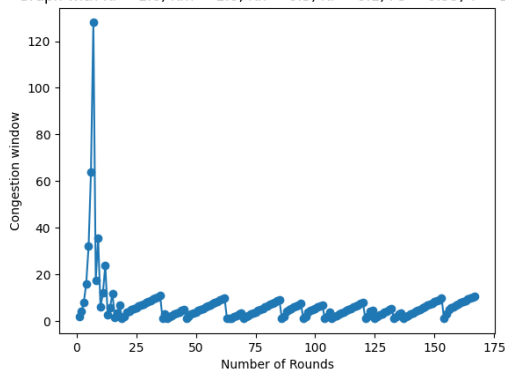
**Initial exponential phase will be quicker**

The initial CW is given by  **$CW_{new} = K_i * MSS$**

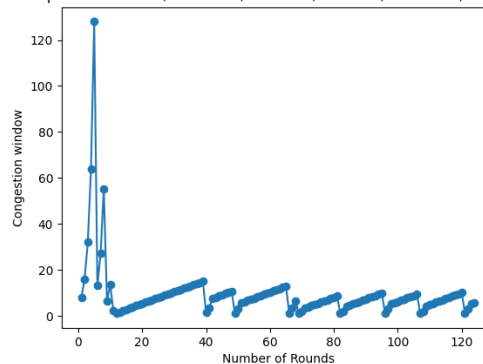
The value of CW starts from a proportionately larger value for larger values.

The initial starts at higher in case of larger  $K_i$  from below diagrams.

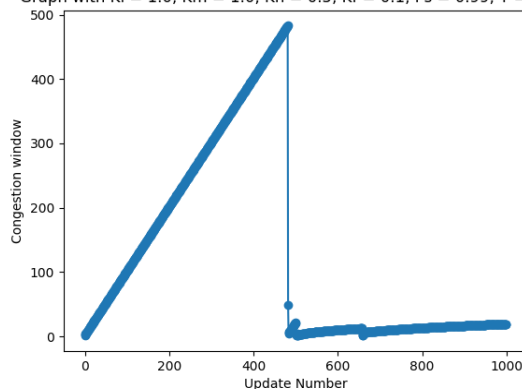
Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



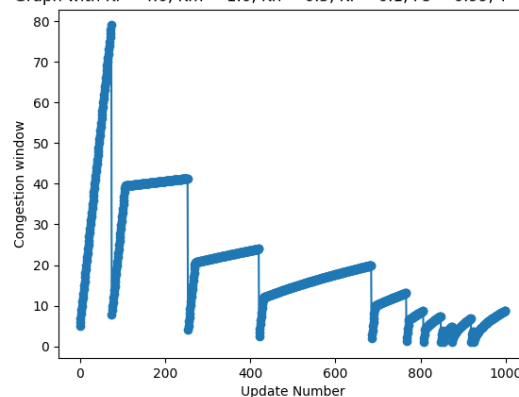
Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



## Km : Multiplier of the CW during exponential growth

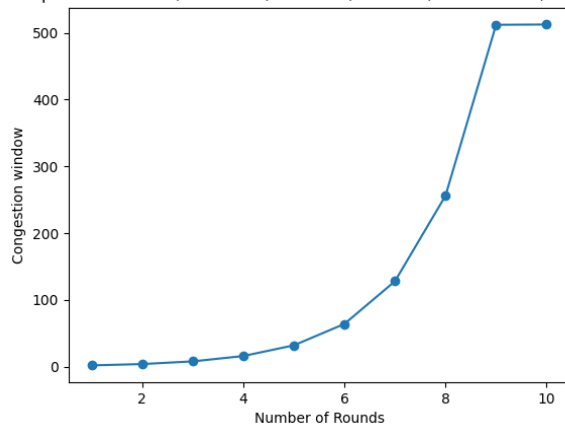
**Steeper will be the exponential phases**

When a segment's ACK is successfully received, CW is updated as follows

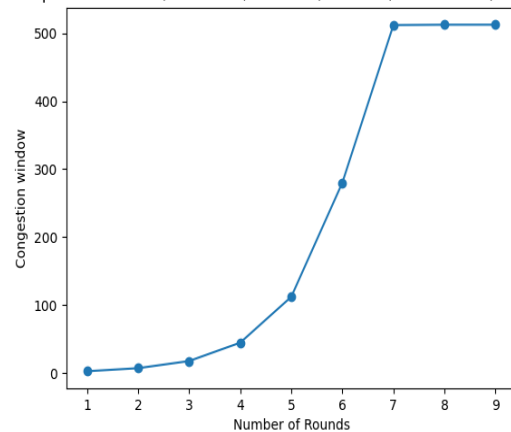
$$CW_{new} = \min(CW_{old} + Km * MSS, RWS)$$

Km affects the increment of congestion window size in the exponential growth phase. This can be observed that  $km = 1.5$  (5 rounds) reaches 100 faster when compared  $km = 1.0$  (6 rounds).

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$

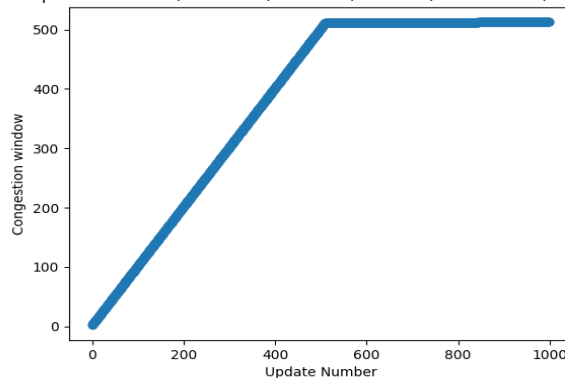


Graph with  $K_i = 1.0$ ,  $K_m = 1.5$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$

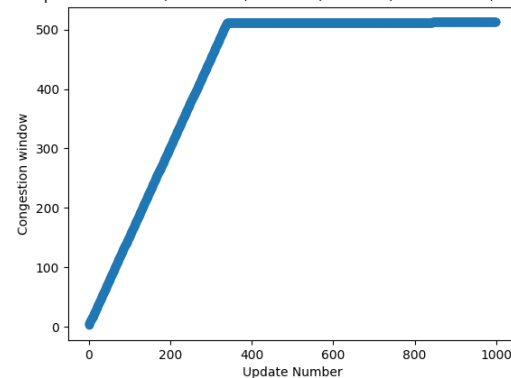


## Plots for different values of Km w.r.to rounds

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$



Graph with  $K_i = 1.0$ ,  $K_m = 1.5$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$



## Plots for different values of Km w.r.to update Number

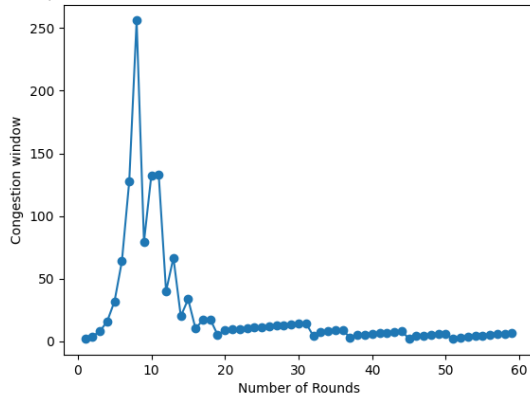
**Kn : Multiplier of the CW during linear growth**

**Slope of linear phase will be high**

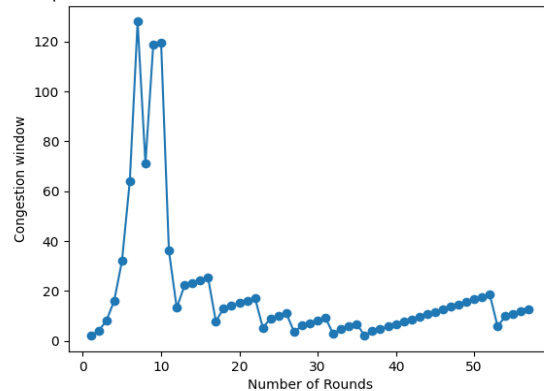
When a segment's ACK is successfully received, CW is updated as follows  
 **$CW_{new} = \min(CW_{old} + Kn * MSS * MSS_{CW_{old}}, RWS)$**

*Kn* affects the increment of congestion window in linear growth phase. This can be seen from the diagrams that slopes are higher at regions around 30 rounds in case of higher *Kn*.

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$

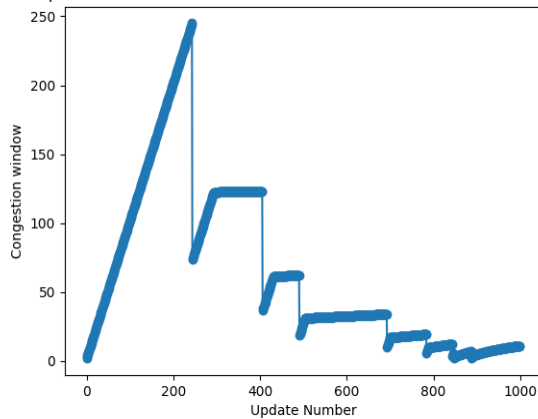


Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 1.0$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$

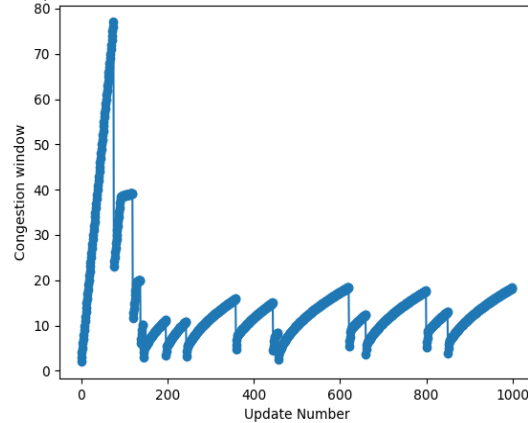


**Plots for different values of *Kn* w.r.to rounds**

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$



Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 1.0$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$



**Plots for different values of *Kn* w.r.to update Number**

**Kf : Multiplier when a timeout occurs**

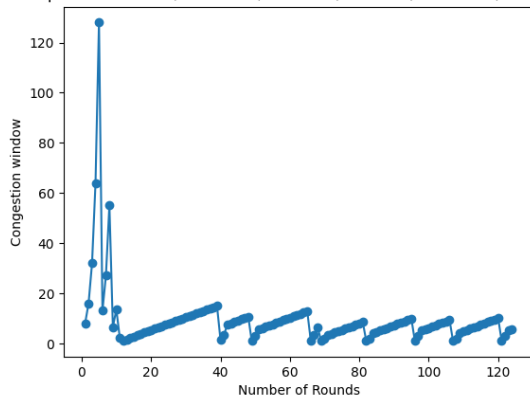
**Drop at time out will be higher**

Updating the CW when timeout occurs

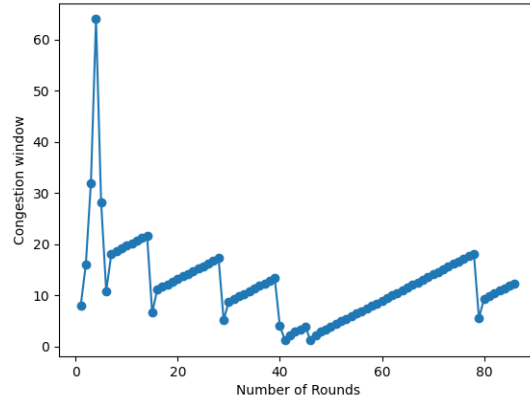
$$CW_{new} = \max(1, K_f * CW_{old})$$

The sudden drop in congestion window size implies that there is an occurrence of timeout, which can be seen that drop is higher in case of  $k_f = 0.3$  when compared  $k_f = 0.1$ .

Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$

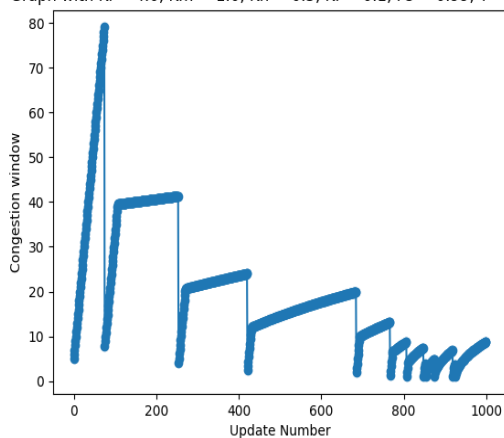


Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$

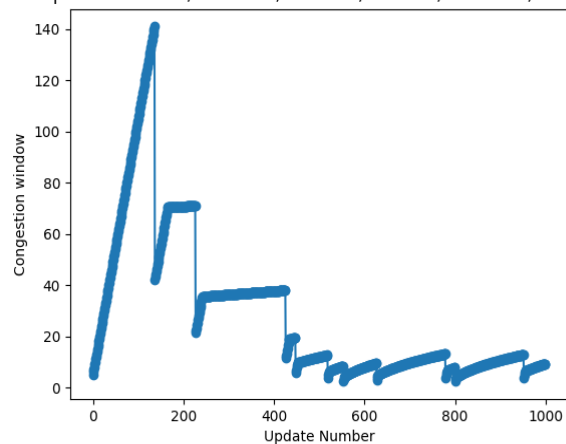


**Plots for different values of Kf w.r.to rounds**

Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



Graph with  $K_i = 4.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.3$ ,  $P_s = 0.99$ ,  $T = 1000$



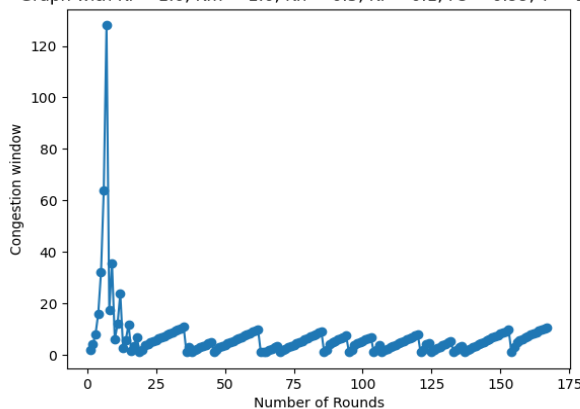
**Plots for different values of kf w.r.to update number**

## Ps : Probability of ACK before timeout

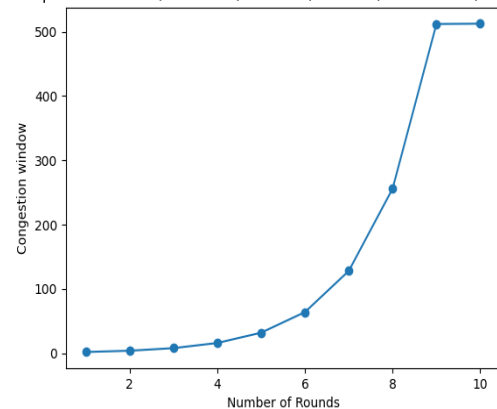
**Less chance of getting timeout,so less rounds to send T segments**

Ps affects the probability of timeout. When  $P_s = 0.99$ , we can observe the drastic decrement of  $cw$  value at many instants compared to that of  $P_s = 0.9999$ . When  $P_s = 0.9999$ , the chance for timeout is very less which can be easily observed from below diagrams.

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$

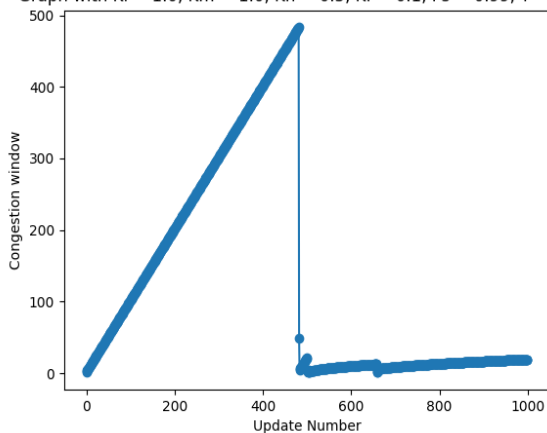


Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$

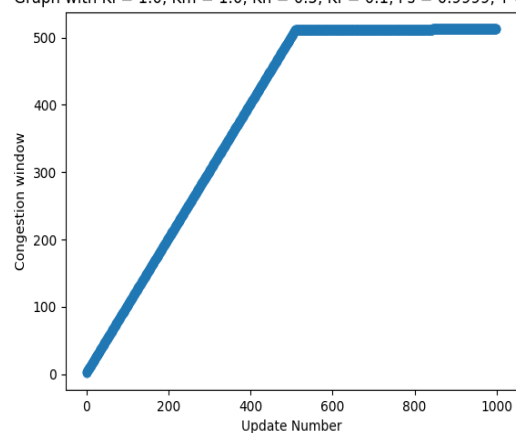


## Plots for different values of Ps w.r.to rounds

Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.99$ ,  $T = 1000$



Graph with  $K_i = 1.0$ ,  $K_m = 1.0$ ,  $K_n = 0.5$ ,  $K_f = 0.1$ ,  $P_s = 0.9999$ ,  $T = 1000$



## Plot for different values of Ps w.r.to update number