Project 5

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Abstract

This project is to demonstrate the capabilities of implementing constructing and deconstructing HOL Terms using the tools and techniques - LATEX, AcuTeX, emacs and ML.

Each chapter documents the given problems with a structure of:

- 1. Problem Statement
- 2. Relevant Code
- 3. Execution Transcripts
- 4. Explanation of results

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Executive Summary

All requirements for this project are statisfied specifically, and by using HOL proved the below theorems:

Exercise 11.6.1

2.1 Problem Statement

In this exercise we need to prove the theorem: $\vdash \forall l_1 \ l_2$. LENGTH (APP $l_1 \ l_2$) = LENGTH l_1 + LENGTH l_2

2.2 Relevant Code

```
val LENGTH_APP=
TAC_PROOF(
([],
    ''!(11: 'a list)(12: 'a list).
    (LENGTH (APP 11 12)) =(LENGTH 11 + LENGTH 12)''),
Induct_on 'l1' THEN
ASM_REWRITE_TAC[ADD_CLAUSES, APP_def, LENGTH])
val _ = save_thm("LENGTH_APP", LENGTH_APP);
```

2.3 Execution Transcripts

2.3.1 Explanation of Results

The above results shows that the requirements are satisfied.

Exercise 11.6.2

3.1 Problem Statement

In this exercise we need to prove the theorem: $\vdash \operatorname{Map} f (\operatorname{APP} l_1 \ l_2) = \operatorname{APP} (\operatorname{Map} f \ l_1) (\operatorname{Map} f \ l_2)$

3.2 Relevant Code

```
val Map_def =
Define
'(Map f [] = []) /\ (Map f (x::f1) = f x::Map f (f1))';

val Map_APP =
TAC_PROOF(
([],
''Map f (APP l1 l2) = APP (Map f l1) (Map f l2) ''),
Induct_on 'l1' THEN
ASM_REWRITE_TAC[ADD_CLAUSES, Map_def, APP_def])

val _ = save_thm("Map_APP", Map_APP);
```

3.3 Execution Transcripts

3.3.1 Explanation of Results

The above results shows that the requirements are satisfied.

Exercise 11.6.3

4.1 Problem Statement

In this exercise we need to prove the following:

```
\vdash \forall f. \text{ nexpVal (Add (Num 0) } f) = \text{nexpVal } f \vdash \forall f_1 \ f_2. \text{ nexpVal (Add } f_1 \ f_2) = \text{nexpVal (Add } f_2 \ f_1) \vdash \forall f. \qquad \qquad \text{(nexpVal (Sub (Num 0) } f) = 0) \ \land \qquad \qquad \text{(nexpVal (Sub } f \text{ (Num 0))} = \text{nexpVal } f) \vdash \forall f_1 \ f_2 \ f_3. \qquad \qquad \text{nexpVal (Mult } f_1 \text{ (Mult } f_2 \ f_3)) = \text{nexpVal (Mult (Mult } f_1 \ f_2) \ f_3)
```

4.2 Relevant Code

```
(* Introduce the syntax of natural number expression nexp
val _ = Datatype
'nexp = Num num | Add nexp nexp | Sub nexp nexp | Mult nexp nexp';
(* Prove nexp Val (Add (Num 0) f) = nexp Val f
val Add_0 =
TAC.PROOF(([], ``!(f:nexp).nexpVal(Add (Num 0) f) = nexpVal(f)``),
Induct_on 'f' THEN
ASM_REWRITE_TAC[nexpVal_def] THEN
ASM_REWRITE_TAC[ADD_CLAUSES]);
val = save\_thm("Add\_0",Add\_0)
(* Prove nexp Val (Add f1 f2) = nexp Val (Add f2 f1)
val Add\_SYM =
TAC.PROOF(([], ``!f1 f2.nexpVal (Add f1 f2) = nexpVal (Add f2 f1)``),
```

```
Induct_on 'f1' THEN
PROVE_TAC [nexpVal_def, ADD_SYM]);
val _ = save_thm("Add_SYM", Add_SYM)
(* Prove (nexp Val (Sub (Num 0) f) = 0)
* )
      (nexp Val (Sub f (Num 0)) = nexp Val f)
(*
val Sub_0 =
TAC_PROOF(([],
", "! f.(\text{nexpVal} (\text{Sub} (\text{Num } 0) f)) = 0)
 (nexpVal (Sub f (Num 0)) = nexpVal f ) ``),
STRIP_TAC THEN
ASM_REWRITE_TAC[nexpVal_def] THEN
PROVE_TAC[nexpVal_def, SUB_0]);
val = save_thm("Sub_0", Sub_0)
(* Prove f 1 f 2 f 3.
*)
(*
     nexp Val (Mult f1 (Mult f2 f3)) =
                                                        *)
     nexpVal (Mult (Mult f1 f2) f3)
                                                        *)
val Mult_ASSOC =
TACPROOF(([], ``!f1 f2 f3.
  nexpVal (Mult f1 (Mult f2 f3)) =
  nexpVal (Mult (Mult f1 f2) f3) ''),
REPEAT STRIP_TAC THEN
PROVE_TAC[nexpVal_def, MULT_ASSOC]);
val _ = save_thm("Mult_ASSOC",Mult_ASSOC)
```

4.3 Execution Transcripts

```
1
> # <<HOL message: Defined type: "nexp">>
> # val nexp_one_one =
   |- (!(a :num) (a' :num). (Num a = Num a') <=> (a = a')) /\
   (!(a0 :nexp) (a1 :nexp) (a0' :nexp) (a1' :nexp).
      (Add a0 a1 = Add a0' a1') <=> (a0 = a0') /\ (a1 = a1')) /\
   (!(a0 :nexp) (a1 :nexp) (a0' :nexp) (a1' :nexp).

(Sub a0 a1 = Sub a0' a1') <=> (a0 = a0') / (a1 = a1')) / (a0 :nexp) (a1 :nexp) (a0' :nexp) (a1' :nexp).
     (Mult a0 a1 = Mult a0' a1') <=> (a0 = a0') /\ (a1 = a1'):
   thm
> # val nexp_distinct_clauses =
   |- (!(a1 :nexp) (a0 :nexp) (a :num). Num a <> Add a0 a1) /\
   (!(a1 :nexp) (a0 :nexp) (a :num). Num a <> Sub a0 a1) /\
(!(a1 :nexp) (a0 :nexp) (a :num). Num a <> Mult a0 a1) /\
   (!(a1' :nexp) (a1 :nexp) (a0' :nexp) (a0 :nexp).
Add a0 a1 <> Sub a0' a1') /\
   (!(a1':nexp) (a1 :nexp) (a0':nexp) (a0 :nexp). Add a0 a1 <> Mult a0' a1') /\
   !(a1' :nexp) (a1 :nexp) (a0' :nexp) (a0 :nexp).
    Sub a0 a1 <> Mult a0' a1':
   thm
> # # # # # Definition has been stored under "nexpVal_def"
val nexpVal_def =
   |- (!(f :num). nexpVal (Num f) = f) /
   (!(f1 : nexp) (f2 : nexp).
      nexpVal (Add f1 f2) = nexpVal f1 + nexpVal f2) /\
   (!(f1 :nexp) (f2 :nexp).
      nexpVal (Sub f1 f2) = nexpVal f1 - nexpVal f2) /
   !(f1 :nexp) (f2 :nexp).
     nexpVal (Mult f1 f2) = nexpVal f1 * nexpVal f2:
   thm
> # # # # val Add_0 =
   |- !(f :nexp). nexpVal (Add (Num (0 :num)) f) = nexpVal f:
   thm
> > # > # # # Meson search level: .......
Meson search level: .....
Meson search level: .....
Meson search level: .....
val Add_SYM =
   |- !(f1 :nexp) (f2 :nexp). nexpVal (Add f1 f2) = nexpVal (Add f2 f1):
> > # # # # # Meson search level: ....
val Sub_0 =
   |- !(f :nexp).
     (nexpVal (Sub (Num (0 :num)) f) = (0 :num)) /\
     (nexpVal (Sub f (Num (0 :num))) = nexpVal f):
> > + # # # # Meson search level: ......
val Mult_ASSOC =
   |-!(f1 :nexp) (f2 :nexp) (f3 :nexp).
     nexpVal (Mult f1 (Mult f2 f3)) = nexpVal (Mult (Mult f1 f2) f3):
```

4.3.1 Explanation of Results

The above results shows that the requirements are satisfied.

Appendix A: Exercise 11.6.1 and 11.6.2

The following code is from the file exTypeScript.sml (* Exercise: Chapter 11 Author: Bharath Karumudi *) Date: Aug 10, 2019 structure exTypeScript = struct open HolKernel Parse boolLib bossLib; open listTheory TypeBase arithmeticTheory (* = = interactive mode = = =map load ["boolTheory", "TypeBase", "bexpTheory"]; open boolTheory TypeBase bexpTheory ==== end interactive mode ===== *)val _ = new_theory "exType"; $val APP_def =$ Define (APP [] (1: 'a list) = 1) /(APP (h :: (11 : 'a list)) (12: 'a list) = h :: (APP 11 12))Exercise: 11.6.1l 1 l 2 . LENGTH (APP l 1 l 2) = LENGTH l 1 + LENGTH l 2 (* val LENGTH_APP= TAC_PROOF(([]]"!(l1: 'a list)(l2: 'a list). (LENGTH (APP 11 12)) = (LENGTH 11 + LENGTH 12), Induct_on 'l1' THEN ASM_REWRITE_TAC [ADD_CLAUSES, APP_def, LENGTH]) val _ = save_thm("LENGTH_APP", LENGTH_APP);

```
(* Exercise: 11.6.2
(* 'Map f (APP l 1 l 2 ) = APP (Map f l 1 ) (Map f l 2 )
val Map_def =
Define
'(Map f [] = []) /\ (Map f (x:: f1) = f x:: Map f (f1))';
val Map\_APP =
TAC_PROOF(
([]]
"
(Map f (APP 11 12) = APP (Map f 11) (Map f 12) "
(Map f 12) "
Induct_on 'l1' THEN
ASM.REWRITE_TAC[ADD_CLAUSES, Map_def, APP_def])
val _ = save_thm("Map_APP", Map_APP);
(* Exporting Theory
                                                     *)
val _ = export_theory();
val = print_theory "-";
end (* Structure *)
```

Appendix B: Exercise 11.6.3

The following code is from the file nexpScript.sml

```
(* Exercise: Chapter 11.6.3
 Author: Bharath Karumudi
 Derived\ from:\ bexpScript.sml
(* Date: Aug 10, 2019
structure nexpScript = struct
open HolKernel Parse boolLib bossLib;
open TypeBase boolTheory arithmeticTheory
(* = = interactive mode = = = 
map load ["boolTheory", "TypeBase", "nexpTheory"];
open boolTheory TypeBase nexpTheory
==== end interactive mode ===== *)
val _ = new_theory "nexp";
(* Introduce the syntax of natural number expression nexp
val _ = Datatype
'nexp = Num num | Add nexp nexp | Sub nexp nexp | Mult nexp nexp';
(* Prove that identical nexps have identical components
val nexp_one_one = one_one_of ':nexp''
val _ = save_thm("nexp_one_one", nexp_one_one)
(* Prove that the different forms of bexp expressions are distinct
val nexp_distinct_clauses = distinct_of ': nexp''
val _ = save_thm("nexp_distinct_clauses", nexp_distinct_clauses)
(* Define the semantics of nexp expressions
val nexpVal_def =
Define
(\text{nexpVal }(\text{Num }f) = f)/\
```

```
(\text{nexpVal } (\text{Add } f1 \text{ } f2) = (\text{nexpVal } f1) + (\text{nexpVal } f2)) / (
(\text{nexpVal (Sub f1 f2}) = (\text{nexpVal f1}) - (\text{nexpVal f2})) / 
(\text{nexpVal } (\text{Mult } \text{f1 } \text{f2}) = (\text{nexpVal } \text{f1}) * (\text{nexpVal } \text{f2}))
(* Prove nexp Val (Add (Num 0) f) = nexp Val f
val Add_0 =
TAC.PROOF(([], ``!(f:nexp).nexpVal(Add (Num 0) f) = nexpVal(f)``),
Induct_on 'f' THEN
ASM_REWRITE_TAC[nexpVal_def] THEN
ASM_REWRITE_TAC[ADD_CLAUSES]);
val = save_thm("Add_0", Add_0)
(* Prove nexpVal (Add f1 f2) = nexpVal (Add f2 f1)
val Add_SYM =
TAC.PROOF(([], 'if1 f2.nexpVal (Add f1 f2) = nexpVal (Add f2 f1)'),
Induct\_on\ `f1\ `THEN"
PROVE_TAC [nexpVal_def, ADD_SYM]);
val _ = save_thm("Add_SYM", Add_SYM)
(* Prove (nexp Val (Sub (Num 0) f) = 0)
* )
      (nexp Val (Sub f (Num 0)) = nexp Val f)
(*
val Sub_0 =
TACPROOF(([],
''! f. (\text{nexpVal (Sub (Num 0) f }) = 0)
 (nexpVal (Sub f (Num 0)) = nexpVal f),
STRIP_TAC THEN
ASM_REWRITE_TAC[nexpVal_def] THEN
PROVE_TAC[nexpVal_def, SUB_0]);
val = save_thm("Sub_0", Sub_0)
(* Prove f1 f2 f3.
*)
(*
      nexpVal (Mult f1 (Mult f2 f3)) =
                                                         * )
     nexp \ Val \ (Mult \ (Mult \ f1 \ f2) \ f3)
                                                         * )
```