

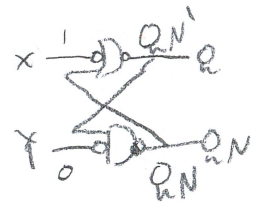
Bryant Har

bjh254

Q1 a) Yes

$X=1$ $Y=0$
to hold state

X	Y	Q_n	Q_{n+1}	$Q_n N_f$
1	0	Q	$Q N'$	$Q N$
$= Q$				



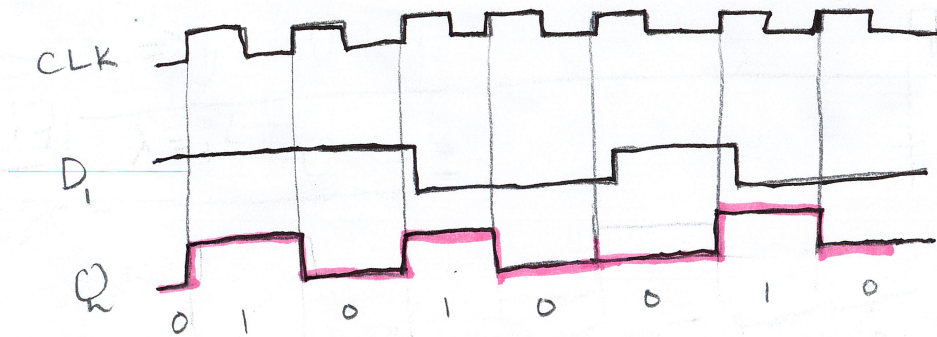
b) Yes

$X=0$ $Y=$ Don't Care

X	Y	Q_n	Q_{n+1}	$Q_n N_f$
0	1	Q	0	1
0	0	Q	0	1

c) No. $X=0$ for Q to possibly be 1, and $Y=0$ for Q_N to possibly be 0. But even with these inputs, $Q=1$ for the NAND output Q_N to be 0, $Q=0$ will mean Q_N will be 1. In other words, no input can always set $Q=1$, $Q_N=0$ independent of the previous state.

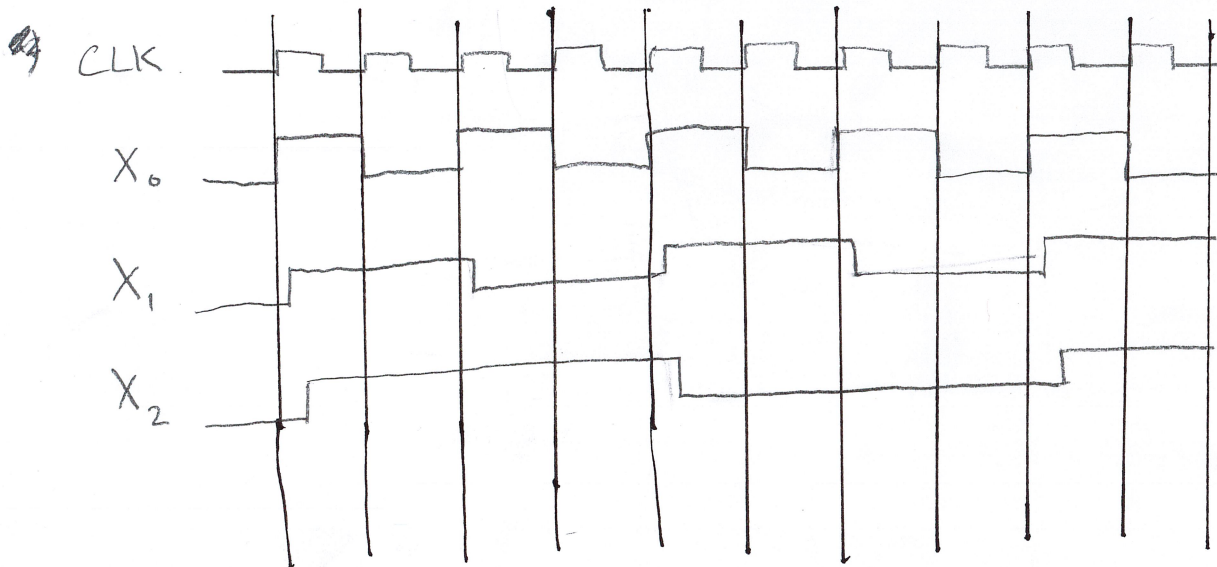
Q2



Q3)

a)

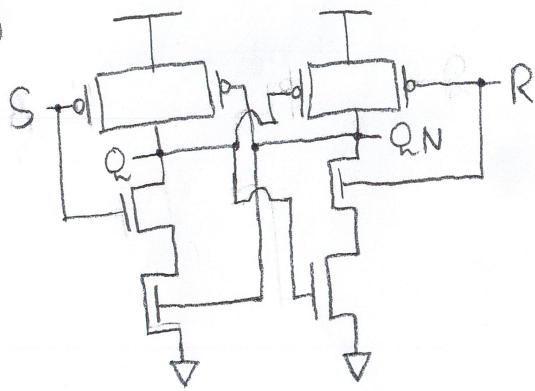
Cycle	0	1	2	3	4	5	6	7	8	9	10
X ₀	0	1	0	1	0	1	0	1	0	1	0
X ₁	0	1	1	0	0	1	1	0	0	1	1
X ₂	0	1	1	1	1	0	0	0	0	1	1



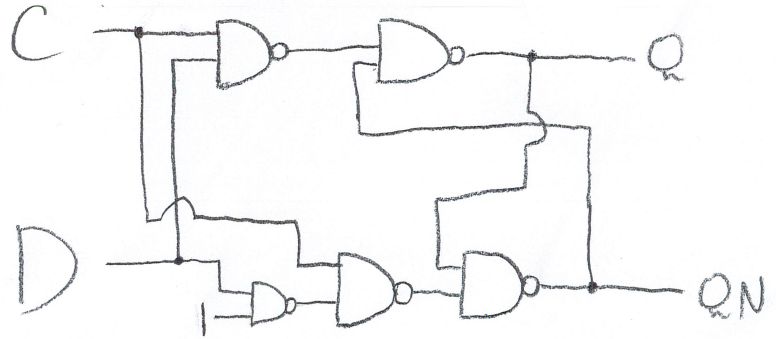
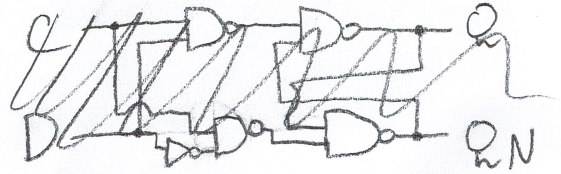
b) We made a down counter (counts in reverse)

c) Every X₀, X₁, X₂ will now read from Q_n' instead of Q_n. So, it will go, starting with all Qs being 0 or all Q_ns being 1: 111, 000, 001, 010, 011, 100, ... We have now made an up counter.

Q4) a)



b) C is Enable, so $C=0 \Rightarrow S=R=1$
 If $C=1$ $\begin{cases} D=0: S=1, R=0 \\ D=1: S=0, R=1 \end{cases}$



Q5)

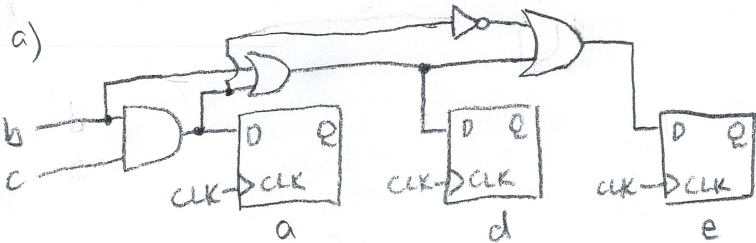
First = ? \Rightarrow if $i=1xxx$, $y=11$. Otherwise, if $i=x1xx$, $y=10$.
 Otherwise, if $i=xx1x$, $y=01$. otherwise if $i=xxx1$, $y=00$.

i	y
$1xxx$	00
$x1xx$	01
$xx1x$	10
$xxx1$	11

This is clearly

a 4-to-2 encoder

Q6)



b)

