

~~Sorting~~
 * To check an array is sorted or not through recursion

bool isSorted (vector<int> &arr, int i) {

// base case

if (i == arr.size() - 1) {

return true;

}

if (arr[i] < arr[i+1]) {

return false;

}

return isSorted (arr, i+1);

}

int main () {

int size;

cin >> size;

vector<int> arr (size);

for (int i = 0; i < arr.size(); i++) {

cin >> arr[i];

}

int i = 0;

if (isSorted (arr, i)) {

cout << "Array is Sorted";

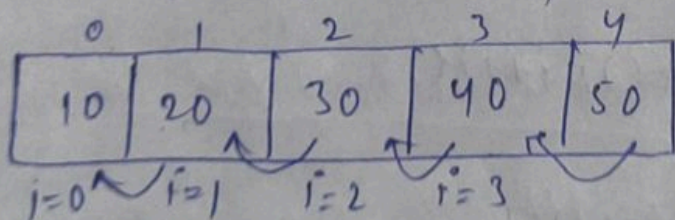
}

else {

cout << "Array is not Sorted";

}

* As we have passed an array by reference but any changes will be made inside the original array no copy of an array will be created.



10 > 20 T 30 > 20 T 40 > 30 T 50 > 40 T

if (arr[i] > arr[i+1]) {

return true;

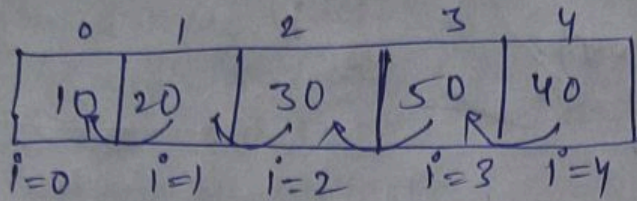
else false;

output

10 20 30 40 50
 Array is sorted

10 20 30 50 40
 Array is Not sorted

if i reach to last index after traversing to each element then we can say that i have checked all i+1 index
 \therefore array is in sorted order. return true.



10 < 20 T \rightarrow 30 > 20 T \rightarrow 50 > 30 T \rightarrow 40 > 50 F (return false)

* Binary search with Recursion:

```
int BinarySearch (vector<int> arr, int start, int end, int &key)
{
    if (start > end) {
        return -1;
    }
    int mid = start + (end - start) / 2;
    if (arr[mid] == key) {
        return mid;
    }
    if (arr[mid] < key) {
        return BinarySearch (arr, mid + 1, end, key);
    }
    else {
        return BinarySearch (arr, start, mid - 1, key);
    }
}
```

```
int main () {
    int size;
    cin >> size;
    vector<int> arr (size);
    for (int i = 0; i < arr.size(); i++) {
        cin >> arr[i];
    }
    int start = 0, end = arr.size() - 1;
    int key;
    cin >> key;
    int index = BinarySearch (arr, start, end, key);
    cout << index;
}
```


output

→ 10 20 30 40 50 60 → Input

60 → key.

5 → index found.

→ 10 20 30 40 50 60 → Input

70 key.

-1 → index not found.

* subsequence of a string:

① I/P: → "abc"

O/P: → print all subsequence. i.e.

$\langle - \rangle, \langle a \rangle, \langle b \rangle, \langle c \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, c \rangle, \langle a, b, c \rangle = 8$

→
Power set

The Question uses famous pattern named include - exclude pattern.

② I/P → "xy"

O/P → $\langle - \rangle, \langle x \rangle, \langle y \rangle, \langle x, y \rangle = 4$

We can observe that it is power set i.e. 2^n
n here stands for no. of length of string.

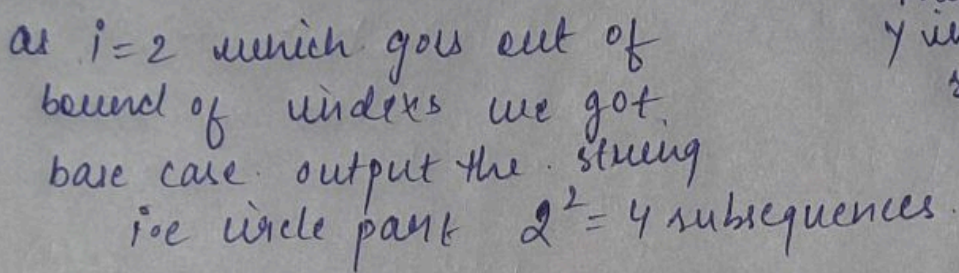
for ① length of string = 3

power set = $2^3 = 8$ → Total no. of subsequence strings

for ② length of string = 2

power set = $2^2 = 4$ → Total no. of subsequence strings

* Include-Exclude pattern means we will have an empty string which means we are creating extra space for new string. One time we include a character to empty string & one time we will exclude that character in an empty string. We will iterate through $i=0$ for every character in a string once we include-exclude move forward i.e. $i=i+1$;



```

// exclude
printSubsequences (str, ans, i+1);

```

```

// include
// print subsequences (str, ans)
// ans.push_back(str[i]);
// print subsequences (str, ans, i+1);

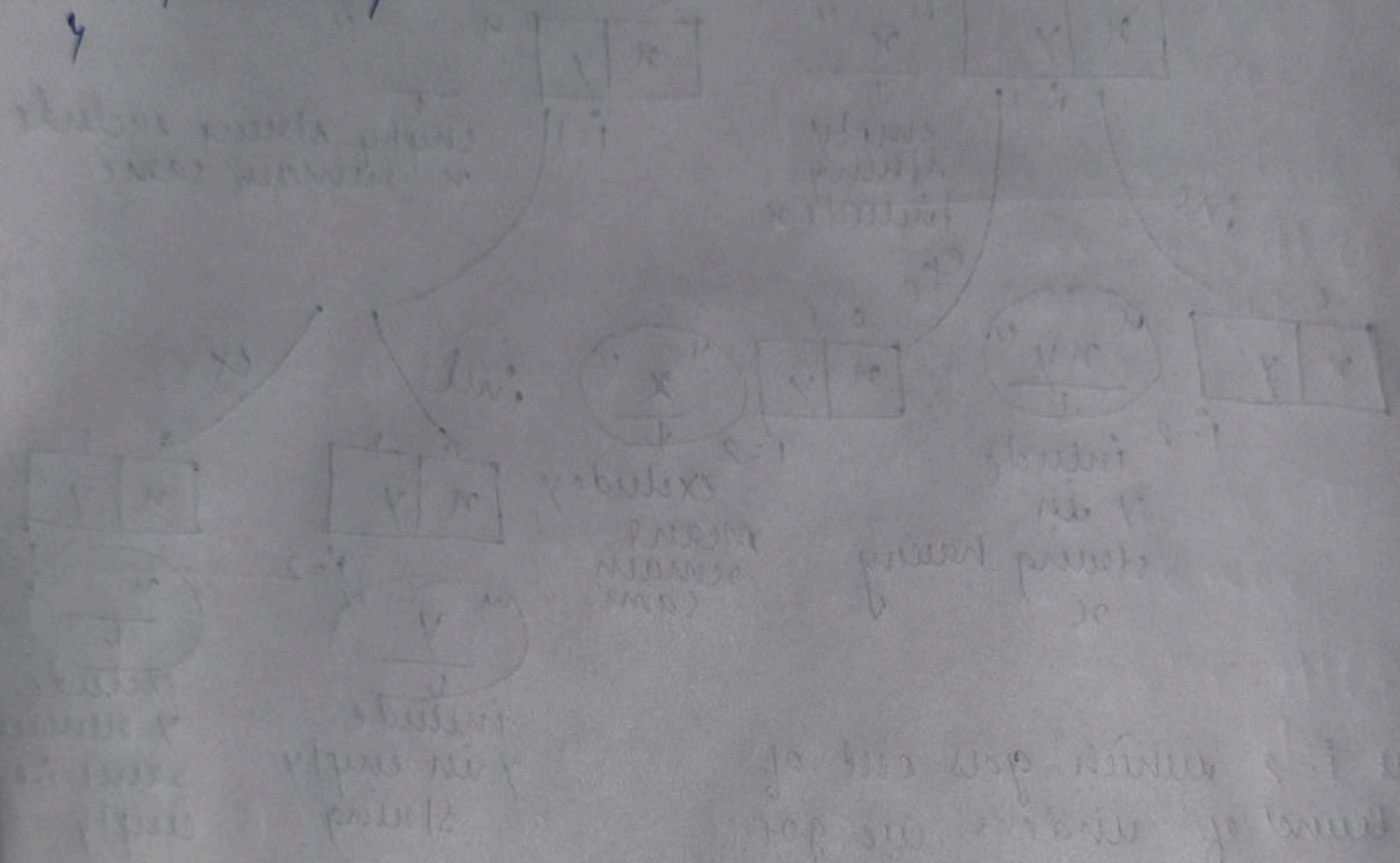
```



```

int main() {
    string str;
    getline(cin, str);
    int i = 0;
    string ans = "";
    printSubsequences(str, ans, i);
}

```



as the string is empty, we print the string and return. This is the base case of the recursion.

In the recursive step, we call the function twice for each string. The first call is for the string with the first character removed, and the second call is for the string with the first character added to the current string.

The function returns the list of all possible subsequences of the input string.

The time complexity of this function is $O(2^n)$, where n is the length of the string. This is because there are 2^n possible subsequences of a string of length n .

The space complexity of this function is $O(n)$, as the maximum depth of the recursion stack is n .