

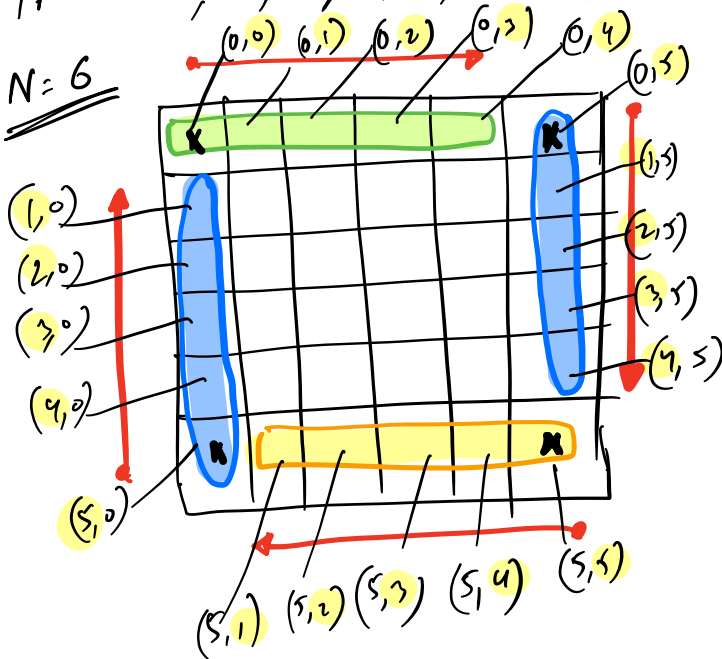
~~Q~~ Given a  $N \times N$  matrix.  
Print the boundary in clockwise fashion!

I/P

1	2	3	4	5
6	7	8	9	10
11	12	13	14	15
16	17	18	19	20
21	22	23	24	25

O/P: 1, 2, 3, 4, 5, 10, 15, 20, 25, 24, 23, 22, 21, 16, 11, 6

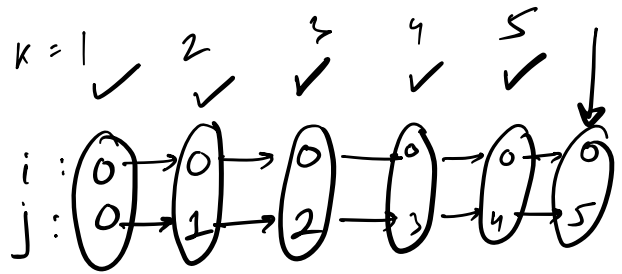
N = 6



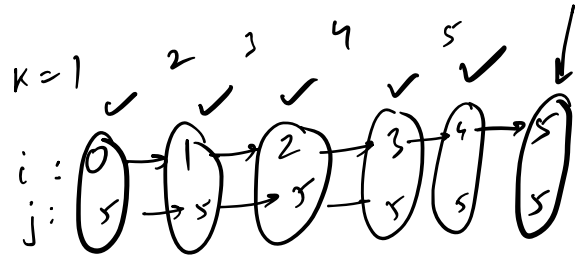
N-1

1. 5 elements to right
2. 5 ——— down
- 5 ——— left
- 5 ——— up

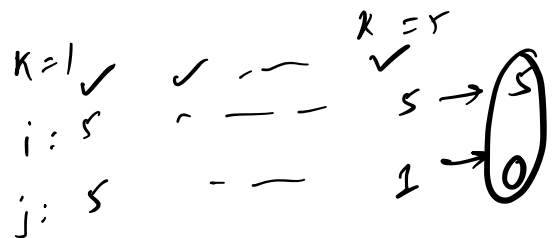
$i = 0, j = 0;$   
 $\rightarrow O(N)$   
 $\{ (k=1; k < N; k++) \{$   
 $\quad \text{print}(A[i][j]);$   
 $\quad j++;$   
 $\}$



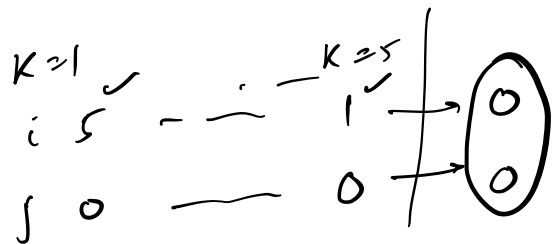
$\rightarrow O(N)$   
 $\{ (k=1; k < N; k++) \{$   
 $\quad \text{print}(A[i][j]);$   
 $\quad i++;$



$\rightarrow O(N)$   
 $\{ (k=1; k < N; k++) \{$   
 $\quad \text{print}(A[i][j]);$   
 $\quad j--;$



$\rightarrow O(N)$   
 $\{ (k=1; k < N; k++) \{$   
 $\quad \text{print}(A[i][j]);$   
 $\quad i--;$

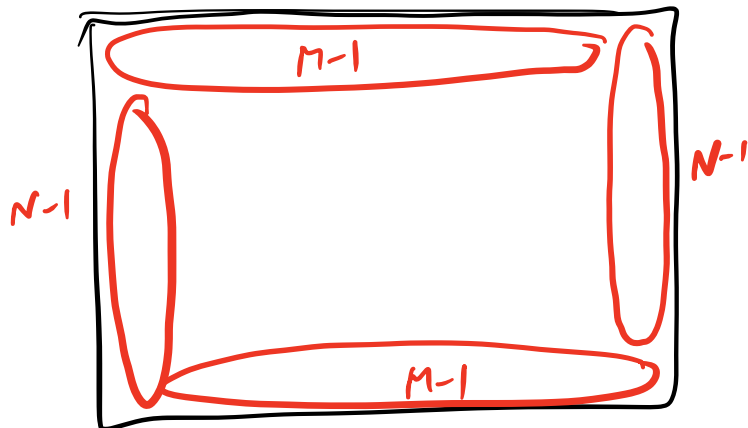


$$TC = 4 \times O(N)$$

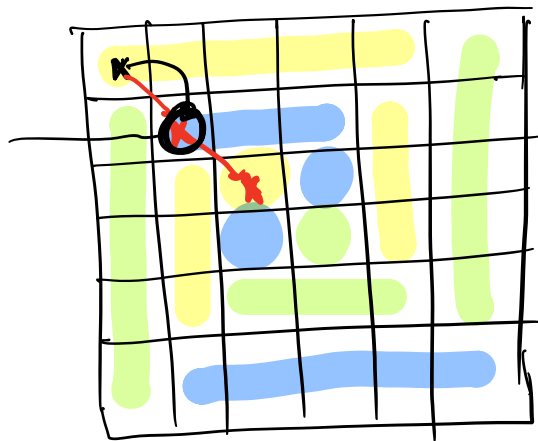
$$TC: O(N)$$

$$SC: O(1)$$

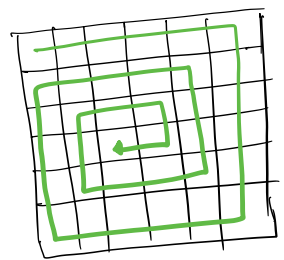
#  $N \times M$  matrix!



Q Given a  $N \times N$  matrix.  
Print it in **SPIRAL** fashion!



$i \quad j$   
 $N=6 \quad 0 \quad 0$   
 $N=4 \quad 1 \quad 1$   
 $N=2 \quad 2 \quad 2$   
 $N=0 \quad \text{STOP!}$



$i = 0, j = 0;$

$\text{while}(N > 1) \{$

$\{ (k=1; k < N; k++) \{$   
 $\text{print}(A[i][j]);$   
 $j++;$   
 $\}$

$\{ (k=1; k < N; k++) \{$   
 $\text{print}(A[i][j]);$   
 $i++;$

$\}$   
 $\{ (k=1; k < N; k++) \{$   
 $\text{print}(A[i][j]);$   
 $j--;$

$\}$   
 $\{ (k=1; k < N; k++) \{$   
 $\text{print}(A[i][j]);$   
 $i--;$

$\}$

$N = N - 2;$

$i++;$

$j++;$

$\}$   
 $\text{if}(N == 1) \{$   
 $\text{print}(A[i][j]);$

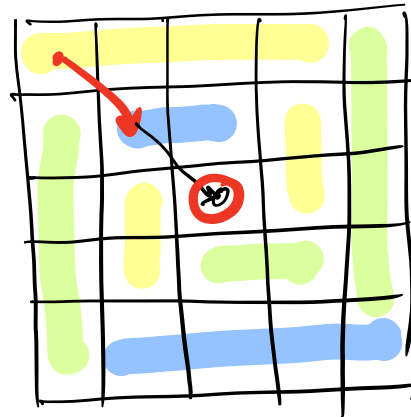
$\}$

$TC = O(N^2)$

$SC = O(1)$

$N \rightarrow 0 \text{ DP}$

$N = 5$

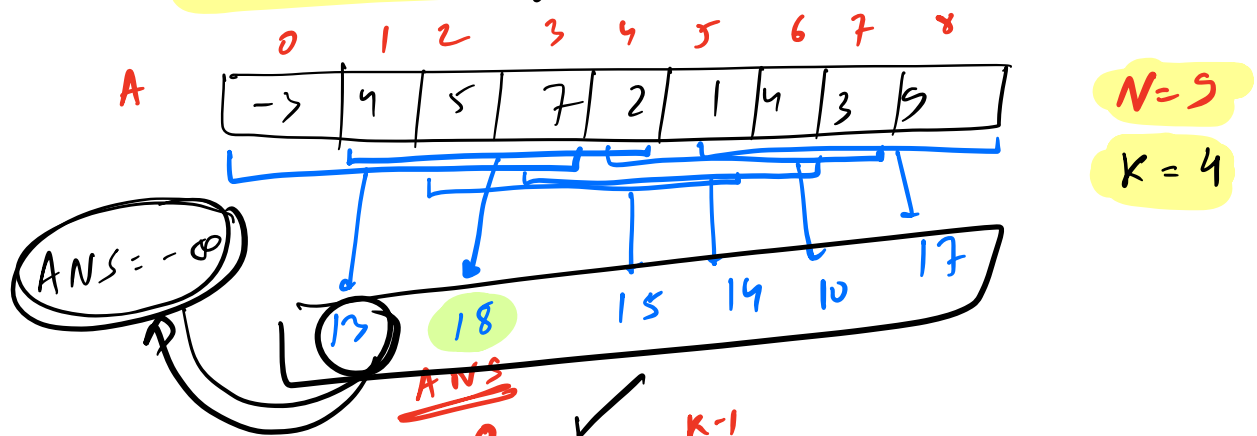


$N \quad i \quad j$   
 $\checkmark 5 \quad 0 \quad 0$   
 $\checkmark 3 \quad 1 \quad 1$

$\textcircled{1} \quad \textcircled{2 \quad 2}$

$N = 1$   
 $N - 1 \rightarrow 0$

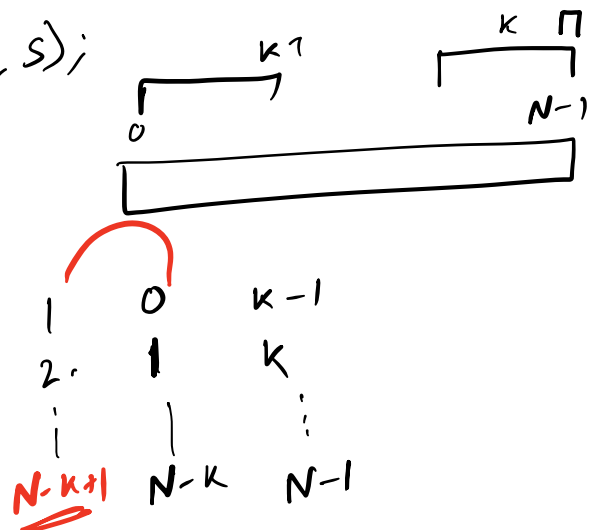
Given  $N$  elements, return the Maximum Subarray Sum of len =  $K$ !



1) BF

$L = 0, R = K - 1$   
 $ANS = -\infty$   
 while ( $R < N$ ) {  
   //  $[L, R]$   
    $S = 0$ ;  
   for ( $i = L; i \leq R; i++$ ) {  
      $S = S + A[i]$ ;  
   }  
    $ANS = \max(ANS, S)$ ;  
    $L++$ ,  $R++$ ;  
 }

Complexity:  $O(N \cdot K)$



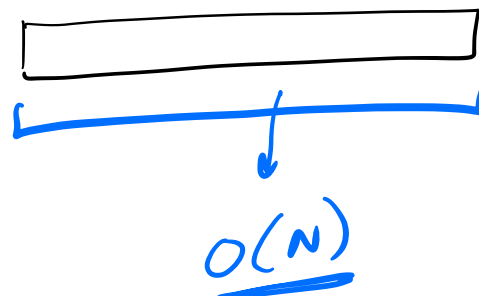
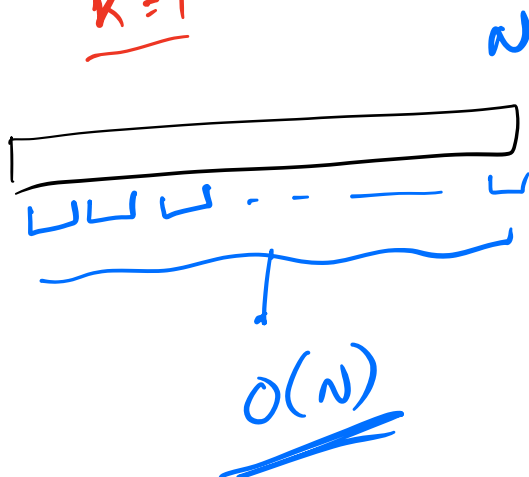
# SA of size  $K$  in an array of size  $N$   
 $= N - K + 1$

$$TC = K \times (N - K + 1)$$

$$TC = O(K(N - K))$$

$K = N$

$K = 1$



$K = N/2$



#SA  $\approx K$

$N - K = K$

$N = 2K$

$K = N/2$

#SA  $\rightarrow \frac{N}{2} \times \frac{N}{2} = \frac{N^2}{4}$

$TC: O(N^2)$

$SC = O(1)$

## II) Prefix Sum

1. Build the ps Array!  $\rightarrow$   $TC: O(N)$   
 $SC: O(N)$

2.  $L=0, R=k-1$   
 $ANS = -\infty$

while ( $R < N$ ) {

//  $[L, R]$

$S = ps[R] - ps[L-1]$   $\rightarrow O(1)$

$ANS = \max(ANS, S);$

$L++, R++;$

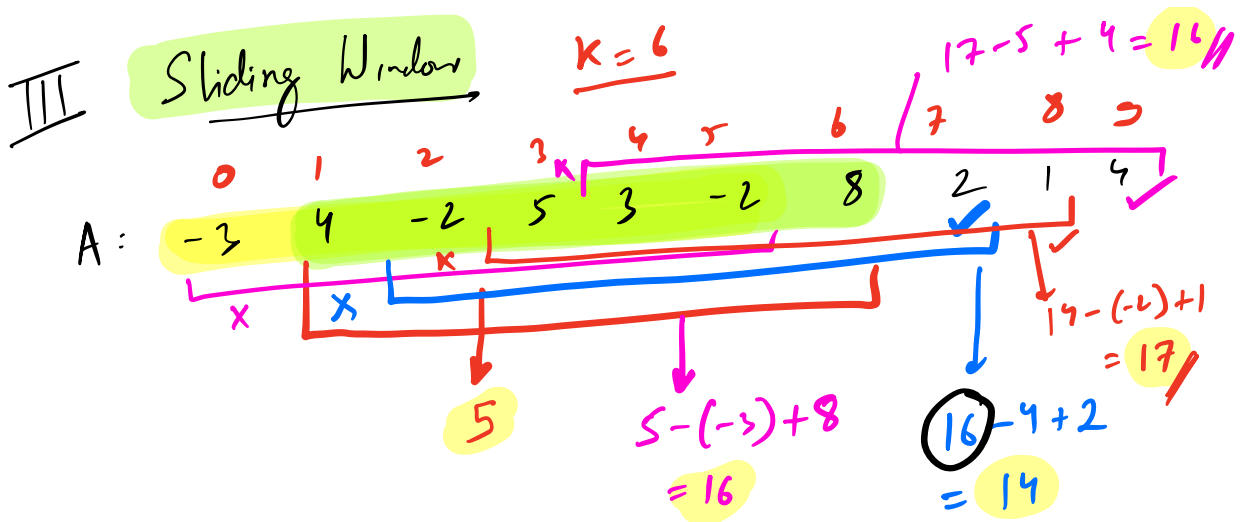
}  
return ANS //

$TC = O(N + N - K)$

$TC = O(N - K)$

$O(N)$

$SC = O(N)$



**ANS = 17**

$K=6, N=10$

L	R
0	5
1	6
2	7
3	8
4	9

ANS

Iterate over them & find the Sum

Sum = Sum - A[0] + A[6]

Sum = Sum - A[1] + A[7]

Sum = Sum - A[2] + A[8]

Sum = Sum - A[3] + A[9]

**Sum = Sum - A[L-1] + A[R]**



ANS =  $-\infty$

sum = 0

$\{ (i=0; i < k; i++) \}$   $\rightarrow O(k)$   
    sum = sum + A[i];

}

ANS = sum

L = 1, R = k;

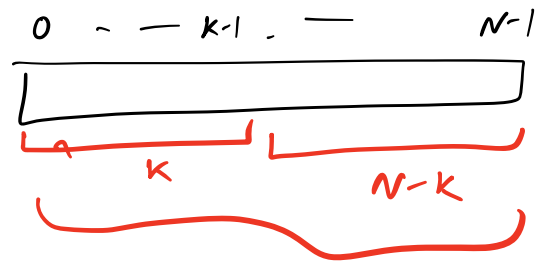
while (R < N)  $\rightarrow O(N-k)$   
    sum = sum - A[L-1] + A[R];

ANS = max(ANS, sum);

L++, R++;

}

ret ANS

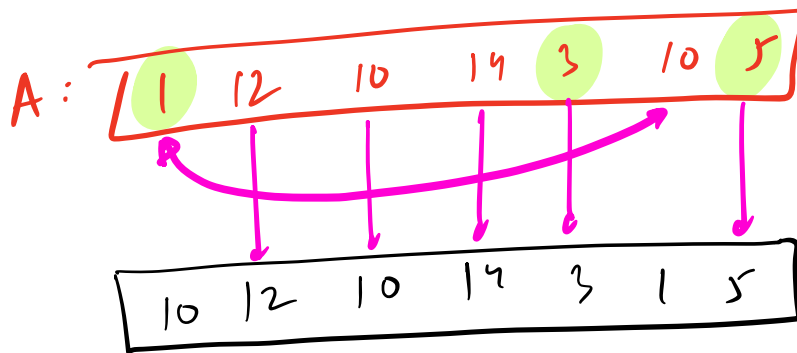


TC =  $O(k + N - k)$

**TC =  $O(N)$**

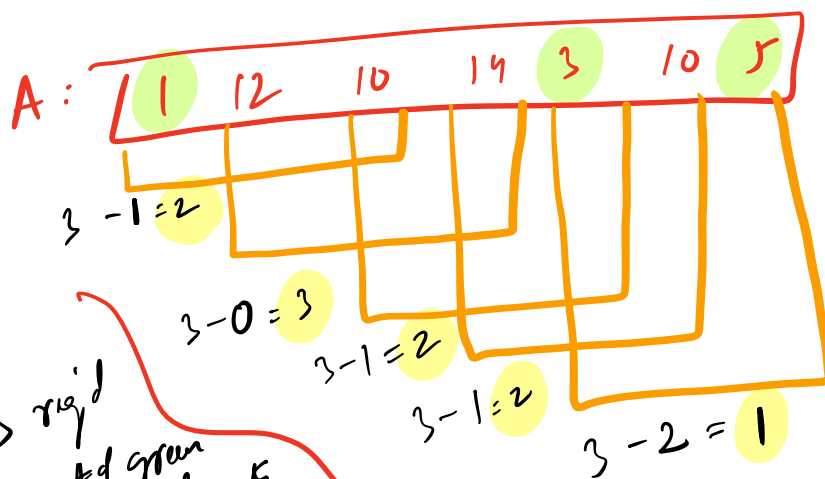
**SC =  $O(1)$**

Q Given an array A & an int B  
find the MIN no. of SWAPS required to bring  
all the elements  $\leq B$  together!



B = 8

ANS  
1



B = 8

Count of elements  
 $\leq B = 3$   
k = 3

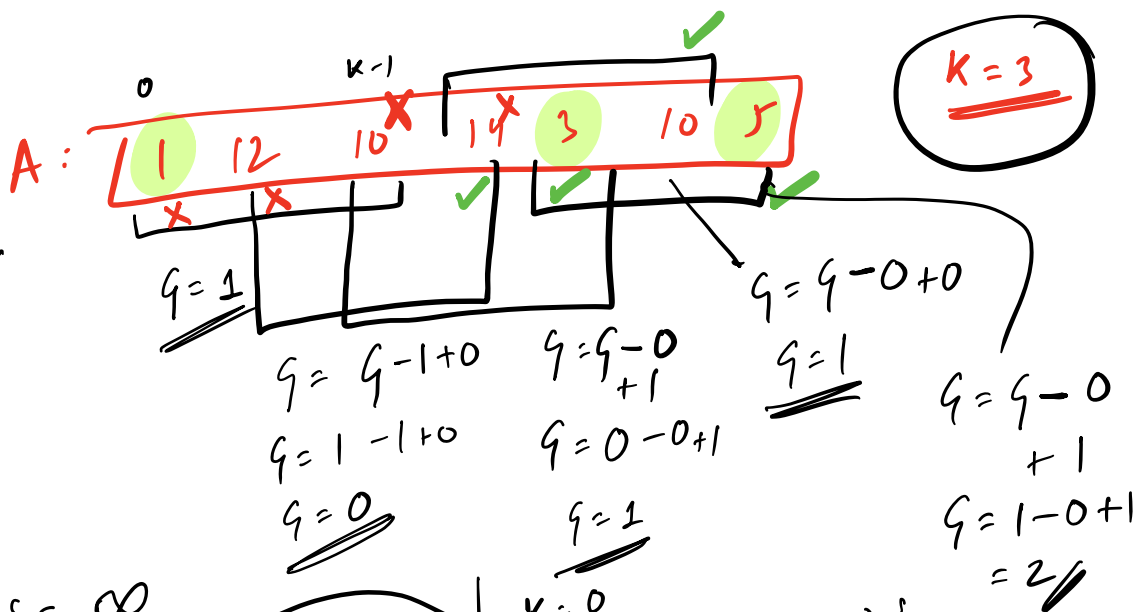
Swaps req'd  
= k - # of green elements

MIN = 1  $\rightarrow$  ANS!

1) BF  $\rightarrow$  TC:  $O(N^2)$

II

Keep track  
of group  
elements



ANS =  $\infty$

$q = 0$

$\{ (i=0; i < K; i++) \}$   
 $\{ if (A[i] <= B) q++; \}$

}

ANS =  $K - q$

$L = 1, R = K;$

while ( $R < N$ ) {

if ( $A[L-1] <= B$ )  $q--;$

if ( $A[R] <= B$ )  $q++;$

ANS = min (ANS,  $K - q$ );

$L++, R++;$

}

ret ANS

$K = 0$

$\{ (i=0; i < N; i++) \}$   
 $\{ if (A[i] <= B) K++; \}$

}

$TC = O(N)$

$SC = O(1)$