# National Institute of Technology Karnataka Surathkal



Compiler Design Lab Report - 1

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# Introduction

A compiler is a computer program that translates instruction text (source language) into a different language (target language) of instruction text. Typically, the source language is a high-level language, while the target code is machine code. Compilers normally converts into machine code for the machine it compiles in, but special compilers known as cross-compiler.

User writes the program in C and gives it to the compiler. Then the compiler will compile the code i.e. it converts the high-level language to machine level language. That converted code is executed to get the desired output.

# Pre-processor:

A pre-processor is considered as a part of the compiler, it is a tool that produces input for the compilers.

# Interpreter:

An interpreter translates high level language to machine level language. The difference lies in the way they read the source code. A compiler will read the whole source code at once, creates tokens, check for syntax, etc... while the interpreter will read the statement from the input converts it to an intermediate code, executes it and then takes the next statement for execution. If error occurs interpreter will stop execution and reports the error.

#### Linker:

Linker is a computer program that links and merges various object files together in order to make the file executable. The major task of a linker is to search and locate referenced module/routines in a program and to determine the memory location where these codes will be loaded.

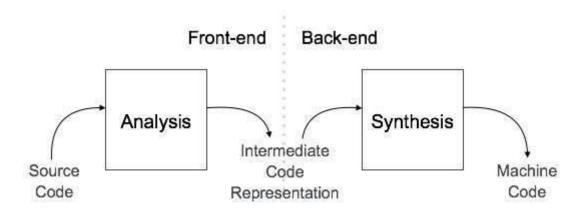
#### Loader:

Loader is a part of operating system and is responsible for loading executable files into memory and execute them. It calculates the size of a program (instructions and data) and creates memory space for it.

# **Cross – Compiler:**

A compiler that runs on one platform and the generated code is capable of executed in another platform then it is called cross-compiler.

# Phases Involved in a Compiler



There may be many phases involved in a compiler, however they can be assigned to one of the two stages:

- The front end: In this phase, the compiler verifies syntax and semantics according to the specific source language. In case of errors, warnings and highlighted erroneous code are shown. The input here is the source code and the output is an intermediate representation (IR) for further processing. More specifically, the intermediate outputs are token stream of the source code and sentences of the program.
- 2. **The back end**: This phase performs further analysis, transformations and optimizations specific to the target CPU architecture. As a result, the target assembly code and register allocation is performed. Further optimization to ensure maximum utilization of the target hardware is done. The input here is the optimized IR and the output is the target code.

**Pass:** A pass refers to the traversal of a compiler through the entire program.

**Phase:** A phase of a compiler is a distinguishable stage, which takes input from the previous stage, processes and yields output that can be used as input for the next stage.

The compilation process is a sequence of various phases. Each phase takes input from its previous stage, has its own representation of source program, and feeds its output to the next phase of the compiler.

In frontend there are four stages they are lexical analyser, syntax analyser, semantic analyser, intermediate code generator, machine independent code optimiser, machine dependent code optimiser.

# Lexical Analysis:

A lexical analyser or scanner performs tokenization on a sequence of characters such as a computer program. Tokenization is a process of converting a sequence of characters into a sequence of tokens. A token is a string assigned with some meaning in a language. Lexical Analysis is the first stage in the front end of a compiler. A scanner also does other tasks, such as removing whitespaces and comments from the source program.

Lexical Analysis can be further divided into two stages - scanning and evaluating. In scanning, the source code is converted into syntactical units called lexemes and categorized into token classes. A lexeme is a sequence of characters that matches a pattern for a token. Types of tokens include identifiers, keywords, separators, operators, comments, etc. In evaluating, the lexemes are converted into processed values. This process is considered simpler when compared to parsing, etc. This can be performed using a lexer generator such as lex. A free implementation of lex found in Unix OS is called flex.

# Syntax Analysis

The next phase is called the syntax analysis or **parsing**. It takes the token produced by lexical analysis as input. For syntax analysis, context-free grammars and the associated parsing techniques are to be used. In this phase, token arrangements are checked against the source code

grammar, i.e. the parser checks if the expression made by the tokens is syntactically correct.

The parse trees are used to show the structure of the sentence. There are many techniques for parsing the algorithms, the main techniques are top-down and bottom-up parsing.

# Semantic Analysis

Semantic analysis checks whether the parse tree constructed follows the rules of language. For example, assignment of values is between compatible data types, and adding string to an integer. Also, the semantic analyser keeps track of identifiers, their types and expressions; whether identifiers are declared before use or not etc. The semantic analyser produces an annotated syntax tree as an output.

# Intermediate code generation:

After semantic analysis the compiler generates an intermediate code of the source code for the target machine. It represents a program for some abstract machine. The code generated in here should be such a way that it would be easier to be translated into target code.

# Code optimization:

In this phase, the code is optimized. It removes the unwanted code lines, and arranges the sequence of statements in order to speed up the execution of the program.

# Machine code generation:

The code generator takes the optimized representation of the intermediate code and then it converts it to target machine language.

#### Tokens:

Lexemes are said to be a sequence characters in tokens. There are some rules where every lexeme is identified. These rules are defined by means of a pattern. A pattern explains what can be a token and these are represented by regular expressions.

Keywords, constant, identifiers, numbers, operators and punctuations symbols are possible tokens to be identified.

# **Issues in Lexical Analyser:**

It has two errors. They are lookahead, ambiguities. Lookahead is required to decide when one token will end and the next token will begin. The lexical analysis programs written with lex accept ambiguous specifications and choose the longest match possible at each input point.

# **Lexical errors:**

A character sequences that cannot be scanned into any valid token is a lexical error.

Lexical errors are uncommon, but they still must be handled by scanner.

# Code for Lexical Analyzer:

```
◆
       p3.I
     %{
     int comm=0;
     #include<stdio.h>
  4 #include<string.h>
  5 #include<stdlib.h>
     int var_cnt=0;
     void insert(char *yytext,char t,int 1);
     struct Storage
      char *name, token[20];
      int val;
     int li;
     struct Storage *next;
     }*st,*head
     int braces=0;
 16 int opbrc=0; //open braces ( )
17 int quo=0;
     int line=1;
     %}
     D [0-9]
     L [a-zA-Z]
 23 0 [-+*/=<>]
 24 op "-="|"+="|"*="|"/="|"<<"|">>>"|"<="|">="|"=="
 25 WS [ \t\n|"\\"]
 26 E [Ee][-+]?{D}+
     PUN [,.;]
     U [_]
     Qu ["']
     per [%&]
 31 Keyword "break"|"case"|"continue"|"default"|"enum"|"register"|"return"|"sizeof"|"static"|"typedef"|"void"|"volatile"
Datatype "char"|"double"|"float"|"int"|"union"|"struct"
Loop "do"|"for"|"while"|"goto"
     Condition "if" | "else" | "switch"
     EDT "auto"|"short"|"long"|"const"|"extern"|"signed"|"unsigned"
     FS "d"|"f"|"c"|"s"|"ld"|"lf"
     SB "["|"]"
     ws2 "\\"{Qu}
     headerFunc "printf" | "scanf"
     header "#include<"
 43 close ".h>"
     def "#define" | "# define"
 45 String {Qu}({L}|{D}|{per}|{WS}|{0}|{ws2})*{Qu}
```

```
◆▶
       p3.1
      String {Qu}({L}|{D}|{per}|{WS}|{0}|{ws2})*{Qu}
      IFR ({L}|{U})({L}|{D})*
      array {IFR}{SB}{D}*{SB}
      wl {D}({L}|{D})*
      "\n" line++;
      "/*" { if(comm==0) comm++; }
      "*/" if(comm==1){ comm--; } else { printf("Error: Encountered a */ before /*\n"); }
                                                     { insert(yytext, 'P', line);}
      {header}{L}*{close}
      {def}
                                                     { insert(yytext, 'P', line);}
                                                     {if(comm==0) { insert(yytext, 'k', line);}}
      {Keyword}
                                                     {if(comm==0){ insert(yytext, 'd', line);}
      {Datatype}
                                                     {if(comm==0){ insert(yytext,'l',line);}
      {Loop}
      {Condition}
                                                     {if(comm==0){ insert(yytext,'c',line);} }
      {EDT}+" "+{EDT}*+" "+{Datatype}
                                                     {if(comm==0){ insert(yytext, 'e', line);} }
      {headerFunc}
                                                     {if(comm==0){ insert(yytext, 'h', line);}}
      ({D}+)+"."+{D}*
                                                     {if(comm==0){ insert(yytext, 'C', line);} }
      {D}*
                                                     {if(comm==0){ insert(yytext,'C',line);}}
      {String}
                                                     {if(comm==0){ insert(yytext,'s',line);} }
      {Qu}+{IFR} { {printf(" %s Error Quotes not ended at line %d \n",yytext,line);}}
      {IFR}+{Qu}+{IFR} { {printf(" %s Error Wrong usage of Quotes at line %d\n",yytext,line);}}
      {IFR}+{Qu} { {printf(" %s Error Quotes not started at line %d\n",yytext,line);}}
      {wl} { {printf(" %s Wrong language used\n",yytext); } }
      {IFR}
                                                      {if(comm==0&&quo==0){ insert(yytext,'v',line);}}
      {IFR}+"."+{IFR}
                                                      {if(comm==0){ insert(yytext,'v',line);} }
                                                      { if(comm==0) insert(yytext, 'p', line);}
      {PUN}
      "{"
                                                      { braces++; insert(yytext, 'p', line);}
      "}"
                                                      { braces--; insert(yytext,'p',line);}
      "("
                                                      { opbrc++; insert(yytext,'p',line);}
      ")"
                                                      { opbrc--; insert(yytext, 'p', line);}
      {Qu} {}
      0__0
 85
                                                      {if(comm==0){ insert(yytext, 'o', line);} }
      "++"
                                                      {if(comm==0){ insert(yytext, 'o', line);} }
                                                      {if(comm==0){ insert(yytext,'o',line);}}
      {per}
                                                       {if(comm==0){ insert(yytext, 'o', line);}}
    {op}
```

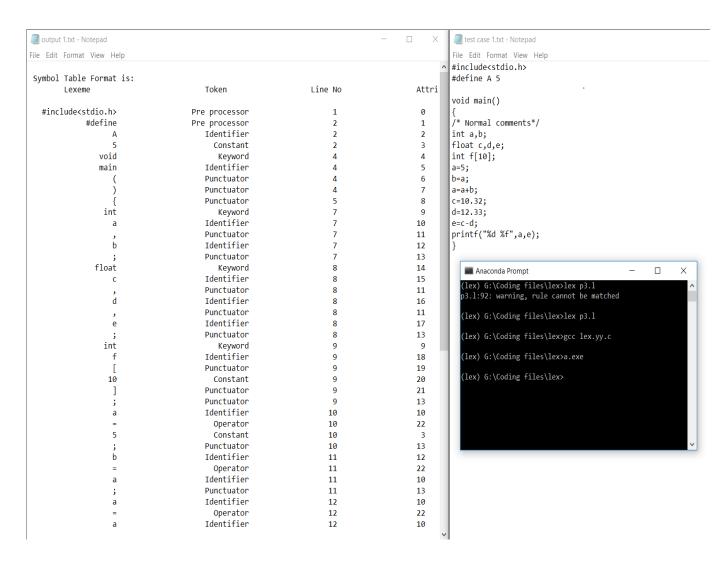
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₩
       p3.l
                                                       {if(comm==0){ insert(yytext,'o',line);}}
      {op}
                                                       {if(comm==0){ insert(yytext, 'o', line);}}
      {0}
      {WS} {}
      {SB} {if(comm==0){ insert(yytext,'p',line);} }
      . { printf(" %s Errors at line %d\n",yytext,line);}
      %%
      int main()
     yyin=fopen("test case 5.txt","r");
      yyout=fopen("output 5.txt","w");
      fprintf(yyout, "\n Symbol Table Format is:\n \tLexeme\t\t\tToken\t\tLine No\t\tAttribute Value\n");
      yylex();
      if(comm>0) { printf("Comments does not end"); }
      int yywrap()
      return(1);
      void insert(char *yytext,char t,int 1)
        int len1 = strlen(yytext);
        int i;
          char token[20];
          struct Storage *symbol,*temp,*nextptr;
          nextptr=head;
          switch(t)
              case 'k':
              case 'd':
120
              case 'l':
122
              case 'e':
              case 'c':
                  strcpy(token, "Keyword");
124
                  break;
126
              case 'C':
                  strcpy(token, "Constant");
128
129
                  break;
130
              case 's':
                  strcpy(token, "String");
                  break;
```

```
\blacktriangleleft \blacktriangleright
        p3.1
132
                    strcpy(token, "String");
133
                    break:
134
                case 'a':
                case 'v':
135
                case 'I':
136
                 case 'h':
137
                    strcpy(token, "Identifier");
138
139
                    break:
140
                case 'u':
141
142
                    strcpy(token, "User defined function");
143
                    break:
144
                case 'p':
145
146
                    strcpy(token, "Punctuator");
147
                    break;
                case 'o':
149
150
                    strcpy(token, "Operator");
151
                    break:
152
153
                case 'P':
154
                    strcpy(token, "Pre processor");
155
                    break;
156
           }
157
158
                for(i=0;i<var_cnt;i++,nextptr=nextptr->next)
159
                    symbol = nextptr;
                    if(strcmp(symbol->name,yytext)==0)
162
                         break:
163
164
                if(i==var_cnt)
165
                    temp = (struct Storage*)malloc(sizeof(struct Storage));
                    temp->name = (char*)malloc((len1+1)*sizeof(char));
167
                    strcpy(temp->name,yytext);
                    strcpy(temp->token,token);
                    temp->val = i;
170
171
                    temp->li = 1;
172
                    temp->next = NULL;
                    if(var_cnt==0)
173
174
                         head = temp;
175
                         symbol->next = temp;
176
Line 1/18 Column 1
```

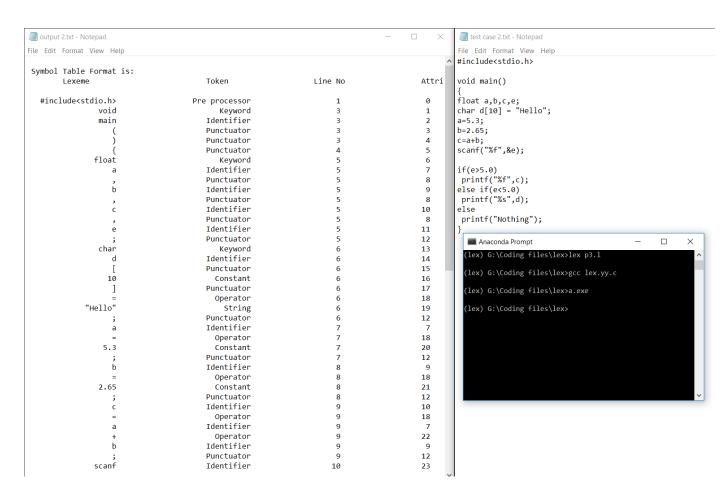
```
4▶
       p3.I
                         ×
159
               {
                   symbol = nextptr;
160
                   if(strcmp(symbol->name,yytext)==0)
161
162
                       break:
163
               }
               if(i==var_cnt)
164
165
                   temp = (struct Storage*)malloc(sizeof(struct Storage));
166
                   temp->name = (char*)malloc((len1+1)*sizeof(char));
167
                   strcpy(temp->name,yytext);
168
                   strcpy(temp->token,token);
169
                   temp->val = i;
170
171
                   temp->li = 1;
                   temp->next = NULL;
172
173
                   if(var_cnt==0)
                       head = temp;
174
175
                       symbol->next = temp;
176
177
                   var cnt++;
178
179
           fprintf(yyout, "\n%20s%30.30s%20d%20d", yytext, token, 1, i);
180
181
      }
182
183
```

# **Outputs:**

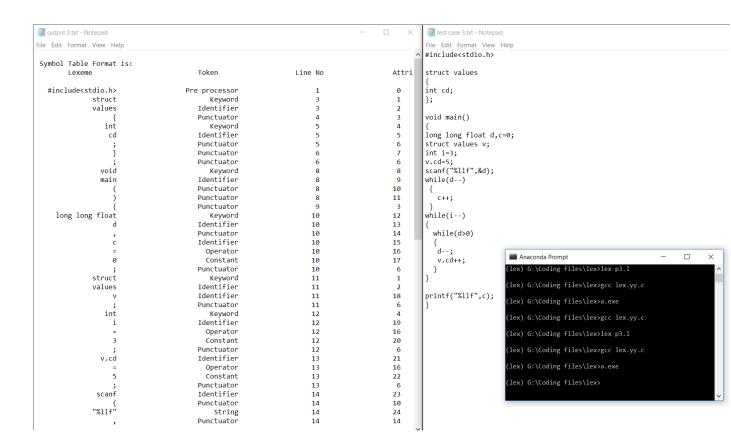
#### Test case 1:



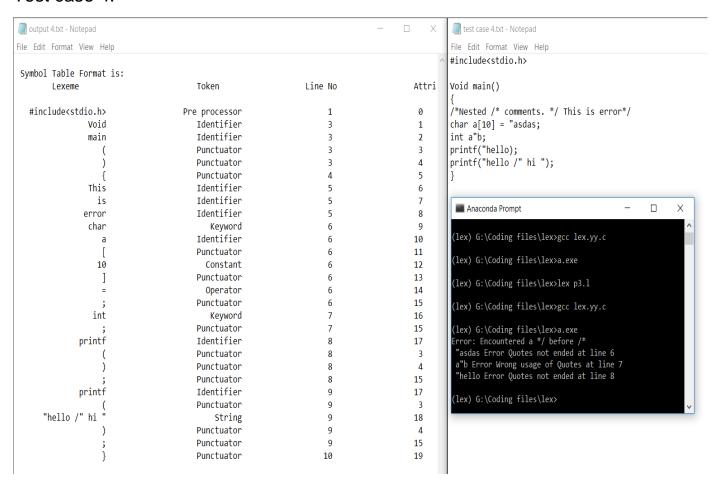
## Test case 2:



## Test case 3:



#### Test case 4:



## Test Case 5:

