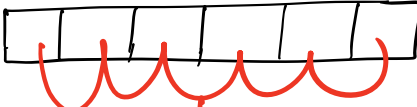
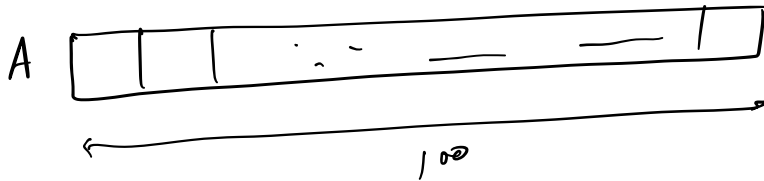


Array →

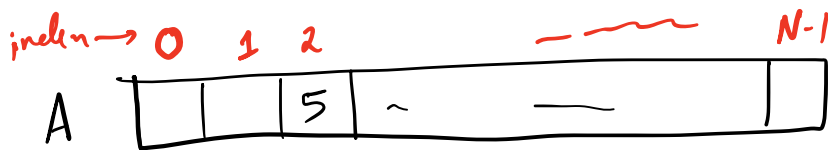
Linear D.S. → 
type is all same!



int A[100];
↑ ↑ ↑
data type Name Size!

→ N

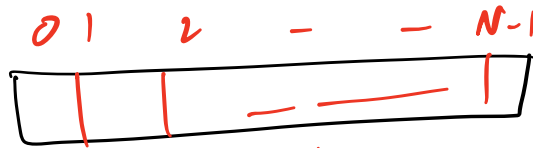
int A[N];



$A[2] = 5$ → put 5 at index 2
print(A[2]) → print val at index 2.



Q Given an array of size N .
Print the entire array!



$N=5$
 \downarrow
 10^5

$\text{print}(A[0])$
 $\text{— } (A[1])$
 $\text{— } (A[2])$
 $\text{— } (A[3])$
 $\text{— } (A[4])$

0
 1
 2
 3
 4

```

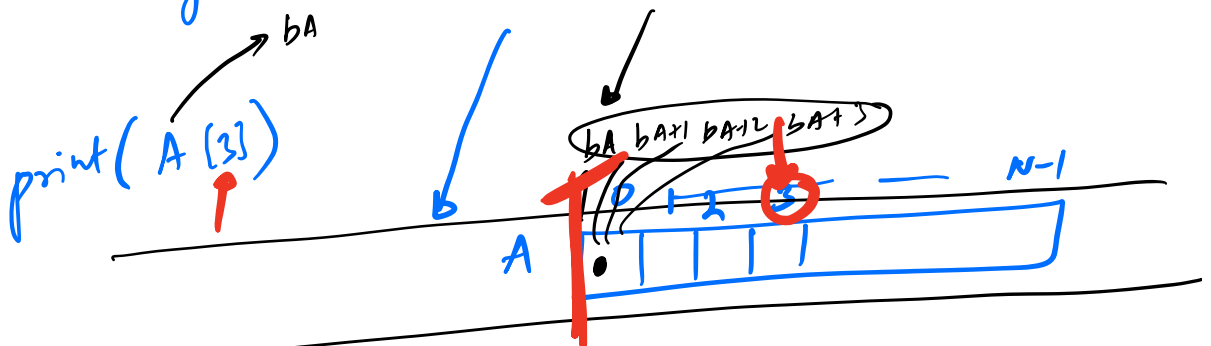
{ (i=0; i < N; i++) {
  print(A[i]);
}

```

$N \times O(1)$

TC: $O(N)$

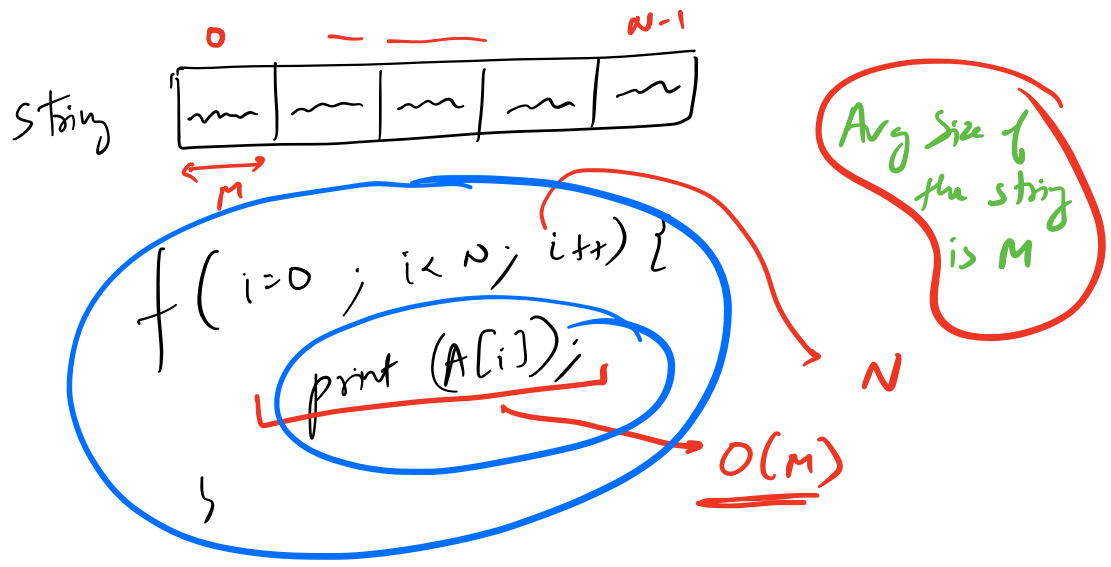
TC for accessing a single element of an array!



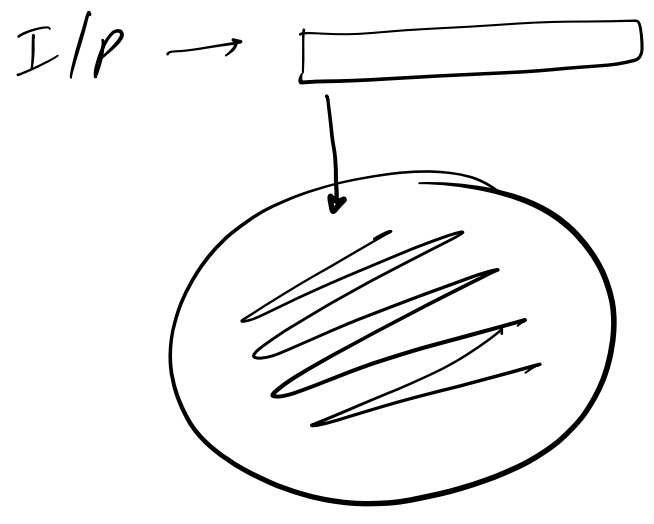
$A[3] \rightarrow bA + 3 \times 4$

$bA+7 \dots bA+7$

$O(1)$



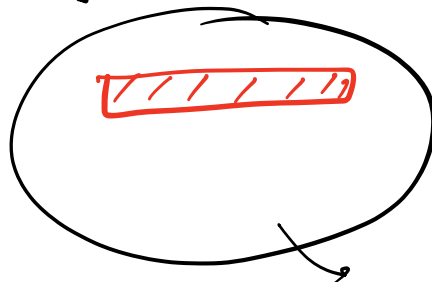
$TC = O(N \cdot M)$



$SC: O(1)$

Q

N



SC: $O(1)$

1 2 ... N

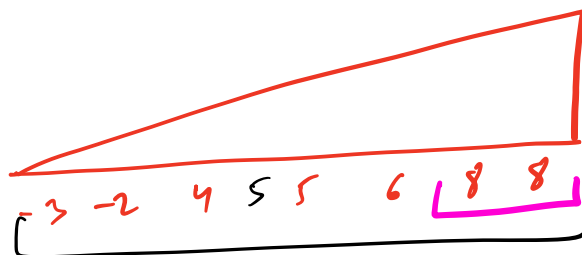
Q

Given N elements \rightarrow Array.

Count the no. of elements having atleast 1 element greater than itself!

$A = [-3, -2, 6, 8, 4, 8, 5]$

Below the array, there are checkmarks and crosses: $\checkmark, \checkmark, \checkmark, \times, \checkmark, \times, \checkmark$. An arrow points from the last element (5) to the right.



ANS \rightarrow $N - \# \text{ of occ. of MAX element}$

2. `int cnt = 0;`
`{ (i=0; i < N; i++) {`
`if (A[i] == mC) {`
`cnt++;`
`}`
`}`

$mC = 5$

3 5 2 2 5 4

i

$\frac{cnt}{\cancel{0} \cancel{1} 2}$

$TC = O(N)$

$SC = O(1)$

3. $ANS \rightarrow \underline{N - cnt}$

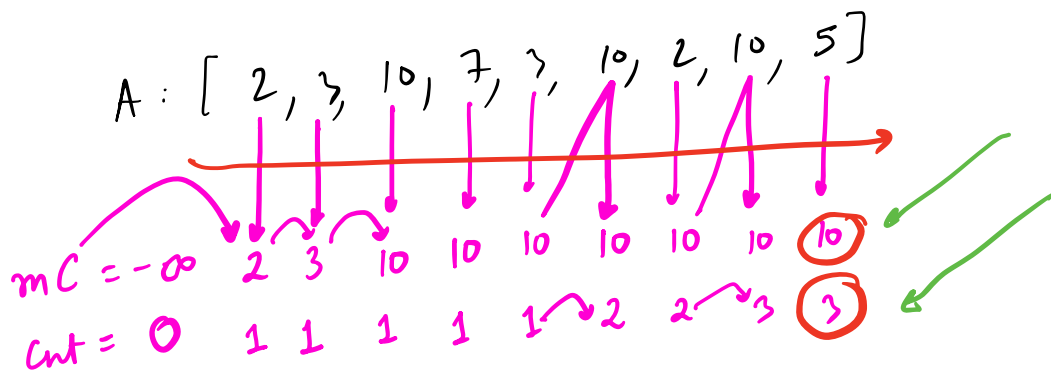
$$TC = O(N) + O(N) = O(2N) = O(N)$$

$$TC = O(N)$$

$$SC = O(1)$$

II

Q



$mc = -\infty$

$cnt = 0$

```
f( i = 0; i < N; i++ ) {  
    if( A[i] > mc ) {  
        mc = A[i];  
        cnt = 1;  
    }  
    else if( A[i] == mc ) {  
        cnt++;  
    }  
}
```

Ans $\rightarrow N - cnt$

$N \times O(1)$

$\boxed{TC: O(N)}$

$\boxed{SC: O(1)}$



Q

Given an array of size N & K

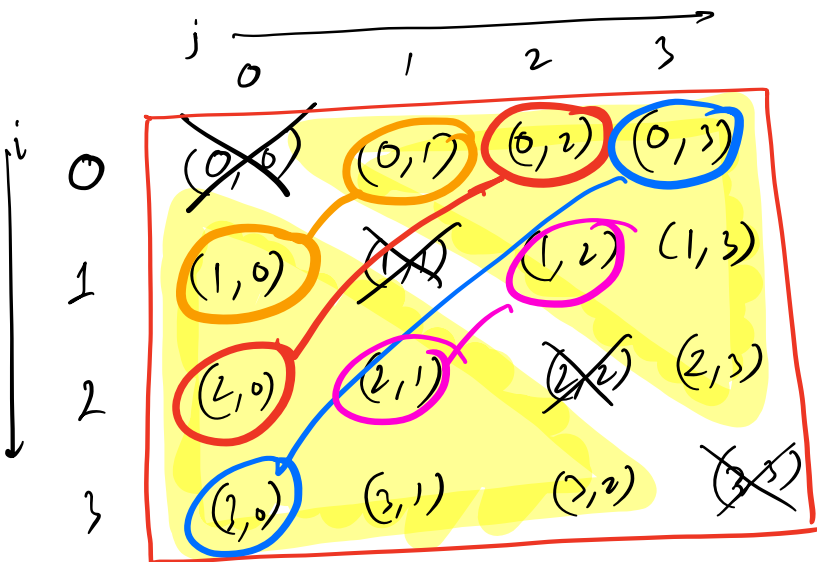
Check if there exists a pair (i, j)

Such that $A[i] + A[j] == K$ & $i \neq j$

$A: [3, -2, 1, 7, 3, 6, 8]$ $K = 10$

$(i, j) \rightarrow (3, 5)$

$A: [2, 4, -3, 7]$ $K = 8$



	$j: [i+1 \dots N-1]$
$i=0$	①, 2, 3
$i=1$	②, 3
$i=2$	3
$i=3$	X

I)

```

f(i=0; i<N; i++){
    f(j=0; j<N; j++){
        if(i==j) continue;
        if(A[i] + A[j] == K){
            print(i, j);
        }
    }
}

```

TC:

$$N \times N \times O(1)$$

$$TC: O(N^2)$$

II

```

f(i=0; i<N-1; i++){
    f(j=i+1; j<N; j++){
        if(A[i] + A[j] == K){
            print(i, j);
        }
    }
}

```

i	j, [i+1..N-1]
0	[1, N-1] → N-1
1	[2, N-1] → N-2
2	[3, N-1] → N-3
3	
.	
N-2	(N-2, N-1) → 2

$$TC = O(N^2)$$

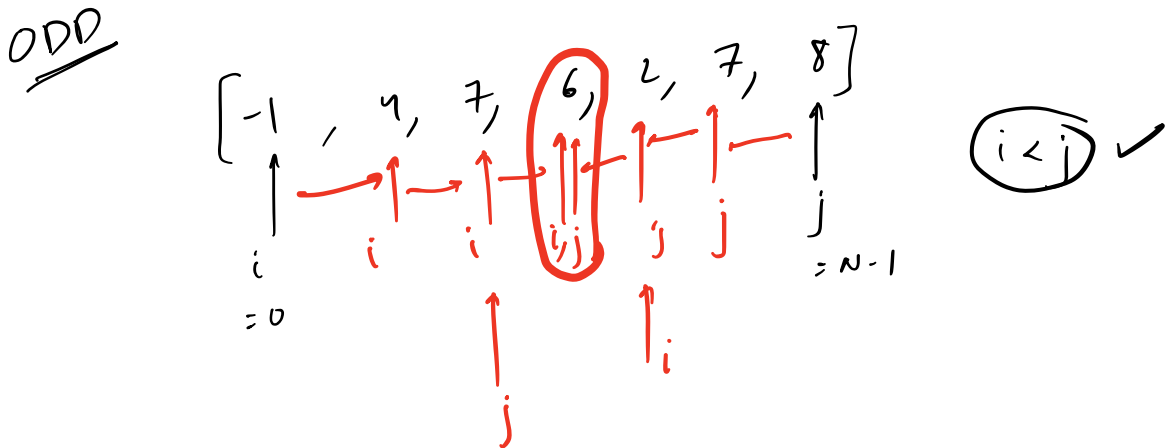
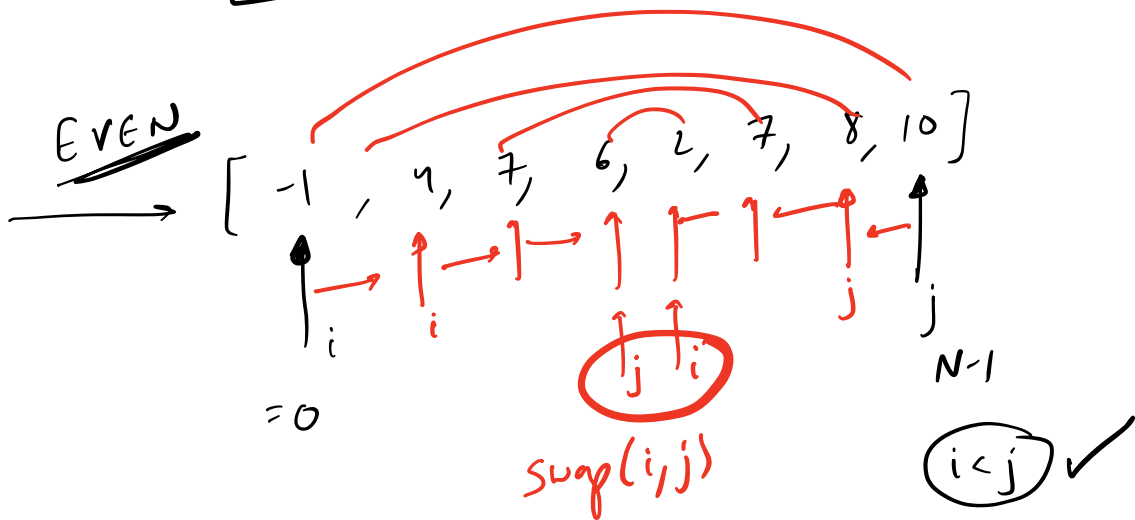
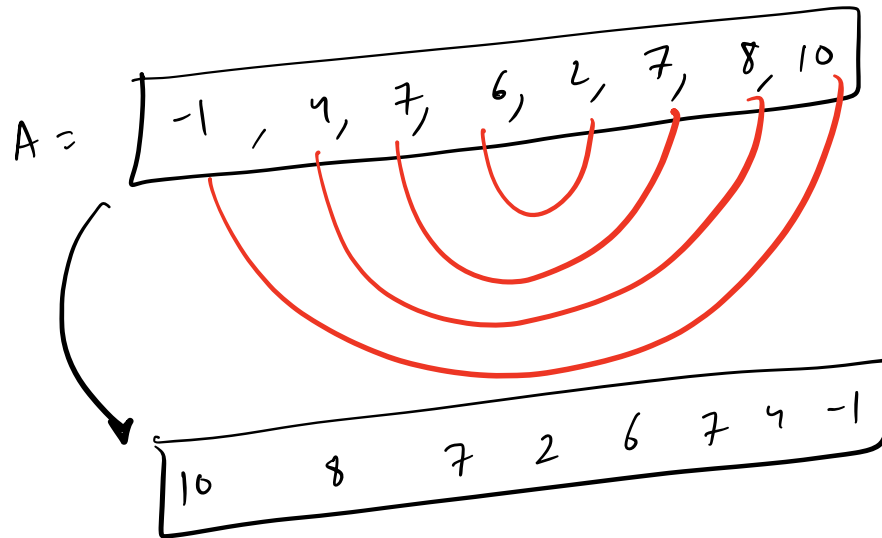
$$SC: O(1)$$

$$2 + 3 + 4 \dots N-1$$

$$\frac{N(N-1)}{2} - 1$$

Q Given an Array. Reverse it!

IN PLACE!



// A[], N

i = 0, j = N-1;

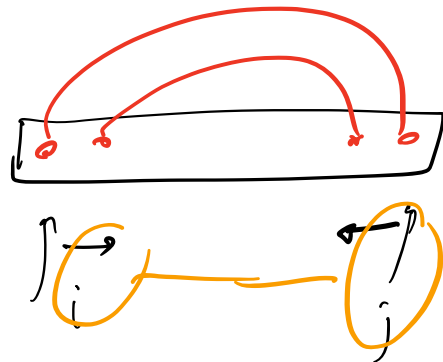
while (i < j) {
 swap(A[i], A[j]);

 i++;

 j--;

}

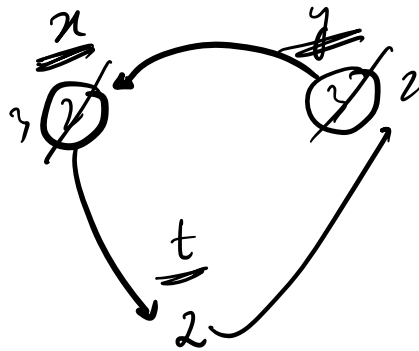
2



↘ O(1)

#it → $N/2$

TC = $O(N)$

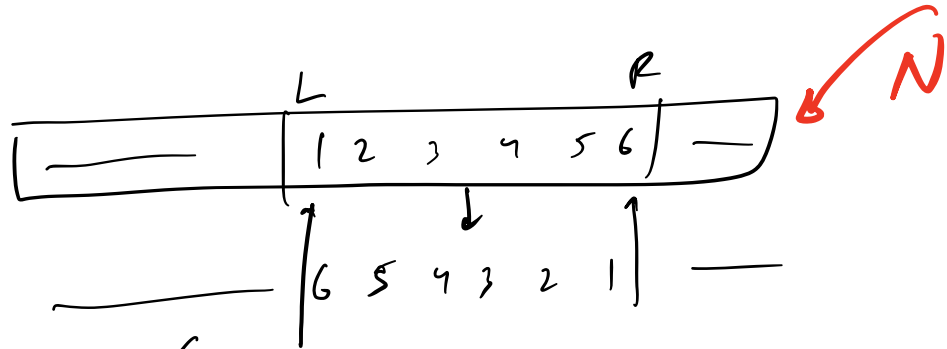


t = x
x = y
y = t

Q

Given an Array.
Reverse the subarray $[L, R]$.

$R-L+1$



sw ($A[]$, L , R) {

$i = L$; $j = R$;

while ($i < j$) {

swp ($A[i]$, $A[j]$);

$i++$;

$j--$;

}

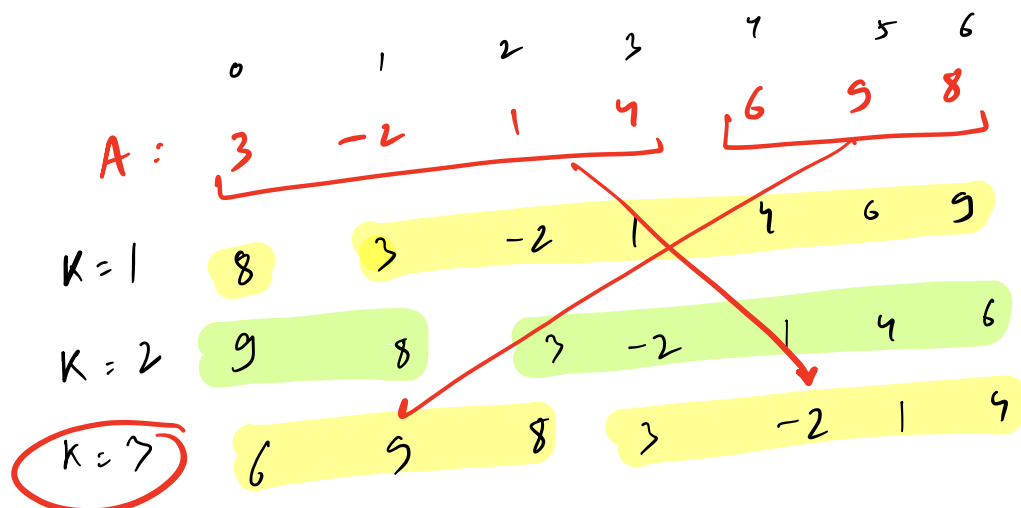
TC : $O(N)$

TC : $O(R-L)$

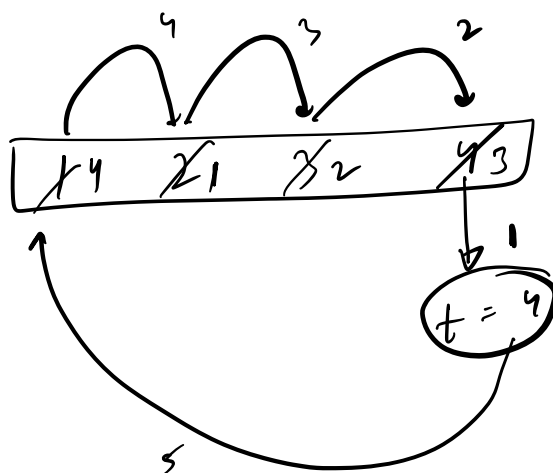
Q

Given an Array.

Rotate the array from last to first K times



1) BF

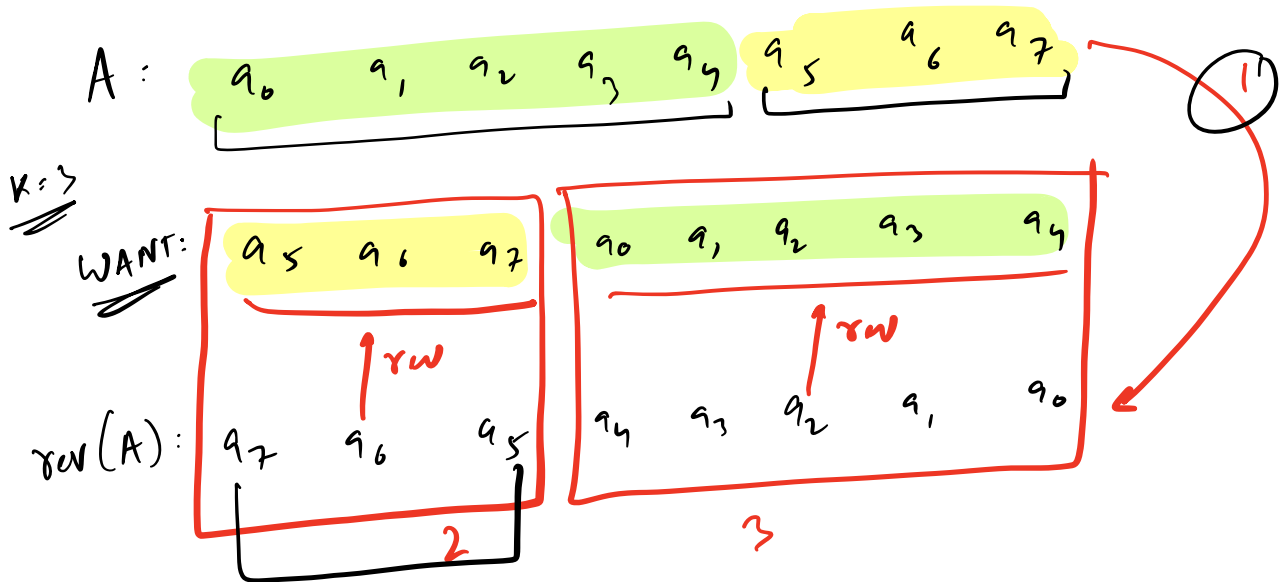
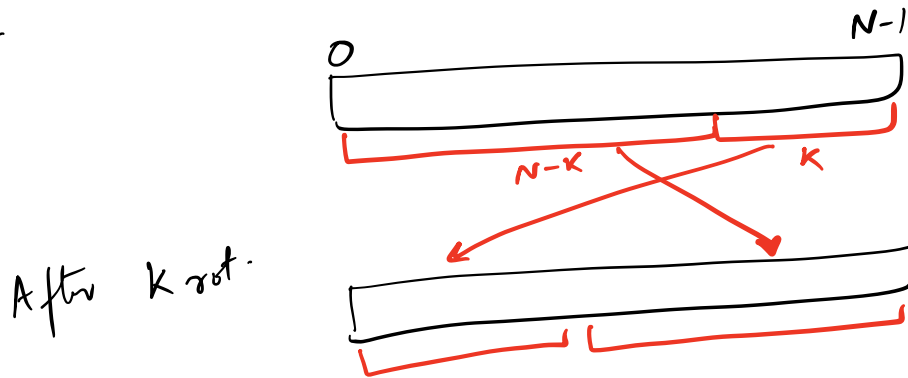


rot 1 time
 \downarrow
 $O(N)$

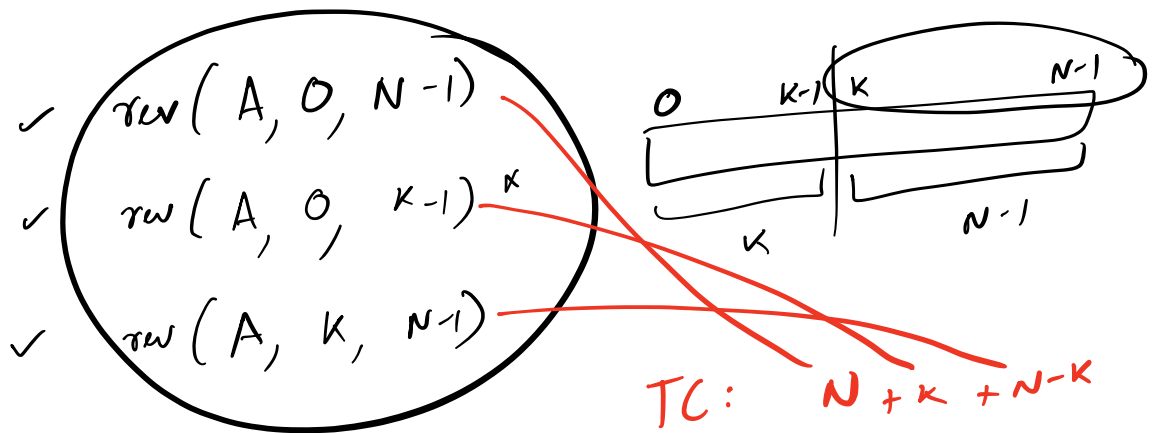
rot K times

$\boxed{TC: O(K \cdot N)}$

II

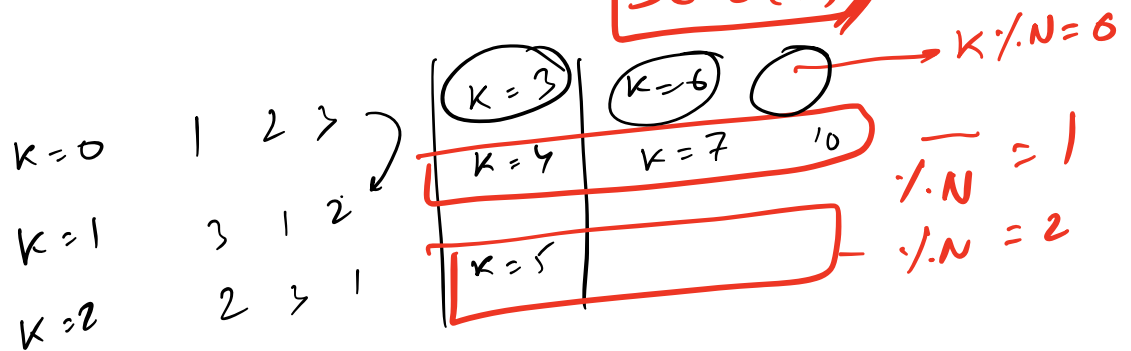


- Steps:
- 1) $A \rightarrow rev(A)$
 - 2) Rev the first K elements
 - 3) Rev the last $N-K$ element



$TC: O(N)$

$SC: O(1)$



$N = 3$
 $K = 10$

$10 \% 3 = 1$
 $K = K \% N$

$K = 0$

